

ELEC6027 - VLSI Design Project : Programmers Guide

Team R4

27th April, 2014

Todo list

■ HSL - this is a bit too short. Surely there is more to say about it? .	6
■ HSL - this doesn't sound right. Maybe "transfer of program flow". Not sure on the use of the word "control" but i know it is the technical term	6
■ Maybe change to IEEE symbols if we have time, AJR: we still have the eagle d-types but I think it would look a bit messy	56
■ A register window could also be done for this section too	58
■ Sim.py needs a fair bit of change. If we have time, this could be altered to use Iains dir structure. This section highlights how to use his script and our assembler.	58
■ Make these more accurate when AJR has finished playing around .	60
■ how to run scan path sim too	60

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1 Introduction

Lorem Ipsum. . .

2 Architecture

Lorem Ipsum. . .

3 Register Description

Lorem Ipsum. . .

4 Instruction Set

The complete instruction set architecture includes a number of instructions for performing calculations on data, memory access, transfer of control within a program and interrupt handling.

HSL - this doesn't sound right. Maybe "transfer of program flow". Not sure on the use of the word "control" but i know it is the technical term

HSL - this is a bit too short. Surely there is more to say about it?

All instructions implemented by this architecture fall into one of 6 groups, categorized as follows:

- Data Manipulation - Arithmetic, Logical, Shifting
- Byte Immediate - Arithmetic, Byte Load
- Data Transfer - Memory Access
- Control Transfer - (Un)conditional Branching
- Stack Operations - Push, Pop
- Interrupts - Enabling, Status Storage, Returning

There is only one addressing mode associated with each instruction, generally following these groupings:

- Data Manipulation - Register-Register, Register-Immediate
- Byte Immediate - Register-Immediate
- Data Transfer - Base Plus Offset
- Control Transfer - PC Relative, Register-Indirect, Base Plus Offset
- Stack Operations - Register-Indirect Preincrement/Postdecrement
- Interrupts - Register-Indirect Preincrement/Postdecrement

4.1 General Instruction Formatting

Instruction Type		Sub-Type	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
A1	Data Manipulation	Register	Opcode						Rd		Ra		Rb		X		X	
A2		Immediate							Rd		Ra		imm4/5					
B	Byte Immediate		Opcode						Rd		imm8							
C	Data Transfer		0	LS	0	0	0	Rd		Ra		imm5						
D1	Control Transfer	Others	1 1 1 1 0						Cond.		imm8							
D2		Jump									Ra		imm5					
E	Stack Operations		0	U	0	0	1	L	X	X	Ra		0	0	0	0	1	
F	Interrupts		1	1	0	0	1	ICond.			1	1	1	X	X	X	X	X

Instruction Field Definitions

Opcode: Operation code as defined for each instruction

Rd: Destination Register

Ra: Source register 1

Rb: Source register 2

immN: Immediate value of length N

Cond.: Branching condition code as defined for branch instructions

ICond.: Interrupt instruction code as defined for interrupt instructions

LS: 0=Load Data, 1=Store Data

U: 1=PUSH, 0=POP

L: 1=Use Link Register, 0=Use GPR

Pseudocode Notation

Symbol	Meaning
\leftarrow, \rightarrow	Assignment
Result[x]	Bit x of result
Ra[$x : y$]	Bit range from x to y of register Ra
$+Ra$	Positive value in Register Ra
$-Ra$	Negative value in Register Ra
$<$	Numerically greater than
$>$	Numerically less than
$<<$	Logical shift left
$>>$	Logical shift right
$>>>$	arithmetic shift right
Mem[val]	Data at memory location with address val
$\{x, y\}$	Contatenation of x and y to form a 16-bit value
$(cond)?$	Operation performed if $cond$ evaluates to true
$!$	Bitwise Negation

Use of the word UNPREDICTABLE indicates that the resultant flag value after operation execution will not be indicative of the ALU result. Instead its value will correspond to the result of an undefined arithmetic operation and as such should not be used.

4.2 ADD

Add Word

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	1	0	Rd			Ra			Rb			X	X

Syntax

ADD Rd, Ra, Rb

eg. ADD R5, R3, R2

Operation

$Rd \leftarrow Ra + Rb$

$N \leftarrow \text{if } (\text{Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if } (\text{Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{if } (Ra > 0 \text{ and } Rb > 0 \text{ and } \text{Result} < 0) \text{ or}$
 $(Ra < 0 \text{ and } Rb < 0 \text{ and } \text{Result} > 0) \text{ then } 1, \text{ else } 0$

$C \leftarrow \text{if } (\text{Result} > 2^{16} - 1) \text{ or}$
 $(\text{Result} < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The 16-bit word in GPR[Ra] is added to the 16-bit word in GPR[Rb] and the result is placed into GPR[Rd].

Addressing Mode: Register-Register.

4.3 ADDI

Add Immediate

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	0	Rd			Ra			imm5				

Syntax

ADDI Rd, Ra, #imm5

eg. ADDI R5, R3, #7

Operation

$Rd \leftarrow Ra + \#imm5$

$N \leftarrow \text{if } (\text{Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if } (\text{Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{if } (Ra > 0 \text{ and } \#imm5 > 0 \text{ and } \text{Result} < 0) \text{ or}$

$(Ra < 0 \text{ and } \#imm5 < 0 \text{ and } \text{Result} > 0) \text{ then } 1, \text{ else } 0$

$C \leftarrow \text{if } (\text{Result} > 2^{16} - 1) \text{ or}$

$(\text{Result} < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The 16-bit word in GPR[Ra] is added to the sign-extended 5-bit value given in the instruction and the result is placed into GPR[Rd].

Addressing Mode: Register-Immediate.

4.4 ADDIB

Add Immediate Byte

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	1	1	Rd			imm8							

Syntax

ADDIB Rd, #imm8

eg. ADDIB R5, #93

Operation

$Rd \leftarrow Rd + \#imm8$

$N \leftarrow \text{if } (\text{Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if } (\text{Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{if } (Rd > 0 \text{ and } \#imm8 > 0 \text{ and } \text{Result} < 0) \text{ or}$

$(Rd < 0 \text{ and } \#imm8 < 0 \text{ and } \text{Result} > 0) \text{ then } 1, \text{ else } 0$

$C \leftarrow \text{if } (\text{Result} > 2^{16} - 1) \text{ or}$

$(\text{Result} < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The 16-bit word in GPR[Rd] is added to the sign-extended 8-bit value given in the instruction and the result is placed into GPR[Rd].

Addressing Mode: Register-Immediate.

4.5 ADC

Add Word With Carry

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	0	Rd			Ra			Rb		X	X	

Syntax

ADC Rd, Ra, Rb

eg. ADC R5, R3, R2

Operation

$Rd \leftarrow Ra + Rb + C$

$N \leftarrow \text{if } (\text{Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if } (\text{Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{if } (Ra > 0 \text{ and } (Rb + CFlag) > 0 \text{ and } \text{Result} < 0) \text{ or}$

$(Ra < 0 \text{ and } (Rb + CFlag) < 0 \text{ and } \text{Result} > 0) \text{ then } 1, \text{ else } 0$

$C \leftarrow \text{if } (\text{Result} > 2^{16} - 1) \text{ or}$

$(\text{Result} < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The 16-bit word in GPR[Ra] is added to the 16-bit word in GPR[Rb] with the added carry in set according to the Carry flag from previous operation, and the result is placed into GPR[Rd].

Addressing Mode: Register-Register.

4.6 ADCI

Add Immediate With Carry

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0 0 1 0 1					Rd			Ra			imm5				

Syntax

ADCI Rd, Ra, #imm5

eg. ADCI R5, R4, #7

Operation

$Rd \leftarrow Ra + \#imm5 + C$

$N \leftarrow \text{if } (\text{Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if } (\text{Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{if } (Ra > 0 \text{ and } (\#imm5 + CFlag) > 0 \text{ and } \text{Result} < 0) \text{ or}$
 $(Ra < 0 \text{ and } (\#imm5 + CFlag) < 0 \text{ and } \text{Result} > 0) \text{ then } 1, \text{ else } 0$

$C \leftarrow \text{if } (\text{Result} > 2^{16} - 1) \text{ or}$
 $(\text{Result} < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The 16-bit word in GPR[Ra] is added to the sign-extended 5-bit value given in the instruction with carry in set according to the Carry flag from previous operation, and the result is placed into GPR[Rd].

Addressing Mode: Register-Immediate.

4.7 NEG

Negate Word

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	1	0	Rd			Ra			X	X	X	X	X

Syntax

NEG Rd, Ra

eg. NEG R5, R3

Operation

$Rd \leftarrow 0 - Ra$

$N \leftarrow \text{if } (\text{Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if } (\text{Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{if } (Ra > 0 \text{ and } Rb > 0 \text{ and } \text{Result} < 0) \text{ or}$

$(Ra < 0 \text{ and } Rb < 0 \text{ and } \text{Result} > 0) \text{ then } 1, \text{ else } 0$

$C \leftarrow \text{if } (\text{Result} > 2^{16} - 1) \text{ or}$

$(\text{Result} < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The 16-bit word in GPR[Ra] is added to the 16-bit word in GPR[Rb] and the result is placed into GPR[Rd].

Addressing Mode: Register-Register.

4.8 SUB

Subtract Word

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	Rd			Ra			Rb		X	X	

Syntax

SUB Rd, Ra, Rb

eg. SUB R5, R3, R2

Operation

$Rd \leftarrow Ra - Rb$

$N \leftarrow \text{if } (\text{Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if } (\text{Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{if } (Ra > 0 \text{ and } Rb > 0 \text{ and } \text{Result} < 0) \text{ or}$
 $(Ra < 0 \text{ and } Rb < 0 \text{ and } \text{Result} > 0) \text{ then } 1, \text{ else } 0$

$C \leftarrow \text{if } (\text{Result} > 2^{16} - 1) \text{ or}$
 $(\text{Result} < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The 16-bit word in GPR[Rb] is subtracted from the 16-bit word in GPR[Ra] and the result is placed into GPR[Rd].

Addressing Mode: Register-Register.

4.9 SUBI

Subtract Immediate

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	Rd			Ra			imm5				

Syntax

SUBI Rd, Ra, #imm5

eg. SUBI R5, R3, #7

Operation

$Rd \leftarrow Ra - \#imm5$

$N \leftarrow \text{if } (\text{Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if } (\text{Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{if } (Ra > 0 \text{ and } \#imm5 > 0 \text{ and } \text{Result} < 0) \text{ or}$

$(Ra < 0 \text{ and } \#imm5 < 0 \text{ and } \text{Result} > 0) \text{ then } 1, \text{ else } 0$

$C \leftarrow \text{if } (\text{Result} > 2^{16} - 1) \text{ or}$

$(\text{Result} < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The sign extended 5-bit value given in the instruction is subtracted from the 16-bit word in GPR[Ra] and the result is placed into GPR[Rd].

Addressing Mode: Register-Immediate.

4.10 SUBIB

Subtract Immediate Byte

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	1	Rd			imm8							

Syntax

SUBIB Rd, #imm8

eg. SUBIB R5, #93

Operation

$Rd \leftarrow Rd - \#imm8$

$N \leftarrow \text{if } (\text{Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if } (\text{Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{if } (Rd > 0 \text{ and } \#imm8 > 0 \text{ and } \text{Result} < 0) \text{ or}$

$(Rd < 0 \text{ and } \#imm8 < 0 \text{ and } \text{Result} > 0) \text{ then } 1, \text{ else } 0$

$C \leftarrow \text{if } (\text{Result} > 2^{16} - 1) \text{ or}$

$(\text{Result} < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The 8-bit immediate value given in the instruction is subtracted from the 16-bit word in GPR[Rd] and the result is placed into GPR[Rd].

Addressing Mode: Register-Immediate.

4.11 SUC

Subtract Word With Carry

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	Rd			Ra			Rb		X	X	

Syntax

SUC Rd, Ra, Rb

eg. SUC R5, R3, R2

Operation

$Rd \leftarrow Ra - Rb - C$

$N \leftarrow \text{if } (\text{Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if } (\text{Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{if } (Ra > 0 \text{ and } (Rb - CFlag) > 0 \text{ and } \text{Result} < 0) \text{ or}$
 $(Ra < 0 \text{ and } (Rb - CFlag) < 0 \text{ and } \text{Result} > 0) \text{ then } 1, \text{ else } 0$

$C \leftarrow \text{if } (\text{Result} > 2^{16} - 1) \text{ or}$
 $(\text{Result} < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The 16-bit word in GPR[Rb] is subtracted from the 16-bit word in GPR[Ra] with the subtracted carry in set according to the Carry flag from previous operation, and the result is placed into GPR[Rd].

Addressing Mode: Register-Register.

4.12 SUCI

Subtract Immediate With Carry

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	Rd			Ra			imm5				

Syntax

SUCI Rd, Ra, #imm5

eg. SUCI R5, R4, #7

Operation

$Rd \leftarrow Ra - \#imm5 - C$

$N \leftarrow \text{if } (Result < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if } (Result = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{if } (Ra > 0 \text{ and } (\#imm5 - CFlag) > 0 \text{ and } Result < 0) \text{ or}$

$(Ra < 0 \text{ and } (\#imm5 - CFlag) < 0 \text{ and } Result > 0) \text{ then } 1, \text{ else } 0$

$C \leftarrow \text{if } (Result > 2^{16} - 1) \text{ or}$

$(Result < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The 5-bit immediate value in instruction is subtracted from the 16-bit word in GPR[Ra] with the subtracted carry in set according to the Carry flag from previous operation, and the result is placed into GPR[Rd].

Addressing Mode: Register-Immediate.

4.13 CMP

Compare Word

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	1	X	X	X	Ra			Rb		X	X	

Syntax

CMP Ra, Rb

eg. CMP R3, R2

Operation

Ra - Rb

$N \leftarrow \text{if (Result} < 0 \text{) then } 1, \text{ else } 0$

$Z \leftarrow \text{if (Result} = 0 \text{) then } 1, \text{ else } 0$

$V \leftarrow \text{if (Ra} > 0 \text{ and Rb} > 0 \text{ and Result} < 0 \text{) or}$

$(\text{Ra} < 0 \text{ and Rb} < 0 \text{ and Result} > 0 \text{) then } 1, \text{ else } 0$

$C \leftarrow \text{if (Result} > 2^{16} - 1 \text{) or}$

$(\text{Result} < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The 16-bit word in GPR[Rb] is subtracted from the 16-bit word in GPR[Ra] and the status flags are updated without saving the result.

Addressing Mode: Register-Register.

4.14 CMPI

Compare Immediate

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	1	X	X	X	Ra			imm5				

Syntax

CMPI Ra, #imm5

eg. CMPI R3, #7

Operation

Ra - #imm5

$N \leftarrow \text{if (Result} < 0 \text{) then } 1, \text{ else } 0$

$Z \leftarrow \text{if (Result} = 0 \text{) then } 1, \text{ else } 0$

$V \leftarrow \text{if (Ra} > 0 \text{ and \#imm5} > 0 \text{ and Result} < 0 \text{) or}$

$(\text{Ra} < 0 \text{ and \#imm5} < 0 \text{ and Result} > 0) \text{ then } 1, \text{ else } 0$

$C \leftarrow \text{if (Result} > 2^{16} - 1 \text{) or}$

$(\text{Result} < -2^{16}) \text{ then } 1, \text{ else } 0$

Description

The sign extended 5-bit value given in the instruction is subtracted from the 16-bit word in GPR[Ra] and the status flags are updated without saving the result.

Addressing Mode: Register-Immediate.

4.15 AND

Logical AND

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1 0 0 0 0					Rd			Ra			Rb			X X	

Syntax

AND Rd, Ra, Rb

eg. AND R5, R3, R2

Operation

$Rd \leftarrow Ra \text{ AND } Rb$

$N \leftarrow \text{if (Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if (Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{UNPREDICTABLE}$

$C \leftarrow \text{UNPREDICTABLE}$

Description

The logical **AND** of the 16-bit words in GPR[Ra] and GPR[Rb] is performed and the result is placed into GPR[Rd].

Addressing Mode: Register-Register.

4.16 OR

Logical OR

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	Rd			Ra			Rb		X	X	

Syntax

OR Rd, Ra, Rb

eg. OR R5, R3, R2

Operation

$Rd \leftarrow Ra \text{ OR } Rb$

$N \leftarrow \text{if (Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if (Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{UNPREDICTABLE}$

$C \leftarrow \text{UNPREDICTABLE}$

Description

The logical OR of the 16-bit words in GPR[Ra] and GPR[Rb] is performed and the result is placed into GPR[Rd].

Addressing Mode: Register-Register.

4.17 XOR

Logical XOR

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	1	1	Rd			Ra			Rb		X	X	

Syntax

XOR Rd, Ra, Rb

eg. XOR R5, R3, R2

Operation

$Rd \leftarrow Ra \text{ XOR } Rb$

$N \leftarrow \text{if (Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if (Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{UNPREDICTABLE}$

$C \leftarrow \text{UNPREDICTABLE}$

Description

The logical **XOR** of the 16-bit words in GPR[Ra] and GPR[Rb] is performed and the result is placed into GPR[Rd].

Addressing Mode: Register-Register.

4.18 NOT

Logical NOT

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	1	0	Rd			Ra			X	X	X	X	X

Syntax

NOT Rd, Ra

eg. NOT R5, R3

Operation

$Rd \leftarrow \text{NOT } Ra$

$N \leftarrow \text{if (Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if (Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{UNPREDICTABLE}$

$C \leftarrow \text{UNPREDICTABLE}$

Description

The logical NOT of the 16-bit word in GPR[Ra] is performed and the result is placed into GPR[Rd].

Addressing Mode: Register-Register.

4.19 NAND

Logical NAND

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	0	Rd			Ra			Rb		X	X	

Syntax

NAND Rd, Ra, Rb

eg. NAND R5, R3, R2

Operation

$Rd \leftarrow Ra \text{ NAND } Rb$

$N \leftarrow \text{if (Result} < 0 \text{) then } 1, \text{ else } 0$

$Z \leftarrow \text{if (Result} = 0 \text{) then } 1, \text{ else } 0$

$V \leftarrow \text{UNPREDICTABLE}$

$C \leftarrow \text{UNPREDICTABLE}$

Description

The logical NAND of the 16-bit words in GPR[Ra] and GPR[Rb] is performed and the result is placed into GPR[Rd].

Addressing Mode: Register-Register.

4.20 NOR

Logical NOR

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	Rd			Ra			Rb		X	X	

Syntax

NOR Rd, Ra, Rb

eg. NOR R5, R3, R2

Operation

$Rd \leftarrow Ra \text{ NOR } Rb$

$N \leftarrow \text{if (Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if (Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{UNPREDICTABLE}$

$C \leftarrow \text{UNPREDICTABLE}$

Description

The logical NOR of the 16-bit words in GPR[Ra] and GPR[Rb] is performed and the result is placed into GPR[Rd].

Addressing Mode: Register-Register.

4.21 LSL

Logical Shift Left

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	1	Rd				Ra		0		imm4		

Syntax

LSL Rd, Ra, #imm4

eg. LSL R5, R3, #7

Operation

$Rd \leftarrow Ra \ll \#imm4$

$N \leftarrow \text{if (Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if (Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{UNPREDICTABLE}$

$C \leftarrow \text{UNPREDICTABLE}$

Description

The 16-bit word in GPR[Ra] is shifted left by the 4-bit amount specified in the instruction, shifting in zeros, and the result is placed into GPR[Rd].

Addressing Mode: Register-Immediate.

4.22 LSR

Logical Shift Right

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	1	Rd			Ra			0	imm4			

Syntax

LSR Rd, Ra, #imm4

eg. LSR R5, R3, #7

Operation

$Rd \leftarrow Ra \gg \#imm4$

$N \leftarrow \text{if (Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if (Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{UNPREDICTABLE}$

$C \leftarrow \text{UNPREDICTABLE}$

Description

The 16-bit word in GPR[Ra] is shifted right by the 4-bit amount specified in the instruction, shifting in zeros, and the result is placed into GPR[Rd].

Addressing Mode: Register-Immediate.

4.23 ASR

Arithmetic Shift Right

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	0	Rd			Ra			0	imm4			

Syntax

ASR Rd, Ra, #imm4

eg. ASR R5, R3, #7

Operation

$Rd \leftarrow Ra \ggg \#imm4$

$N \leftarrow \text{if (Result} < 0) \text{ then } 1, \text{ else } 0$

$Z \leftarrow \text{if (Result} = 0) \text{ then } 1, \text{ else } 0$

$V \leftarrow \text{UNPREDICTABLE}$

$C \leftarrow \text{UNPREDICTABLE}$

Description

The 16-bit word in GPR[Ra] is shifted right by the 4-bit amount specified in the instruction, shifting in the sign bit of Ra, and the result is placed into GPR[Rd].

Addressing Mode: Register-Immediate.

4.24 LDW

Load Word

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	Rd			Ra			imm5				

Syntax

LDW Rd, [Ra, #imm5]

eg. LDW R5, [R3, #7]

Operation

$Rd \leftarrow \text{Mem}[Ra + \#imm5]$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Data is loaded from memory at the resultant address from addition of GPR[Ra] and the 5-bit immediate value specified in the instruction, and the result is placed into GPR[Rd].

Addressing Mode: Base Plus Offset.

4.25 STW

Store Word

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	0	0	Rd			Ra			imm5				

Syntax

STW Rd, [Ra, #imm5]

eg. STW R5, [R3, #7]

Operation

$\text{Mem}[\text{Ra} + \#\text{imm5}] \leftarrow \text{Rd}$

$\text{N} \leftarrow \text{N}$

$\text{Z} \leftarrow \text{Z}$

$\text{V} \leftarrow \text{V}$

$\text{C} \leftarrow \text{C}$

Description

Data in GPR[Rd] is stored to memory at the resultant address from addition of GPR[Ra] and the 5-bit immediate value specified in the instruction.

Addressing Mode: Base Plus Offset.

4.26 LUI

Load Upper Immediate

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	Rd			imm8							

Syntax

LUI Rd #imm8

eg. LUI R5, #93

Operation

$Rd \leftarrow \{\#imm8, 0\}$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

The 8-bit immediate value provided in the instruction is loaded into the top half in GPR[Rd], setting the bottom half to zero.

Addressing Mode: Register-Immediate.

4.27 LLI

Load Lower Immediate

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	Rd			imm8							

Syntax

LLI Rd #imm8

eg. LLI R5, #93

Operation

$Rd \leftarrow \{Rd[15:8], \#imm8\}$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

The 8-bit immediate value provided in the instruction is loaded into the bottom half in GPR[Rd], leaving the top half unchanged.

Addressing Mode: Register-Immediate.

4.28 BR

Branch Always

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	0	0	0	0	imm8							

Syntax

BR LABEL

eg. BR .loop

Operation

$PC \leftarrow PC + \#imm8$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Unconditionally branch to the resultant address from addition of PC and the 8-bit immediate value specified in the instruction. LABEL can be both a symbolic name or a numeric value, and is capable of jumping forwards or backwards.

Addressing Mode: PC Relative.

4.29 BNE

Branch If Not Equal

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	0	1	1	0	imm8							

Syntax

BNE LABEL

eg. BNE .loop

Operation

if (z=0) $PC \leftarrow PC + \#imm8$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Conditionally branch to the resultant address from addition of PC and the 8-bit immediate value specified in the instruction if zero status flag (Z) equals zero. LABEL can be both a symbolic name or a numeric value, and is capable of jumping forwards or backwards.

Addressing Mode: PC Relative.

4.30 BE

Branch If Equal

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	0	1	1	1	imm8							

Syntax

BE LABEL

eg. BE .loop

Operation

if (z=1) $PC \leftarrow PC + \#imm8$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Conditionally branch to the resultant address from addition of PC and the 8-bit immediate value specified in the instruction if zero status flag (Z) equals one. LABEL can be both a symbolic name or a numeric value, and is capable of jumping forwards or backwards.

Addressing Mode: PC Relative.

4.31 BLT

Branch If Less Than

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	0	1	0	0	imm8							

Syntax

BLT LABEL

eg. BLT .loop

Operation

if (n&!v OR !n&v) $PC \leftarrow PC + \#imm8$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Conditionally branch to the resultant address from addition of PC and the 8-bit immediate value specified in the instruction if negative status flag and overflow status flag are not equivalent. LABEL can be both a symbolic name or a numeric value, and is capable of jumping forwards or backwards.

Addressing Mode: PC Relative.

4.32 BGE

Branch If Greater Than Or Equal

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	0	1	0	1	imm8							

Syntax

BGE LABEL

eg. BGE .loop

Operation

if (n&v OR !n&!v) $PC \leftarrow PC + \#imm8$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Conditionally branch to the resultant address from addition of PC and the 8-bit immediate value specified in the instruction if negative status flag and overflow status flag are equivalent. LABEL can be both a symbolic name or a numeric value, and is capable of jumping forwards or backwards.

Addressing Mode: PC Relative.

4.33 BWL

Branch With Link

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	0	0	1	1	imm8							

Syntax

BWL LABEL

eg. BWL .loop

Operation

$LR \leftarrow PC + 1; PC \leftarrow PC + \#imm8$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Save the current program counter (PC) value plus one to the link register. Then unconditionally branch to the resultant address from addition of PC and the 8-bit immediate value specified in the instruction. LABEL can be both a symbolic name or a numeric value, and is capable of jumping forwards or backwards.

Addressing Mode: PC Relative.

4.34 RET

Return

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	0	0	1	0	imm8							

Syntax

RET

eg. RET

Operation

$PC \leftarrow LR$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Unconditionally branch to the address stored in the link register (LR).

Addressing Mode: Register-Indirect.

4.35 JMP

Jump

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	0	0	0	1	imm8							

Syntax

JMP Ra, #imm5

eg. JMP R3, #7

Operation

$PC \leftarrow Ra + \#imm5$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Unconditionally jump to the resultant address from the addition of GPR[Ra] and the 5-bit immediate value specified in the instruction.

Addressing Mode: Base Plus Offset.

4.36 PUSH

Push From Stack

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	0	1	L	X	X	Ra			0	0	0	0	1

Syntax

PUSH Ra

eg. PUSH R3

PUSH LR

eg. PUSH LR

Operation

$\text{Mem}[\text{R7}] \leftarrow \text{reg}; \text{R7} \leftarrow \text{R7} - 1$

$\text{N} \leftarrow \text{N}$

$\text{Z} \leftarrow \text{Z}$

$\text{V} \leftarrow \text{V}$

$\text{C} \leftarrow \text{C}$

Description

‘reg’ corresponds to either a GPR or the link register, the contents of which are stored to the stack using the address stored in the stack pointer (R7). Then Decrement the stack pointer by one.

Addressing Modes: Register-Indirect, Postdecrement.

4.37 POP

Pop From Stack

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	L	X	X	Ra			0	0	0	0	1

Syntax

POP Ra

POP LR

eg. POP R3

eg. POP LR

Operation

$R7 \leftarrow R7 + 1$; $\text{Mem}[R7] \leftarrow \text{reg}$;

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Increment the stack pointer by one. Then ‘reg’ corresponds to either a GPR or the link register, the contents of which are retrieved from the stack using the address stored in the stack pointer (R7).

Addressing Modes: Register-Indirect, Preincrement.

4.38 RETI

Return From Interrupt

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	0	1	0	0	0	1	1	1	X	X	X	X	X

Syntax

RETI

eg. RETI

Operation

$PC \leftarrow \text{Mem}[R7]$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Restore program counter to its value before interrupt occurred, which is stored on the stack, pointed to be the stack pointer (R7). This must be the last instruction in an interrupt service routine.

Addressing Mode: Register-Indirect.

4.39 ENAI

Enable Interrupts

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	0	1	0	0	1	1	1	1	X	X	X	X	X

Syntax

ENAI

eg. ENAI

Operation

Set Interrupt Enable Flag

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Turn on interrupts by setting interrupt enable flag to true (1).

4.40 DISI

Disable Interrupts

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	0	1	0	1	0	1	1	1	X	X	X	X	X

Syntax

DISI

eg. DISI

Operation

Reset Interrupt Enable Flag

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Turn off interrupts by setting interrupt enable flag to false (0).

4.41 STF

Store Status Flags

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	0	1	0	1	1	1	1	1	X	X	X	X	X

Syntax

STF

eg. STF

Operation

$\text{Mem}[\text{R7}] \leftarrow \{12\text{-bit } 0, Z, C, V, N\}; \text{R7} \leftarrow \text{R7} - 1;$

$N \leftarrow N$

$Z \leftarrow Z$

$V \leftarrow V$

$C \leftarrow C$

Description

Store contents of status flags to stack using address held in stack pointer (R7). Then decrement the stack pointer (R7) by one.

Addressing Modes: Register-Indirect, Postdecrement.

4.42 LDF

Load Status Flags

Format

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	0	1	1	0	0	1	1	1	X	X	X	X	X

Syntax

LDF

eg. LDF

Operation

$R7 \leftarrow R7 + 1$

$N \leftarrow \text{Mem}[R7][0]$

$Z \leftarrow \text{Mem}[R7][3]$

$V \leftarrow \text{Mem}[R7][1]$

$C \leftarrow \text{Mem}[R7][2]$

Description

Increment the stack pointer (R7) by one. Then load content of status flags with lower 4 bits of value retrieved from stack using address held in stack pointer (R7).

Addressing Modes: Register-Indirect, Preincrement.

5 Programming Tips

Lorem Ipsum...

6 Assembler

The current instruction set architecture includes an assembler for converting assembly language into hex. This chapter outlines the required formatting and available features of this assembler.

6.1 Instruction Formatting

Each instruction must be formatted using the following syntax, here “[...]” indicates an optional field:

```
[.LABELNAME] MNEMONIC, OPERANDS, ..., :[COMMENTS]
```

eg. `.loop ADDI, R5, R3, #5 :Add 5 to R3`

Comments may be added by preceding them with either `:` or `;`

Accepted general purpose register values are: R0, R1, R2, R3, R4, R5, R6, R7, SP. These can be upper or lower case and SP is equivalently evaluated to R7.

Branch instructions can take either a symbolic or numeric value. Where a numeric must be relative and between -32 and 31 for a JMP instruction, or between -128 and 127 for any other branch type. If the branch exceeds the accepted range, the assembler will flag an error message.

All label names must begin with a ‘.’ while `.ISR/.isr` and `.define` are special cases used for the interrupt service routine and variable definitions respectively.

Instruction-less or comments only lines are allowed within the assembly file.

Special Case Label

The `.ISR/.isr` label is reserved for the interrupt service routine and may be located anywhere within the file but must finish with a `'RETI'` instruction and be no longer than 126 lines of code. Branches may occur within the ISR, but are not allowed into this subroutine with the exception of a return from a separate subroutine.

6.2 Assembler Directives

Symbolic label names are supported for branch-type instructions. Following the previous syntax definition for `'LABELNAME'`, they can be used instead of numeric branching provided they branch no further than the maximum distance allowed for the instruction used. Definitions are supported by the assembler. They are used to assign meaningful names to the GPRs to aid with programming. Definitions can occur at any point within the file and create a mapping from that point onwards. Different names can be assigned to the same register, but only one is valid at a time.

The accepted syntax for definitions is:

```
.define NAME REGISTER
```

6.3 Running The Assembler

The assembler is a python executable and is run by typing `“./assemble.py”`. Alternatively, the assembler can be placed in a folder on the users path and executed by running `“assemble.py”`. It supports Python versions 2.4.3 to 2.7.3. A help prompt is given by the script if the usage is not correct, or given a `-h` or `--help` argument.

By default, the script will output the assembled hex to a file with the same name, but with a `‘.hex’` extension in the same directory. The user can specify a different file to use by using a `-o filename.hex` or `--output=filename.hex` argument to the script. The output file can also be a relative or absolute path to a different directory.

The full usage for the script is seen in listing 1. This includes the basic rules for writing the assembly language and a version log.

Listing 1: Assembler help prompt

```

1 $> assemble.py
2 Usage: assemble.py [-o outfile] input
3
4 —Team R4 Assembler Help—
5 —Version:
6 1 (CMPI addition onwards)
7 2 (Changed to final ISA, added special case I's and error
   checking
8 3 (Ajr changes – Hex output added, bug fix)
9 4 (Added SP symbol)
10 5 (NOP support added, help added) UNTESTED
11 6 (Interrupt support added [ENAI, DISI, RETI])
12 7 (Checks for duplicate Labels)
13 8 (Support for any ISR location & automated startup code entry)
14 9 (Support for .define)
15 10 (Changed usage)
16 Current is most recent iteration
17 Commenting uses : or ;
18 Labels start with '.': SPECIAL .ISR/.isr -> Interrupt Service
   Routine)
19
   SPECIAL .define -> define new name for
   General Purpose Register, .define NAME R0-R7/SP
20 Instruction Syntax: .[LABELNAME] MNEUMONIC, OPERANDS, ..., :[
   COMMENTS]
21 Registers: R0, R1, R2, R3, R4, R5, R6, R7==SP
22 Branching: Symbolic and Numeric supported
23
24 Notes:
25 Input files are assumed to end with a .asm extension
26 Immediate value sizes are checked
27 Instruction-less lines allowed
28 .ISR may be located anywhere in file
29 .define may be located anywhere, definition valid from location
   in file onwards, may replace existing definitions
30
31
32 Options:
33 -h, --help show this help message and exit
34 -o FILE, --output=FILE
35 output file for the assembled output

```

6.4 Error Messages

Code	Description
ERROR1	Instruction mnemonic is not recognized
ERROR2	Register code within instruction is not recognized
ERROR3	Branch condition code is not recognised
ERROR4	Attempting to branch to undefined location
ERROR5	Instruction mnemonic is not recognized
ERROR6	Attempting to shift by more than 16 or perform a negative shift
ERROR7	Magnitude of immediate value for ADDI, ADCI, SUBI, SUCI, LDW or STW is too large
ERROR8	Magnitude of immediate value for CMPI or JMP is too large
ERROR9	Magnitude of immediate value for ADDIB, SUBIB, LUI or LLI is too large
ERROR10	Attempting to jump more than 127 forward or 128 backwards
ERROR11	Duplicate symbolic link names
ERROR12	Illegal branch to ISR
ERROR13	Multiple ISRs in file
ERROR14	Invalid formatting for .define directive

7 Programs

Every example program in this section uses R7 as a stack pointer which is initialised to the by the program to 0x07D0 using the LUI and LLI instructions. The testbench contains an area of an area of memory with 2048 locations and memory mapped devices. 16 switches at location 0x0800, 16 LEDs at location 0x0801 and a serial io device which can be read from location 0xA000 and has a control register at location 0xA001.

7.1 Multiply

The code for the multiply program is held in Appendix A.1 listing 7. A sixteen bit number is read from input switches, split in to lower and upper bytes which are then multiplied. The resulting sixteen bit word is written to the LEDs before reaching a terminating loop. Equation (1) formally describes the algorithm disregarding physical limitations.

$$A = M \times Q = \sum_{i=0}^{\infty} 2^i M_i Q \text{ where } M_i \in \{0,1\} \quad (1)$$

The subroutine operation is described using C in listing 2. If the result is greater than or equal to 2^{16} the subroutine will fail and return zero. The lowest bit of the multiplier controls the accumulator and the overflow check. The multiplier is shifted right and the quotient is shifted left at every iteration. An unconditional branch is used to keep the algorithm in a while loop. The state of the multiplier is compared at every iteration against zero when the algorithm is finished. As size of the multiplier controls the number of iterations a comparison is made on entry to use the smallest operand.

Listing 2: Multiply Subroutine

```
1  uint16_t multi(uint16_t op1, op2){
2      uint16_t A,M,Q;
3      A = 0;
4      if(op1 < op2){                // Make M small, less loops
5          M = op1; Q = op2;
6      }else{
7          M = op2; Q = op1;
8      }
9      while(1){                    // No loop counter
10         if(M & 0x0001){           // LSb
11             A = A + Q;
12             if(A > 0xFFFF){       // Using carry flag
13                 return 0;         // Overflow - fail
14             }
15         }
16         M = M >> 1;
17         if(0 == M){
18             return A;             // Finished - pass
19         }
20         if(Q & 0x8000){
21             return 0;             // Q >= 2^16 - fail
22         }
23     }
```

```

23     Q = Q << 1;
24 }
25 }

```

7.2 Factorial

The code for the factorial program is held in Appendix A.2 listing 8. It is possible to calculate the factorial of any integer value between 0 and 8 inclusive. The subroutine is called which in turn calls the multiply subroutine discussed in section 7.1. The factorial subroutine does no parameter checking but the multiply code does so if overflow does occur zero is propagated and returned; zero is not a possible factorial. The result is calculated recursively as described using C in listing 3. Large values can cause stack overflow the main body of code makes sure inputs, read from the switches, are sufficiently small.

Listing 3: Recursive Factorial Subroutine

```

1  uint16_t fact(uint16_t x){
2      if(x == 0){
3          return 1;           // 0! = 1
4      }
5      return multi(x, fact(x-1)); // Recursive
6  }

```

7.3 Random

The code for the random program is held in Appendix A.3 listing 9. A random series of numbers is achieved by simulating the 16 bit linear feedback shift register in Figure 1. This produces a new number every 16 sixteen clock cycles so in this case a simulation subroutine is called 16 times. A seed taken from switches and passed to the first subroutine call via the stack is altered and passed to the next subroutine call. No more stack operations are performed. A load from the stack pointer is used write a new random number to LEDs. All contained within an unconditional branch but a loop counter is used control write and reset.

A two input XOR gate is simulated using the XOR operation along with shifting to compare bits in different locations. Bits 2 and 4 are used as inputs so a logical shift left by two is used to align them at the bit 4 position.

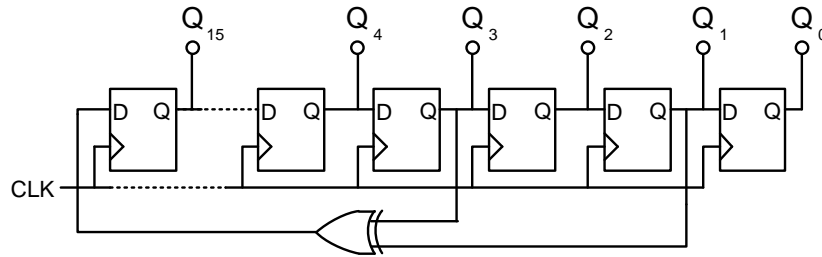


Figure 1: 16 Bit Linear Feedback Shift Register.

Maybe change to IEEE symbols if we have time, AJR: we still have the eagle d-types but I think it would look a bit messy

Masking the output value is used feedback to the top bit. This is described using C in listing 4.

Listing 4: Linear Feedback Shift Register Subroutine

```

1  uint16_t rand(uint16_t last){
2      uint16_t next, test;
3      next = last >> 1;           // The shift
4      test = last << 2;           // Compare different bits
5      test = test ^ last;
6      if(test & 0x0008){          // Feedback to top
7          return (next | 0x8000);
8      }
9      return next;
10 }
```

7.4 Interrupt

The code for the interrupt program is held in Appendix A.4 listing 10. This is the most complex example and makes use of both the multiply and factorial subroutines in sections 7.1 and 7.2 respectively. The interrupt services the serial device by writing data to a 4 byte circular buffer. A main program check to see if data is in the buffer then and if so calculates the factorial writing the result to the LEDs. The buffer is purposefully small to test overflow.

Listing 5: Serial Device Interrupt Service Request

```

1  #define TOP      0x0206
```



```

2 #define BOTTOM 0x0202
3 #define WRITE 0x0201
4 #define READ 0x0200
5 #define SERIAL 0xA000
6
7 isr () {
8     uint16_t data, readPtr, writePtr;
9     asm("DISI"); // critical op
10    data = read(SERIAL);
11    asm("ENAI"); // nested ints
12    readPtr = read(READ);
13    writePtr = read(WRITE);
14    if (((readPtr-1) == writePtr) ||
15        (readPtr == BOTTOM) ||
16        (writePtr == (TOP-1))) {
17        asm("RETI"); // full, don't write
18    }
19    if (readPtr == BOTTOM)
20        write(readPtr, data); // write to buffer
21    writePtr++;
22    if (writePtr == TOP) {
23        writePtr = BOTTOM;
24    } else {
25        writePtr++;
26    }
27    write(WRITE, writePtr);
28    asm("RETI");
29 }
30
31 void main () {
32     uint16_t readPtr, writePtr, data;
33     do {
34         readPtr = read(READ);
35         writePtr = read(WRITE);
36     } while (readPtr == writePtr)
37     data = read(readPtr)
38
39     fact ();
40 }

```

8 Simulation

8.1 Running the simulations

A register window could also be done for this section too

Sim.py needs a fair bit of change. If we have time, this could be altered to use Iains dir structure. This section highlights how to use his script and our assembler.

How to run for each of the behavioural, extracted and mixed

Before the simulator is invoked, the assembler should be invoked. This is discussed in section 6.3. It can be done from within the programs directory (`/design/fcde/verilog/programs`) by running, for example `assemble.py` multiply

The script “simulate” is an executable shell script. It is run from the terminal in the directory `/design/fcde/verilog`. This supports running simulations of a full verilog model, cross simulation and a fully extracted simulation. Usage is as follows:

```
./simulate type program [definitions]
```

The ‘type’ can be one of the following: *behavioural*, *mixed*, *extracted*. ‘Program’ is a relative path to the assembled hex file, usually located in the programs folder. Extra definitions can also be included to set the switch value or serial data input.

The serial data file used is located in the programs directory. This is a hex file with white space separated values of the form “ time data”. The data is then sent at the time to the processor by the serial module. An example serial data hex file is shown in listing 6.

Listing 6: Example serial data file

```
1 // Hex file to specify serial data input
2 //
3 //
4 248    7
5 48F    6
6 6D6    5
7 91D    4
8 B64    7
9 DAB    5
```

10	36B1	3
11	6D61	2

Below is a complete list of commands to run all programs on all versions of the processor. 'Number' is a user defined decimal value to set the switches.

- `./assembler/assemble.py programs/multiply.asm`
`./simulate behavioural programs/multiply.hex +define+switch_value=number`
- `./assembler/assemble.py programs/multiply.asm`
`./simulate mixed programs/multiply.hex +define+switch_value=number`
- `./assembler/assemble.py programs/multiply.asm`
`./simulate extracted programs/multiply.hex +define+switch_value=number`
- `./assembler/assemble.py programs/random.asm`
`./simulate behavioural programs/random.hex +define+switch_value=number`
- `./assembler/assemble.py programs/random.asm`
`./simulate mixed programs/random.hex +define+switch_value=number`
- `./assembler/assemble.py programs/random.asm`
`./simulate extracted programs/random.hex +define+switch_value=number`
- `./assembler/assemble.py programs/factorial.asm`
`./simulate behavioural programs/factorial.hex +define+switch_value=number`
- `./assembler/assemble.py programs/factorial.asm`
`./simulate mixed programs/factorial.hex +define+switch_value=number`
- `./assembler/assemble.py programs/factorial.asm`
`./simulate extracted programs/factorial.hex +define+switch_value=number`
- `./assembler/assemble.py programs/interrupt.asm`
`./simulate behavioural programs/interrupt.hex programs/serial_data.hex`
- `./assembler/assemble.py programs/interrupt.asm`
`./simulate mixed programs/interrupt.hex programs/serial_data.hex`

Table 1: Clock cycles required for each program to run

Make these more accurate when AJR has finished playing around

Program	Clock Cycles
Multiply	900
Factorial	6000
Random	
Interrupt	30000

- `./assembler/assemble.py programs/interrupt.asm`
`./simulate extracted programs/interrupt.hex programs/serial_data.hex`

how to run scan path sim too

The number of clock cycles for each program to fully run is shown in table 1. Factorial run time is given for an input of 8 and is the worst case. Interrupt is dependant on the serial data input and the time is given for the serial data file mentioned above.

A dissembler is also implemented in System Verilog to aid debugging. It is an ASCII formatted array implemented at the top level of the simulation. It is capable of reading the instruction register with in the design, and reconstructing the assembly language of the instruction and is supported in behavioural, mixed and extracted simulations. It will show the opcode, register addresses and immediate values. It is automatically included by the TCL script. The TCL script also opens a waveform window and adds important signals.

A Code Listings

All code listed in this section is passed to the assembler *as is* and has been verified using the final design of the processor.

A.1 Multiply

Listing 7: multiply.asm

```
1      LUI SP, #7      ; Init SP
2      LLI SP, #208
3      LUI R0, #8      ; SWs ADDR
4      LLI R0, #0
5      LDW R0, [R0, #0] ; READ SWs
6      LUI R1, #0
7      LLI R1, #255    ; 0x00FF in R1
8      AND R1, R0, R1   ; Lower byte SWs in R1
9      LSR R0, R0, #8   ; Upper byte SWs in R0
10     SUB R2, R2, R2    ; Zero required
11     PUSH R0           ; Op1
12     PUSH R1           ; Op2
13     PUSH R2           ; Place holder is zero
14     BWL .multi        ; Run Subroutine
15     POP R1            ; Result
16     ADDIB SP, #2      ; Dummy pop
17     LUI R4, #8
18     LLI R4, #1        ; Address of LEDS
19     STW R1, [R4, #0]  ; Result on LEDS
20 .end BR .end          ; Finish loop
21 .multi PUSH R0
22     PUSH R1
23     PUSH R2
24     PUSH R3
25     PUSH R4
26     PUSH R5
27     PUSH R6
28     LDW R0, [SP, #8]  ; R0 - Multiplier
29     LDW R1, [SP, #9]  ; R1 - Quotient
30     SUB R2, R2, R2    ; R2 - Accumulator
31     ADDI R3, R2, #1   ; R3 - Compare 1/0
32     SUB R4, R4, R4    ; R4 - Loop counter
33 .lpMul AND R6, R0, R3 ; R6 - Cmp var
34     CMPI R6, #0
35     BE .sh
```

```

36      SUB R3,R3,R3
37      ADD R2,R2,R1      ; A = A + Q
38      ADCI R3,R3,#1
39      CMPI R3,#2
40      BE .over          ; OV
41 .sh      LUI R5,#128
42          LLI R5,#0      ; 0x8000
43          AND R5,R5,R1
44          CMPI R5,#0
45          BE .shift
46          CMPI R0,#0
47          BNE .over      ; And M != 0
48 .shift   LSL R1,R1,#1   ; Q = Q << 1
49          LSR R0,R0,#1   ; M = M >> 1
50          ADDIB R4,#1    ; i++
51          CMPI R4,#15
52          BNE .lpMul
53 .done    STW R2,[SP,#7] ; Res on stack frame
54          POP R6
55          POP R5
56          POP R4
57          POP R3
58          POP R2
59          POP R1
60          POP R0
61          RET
62 .over    SUB R2,R2,R2    ; OV – RET 0
63          BR .done

```

A.2 Factorial

Listing 8: factorial.asm

```

1      LUI R7, #7
2      LLI R7, #208
3      LUI R0, #8        ; Address in R0
4      LLI R0, #0
5      LDW R0,[R0,#0]    ; Read switches into R0
6      LUI R1,#0         ; Calculate only 8 or less
7      LLI R1,#8
8      CMP R1,R0
9      BE .do
10     SUBIB R1,#1
11     AND R0,R0,R1

```

```

12 .do      PUSH R0          ; Pass para
13         BWL .fact        ; Run Subroutine
14         POP R0           ; Para overwritten with result
15         LUI R4, #8
16         LLI R4, #1       ; Address of LEDS
17         STW R0,[R4,#0]   ; Result on LEDS
18 .end     BR .end         ; finish loop
19 .fact    PUSH R0
20         PUSH R1
21         PUSH LR
22         LDW R1,[SP,#3]   ; Get para
23         ADDIB R1,#0
24         BE .retOne      ; 0! = 1
25         SUBI R0,R1,#1
26         PUSH R0         ; Pass para
27         BWL .fact        ; The output remains on the stack
28         PUSH R1         ; Pass para
29         SUBIB SP,#1      ; Placeholder
30         BWL .multi
31         POP R1           ; Get res
32         ADDIB SP,#2      ; POP x 2
33         STW R1,[SP,#3]
34         POP LR
35         POP R1
36         POP R0
37         RET
38 .retOne  ADDIB R1,#1     ; Avoid jump checking
39         STW R1,[SP,#3]
40         POP LR
41         POP R1
42         POP R0
43         RET
44 .multi   PUSH R0
45         PUSH R1
46         PUSH R2
47         PUSH R3
48         PUSH R4
49         PUSH R5
50         PUSH R6
51         LDW R2,[SP,#8]   ; R2 - Multiplier
52         LDW R3,[SP,#9]   ; R3 - Quotient
53         SUB R4,R4,R4     ; R4 - Accumulator
54         ADDI R6,R4,#1    ; R6 - Constant 1
55         SUB R5,R5,R5     ; R5 - Constant 0
56         SUB R0,R0,R0     ; R0 - C check

```

```

57      AND R1,R2,R6      ; Stage 1, R1 – cmp
58      CMPI R1,#0        ; LSb ?
59      BE .sh1
60      ADD R4,R4,R3      ; (LSb == 1)?
61 .sh1  LSL R3,R3,#1
62      LSR R2,R2,#1
63      AND R1,R2,R6      ; Stage 2
64      CMPI R1,#0
65      BE .sh2
66      ADD R4,R4,R3
67 .sh2  LSL R3,R3,#1
68      LSR R2,R2,#1
69      AND R1,R2,R6      ; Stage 3
70      CMPI R1,#0
71      BE .sh3
72      ADD R4,R4,R3
73 .sh3  LSL R3,R3,#1
74      LSR R2,R2,#1
75      AND R1,R2,R6      ; Stage 4
76      CMPI R1,#0
77      BE .sh4
78      ADD R4,R4,R3
79 .sh4  LSL R3,R3,#1
80      LSR R2,R2,#1
81      AND R1,R2,R6      ; Stage 5
82      CMPI R1,#0
83      BE .sh5
84      ADD R4,R4,R3
85 .sh5  LSL R3,R3,#1
86      LSR R2,R2,#1
87      AND R1,R2,R6      ; Stage 6
88      CMPI R1,#0
89      BE .sh6
90      ADD R4,R4,R3
91 .sh6  LSL R3,R3,#1
92      LSR R2,R2,#1
93      AND R1,R2,R6      ; Stage 7
94      CMPI R1,#0
95      BE .sh7
96      ADD R4,R4,R3
97 .sh7  LSL R3,R3,#1
98      LSR R2,R2,#1
99      AND R1,R2,R6      ; Stage 8
100     CMPI R1,#0
101     BE .sh8

```



```

102      ADD R4,R4,R3
103 .sh8  LSL R3,R3,#1
104      LSR R2,R2,#1
105      AND R1,R2,R6      ; Stage 9
106      CMPI R1,#0
107      BE .sh9
108      ADD R4,R4,R3
109      ADCI R0,R5,#0
110      CMPI R0,#0
111      BNE .over
112 .sh9  LSL R3,R3,#1
113      LSR R2,R2,#1
114      AND R1,R2,R6      ; Stage 10
115      CMPI R1,#0
116      BE .sh10
117      ADD R4,R4,R3
118      ADCI R0,R5,#0
119      CMPI R0,#0
120      BNE .over
121 .sh10 LSL R3,R3,#1
122      LSR R2,R2,#1
123      AND R1,R2,R6      ; Stage 11
124      CMPI R1,#0
125      BE .sh11
126      ADD R4,R4,R3
127      ADCI R0,R5,#0
128      CMPI R0,#0
129      BNE .over
130 .sh11 LSL R3,R3,#1
131      LSR R2,R2,#1
132      AND R1,R2,R6      ; Stage 12
133      CMPI R1,#0
134      BE .sh12
135      ADD R4,R4,R3
136      ADCI R0,R5,#0
137      CMPI R0,#0
138      BNE .over
139 .sh12 LSL R3,R3,#1
140      LSR R2,R2,#1
141      AND R1,R2,R6      ; Stage 13
142      CMPI R1,#0
143      BE .sh13
144      ADD R4,R4,R3
145      ADCI R0,R5,#0
146      CMPI R0,#0

```

```

147      BNE .over
148 .sh13  LSL R3,R3,#1
149      LSR R2,R2,#1
150      AND R1,R2,R6      ; Stage 14
151      CMPI R1,#0
152      BE .sh14
153      ADD R4,R4,R3
154      ADCI R0,R5,#0
155      CMPI R0,#0
156      BNE .over
157 .sh14  LSL R3,R3,#1
158      LSR R2,R2,#1
159      AND R1,R2,R6      ; Stage 15
160      CMPI R1,#0
161      BE .sh15
162      ADD R4,R4,R3
163      ADCI R0,R5,#0
164      CMPI R0,#0
165      BNE .over
166 .sh15  LSL R3,R3,#1
167      LSR R2,R2,#1
168      AND R1,R2,R6      ; Stage 16
169      CMPI R1,#0
170      BE .sh16
171      ADD R4,R4,R3
172      ADCI R0,R5,#0
173      CMPI R0,#0
174      BNE .over
175 .sh16  STW R4,[SP,#7]  ; Res on stack frame
176      POP R6
177      POP R5
178      POP R4
179      POP R3
180      POP R2
181      POP R1
182      POP R0
183      RET
184 .over  SUB R4,R4,R4
185      STW R4,[SP,#7]  ; Res on stack frame
186      POP R6
187      POP R5
188      POP R4
189      POP R3
190      POP R2
191      POP R1

```

```

192     POP R0
193     RET

```

A.3 Random

Listing 9: random.asm

```

1      LUI    SP,#7           ; Init SP
2      LLI    SP,#208
3      LUI    R0,#8           ; SW Address in R0
4      LLI    R0,#0
5      LDW    R1,[R0,#0]      ; Read switches into R1
6      ADDIB  R0,#1           ; Address of LEDS in R0
7      PUSH   R1
8  .reset SUB   R4,R4,R4       ; Reset Loop counter
9  .loop   BWL    .rand
10         CMPI   R4,#15
11         BE     .write
12         ADDIB  R4,#1        ; INC loop counter
13         BR     .loop
14  .write LDW    R1,[SP,#0]    ; No POP as re-run
15         STW    R1,[R0,#0]   ; Result on LEDS
16         BR     .reset
17  .rand   PUSH   R0           ; LFSR Sim
18         PUSH   R1           ; Protect regs
19         PUSH   R2
20         LDW    R0,[SP,#3]    ; Last reg value
21         LSL    R1,R0,#2      ; Shift Bit 4 <- 2
22         XOR    R1,R0,R1      ; XOR Gate
23         LSR    R0,R0,#1      ; Shifted reg
24         LUI    R2,#0
25         LLI    R2,#8
26         AND    R1,R2,R1      ; Mask off Bit 4
27         CMPI   R1,#0
28         BNE    .done
29         LUI    R1,#128
30         LLI    R1,#0
31         OR     R0,R0,R1      ; OR with 0x8000
32  .done  STW    R0,[SP,#3]
33         POP    R2
34         POP    R1
35         POP    R0
36         RET

```

A.4 Interrupt

Listing 10: interrupt.asm

```
1      DISI                ; Reset is off anyway
2      LUI R7, #7
3      LLI R7, #208
4      LUI R0, #2          ; R0 is read ptr    0x0200
5      LLI R0, #0
6      ADDI R1,R0,#2       ; 0x0202
7      STW R1,[R0,#0]      ; Read ptr set to   0x0202
8      STW R1,[R0,#1]      ; Write ptr set to  0x0202
9      LUI R0,#160         ; Address of Serial control reg
10     LLI R0,#1
11     LUI R1,#0
12     LLI R1,#1          ; Data to enable ints
13     STW R1,[R0,#0]      ; Store 0x001 @ 0xA001
14     ENAI
15     BR .main
16 .isr DISI
17     STF                ; Keep flags
18     PUSH R0             ; Save only this for now
19     LUI R0,#160
20     LLI R0,#0
21     LDW R0,[R0,#0]      ; R1 contains read serial data
22     ENAI                ; Don't miss event
23     PUSH R1
24     PUSH R2
25     PUSH R3
26     PUSH R4
27     LUI R1,#2
28     LLI R1,#0
29     LDW R2,[R1,#0]      ; R2 contains read ptr
30     ADDI R3,R1,#1
31     LDW R4,[R3,#0]      ; R4 contain the write ptr
32     SUBIB R2,#1         ; Get out if W == R - 1
33     CMP R4,R2
34     BE .isrOut
35     ADDIB R2,#1
36     LUI R1,#2
37     LLI R1,#2
38     CMP R2,R1
39     BNE .write
40     ADDIB R1,#3
41     CMP R4,R1
42     BE .isrOut
```

```

43 .write STW R0,[R4,#0] ; Write to buffer
44 ADDIB R4,#1
45 LUI R1,#2
46 LLI R1,#6
47 CMP R1,R4
48 BNE .wrapW
49 SUBIB R4,#4
50 .wrapW STW R4,[R3,#0] ; Inc write ptr
51 .isrOut POP R4
52 POP R3
53 POP R2
54 POP R1
55 POP R0
56 LDF
57 RETI
58 .main LUI R0, #2 ; Read ptr address in R0
59 LLI R0, #0
60 LDW R2,[R0,#0] ; Read ptr in R2
61 LDW R3,[R0,#1] ; Write ptr in R3
62 CMP R2,R3
63 BE .main ; Jump back if the same
64 LDW R3,[R2,#0] ; Load data out of buffer
65 ADDIB R2,#1 ; Inc read ptr
66 SUB R0,R0,R0
67 LUI R0,#2
68 LLI R0,#6
69 SUB R0,R0,R2
70 BNE .wrapR
71 SUBIB R2,#4
72 .wrapR LUI R0, #2 ; Read ptr address in R0
73 LLI R0, #0
74 STW R2,[R0,#0] ; Store new read pointer
75 SUB R4,R4,R4
76 LLI R4,#15
77 AND R3,R4,R3
78 CMPI R3,#8
79 BE .do
80 LLI R4,#7
81 AND R3,R3,R4
82 .do PUSH R3
83 BWL .fact
84 POP R3
85 LUI R4,#8
86 LLI R4,#1 ; Address of LEDs
87 STW R3,[R4,#0] ; Put factorial on LEDs

```

```

88      BR .main          ; look again
89 .fact  PUSH R0
90      PUSH R1
91      PUSH LR
92      LDW R1,[SP,#3]    ; Get para
93      ADDIB R1,#0
94      BE .retOne       ; 0! = 1
95      SUBI R0,R1,#1
96      PUSH R0          ; Pass para
97      BWL .fact        ; The output remains on the stack
98      PUSH R1          ; Pass para
99      SUBIB SP,#1      ; Placeholder
100     BWL .multi
101     POP R1            ; Get res
102     ADDIB SP,#2      ; POP x 2
103     STW R1,[SP,#3]
104     POP LR
105     POP R1
106     POP R0
107     RET
108 .retOne ADDIB R1,#1   ; Avoid jump checking
109     STW R1,[SP,#3]
110     POP LR
111     POP R1
112     POP R0
113     RET
114 .multi  PUSH R0
115     PUSH R1
116     PUSH R2
117     PUSH R3
118     PUSH R4
119     PUSH R5
120     PUSH R6
121     LDW R2,[SP,#8]    ; R2 - Multiplier
122     LDW R3,[SP,#9]    ; R3 - Quotient
123     SUB R4,R4,R4      ; R4 - Accumulator
124     ADDI R6,R4,#1     ; R6 - Constant 1
125     SUB R5,R5,R5      ; R5 - Constant 0
126     SUB R0,R0,R0      ; R0 - C check
127     AND R1,R2,R6      ; Stage 1, R1 - cmp
128     CMPI R1,#0        ; LSb ?
129     BE .sh1
130     ADD R4,R4,R3      ; (LSb == 1)?
131 .sh1   LSL R3,R3,#1
132     LSR R2,R2,#1

```

```

133      AND R1,R2,R6      ; Stage 2
134      CMPI R1,#0
135      BE .sh2
136      ADD R4,R4,R3
137 .sh2  LSL R3,R3,#1
138      LSR R2,R2,#1
139      AND R1,R2,R6      ; Stage 3
140      CMPI R1,#0
141      BE .sh3
142      ADD R4,R4,R3
143 .sh3  LSL R3,R3,#1
144      LSR R2,R2,#1
145      AND R1,R2,R6      ; Stage 4
146      CMPI R1,#0
147      BE .sh4
148      ADD R4,R4,R3
149 .sh4  LSL R3,R3,#1
150      LSR R2,R2,#1
151      AND R1,R2,R6      ; Stage 5
152      CMPI R1,#0
153      BE .sh5
154      ADD R4,R4,R3
155 .sh5  LSL R3,R3,#1
156      LSR R2,R2,#1
157      AND R1,R2,R6      ; Stage 6
158      CMPI R1,#0
159      BE .sh6
160      ADD R4,R4,R3
161 .sh6  LSL R3,R3,#1
162      LSR R2,R2,#1
163      AND R1,R2,R6      ; Stage 7
164      CMPI R1,#0
165      BE .sh7
166      ADD R4,R4,R3
167 .sh7  LSL R3,R3,#1
168      LSR R2,R2,#1
169      AND R1,R2,R6      ; Stage 8
170      CMPI R1,#0
171      BE .sh8
172      ADD R4,R4,R3
173 .sh8  LSL R3,R3,#1
174      LSR R2,R2,#1
175      AND R1,R2,R6      ; Stage 9
176      CMPI R1,#0
177      BE .sh9

```

```

178      ADD R4,R4,R3
179      ADCI R0,R5,#0
180      CMPI R0,#0
181      BNE .over
182 .sh9   LSL R3,R3,#1
183       LSR R2,R2,#1
184       AND R1,R2,R6      ; Stage 10
185       CMPI R1,#0
186       BE .sh10
187      ADD R4,R4,R3
188      ADCI R0,R5,#0
189      CMPI R0,#0
190      BNE .over
191 .sh10  LSL R3,R3,#1
192       LSR R2,R2,#1
193       AND R1,R2,R6      ; Stage 11
194       CMPI R1,#0
195       BE .sh11
196      ADD R4,R4,R3
197      ADCI R0,R5,#0
198      CMPI R0,#0
199      BNE .over
200 .sh11  LSL R3,R3,#1
201       LSR R2,R2,#1
202       AND R1,R2,R6      ; Stage 12
203       CMPI R1,#0
204       BE .sh12
205      ADD R4,R4,R3
206      ADCI R0,R5,#0
207      CMPI R0,#0
208      BNE .over
209 .sh12  LSL R3,R3,#1
210       LSR R2,R2,#1
211       AND R1,R2,R6      ; Stage 13
212       CMPI R1,#0
213       BE .sh13
214      ADD R4,R4,R3
215      ADCI R0,R5,#0
216      CMPI R0,#0
217      BNE .over
218 .sh13  LSL R3,R3,#1
219       LSR R2,R2,#1
220       AND R1,R2,R6      ; Stage 14
221       CMPI R1,#0
222       BE .sh14

```



```

223      ADD R4,R4,R3
224      ADCI R0,R5,#0
225      CMPI R0,#0
226      BNE .over
227 .sh14  LSL R3,R3,#1
228      LSR R2,R2,#1
229      AND R1,R2,R6      ; Stage 15
230      CMPI R1,#0
231      BE .sh15
232      ADD R4,R4,R3
233      ADCI R0,R5,#0
234      CMPI R0,#0
235      BNE .over
236 .sh15  LSL R3,R3,#1
237      LSR R2,R2,#1
238      AND R1,R2,R6      ; Stage 16
239      CMPI R1,#0
240      BE .sh16
241      ADD R4,R4,R3
242      ADCI R0,R5,#0
243      CMPI R0,#0
244      BNE .over
245 .sh16  STW R4,[SP,#7]  ; Res on stack frame
246      POP R6
247      POP R5
248      POP R4
249      POP R3
250      POP R2
251      POP R1
252      POP R0
253      RET
254 .over  SUB R4,R4,R4
255      STW R4,[SP,#7]  ; Res on stack frame
256      POP R6
257      POP R5
258      POP R4
259      POP R3
260      POP R2
261      POP R1
262      POP R0
263      RET

```