**1. a)** On a single-CPU system, under what circumstances does a multithreaded program using kernel threads provide better performance (such as faster execution time) compared to a single threaded solution (that does not use asynchronous or event-based programming)? Explain with general principles. Give TWO example applications.

With the use of a multi-threaded program, it provides better performance by having the threads share memory and message passing. This allows address space to be shared allowing one application as data is being transferred. This is because multithreading has multiple threads that hold the code, data, and files with their own registers and stack.

Examples of applications include document editors such as Google Docs and video games such as Mario Kart. Google Docs is checking for spelling errors, text alignment, font type, and auto-saving for document files. Mario Kart has different sound effects from the characters, items, and karts while the race tracks have their hazards and background music enabled.

**b)** A new operating system provides a synchronization API and a library for user-level programs (i.e. like the pthread) for which the mutex lock and unlock operation are implemented with **test\_and\_set** like this:

Is this implementation of mutex synchronization correct for use in general-purpose user-level applications? What could go wrong? It helps to think of an example application, like the bounded-buffer problem or the dining philosophers.

That implementation of mutex synchronization is not correct because the mutex would be locked first and has the test\_and\_set shared as well as the restrictions such as 1 semaphore and 1 condvar. What could go wrong is the different threads having access to the mutex and not traversing. This can include bounded waiting, which happens when no process is waiting for a resource for an infinite amount of time.

When the process is done, there's no guarantee that another process will happen which will cause the 1st greedy process to hog the lock and go through the process of

releasing and taking. This will cause the whole process to go through starvation over time.

**c)** On a running Linux kernel (version > 2.6) at some point the thread\_info.preempt\_count field for a kernel task we call *A* is equal to 2. (Linux kernel synchronization is discussed in the textbook).

Answer these questions:

**c1)** Is task A currently preemptable? Explain.

Task A isn't preemptable because thread\_info.preempt\_count is considered a non-zero, which would make it hold 2 locks.

**c2)** What is new value of thread\_info.preempt\_count field for task A after it acquires a new lock? Explain.

thread\_info.preempt\_count will have 3 as the new value if task A acquires a new lock because having the new lock available will have thread\_info.preempt\_count increment by 1.

c3) What is the condition for kernel task A to be safely interruptible?

Kernel task A becomes interruptible by having the lock not activated, which is done by having the thread\_info.preempt\_count value checked.

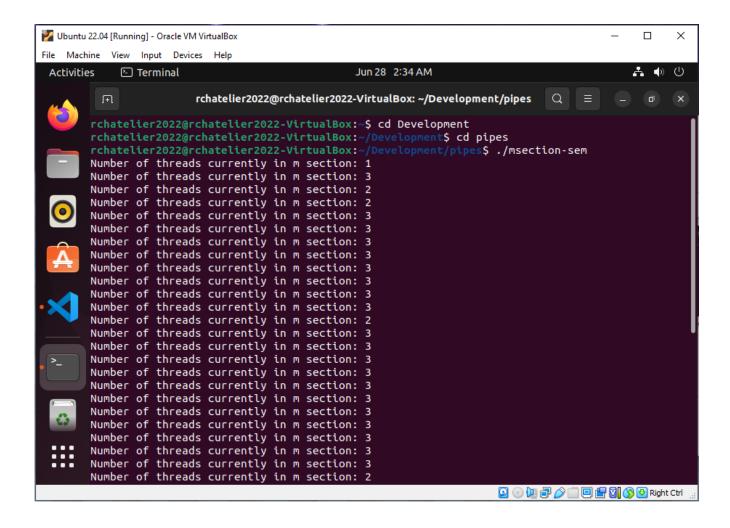
- **c4)** Assuming that all locks held by task A are spinlocks, how many CPUs are on that computer?
- 9 CPUs would be on that computer since there are threads blocked by the spinlock. This will cause the resources in a loop that doesn't stop, which means that the CPU will require more overall.
- **2. a)** Write a C program called **msection-sem.c** using the *pthread* library that implements the algorithm above.

```
#include <string.h>
#include <pthread.h>
#include <stdbool.h>
#include <semaphore.h>
```

```
#include <stdlib.h>
#include <unistd.h>
#include <stdio.h>
void enter();
void leave();
void doWork();
void doCriticalWork();
// amount of threads
int m = 3;
int counter = 0;
// semaphore object
sem_t semObject;
int main(int argc, char** argv) {
sem init (&semObject, 0, 3);
long i;
// number of threads
long n = 10;
pthread attr t attr;
pthread t *vtid = (pthread t *) calloc(n, sizeof(pthread t));
if (vtid == NULL) {
 printf("Error \n\n");
 exit(1);
```

```
pthread attr init(&attr);
pthread attr setdetachstate(&attr, PTHREAD CREATE JOINABLE);
  int r = pthread create(&vtid[i], &attr, doWork, (void*)i);
   printf("Error %d\n", r);
   exit(2);
 pthread join(vtid[i], NULL);
printf("all threads finished\ncounter=%ld\n", counter);
pthread attr destroy(&attr);
free (vtid);
return 0;
void doWork() {
```

```
enter (&semObject);
     sleep(69);
   leave (&semObject);
void doCriticalWork() {
 pthread_t vT = pthread_self();
 printf("Number of threads currently in m section: %ld\n", counter);
void enter(sem t *semObject) {
 sem wait (semObject);
 counter++;
void leave(sem t *semObject) {
  counter--;
 sem_post(semObject);
```



**b)** Write a C program called **msection-condvar.c** using the *pthread* library that implements the algorithm above. The *enter()* and *leave()* functions must use only **one condition variable** and one or more **mutex**es for synchronization.

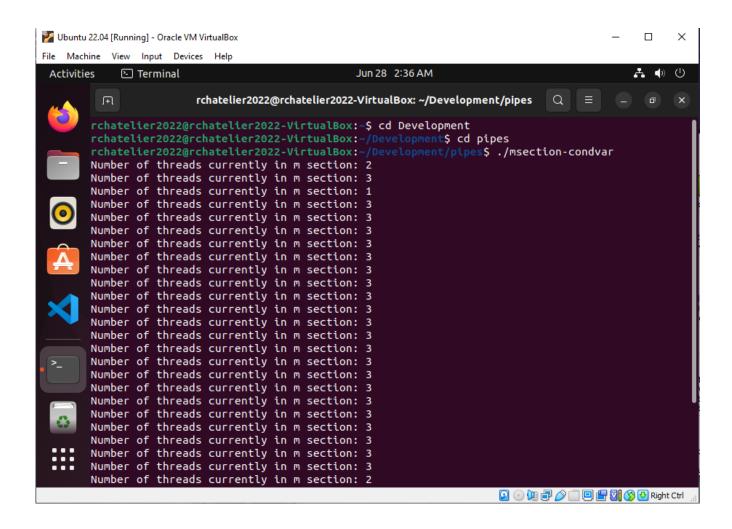
```
#include <string.h>
#include <pthread.h>
#include <stdbool.h>
#include <stdlib.h>
#include <unistd.h>
#include <stdio.h>

void enter();
void leave();
```

```
void doWork();
void doCriticalWork();
int m = 3;
int counter = 0;
pthread mutex t cvmutex;
pthread cond t condvar;
int main (int argc, char** argv) {
  pthread mutex init (&cvmutex, NULL);
  pthread cond init (&condvar, NULL);
pthread t *vtid = (pthread t *) calloc(n, sizeof(pthread t));
if (vtid == NULL) {
  printf("Error \n\n");
  exit(1);
 pthread attr init(&attr);
pthread attr setdetachstate(&attr, PTHREAD CREATE JOINABLE);
```

```
int r = pthread create(&vtid[i], &attr, doWork, (void*)i);
    printf("Error %d\n", r);
    exit(2);
pthread join(vtid[i], NULL);
printf("all threads finished\ncounter=%d\n", counter);
pthread attr destroy(&attr);
 free(vtid);
void doWork() {
   doCriticalWork();
       sleep(69);
   leave (&cvmutex, &condvar);
```

```
void doCriticalWork() {
  pthread t vT = pthread self();
  printf("Number of threads currently in m section: %d\n", counter);
void enter(pthread cond t *cThread, pthread mutex t *muTex ) {
 pthread mutex lock(muTex);
  while(counter >= m) {
      pthread cond wait(cThread, muTex);
  pthread mutex unlock(muTex);
  counter++;
void leave(pthread cond t *cThread, pthread mutex t *muTex) {
  counter--;
  pthread mutex unlock(muTex);
  pthread cond signal(cThread);
  pthread mutex unlock(muTex);
```



**3. a)** Implement a *barrier* for pthreads in C++ using pthread condition variables and mutexes in a file called *barrier.cc*.

```
#include <string.h>
#include <pthread.h>
#include <stdbool.h>
#include <stdlib.h>
#include <unistd.h>
#include <stdio.h>

woid thread();

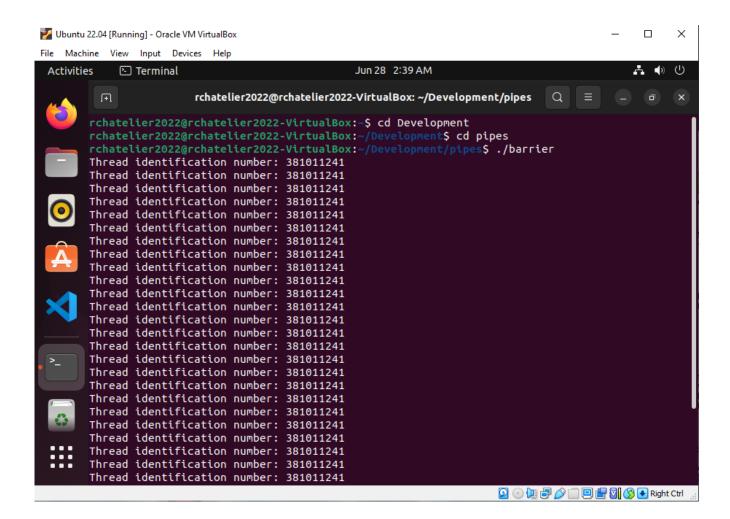
void *thread_fun(void *param);
```

```
pthread mutex t cvmutex;
pthread cond t condvar;
class Barrier {
 Barrier(int n)
     pthread_mutex_init (&cvmutex, NULL);
     pthread cond init(&condvar, NULL);
void wait() {
  if(counter < 0) {</pre>
      pthread_mutex_lock(&cvmutex);
     pthread_cond_wait(&condvar, &cvmutex);
      pthread_mutex_unlock(&cvmutex);
      pthread_mutex_lock(&cvmutex);
      pthread cond broadcast(&condvar);
      pthread_mutex_unlock(&cvmutex);
```

```
private:
int x;
pthread_mutex_t cvmutex;
pthread cond t condvar;
};
Barrier shield(7);
int main(int argc, char** argv) {
long i;
// number of threads
long n = 10;
pthread attr t attr;
pthread_t *vtid = (pthread_t *) calloc(n, sizeof(pthread_t));
if (vtid == NULL) {
  printf("Error \n\n");
  exit(1);
pthread attr init(&attr);
// sets the detach state attribute
pthread attr setdetachstate(&attr, PTHREAD CREATE JOINABLE);
  int r = pthread_create(&vtid[i], &attr, thread_fun, (void*)i);
```

```
printf("Error %d\n", r);
    exit(2);
  pthread join(vtid[i], NULL);
printf("all threads finished\ncounter=%d\n", shield.counter);
pthread_attr_destroy(&attr);
free(vtid);
return 0;
void *thread_fun(void *param) {
  while(true) {
     thread();
void thread(){
pthread_t results = pthread_self();
```

```
printf("Thread identification number: %d\n", thread);
sleep(69);
```



**b)** Write a thread function in file barrier.cc that demonstrates in a meaningful way how your Barrier object works. It should look like the code in function *thread fun* above.

```
#include <string.h>
#include <pthread.h>
```

```
#include <stdbool.h>
#include <stdlib.h>
#include <unistd.h>
#include <stdio.h>
void thread();
void *thread fun(void *param);
pthread mutex t cvmutex;
pthread cond t condvar;
class Barrier {
 int counter;
  Barrier(int n)
      pthread_mutex_init (&cvmutex, NULL);
      pthread_cond_init(&condvar, NULL);
void wait() {
  if(counter < 0) {</pre>
      pthread_mutex_lock(&cvmutex);
      counter++;
      pthread cond wait(&condvar, &cvmutex);
     pthread_mutex_unlock(&cvmutex);
```

```
pthread_mutex_lock(&cvmutex);
     pthread cond broadcast(&condvar);
      pthread_mutex_unlock(&cvmutex);
private:
int x;
pthread_mutex_t cvmutex;
pthread cond t condvar;
};
Barrier shield(7);
int main (int argc, char** argv) {
long i;
long n = 10;
pthread_attr_t attr;
pthread_t *vtid = (pthread_t *) calloc(n, sizeof(pthread_t));
if (vtid == NULL) {
  printf("Error \n\n");
  exit(1);
```

```
pthread attr init(&attr);
pthread attr setdetachstate(&attr, PTHREAD CREATE JOINABLE);
 for (i=0; i<n; i++) {
  int r = pthread_create(&vtid[i], &attr, thread_fun, (void*)i);
    printf("Error %d\n", r);
   exit(2);
 pthread join(vtid[i], NULL);
printf("all threads finished\ncounter=%d\n", shield.counter);
pthread attr destroy(&attr);
free (vtid);
return 0;
void *thread fun(void *param) {
 while(true) {
    shield.wait();
```

```
thread();
void thread(){
pthread_t results = pthread_self();
 while(adding <= 10) {</pre>
     printf("Threads printed: %d\n", bundle);
     printf("Thread identification number: %d\n", thread);
     sleep(69);
     adding++;
     bundle++;
```

