

# Assignment - 4

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1. Bidder 1 -  $b_1 \rightarrow s_1$  wallet

$b_2 \rightarrow s_2$

Prize :-  $s_1 + s_2$

one bidder quits and other bidder pays  $P$  & receives  $s_1 + s_2$

The equilibrium is the case where both players has to bid double price they have.  
(individual signal)

winner payment ( $P$ ) < maximum he should pay

A bidder who wins at price  $P$  knows the

actual value is  $v = \frac{s_1}{2} + \frac{P}{2} > P$

if  $P < 25$ ; So 1 is pleased to

be a winner at any price up to  $25$  but would

lose money if he "won" the auction at any

higher price. this is unique symmetric equilibrium.

if  $b_1$  cons at  $p$ , then  $v = \overset{S_1}{\cancel{b_1}} + \frac{9p}{10} > p \Leftrightarrow p < 10 \cancel{b_1} S_1$

if  $b_2$  cons at  $-p$  then  $v = \underset{S_2}{\cancel{b_2}} + \frac{p}{10} > p \Leftrightarrow p < \frac{10}{9} \cancel{b_2} S_2$

2. Analyze the allocation algorithm for downward sloping valuation for following  $v_i(k)$ , values for buyer  $i$ .

$$v_i(k) = v_{i1} + v_{i2} + \dots + v_{ik}$$

$v_i(k)$	0	1	2	3	4	5	6	7	8	
1	0	50	100	140	180	210	230	240	250	$\rightarrow v_i(k)$
2	0	50	50	40	40	30	20	10	10	
3	0	70	130	181	216	249	279	304	314	
	70	60	51	35	33	30	25	10		
	0	60	110	150	190	225	251	276	301	
	60	50	40	40	35	26	25	25		

$$n=3, m=8$$

$$v_i(m_i) - v_i(m_i-1) \geq p > v_i(m_i+1) - v_i(m_i)$$

if  $\sum m_i > m$  then  $p$  is too low

$\sum m_i < m$  then  $p$  is too high

otherwise we have found the right  $p$



cleaning price in  $[0, v]$  is  $P$

$P$	$m_1$	$m_2$	$m_3$	$\Sigma$
10	8	8	8	24
20	6	7	8	21
30	5	6	5	16
40	4	4	4	12
50	2	3	2	7
60	0	2	1	3
70	0	1	0	1

we calculate the  $m_1, m_2, m_3$  by

if  $P=10$   $i=1$

$$50-0 \geq 10 > (100-50) \times$$

$$(140-100) < 10 \leq (100-50) \times$$

$$140-100 \geq 10 > (180-140) \times$$

$$(210-180) \geq 10 \leq (180-140) \times$$

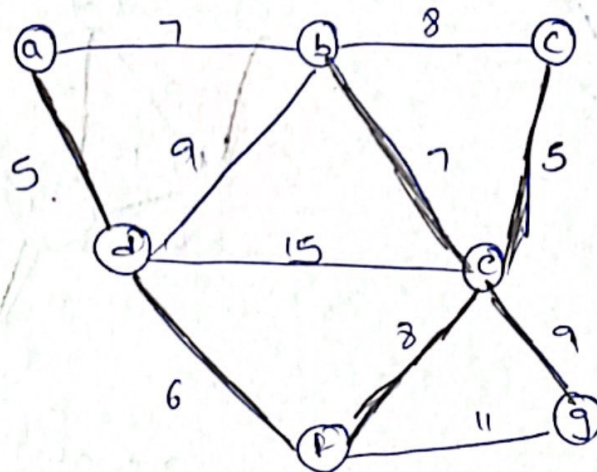
$$210-180 \geq 10 > 230-210 \times$$

$$240-230 < 10 \leq 230-210 \times$$

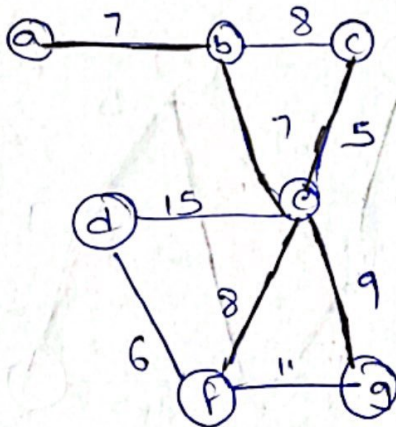
$$250-240 \geq 10 > 240-30 \times$$

so it's 8

3. Minimum Spanning tree

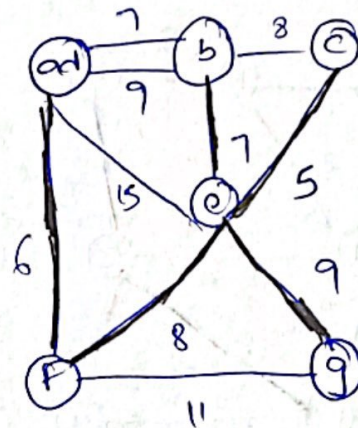


AD MST



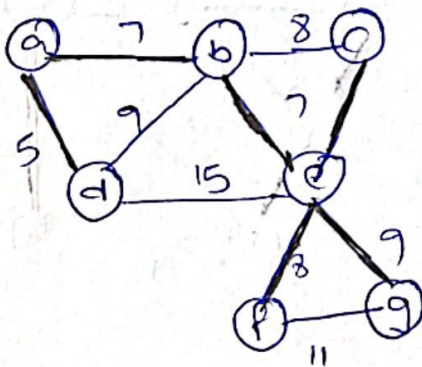
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Payments

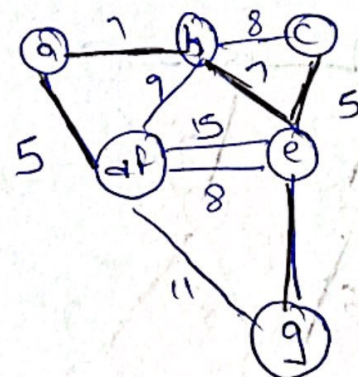


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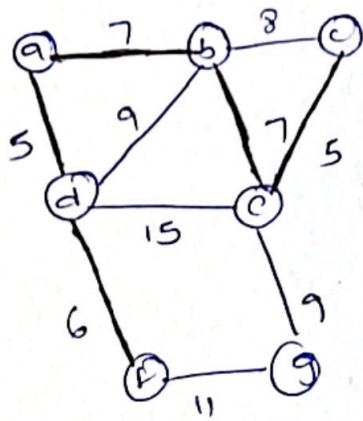
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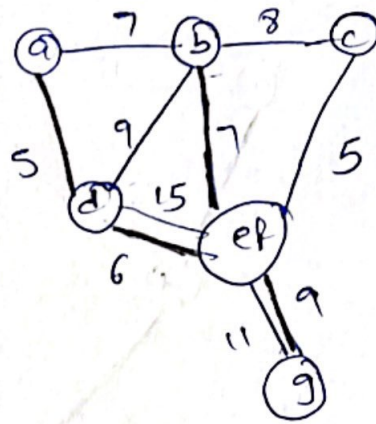
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PE :

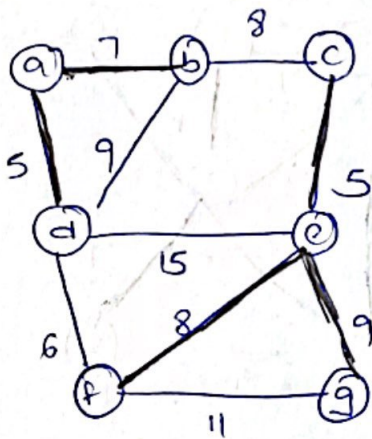


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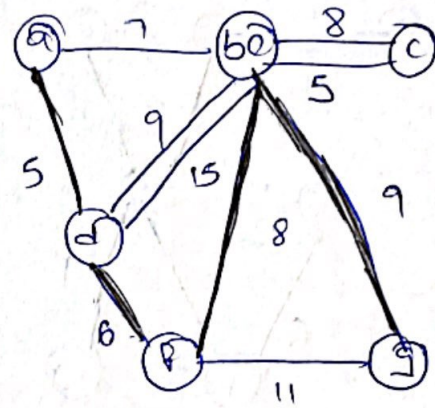


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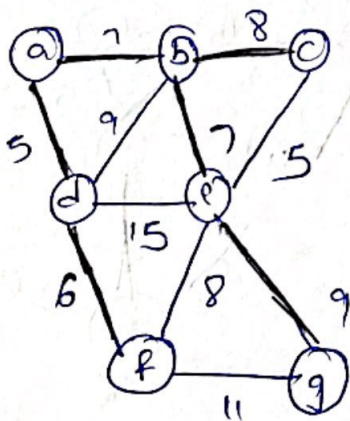


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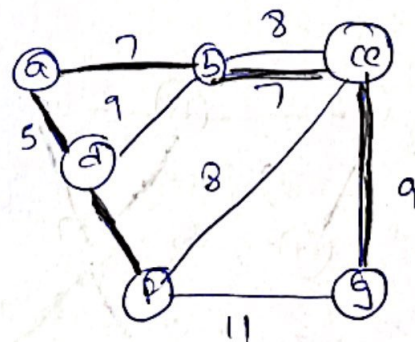


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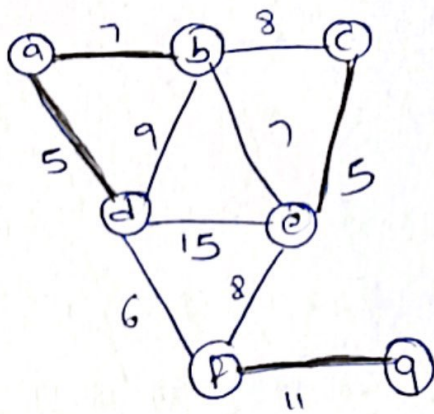


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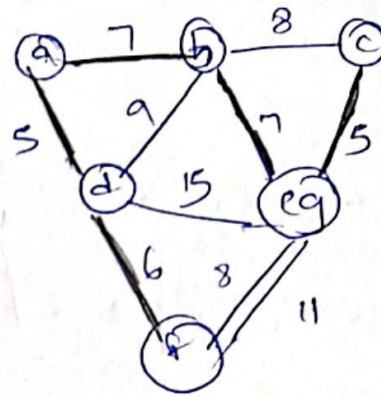


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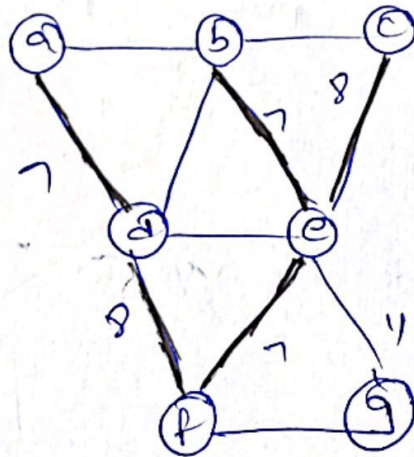
EG :-



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is the VCG payments for the minimum spanning tree for the above graph.

4.

i	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
b	10	9	8	6	6	5	5	5	5	5	5	4	4	2	2	2	2	1	1	1
Revenue	10	18	24	24	30	30	35	40	45	50	55	48	52	28	30	32	34	18	19	20

Conceptually order bids descending for obtaining optimal

fixed-price and maximize price to fit first  $i$  bidders

11  
5  
55 is the point where the revenues goes down