

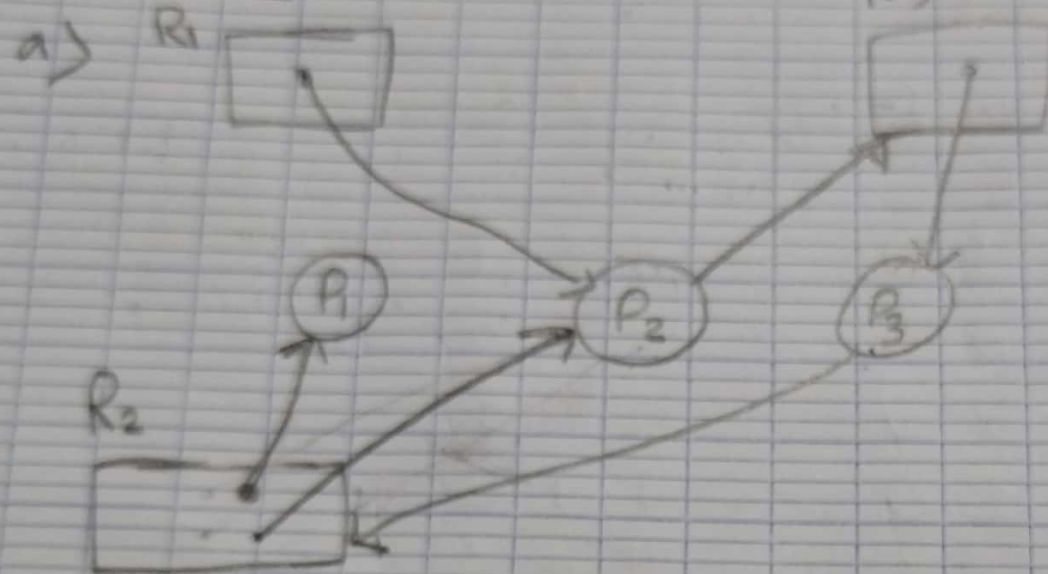
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Question 1

$S_1=0, S_2=0, S_3=0, S_4=0$

$P_1$	$P_2$	$P_3$	$P_4$
wait( $S_1$ ) statement(1)	statement 3; signal( $S_3$ )	statement 5; signal( $S_4$ )	statement 7;
wait( $S_4$ ) statement 2;	statement 4; signal( $S_1$ )	wait( $S_3$ ) statement 6; signal( $S_2$ )	wait( $S_2$ ) statement 8;

Question (2)



b)  $R_2 \rightarrow P_2 \rightarrow R_3 \rightarrow P_3$

c)  $P_3$  can be executed, because  $R_2$  is available

$P_2$  can be executed, because  $R_3$  is available  
Then there is no deadlock



# Question (3)

Allocation A (5 instances)

	A	B	C
P <sub>0</sub>	0	1	0
P <sub>1</sub>	1	0	0
P <sub>2</sub>	2	0	0
P <sub>3</sub>	0	1	0
P <sub>4</sub>	0	0	1

B (3 instances)

	A	B	C
P <sub>0</sub>	4	3	4
P <sub>1</sub>	2	2	5
P <sub>2</sub>	2	1	3
P <sub>3</sub>	1	1	2
P <sub>4</sub>	4	0	1

C (6 instances)

Available

a) Available

2 1 5

b) Need

P <sub>0</sub>	4	2	4	✓
P <sub>1</sub>	1	2	5	✓
P <sub>2</sub>	0	1	3	✓
P <sub>3</sub>	1	0	2	✓
P <sub>4</sub>	4	0	0	✓

c)

P<sub>2</sub> 4 1 5

P<sub>3</sub> 4 2 5

P<sub>0</sub> 4 3 5

P<sub>1</sub> 5 3 5

P<sub>4</sub> 5 3 6

⇒ Safe

Order (P<sub>2</sub> → P<sub>3</sub> → P<sub>0</sub> → P<sub>1</sub> → P<sub>4</sub>)

d) (1, 1, 1) by P<sub>2</sub>

req. > need.

↳ can't be granted by P<sub>2</sub> 1 1 3 > 0 1 3

e) (2, 2, 2) P<sub>0</sub> need 4 2 4

(2, 2, 2) < 4, 2, 4

req. < need ✓

req. & Available

2 2 2 < 2 1 5

⇒ req. > Available

so The request can't be granted.

f) 1, 0, 2

Need of  $P_3$

1 0 2 ✓

Available

2 1 5 ✓

New allocation  $\rightarrow$   
old  
request + Allocation

(1 0 2)

We need to check  $P_3$ .

new Available  $\rightarrow$   
old Available - request

Allocation

Need

A B C

ABC

Available

$P_0$  0 1 0

$P_0$  4 2 4 ✓

2 1 5

new need  $\rightarrow$

$P_1$  1 0 0

$P_1$  1 2 5 ✓

1 0 2

old need - req.

$P_2$  2 0 0

$P_2$  0 1 3 ✓

1 1 3

$P_3$  1 1 2

$P_3$  0 0 0 ✓

$P_4$  0 0 1

$P_4$  4 0 0

Work

1 1 3

$P_2$  3 1 3

$P_3$  4 2 5

$P_0$  4 3 5

$P_1$  5 3 5

$P_4$  5 3 6 ✓