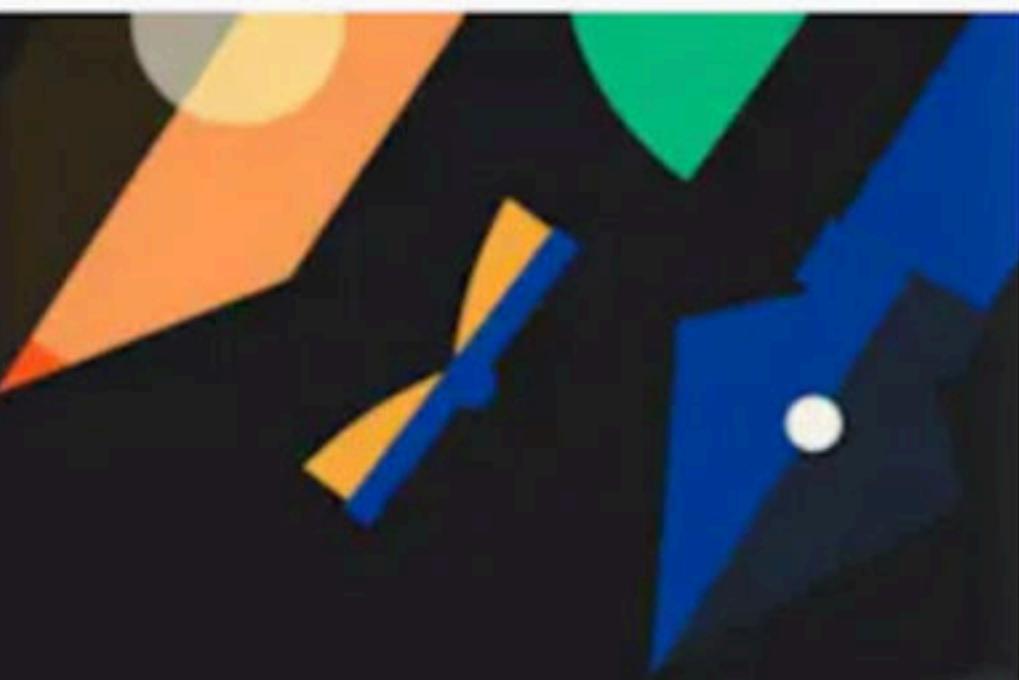


Queue  
==

Not finished  
=====

# Doubt Class by Lakshay Bhaiya

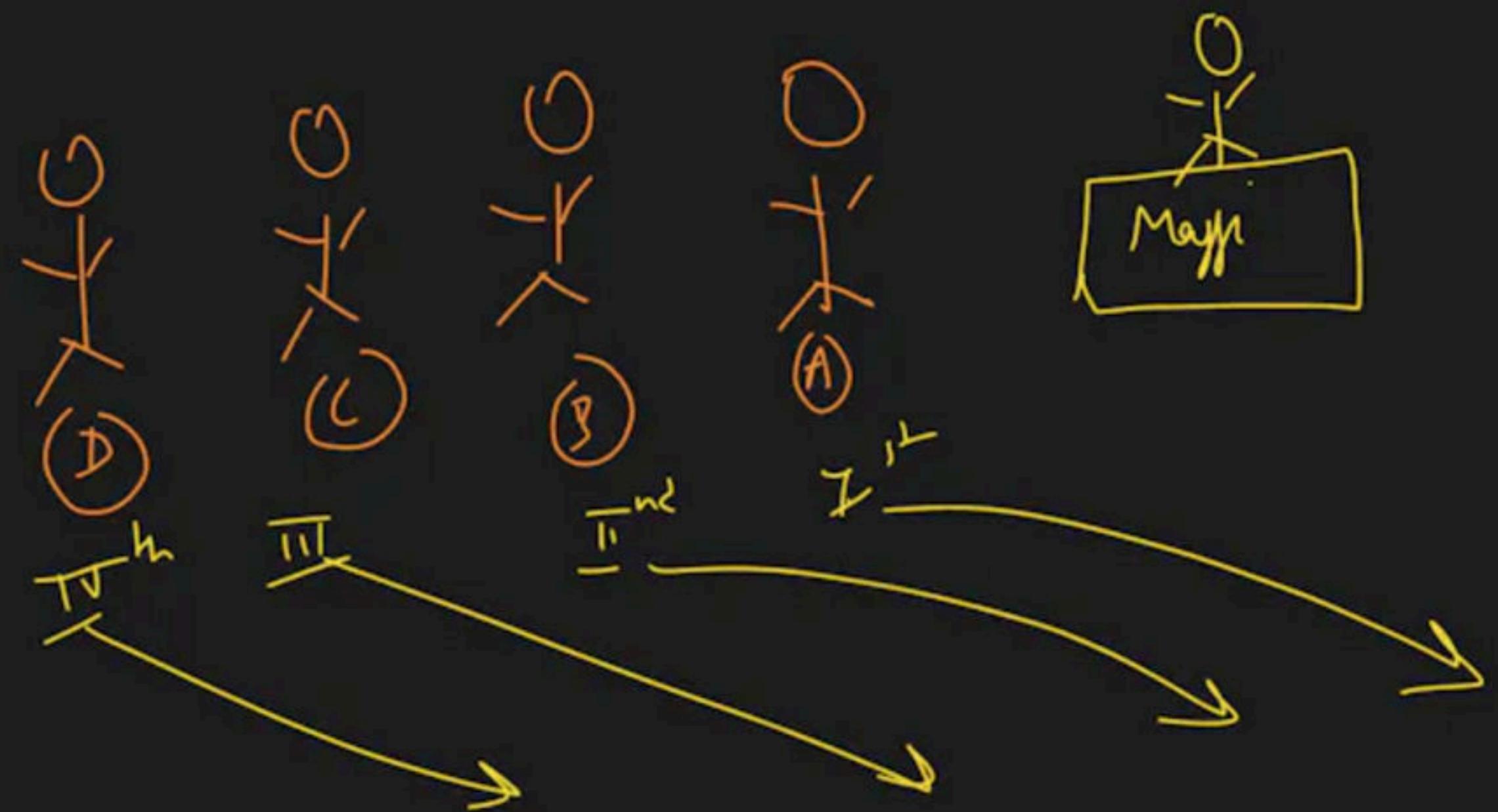
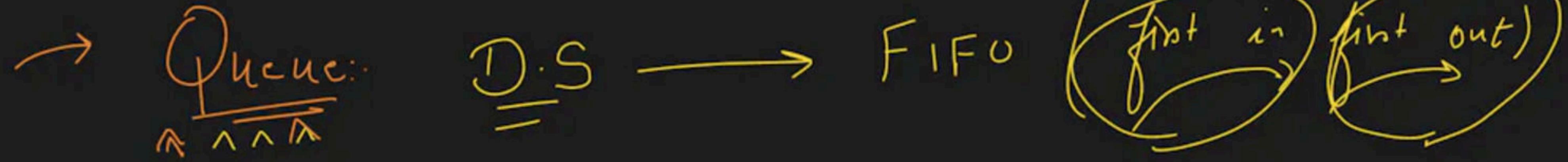
Special class

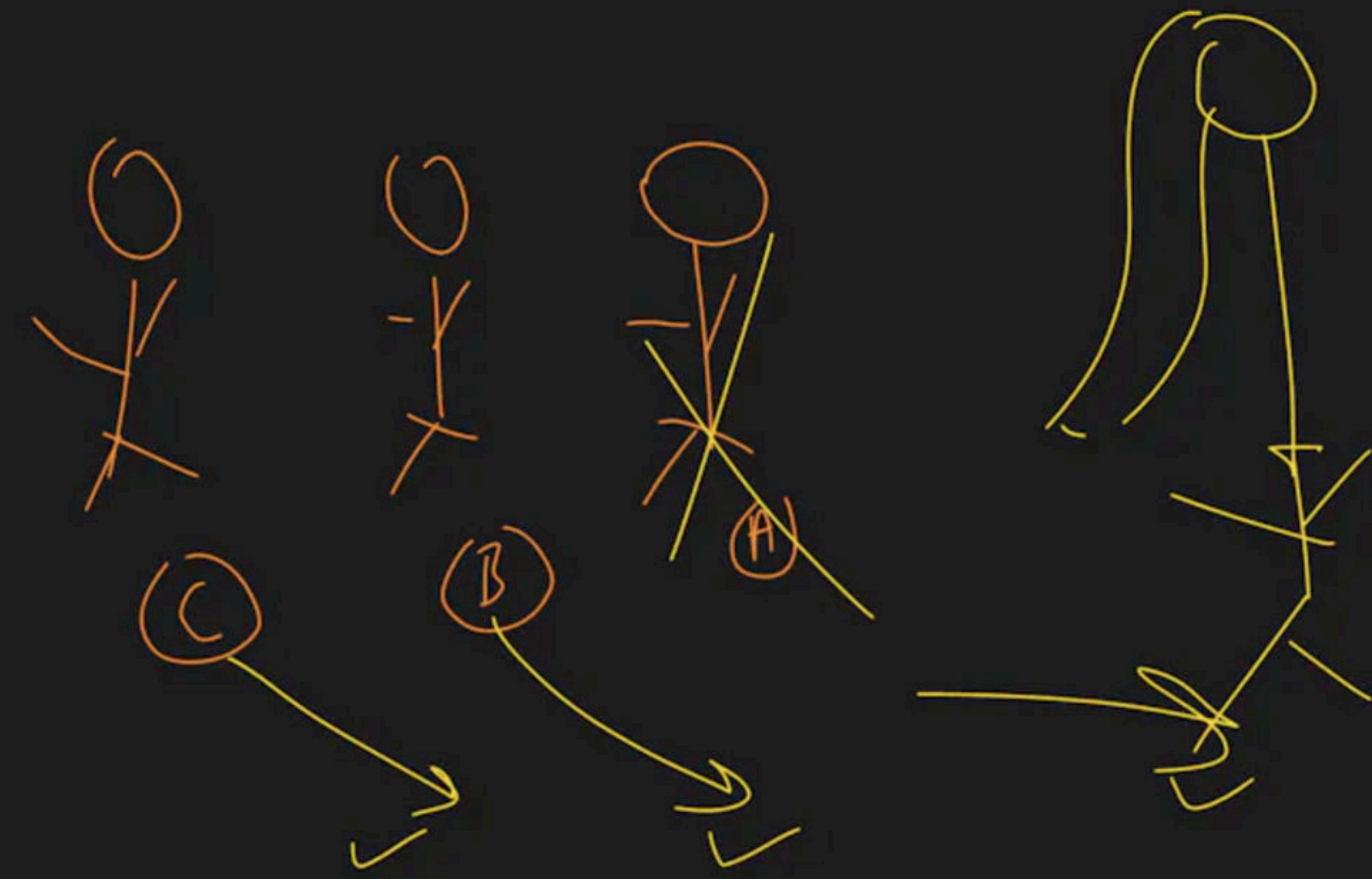


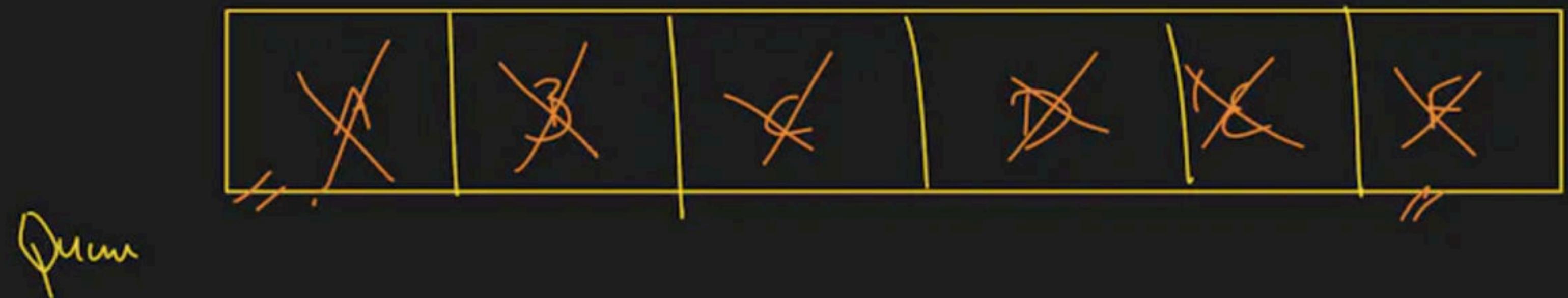
# Queue Class - 1 [Join Here]

Special class

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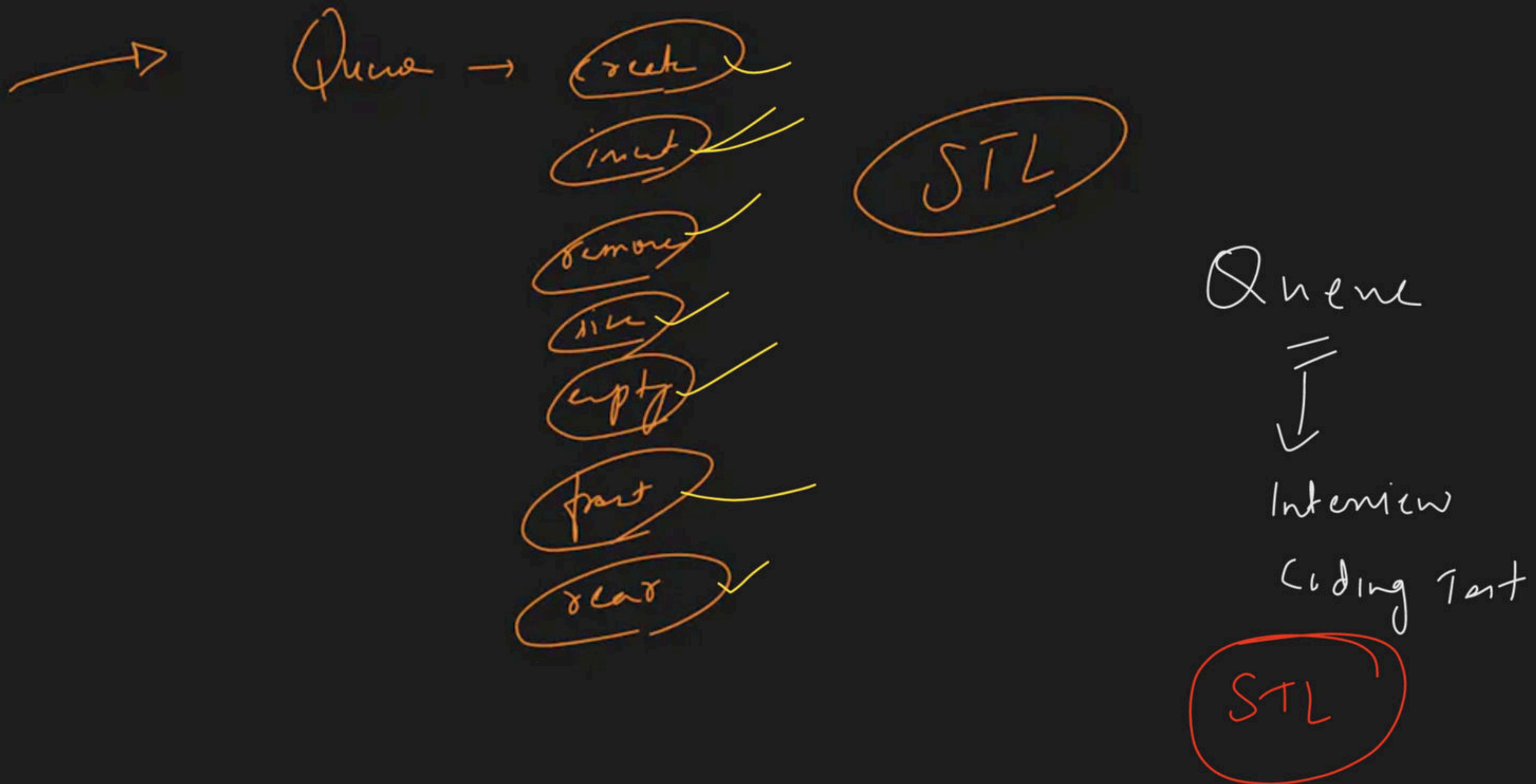


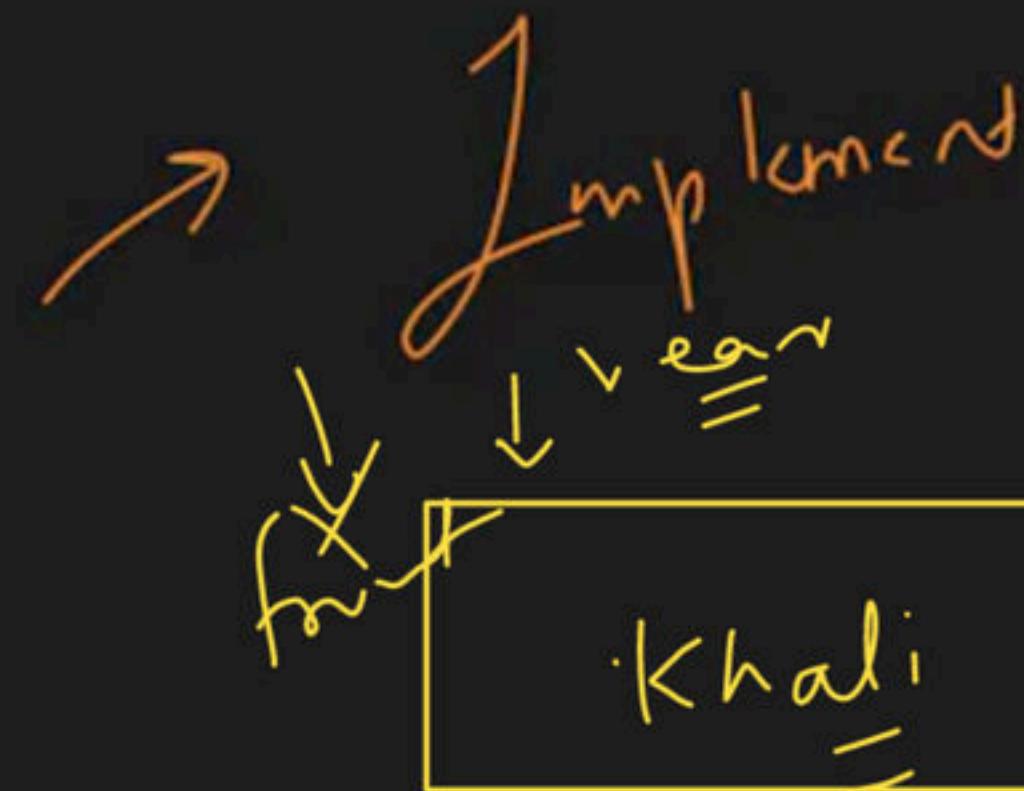




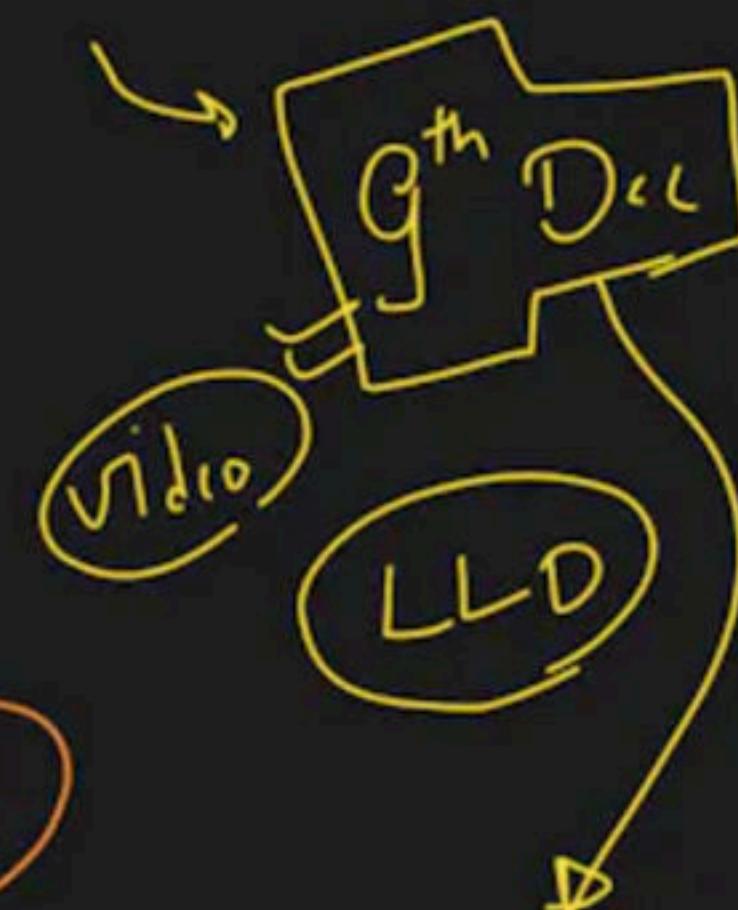
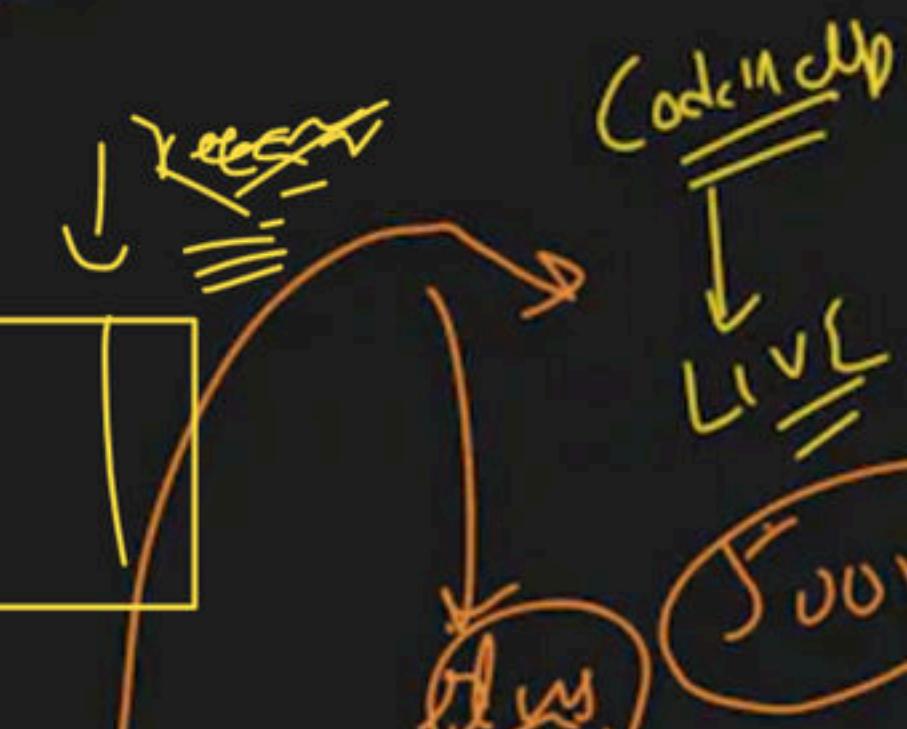
→ Stack → LIFO

→ Queue → FIFO





Queue from  
Scratch



- 1
- 2 Circular Q

1 DDL

No Pkg

LLD  
Cmpt

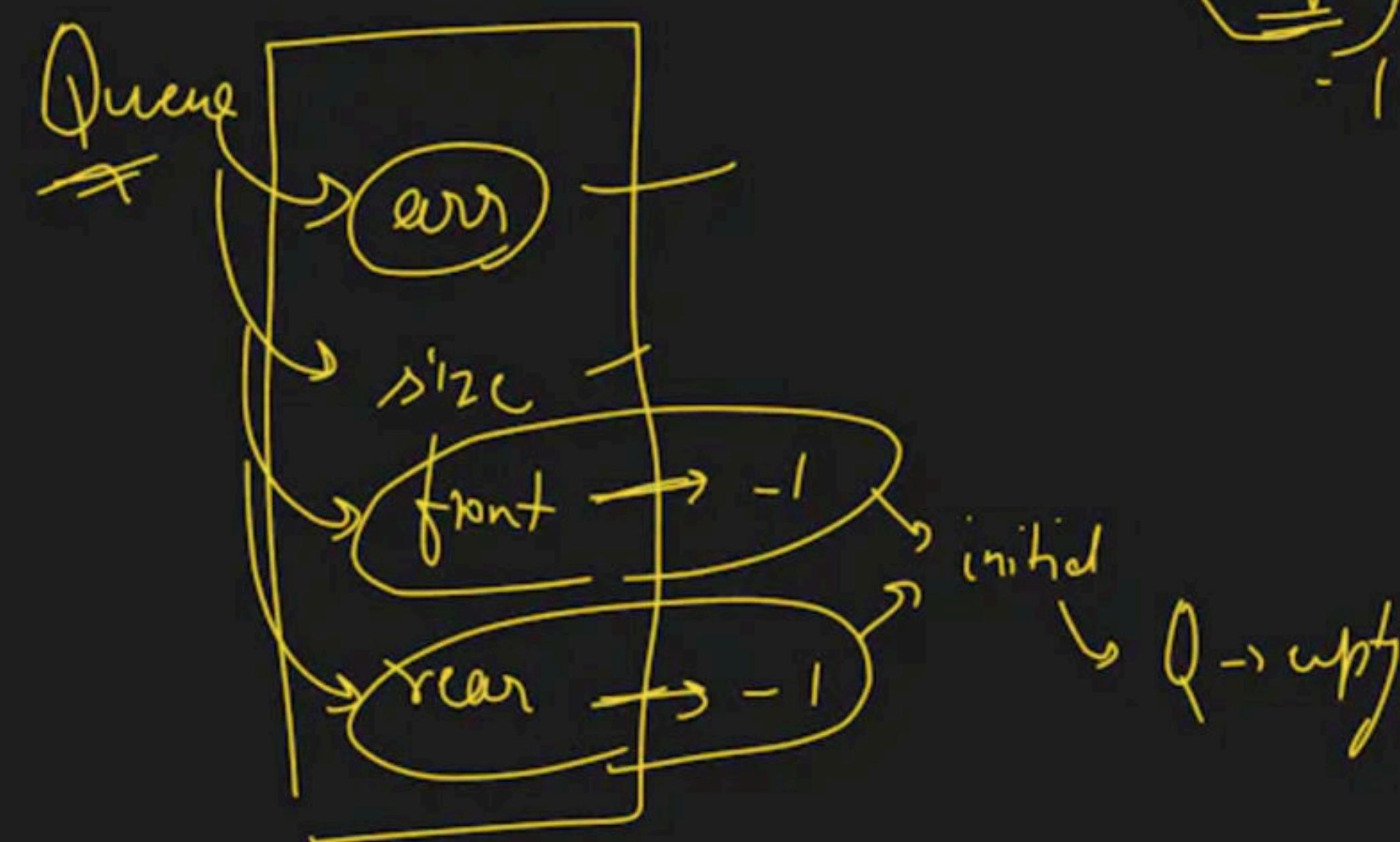
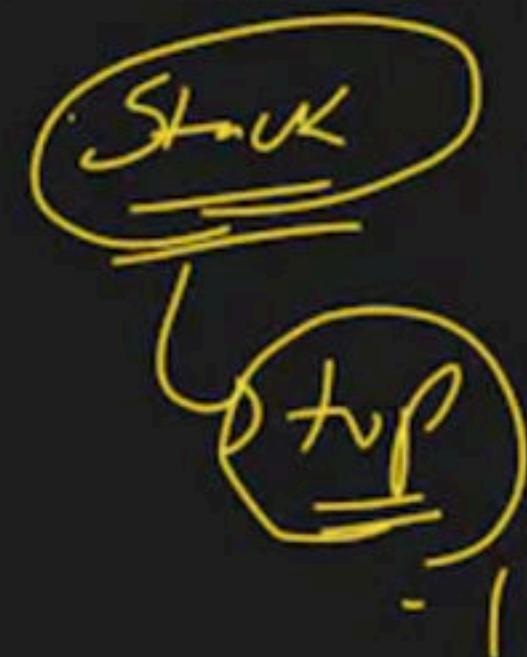
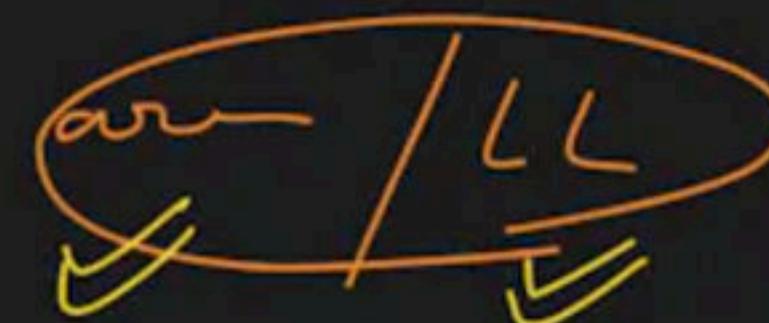
Java

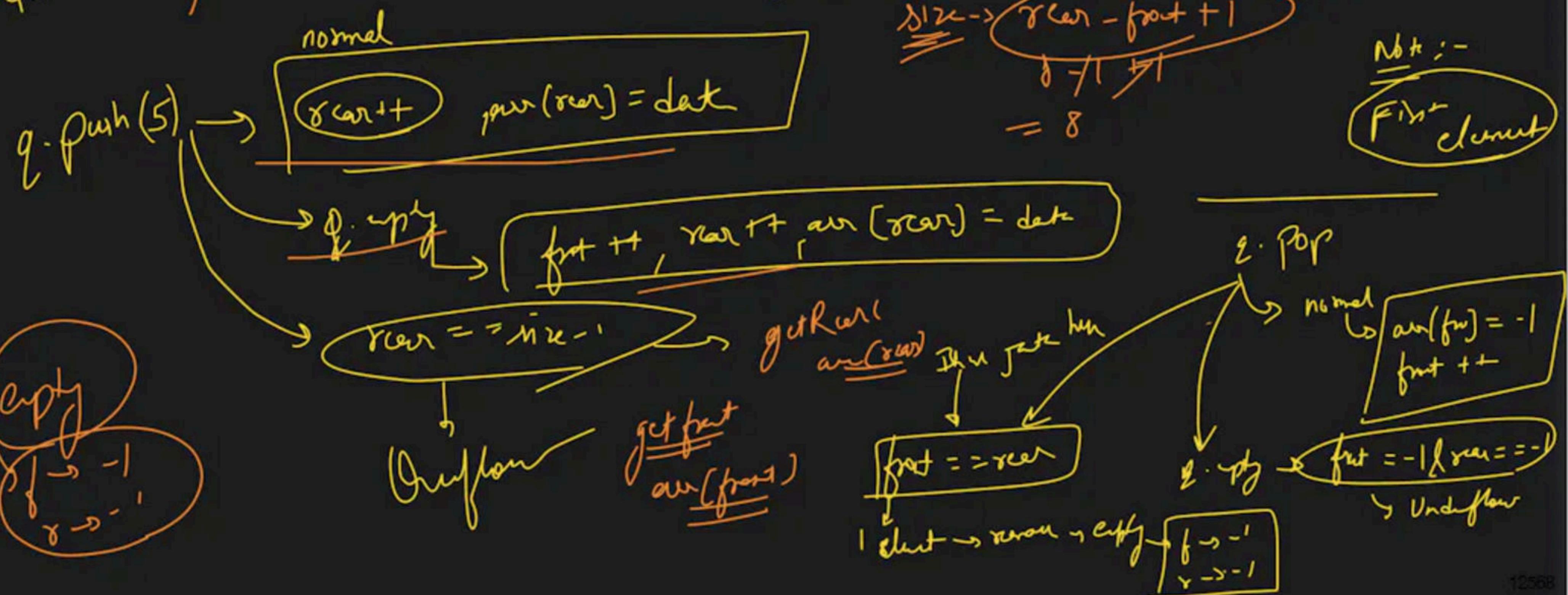
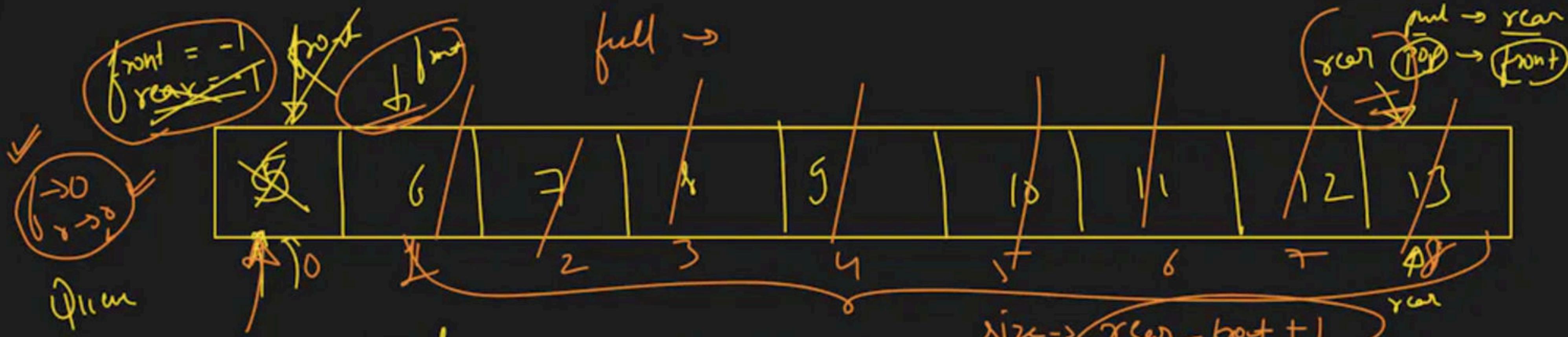
OLD  
Student  
5/10/11  
1015  
Sohail

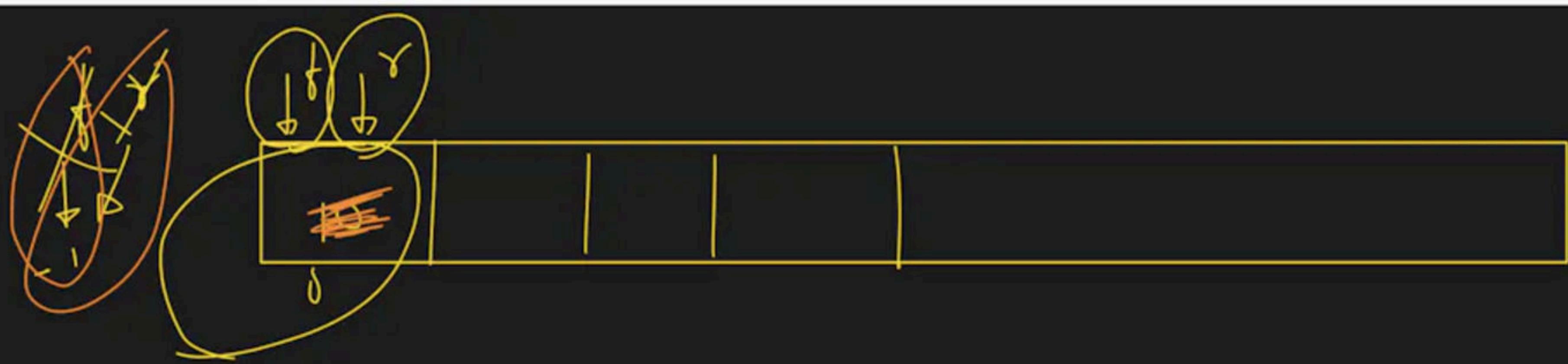
Implement



Queue from Scratch







$q.push(10)$

$q.push(1)$

$f = r = 0$  → single element

$f = r$  → removed →  $q.empty()$

$f \rightarrow -1$   
 $r \rightarrow -1$

$Q$  → push → if  
char in  $q$  → } →  $O(1)$

pop } } → } →  $O(1)$

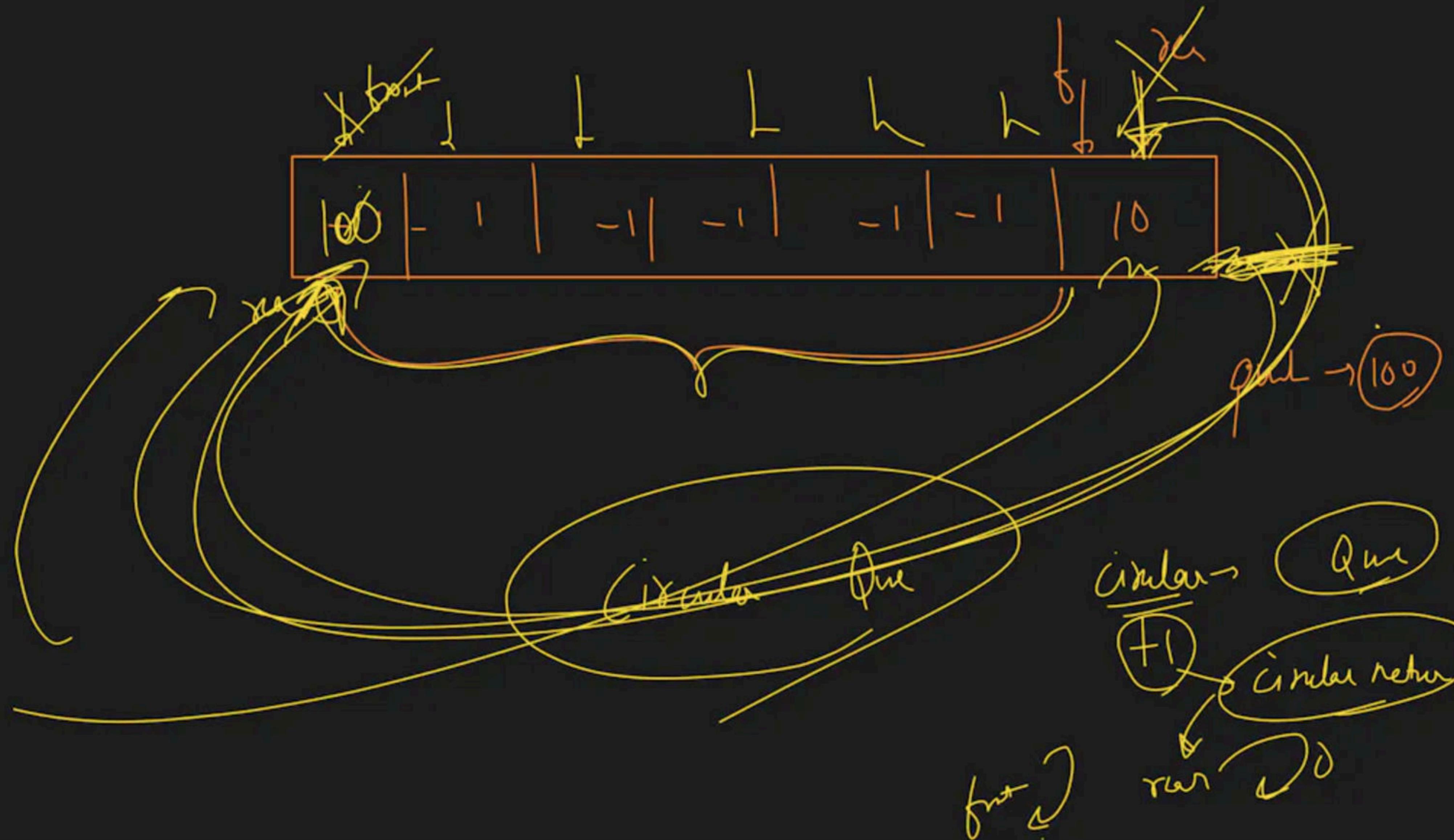
size } → } →  $O(1)$

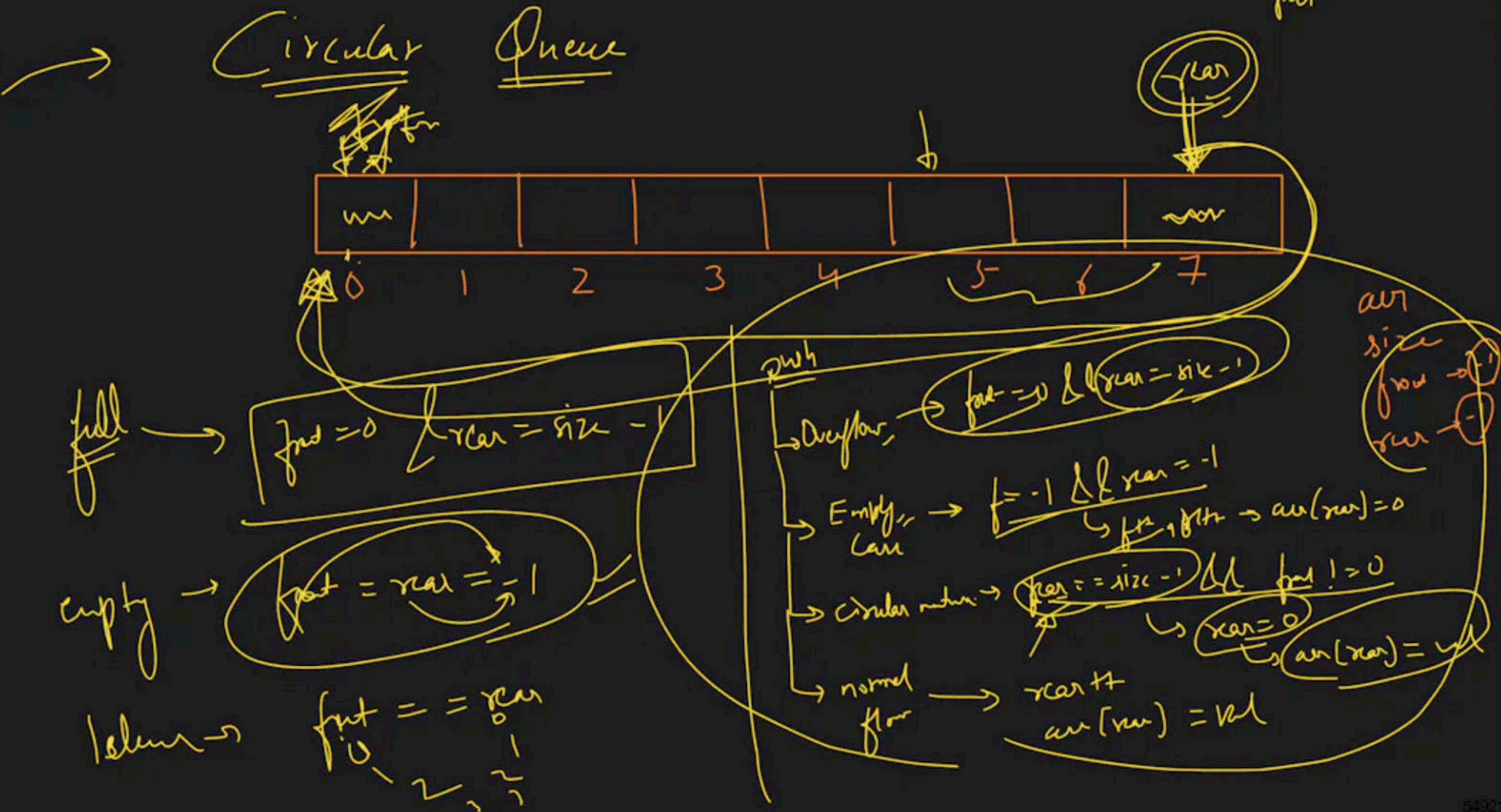
empty } → } →  $O(1)$

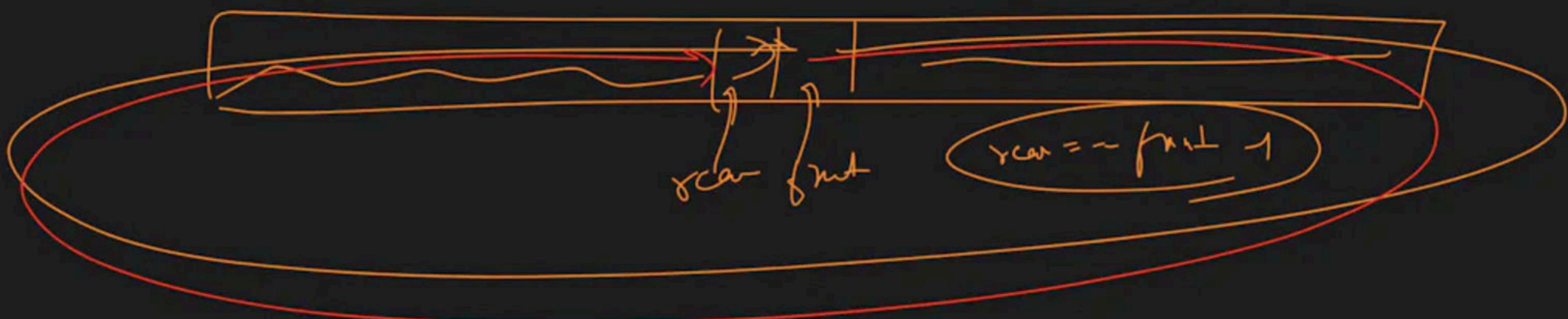
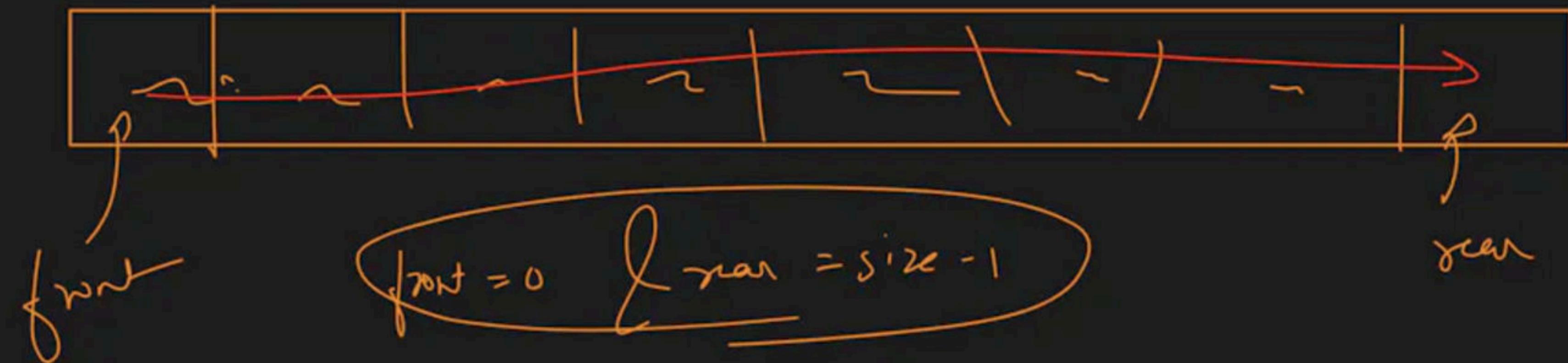
get front / get back } →  $O(1)$

Queue <int>  $q$ ;

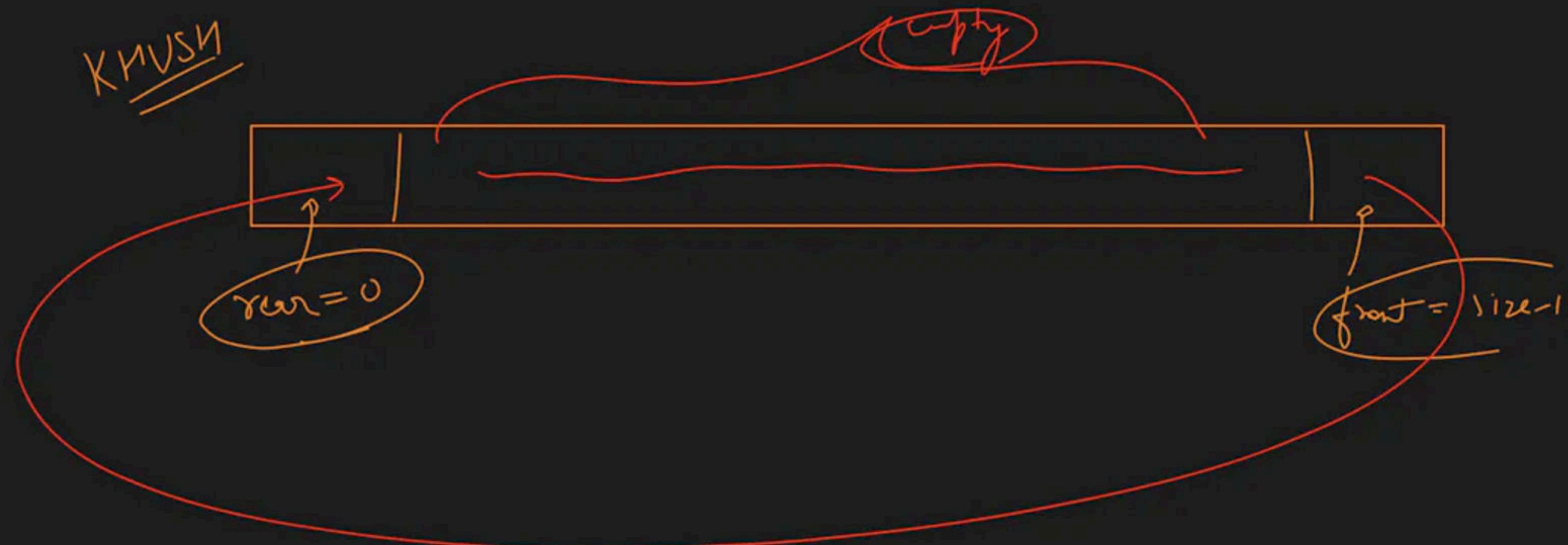
=  
↓  
Circular  
=

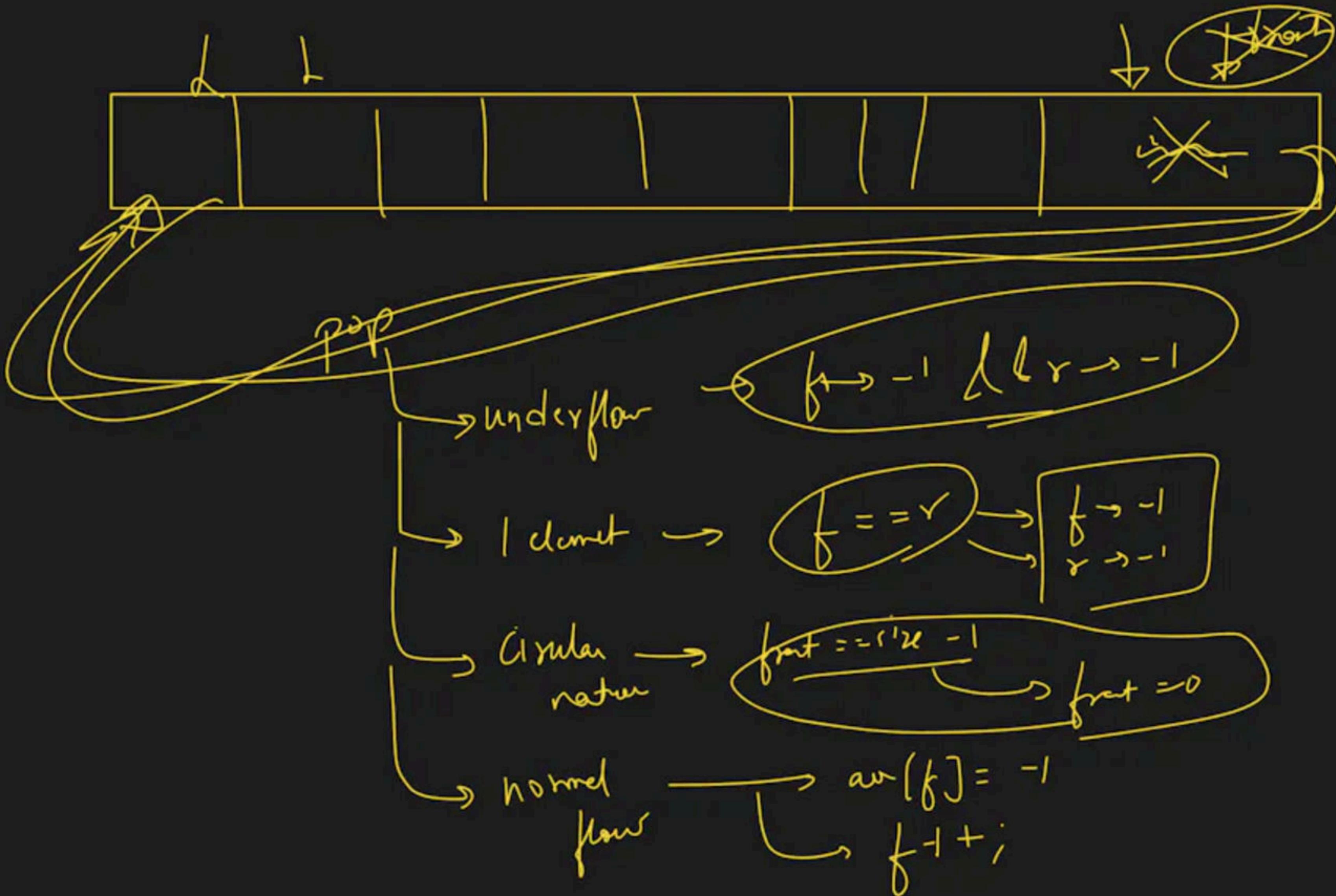


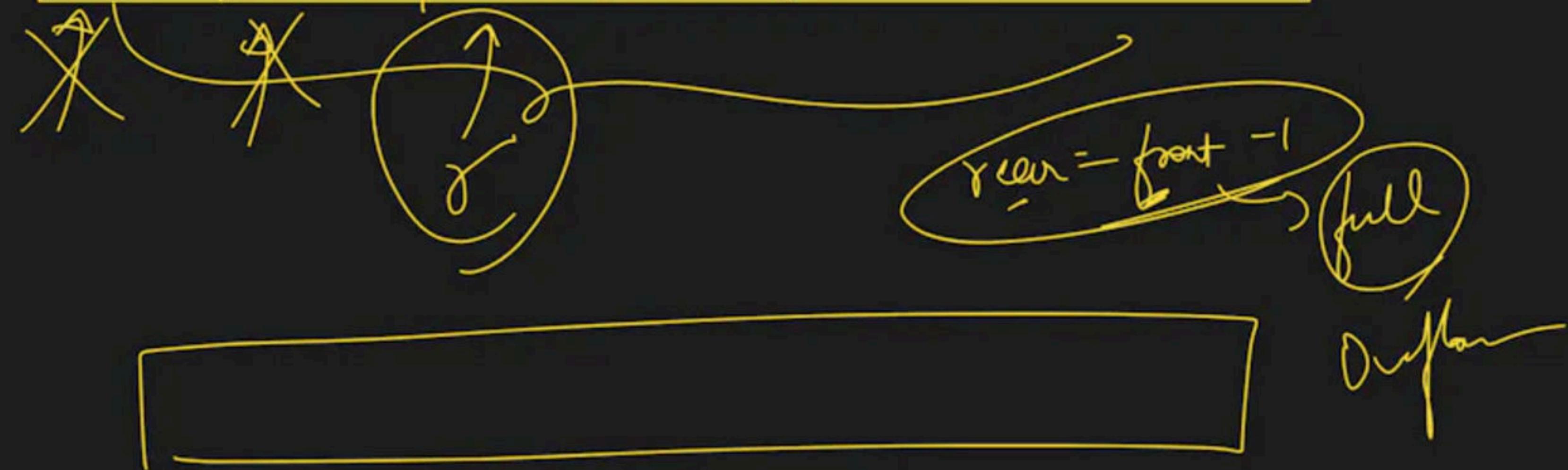
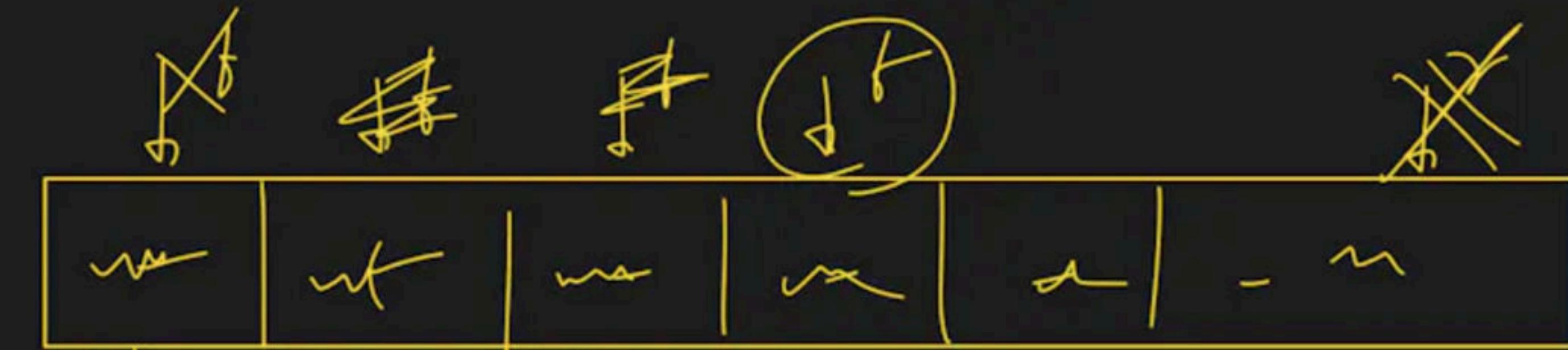




KMUSI







$0^{\circ}/.5^-$

$\theta$

$0^{\circ}/.5^- = 0$

$1^{\circ}/.5^- = 1$

$2^{\circ}/.5^- \rightarrow 2$

$3^{\circ}/.5^- \rightarrow 3$

$4^{\circ}/.5^- \rightarrow 4$

$5^{\circ}/.5^- \rightarrow 5$

$6^{\circ}/.5^- \rightarrow 6$

$7^{\circ}/.5^- \rightarrow 7$

$8^{\circ}/.5^- \rightarrow 8$

$9^{\circ}/.5^- \rightarrow 9$

$10^{\circ}/.5^- \rightarrow 0$

$11^{\circ}/.5^- \rightarrow 1$

$12^{\circ}/.5^- \rightarrow 2$

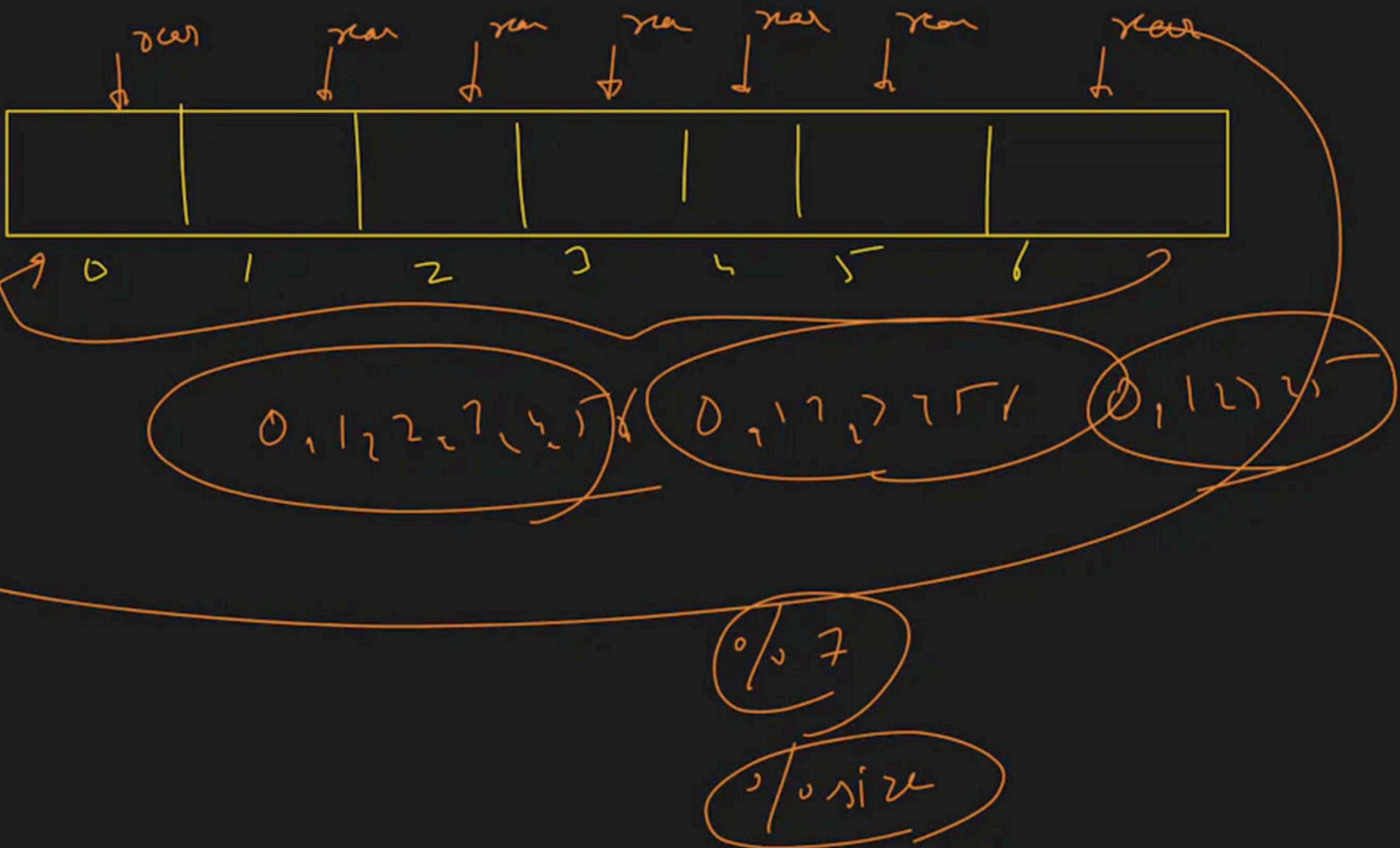
$13^{\circ}/.5^- \rightarrow 3$

$14^{\circ}/.5^- \rightarrow 4$

$15^{\circ}/.5^- \rightarrow 5$

$0, 1, 2, 3, \dots, (n-1)$

$[0, 1, 2, 3, \dots, (n-1)]$





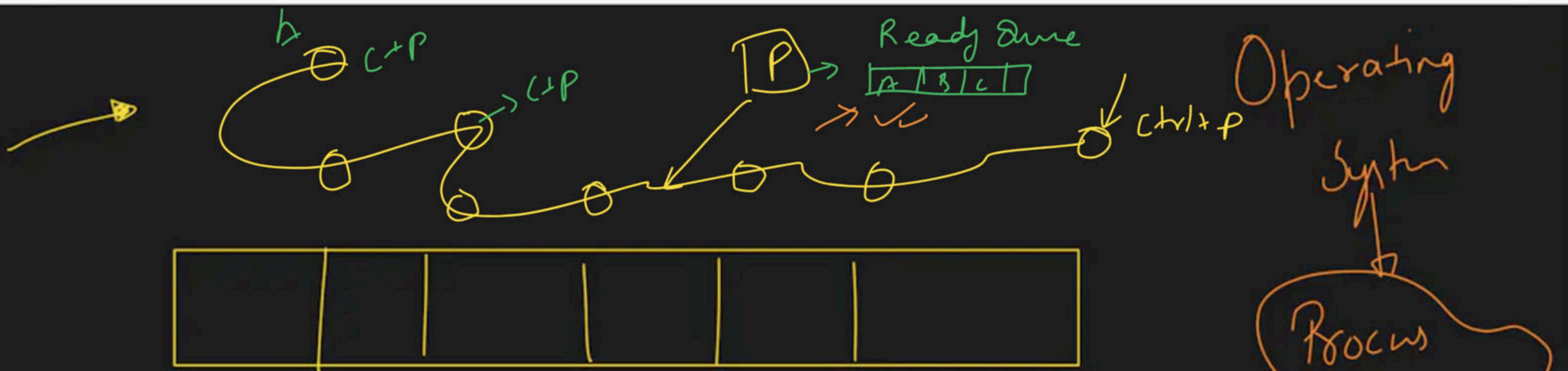
→ Doubly Ended Queue

push-front  
pop-front



deque

dequeue  
push-back( )  
pop-back( )



$\frac{600}{-F}$     120    110    100    10    20    200    500    1000

pull front (10)  
 pull (110)

R/B

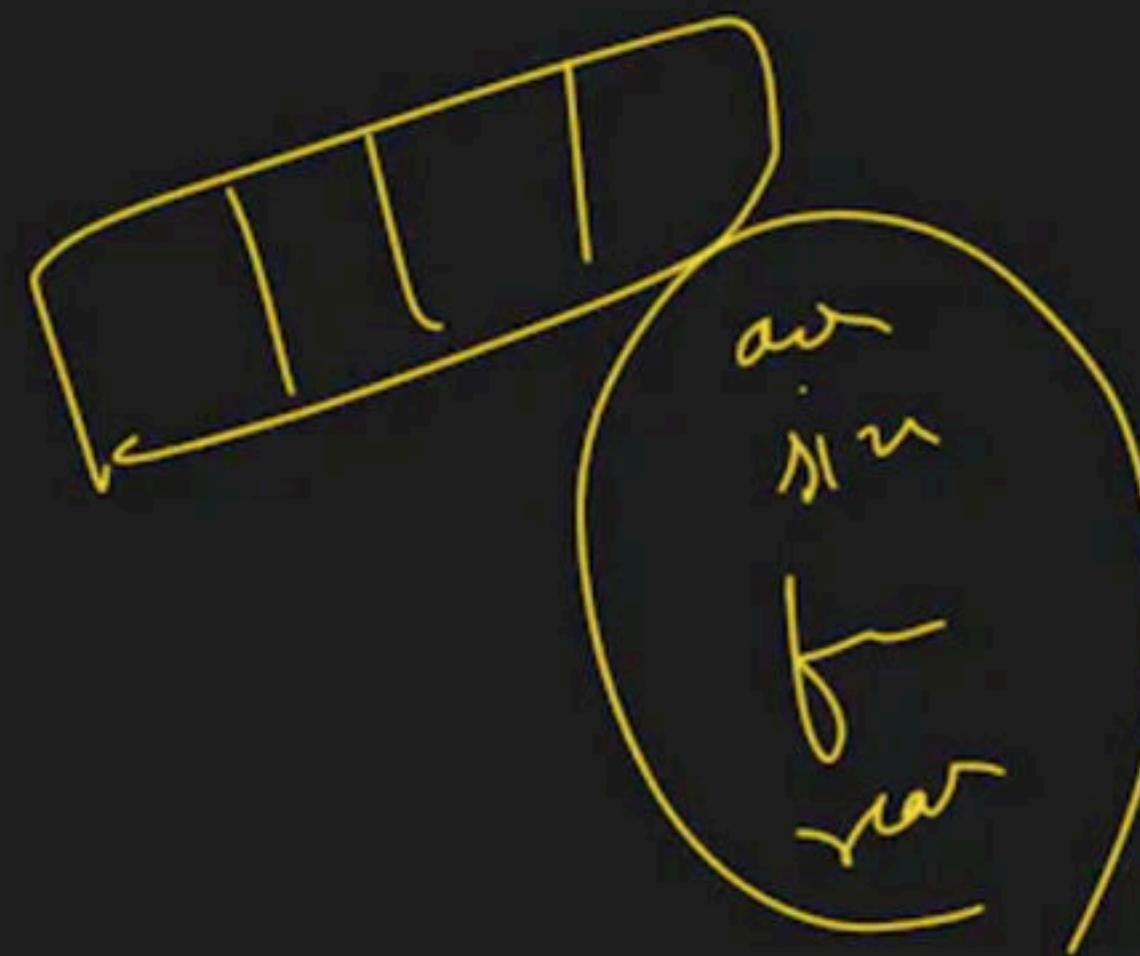
pull front (10)

pullback (20)

Jplant

Degm

Normal



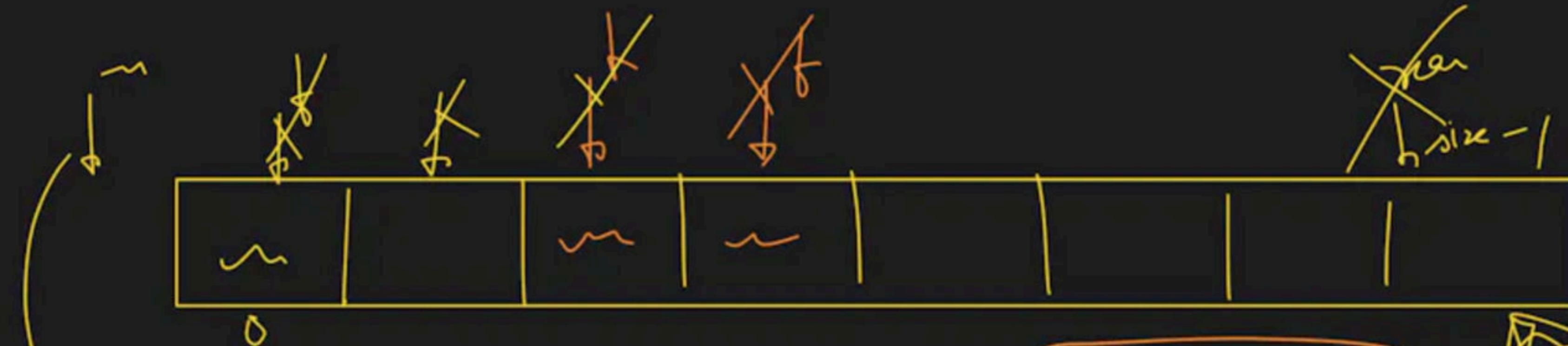
from

~~Scratch~~

→ pink foot  
→ pink back  
→ pop-foot  
→ pop-back

H/wr

Circular degm



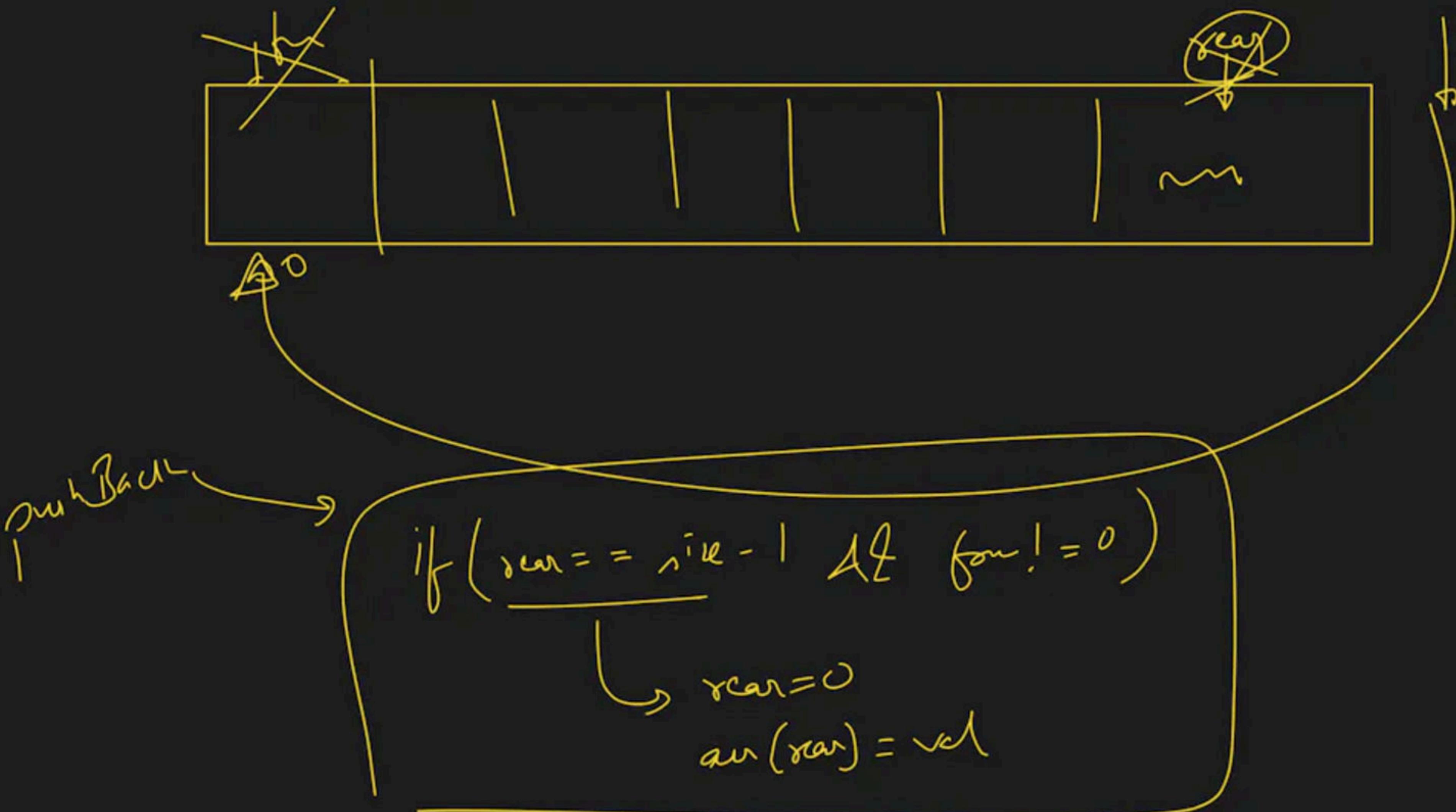
normal  
case  $\rightarrow$   $f = -$   
 $\text{arr}[f] = \text{val}$

$\text{front} = 0$   $\text{size} = 4$

$f = \text{size} - 1$

Diagram illustrating the state of the circular queue after an insertion:

- The array has 10 slots, indexed from 0 to 9.
- Slot 0 contains the character 'm'.
- Slot 4 contains the character 'm'.
- Slot 9 contains the value 4, labeled as  $\text{size} - 1$ .
- A pointer  $\text{front}$  points to index 0.
- A pointer  $f$  points to index 9, labeled as  $\text{size} - 1$ .
- A callout bubble labeled "Circular" points to the array.





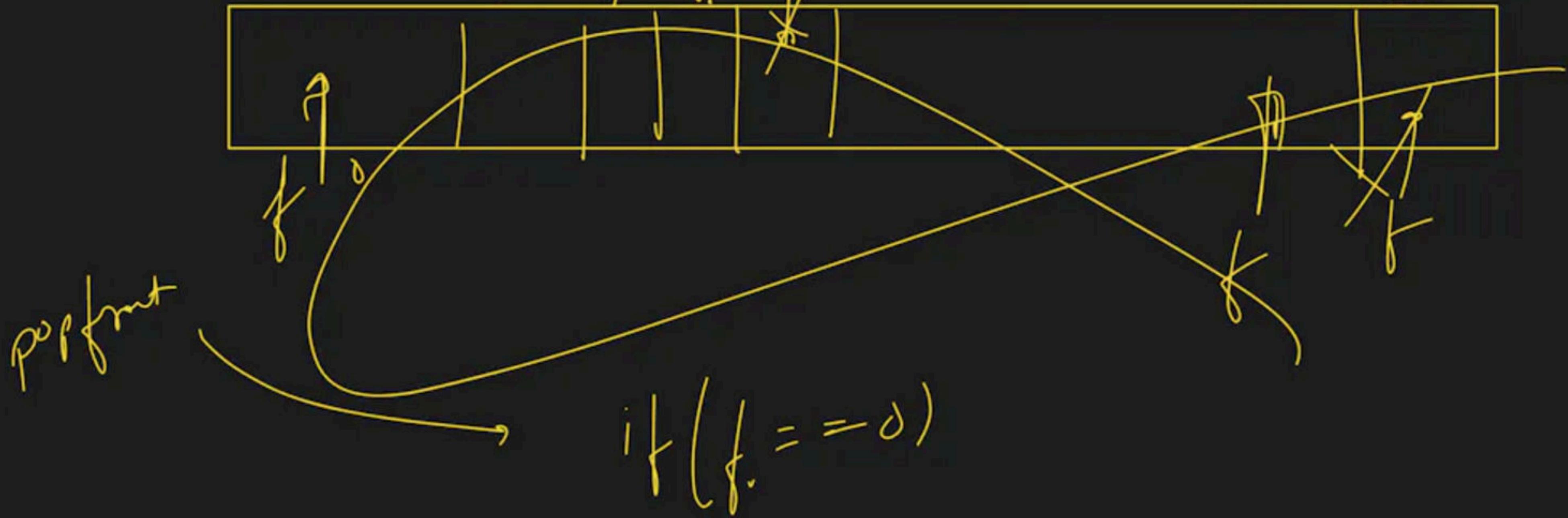
0

if ( $x_{car} == 0$ )

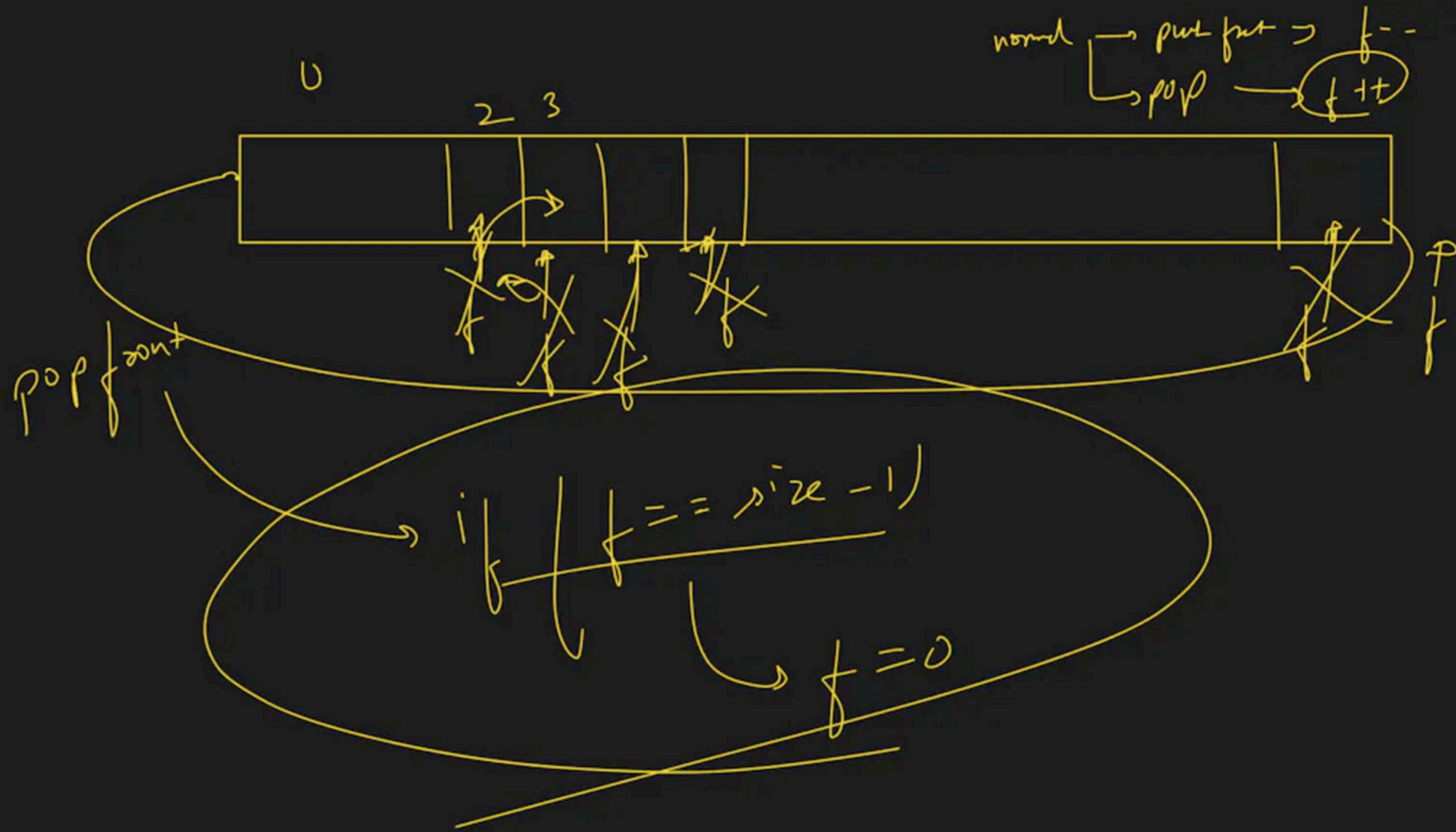
$x_{car} = size - 1$



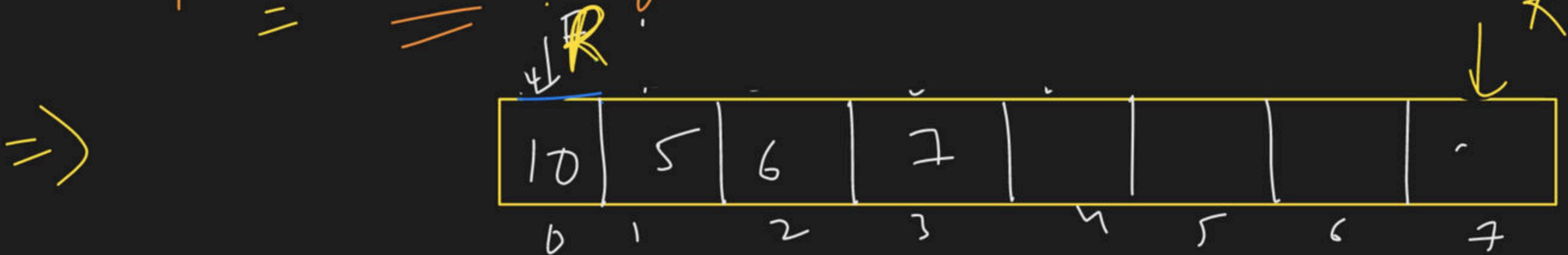
$\rho_{\text{NL}} \rightarrow f^{++}$   
 $\text{normal} \rightarrow \rho D \bar{D} \rightarrow f^{--}$



normal  $\rightarrow$  push first  $\rightarrow$   $t^-$   
 $\rightarrow$   $pop$   $\rightarrow$   $t^+$



Implement  
= = Dequeue



⇒ size = size  
is full => front == 0 & rear  
= size - 1

arry = new int [size]

front = -1

rear = -1

is Empty => front == -1 ; ~  
& rear == -1 ✓

insert front()

=

① Overflow => isFull()

② empty       $F = -1 \ \& \ R = -1$

$F++$ ,  $R++;$

$a[F] = val;$

③ Normal Case

$F--;$

$a[F] = val;$

④ Circular Case

( $front = 0 \ \& \ R[1] = s-1$ )

$front = s-1$ ,  $a[F] = val.$

⑦ Insert Last  
=

⑥ Is Full

⑦ empty  $\Rightarrow F = -1 \neq R = -1$

⑧ Circular =  $R == \text{size} - 1 \wedge F \neq 0$

⑨  $R++ \quad arr[R] = \text{val};$

③ del Front      ① isEmpty()

②  $F \equiv R$ , single element ~~Ans~~

$a[F] = -1$ ;  $F = R = -1$

③ Circular can  
= =

$F == size - 1$

$\Rightarrow F = 0$  ~~Ans~~

④ Normal  $\Rightarrow$

$a[F] = -j$ ;  $F++$

④ del Last()

=

② front = R

① is S

F = R - 1

③ R == 0, Circular S

R = S - 1;

④ a[R] = -1, R = -;

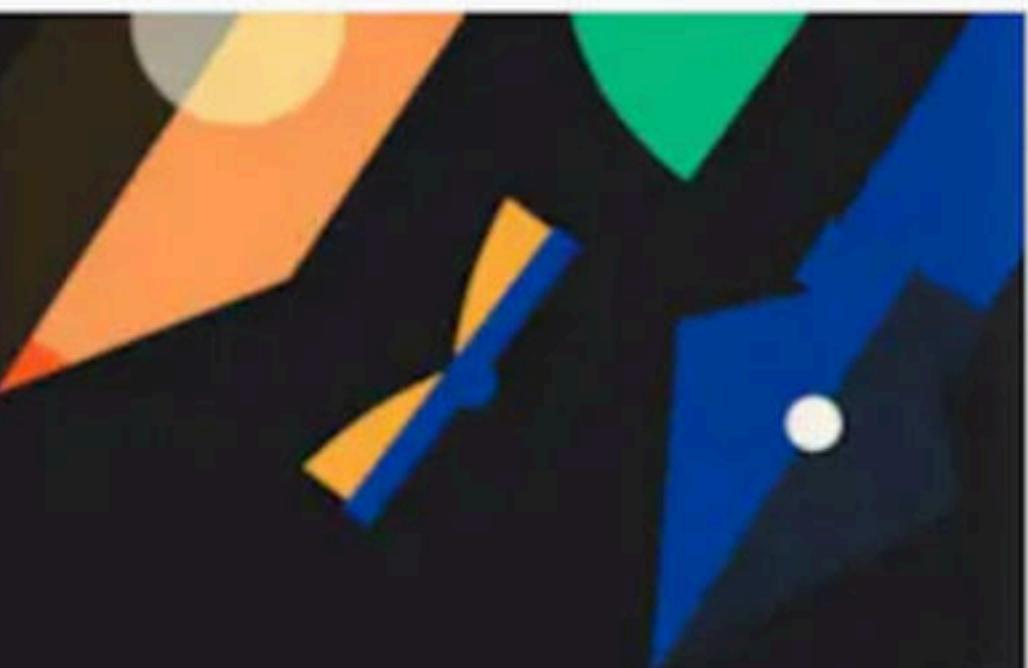
→











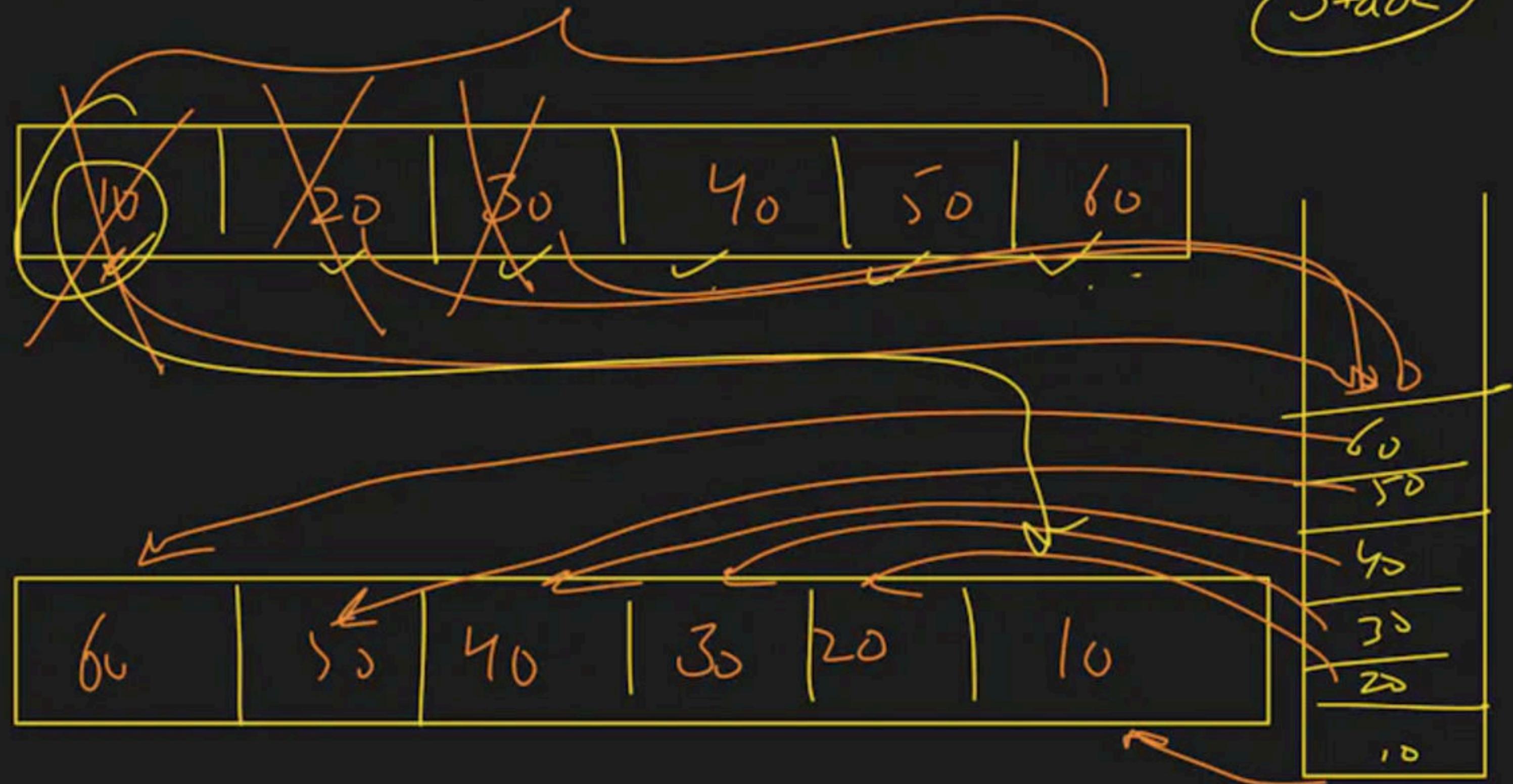
# Queue Class - 2

Special class

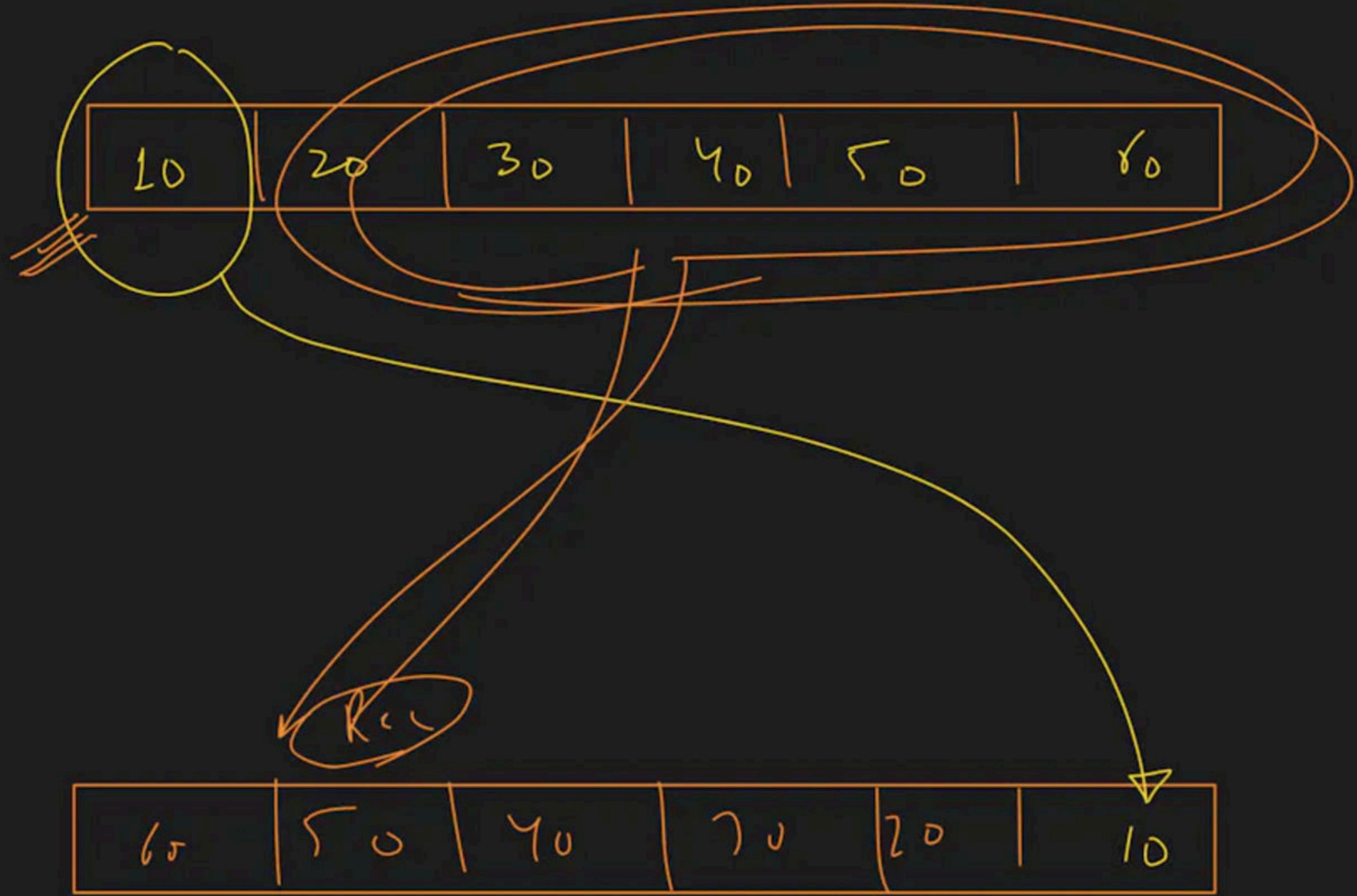
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→ Reverse a Queue

i/p →



o/p →



20	30	40	50	60
----	----	----	----	----

$\text{int } \& p = q \cdot \text{front}()$

$q \cdot \text{pop}()$

$\text{return } (q)$

$\& \cdot \text{push}( \& p )$

60	50	40	30	20
----	----	----	----	----

10



$R_{cc}$

```

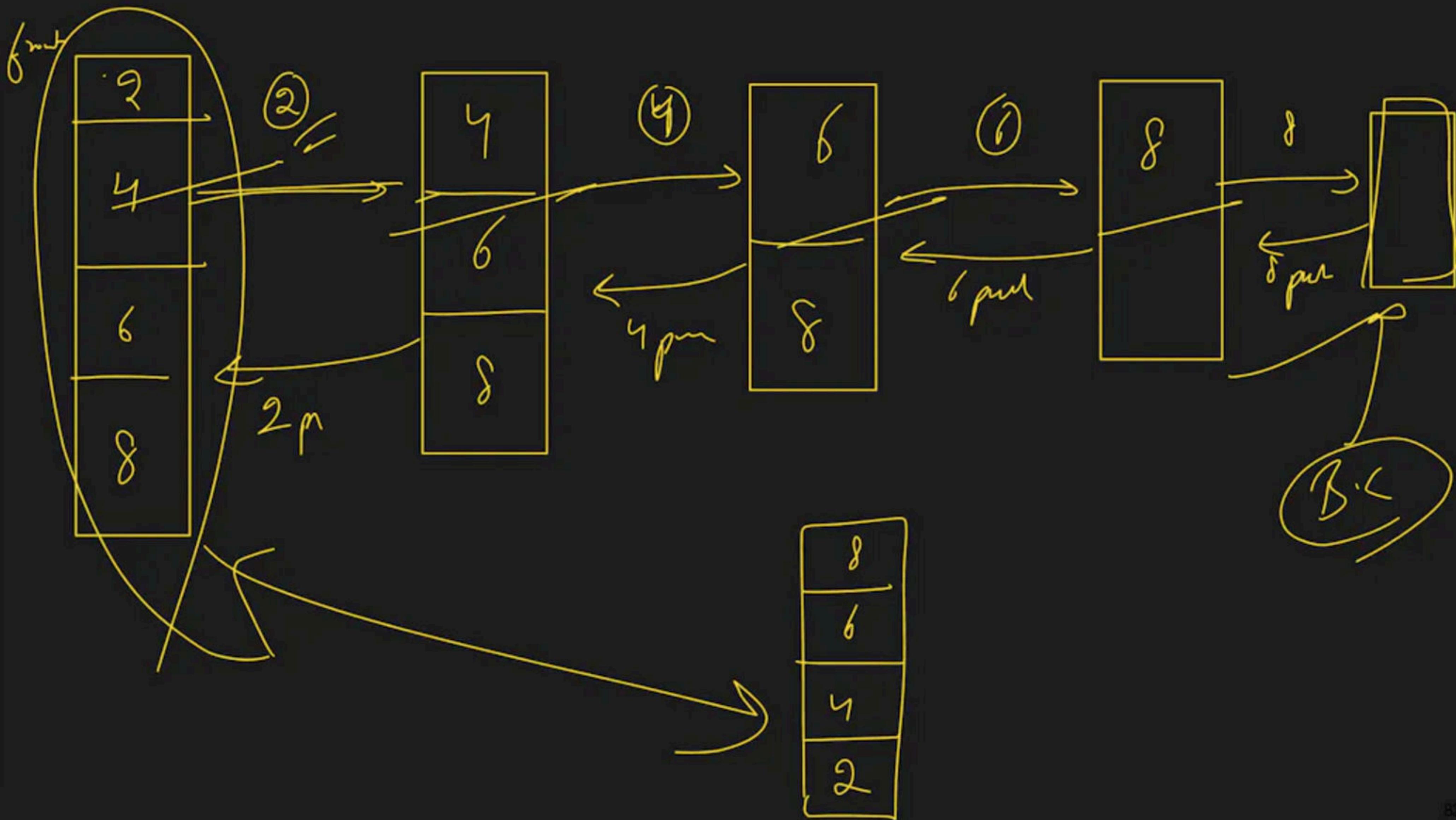
int element = q.front()
q.pop()
    
```

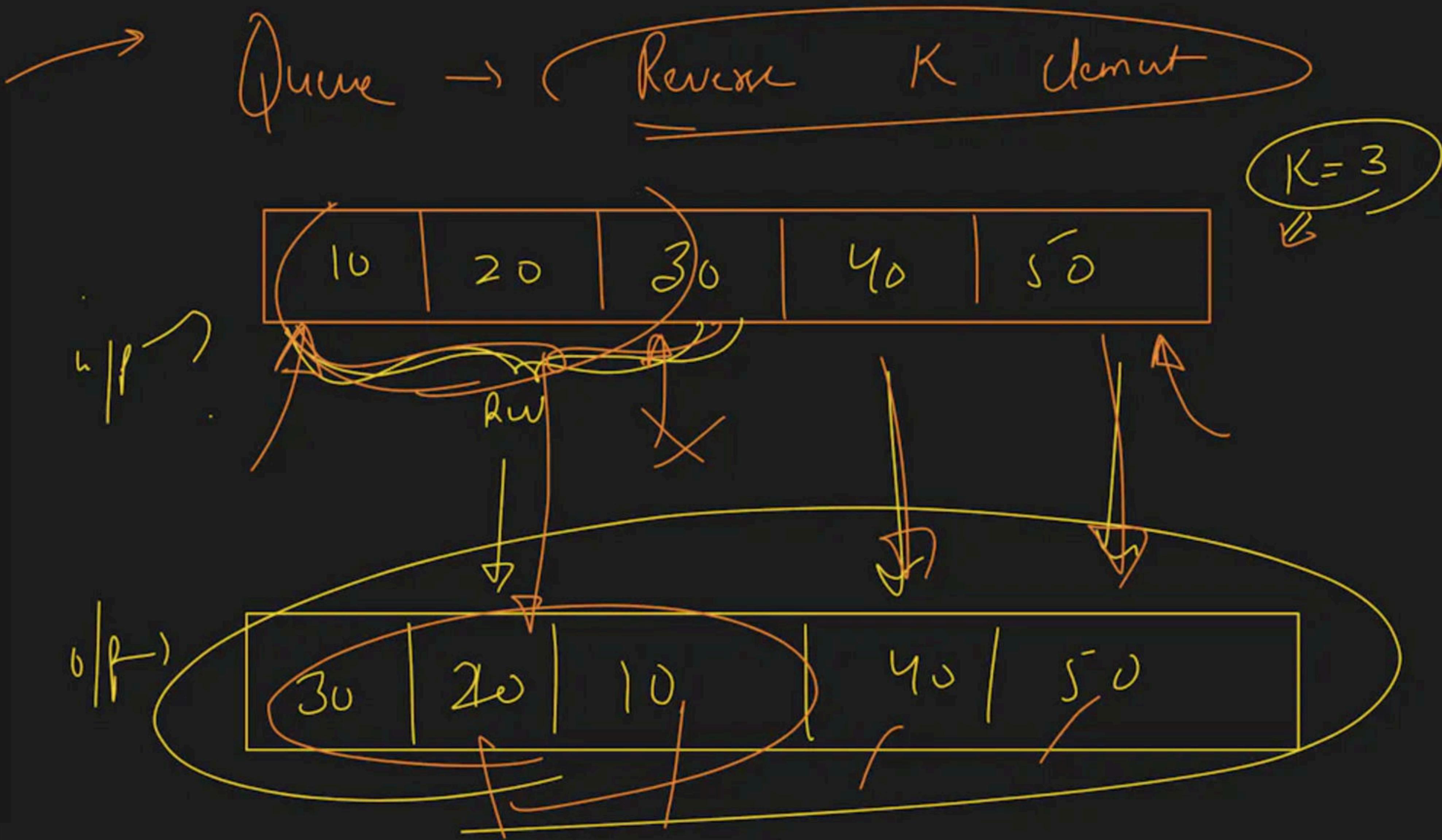
$q$   
reverse( $q$ )

$q.push(element)$



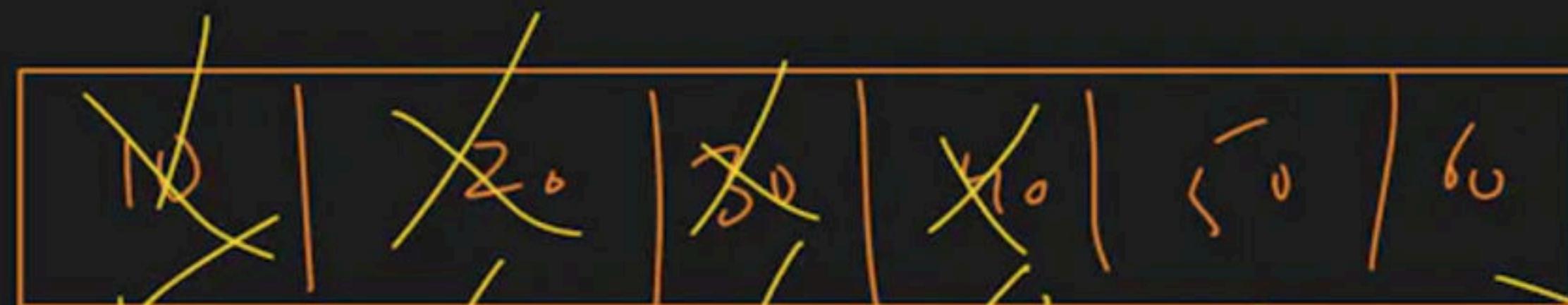
$element = \boxed{2}$





- #
- (1)  $\text{push first } k \text{ elms from } S \rightarrow Q$
  - (2)  $\text{push all } k \text{ elms from } S \rightarrow Q$
  - (3)  $Q \rightarrow \text{fin } k (n - k)$
- $\text{pop } p \leftarrow \text{pop } Q$

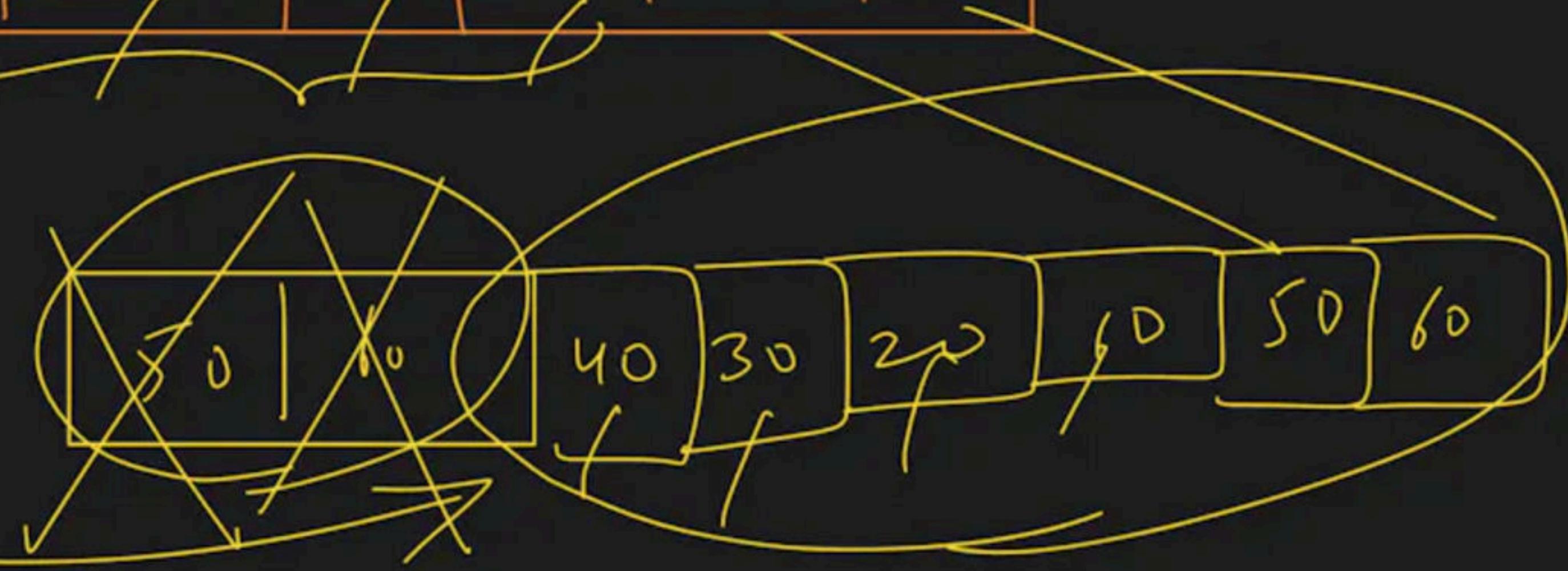
$k = \gamma$



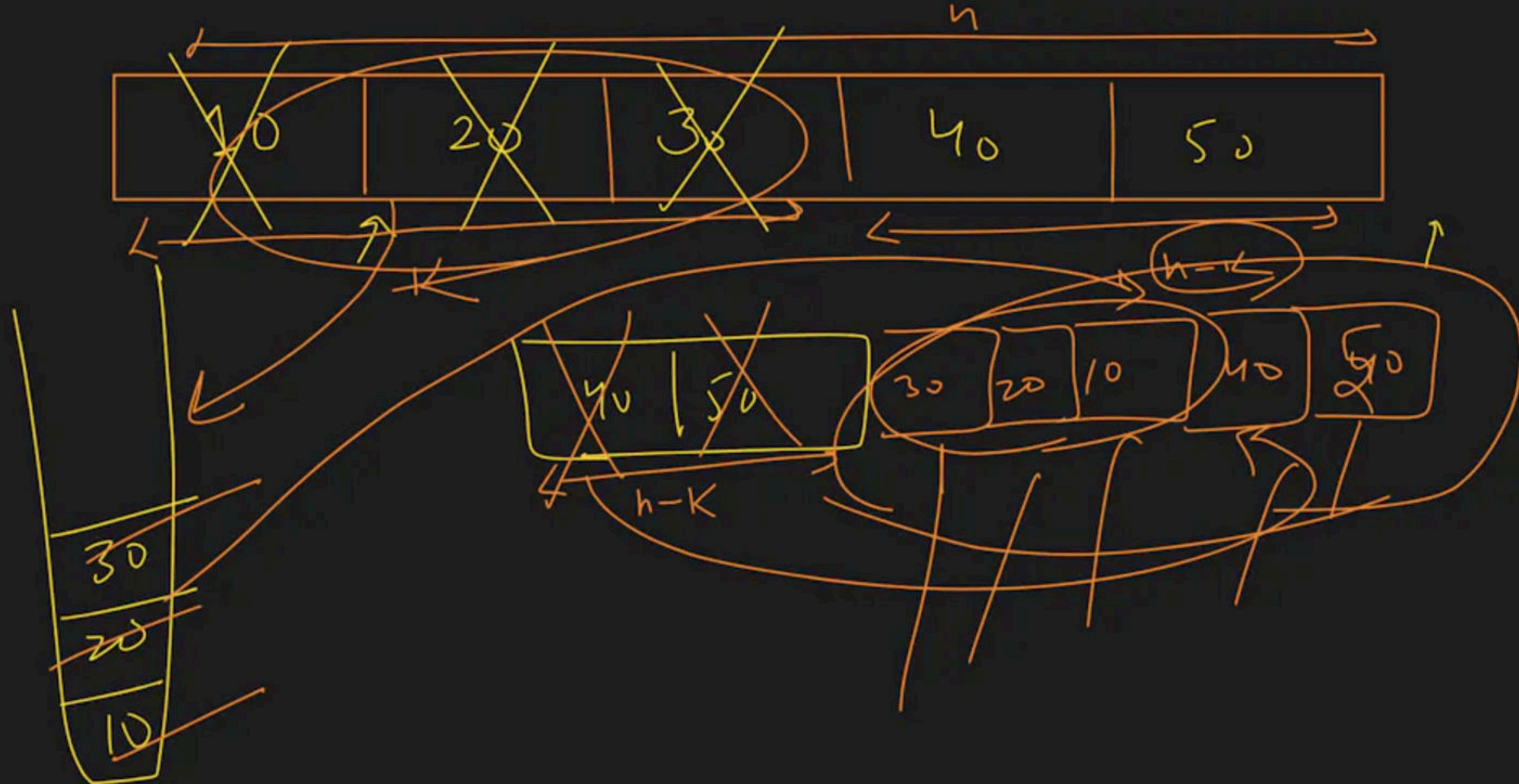
(I)



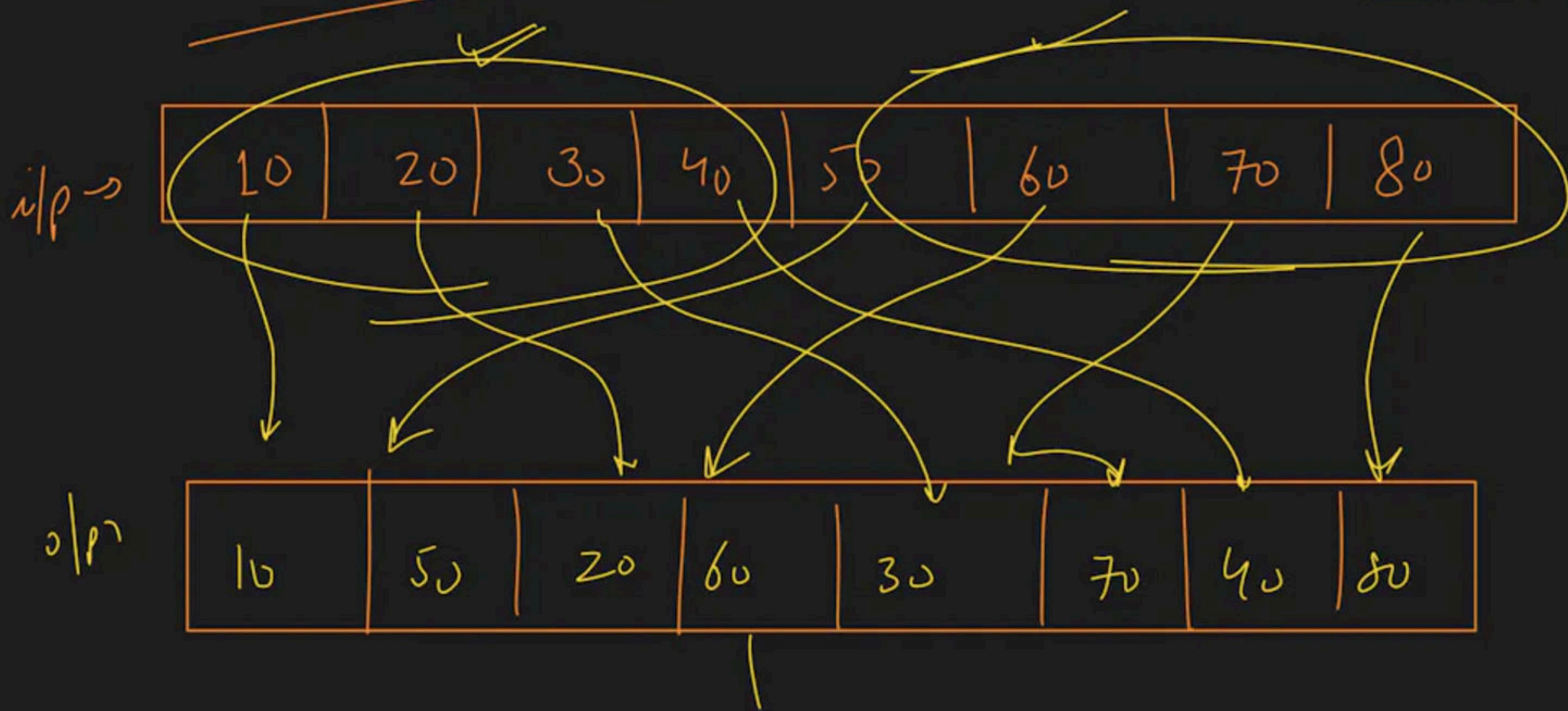
II

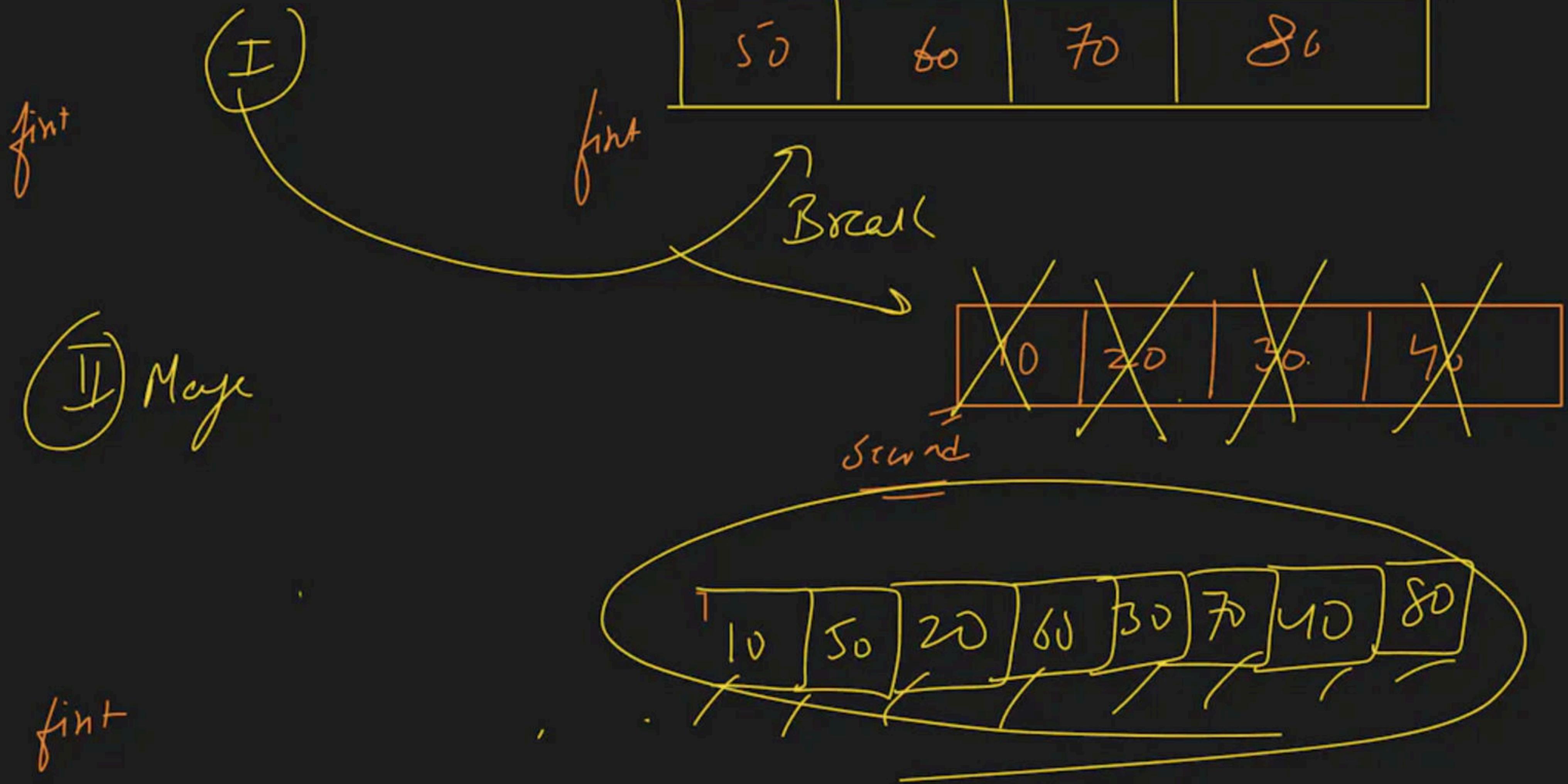


(III)  $(n-k) \rightarrow (-1)^k = 2$



$\rightarrow$  Interval *first* *second* *half of Queue*









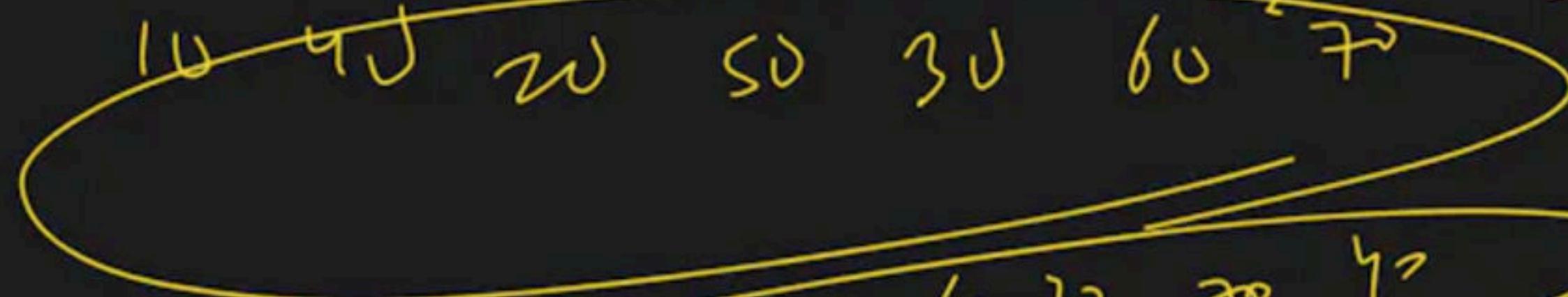
9d2

10	20	30	40	50	60	70
----	----	----	----	----	----	----

$$n = \frac{7}{2} = 3$$

$$n = \frac{7}{2} = 4$$

$$n = 7$$



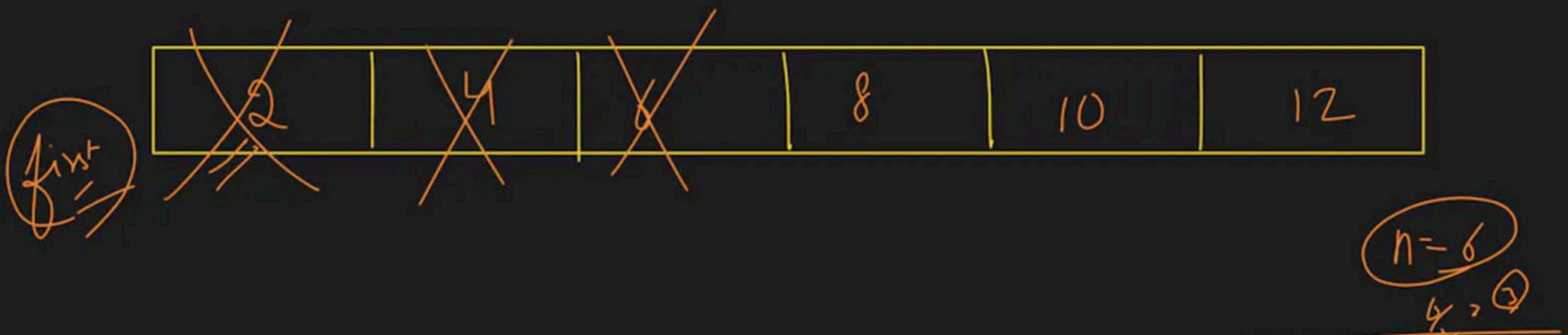
10 50 20 60 30 70

10	20	30	40	50	60	70
----	----	----	----	----	----	----

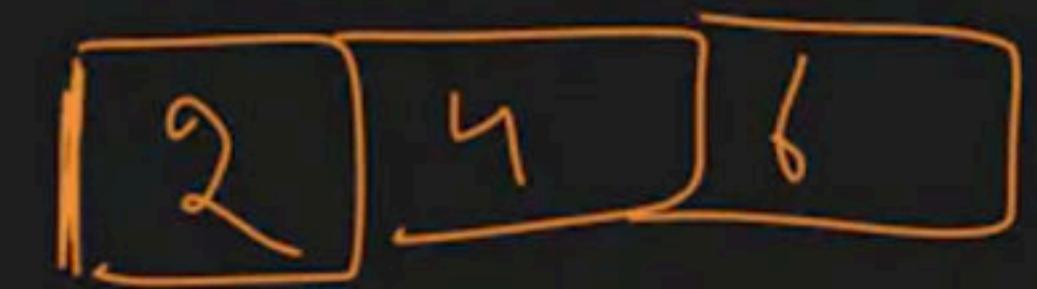
for  
y1  
y2  
if

$$\frac{n}{2} + 1$$

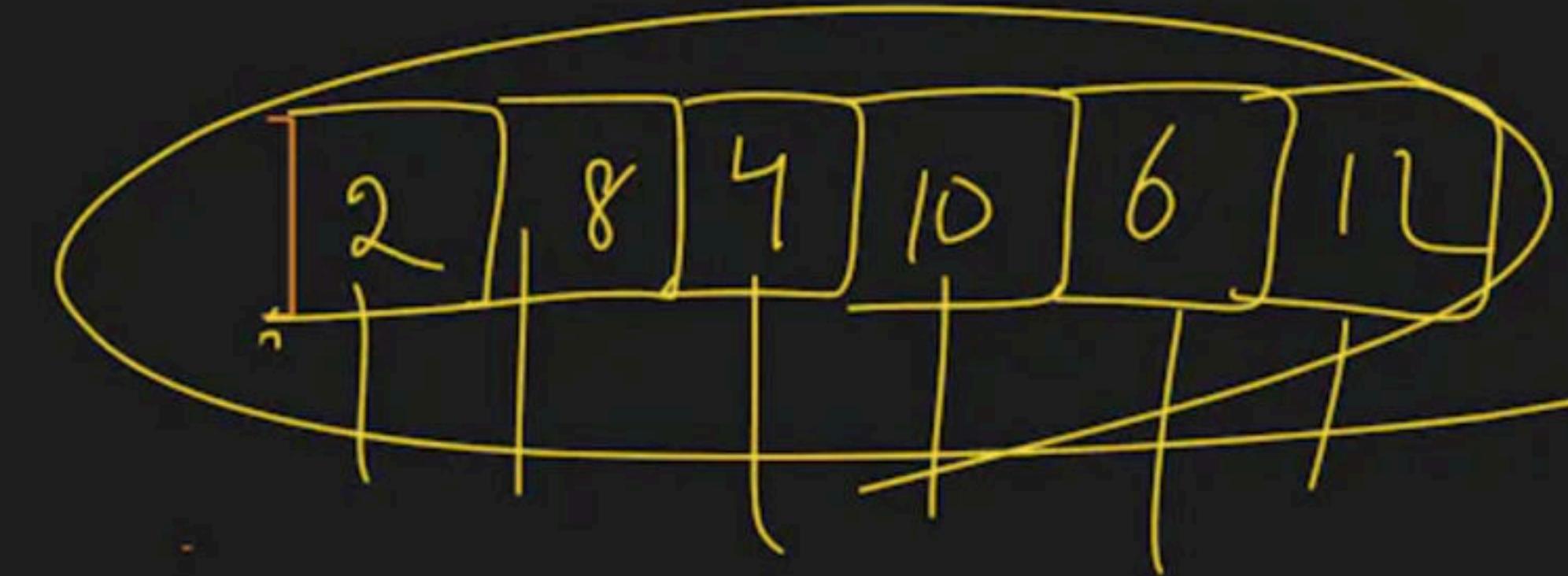
$$\frac{n}{2}$$



second →



```
for (i=0; i<n; i++)
{
    int top = first;
    first = port;
    second.push(top);
}
```



first

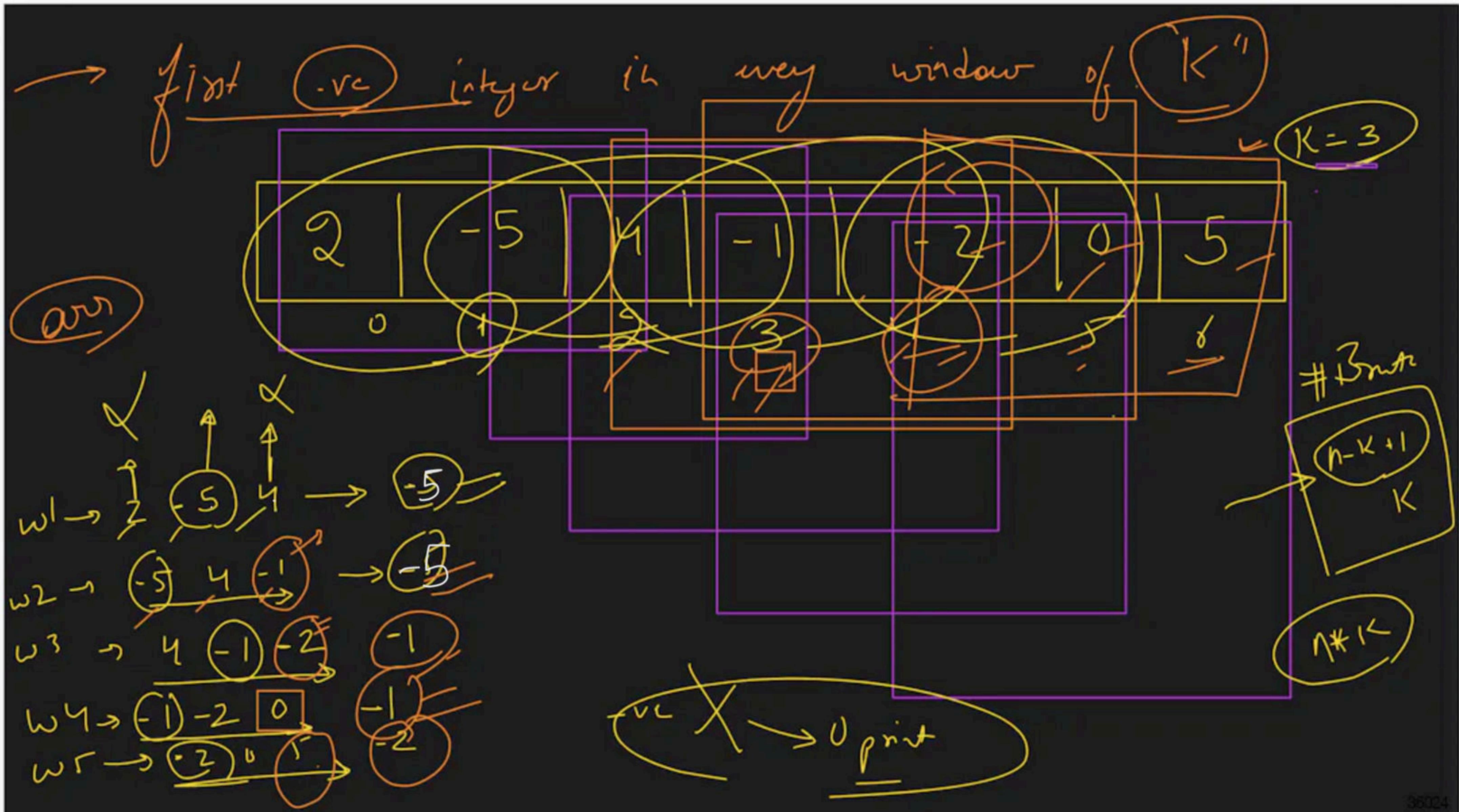
second

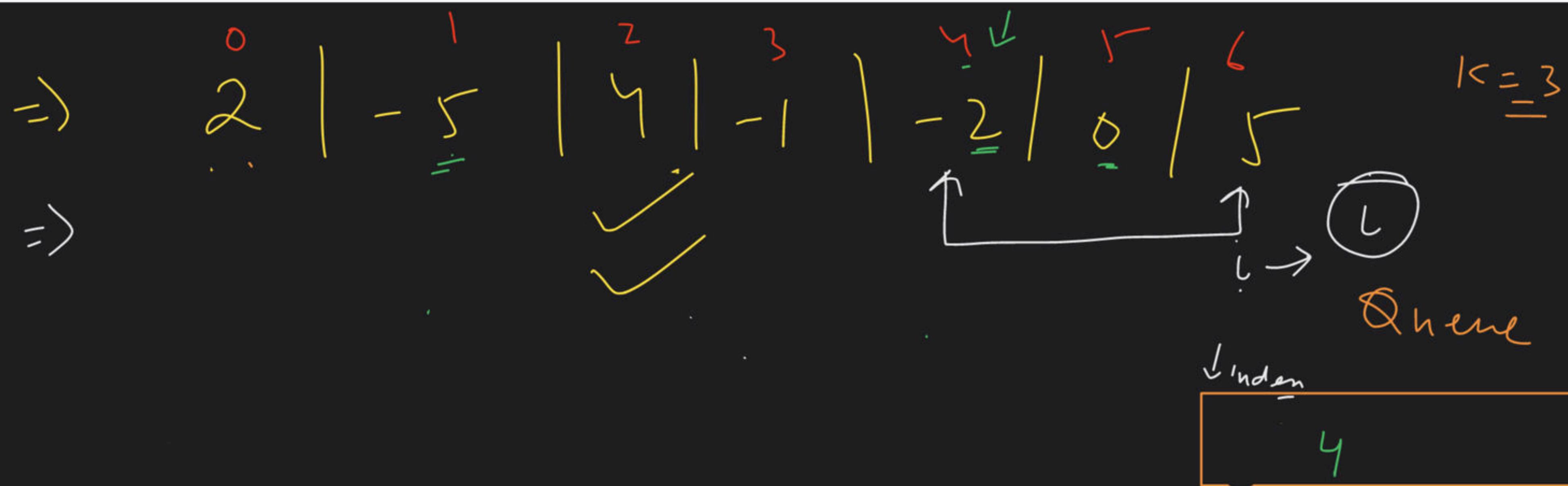


for ( $i=0; i < \frac{n}{2}; i++$ )

tmp =  $\underline{\text{second}} \cdot \text{front}()$   
 $\underline{\text{second}} \cdot \text{push}(\underline{\text{tmp}})$

tmp =  $\underline{\text{first}} \cdot \text{front}()$   
 $\underline{\text{first}} \cdot \text{push}(\underline{\text{tmp}})$





- ① -5
- ② -5
- ③ -1
- ④ -1
- ⑤ -2

$L = qF \geq K$   
 $\hookrightarrow P_{\text{upf}, \text{out}} =$   
 $q_k i \omega_0 = \text{Banyak}$   
 $\text{hai}$   
 $\text{rd window kaha!}$

⇒

c/a.out

main

Fatalk

R<sub>un</sub>

✓

|

int a;

✓ 2

int b;

Runk layers

• 3

— — —



Magic scan states

4

↳ all variables

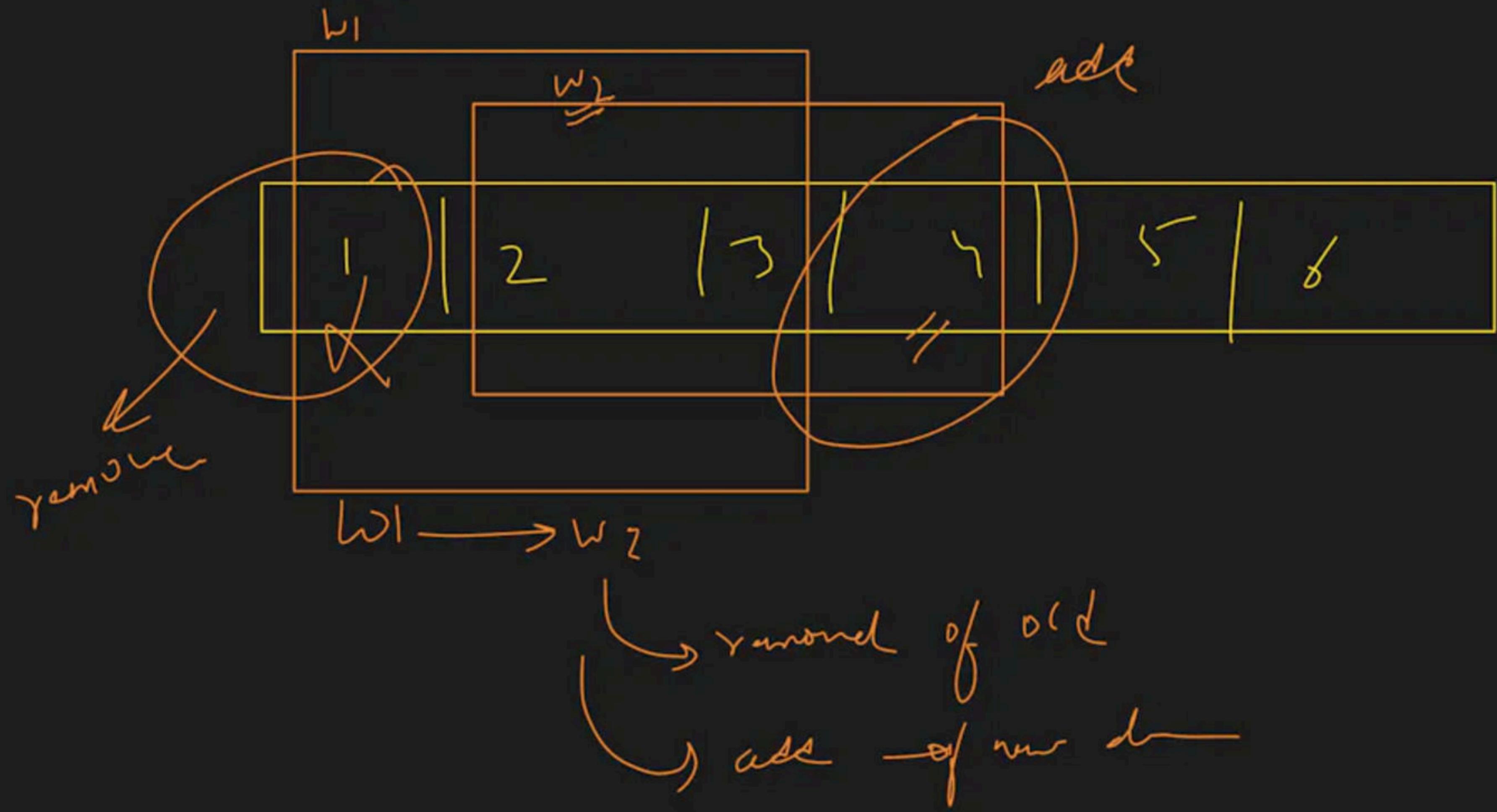
5

with values

6

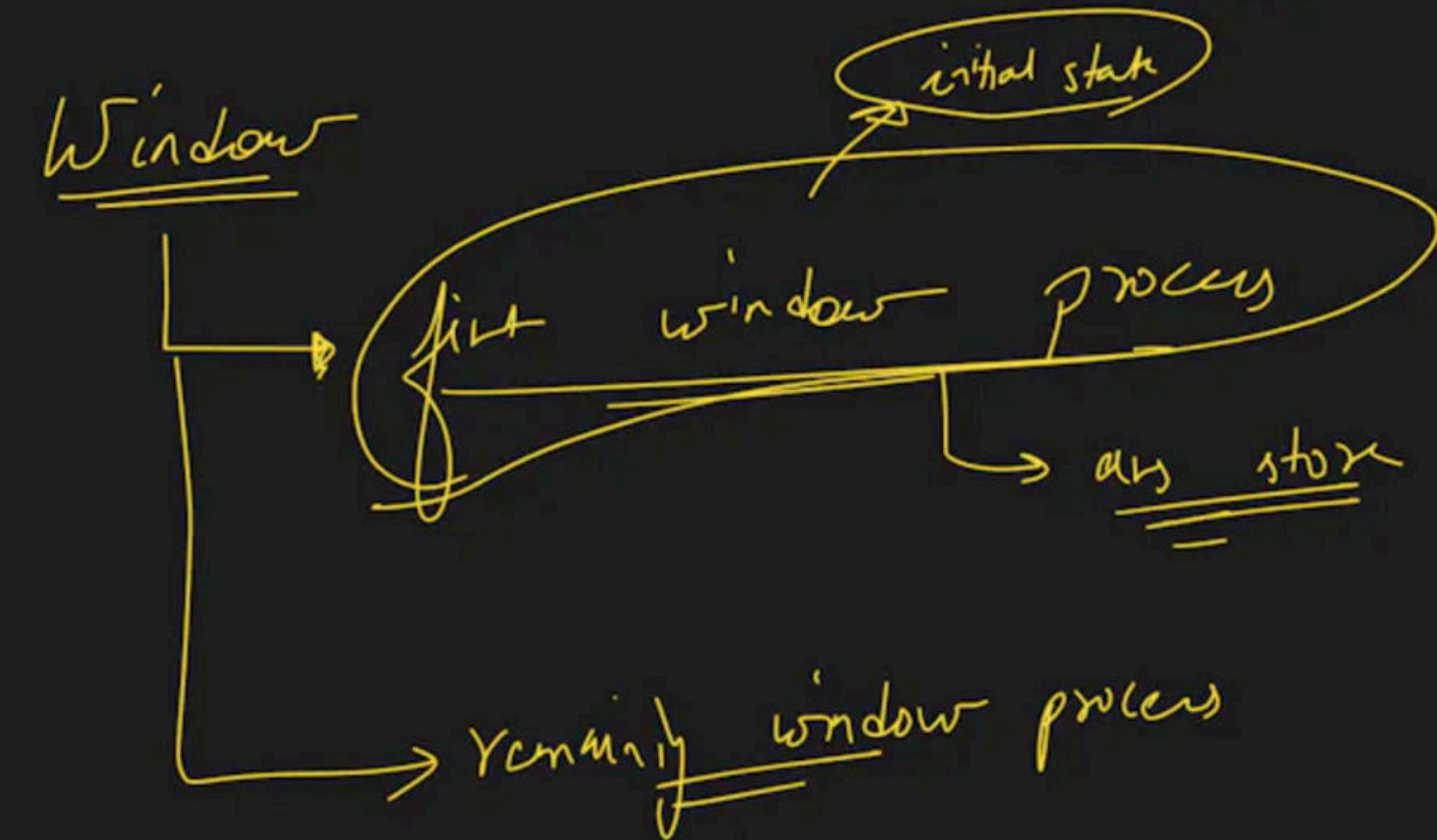
7

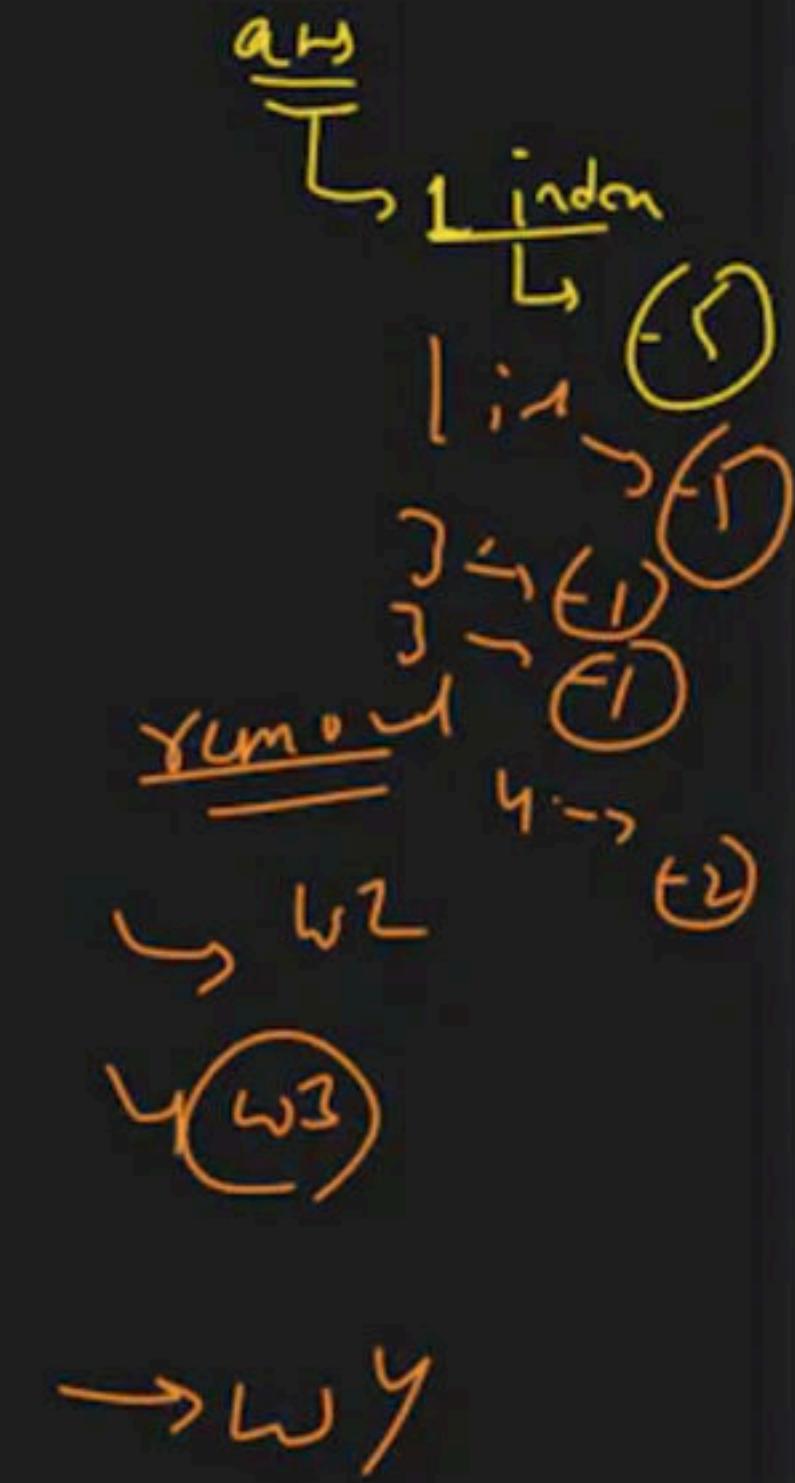
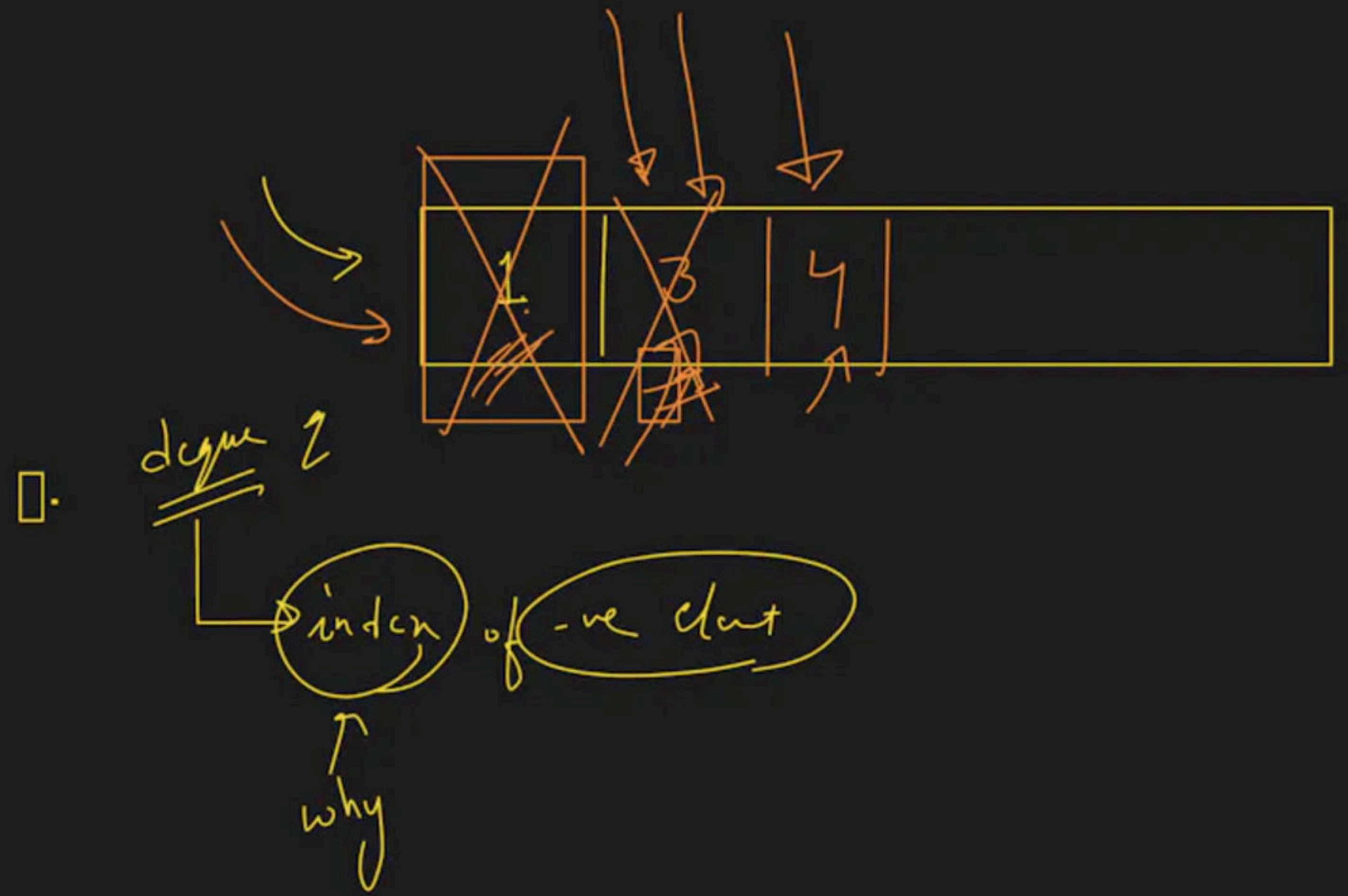




Sliding

dyne





s.w

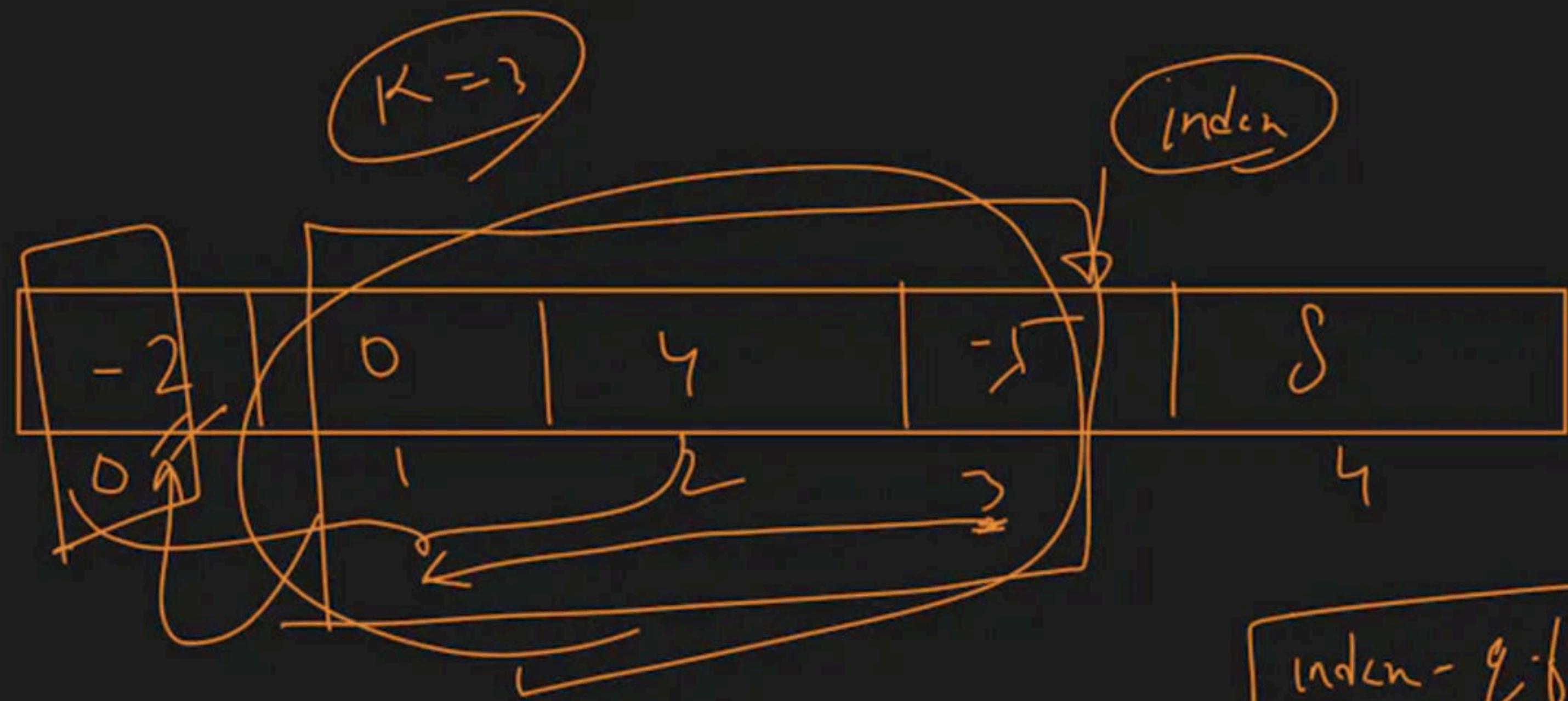
first  $\rightarrow K$  un  $\rightarrow$  first window

) removing

window

removed

) addition



$$\text{index} - \varrho \cdot \text{front} >= K$$

$$3 - 0 >= 3$$

$$\text{index} - \varrho \cdot \text{front} == K$$

$$\text{index} - \varrho \cdot \text{front} < K$$



$d_2$



