



Doubt Class With Lakshay Bhaiya

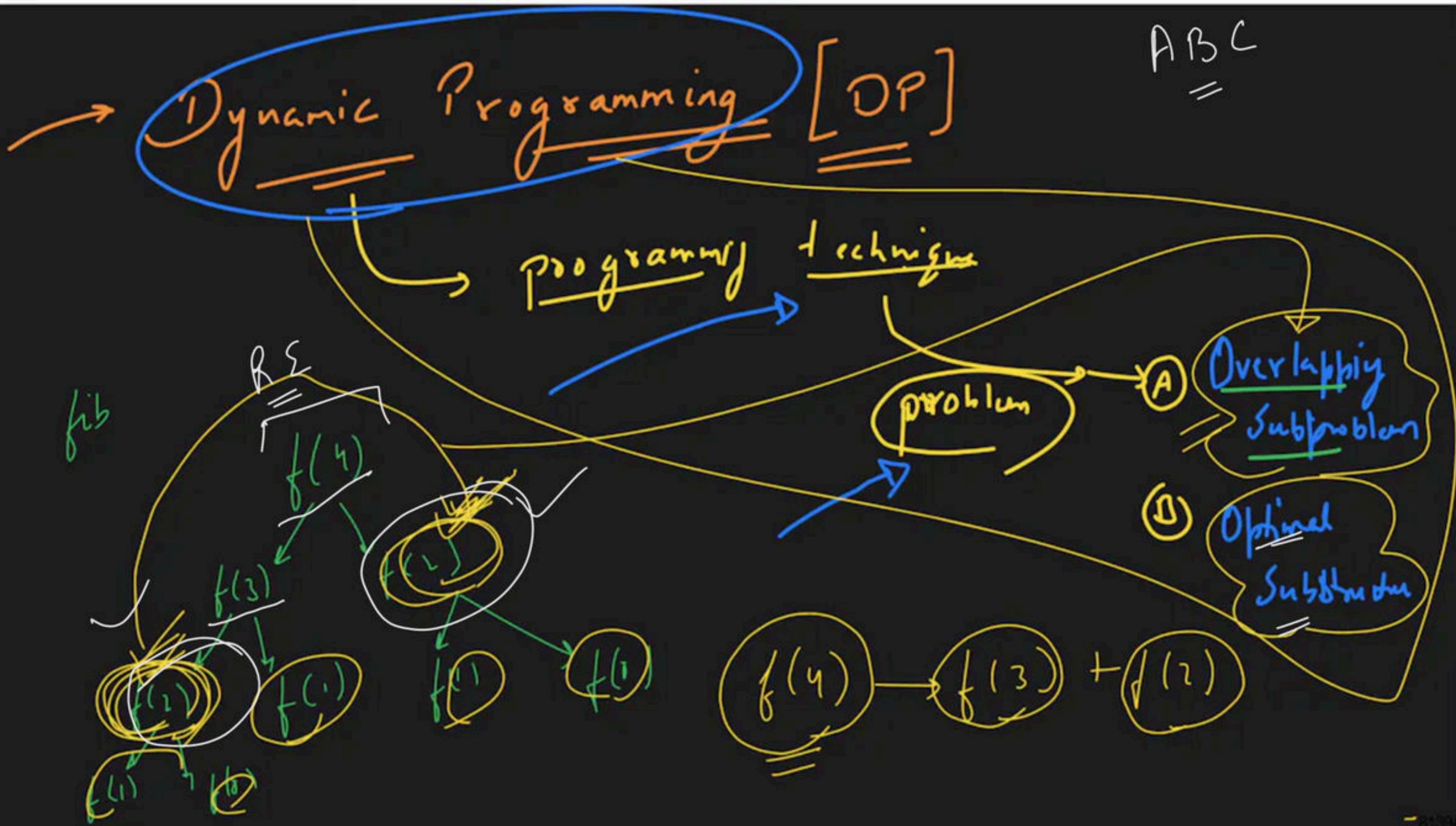
Special class

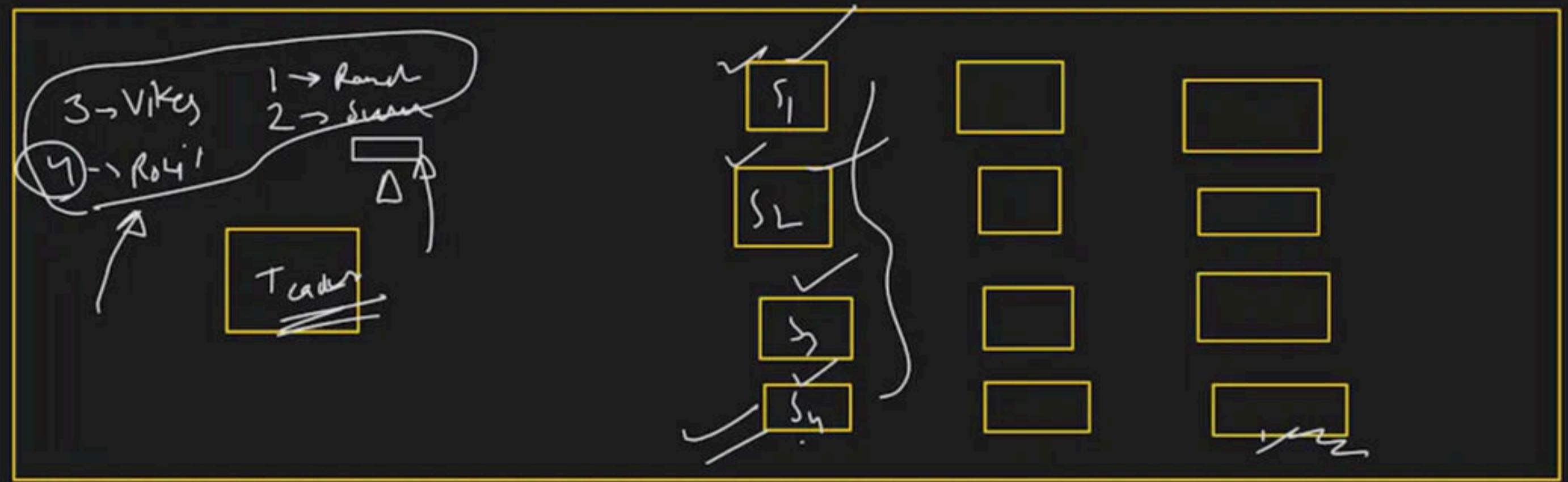
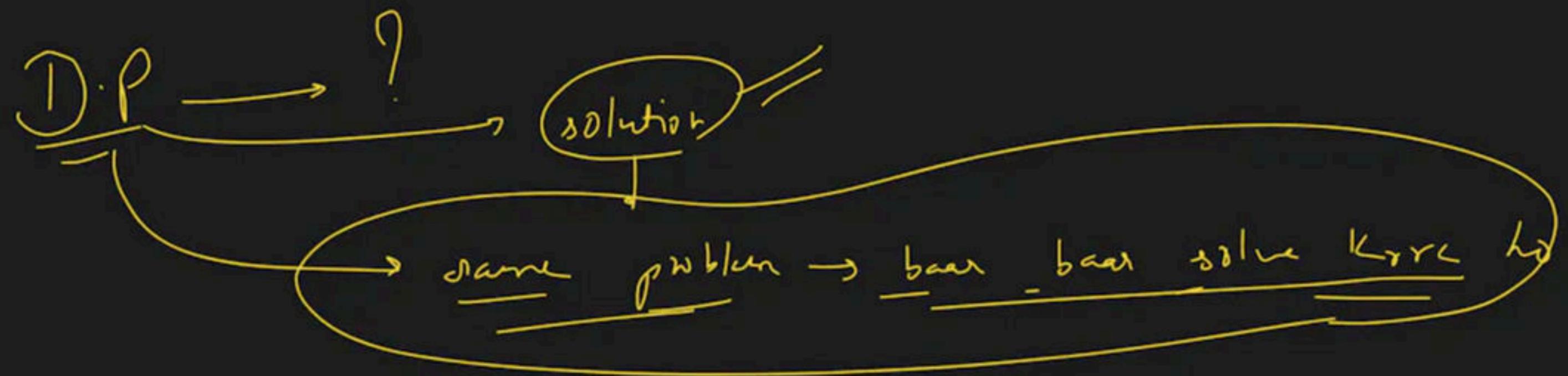


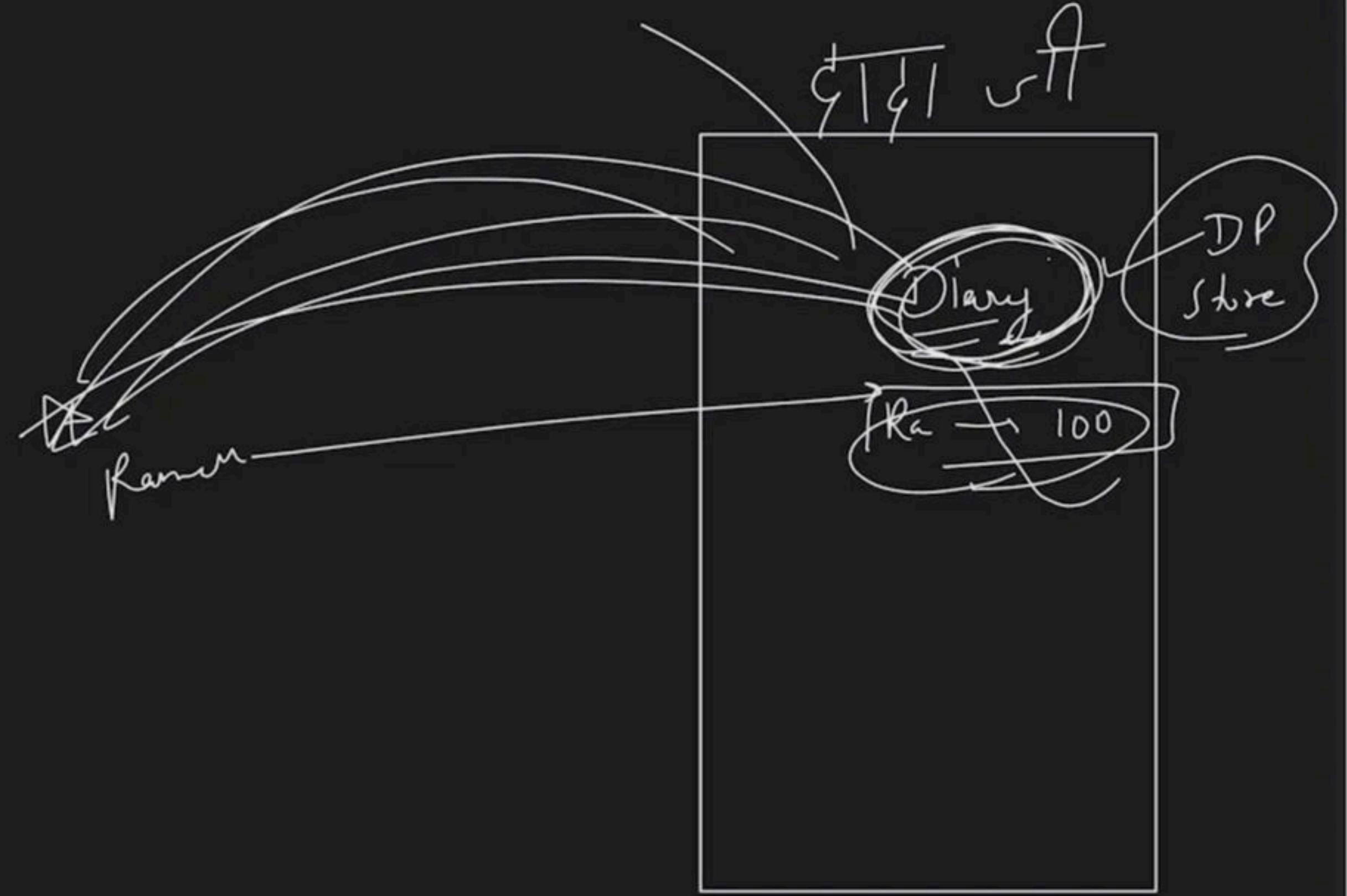
Dynamic Programming Class-1

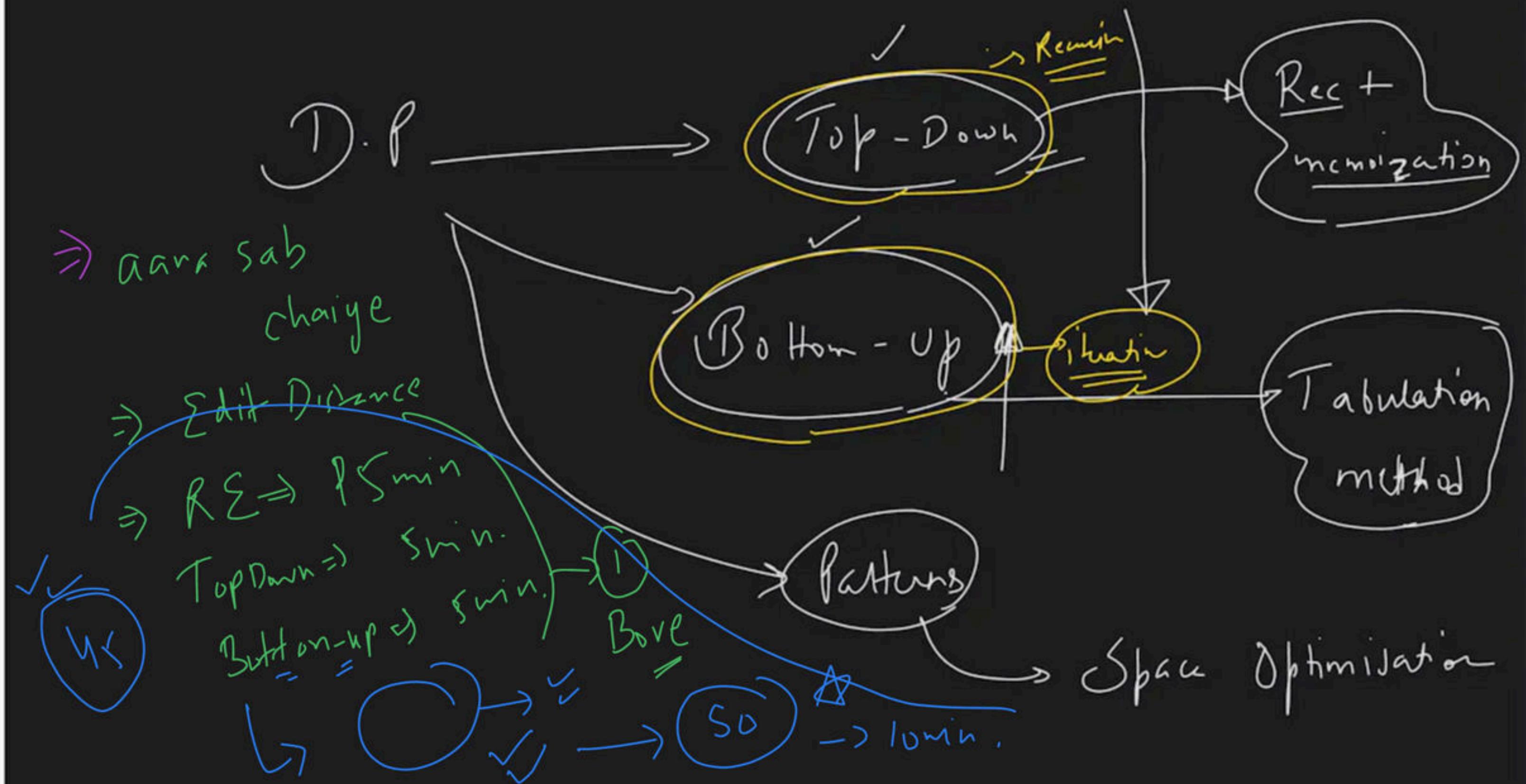
Special class

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fibonacci no:-

assum $\rightarrow R.C$

0, 1, 1, 2, 3, 5, 8, 13, 21, ..., - - -

$$2 = 1 + 1$$

Recursion

R.R \rightarrow

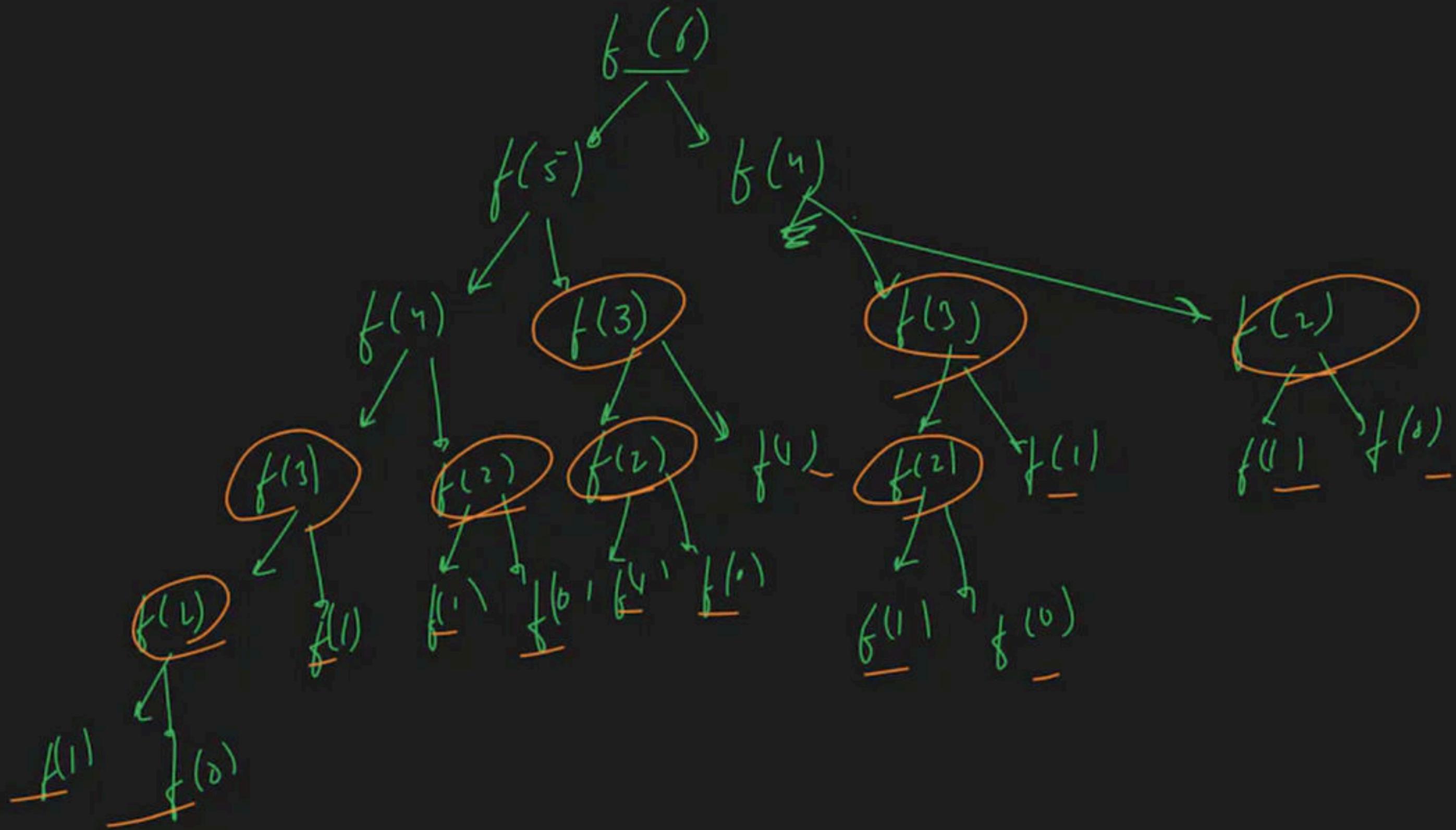
$$f(n) = f(n-1) + f(n-2)$$

$$\underline{n^{\text{th}} \text{ fibonacci no}} = (n-1)^{\text{th}} \text{ fib no} + (n-2)^{\text{th}} \text{ fib no}$$

$$\text{let } n = 6$$

$$f(6) = f(5) + f(4)$$

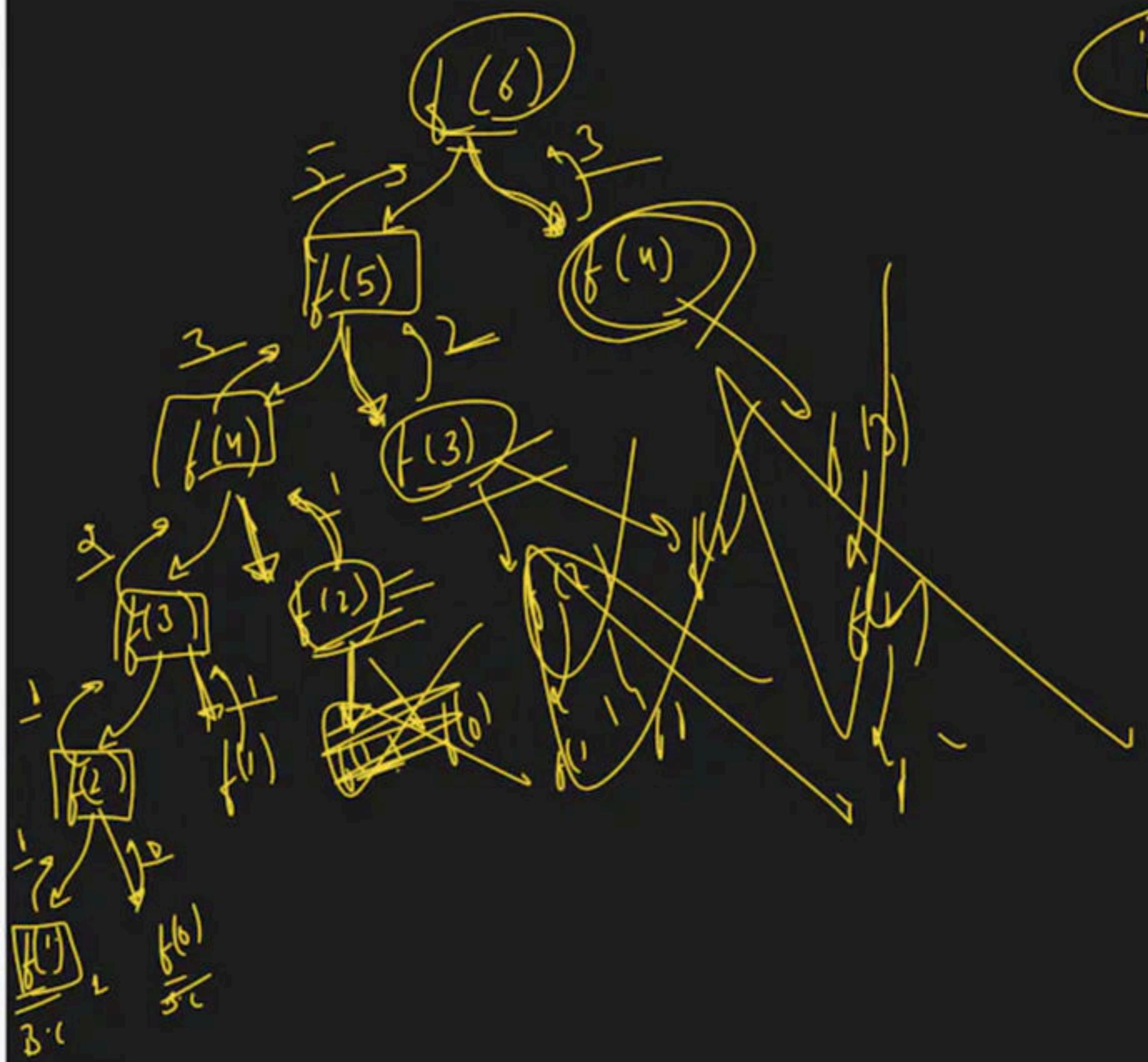
$$\underline{6^{\text{th}} \text{ fib no}} = \underline{5^{\text{th}} \text{ fib no}} + \underline{4^{\text{th}} \text{ fib no}}$$



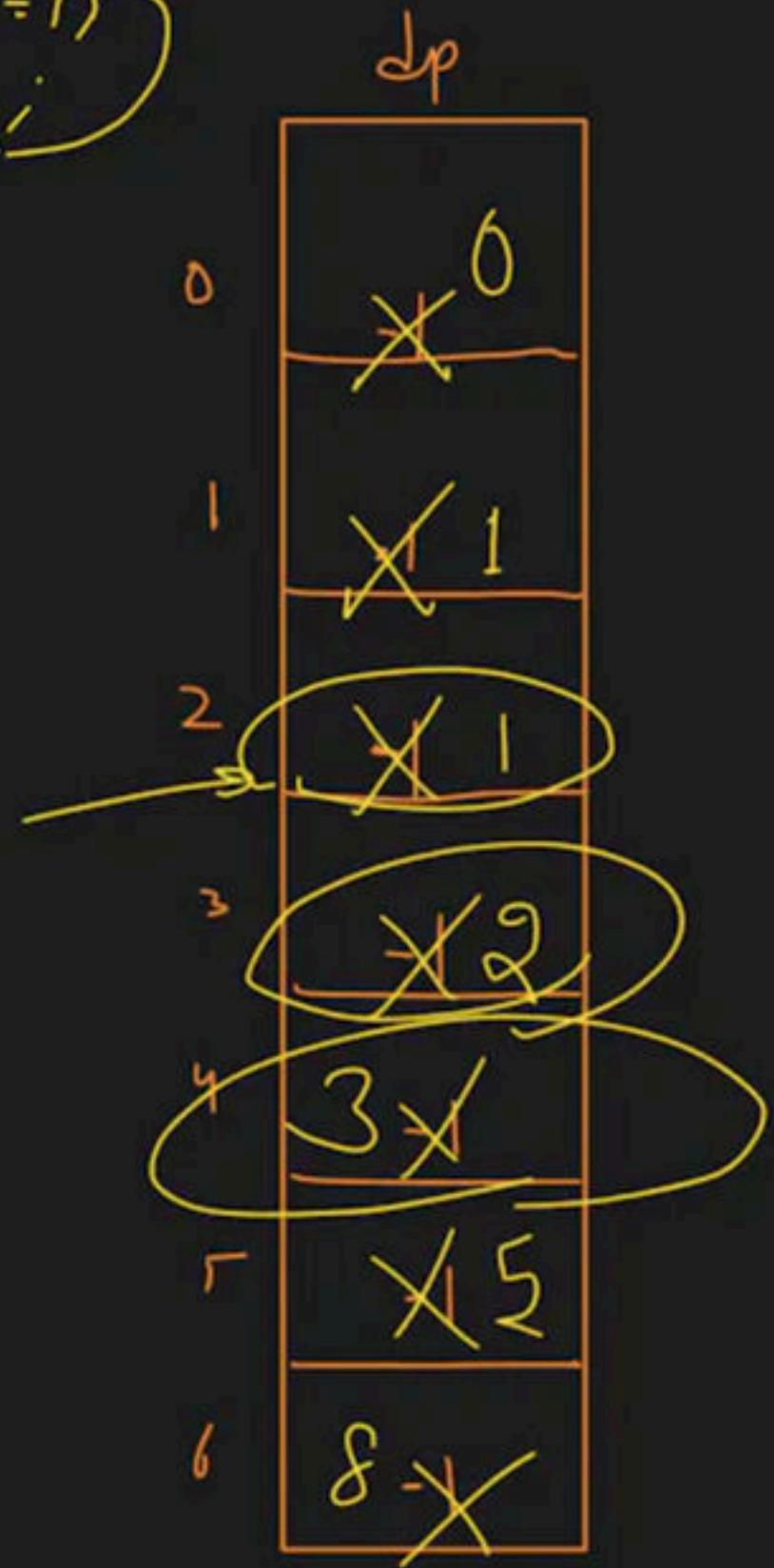
Mon

Step

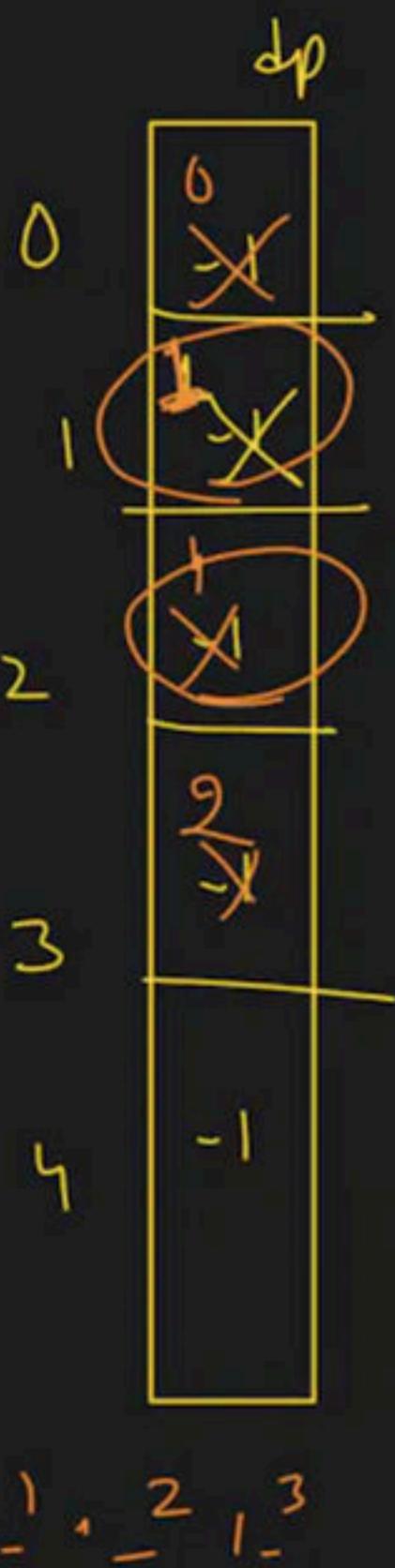
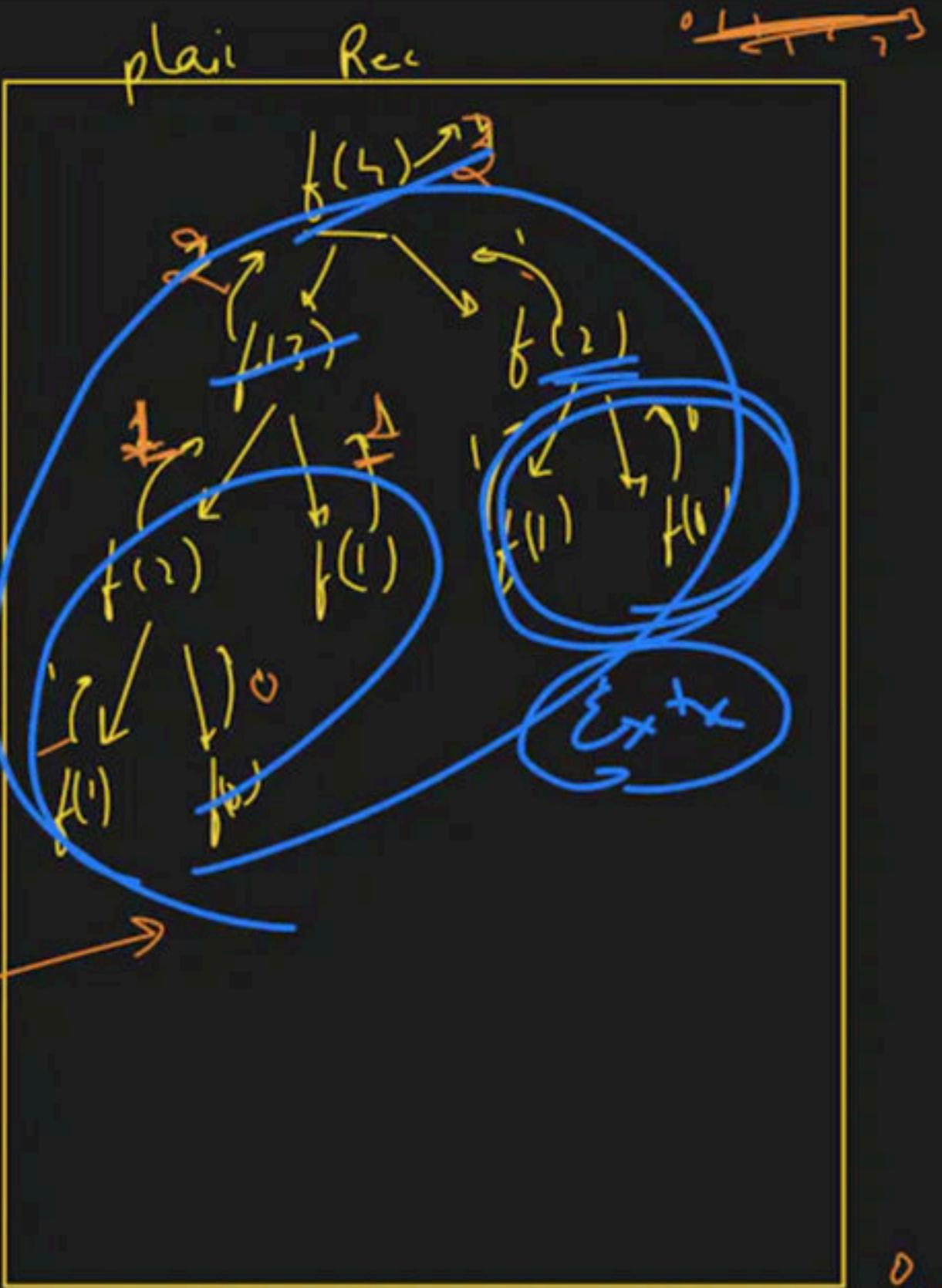
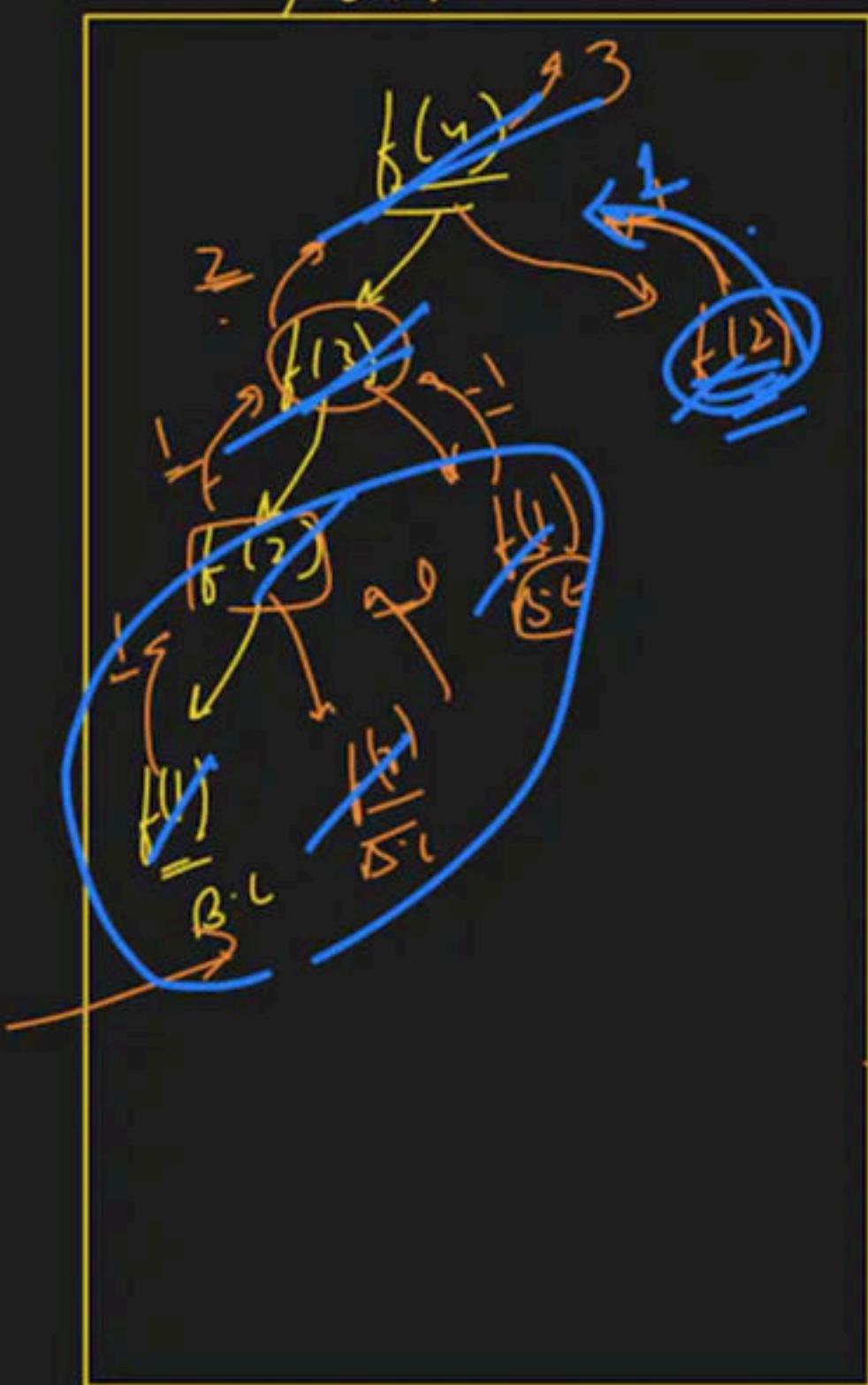
- ① Create a dp array
- ② Store / take an in dp array
- ③ If any already exists in dp array, then return

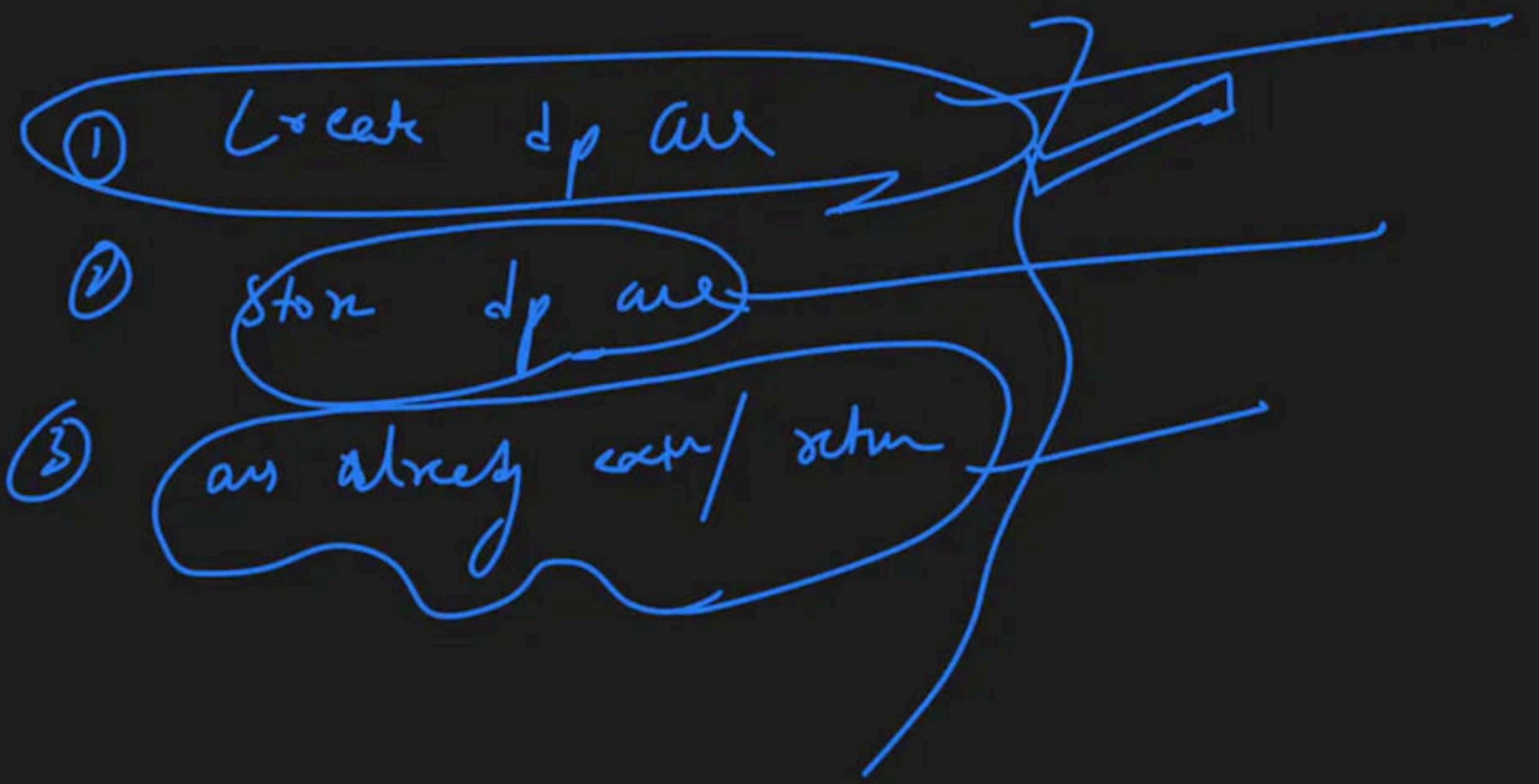


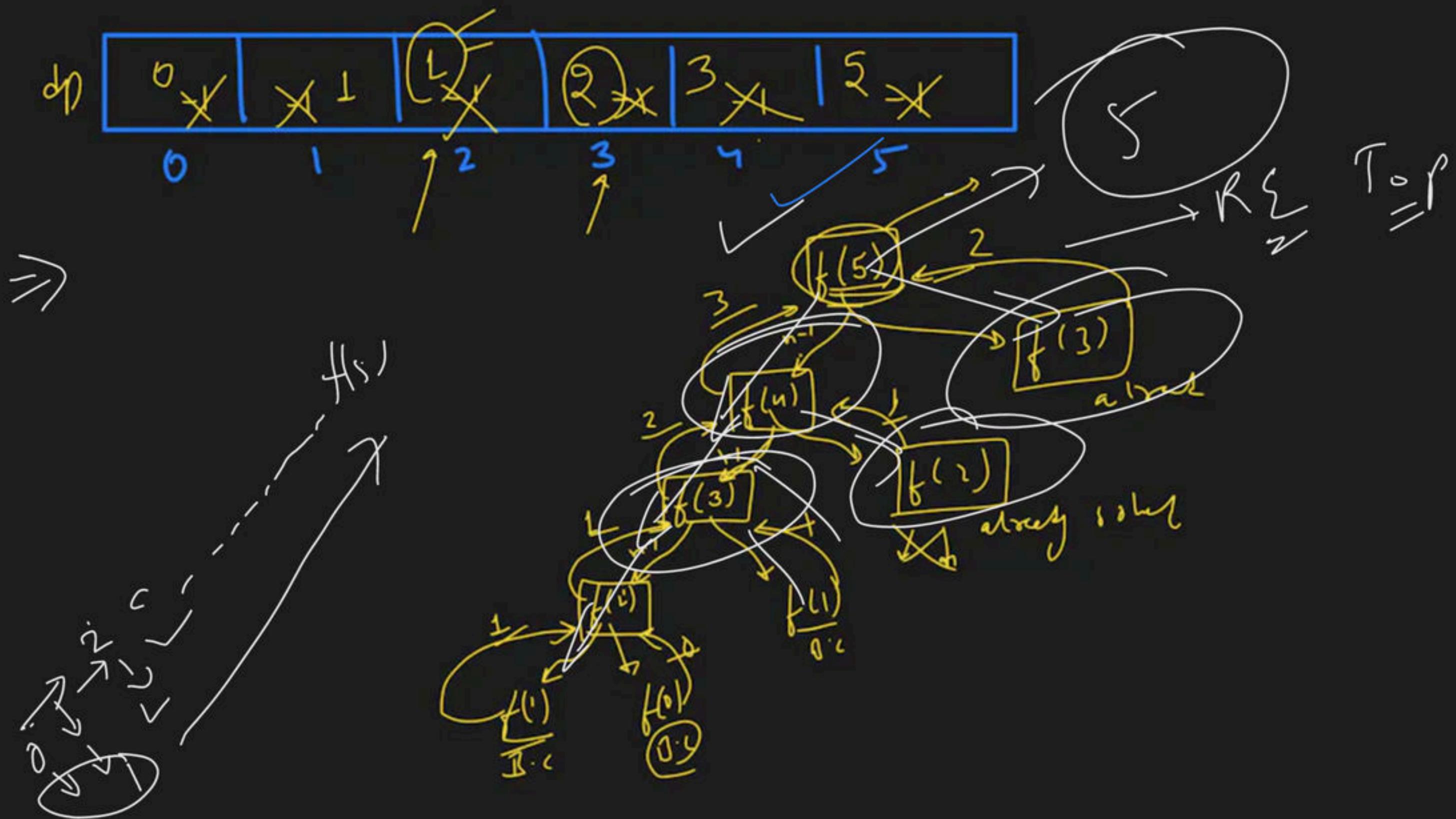
$\text{if } (n = -0 \mid |n| = 1)$
 $\quad \text{return } 2;$



Memoiret

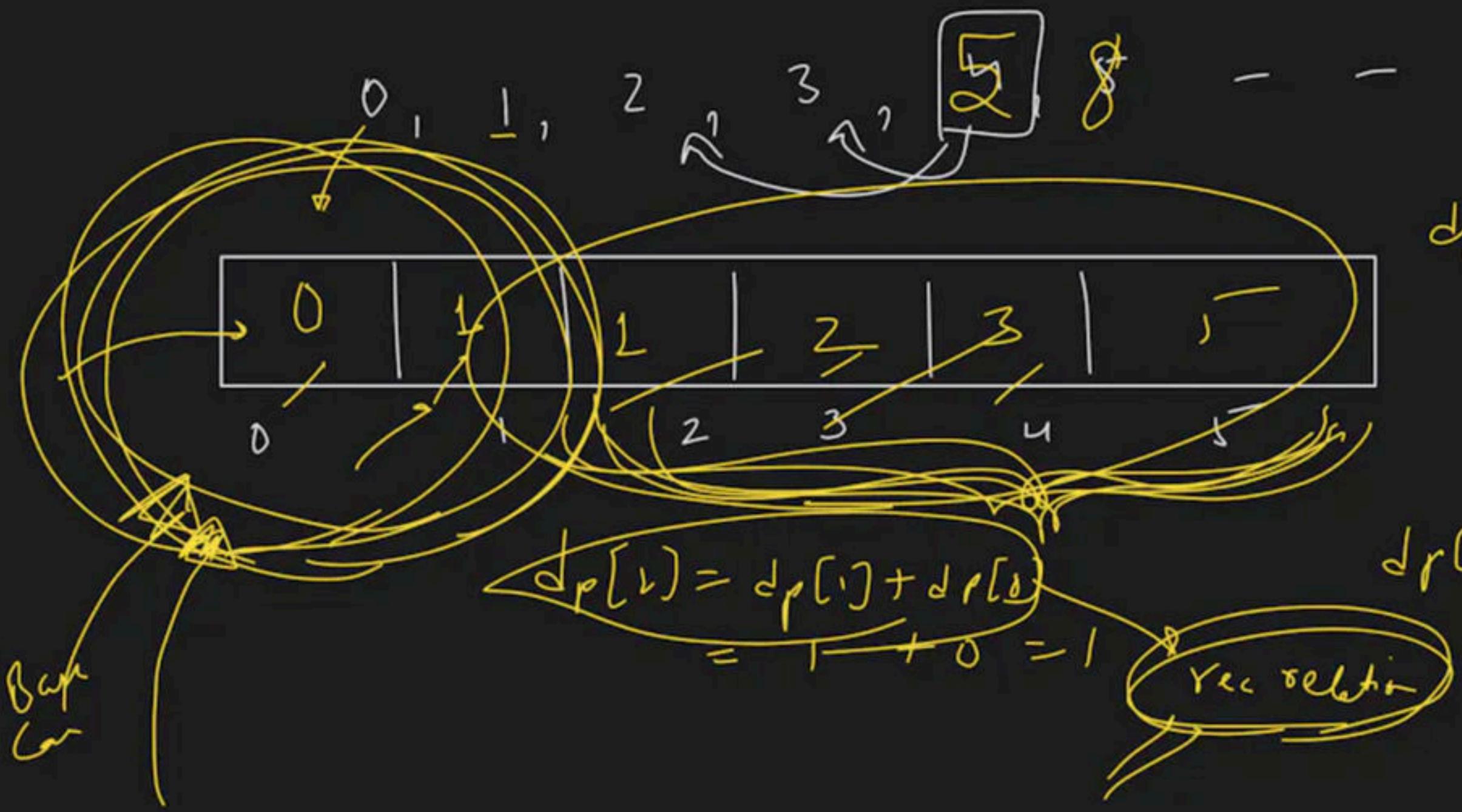






→ Tabulation method (Bottom-up)

$$dp[5] = dp[1] + dp[3]$$
$$= 7 + 1$$



$$dp[1] = dp[0] + dp[0]$$
$$= 1 + 1$$
$$= 2$$

Tabletia

① → check dp array

② → Analyse → Balance → fill dp array according
to logic formula
recursion
relation

③ → fill ternary dp array w/

Span
int solve

Optimisation :-

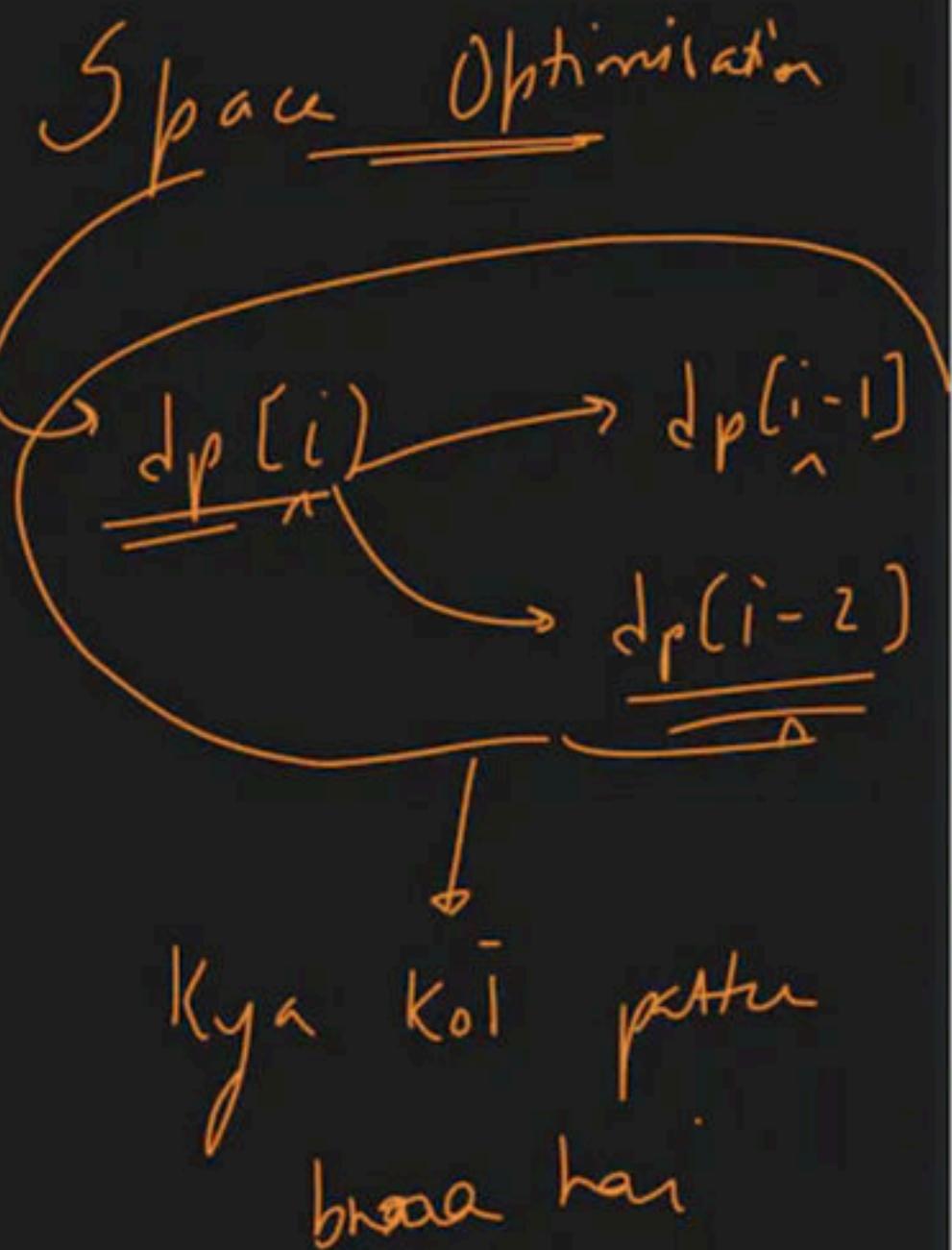


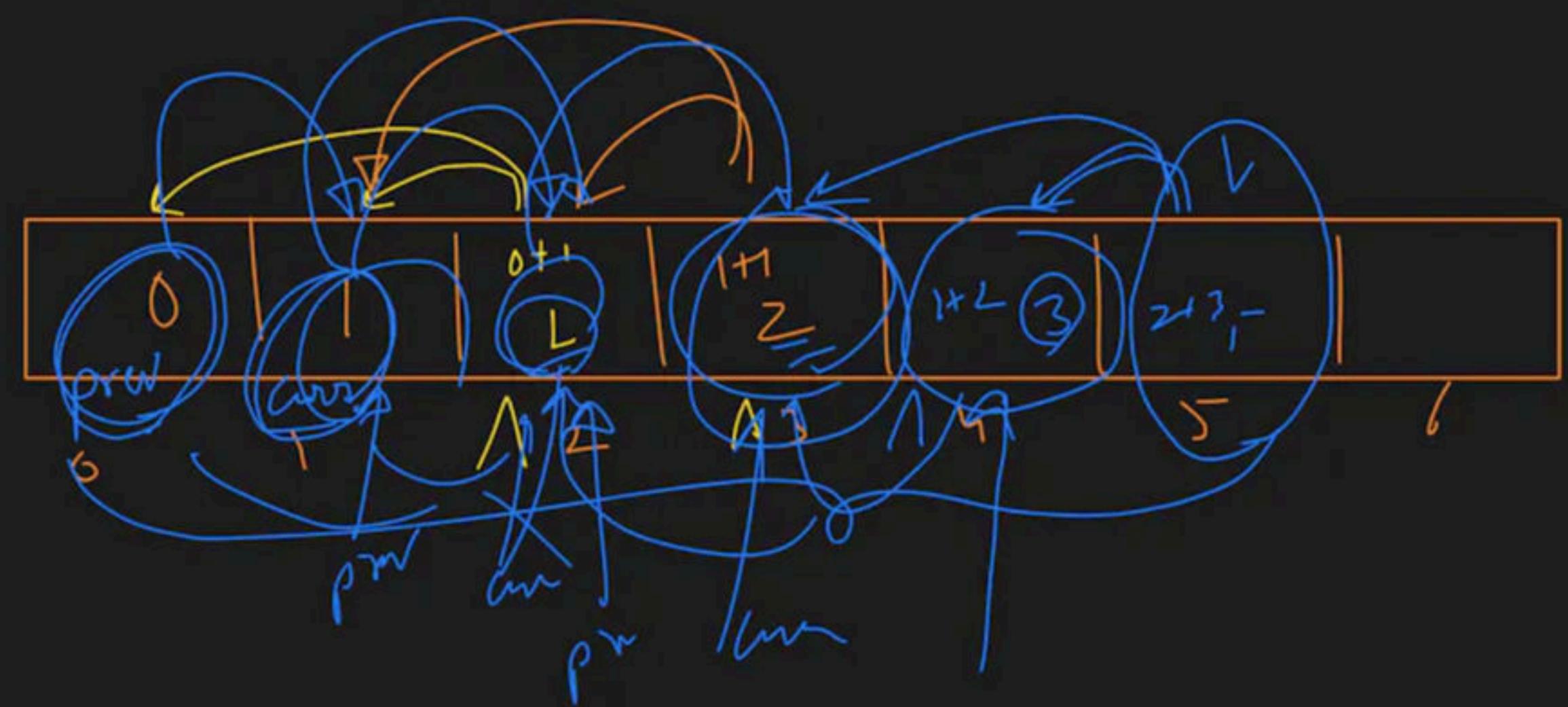
```

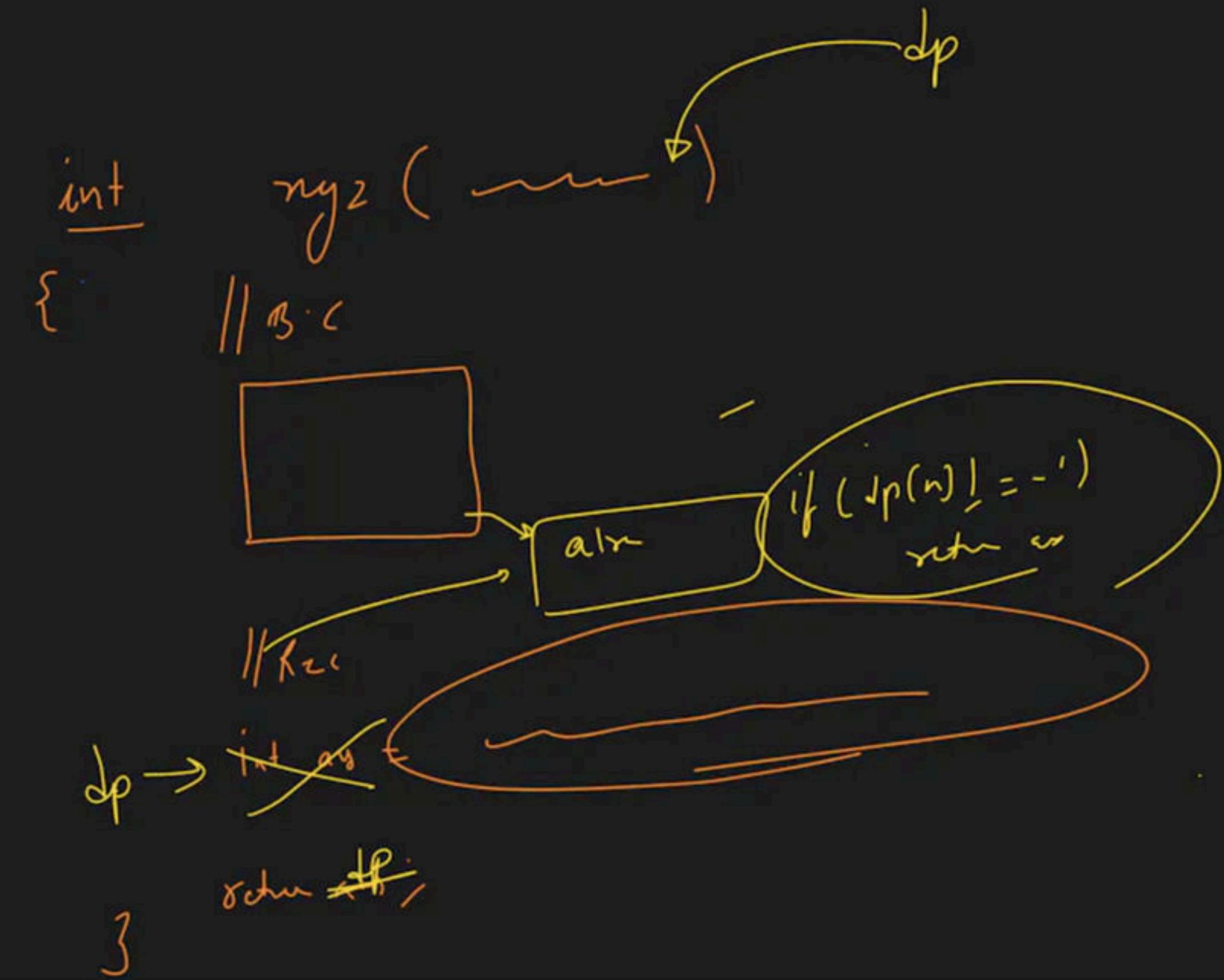
int solveTab(n)
{
    vector<int> dp(n+1, -1);
    dp[0] = 0;
    dp[1] = 1;

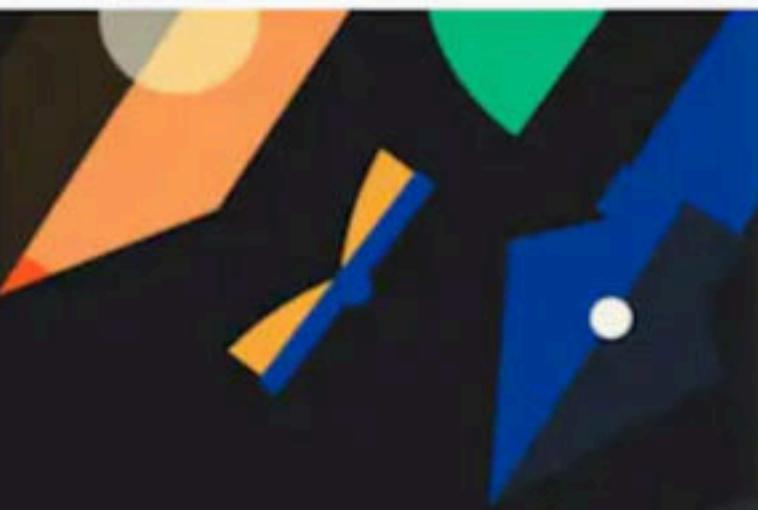
    for (i = 2; i <= n; i++)
    {
        dp[i] = dp[i-1] + dp[i-2];
    }
    return dp[n];
}

```





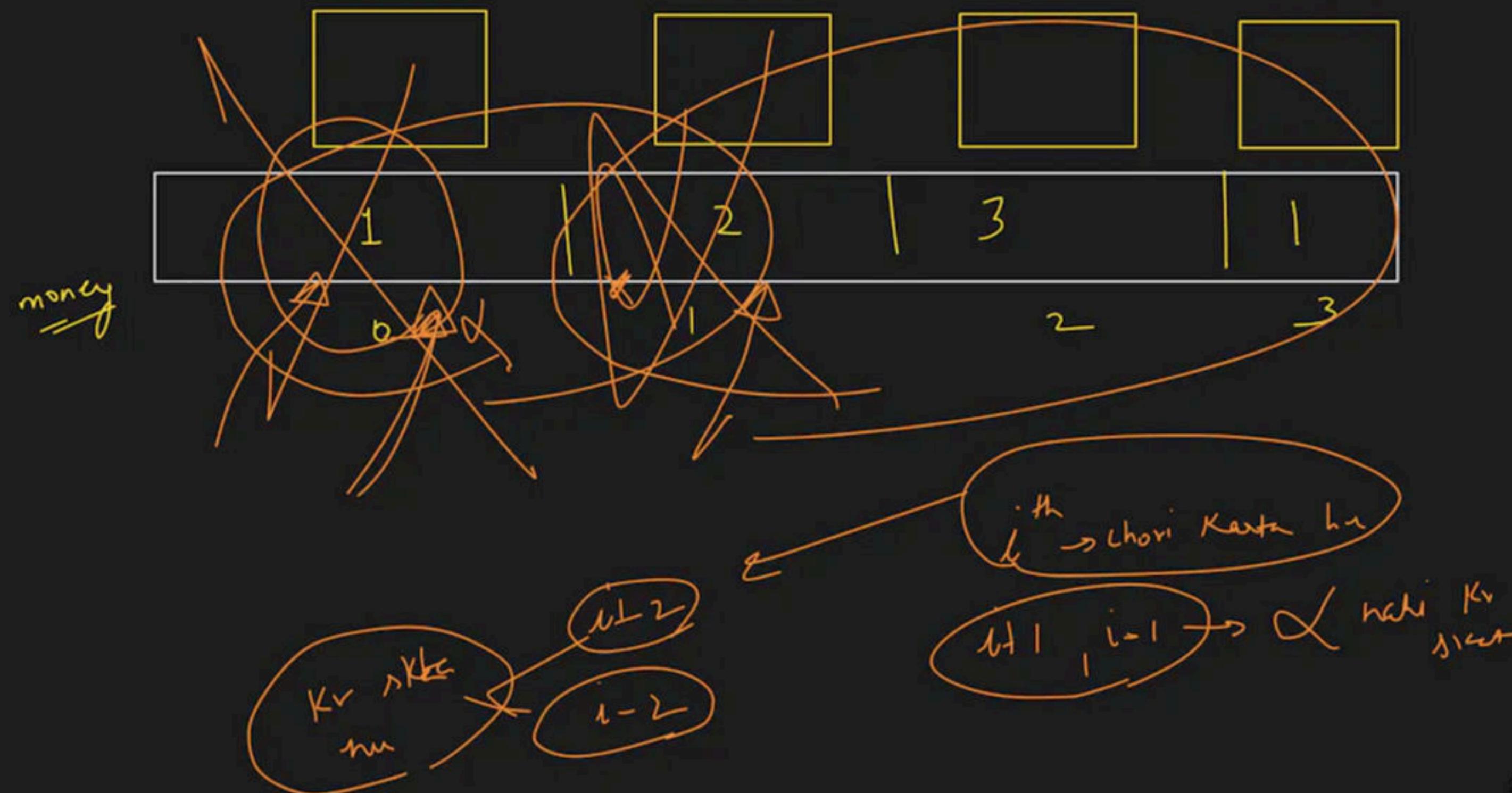


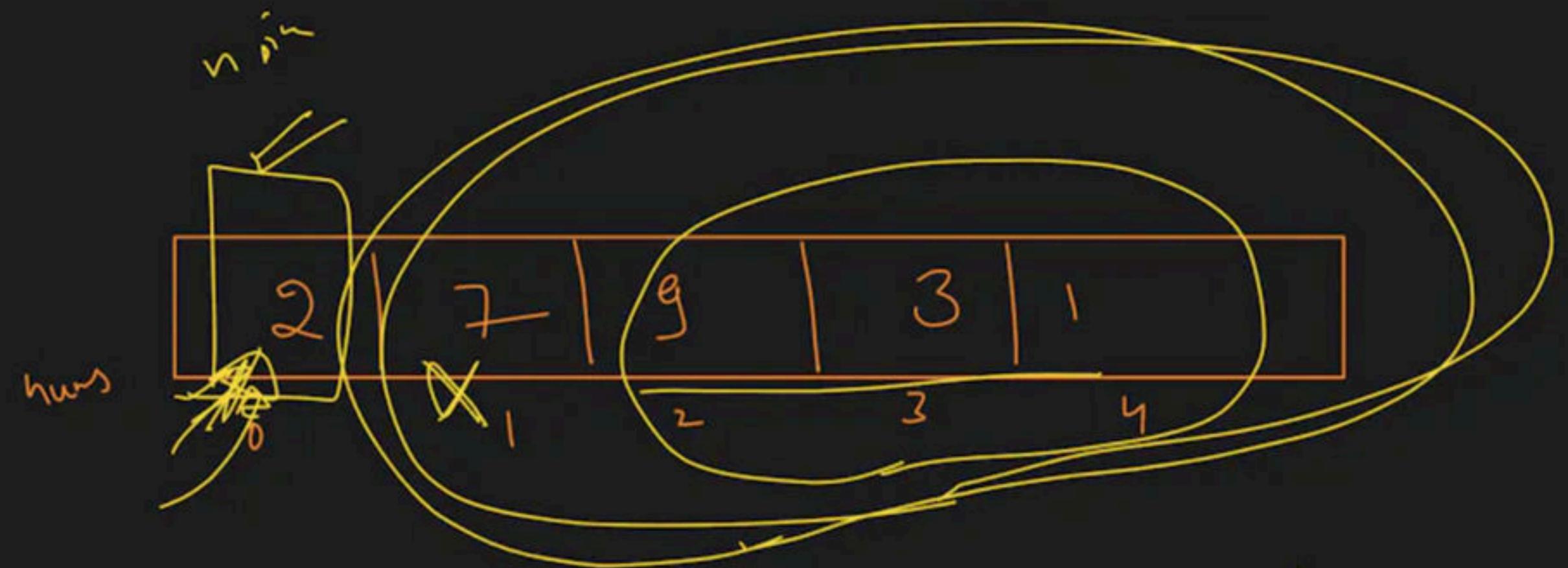


Dynamic Programming Class-2

Special class

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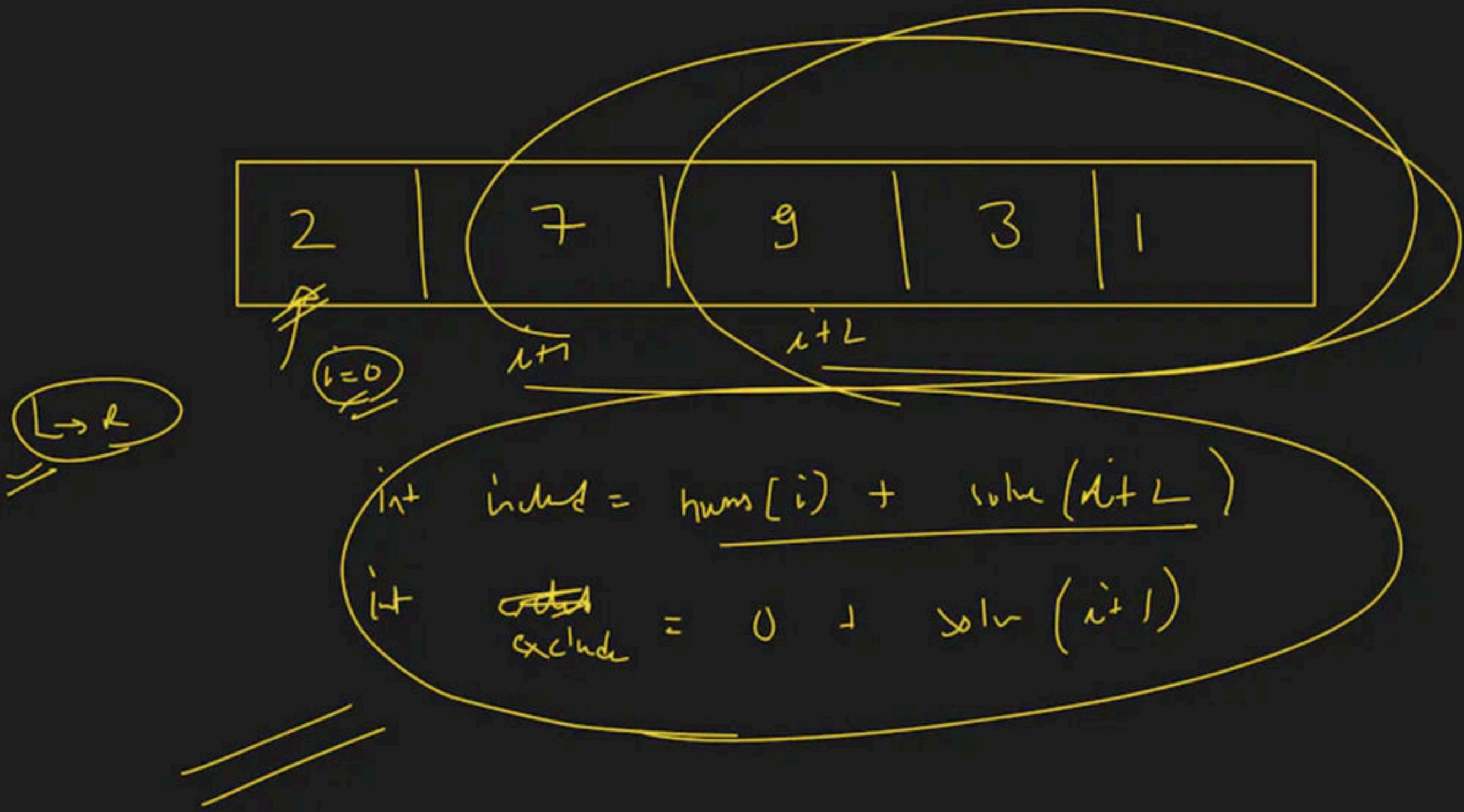


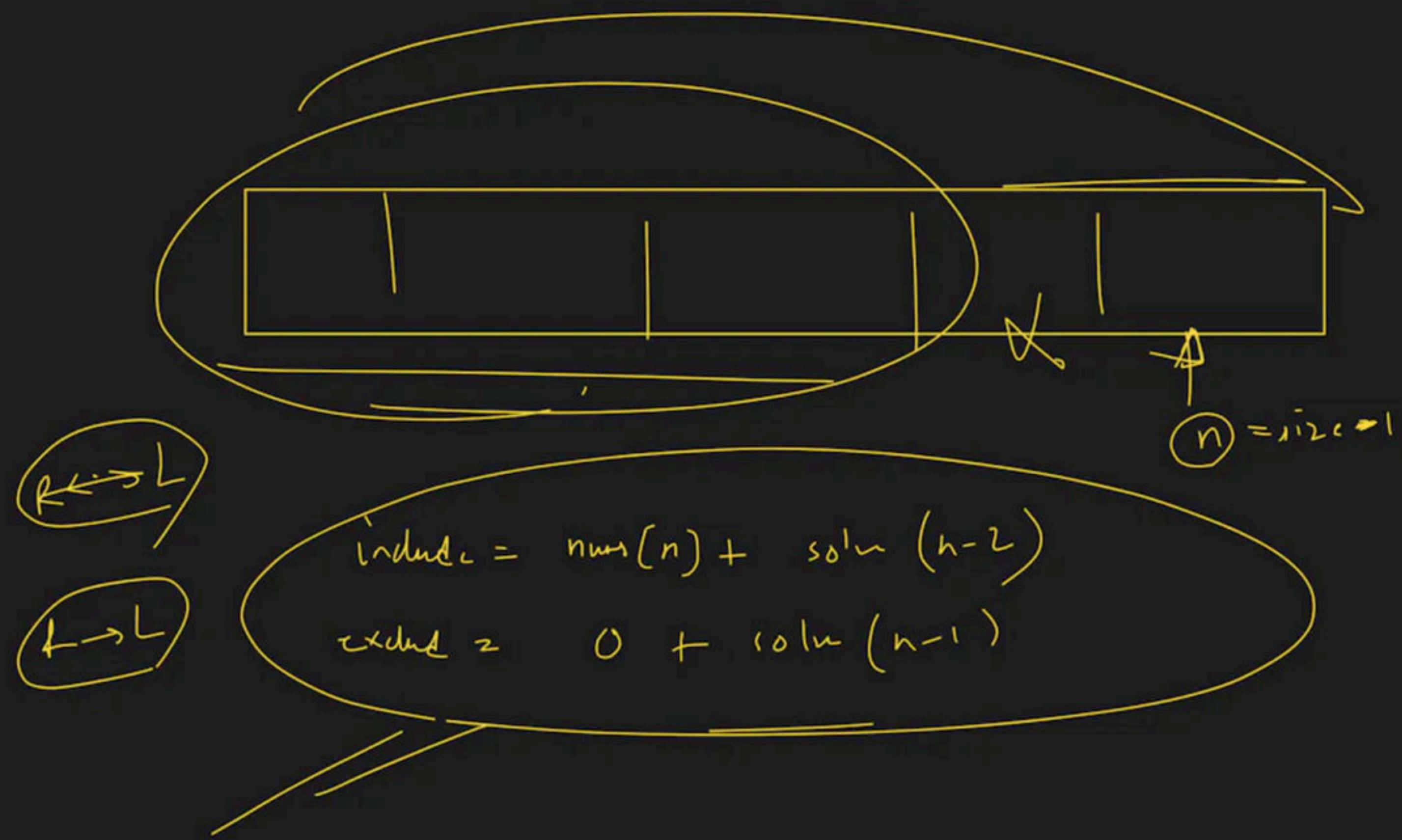


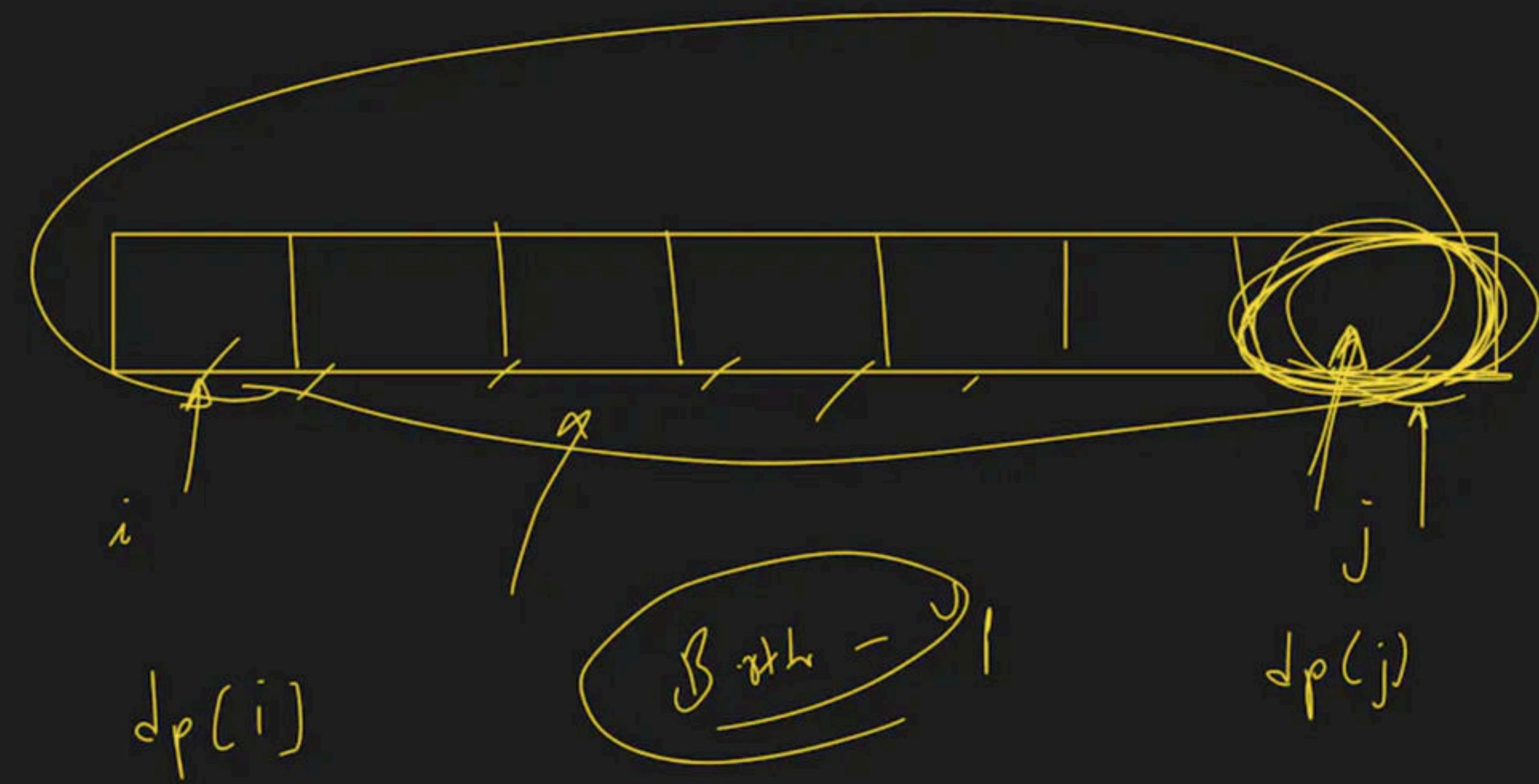
$a_{n+1} \rightarrow$ include $\rightarrow 2 + f^{(n-2)}$

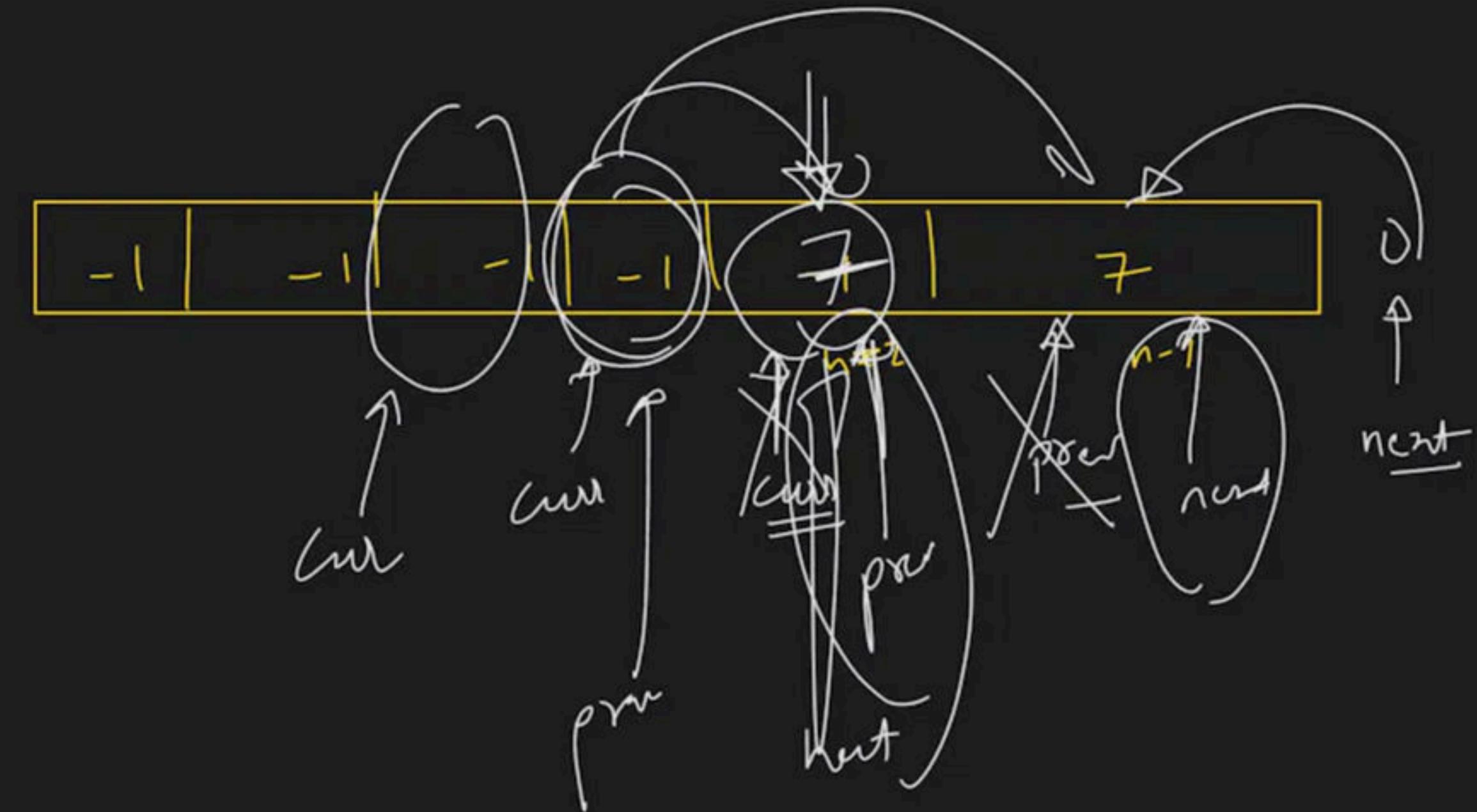
$a_{n+1} \rightarrow$ exclude $\rightarrow 0 + f^{(n-1)}$

find
Ans





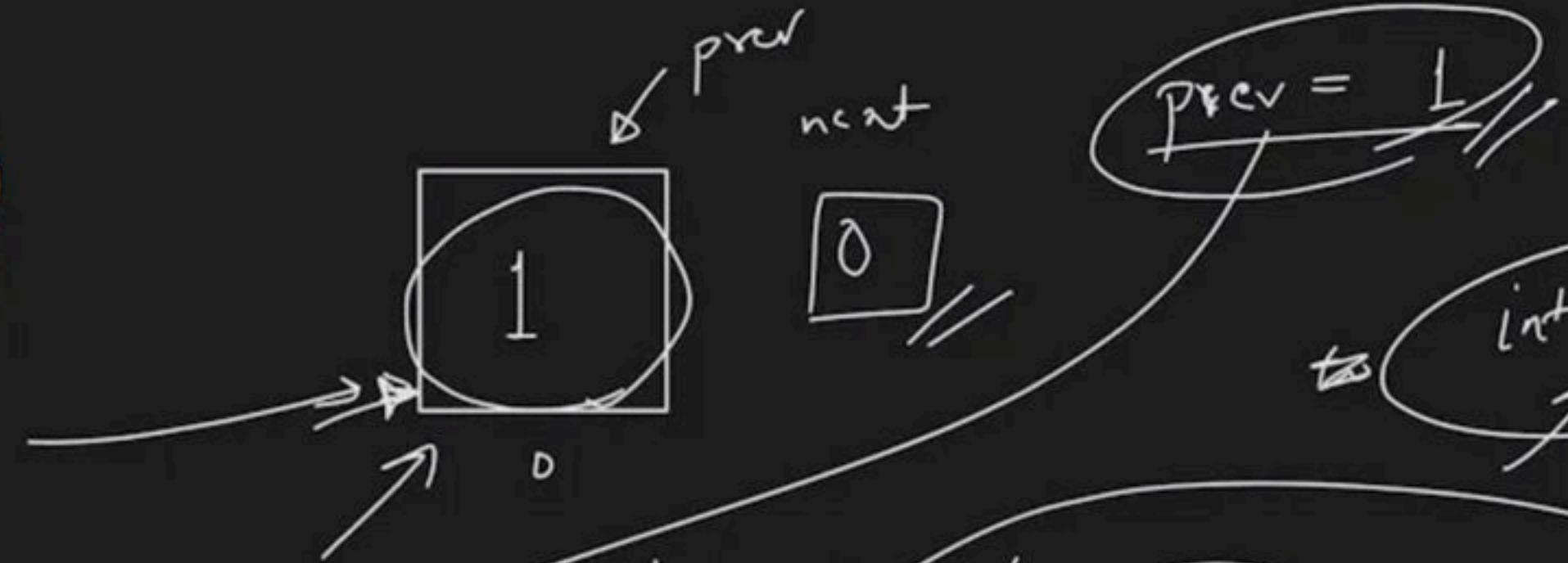




$\epsilon \times inc$

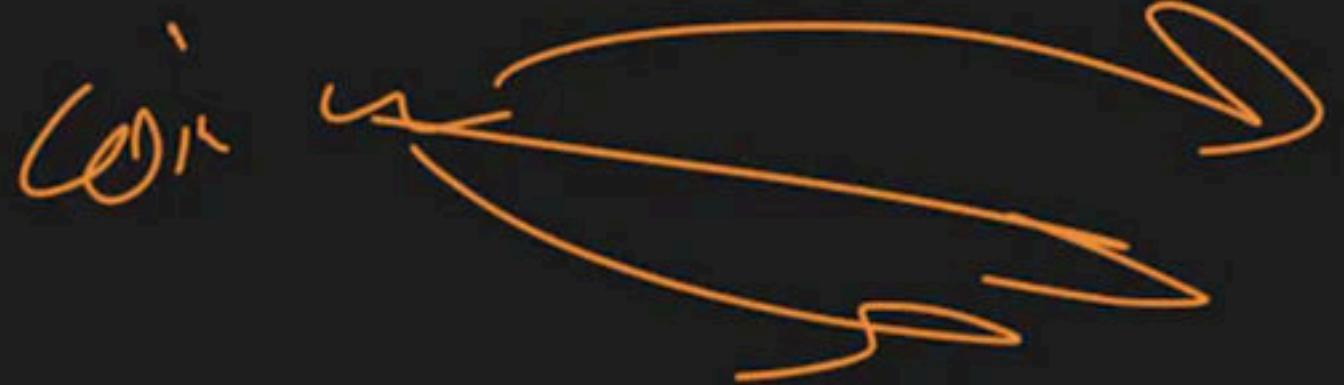
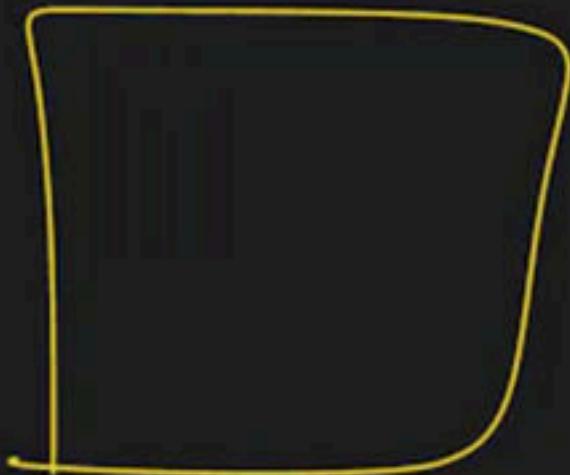
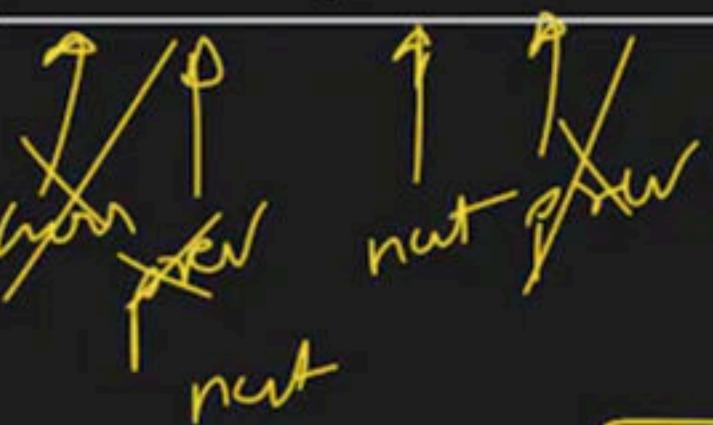
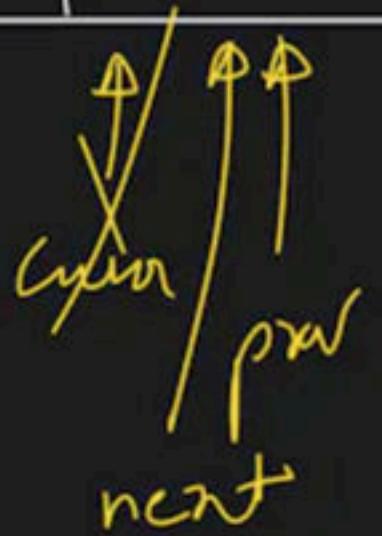
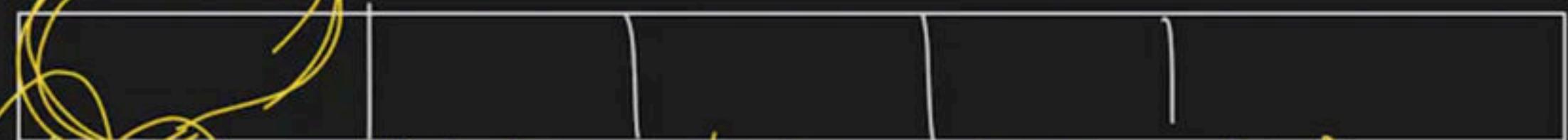
dp

return $prev$



for
 $i=2$ to $h-2$
 $-1 >= 0$

loop
chalgan veli

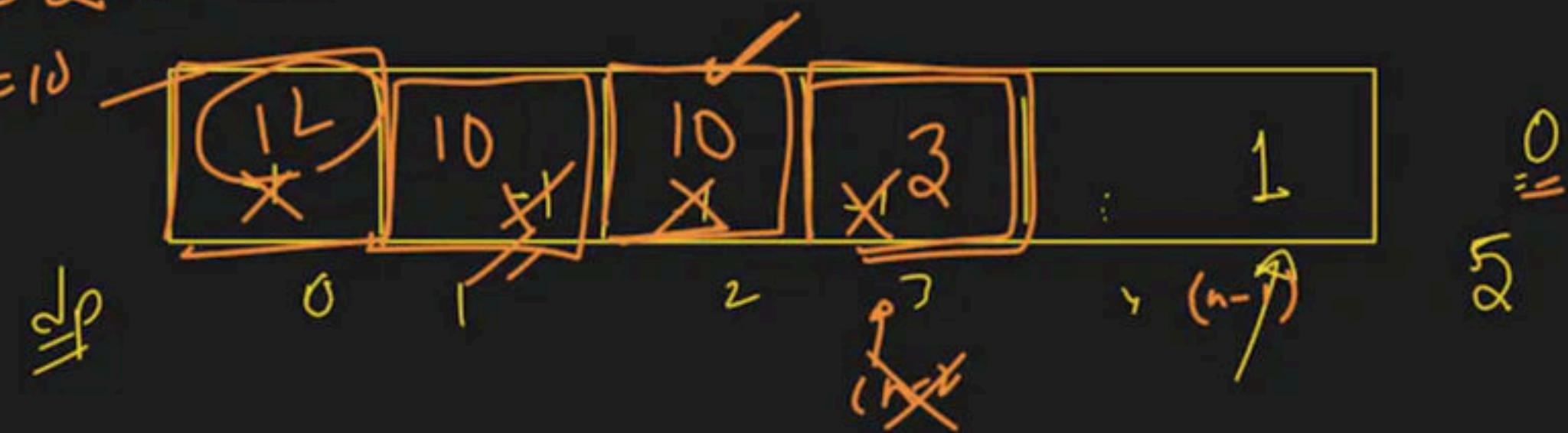


shift



$$\text{inclusion} = 2 + 10 = 12$$

$$exclusion = 10$$



inclusion 10

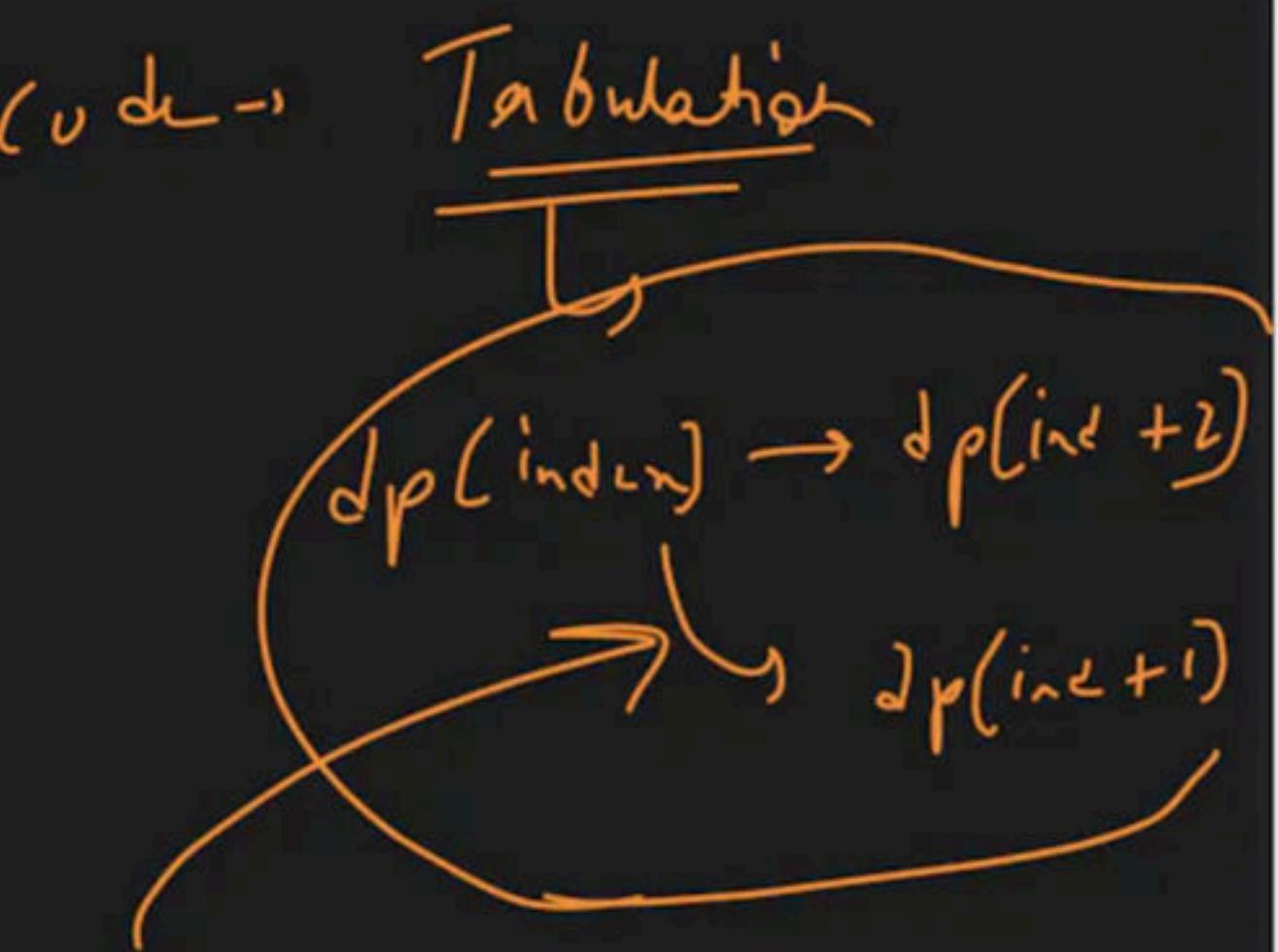
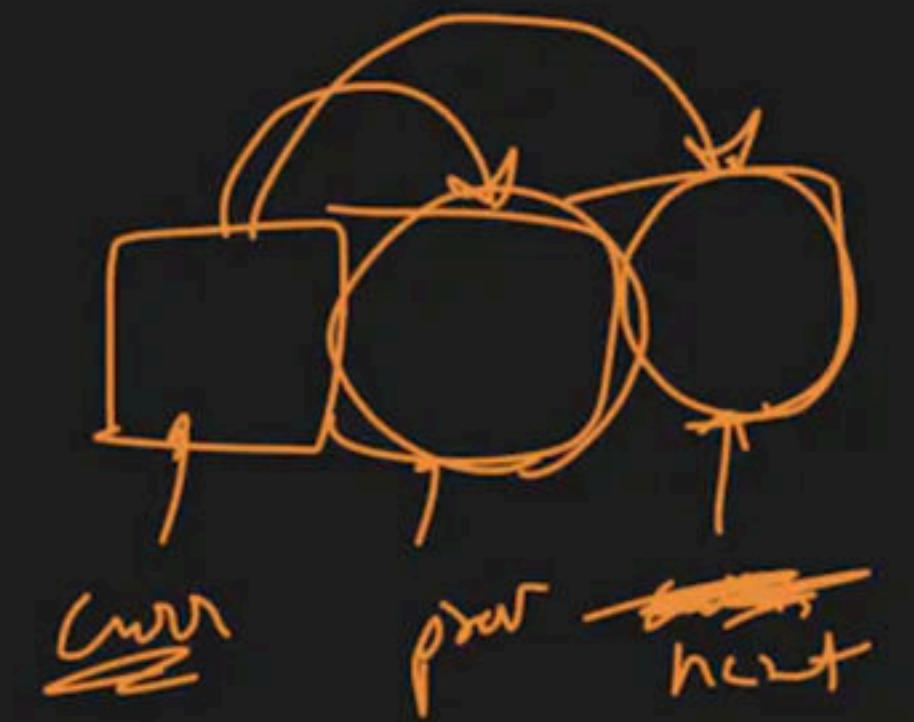
exclusion 3

for ($n - 2$) $\rightarrow 0$

$\text{inclusion} \rightarrow 10 -$
 $\text{exclusion} \rightarrow 10 /$

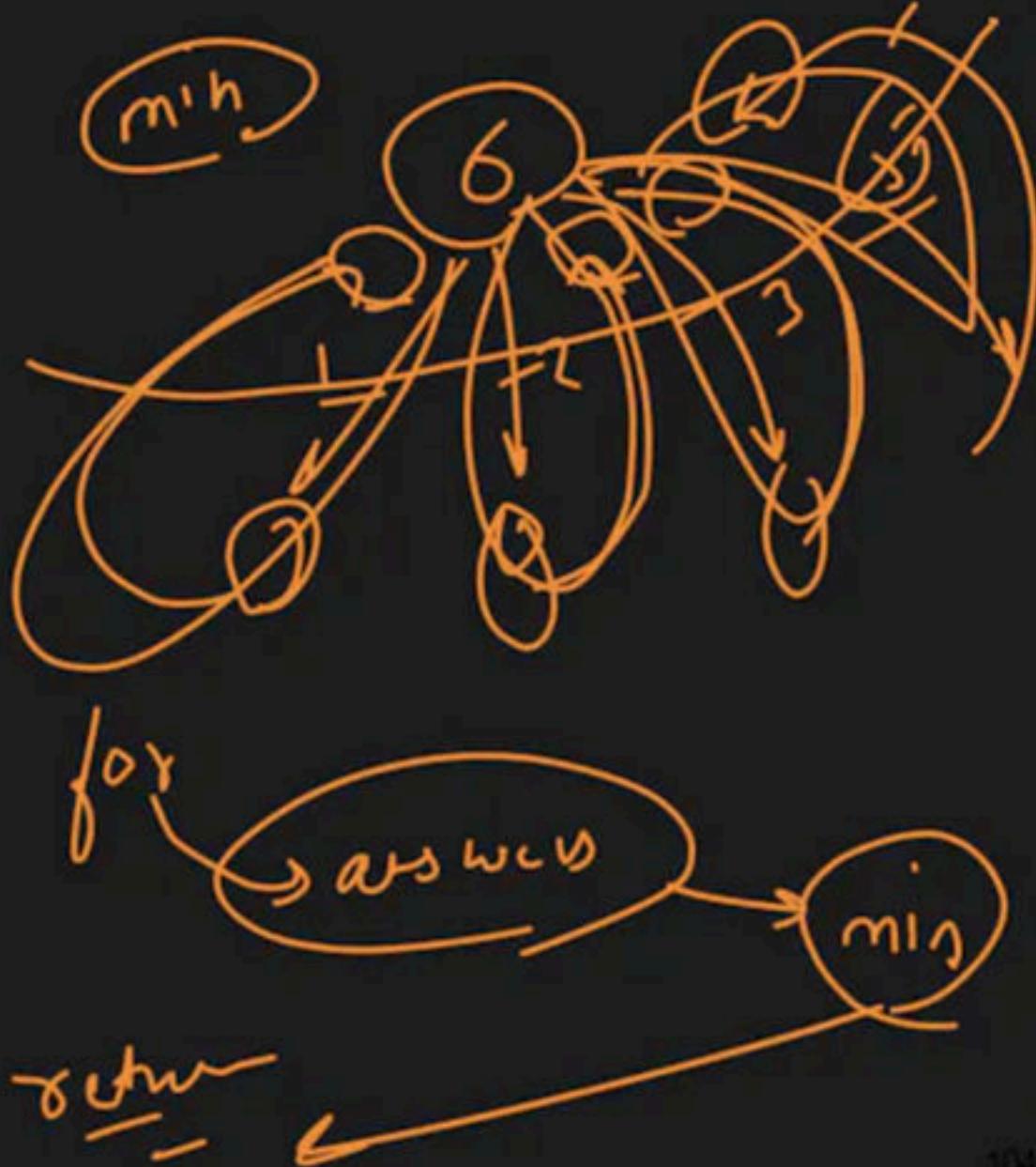
$\text{inclusion} = 3 + 0$
 $- 3$

exclusion 2



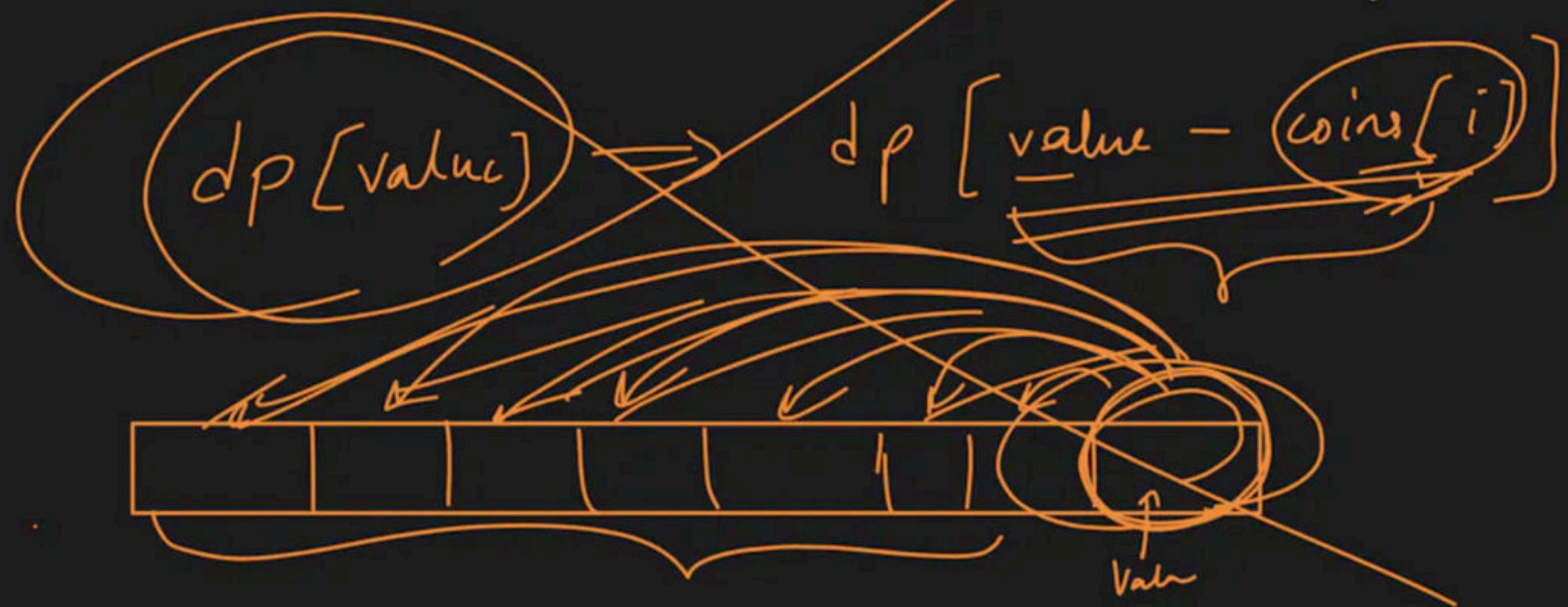


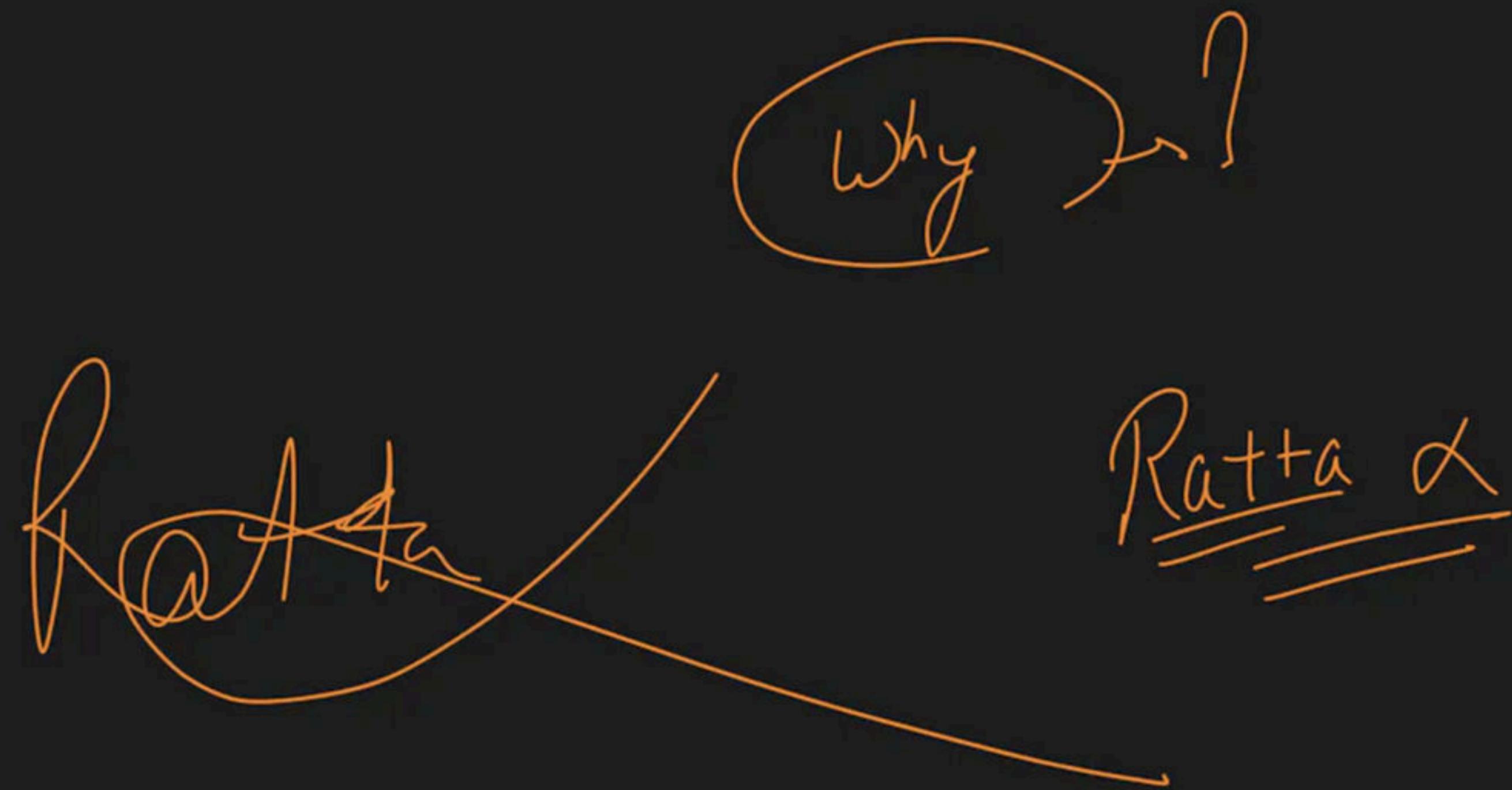
$$\{ \overline{\backslash}, \overline{/}, \overline{\diagup}, \overline{\diagdown} \}$$



$O(\text{amount} * n)$

$O(n^2)$ $\approx O(\sqrt{n})$





$\Rightarrow \Rightarrow \Rightarrow \Rightarrow$

\Rightarrow
1, 2, 5

amount = 11

11 - 0
5 - 5, 1

\Rightarrow
RE



\Rightarrow
if amt = 0:
yell <= 0



Dynamic Programming Class-3

Special class

Painting fence

$n \rightarrow \text{post}$, $K \rightarrow \text{color}$

cond \perp →

no two adjacent
posts have the
same color

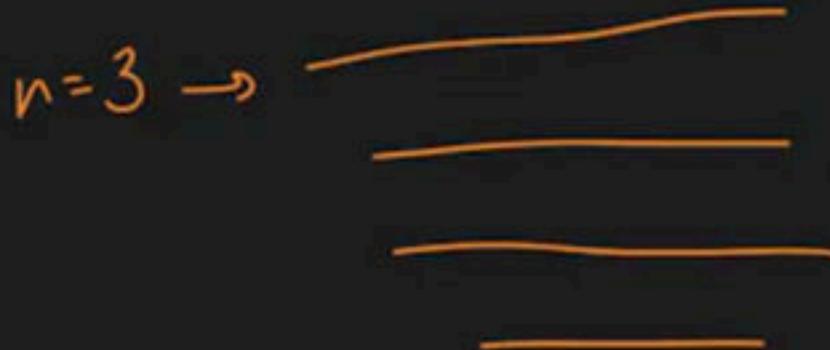
$$n = 1 \\ K = 3 \rightarrow R, G, B$$

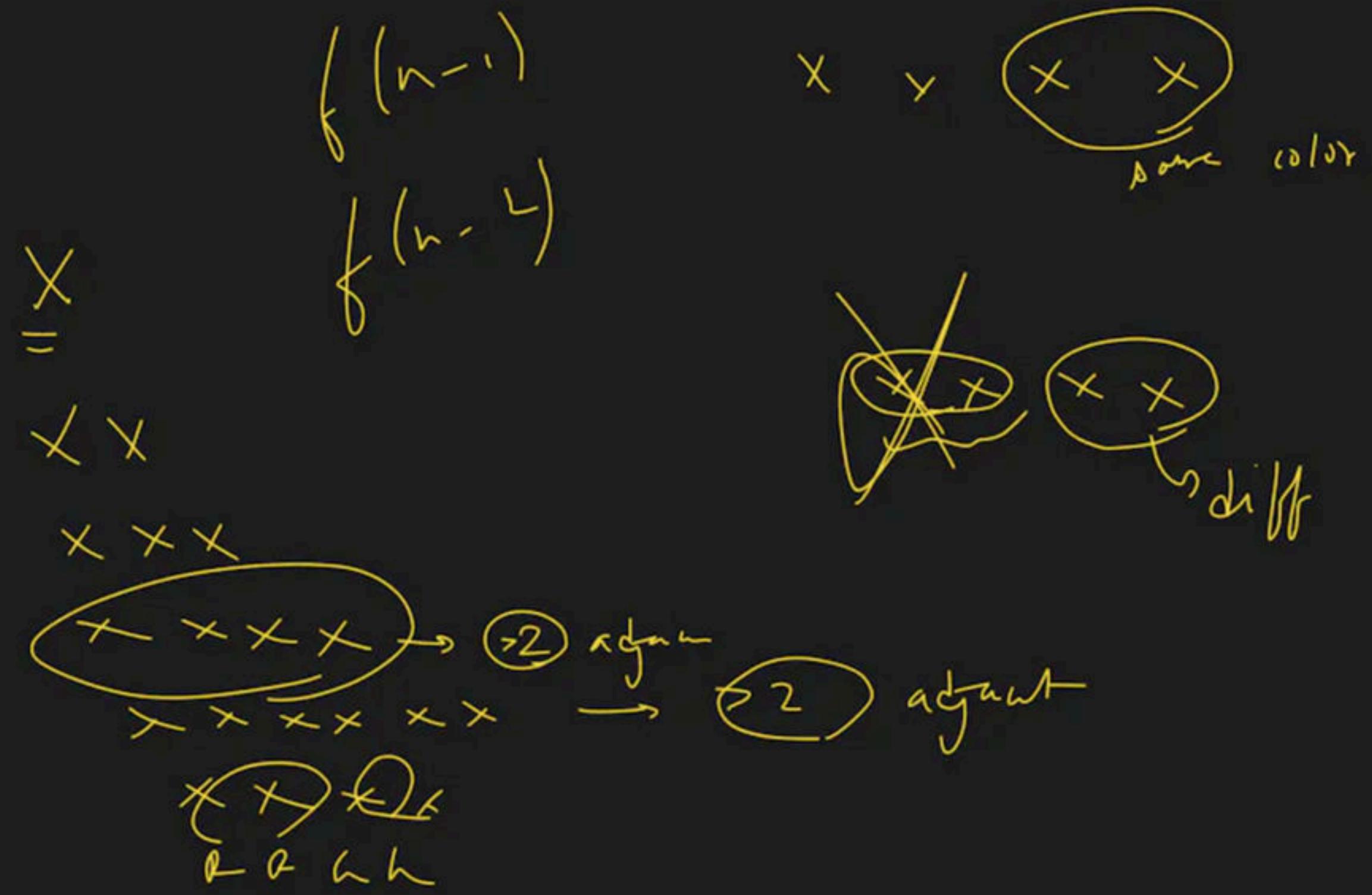


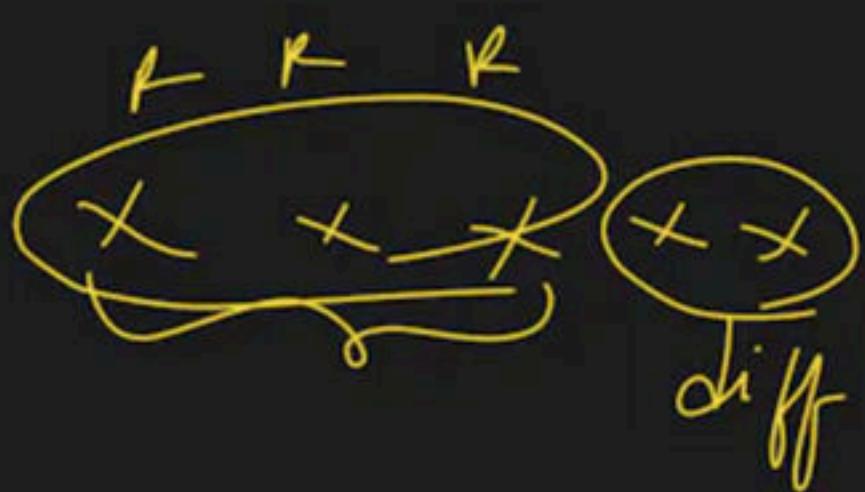
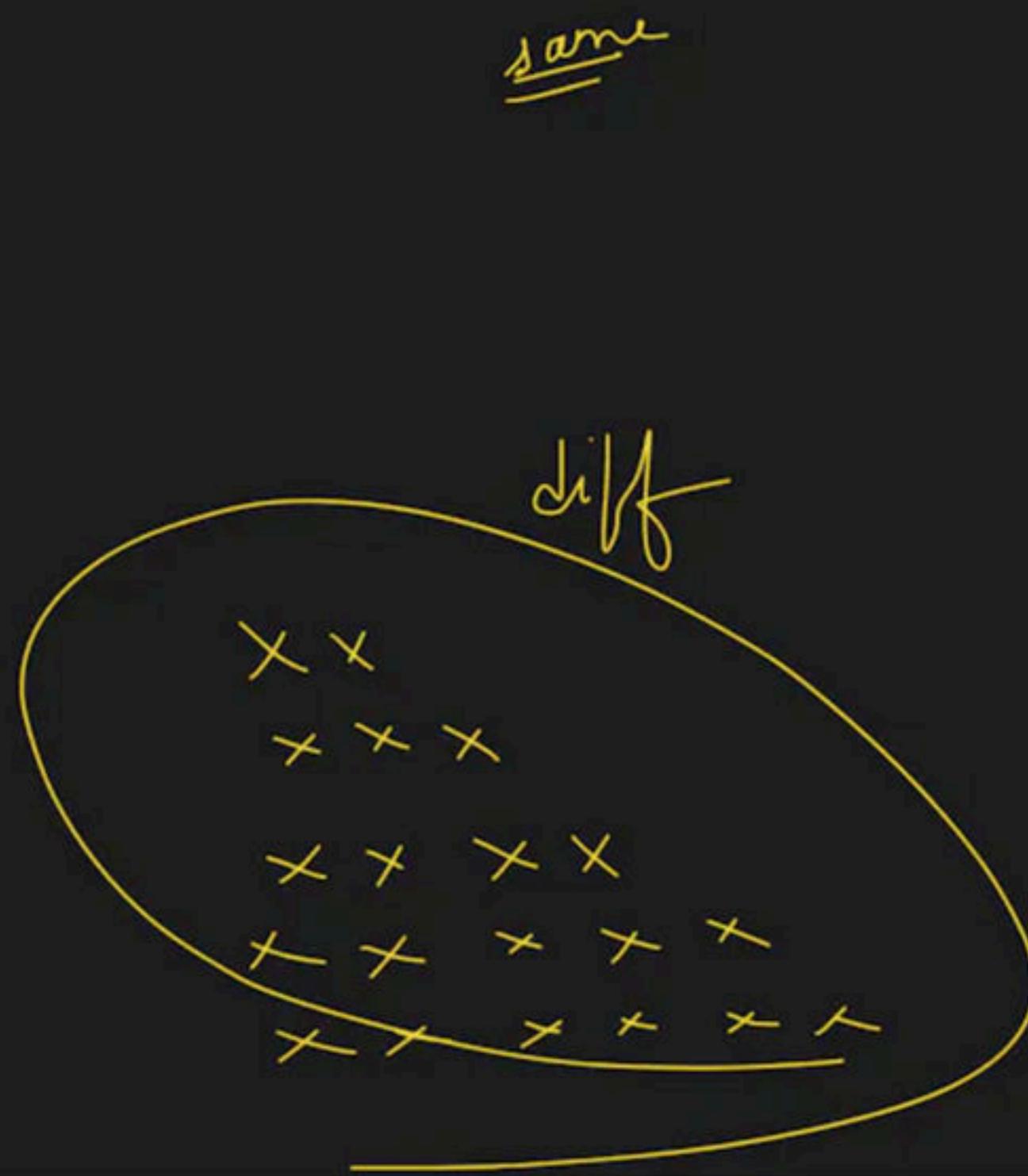
$$n = 2$$

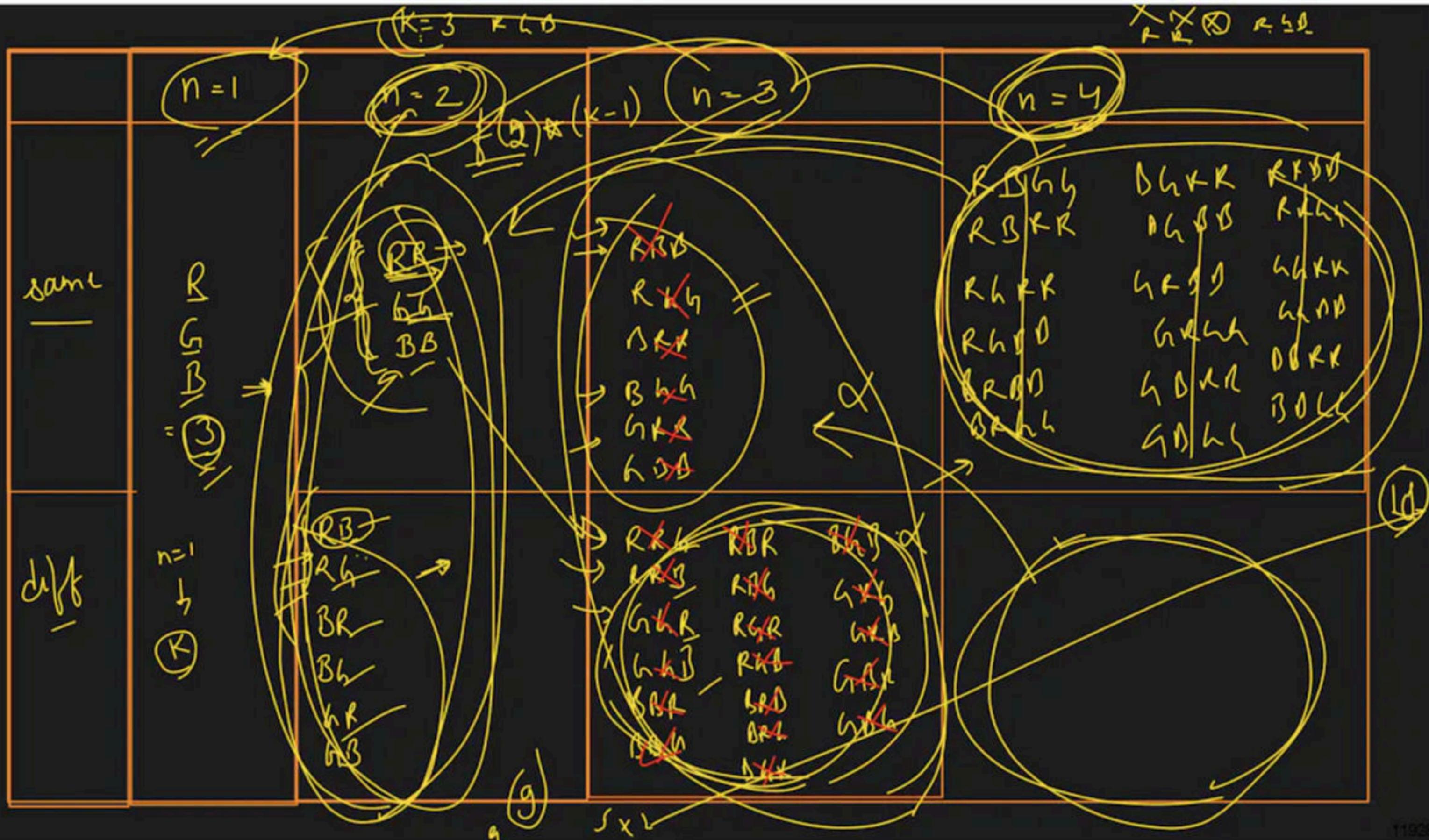
$$\begin{matrix} RR & GG & BB \\ RG & GR & BR \\ RB & GB & BL \end{matrix}$$

$$\rightarrow \text{ans} = 9$$



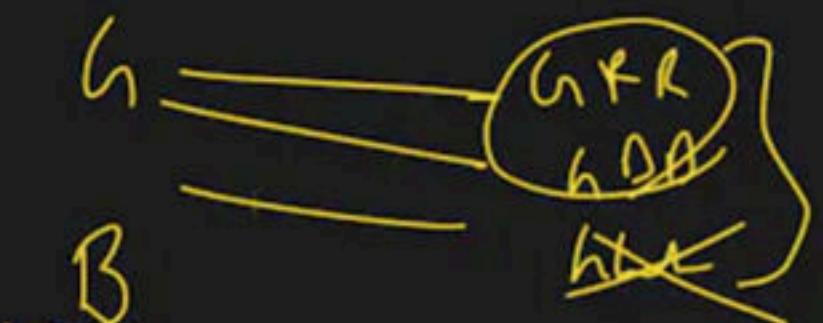
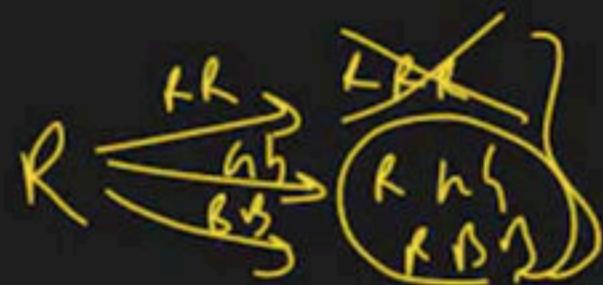


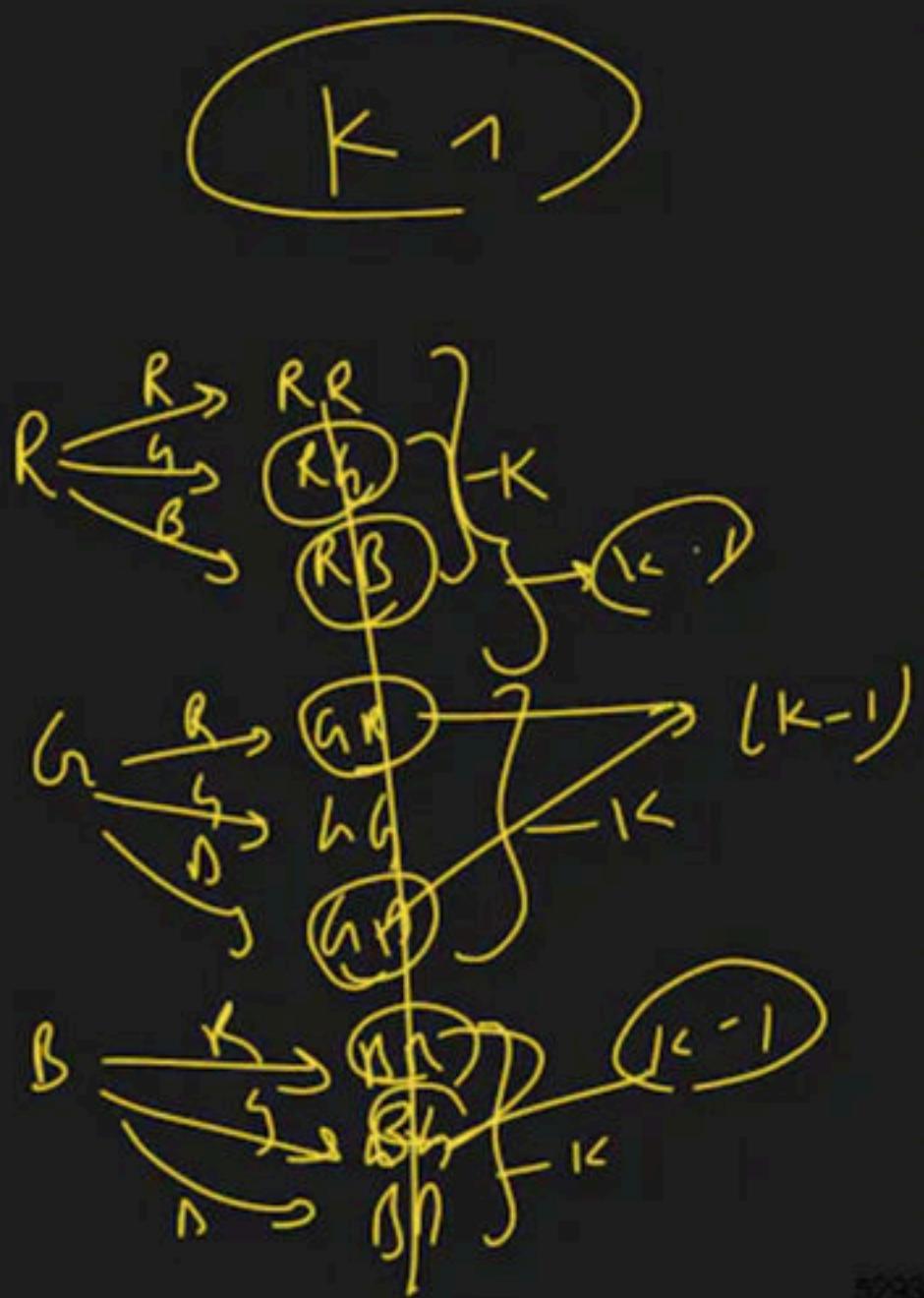
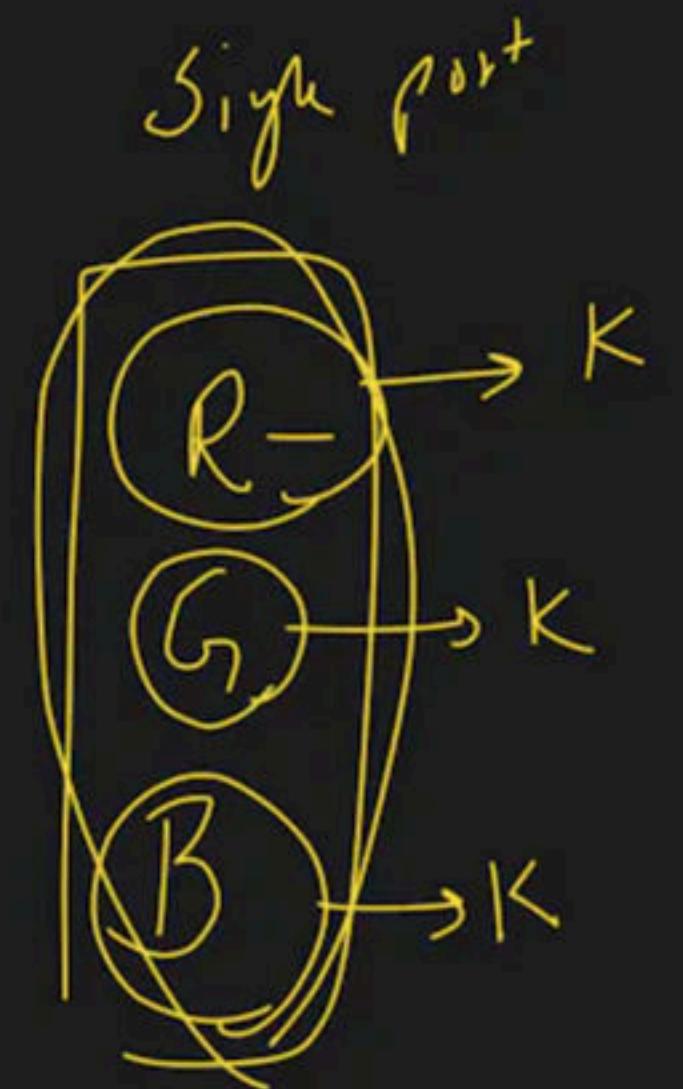




$n=1$

Save ~~XXX~~

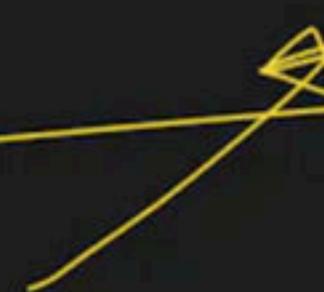




$$f^{(n)} = \text{Same}_n + \underline{\text{diff}_n}$$

$$= f^{(n-2)}(k-1) + f^{(n-1)}(k-1)$$

$$f^{(n)} = (-1) [f^{(n-1)} + f^{(n-2)}]$$



diff

XX
Rh
RD
DK
Dh
LR
hD

same →



$\leftarrow \rightarrow$

$n=1$

$n=2$

δ_{ant}

δ_{eff}

$\text{ans} \in K$

α_1
 α_2
 α_3
 α_4
 α_5
 α_6
 α_7
 α_8

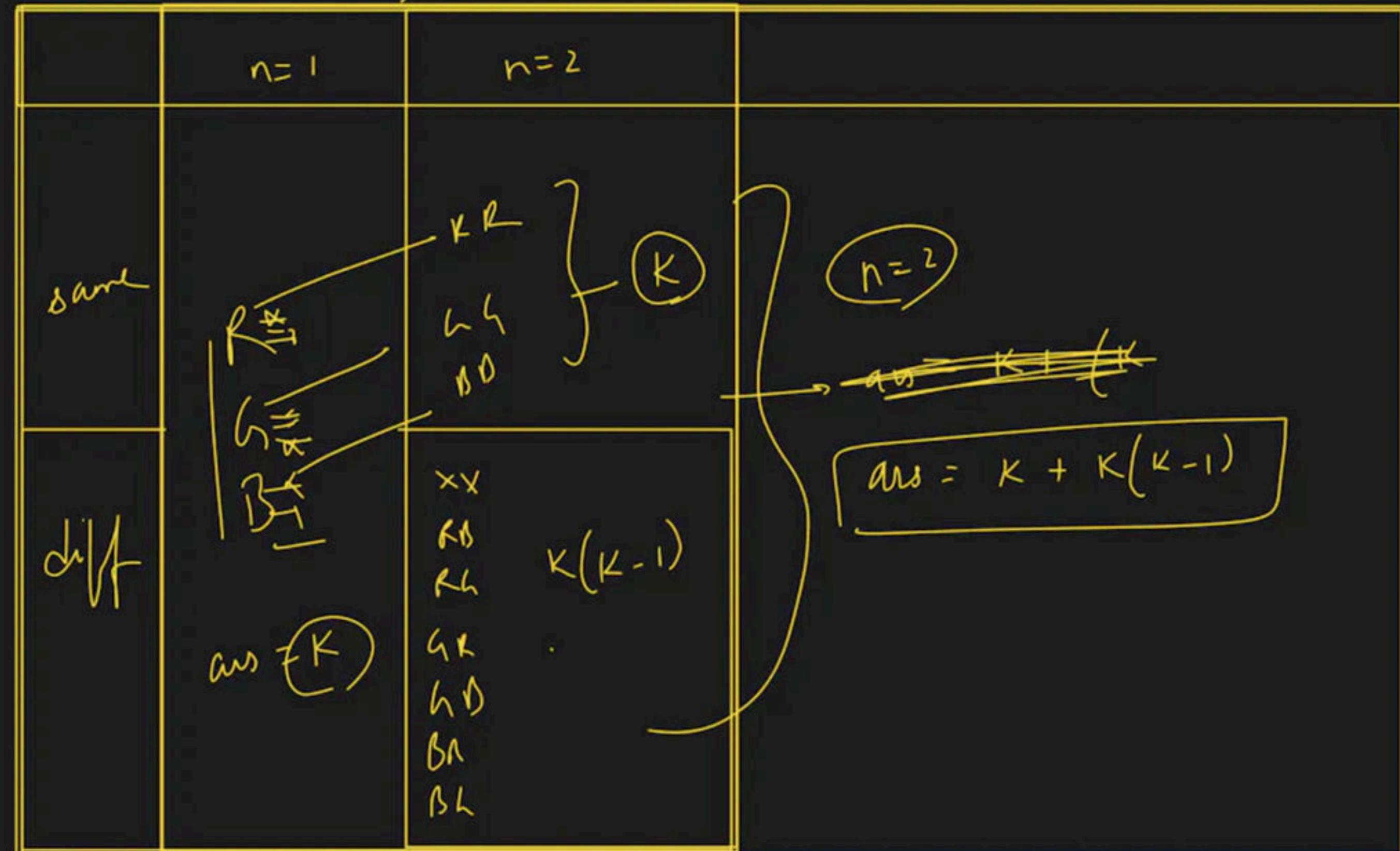
R^*
 R
 β_0
 β_1
 β_2

$K(K-1)$

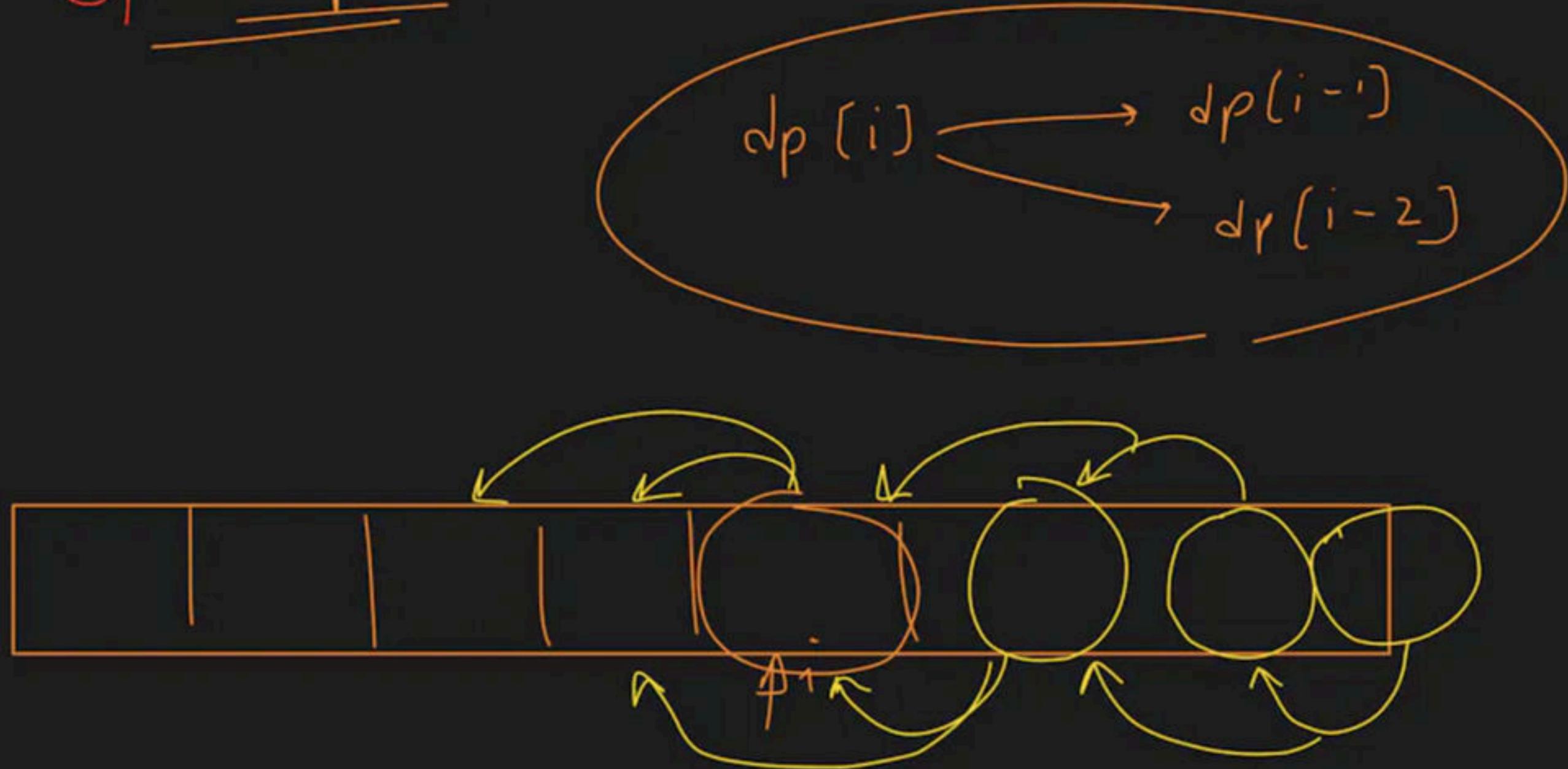
K

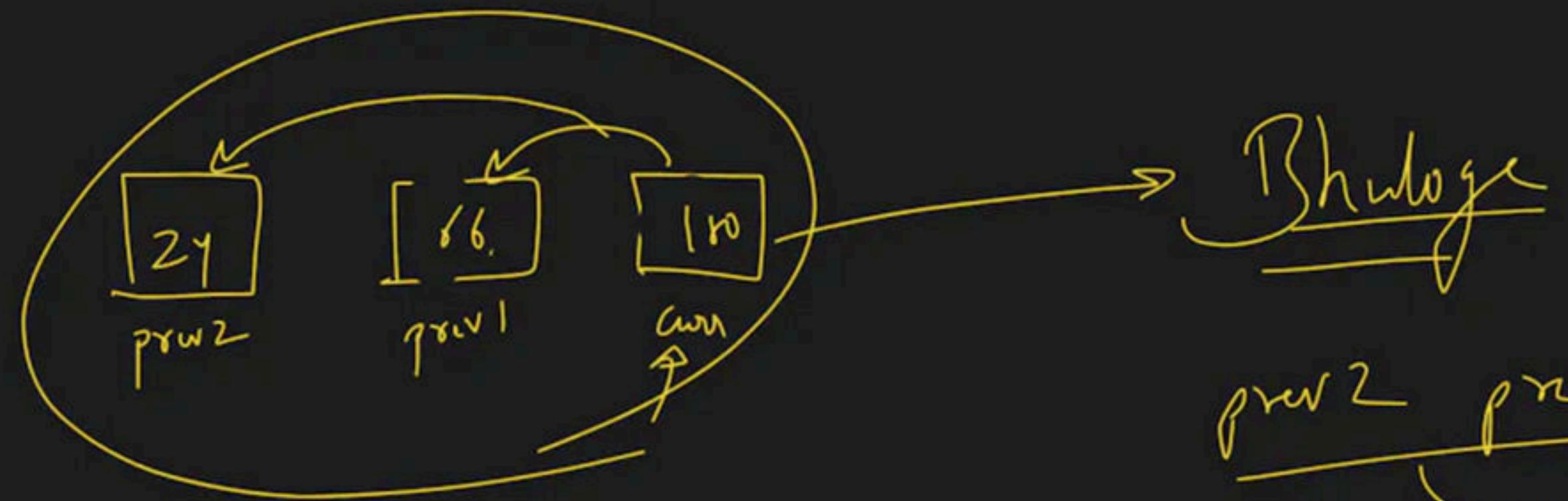
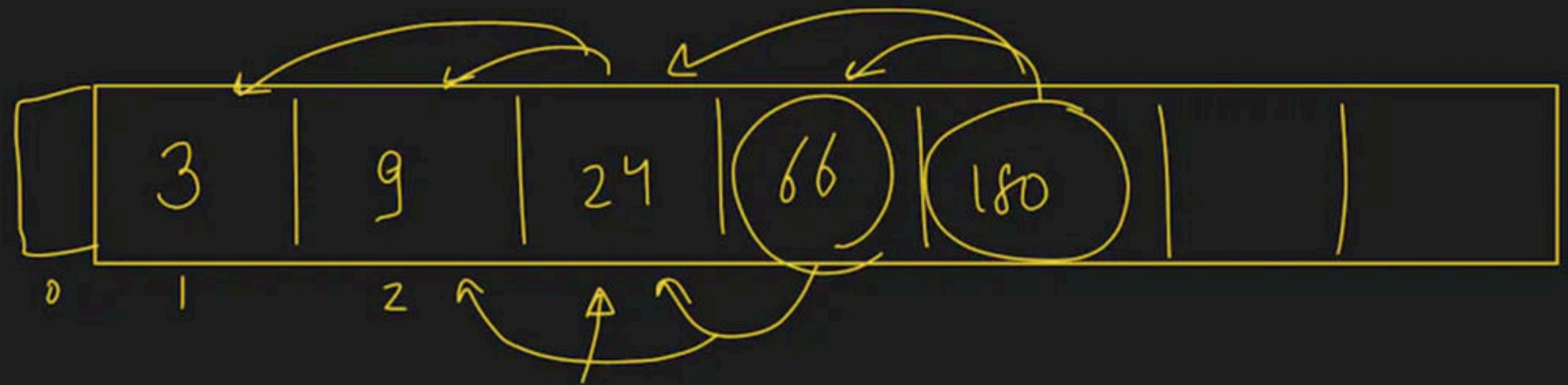
$n=2$

$$\text{ans} = K + K(K-1)$$



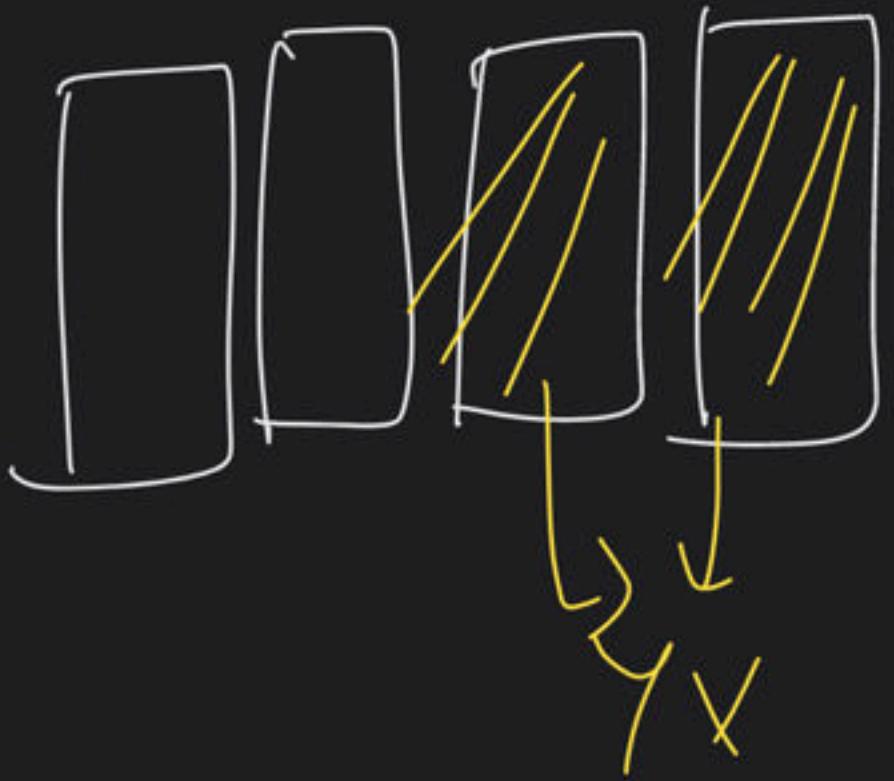
Span Optimisation





prev2 prev
 → shift

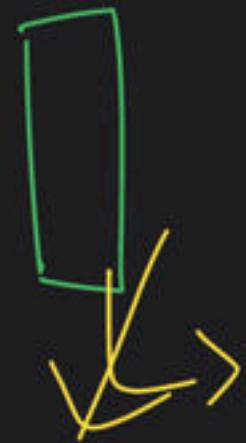
$\Rightarrow Q \Rightarrow n \text{ post}$



\checkmark K colors
Bucket

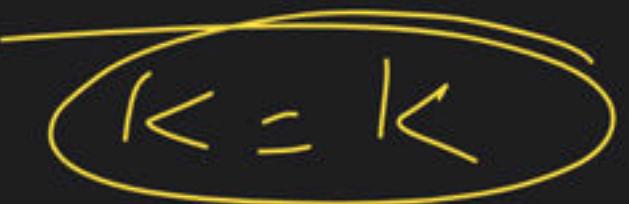
RR , RR
=

$n = 1$



$k = 3$

$k = 1$



$h = 2$

Lattice fun	$n=1$	$n=2$	$n=3$	$n=4$
Same		$\begin{array}{c} RR \\ \text{--} \\ \text{--} \end{array}$ $\xrightarrow{\textcircled{3}} \quad \begin{array}{c} K \\ \equiv \end{array}$	$\begin{array}{c} R \text{---} K \\ \text{B} \text{---} RR \\ \text{R} \text{---} \text{---} \\ \text{B} \text{---} \text{---} \\ \text{R} \text{---} \text{B} \text{---} \\ \text{B} \text{---} \text{B} \end{array}$ $\xrightarrow{\textcircled{6}} \quad K \cdot (K^{-1})$ $f(n-1) * (f-1)$	
diff	K	$\begin{array}{c} RS \\ \text{---} \\ \text{R} \text{---} \text{B} \\ \text{L} \text{---} \text{R} \\ \text{B} \text{---} \text{B} \\ \text{B} \text{---} \text{R} \\ \text{B} \text{---} \text{B} \end{array}$ $\xrightarrow{\textcircled{6}} \quad \begin{array}{c} \leftarrow \\ 3 \cdot (3^{-1}) \end{array}$	$\xrightarrow{\textcircled{18}}$ $f(2) * (1 \cdot K^{-1})$ $f(n-1) * (K^{-1})$	

$$K + K(1 \cdot K^{-1}) \Downarrow \textcircled{9}$$

$\cap \equiv \exists$

(C)

GRR
BRR
RGR
BGG
RBG
GBB

Same

$X X X_j$

diff.

$\begin{array}{l} RR \rightarrow (k-1) \\ GG \rightarrow k \\ BB \rightarrow (k-1) \\ RG \rightarrow R, B \rightarrow (k-1) \times (k-1) \\ RB \rightarrow R, G \rightarrow (k-1) \times (k-1) \\ GR \rightarrow \\ GB \rightarrow \\ BR \rightarrow \\ BG \rightarrow \end{array}$

$\begin{array}{c} X \\ \nearrow \\ R \\ \text{---} \\ B \\ \searrow \end{array}$

\equiv

$q * (3^{-1})$

$$\Rightarrow q \times 2 = 18$$

$$f(n) = f(n-1) * (k-1) + f(n-2) * (k-1)$$

diff

same

0/1

Knap sack

Problem

Level - 1

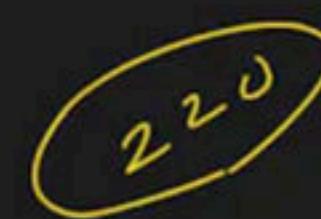
Beginner

capacity

W



$n = 3$

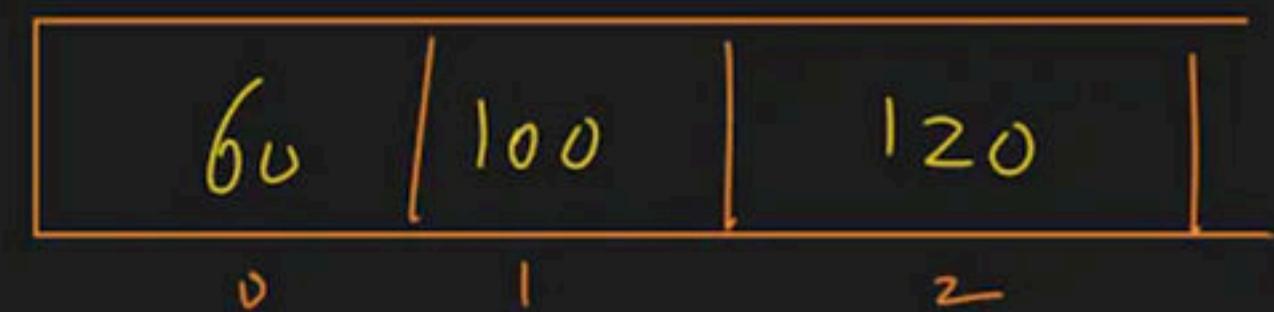


220

$w_{ij} \leq$



Value / Profit



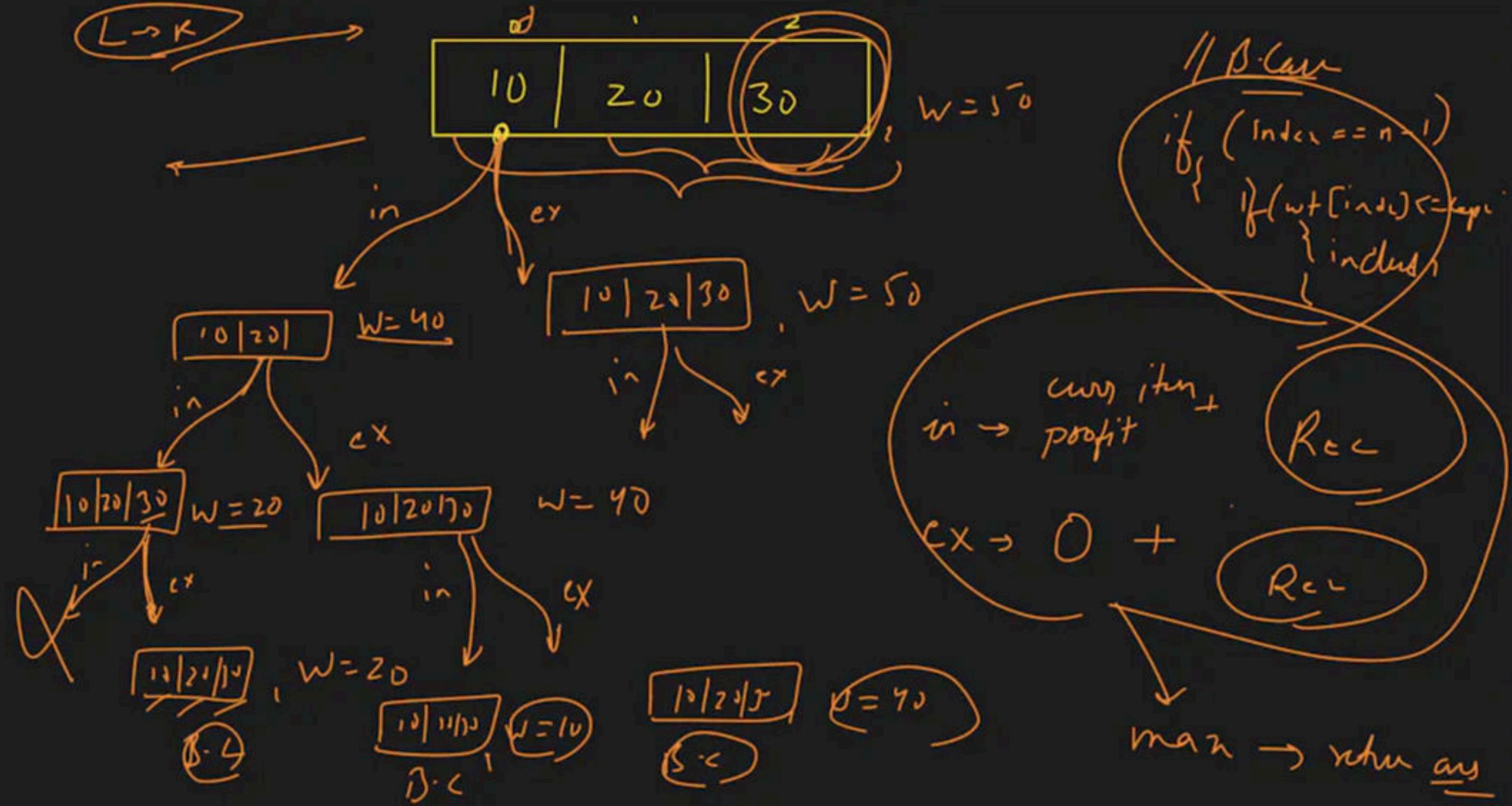
Level = 2

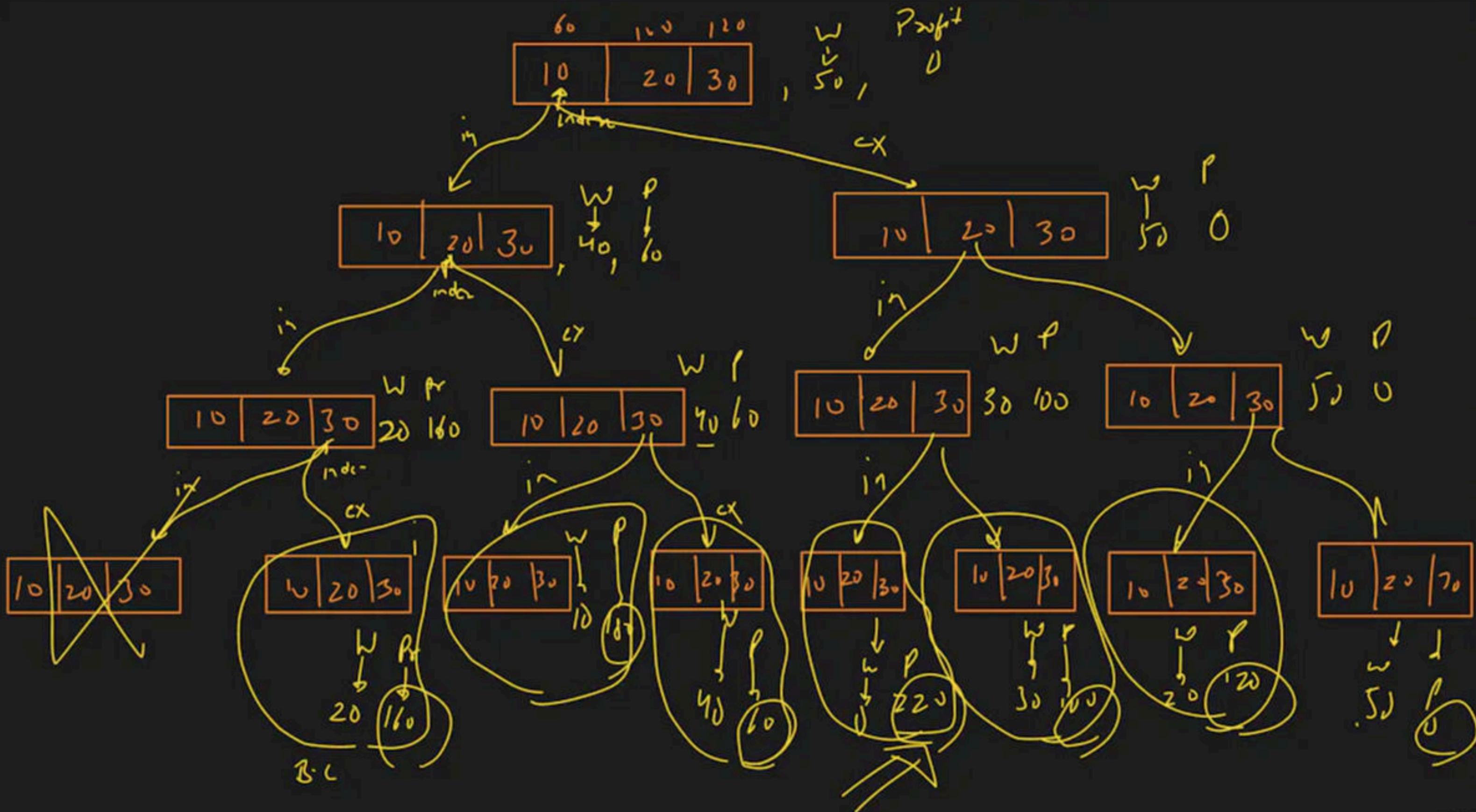
Unbounded knap



$W = 50$

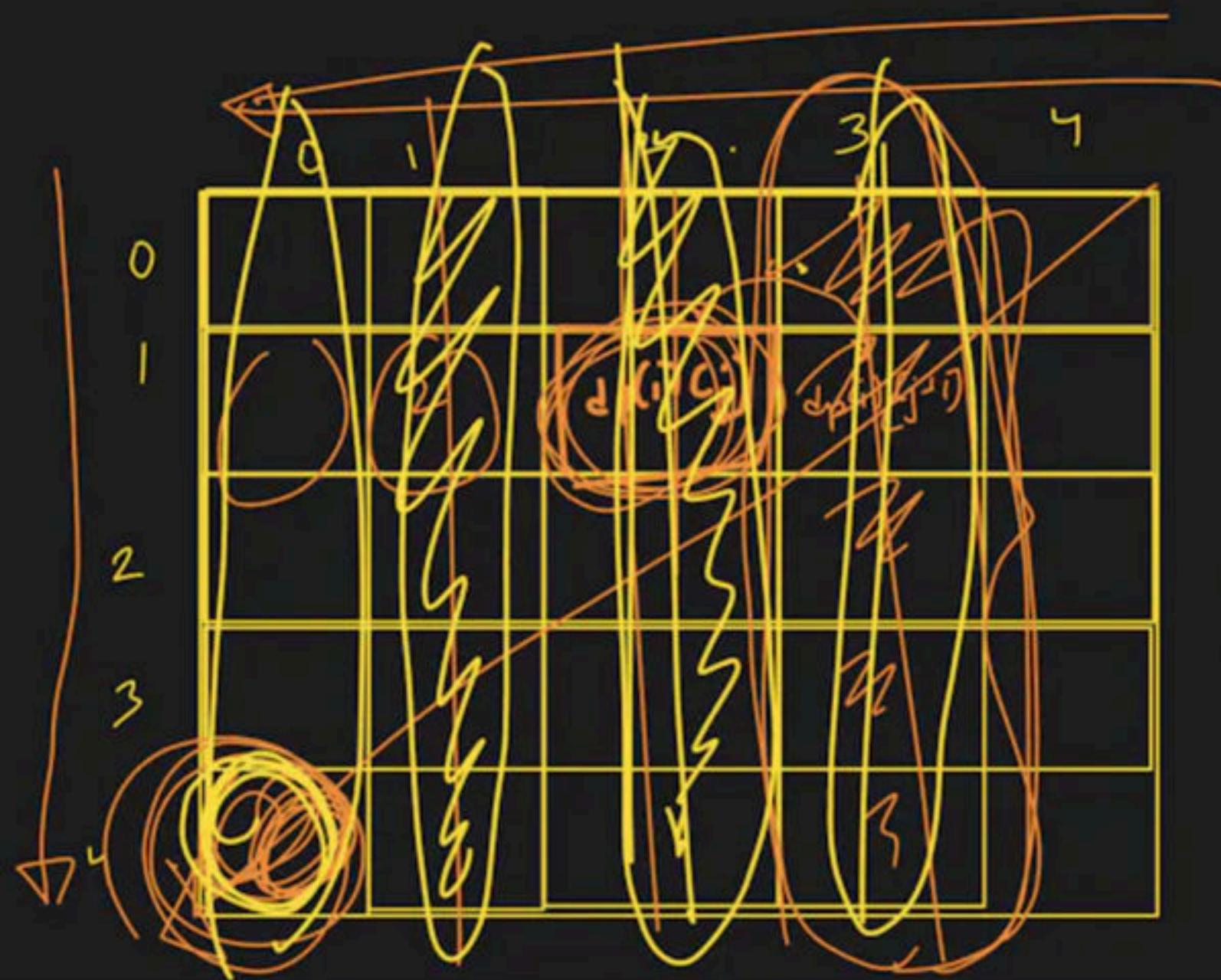
maximum profit





$i = 1$ $j = 2$

$f^{\circ \sqrt{}}$

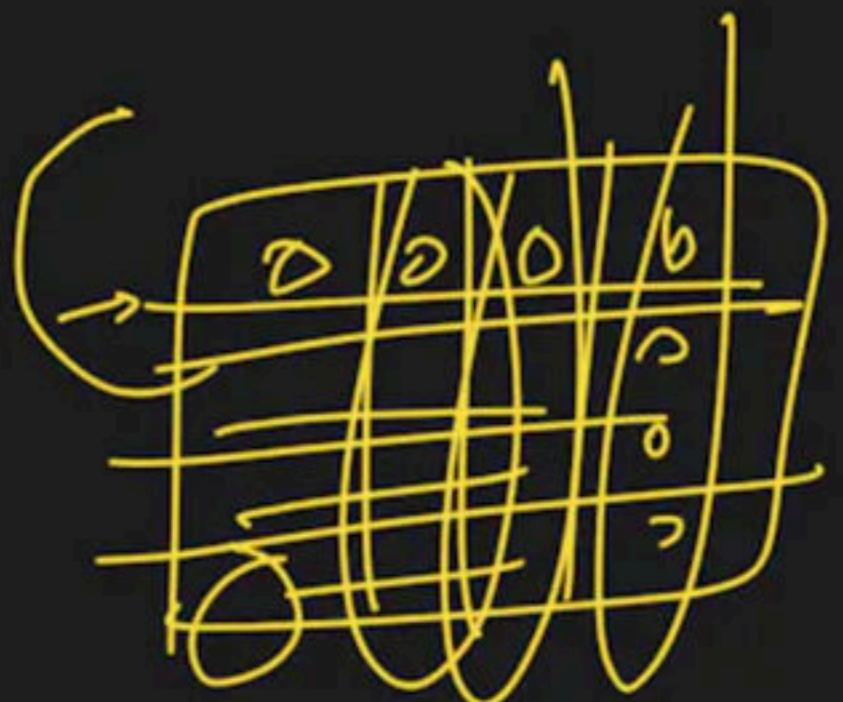


$d\rho(\cdot | j)$

$\overline{d\rho(i | j+1)}$

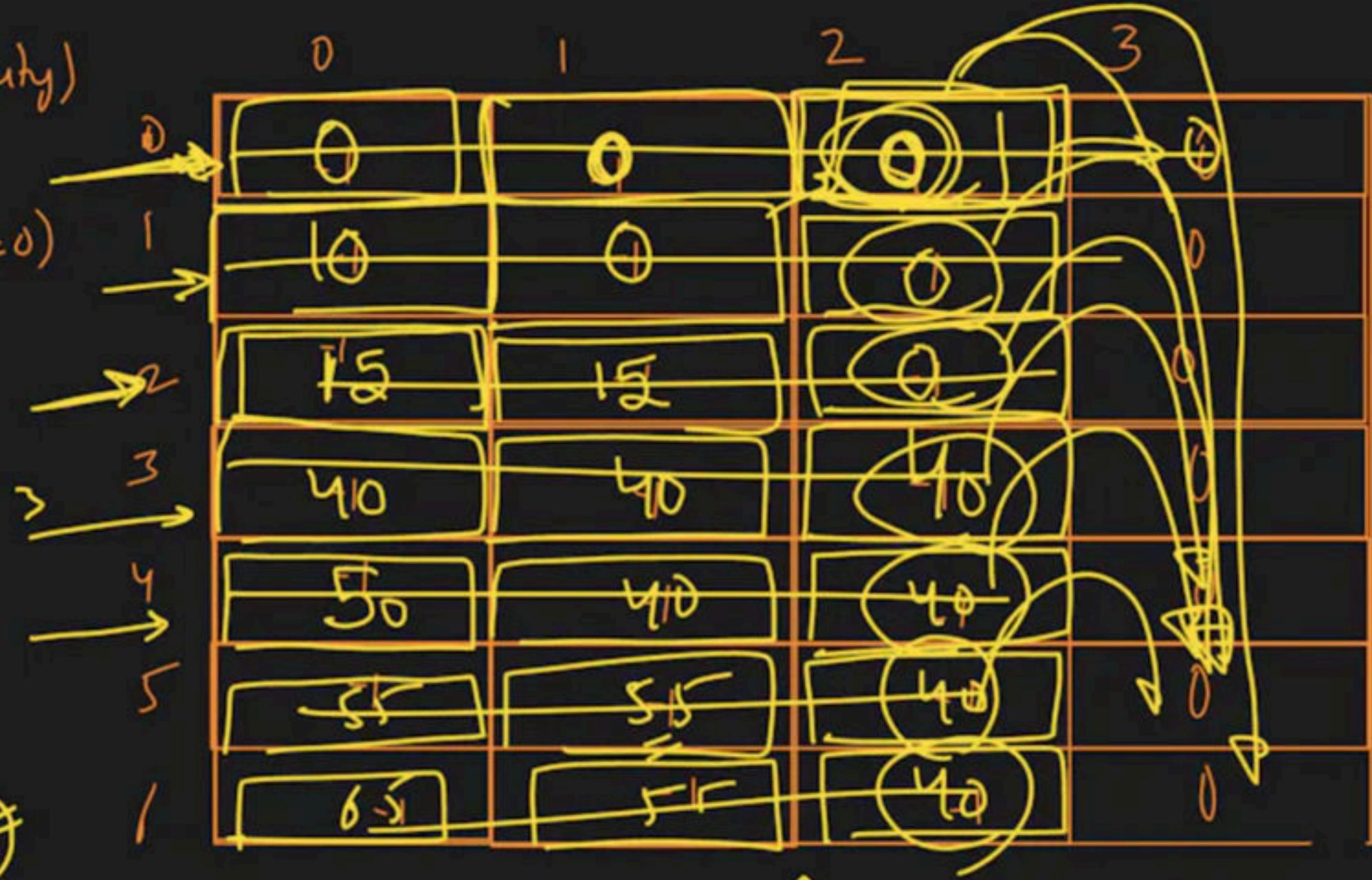
$d\rho(i - w + | j)$

Koi b4
rou h0
siki h 10
nct



$f^v(0 \rightarrow \leq \text{Capacity})$

$\{ f^r(n-1 \rightarrow >= 0)$



Capacity $\rightarrow 0$

wt $\rightarrow 1, 2, 3, 4, 5, 6$

val $\rightarrow 10, 15, 40, 50, 55, 65$

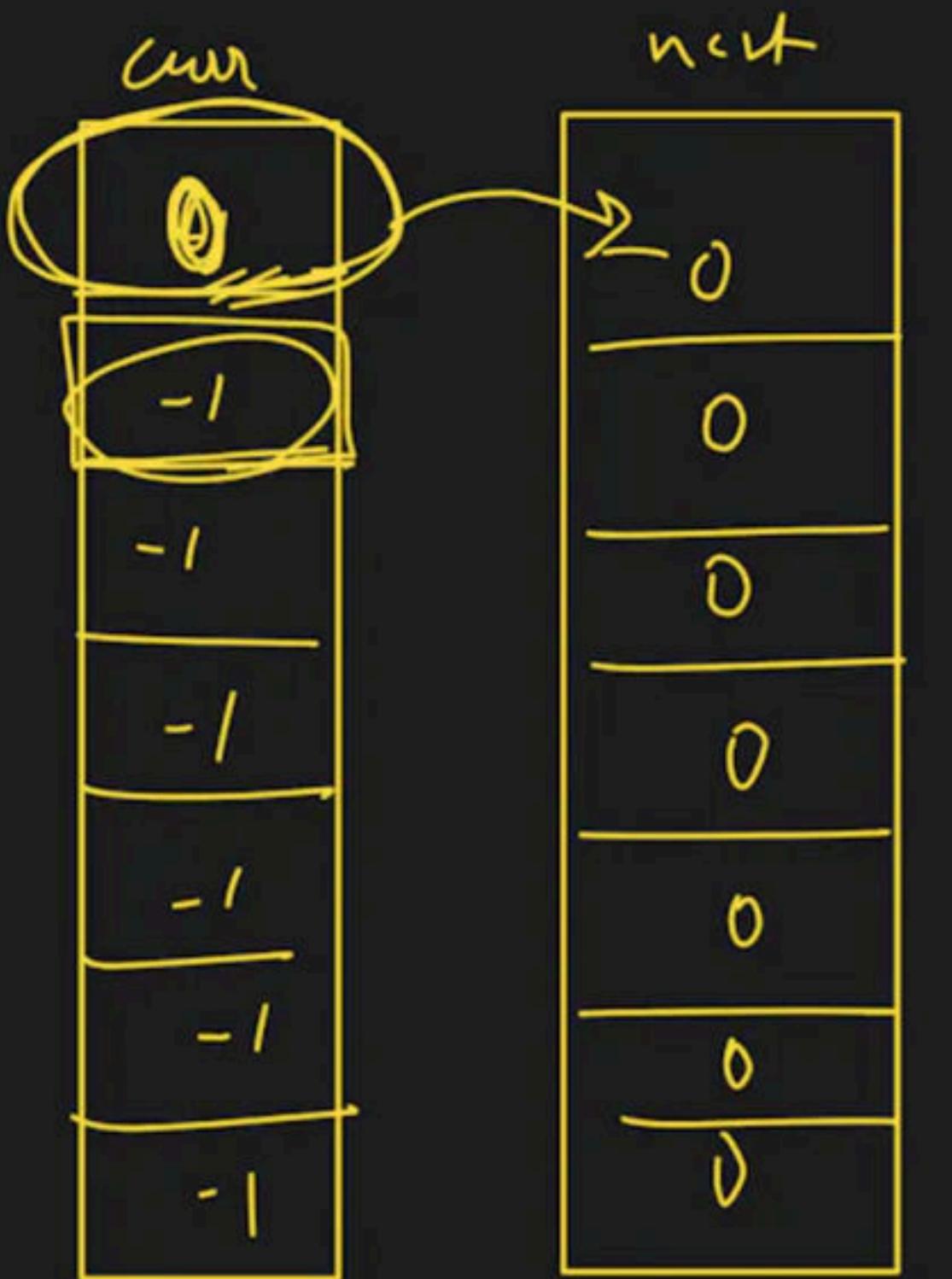
$n=3$

inc = 0
ex = 0+

Capacity = 0

Capacity = 0
induct = 2

$n=3$
 $n-1=2$



~~insert~~
0 -

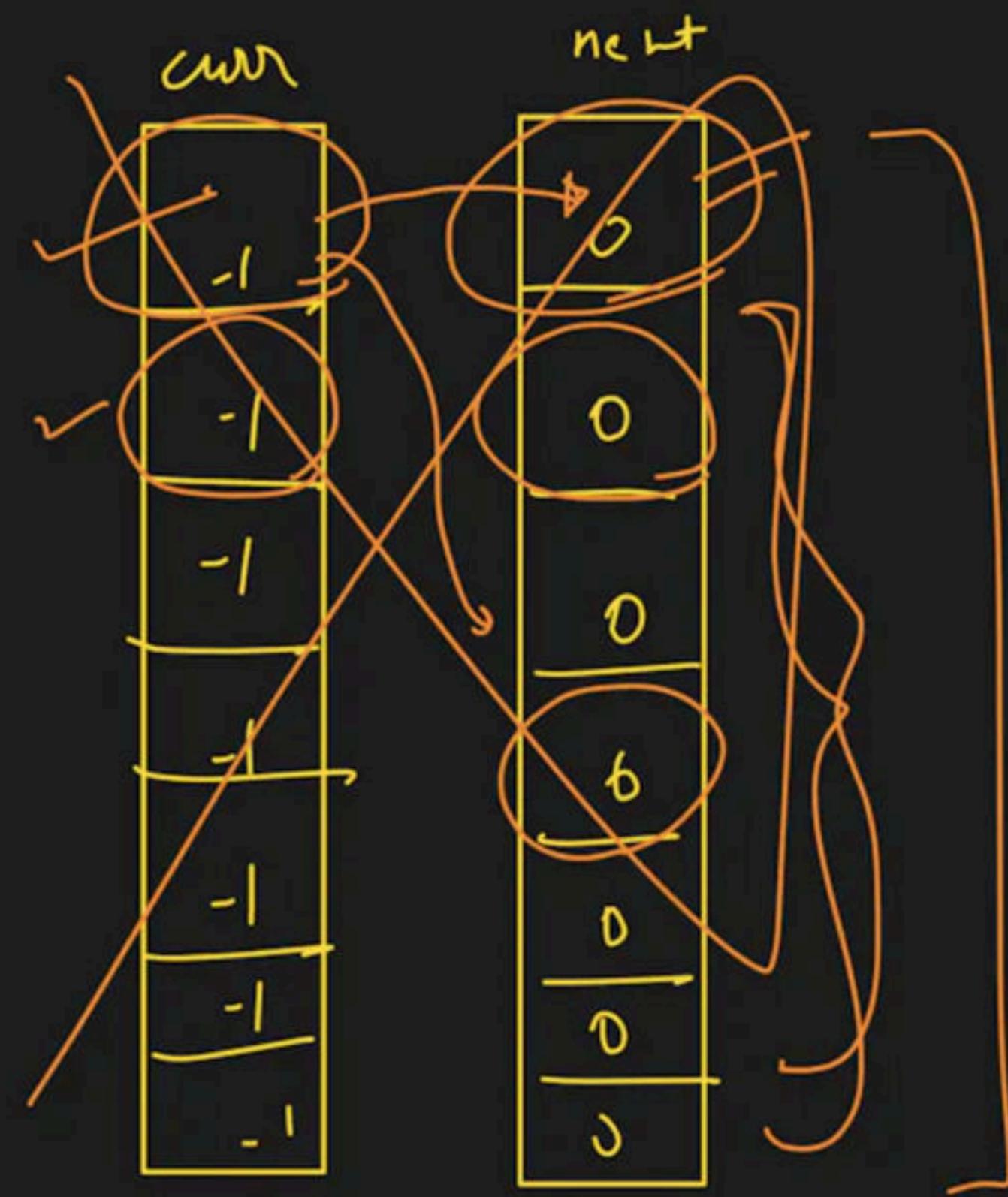
$wt = 2 \rightarrow inc 0$
 $capa = 0 \rightarrow capa + 1$
 $c \leftarrow 0$

$$dp[i][j] \rightarrow dp[i][j+1]$$

capacity $\rightarrow 0$
 $j \rightarrow n-1 \rightarrow 0$

4	2	13
↑	↑	↑

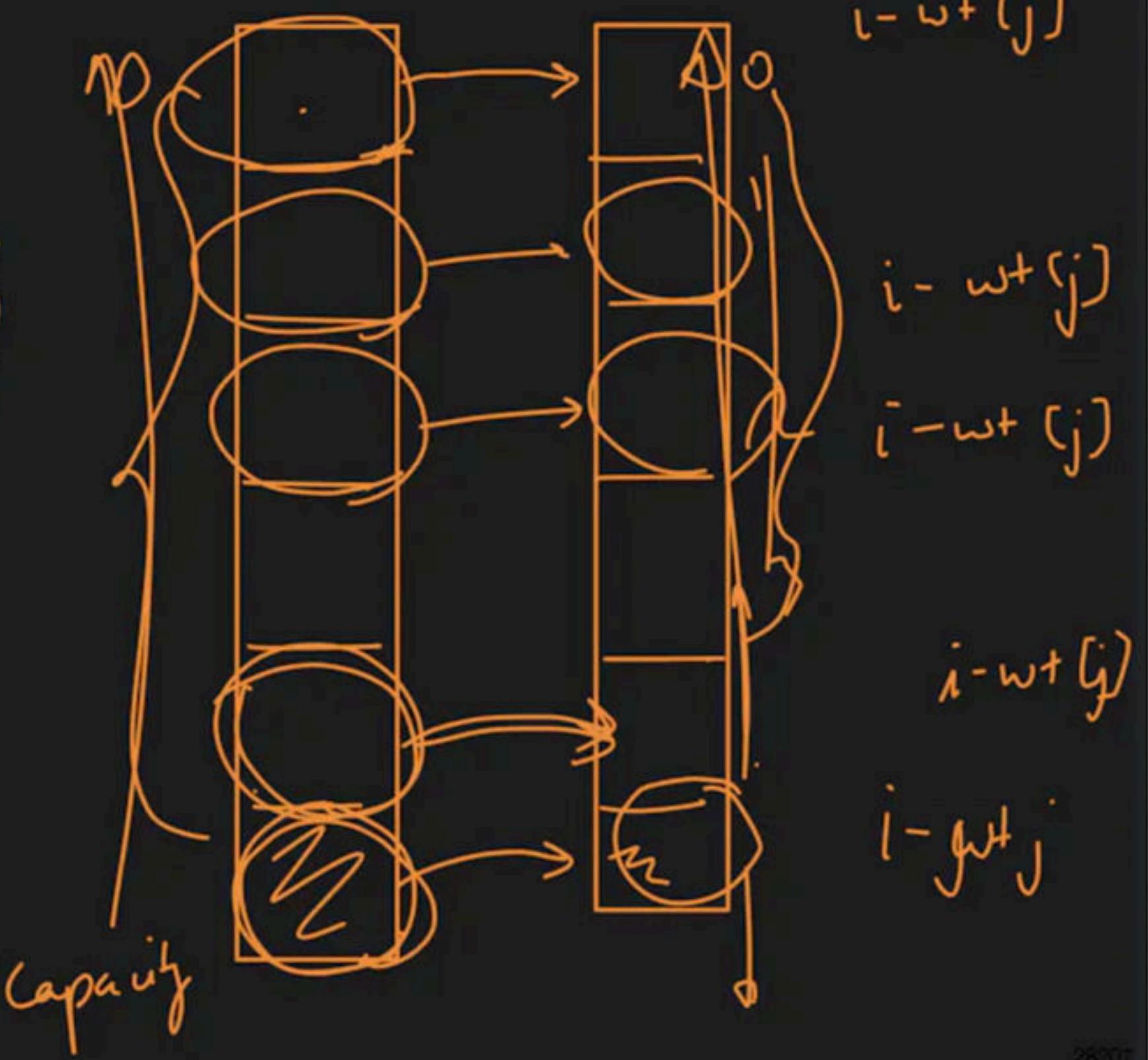
curr	next
0	0
-1	0
-1	0
-1	0
-1	0
-1	0
-1	0
-1	0

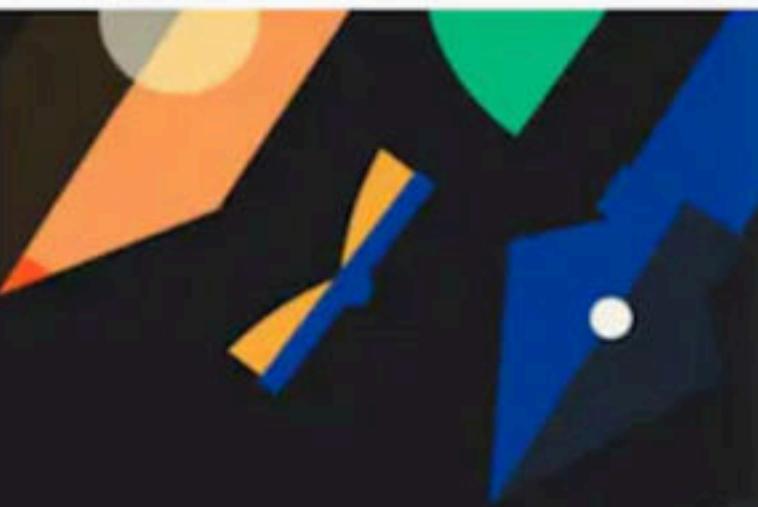


```

for (let i = n - 1; i >= 0; i--) {
    for (let j = 0; j < capacity; j++) {
        if (i - wt[j] >= 0) {
            dp[i][j] = max(dp[i][j], dp[i - wt[j]][j] + v[j]);
        }
    }
}

```





Dynamic Programming Class-4

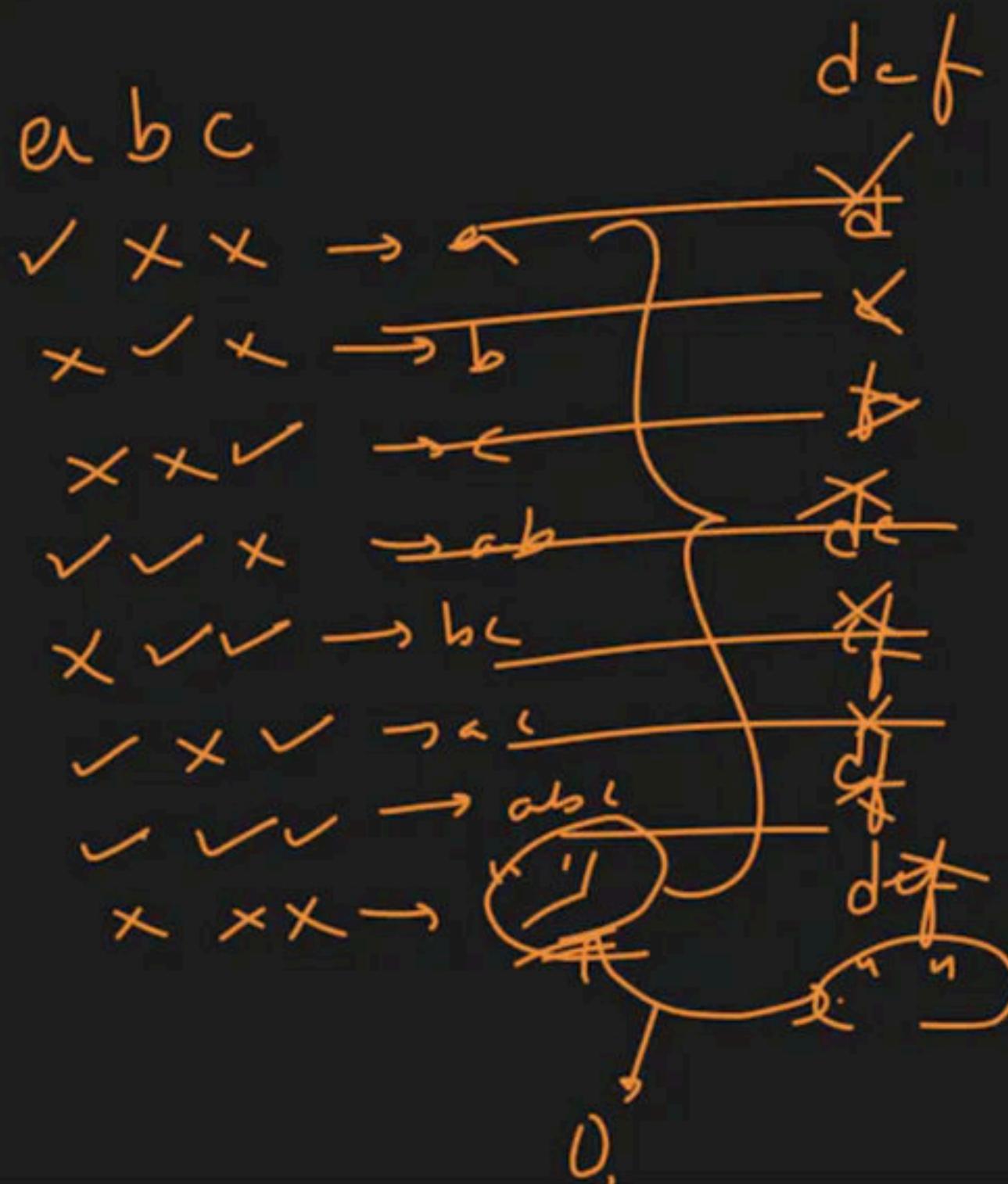
Special class

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→ Longest common Subsequence :-

$S_1 = abc$

$S_2 = def$



$\text{str} \rightarrow abcd$

" "

a

b

c

d

ab

bc

cd

ac

ad

bd

(ab)
bcd

abd

abcd

$\text{str}^2 \rightarrow abc$

a

b

c

ab

bc

ac

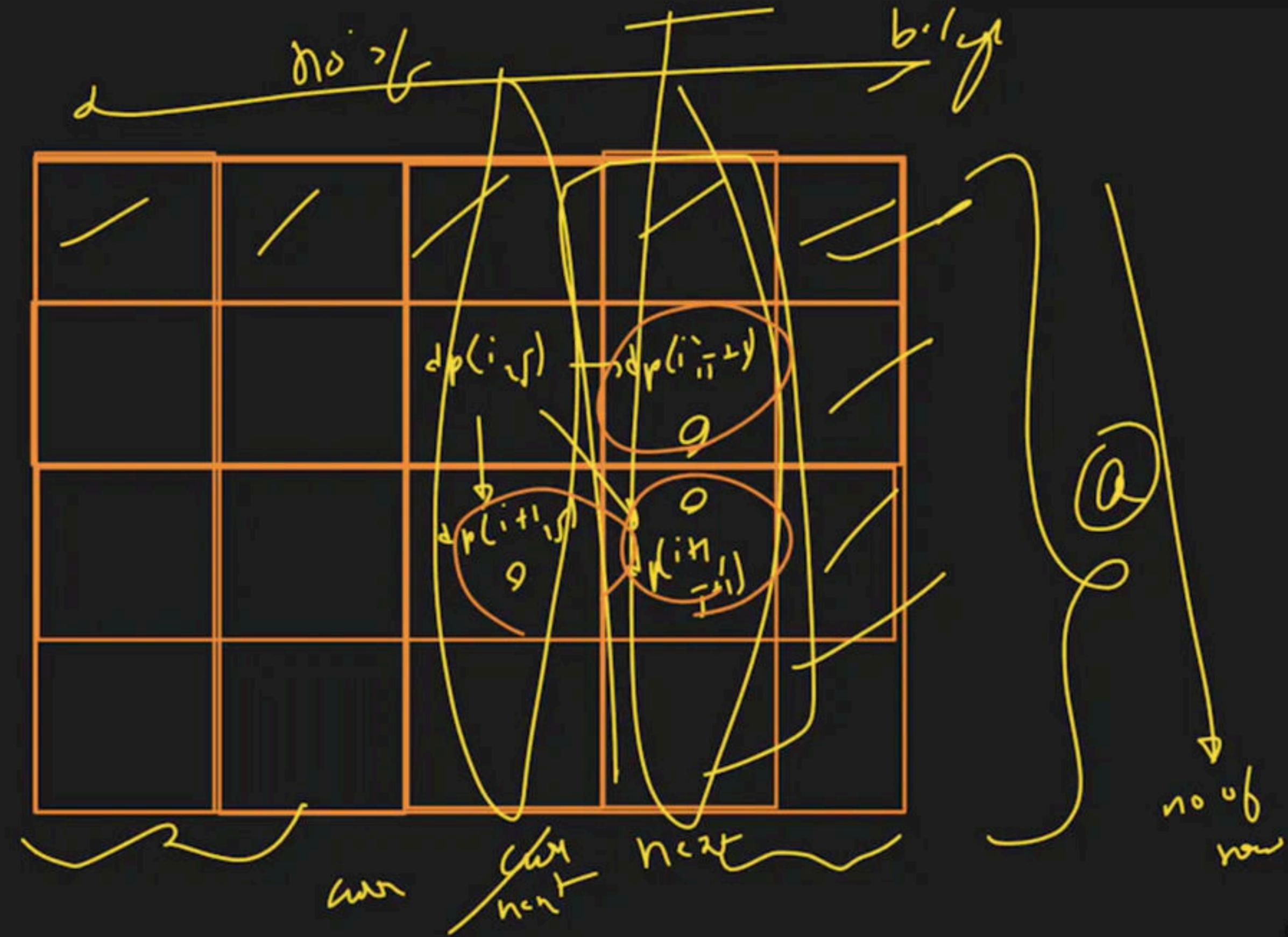
abl

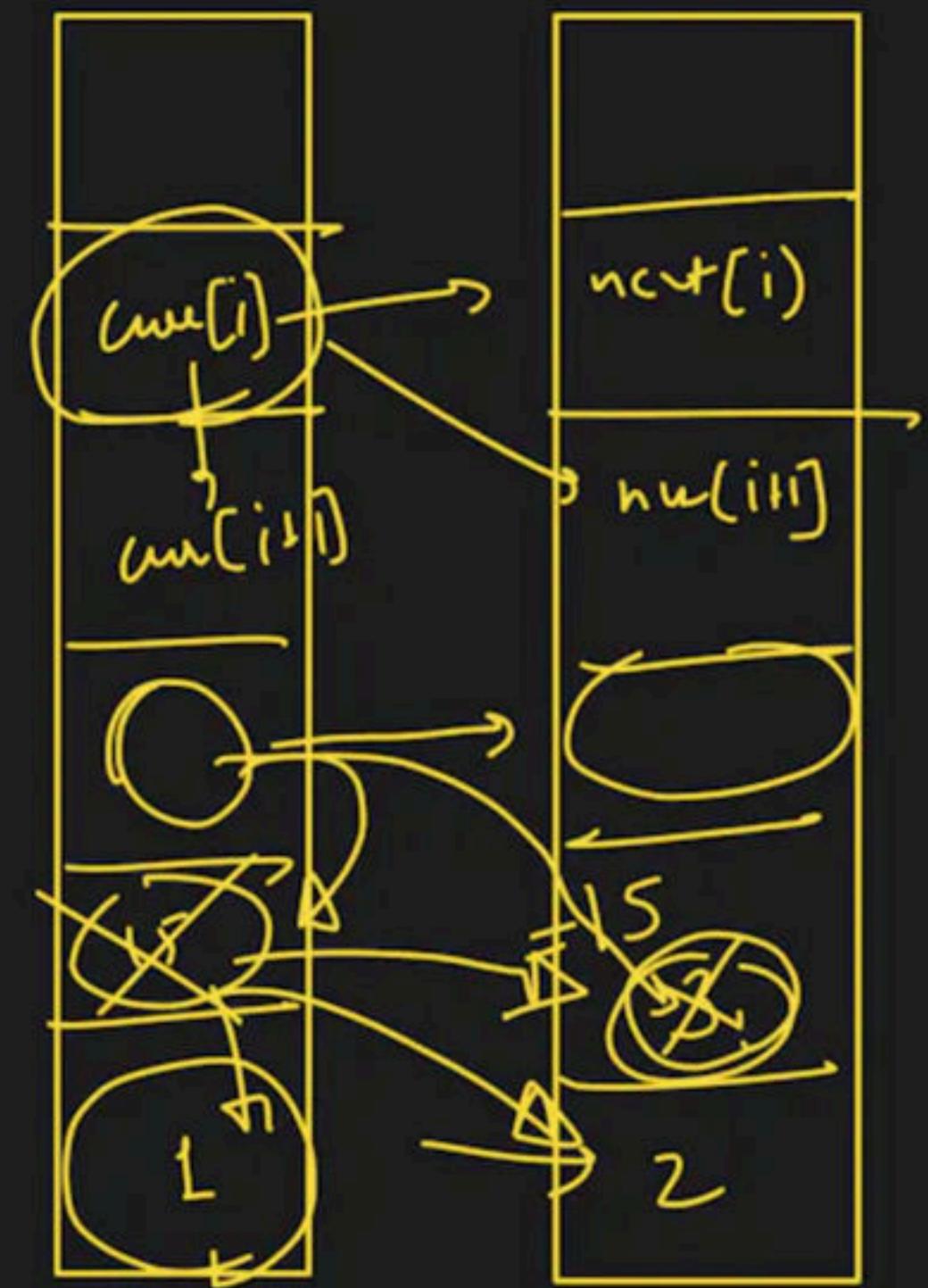
" "

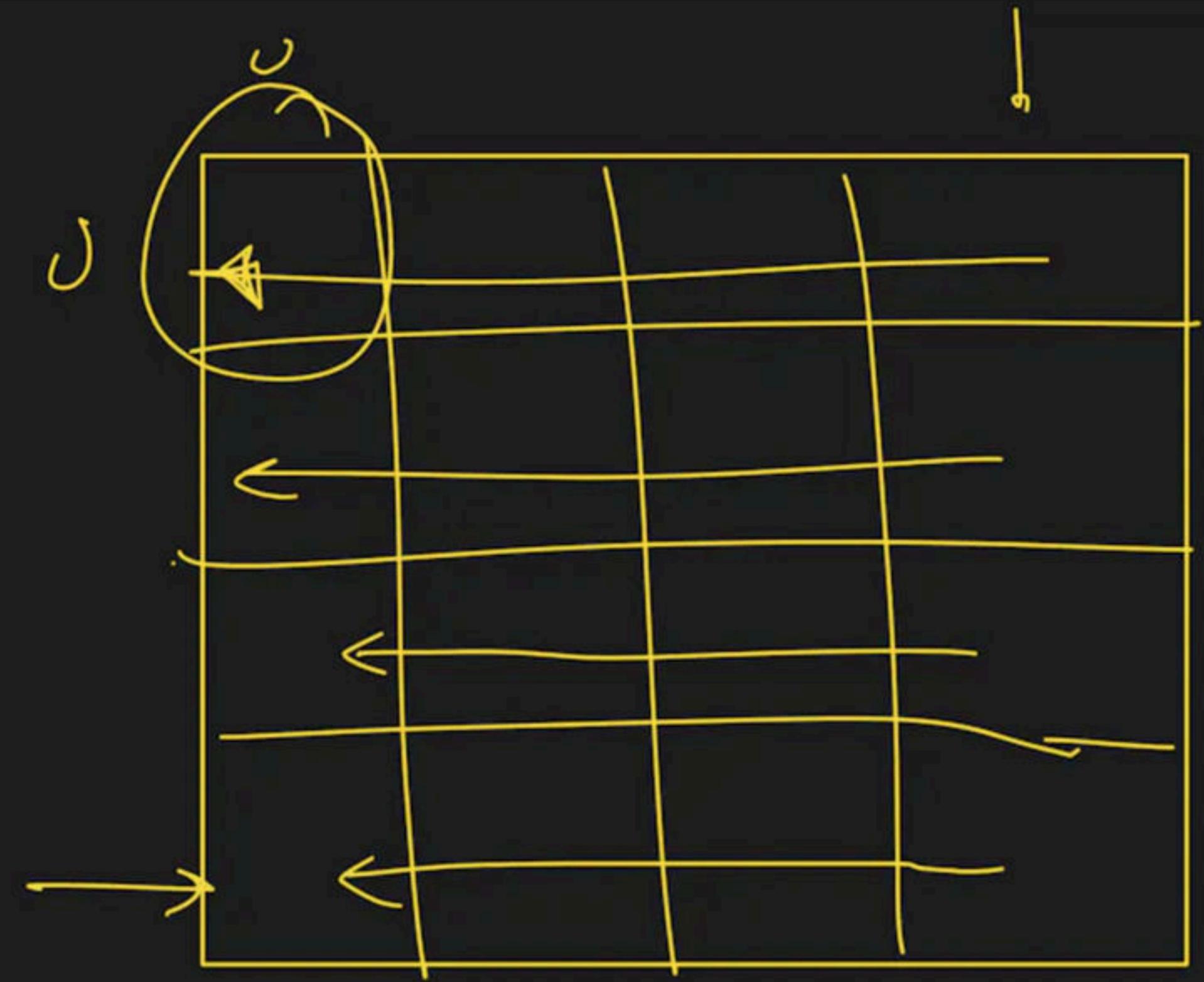
(3)

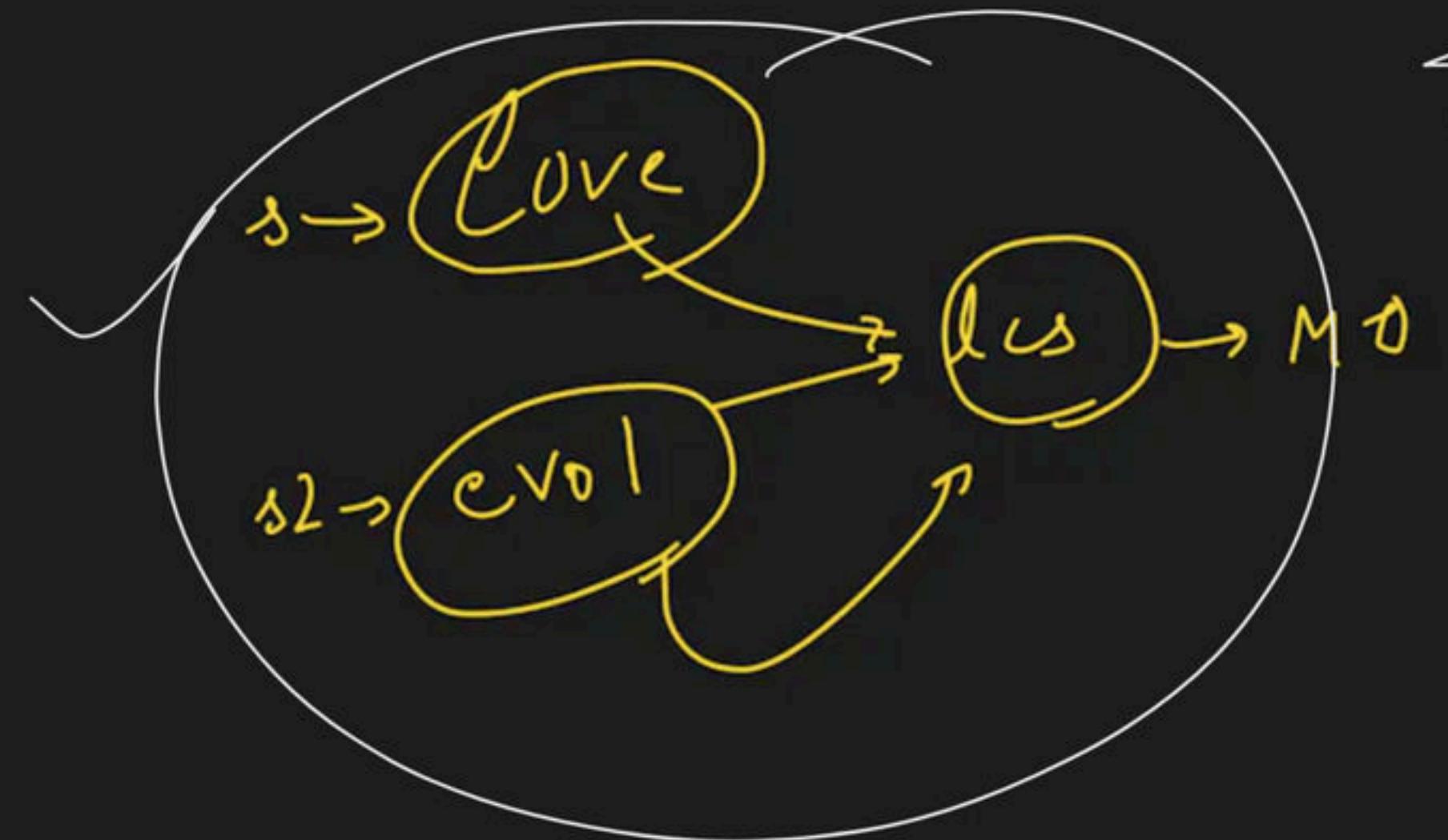
$\begin{array}{c} a \\ \times \\ y^2 \\ \hline A \end{array}$ $\begin{array}{c} b \\ \times \\ z \\ \hline B \end{array}$ $\begin{array}{c} c \\ \times \\ 1 \\ \hline C \end{array}$
 D.L \rightarrow i j Bond
 $\begin{array}{c} a \\ \times \\ y^2 \\ \hline A \end{array}$ $\begin{array}{c} b \\ \times \\ z \\ \hline B \end{array}$ $\begin{array}{c} c \\ \times \\ 1 \\ \hline C \end{array}$

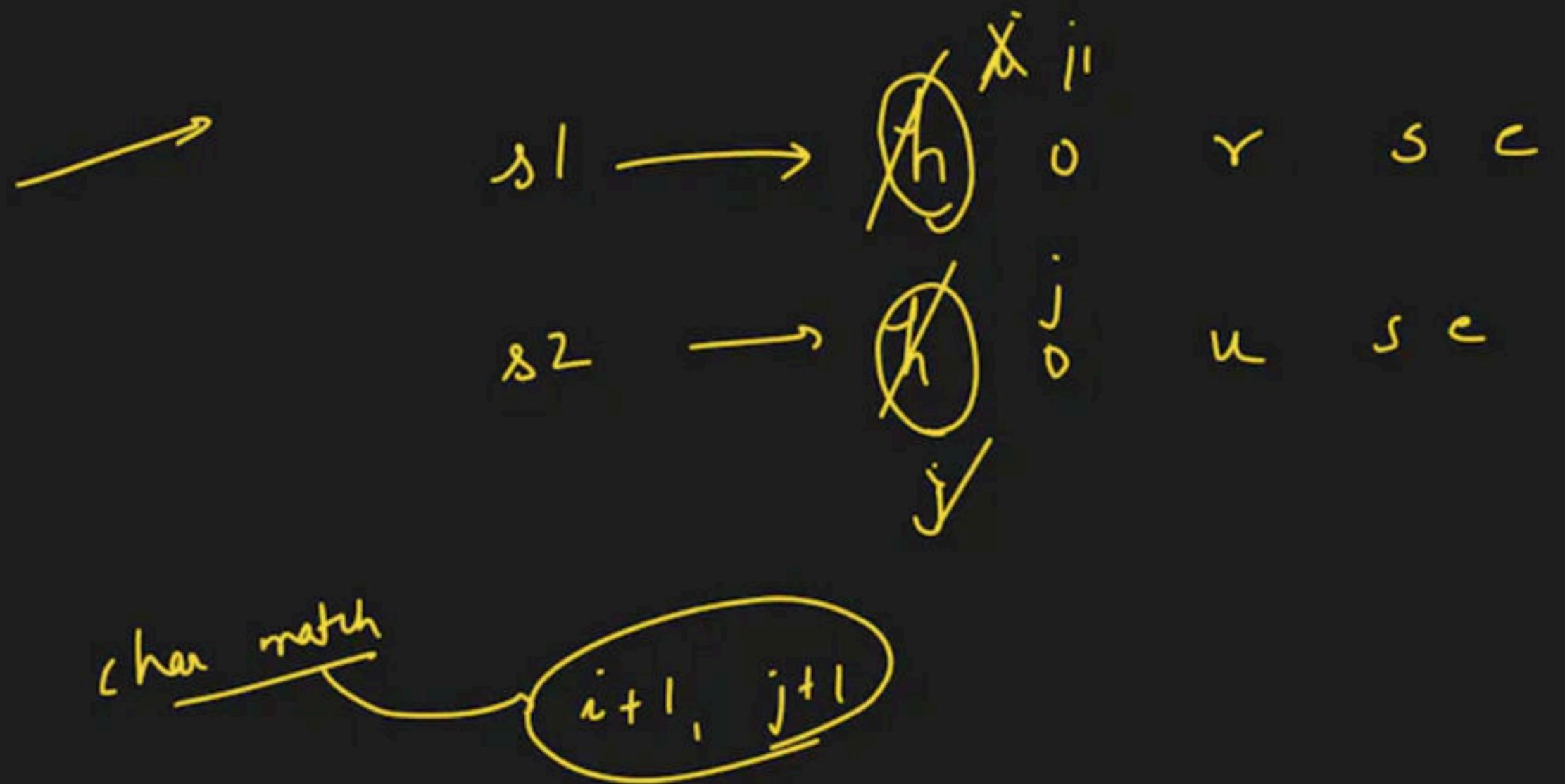
$\begin{array}{c} a \\ \times \\ y^2 \\ \hline A \end{array}$ $\begin{array}{c} b \\ \times \\ z \\ \hline B \end{array}$ $\begin{array}{c} c \\ \times \\ 1 \\ \hline C \end{array}$
 man
 char match
 char dont match
 0 + $\begin{array}{c} \text{Car A} \\ \text{Car B} \\ \hline \text{Am} \end{array}$ \rightarrow man





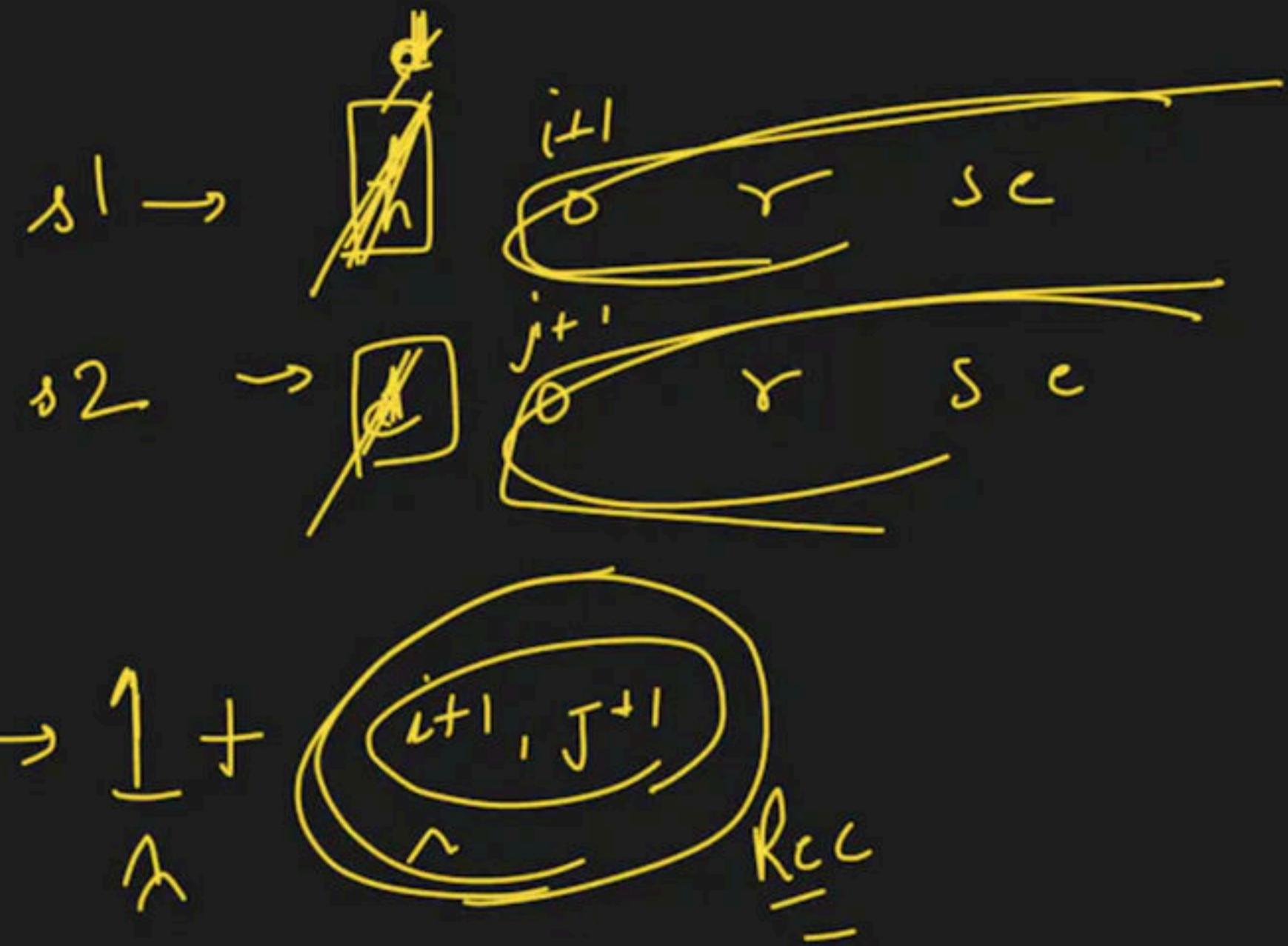






char
not
matching

$\gamma_{\text{replace}} \rightarrow \frac{1}{\lambda} +$



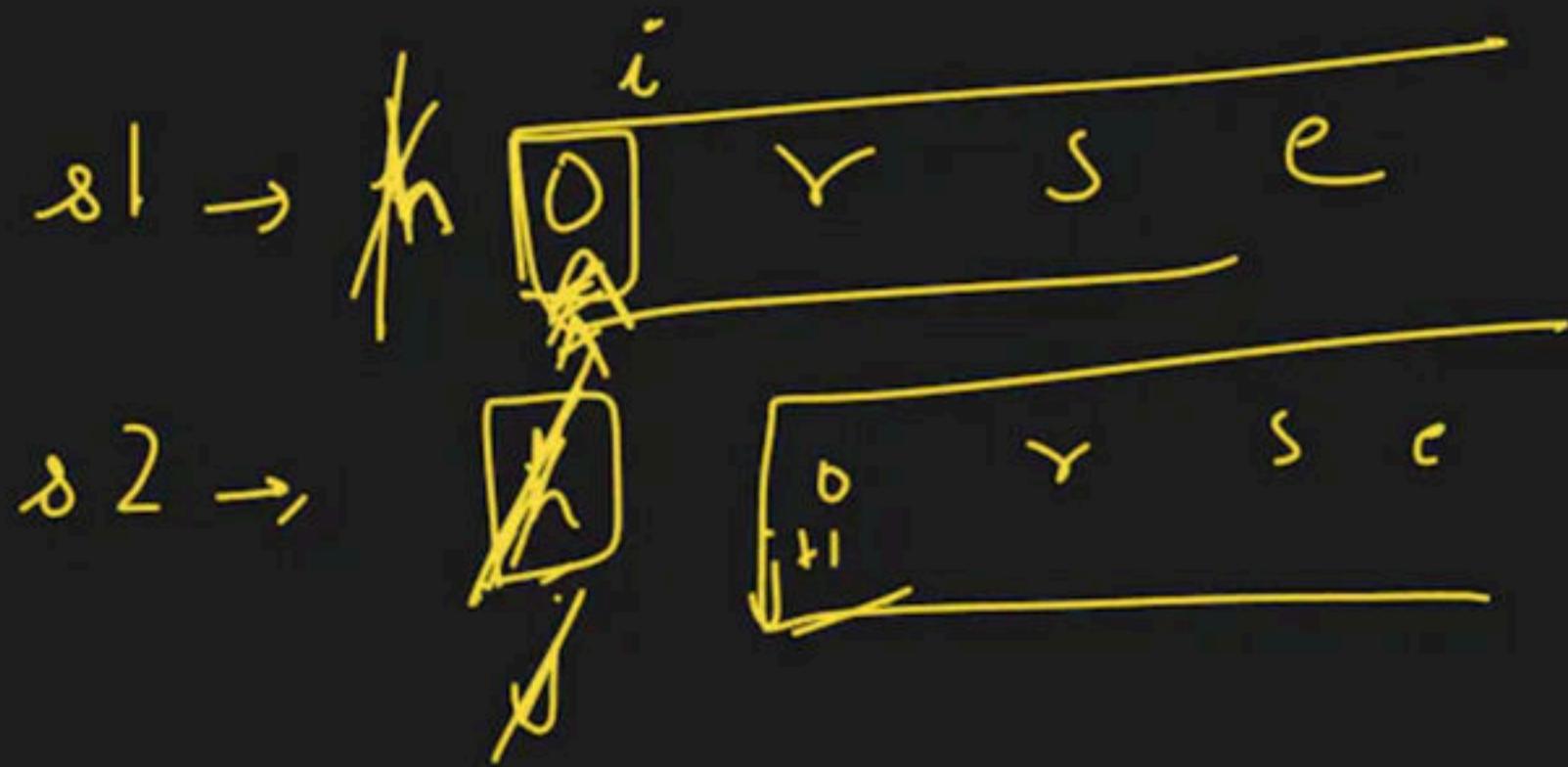
char
not

matching



word \rightarrow | +

$i \mid j \neq l$
Rec



char
not
natural

$$\text{duck} \rightarrow | + \underbrace{\left(\begin{smallmatrix} i+1 & i \\ j & R_{cc} \end{smallmatrix} \right)}$$

$\delta \rightarrow$ 

$\delta^2 \rightarrow$ 

char match

$$R_{\text{cc}} \rightarrow (i+1, j+1)$$

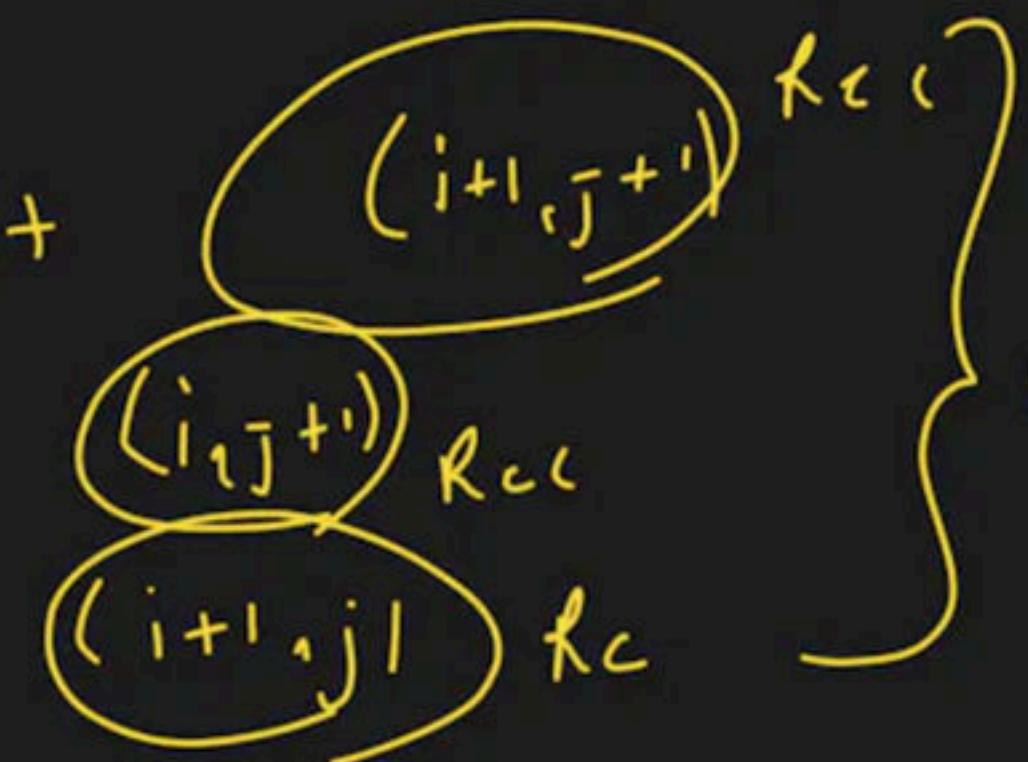
char not match

$$\text{replace} \rightarrow | +$$

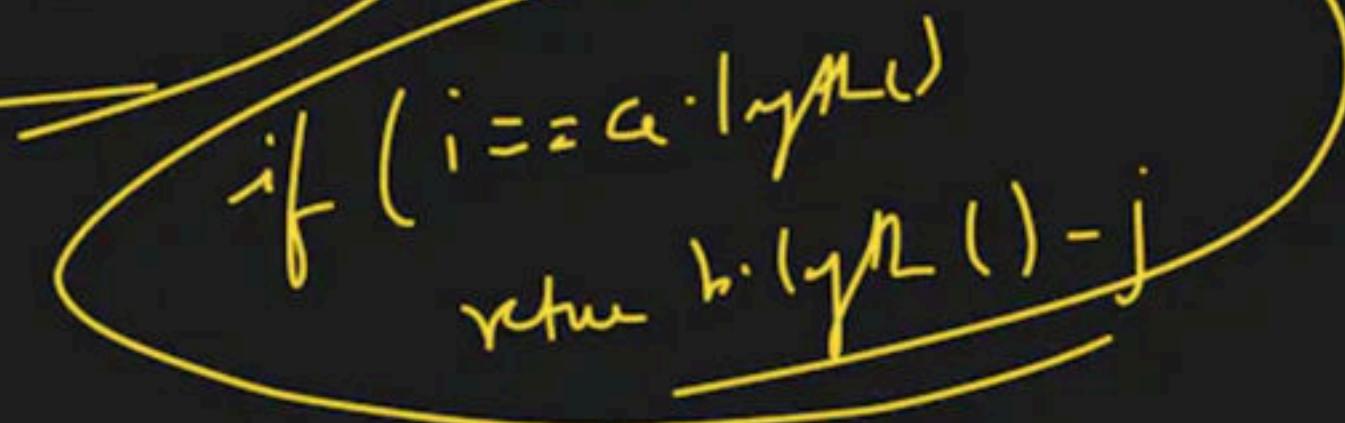
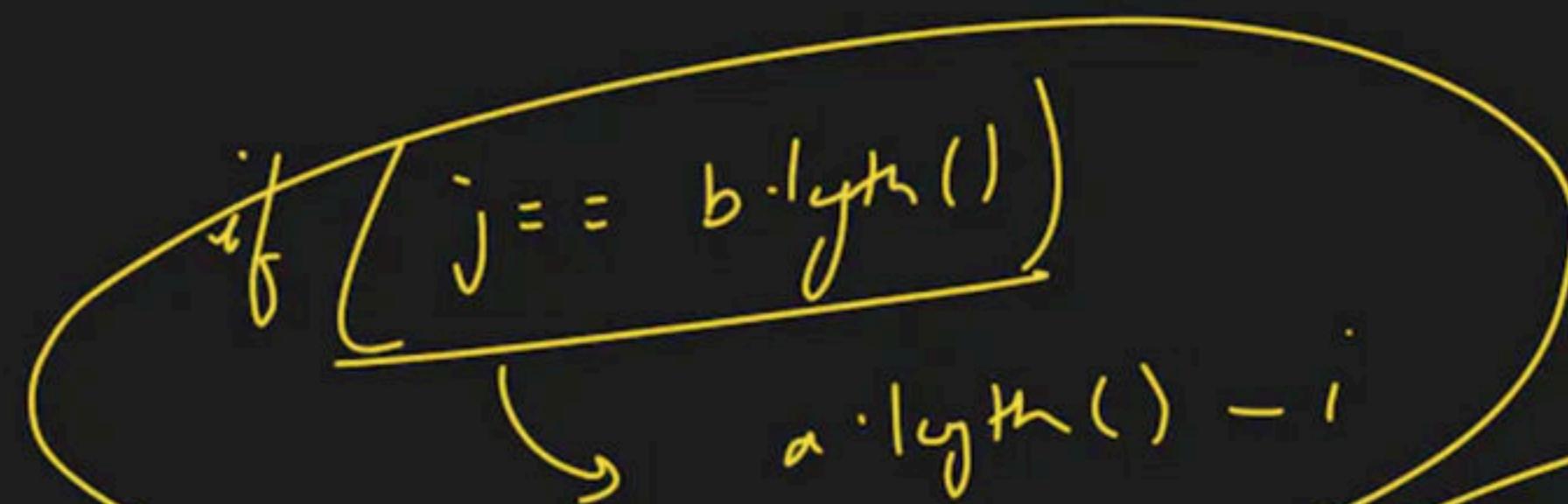
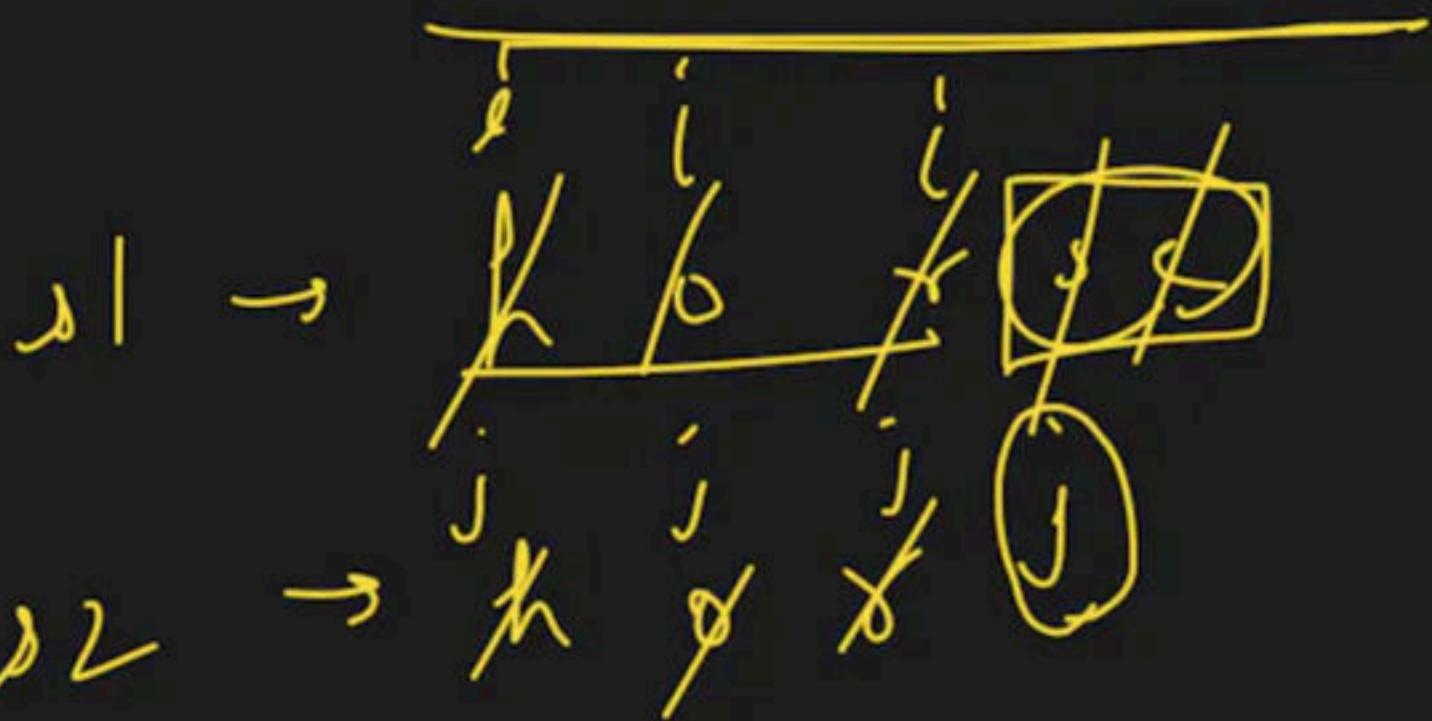
$$\text{input} \rightarrow | +$$

dotstar

$$| +$$



mih



Rec \rightarrow B-C

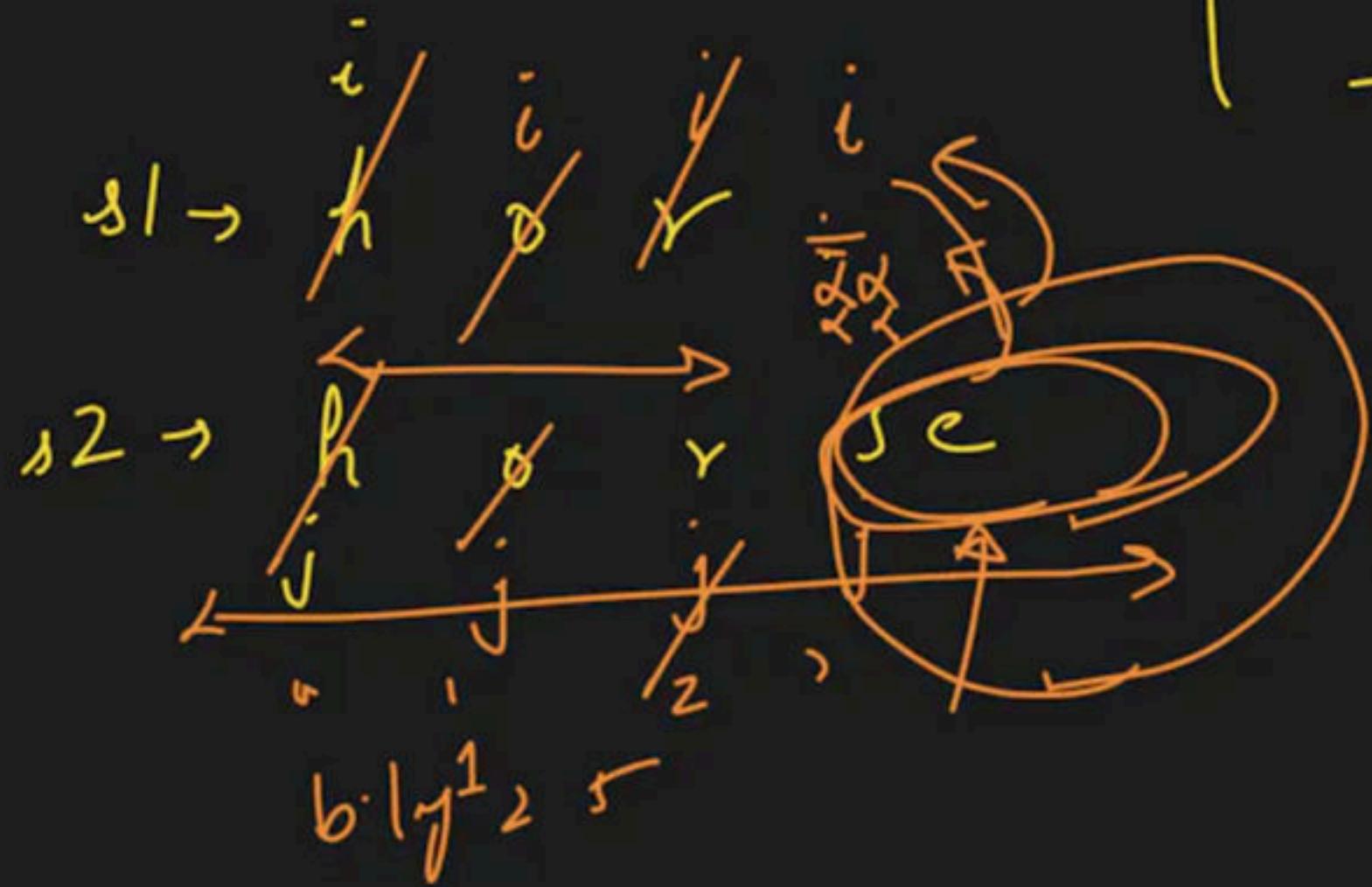
if ($i == a.length()$)

{

return

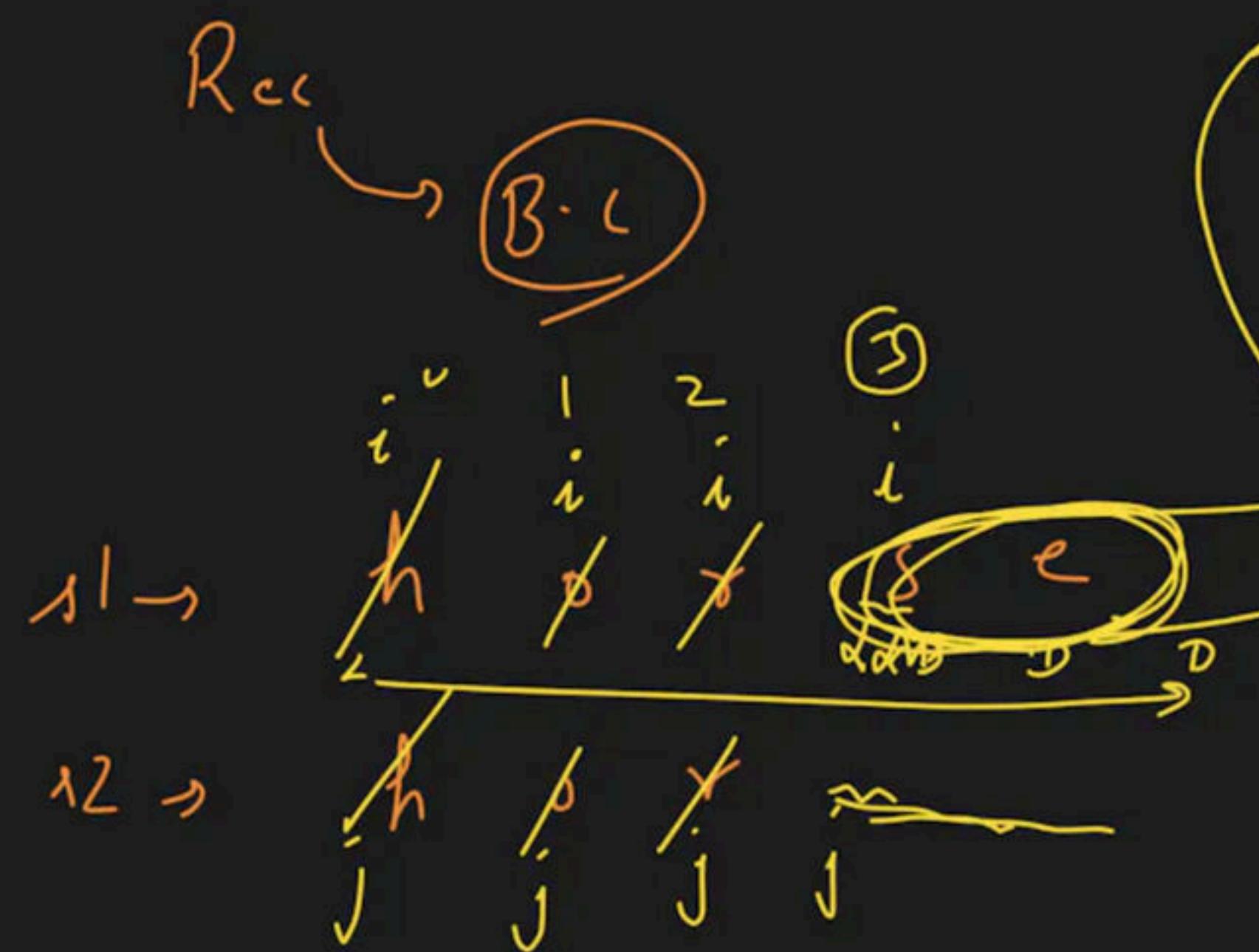
}

b.depth() - j



$$j = 3$$

$$b \cdot \ln - j \\ 5 - 3 = 2$$



$$a \cdot \ln - i$$

$$5 - 3 = 2$$

if

$$i = a \cdot \lfloor n \rfloor$$

return

$$b \cdot \lfloor n \rfloor - j$$

if

$$j = b \cdot \lfloor n \rfloor$$

return

$$a \cdot \lfloor n \rfloor - i$$

for ($i = 0; i <= a \cdot \lfloor n \rfloor; i++$)

$$dp[i][b \cdot \lfloor n \rfloor] = a \cdot \lfloor n \rfloor - i$$

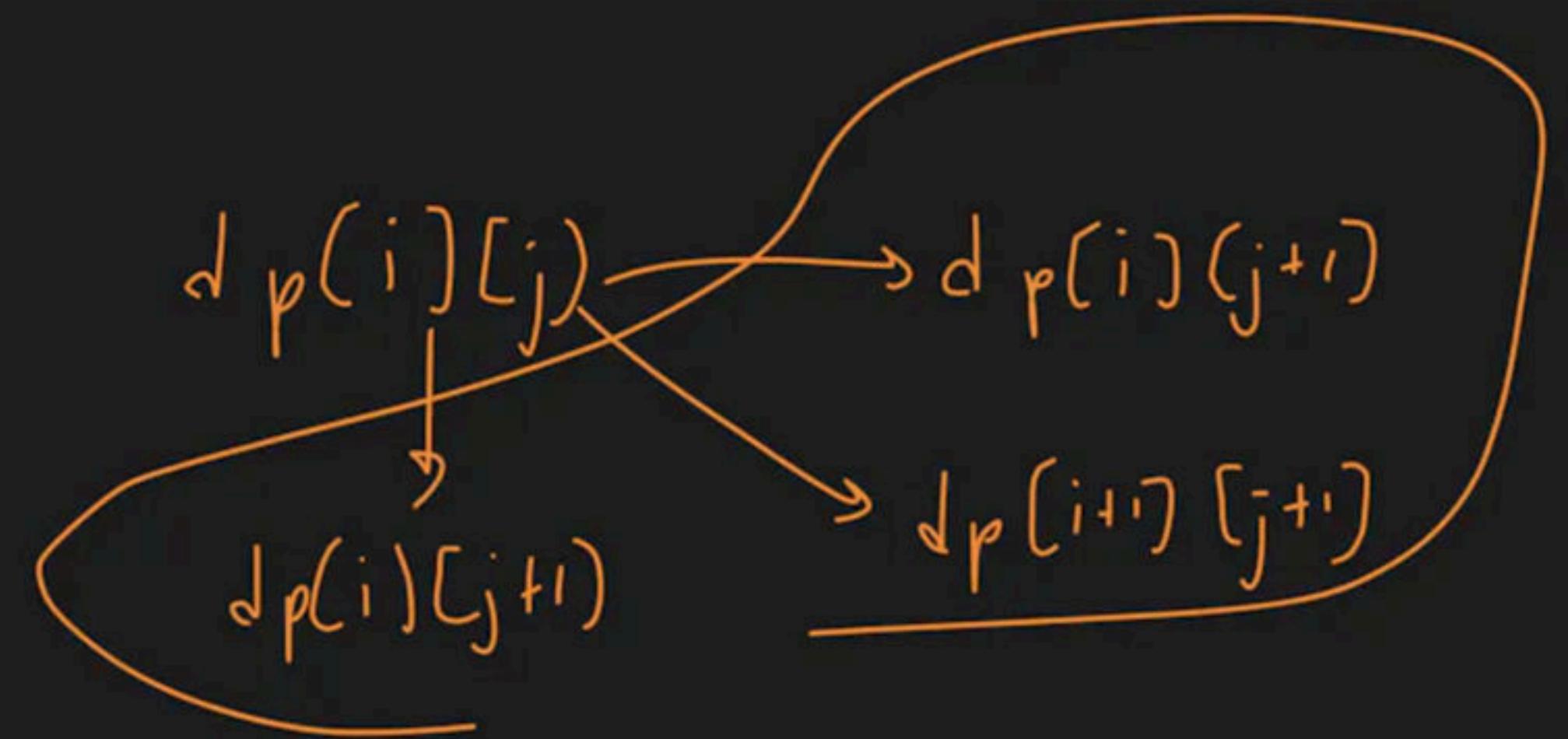
for ($j = 0; j <= b \cdot \lfloor n \rfloor; j++$)

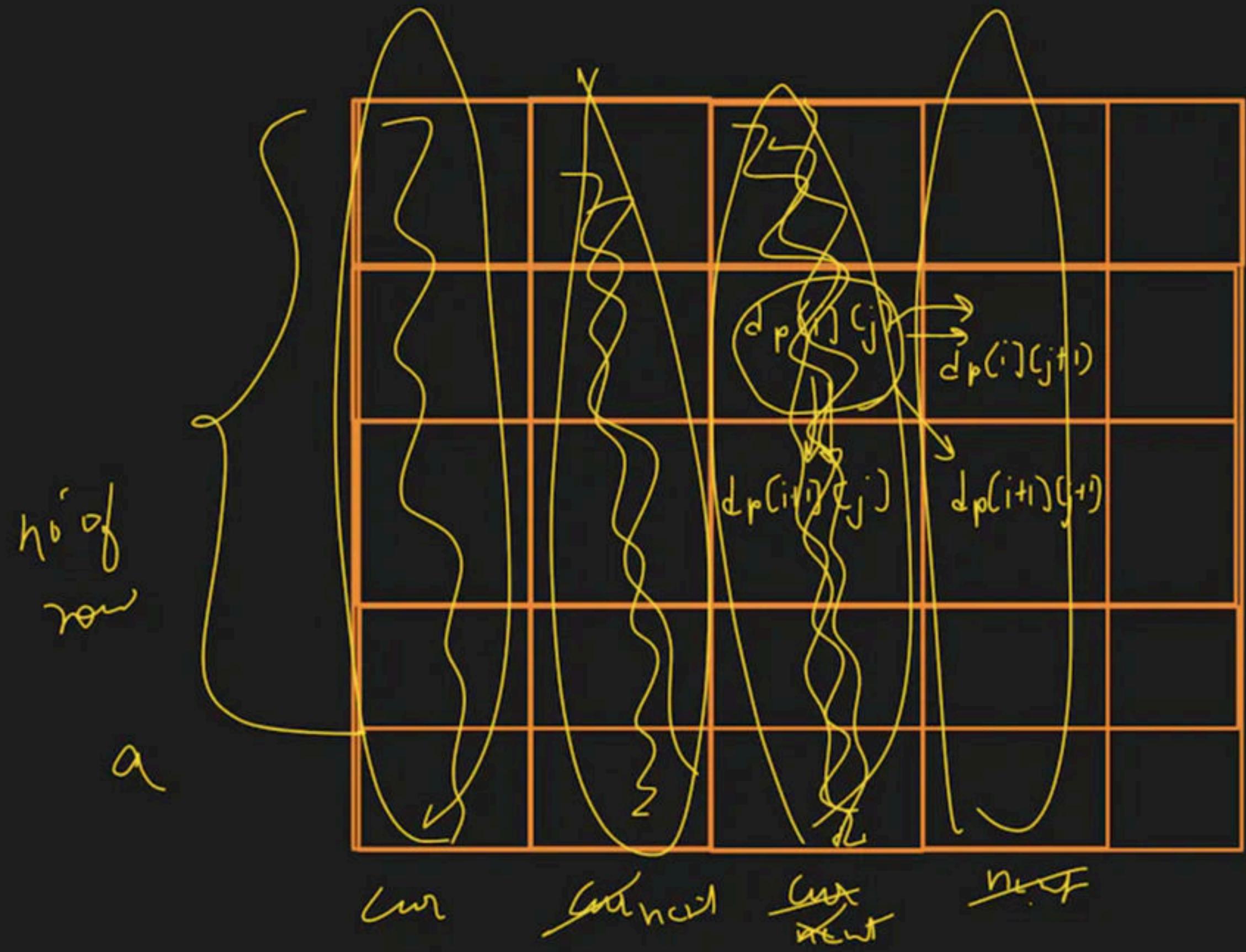
$$dp[a \cdot \lfloor n \rfloor][j] = b \cdot \lfloor n \rfloor - j$$

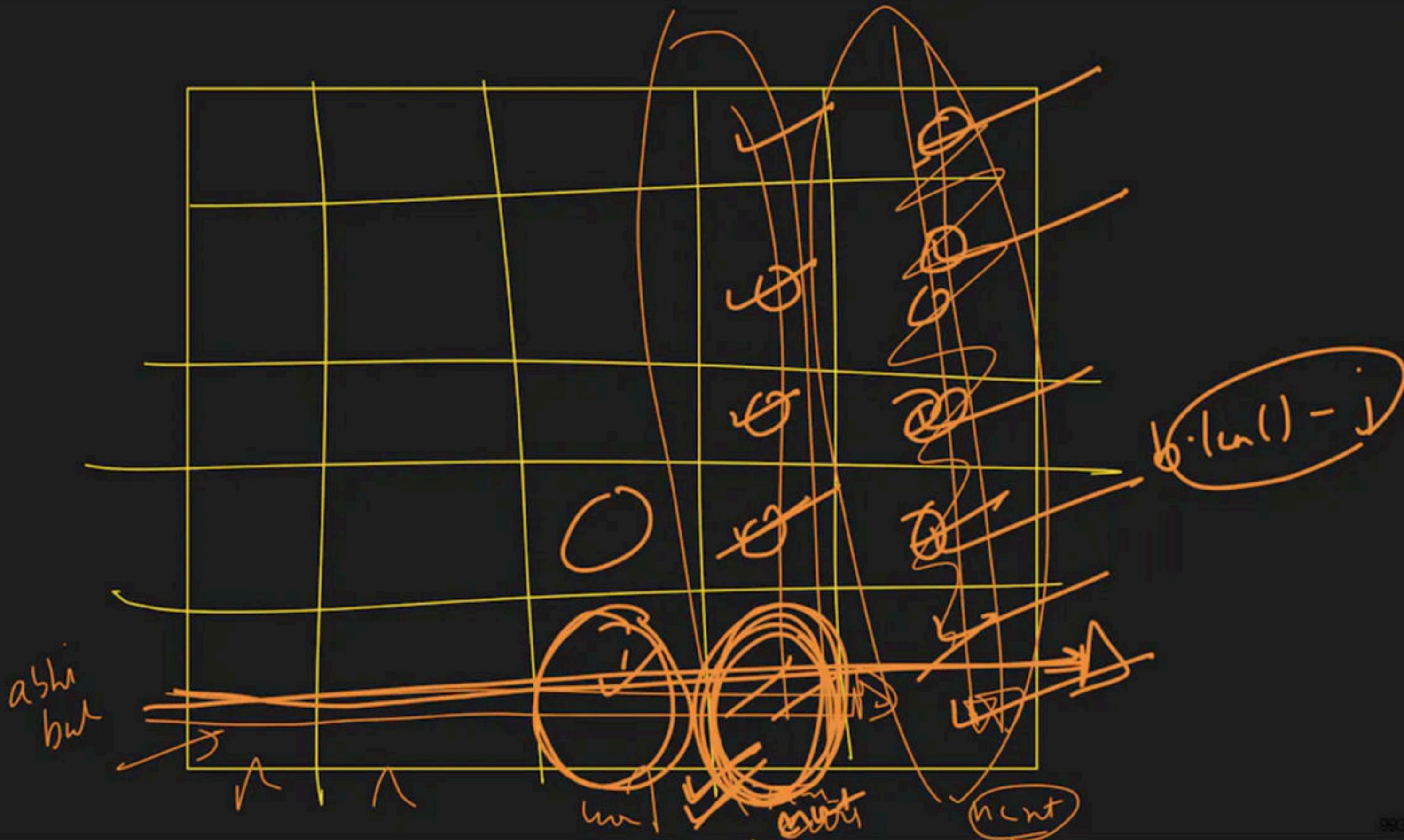
0	1	2	3	4
0	1	2	3	4
0	1	2	3	4
0	1	2	3	4
0	1	2	3	4

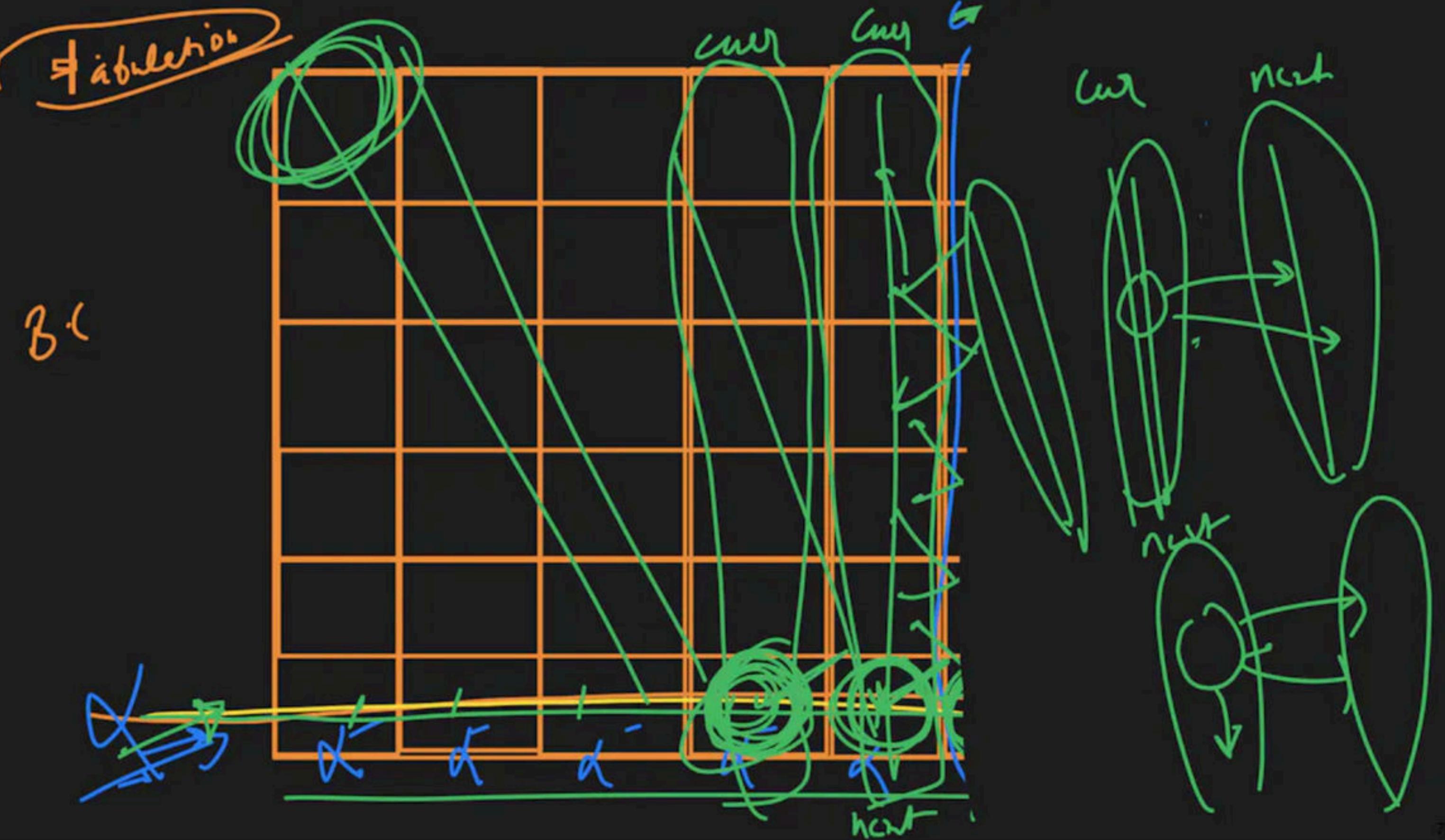
$a+1$

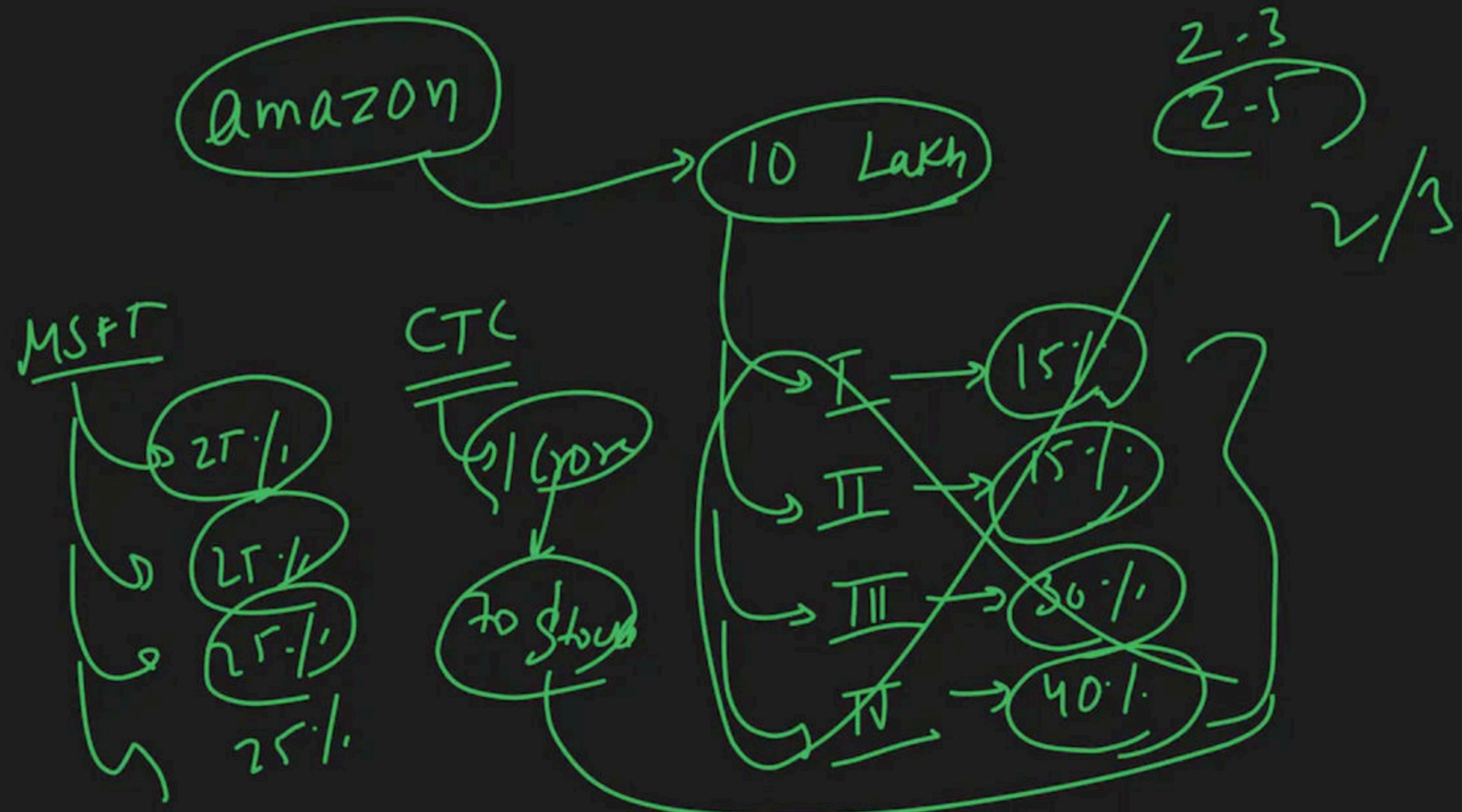
row



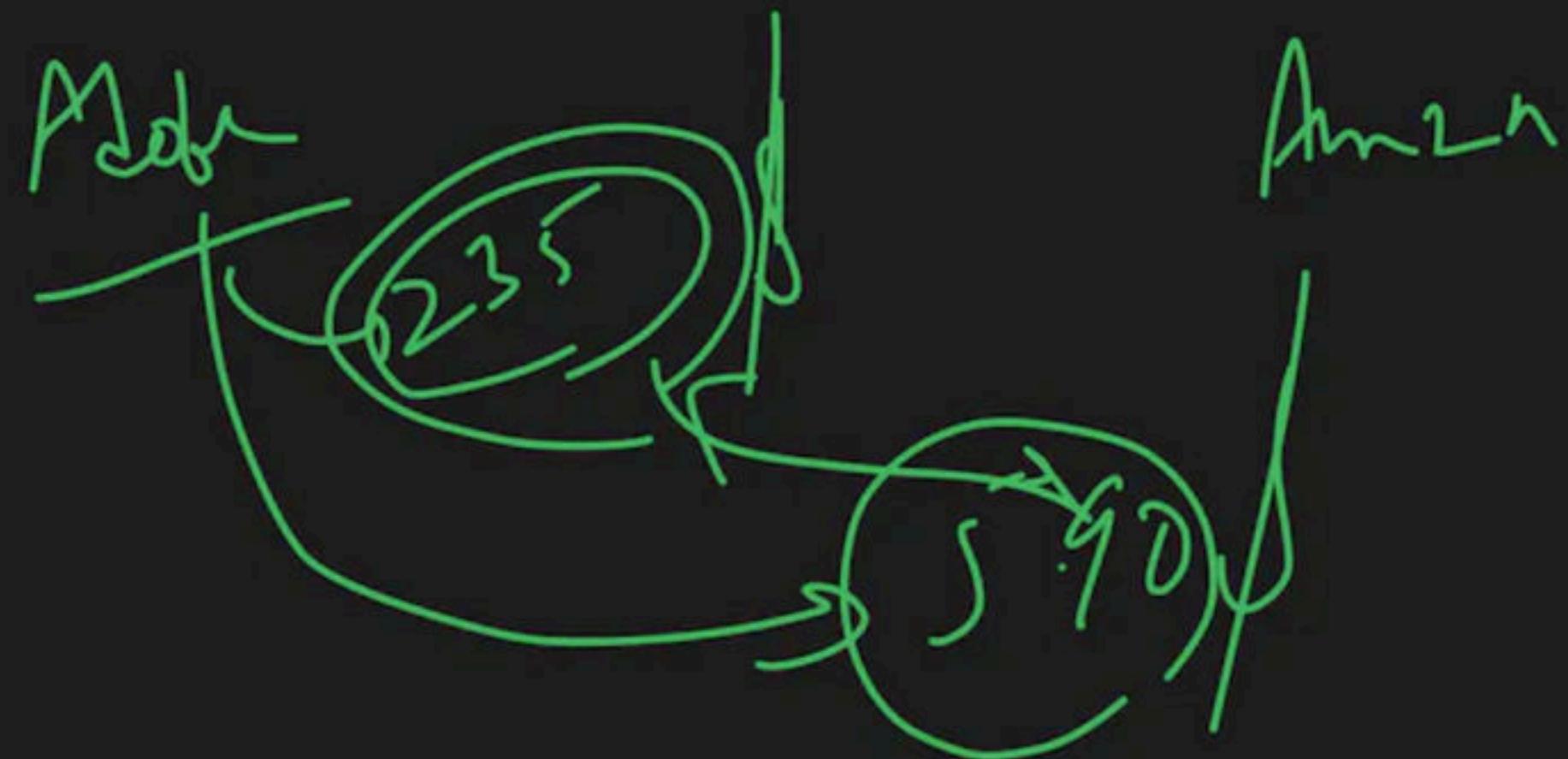


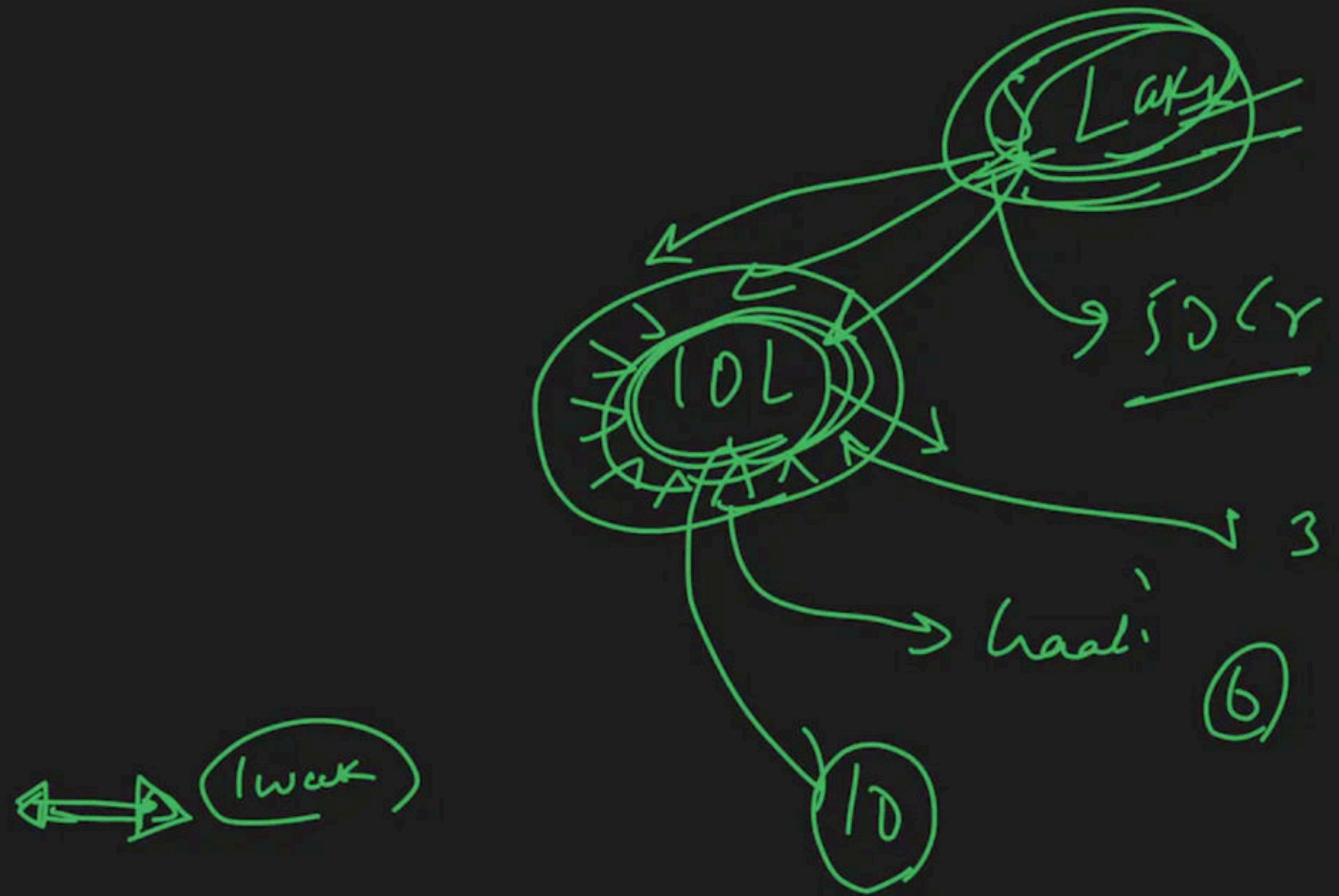


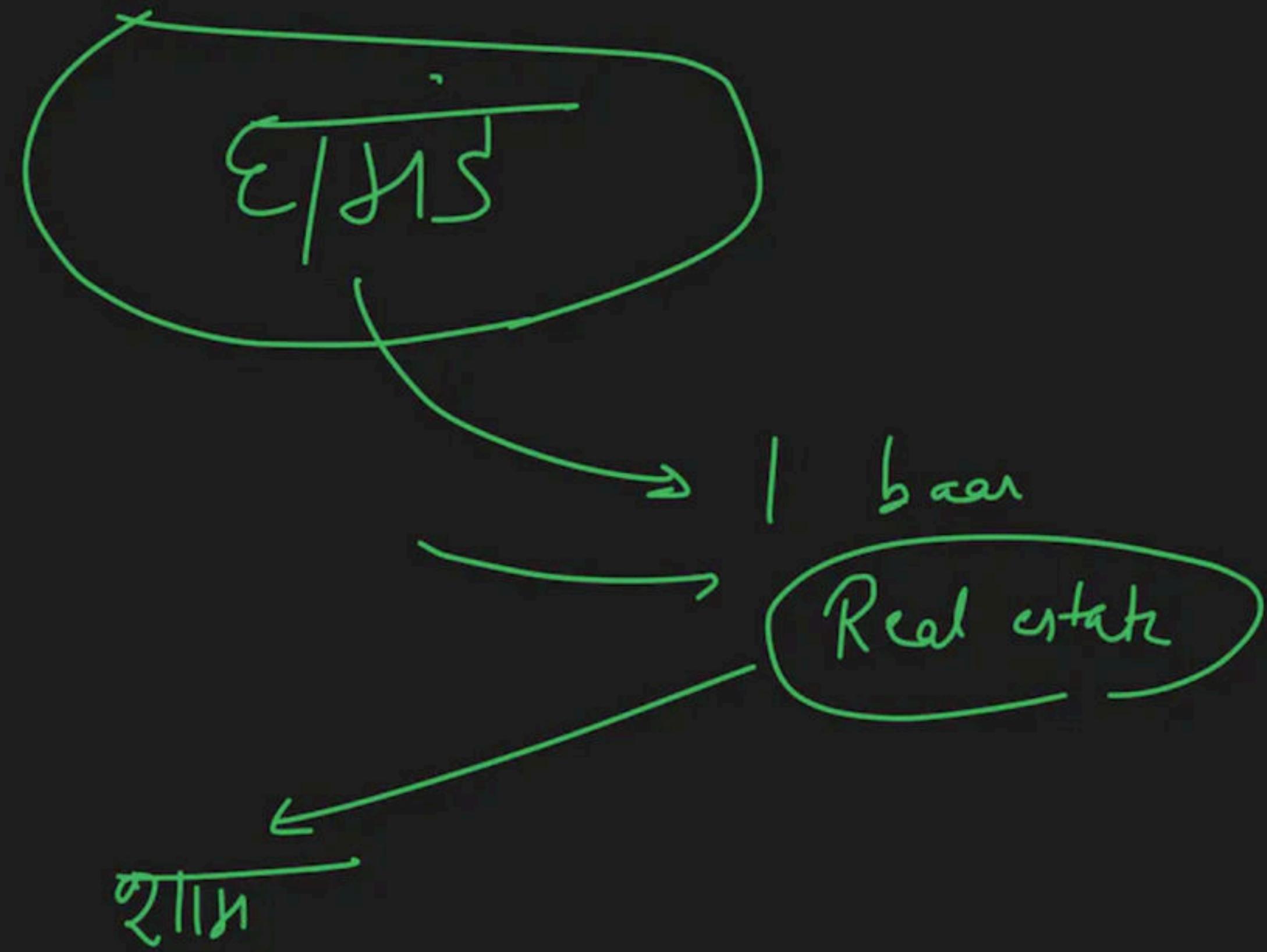




July 2019







①

Knapsack \Rightarrow unbounded

②

B.S. rule

③

$LAP \Rightarrow$

④

$LIS \Rightarrow$

⑤

Cub