ECE374 SP23 HW4

Contributors

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Problem 5

Given an arbitrary array $A[1:n]$, describe an algorithm to determine in $O(n)$ time whether A contains more than $\frac{n}{4}$ copies of any value. **Do not use hashing, or radix sort, or any other method that depends on the precise input values.**

Solution

Intuition

- In a *sorted* array, elements with the same value occupy a contiguous subarray. Such a subarray with length over $\frac{n}{4}$ must cover at least one of the array's three quarters.
- In other words, if there is a value that appears more than $\frac{n}{4}$ times, **it must be among the three quarters**. This also applies to unsorted arrays.

Algorithm

$\mathrm{IfFrequentValueExists}(A, n)$
    for $i \leftarrow 1$ to 3
        $\mathrm{quarter} \leftarrow \mathrm{QuickSelect}(A, \lfloor \frac{in}{4} \rfloor)$
        if $\mathrm{Count}(A, \mathrm{quarter}) > \frac{n}{4}$
            return True
    return False

Runtime Analysis

- $\mathrm{QuickSelect}(A, k)$ (covered in lecture) finds the k^{th} smallest element of A in $O(n)$ time.
- $\mathrm{Count}(A, \mathrm{target})$ scans the *unsorted* A and counts the occurrences of target. Evidently, it runs in $O(n)$ time.
- In total, there are **6 calls** to $O(n)$ algorithms regardless of the array's length. Therefore, $\mathrm{IfFrequentValueExists}(A, n)$ runs in $O(n)$ time.