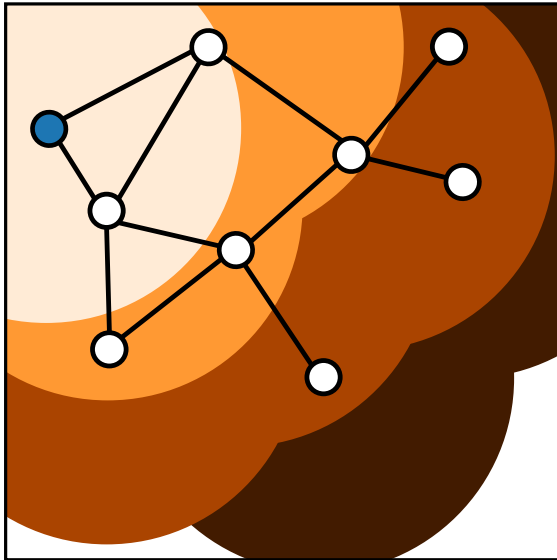


Leveraging Synchronous Transmissions for the Design of Real-time Wireless Cyber-Physical Systems



Romain Jacob

Doctoral Examination

December 17, 2019

Emerging Cyber-Physical Systems (CPS) applications have challenging requirements



Fast and reliable



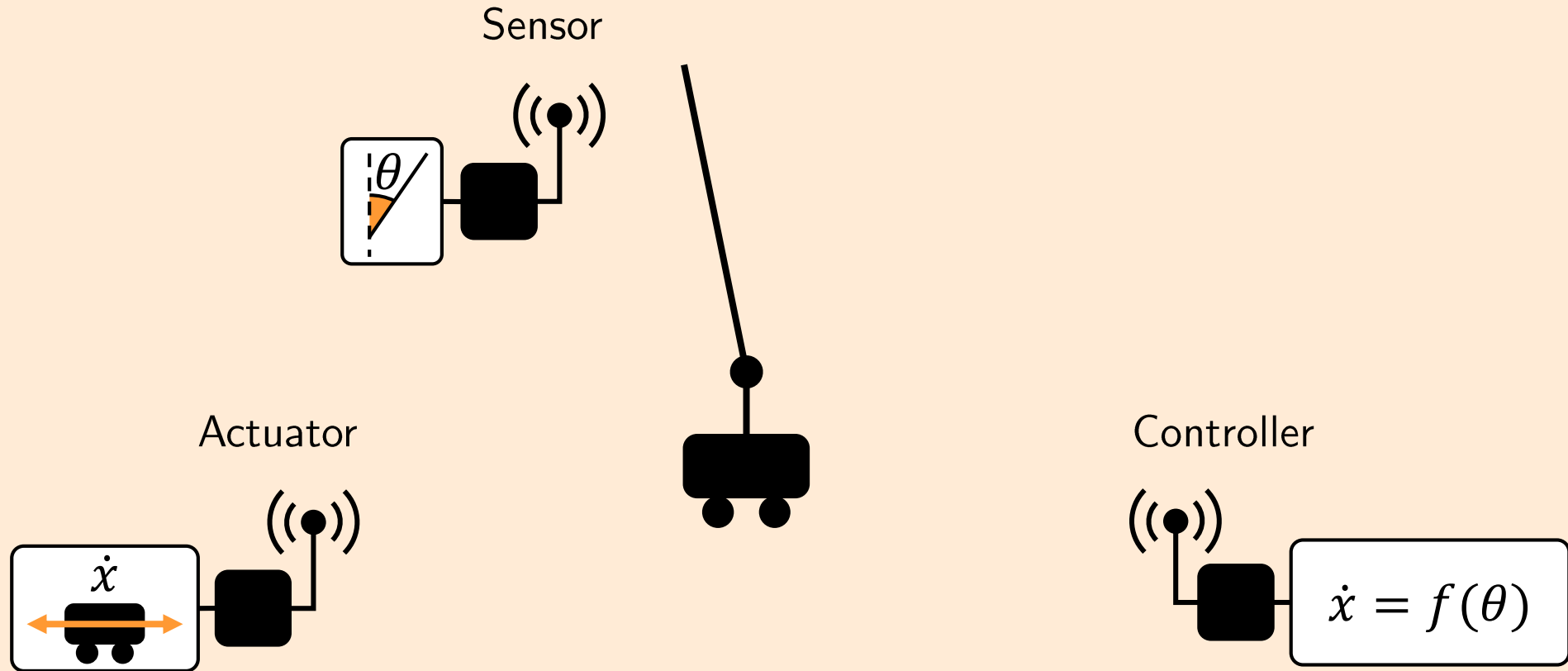
Mobile



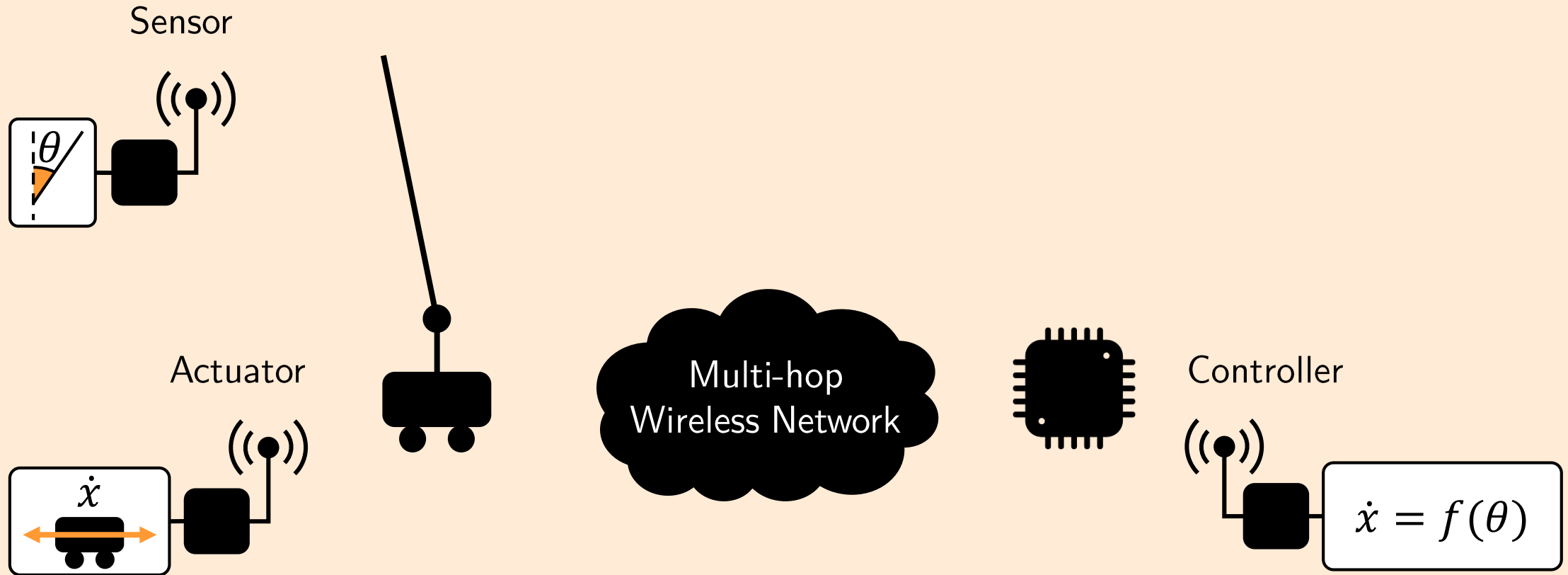
Scalable

Representative example

Stabilizing an inverted pendulum



Things get more complex when the controller is **physically separated** from the pendulum



The application sets the requirements

Primary goals

Predictability

Adaptability

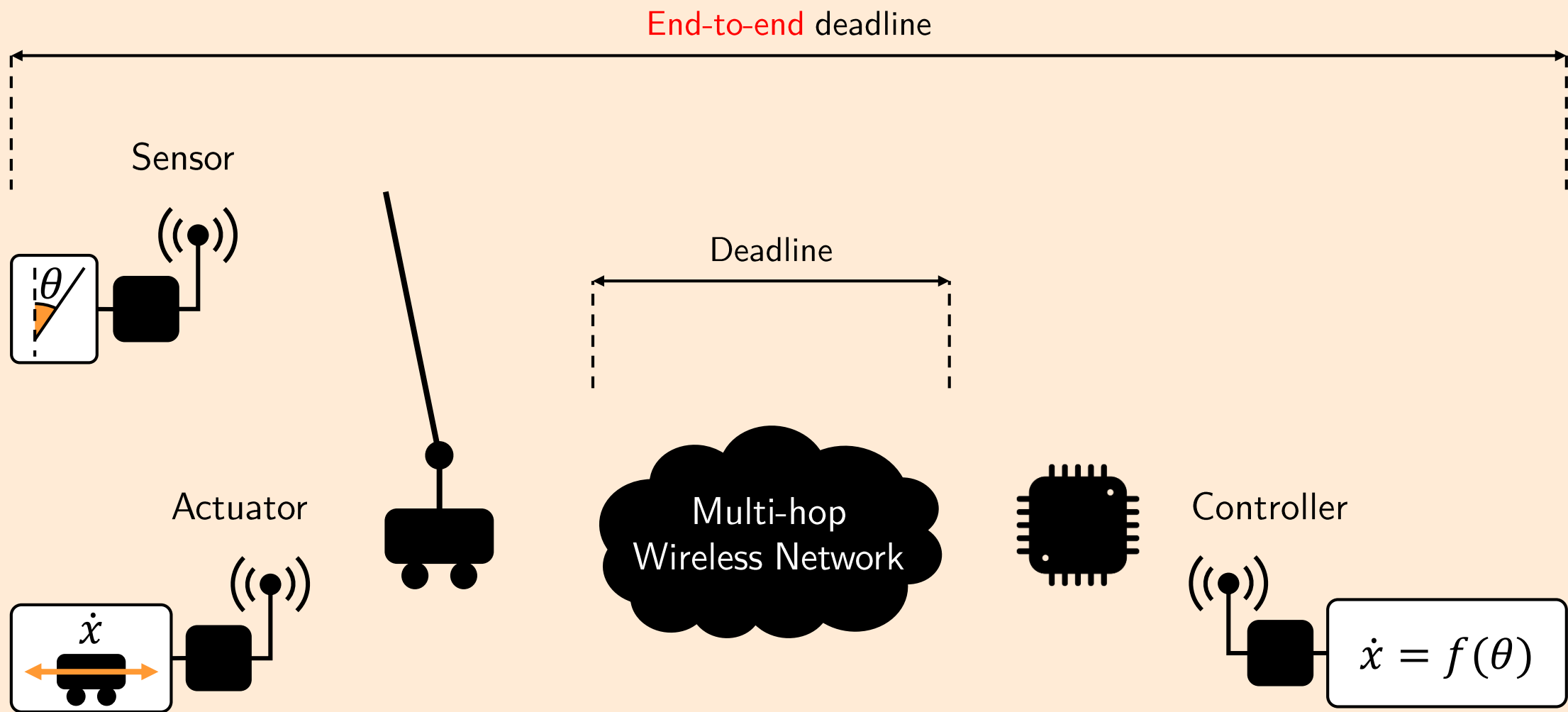
Secondary objectives

Efficiency

System must support

End-to-end real-time guarantees

- Change in traffic demand
- Mobility
- Low latency
- Low energy consumption



Existing wireless system designs are either predictable or adaptive

For example

Adaptive
not Predictable

Splash
Staffeta

Predictable
not Adaptive

WirelessHART
TSCH

and

Nothing proposed
to provide end-to-end
real-time guarantees

Existing wireless system designs are either predictable or adaptive

For example

Adaptive
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Splash
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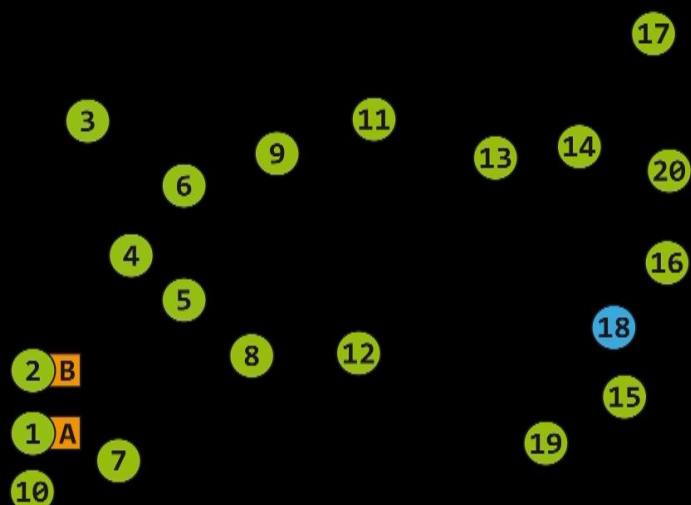
Nothing proposed
to provide **end-to-end**
real-time guarantees

Dissertation objective

Design wireless systems combining
end-to-end predictability and
adaptability while remaining **efficient**

Guaranteed Stability over Low-power Multi-hop Network

Two inverted pendulums



One remote controller

The diagram illustrates a network topology for controlling two inverted pendulums. On the left, two inverted pendulums are shown, each with a blue rod and a black base. On the right, a remote controller is shown as a blue circuit board with a black antenna. In the center, a network of 20 nodes is depicted as green circles with numbers. Nodes 1 and 2 are labeled 'A' and 'B' respectively. The nodes are connected in a multi-hop fashion, forming a path from the remote controller to the pendulums. The nodes are numbered 1 through 20, with 1 and 2 being the closest to the pendulums, and 19 and 20 being the closest to the remote controller. The network is a multi-hop topology, with nodes 1 and 2 connected to node 3, which is connected to node 4, and so on, up to node 19, which is connected to node 20, which is connected to the remote controller.



tiny.cc/WirelessCPSVideo

**Synchronous Transmissions
are well suited to design
real-time wireless CPS.**

Real-time
System Designs

DRP

Chapter 4

TTW

Chapter 5

Tools and
Methods

Baloo

Chapter 3

TriScale

Chapter 2

Synchronous
Transmissions

Networking
in General

Real-time
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DRP

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Chapter 4

DRP: End-to-end Real-time Guarantees in Wireless Cyber-Physical Systems

End-to-End Real-Time Guarantees in Wireless Cyber-Physical Systems

R. Jacob, M. Zimmerling, P. Huang, J. Beutel, L. Thiele, 2016

Real-time
System Designs

DRP

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Methods

Baloo

TriScale

Synchronous
Transmissions

Networking
in General

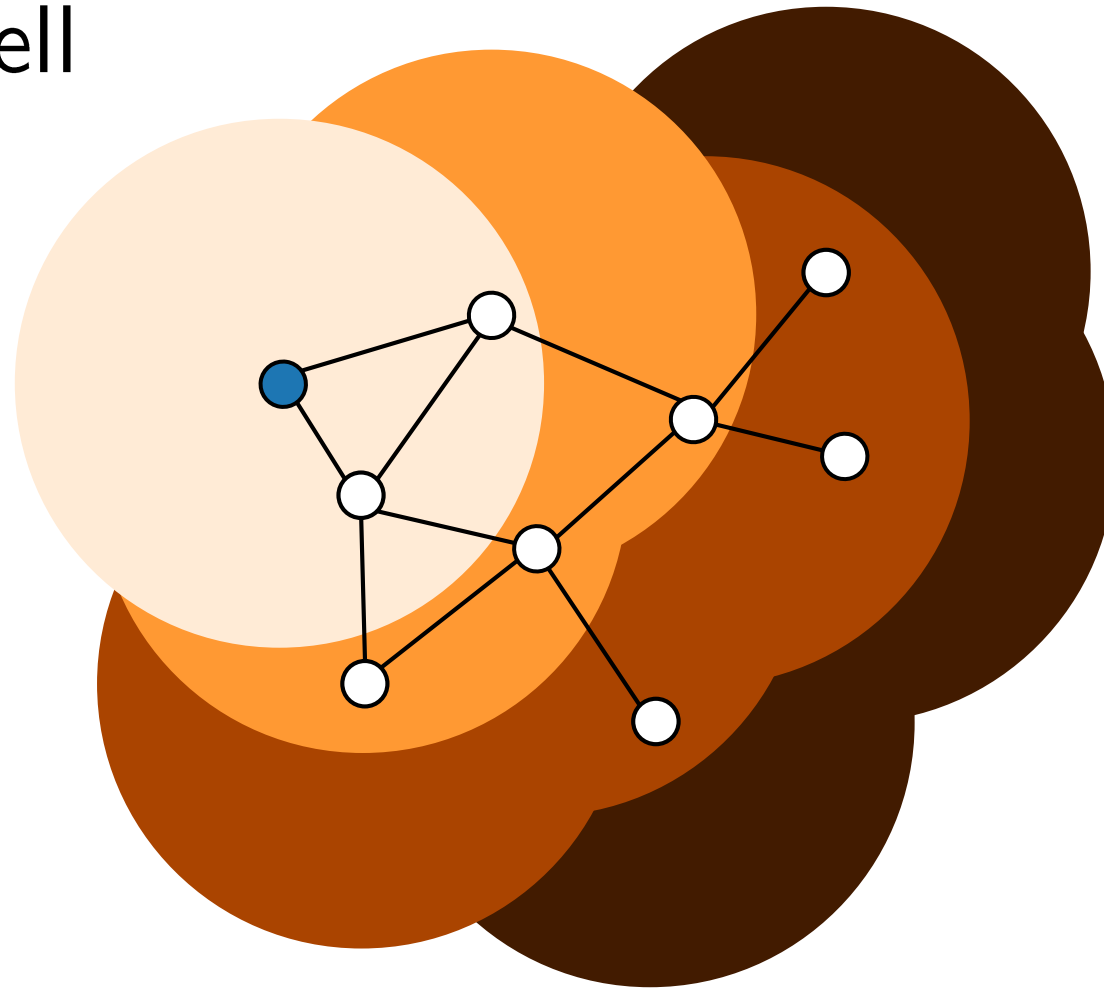
Chapter 5

TTW: A Time-Triggered Design for Wireless Cyber-Physical Systems

TTW: A Time-Triggered Wireless design for CPS

R. Jacob, L. Zhang, M. Zimmerling, J. Beutel, S. Chakraborty, L. Thiele, 2018

Synchronous Transmissions in a nutshell



Multi-hop
broadcast
based on
flooding

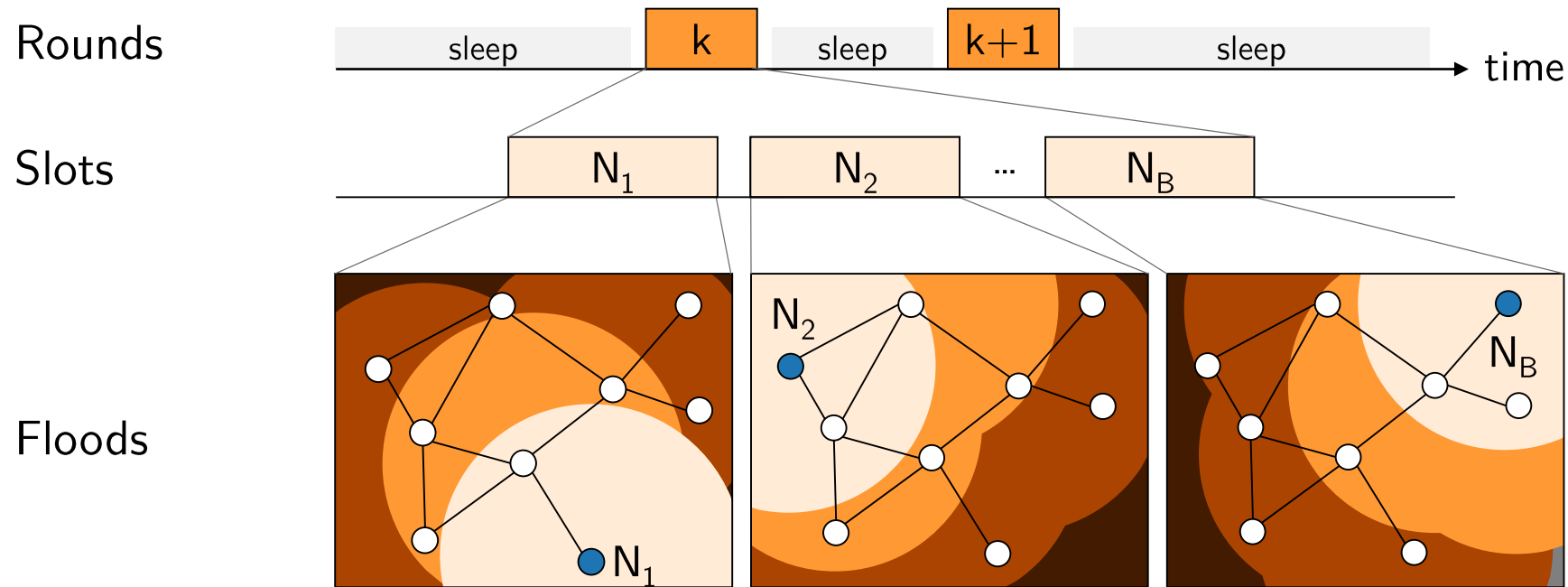
Synchronous Transmissions in a nutshell



Synchronous Transmissions realize
fast **one-to-all communication**
with **predictable timing**

Multi-hop
broadcast
based on
flooding

Typical network stack based on Synchronous Transmissions



Primary
goals

Predictability

Adaptability

Secondary
objectives

Efficiency

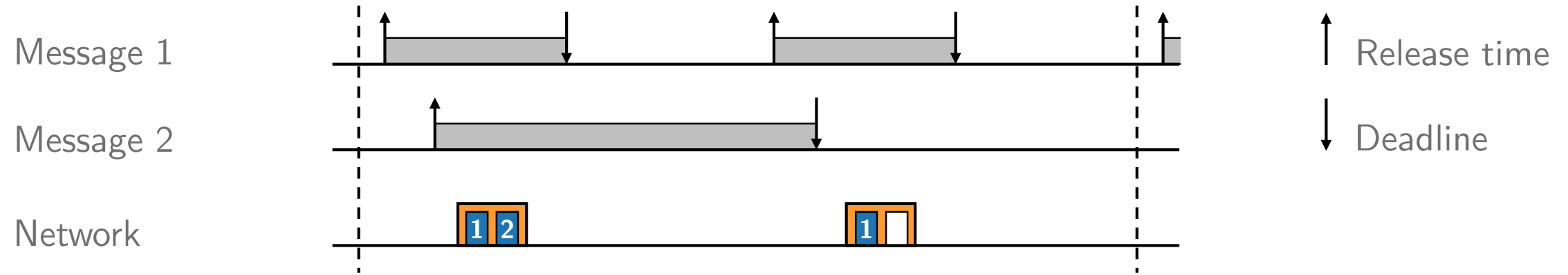
System must support

End-to-end real-time guarantees

- Change in traffic demand
- Mobility

- Low latency
- Low energy consumption

Real-time scheduling allows to meet deadlines



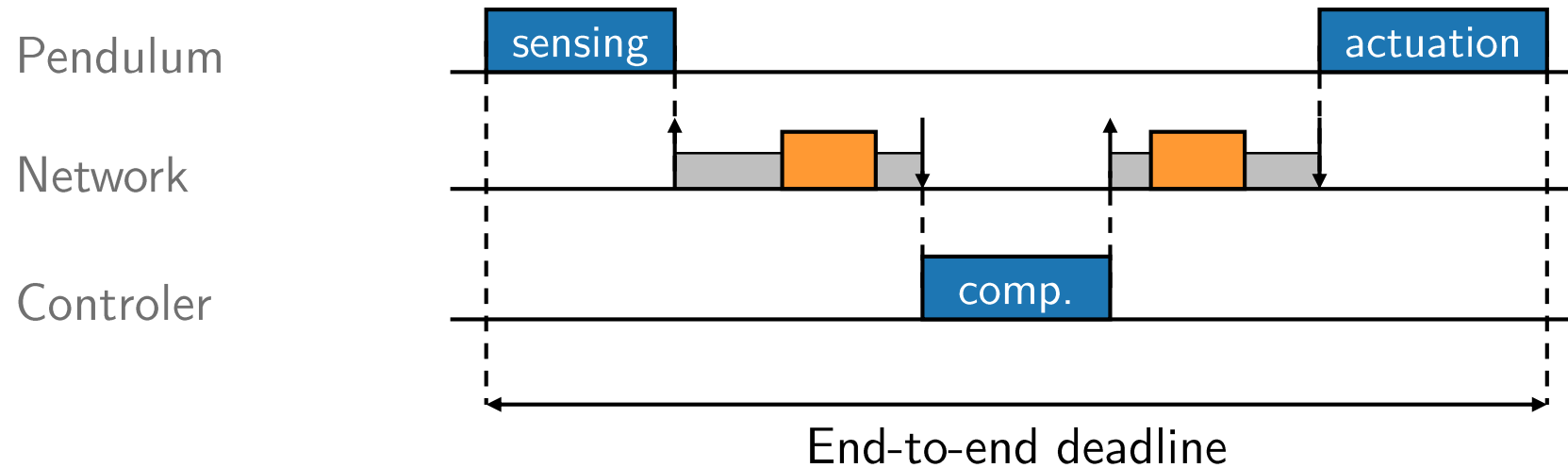
In practice

How do you get
the message deadlines?



From the execution
time of the tasks!

End-to-end guarantees requires **coupling** between the **task** and **message schedules**



End-to-end guarantees requires coupling
between the task and message schedules

Chosen
coupling

TTW statically co-schedules
all tasks and messages

Chosen
coupling

TTW **statically co-schedules**
all tasks and messages

- Schedule **offline**
based on a MILP
- Execute at **runtime**

Inspired by
time-triggered
wired networks
e.g., TTEthernet, FlexRay

Offline scheduling
allows to minimize the
achievable latency at runtime

Rounds save energy but complexify the synthesis

Messages must be served in a round that finishes before their deadline

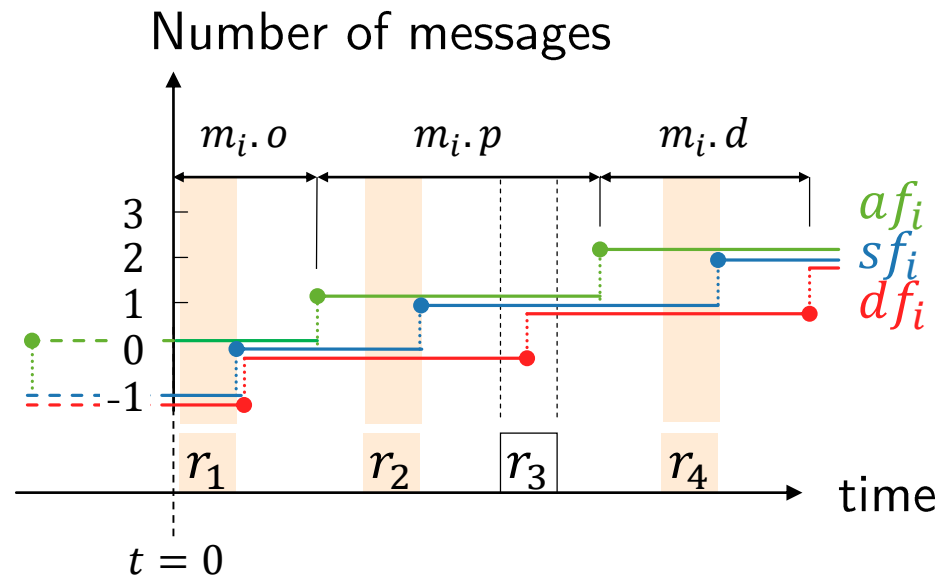
16 B Payload
10 Slots per round \Rightarrow Reduction of
~ 30% in
radio-on time

$$\begin{cases} 1 & \text{if } m_i \text{ is assigned to } r_j \\ 0 & \text{otherwise} \end{cases}$$

$$m_i.d \geq (r_j.o + T_r) * \delta_{ij}$$

Non-linear!

Our solution is inspired by network calculus



Count the number of message instances of m_i that

- af_i have been released
- sf_i have been served
- df_i have passed their deadline

Our solution is inspired by network calculus

$$m_i.d \geq (r_j.o + T_r) * \delta_{ij}$$

$$\Leftrightarrow sf(r_j.o) \geq df(r_j.o + T_r)$$

$$\Leftrightarrow \sum_{k=1}^j \delta_{ik} \geq \left\lceil \frac{r_j.o + T_r - m_i.o - m_i.d}{m_i.p} \right\rceil$$

|
Linear

|
Piecewise-constant
can be handled with
common MILP tricks

Count the number of
message instances of m_i that

af_i have been released

sf_i have been served

df_i have passed their deadline

Static scheduling is nice but **static**

Primary
goals

Predictability

Adaptability

Secondary
objectives

Efficiency

System must support

End-to-end real-time guarantees

- Change in traffic demand
- Mobility
- Low latency
- Low energy consumption

System switches between operation modes at runtime

Well-known approach

- Synthesize schedules for multiple **operation modes**
- Switch between modes at runtime

Challenge

Preserve real-time guarantees across mode changes

- Creates dependencies between modes
- Tackled in TTW while aiming to limit the impact on energy consumption

Primary
goals

Predictability

Adaptability

Secondary
objectives

Efficiency

System must support

End-to-end real-time guarantees

- Change in traffic demand
- Mobility

All great! ... on paper

- Can we implement this?
- Does it really work?

Real-time
System Designs

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TTW

Tools and
Methods

Baloo

TriScale

Synchronous
Transmissions

Networking
in General

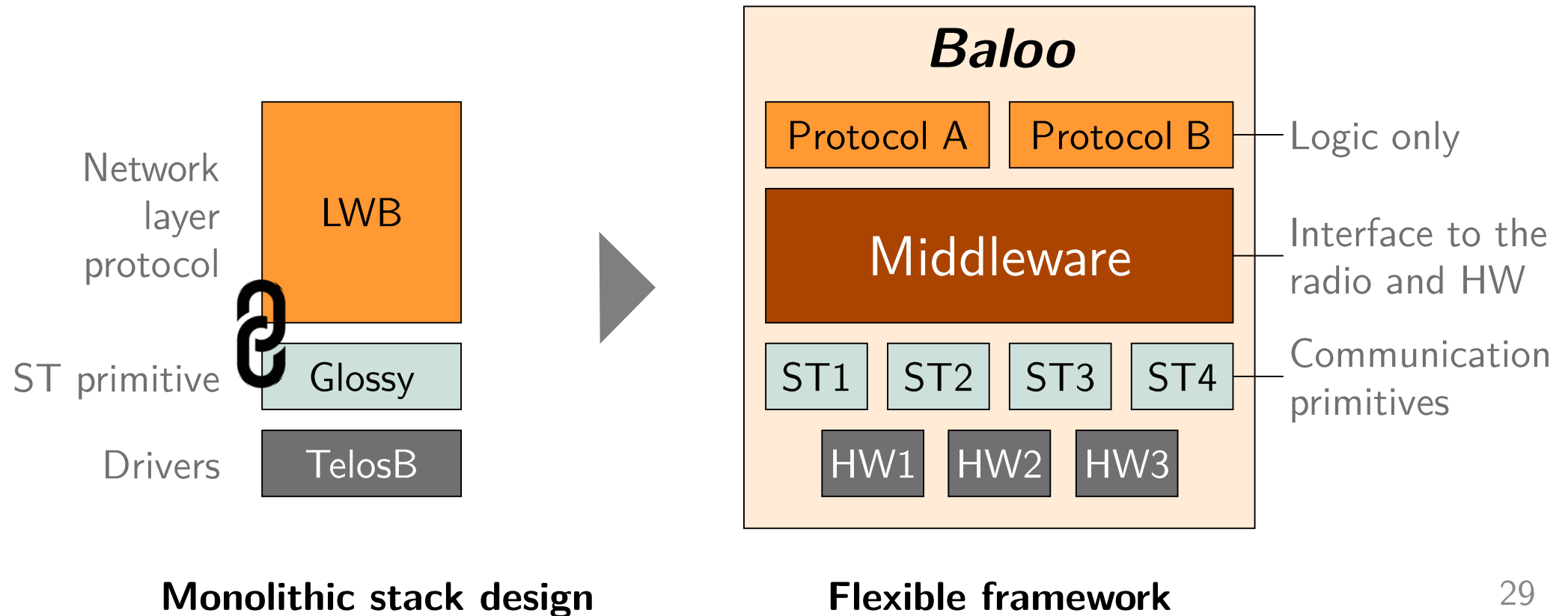
Chapter 3

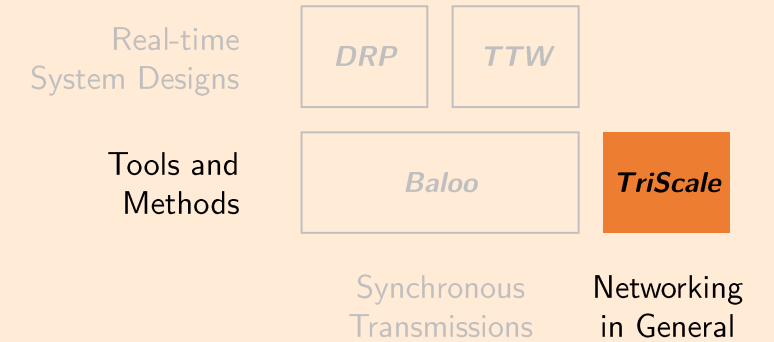
Synchronous Transmissions Made Easy: Design Your Network Stack with Baloo

Synchronous Transmissions Made Easy: Design Your Network Stack with Baloo

R. Jacob, J. Bächli, R. da Forno, L. Thiele, 2019

Baloo facilitates the design of ST-based stacks by decoupling the **protocol logic** and the **radio control**





Chapter 2

TriScale: A Framework Supporting Reproducible Networking Experiments

Towards a Methodology for Experimental Evaluation in Low-Power Wireless Networking

R. Jacob, C. Boano, U. Raza, M. Zimmerling, L. Thiele, 2019

TriScale: A Framework Supporting Reproducible Networking Evaluations

R. Jacob, M. Zimmerling, C. Boano, L. Vanbever, L. Thiele, 2019

What does **reproducibility** mean?

Reproducibility
(unformal)

An experiment is said reproducible if the “same results” are obtained when the experiment is reproduced under the “same conditions”.

Reproducibility (unformal)

An experiment is said reproducible if the “same results” are obtained when the experiment is reproduced under the “**same conditions**”.

- Real RF environment cannot be controlled.
- The RF environment affects performance of protocols in ways that are hard to predict.

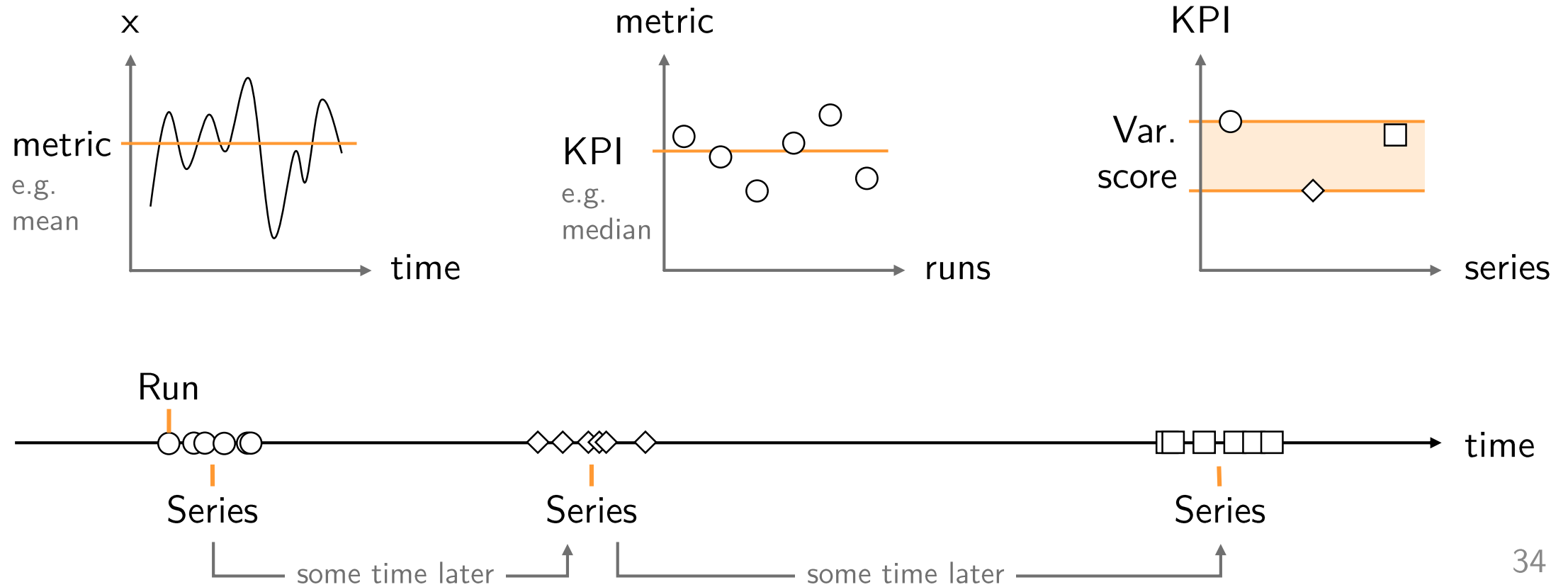
Reproducibility (unformal)

An experiment is said reproducible if the “**same results**” are obtained when the experiment is reproduced under the “**same conditions**”.

- Real RF environment cannot be controlled.
- The RF environment affects performance of protocols in ways that are hard to predict.

⇒ Performance variability is expected.

Performance varies along three different time scales



The four questions of experiment design

How long should a run be?

How many runs in a series?

What time span for a series?

How many series?



The four questions of experiment design

How long should a run be?

How many runs in a series?

What time span for a series?

How many series?

Objective

Find **rational answers** to these questions

Quantify the trade-off between

- experiment effort
- confidence in the results

The four questions of experiment design

How long should a run be?

How many runs in a series?

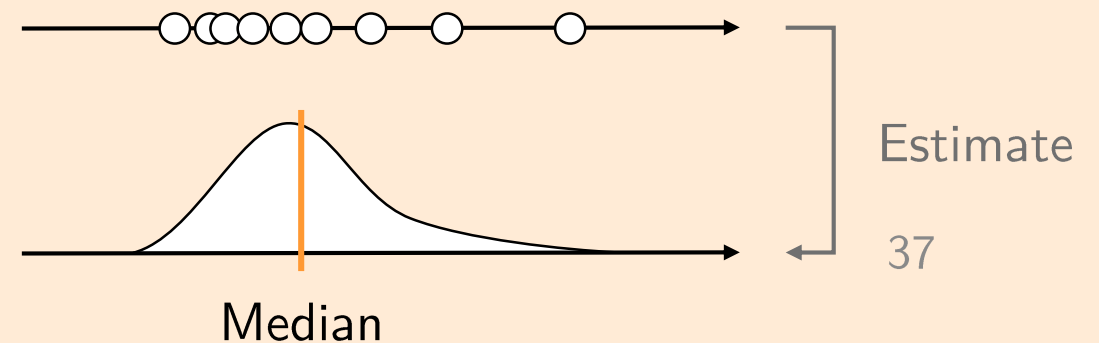
What time span for a series?

How many series?

Objective

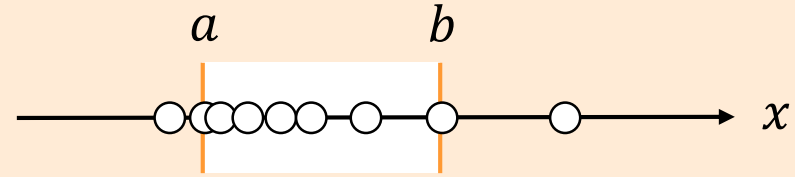
The KPI would be the same
shall more runs be performed

Based on a few runs,
estimates the KPI value of
the underlying distribution



We must use confidence intervals (CI)

Informally Numerical **interval**
in which lies
the **true value**
(which you don't know)
of some parameter
with a certain probability,
called the **confidence level**



$[a, b]$ is a 95% CI for the median of x

which means that

The probability that the true median of x is within $[a, b]$ is larger or equal to 95%.

We must use non-parametric statistics

Makes no assumption on
the nature of the distributions

Common KPIs

Tendency

Mean

Spread

Standard deviation

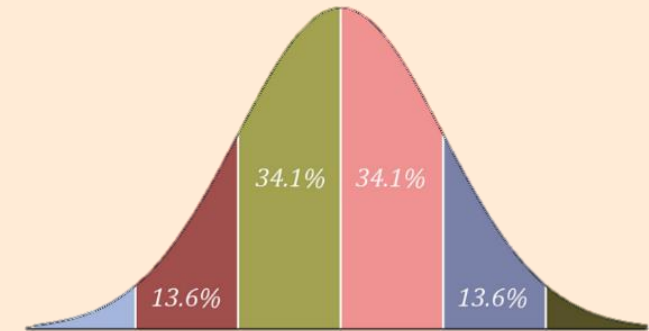
In computer
science

Experimental data
is almost never
normally distributed

Estimate

Only valid for a
normal distribution!!

Normal



Percentage of packets [%]

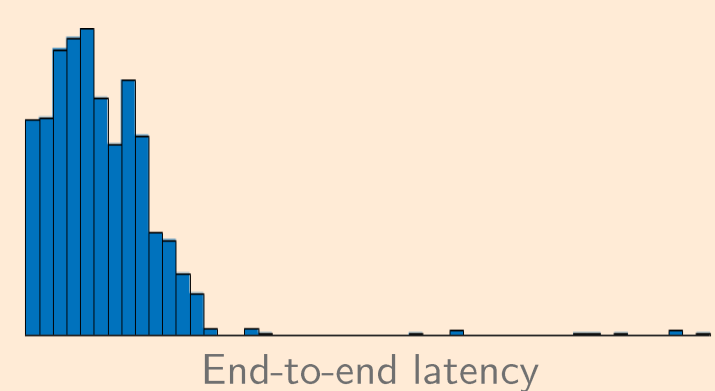
12

0

End-to-end latency

Not normal

39



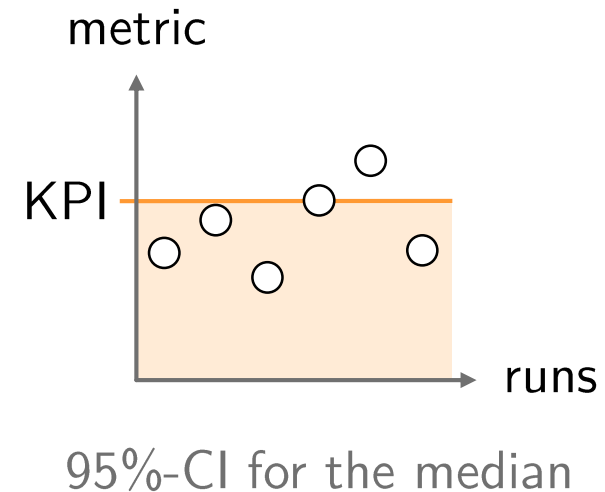
The Thompson's methods provides non-parametric CI for distribution percentiles

In TriScale

Key Performance Indicators (KPIs) are **percentiles** of the **distribution of metric** values

Compute confidence intervals using Thompson's method

- Define upper- and lower-bounds on the true percentile values for a certain confidence level
- The KPIs are defined as one such bounds



The Thompson's methods provides non-parametric CI for distribution percentiles

Probability of any P_p to be between two consecutive samples follows the binomial distribution

$$\mathbb{P} \{x_k \leq P_p \leq x_{k+1}\} = \binom{N}{k} p^k (1-p)^{N-k}$$

Lower-bound

$$\mathbb{P} \{x_m \leq P_p\} =$$

Upper-bound

$$\mathbb{P} \{x_{N-m+1} \geq P_{1-p}\} = 1 - \sum_{k=0}^{m-1} \binom{N}{k} p^k (1-p)^{N-k}$$

The Thompson's methods provides non-parametric CI for distribution percentiles

For any confidence c

For any percentile P_p

$$N \geq \frac{\log(1 - c)}{\log(1 - p)}$$

$$c = 0,95$$

$$p = 0,01$$

1-th percentile

\Rightarrow

$$N \geq 299$$

Minimal number
of runs in a series

We might want to rethink
our performance claims...

The four questions of experiment design

How long should a run be?

How many runs in a series?

What time span for a series?

How many series?

Find **rational answers**
to these questions

Quantify the trade-off between

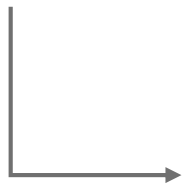
- experiment effort
- confidence in the results

Example – Performance evaluation of TTW

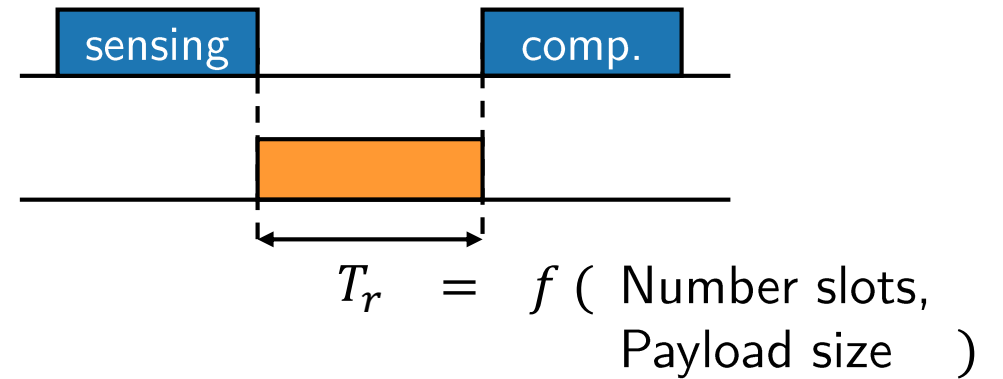
Is the model providing a safe and tight upper-bound?

Metric Maximum measured
round time T_r
among all nodes

KPI 95th percentile
95% confidence level



Aim to upper-bound the
95th percentile of the
maximum round time T_r
with 95% confidence



Round model f

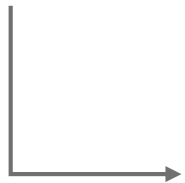
- Must be **safe** No overruns
- Should be **tight** Little wasted time

Metric

Maximum measured
round time T_r
among all nodes

KPI

95th percentile
95% confidence level



Aim to upper-bound the
95th percentile of the
maximum round time T_r
with 95% confidence

Minimal number
of runs in a series

$$c = 0,95$$

$$p = 0,05$$

⇒

$$N \geq 59$$

⇒

Scheduled 60 runs
per series,
randomly distributed
over one week

Slots per round

Payload

5

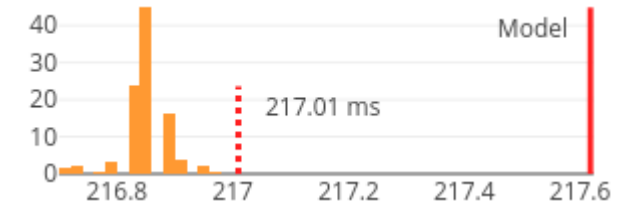
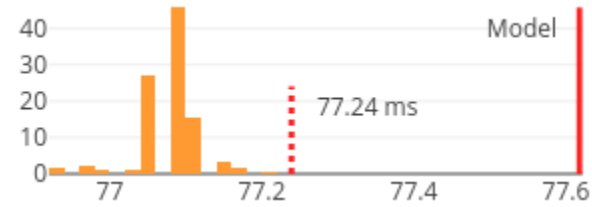
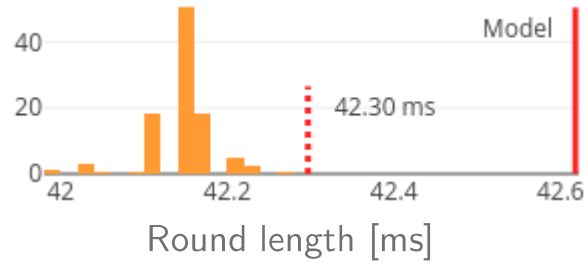
10

30

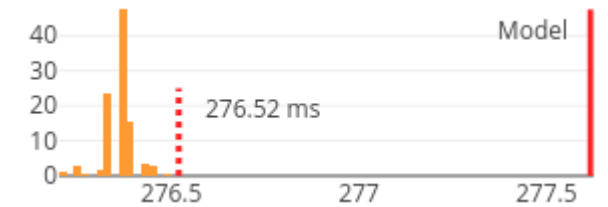
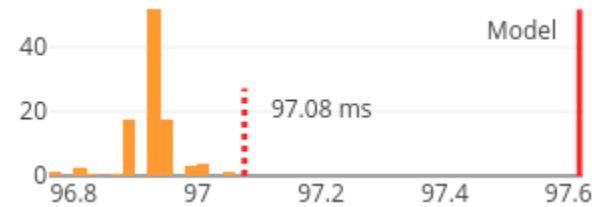
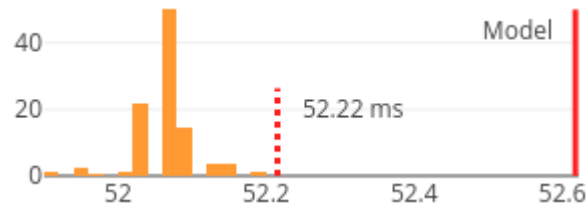
Bytes

Number of samples [%]

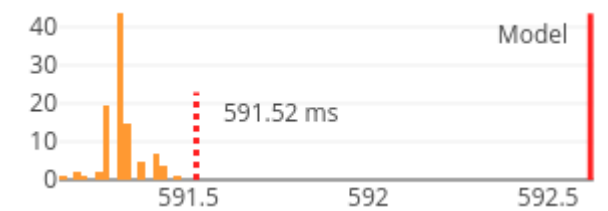
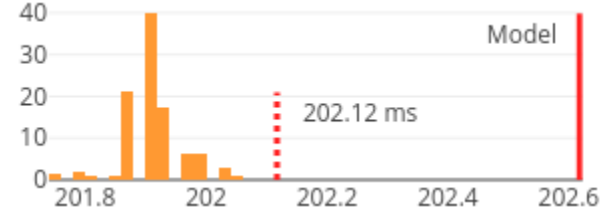
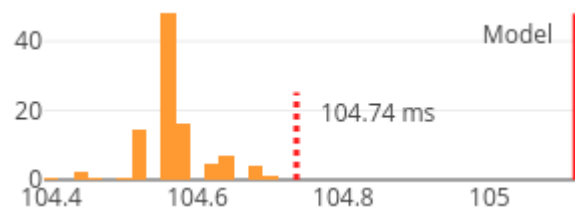
8



16



64



Round length [ms]

Slots per round

Payload

5

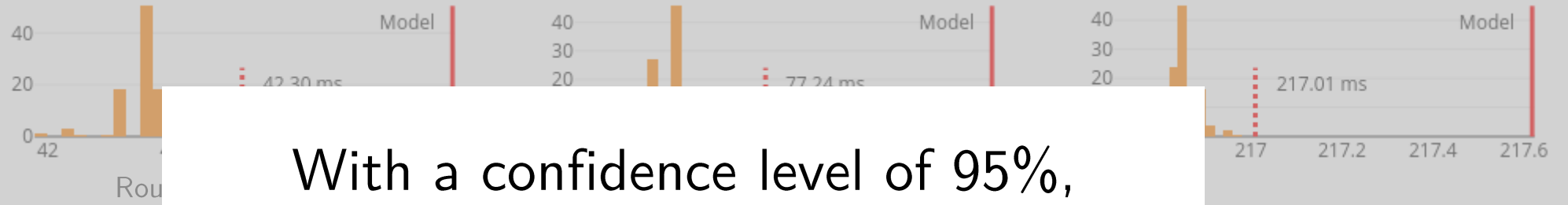
10

30

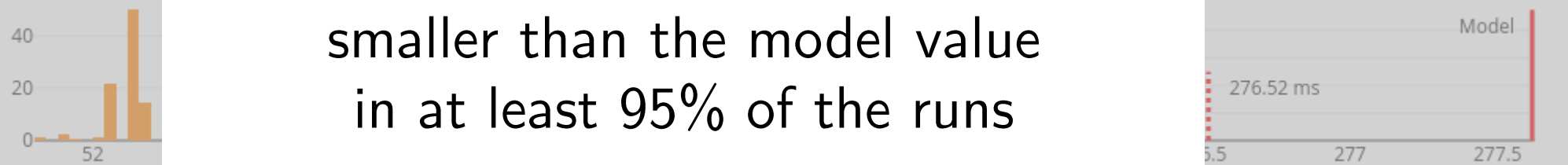
Bytes

Number of samples [%]

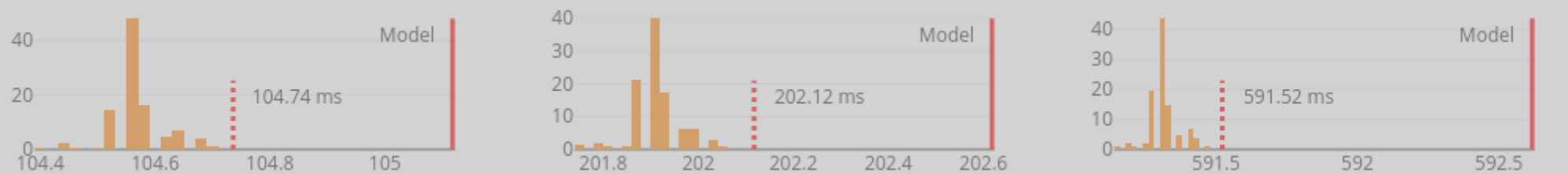
8



16



64



Round length [ms]

With a confidence level of 95%,
the maximum round time will be
smaller than the model value
in at least 95% of the runs



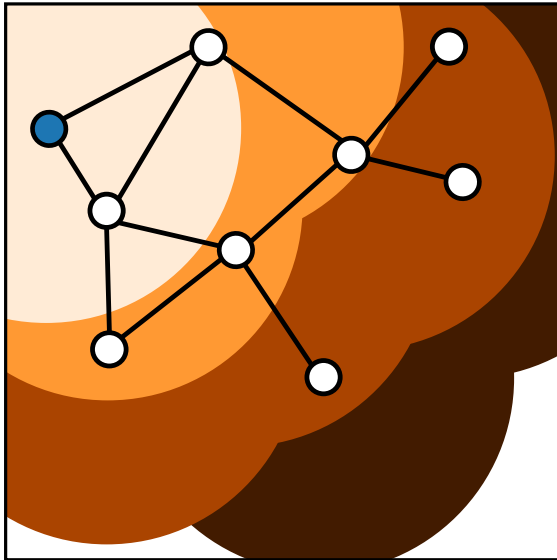
TriScale is a framework helping to design and analyze networking experiments

Implemented as Python package
 Open source

Includes

- Statistical methods and tests
- Experiment design helpers
- Automated data processing
- Dynamic visualizations

Leveraging Synchronous Transmissions for the Design of Real-time Wireless Cyber-Physical Systems



Real-time
System Designs

Tools and
Methods

DRP
Chapter 4

TTW
Chapter 5

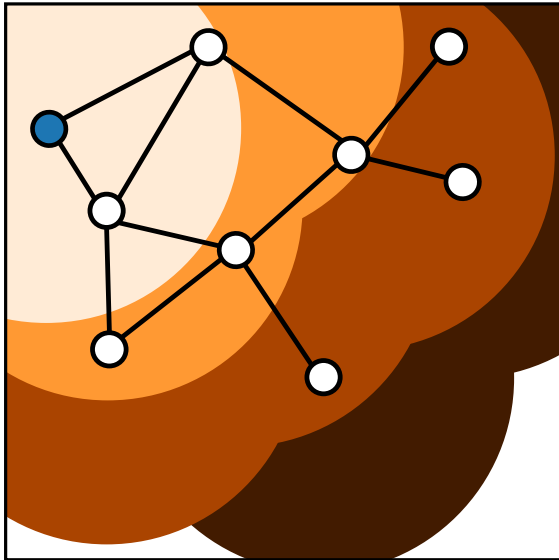
Baloo
Chapter 3

TriScale
Chapter 2

Synchronous
Transmissions

Networking
in General

All documents, software,
data are openly available



Links in the respective chapters



github.com/romain-jacob/doctoral-thesis

zenodo

DOI [10.5281/zenodo.3510184](https://doi.org/10.5281/zenodo.3510184)

Thank you all for your support!

Supervisor
Prof. Thiele

My wife
My friends
My parents

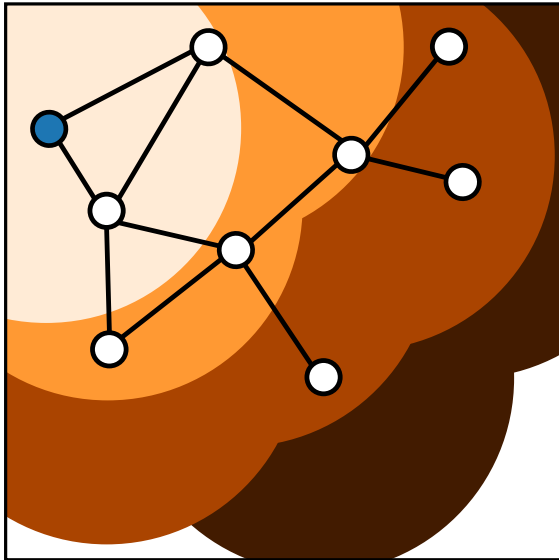
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Leveraging Synchronous Transmissions for the Design of Real-time Wireless Cyber-Physical Systems



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www.romainjacob.net
@RJacobPartner

The four questions of experiment design

How long should a run be?

such that

the metric would be the same
shall the run be longer

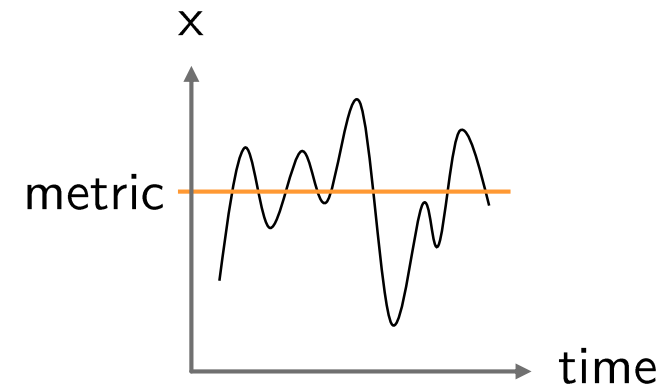
How many runs in a series?

What time span for a series?

How many series?

The Theil-Sen estimator performs robust and non-parametric linear regression

Computing a run's metric is a regression problem

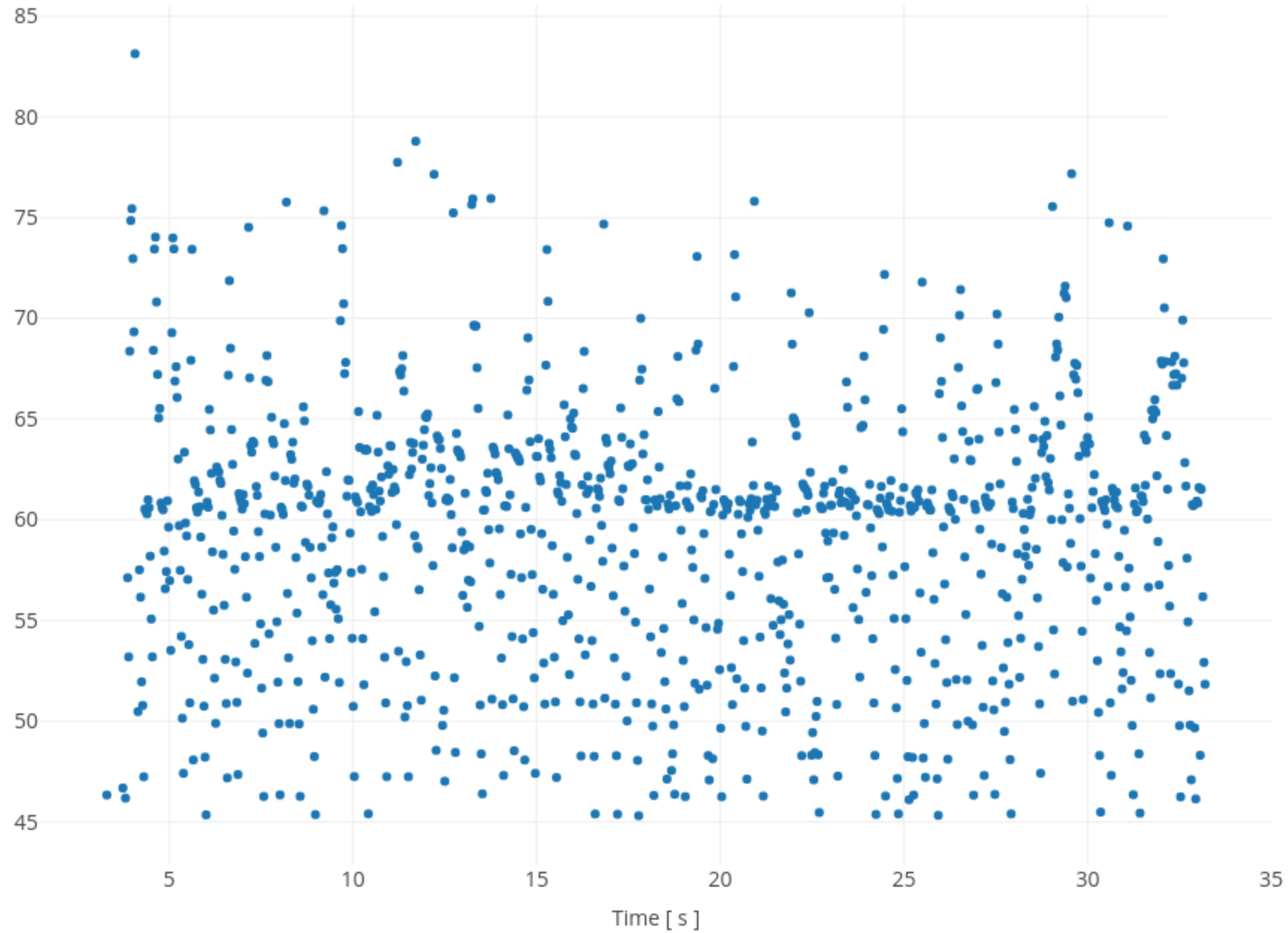


Non-parametric regression

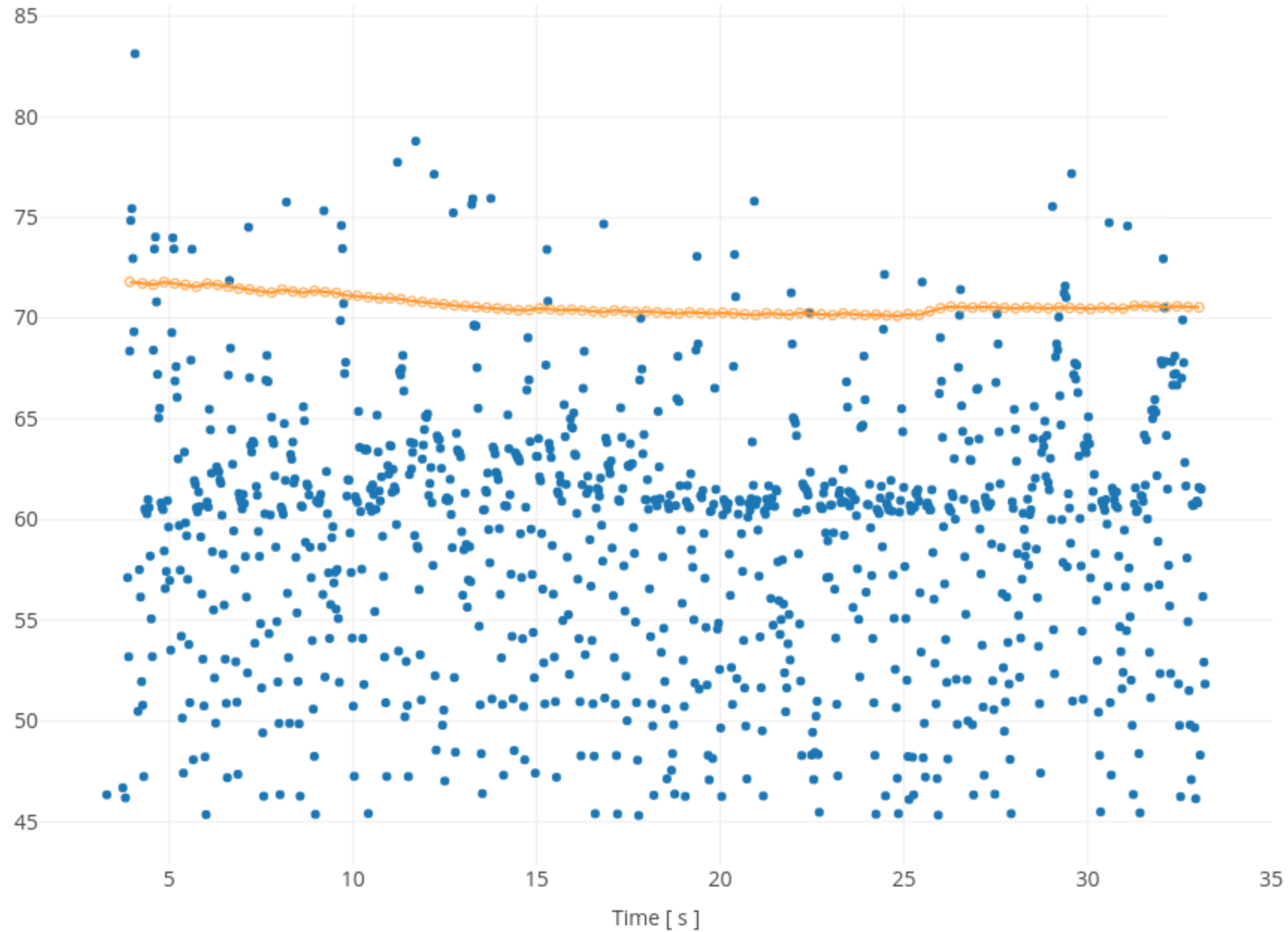
Theil-Sen estimator

- Slope is the median of all slopes between pairs of data points
- CI on the slope using the bootstrapping method

One-way delay [ms]



One-way delay [ms]



Compute the
metric on a
sliding window
of data points

One-way delay [ms]

73

72

71

70

69

68

5

10

15

20

25

30

Time [s]

Compute the
metric on a
sliding window
of data points

... zoom in

One-way delay [ms]

73

72

71

70

69

68

5

10

15

20

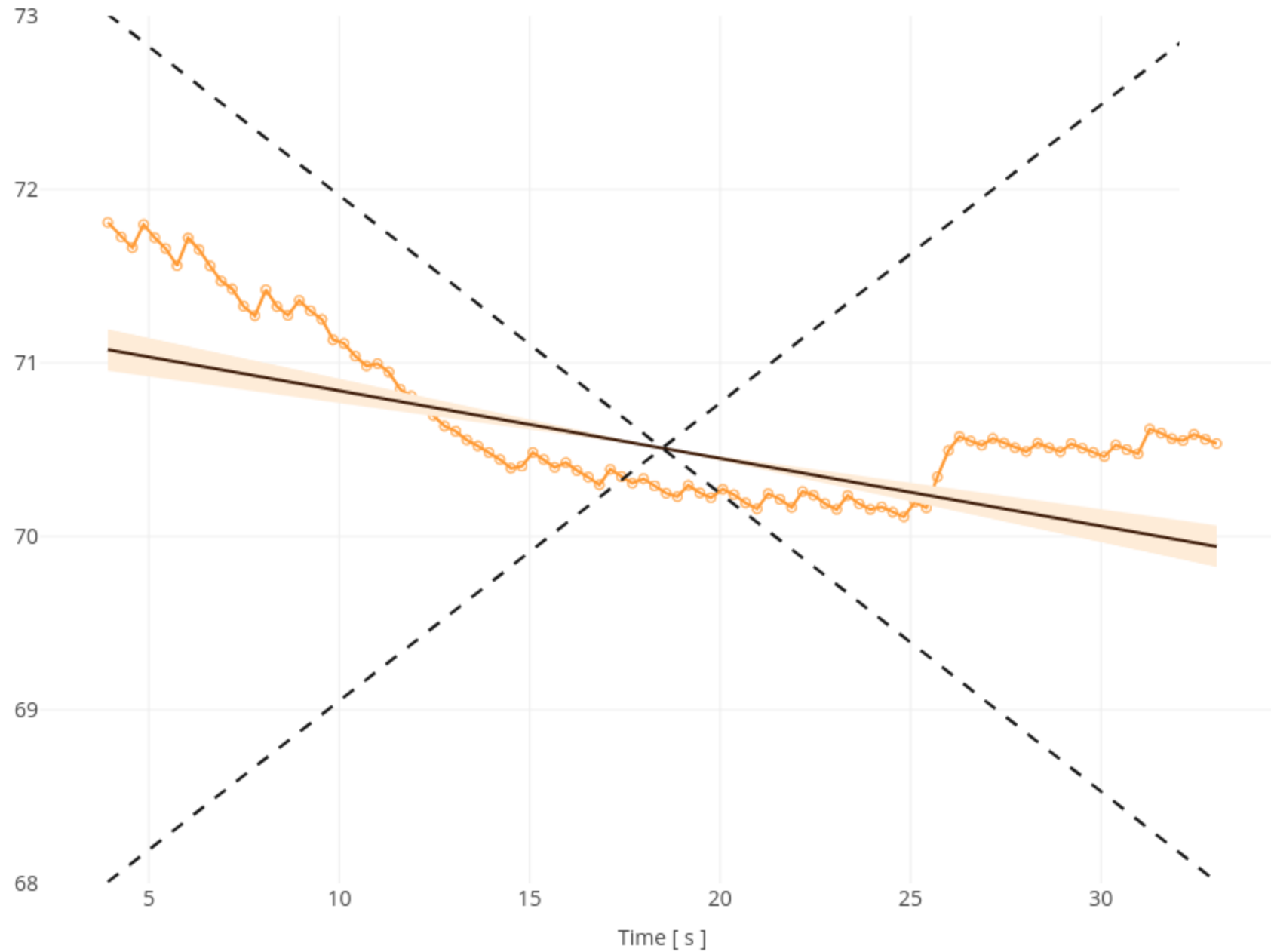
25

30

Time [s]

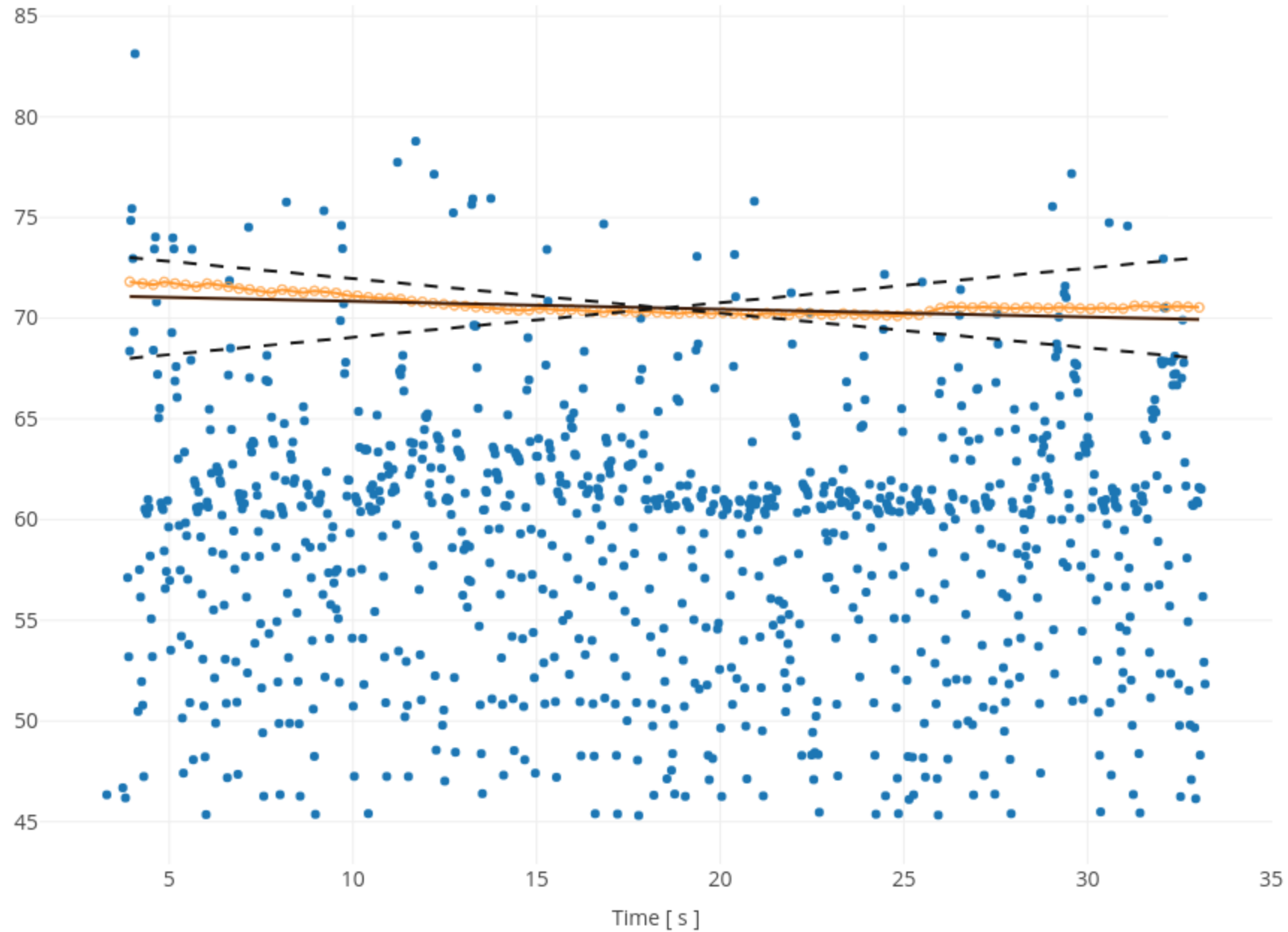
Compute the
Theil-Sen slope
and its 95%-CI

One-way delay [ms]



Compute the
Theil-Sen slope
and its 95%-CI
which must fit
in the tolerance
interval $\pm 5\%$

One-way delay [ms]



Take median of
metric samples
as run's metric

The application sets the requirements

Primary goals

Predictability

Adaptability

System must support

End-to-end real-time guarantees

- Change in traffic demand
- Mobility