MODULE *voldchain*

EXTENDS *Integers*, *Sequences*, *FiniteSets*, *TLC*

CONSTANTS N, C, $STOP$, $FAILNUM$

ASSUME $N - FAILNUM \geq 1 \land STOP < 5 \land \land 0 \leq FAILNUM \land FAILNUM \leq 2$

$Nodes \triangleq 1 \,..\, N$

$Clients \triangleq N + 1 \,..\, N + C$

$Procs \triangleq 1 \,..\, N + C$

$Configurator \triangleq N + C + 1$

```
--algorithm test{
    variable FailNum = FAILNUM, step = 0, LSTWR = -1, CURRD = -1, lsttmp = -1, WRTFLG =
                 msg = [j ∈ Procs ↦ ⟨⟩],
                 up = [n ∈ Nodes ↦ TRUE],
                 db = [n ∈ Nodes ↦ [ver ↦ -1, val ↦ -1, cli ↦ -1]],
                 chain = ⟨⟩ ;

    define
    {
        UpNodes ≜ {i ∈ Nodes : up[i] = TRUE}   Returns the set of up nodes
        ChainNode ≜ {chain[i] : i ∈ 1 .. Len(chain)}   Returns the elements in the chain as set
        ChainNodes ≜ IF (Len(chain) = 0) THEN {-1} ELSE ChainNode   Retruns {-1} if chain is empty else t
        InChain(s) ≜ IF (s ∈ ChainNodes) THEN TRUE ELSE FALSE Check if the element is on chain
        FreeUpNodes ≜ {i ∈ UpNodes : InChain(i) = FALSE ∧ up[i] = TRUE}   Retrun set of upnodes that are not
        GetIndex(s) ≜ CHOOSE i ∈ 1 .. Len(chain) : chain[i] = s   Get the index position of s in chain
        GetNext(s) ≜ IF (GetIndex(s) = Len(chain)) THEN -1 ELSE chain[GetIndex(s) + 1] Get the success
        GetFreeNode ≜ {i ∈ FreeUpNodes : up[i] = TRUE}
        GetTail ≜ chain[Len(chain)]                                   Get the tail of chain
        GetHead ≜ chain[1]
        test(e) ≜ IF (up[e] = TRUE) THEN TRUE ELSE FALSE                                    (
     }

    fair process ( c ∈ Clients )
    variable cntr = 0, hver = -1, lastelm = 0, tail = -1, head = -1, initial = 0 ;
    {
        C0: await (Len(chain) > 0) ;
        CL: while ( cntr ≤ STOP )
        {
          CLR:
          while ( WRTFLG ≠ self )    Iterate through read request
          {
            if ( msg[self] ≠ ⟨⟩ )
            {
              if ( msg[self].val = -1 )
              {
                hver := msg[self].ver + 1 ;
```

```
                CURRD := db[chain[Len(chain)]].ver || LSTWR := lsttmp ;
              if ( WRTFLG = −1 )
              {
                WRTFLG := self ;
              } ;
            }
          } ;
          if ( WRTFLG ≠ self )
          {
            tail := GetTail ;
            msg[tail] := [ver ↦ −1, val ↦ −1, cli ↦ self] ;
          }
        } ;
        CLW :
        while ( WRTFLG = self )   Iterate through write request
        {
          if ( msg[self].val ≠ −1 ∧ msg[self].ver = hver )
          {
            lsttmp := msg[self].ver ;
            cntr := cntr + 1 ;
            WRTFLG := −1 ;
          } ;
          if ( WRTFLG = self )
          {
            head := GetHead ;
            msg[head] := [ver ↦ hver, val ↦ cntr, cli ↦ self] ;
          }
        }
      }
    }
fair process ( n ∈ Nodes )
variable nextnode = −1, clientid = 0 ;
{
  ND : while ( TRUE )
  {
    either
    NM :
    {   react to message
      if ( up[self] = TRUE ∧ msg[self] ≠ ⟨⟩ ∧ InChain(self) = TRUE )
      {
        clientid := msg[self].cli ;
        if ( msg[self].val = −1 ∧ self = GetTail )   handling Read message request
        {
          msg[clientid] := [ver ↦ db[self].ver, val ↦ −1, cli ↦ clientid] || msg[self] := ⟨⟩ ;
        }
```

```
        else if ( msg[self].val ≠ −1 )   handling Write message request
        {
          db[self] := [ver ↦ msg[self].ver, val ↦ msg[self].val, cli ↦ msg[self].cli] ;
          if ( self = GetTail )          check if this is a tail node
          {
            lsttmp := db[self].ver ;
            msg[clientid] := [ver ↦ db[self].ver, val ↦ db[self].val, cli ↦ clientid] || msg[self] := ⟨⟩ ;
          }
          else                           else get the successor
          {
            nextnode := GetNext(self) ;
            msg[nextnode] := msg[self] || msg[self] := ⟨⟩ ;
          }
        } ;
      }
    } or
    NDF:
    {
      if ( FailNum > 0 ∧ up[self] = TRUE )   Storage node can fail
      {
        up[self] := FALSE ;
        FailNum := FailNum − 1 ;
      }
      else if ( up[self] = FALSE )   Or recover as a new node
      {
        up[self]   := TRUE ;
        msg[self] := ⟨⟩     ;
        FailNum := FailNum + 1 ;
      }
    }
  }
}
fair process ( p = Configurator )   Maintain the chain
variable newnode = −1 ;
{
  P: while ( TRUE )
  {
    P1:
    if ( Len(chain) < 3 )   Add a new node
    {
      if ( FreeUpNodes ≠ {} )
      {
        newnode := CHOOSE k ∈ FreeUpNodes : TRUE ;
        if ( Len(chain) = 0 )
        {
```

```
            db[newnode] := [ver ↦ −1, val ↦ −1, cli ↦ 0] ;
            chain := Append(chain, newnode) ;
          }
          else
          {
            db[newnode] := db[chain[Len(chain)]] ;
            chain := Append(chain, newnode) ;
          } ;
        } ;
      }
      else  Delete fail node
      {
        chain := SelectSeq(chain, test) ;
      }
    }
  }
}
```

BEGIN TRANSLATION

VARIABLES *FailNum*, *step*, *LSTWR*, *CURRD*, *lsttmp*, *WRTFLG*, *msg*, *up*, *db*, *chain*, *pc*

define statement

$UpNodes \triangleq \{i \in Nodes : up[i] = \text{TRUE}\}$
$ChainNode \triangleq \{chain[i] : i \in 1 \,..\, Len(chain)\}$
$ChainNodes \triangleq \text{IF } (Len(chain) = 0) \text{ THEN } \{-1\} \text{ ELSE } ChainNode$
$InChain(s) \triangleq \text{IF } (s \in ChainNodes) \text{ THEN TRUE ELSE FALSE}$
$FreeUpNodes \triangleq \{i \in UpNodes : InChain(i) = \text{FALSE} \wedge up[i] = \text{TRUE}\}$
$GetIndex(s) \triangleq \text{CHOOSE } i \in 1 \,..\, Len(chain) : chain[i] = s$
$GetNext(s) \triangleq \text{IF } (GetIndex(s) = Len(chain)) \text{ THEN } -1 \text{ ELSE } chain[GetIndex(s) + 1]$
$GetFreeNode \triangleq \{i \in FreeUpNodes : up[i] = \text{TRUE}\}$
$GetTail \triangleq chain[Len(chain)]$
$GetHead \triangleq chain[1]$
$test(e) \triangleq \text{IF } (up[e] = \text{TRUE}) \text{ THEN TRUE ELSE FALSE}$

VARIABLES *cntr*, *hver*, *lastelm*, *tail*, *head*, *initial*, *nextnode*, *clientid*, *newnode*

$vars \triangleq \langle FailNum, step, LSTWR, CURRD, lsttmp, WRTFLG, msg, up, db, chain, pc, cntr, hver, lastelm, tail, head, initial, nextnode, clientid, newnode\rangle$

$ProcSet \triangleq (Clients) \cup (Nodes) \cup \{Configurator\}$

$Init \triangleq$ Global variables
$\wedge\ FailNum = FAILNUM$
$\wedge\ step = 0$
$\wedge\ LSTWR = -1$

$\land\ CURRD = -1$
$\land\ lsttmp = -1$
$\land\ WRTFLG = -1$
$\land\ msg = [j \in Procs \mapsto \langle\rangle]$
$\land\ up = [n \in Nodes \mapsto \text{TRUE}]$
$\land\ db = [n \in Nodes \mapsto [ver \mapsto -1,\ val \mapsto -1,\ cli \mapsto -1]]$
$\land\ chain = \langle\rangle$
Process c
$\land\ cntr = [self \in Clients \mapsto 0]$
$\land\ hver = [self \in Clients \mapsto -1]$
$\land\ lastelm = [self \in Clients \mapsto 0]$
$\land\ tail = [self \in Clients \mapsto -1]$
$\land\ head = [self \in Clients \mapsto -1]$
$\land\ initial = [self \in Clients \mapsto 0]$
Process n
$\land\ nextnode = [self \in Nodes \mapsto -1]$
$\land\ clientid = [self \in Nodes \mapsto 0]$
Process p
$\land\ newnode = -1$
$\land\ pc = [self \in ProcSet \mapsto \text{CASE } self \in Clients \to \text{“C0”}$
$\qquad\qquad\qquad\qquad\qquad\qquad \square\ \ self \in Nodes \to \text{“ND”}$
$\qquad\qquad\qquad\qquad\qquad\qquad \square\ \ self = Configurator \to \text{“P”}]$

$C0(self) \triangleq\ \land\ pc[self] = \text{“C0”}$
$\qquad\qquad\land\ (Len(chain) > 0)$
$\qquad\qquad\land\ pc' = [pc \text{ EXCEPT } ![self] = \text{“CL”}]$
$\qquad\qquad\land \text{ UNCHANGED } \langle FailNum,\ step,\ LSTWR,\ CURRD,\ lsttmp,\ WRTFLG,\ msg,$
$\qquad\qquad\qquad\qquad up,\ db,\ chain,\ cntr,\ hver,\ lastelm,\ tail,\ head,$
$\qquad\qquad\qquad\qquad initial,\ nextnode,\ clientid,\ newnode\rangle$

$CL(self) \triangleq\ \land\ pc[self] = \text{“CL”}$
$\qquad\qquad\land \text{ IF } cntr[self] \leq STOP$
$\qquad\qquad\qquad\text{THEN } \land\ pc' = [pc \text{ EXCEPT } ![self] = \text{“CLR”}]$
$\qquad\qquad\qquad\text{ELSE } \land\ pc' = [pc \text{ EXCEPT } ![self] = \text{“Done”}]$
$\qquad\qquad\land \text{ UNCHANGED } \langle FailNum,\ step,\ LSTWR,\ CURRD,\ lsttmp,\ WRTFLG,\ msg,$
$\qquad\qquad\qquad\qquad up,\ db,\ chain,\ cntr,\ hver,\ lastelm,\ tail,\ head,$
$\qquad\qquad\qquad\qquad initial,\ nextnode,\ clientid,\ newnode\rangle$

$CLR(self) \triangleq\ \land\ pc[self] = \text{“CLR”}$
$\qquad\qquad\land \text{ IF } WRTFLG \neq self$
$\qquad\qquad\qquad\text{THEN } \land \text{ IF } msg[self] \neq \langle\rangle$
$\qquad\qquad\qquad\qquad\qquad\text{THEN } \land \text{ IF } msg[self].val = -1$
$\qquad\qquad\qquad\qquad\qquad\qquad\qquad\text{THEN } \land\ hver' = [hver \text{ EXCEPT } ![self] = msg[self].ver + 1]$
$\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\land\ \land\ CURRD' = db[chain[Len(chain)]].ver$
$\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\ \land\ LSTWR' = lsttmp$
$\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\land \text{ IF } WRTFLG = -1$

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                                                THEN ∧ WRTFLG′ = self
                                                ELSE ∧ TRUE
                                                     ∧ UNCHANGED WRTFLG
                                          ELSE ∧ TRUE
                                               ∧ UNCHANGED ⟨LSTWR, CURRD,
                                                           WRTFLG, hver⟩
                                     ELSE ∧ TRUE
                                          ∧ UNCHANGED ⟨LSTWR, CURRD, WRTFLG, hver⟩
                                ∧ IF WRTFLG′ ≠ self
                                     THEN ∧ tail′ = [tail EXCEPT ![self] = GetTail]
                                          ∧ msg′ = [msg EXCEPT ![tail′[self]] = [ver ↦ −1, val ↦ −1, cli ↦
                                     ELSE ∧ TRUE
                                          ∧ UNCHANGED ⟨msg, tail⟩
                                ∧ pc′ = [pc EXCEPT ![self] = "CLR"]
                           ELSE ∧ pc′ = [pc EXCEPT ![self] = "CLW"]
                                ∧ UNCHANGED ⟨LSTWR, CURRD, WRTFLG, msg, hver, tail⟩
                     ∧ UNCHANGED ⟨FailNum, step, lsttmp, up, db, chain, cntr,
                                  lastelm, head, initial, nextnode, clientid,
                                  newnode⟩

CLW(self) ≜ ∧ pc[self]       = "CLW"
            ∧ IF WRTFLG = self
                 THEN ∧ IF msg[self].val ≠ −1 ∧ msg[self].ver = hver[self]
                           THEN ∧ lsttmp′ = msg[self].ver
                                ∧ cntr′ = [cntr EXCEPT ![self] = cntr[self] + 1]
                                ∧ WRTFLG′ = −1
                           ELSE ∧ TRUE
                                ∧ UNCHANGED ⟨lsttmp, WRTFLG, cntr⟩
                      ∧ IF WRTFLG′ = self
                           THEN ∧ head′ = [head EXCEPT ![self] = GetHead]
                                ∧ msg′ = [msg EXCEPT ![head′[self]] = [ver ↦ hver[self], val ↦ cntr
                           ELSE ∧ TRUE
                                ∧ UNCHANGED ⟨msg, head⟩
                      ∧ pc′ = [pc EXCEPT ![self] = "CLW"]
                 ELSE ∧ pc′ = [pc EXCEPT ![self] = "CL"]
                      ∧ UNCHANGED ⟨lsttmp, WRTFLG, msg, cntr, head⟩
            ∧ UNCHANGED ⟨FailNum, step, LSTWR, CURRD, up, db, chain, hver,
                         lastelm, tail, initial, nextnode, clientid,
                         newnode⟩

c(self) ≜ C0(self) ∨ CL(self) ∨ CLR(self) ∨ CLW(self)

ND(self) ≜ ∧ pc[self] = "ND"
           ∧ ∨ ∧ pc′ = [pc EXCEPT ![self] = "NM"]
             ∨ ∧ pc′ = [pc EXCEPT ![self] = "NDF"]
           ∧ UNCHANGED ⟨FailNum, step, LSTWR, CURRD, lsttmp, WRTFLG, msg,
```

up, db, $chain$, $cntr$, $hver$, $lastelm$, $tail$, $head$,
$initial$, $nextnode$, $clientid$, $newnode\rangle$

$NM(self) \triangleq$ $\land\ pc[self] = \text{“NM”}$
$\land$ IF $up[self] = \text{TRUE} \land msg[self] \neq \langle\rangle \land InChain(self) = \text{TRUE}$
THEN $\land\ clientid' = [clientid$ EXCEPT $![self] = msg[self].cli]$
$\land$ IF $msg[self].val = -1 \land self = GetTail$
THEN $\land\ msg' = [msg$ EXCEPT $![clientid'[self]] = [ver \mapsto db[self].ver,\ val \mapsto$
$![self] = \langle\rangle]$
$\land$ UNCHANGED $\langle lsttmp,\ db,\ nextnode\rangle$
ELSE $\land$ IF $msg[self].val \neq -1$
THEN $\land\ db' = [db$ EXCEPT $![self] = [ver \mapsto msg[self].ver,\ val \mapsto$
$\land$ IF $self = GetTail$
THEN $\land\ lsttmp' = db'[self].ver$
$\land\ msg' = [msg$ EXCEPT $![clientid'[self]] = [ve$
$![self] = \langle\rangle]$
$\land$ UNCHANGED $nextnode$
ELSE $\land\ nextnode' = [nextnode$ EXCEPT $![self] = Ge$
$\land\ msg' = [msg$ EXCEPT $![nextnode'[self]] = m$
$![self] = \langle\rangle]$
$\land$ UNCHANGED $lsttmp$
ELSE $\land$ TRUE
$\land$ UNCHANGED $\langle lsttmp,\ msg,\ db,$
$nextnode\rangle$
ELSE $\land$ TRUE
$\land$ UNCHANGED $\langle lsttmp,\ msg,\ db,\ nextnode,\ clientid\rangle$
$\land\ pc' = [pc$ EXCEPT $![self] = \text{“ND”}]$
$\land$ UNCHANGED $\langle FailNum,\ step,\ LSTWR,\ CURRD,\ WRTFLG,\ up,\ chain,$
$cntr,\ hver,\ lastelm,\ tail,\ head,\ initial,\ newnode\rangle$

$NDF(self) \triangleq$ $\land\ pc[self] = \text{“NDF”}$
$\land$ IF $FailNum > 0 \land up[self] = \text{TRUE}$
THEN $\land\ up' = [up$ EXCEPT $![self] = \text{FALSE}]$
$\land\ FailNum' = FailNum - 1$
$\land\ msg' = msg$
ELSE $\land$ IF $up[self] = \text{FALSE}$
THEN $\land\ up' = [up$ EXCEPT $![self] = \text{TRUE}]$
$\land\ msg' = [msg$ EXCEPT $![self] = \langle\rangle]$
$\land\ FailNum' = FailNum + 1$
ELSE $\land$ TRUE
$\land$ UNCHANGED $\langle FailNum,\ msg,\ up\rangle$
$\land\ pc' = [pc$ EXCEPT $![self] = \text{“ND”}]$
$\land$ UNCHANGED $\langle step,\ LSTWR,\ CURRD,\ lsttmp,\ WRTFLG,\ db,\ chain,$
$cntr,\ hver,\ lastelm,\ tail,\ head,\ initial,$
$nextnode,\ clientid,\ newnode\rangle$

$n(self) \triangleq ND(self) \lor NM(self) \lor NDF(self)$

$P \triangleq \land pc[Configurator] = \text{“P”}$
$\land pc' = [pc \text{ EXCEPT } ![Configurator] = \text{“P1”}]$
$\land \text{UNCHANGED } \langle FailNum, step, LSTWR, CURRD, lsttmp, WRTFLG, msg, up, db,$
$chain, cntr, hver, lastelm, tail, head, initial, nextnode,$
$clientid, newnode\rangle$

$P1 \triangleq \land pc[Configurator] = \text{“P1”}$
$\land \text{IF } Len(chain) < 3$
$\text{THEN } \land \text{IF } FreeUpNodes \neq \{\}$
$\text{THEN } \land newnode' = (\text{CHOOSE } k \in FreeUpNodes : \text{TRUE})$
$\land \text{IF } Len(chain) = 0$
$\text{THEN } \land db' = [db \text{ EXCEPT } ![newnode'] = [ver \mapsto -1, val \mapsto -1, cli \vdash$
$\land chain' = Append(chain, newnode')$
$\text{ELSE } \land db' = [db \text{ EXCEPT } ![newnode'] = db[chain[Len(chain)]]]$
$\land chain' = Append(chain, newnode')$
$\text{ELSE } \land \text{TRUE}$
$\land \text{UNCHANGED } \langle db, chain, newnode\rangle$
$\text{ELSE } \land chain' = SelectSeq(chain, test)$
$\land \text{UNCHANGED } \langle db, newnode\rangle$
$\land pc' = [pc \text{ EXCEPT } ![Configurator] = \text{“P”}]$
$\land \text{UNCHANGED } \langle FailNum, step, LSTWR, CURRD, lsttmp, WRTFLG, msg, up,$
$cntr, hver, lastelm, tail, head, initial, nextnode,$
$clientid\rangle$

$p \triangleq P \lor P1$

$Next \triangleq p$
$\lor (\exists self \in Clients : c(self))$
$\lor (\exists self \in Nodes : n(self))$

$Spec \triangleq \land Init \land \Box[Next]_{vars}$
$\land \forall self \in Clients : \text{WF}_{vars}(c(self))$
$\land \forall self \in Nodes : \text{WF}_{vars}(n(self))$
$\land \text{WF}_{vars}(p)$

END TRANSLATION

$Invariant1 \triangleq LSTWR = CURRD$
$Invariant2 \triangleq \forall i \in 1 \mathinner{..} (Len(chain) - 1) : db[chain[i]].ver \geq db[chain[i+1]].ver$

* Modification History
* Last modified *Thu Dec* 08 23:49:00 *EST* 2016 by *Rohit*
* Last modified *Thu Dec* 08 20:27:28 *EST* 2016 by *anand*
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Members – *Anand Sankar Bhagavandas* – *UB id* – 50208048

The code given above implements Replicated Storage with Chain Replication.

The invariant we have specified are -

1) The version number of the last write operation is equal to the version number of the current read operation.
2) The second invariant is that the *db.ver* is non-increasing when we traverse from head to tail.

Given below are our observations about this version of the protocol -

Voldchain is capable of tolerating more faults with less number of nodes. This is because in the *project*1 version of the protocol, the condition that the value for *ReadQ* and *WriteQ* should be atleast one greater than the value of *FAILNUM* ineffect also brings up the necessity that the majority of nodes should be up and running.

eg for the case – $ReadQ = 3$, $WriteQ = 3$, $FAILNUM = 2$ and Number of *Nodes* $(N) = 3$, This will bring up an error as there are not enough nodes to satisfy the size requirements for *ReadQ* and *WriteQ* if two nodes are down.

In the *VoldChain* protocol, the chain length only needs to be one greater than *FAILNUM* and the number of nodes can be equal to the chain length. This will ensure that atleast one node with the current value is up to serve the request of the client.

The Read quorum is analogus to the tail of the chain. The client sends read requests to the tail and the tail responds with the value of the version number that is available with it. Since the write messages are constantly put in queue, the tail will eventually possess the highest version number.

The Write quorum is analogus to the whole chain. The client contacts the head which passes on the message to the other nodes. Eventually all the nodes in the chain will possess the updated value.

The previous relation between the quorum and *FAILNUM* holds in this case too. The number of nodes in the chain should be one more than *FAILNUM*.

We have implemented a token based lock system to ensure that the write operation is consistent and that another client does not write while a write operation is in progress. While in possesion of this token, the client keeps sending write message till it gets the confirmation that the write operation is completed and the token is released from the client.

The system runs without the use of the aforementioned token system when the number of clients $C = 1$. But it doesnot satisfy the invariant without the token system when $C = 2$.