## **ECE374 SP23 HW3**

## Contributors

Zhirong Chen (zhirong4)

Ziyuan Chen (ziyuanc3)

## Problem 1

A *finite-state transducer* (FST) gives an output based on the transition instead of the current state. It is defined by a 5-tuple:

$$(\Sigma, \Gamma, Q, \delta, s)$$

The output alphabet of a  $FST_{AR}$  consists of two signals, namely accept and reject ( $\Gamma=\{A,R\}$ ). We say that  $L(FST_{AR})$  represents the language consisting of all strings that end with an accept (A) output signal.

Prove that  $L(FST_{
m AR})$  represents the class of regular languages.

## Solution

**Proof.** We prove that a language is regular (i.e., can be represented by a DFA) *iff* it can be represented by a  $FST_{\rm AR}$ .

*If.* For any given  $FST_{
m AR}$ 

$$M_{T1}=(\Sigma_1,\Gamma,Q_{T1},\delta_{T1},s_{T1})$$

, we construct a DFA

$$M_{D1} = (\Sigma_1, Q_{D1}, \delta_{D1}, s_{D1}, A_{D1})$$

such that  $L(M_{D1}) = L(M_{T1})$ , where

- $Q_{D1} = Q_{T1} \times \Gamma$
- $ullet \ \delta_{D1}((q,b),a) = \delta_{T1}(q,a), \quad orall q \in Q_{T1}, a \in \Sigma_1, b \in \Gamma$
- $s_{D1}=(s_{T1},\mathrm{R})$
- $\bullet \ \ A_{D1}=\{(q,\mathrm{A}) \mid q \in Q_{T1}\}$

Here  $\delta_{T1}:Q_{T1} imes \Sigma_1 o Q_{T1} imes \Gamma$  and  $\delta_{D1}:Q_{D1} imes \Sigma_1 o Q_{D1}.$ 

The core idea is to explicitly encode the last output signal in the states. There is a clear boundary between the *next-state logic* and *output logic*; thus, the expression of  $\delta_{D1}$  is unrelated to b.

We have proven that any language representable by a  $FST_{
m AR}$  is regular.

Only if. For any given regular language L represented by a DFA

$$M_{D2} = (\Sigma_2, Q_2, \delta_{D2}, s_2, A_{D2})$$

, we construct a  $FST_{
m AR}$ 

$$M_{T2}=(\Sigma_2,\Gamma,Q_2,\delta_{T2},s_2)$$

such that  $L(M_{T2}) = L(M_{D2})$ , where

$$oldsymbol{eta}_{T2}(q,a) = egin{cases} (\delta_{D2}(q,a),\mathrm{A}) & ext{if } \delta_{D2}(q,a) \in A_{D2}, \ (\delta_{D2}(q,a),\mathrm{R}) & ext{if } \delta_{D2}(q,a) 
otin A_{D2}, \end{cases} \quad orall q \in Q_2, a \in \Sigma_2$$

The idea is to encode the status of the next state in the arcs pointing to it.

We have proven that any regular language can be represented by a  $FST_{
m AR}.$  **Q.E.D.**