# **Language Translators**

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<sup>\*</sup> Course website: https://www.cs.columbia.edu/ rgu/courses/4115/spring2019

<sup>\*\*</sup> These slides are borrowed from Prof. Edwards.

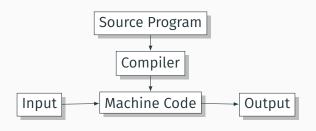
#### What is a Translator?

A programming language is a notation that a person and a computer can both understand.

- It allows you to express what is the **task** to compute
- It allows a computer to **execute** the computation task

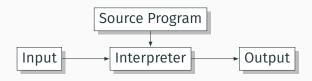
A translator translates what you express to what a computer can execute.

## Compiler



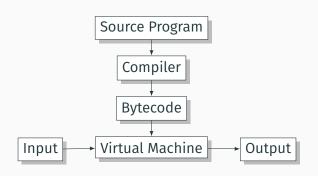
- Pros: translation is done once and for all; optimize code and map identifiers at compile time.
- Cons: long compilation time; hard to port.

#### Interpreter



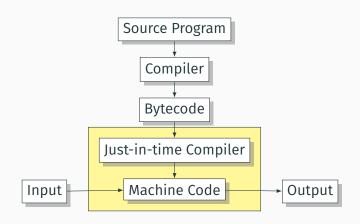
- **Pros**: source code distribution; short development cycle.
- Cons: translation is needed every time a statement is executed; lack optimization; map identifiers repeatedly.

## **Bytecode Interpreter**



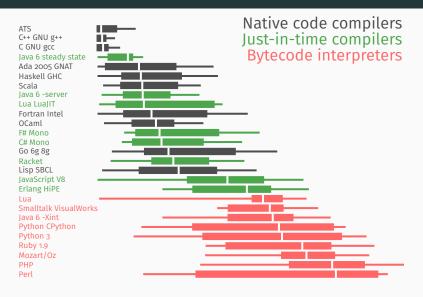
- Pros: bytecode is highly compressed and optimized; bytecode distribution.
- **Cons**: compilation overhead + interpreter overhead.

# Just-In-Time Compiler



- **Pros**: compile and optimize many sections just before the execution; bytecode distribution.
- **Cons**: compilation overhead + warm-up overhead.

#### **Language Speeds Compared**

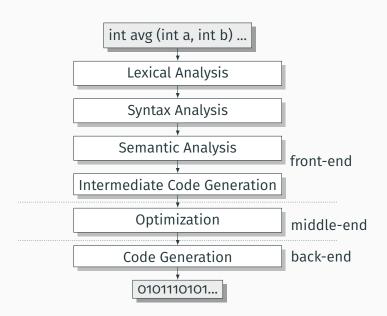


# Compilation Phases

# **Compiling a Simple Program**

```
int avg(int a, int b)
  return (a + b) / 2;
       Compiler
      0101110101...
```

## **Compilation Phases**



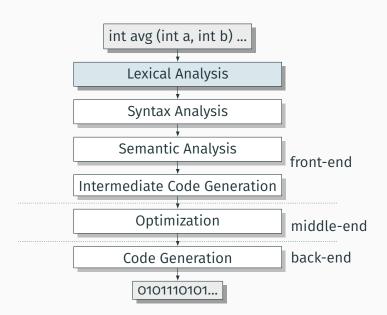
# **What the Compiler Sees**

```
int avg(int a, int b)
{
  return (a + b) / 2;
}
```

```
i n t SP a v g ( i n t SP a , SP i n t SP b ) NL { NL SP SP r e t u r n SP ( a SP + SP b ) SP / SP 2 ; NL } NL
```

Just a sequence of characters

# **Lexical Analysis**



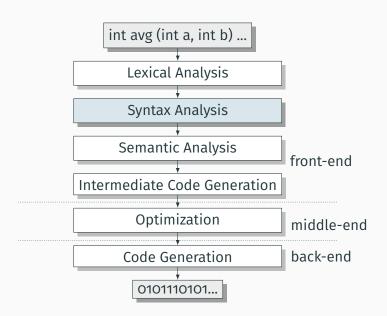
# **Lexical Analysis Gives Tokens**

```
int avg(int a, int b)
{
  return (a + b) / 2;
}

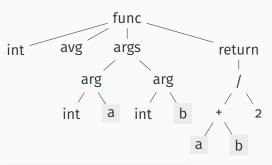
int avg((int a , int b)) { return (a + b)
) / 2;}
```

- · A stream of tokens; whitespace, comments removed.
- Throw errors when failing to create tokens: malformed strings or numbers or invalid characters (such as non-ASCII characters in C).

# **Syntax Analysis**



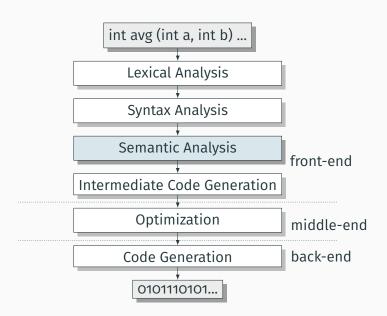
# Syntax Analysis Gives an Abstract Syntax Tree



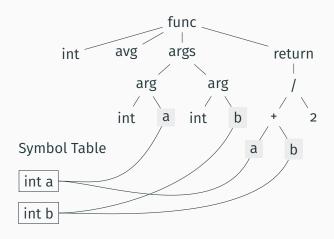
```
int avg(int a, int b)
{
   return (a + b) / 2;
}
```

 Syntax analysis will throw errors if "}" is missing. Lexical analysis will not.

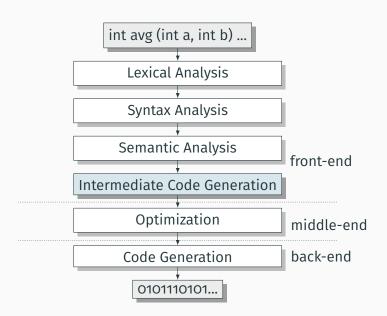
# **Semantic Analysis**



# Semantic Analysis: Resolve Symbols; Verify Types



#### **Intermediate Code Generation**



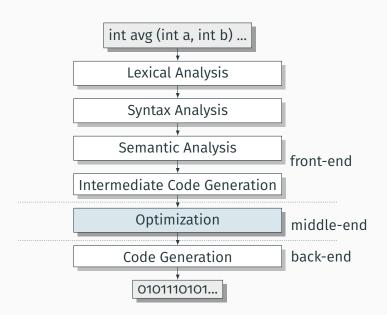
#### **Translation into 3-Address Code**

```
int avg(int a, int b)
{
  return (a + b) / 2;
}
```

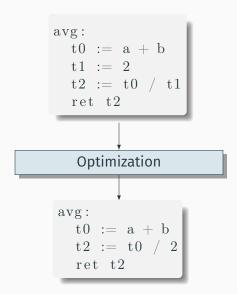
Idealized assembly language w/ infinite registers

```
\begin{array}{l} avg: \\ t0 := a + b \\ t1 := 2 \\ t2 := t0 \ / \ t1 \\ ret \ t2 \end{array}
```

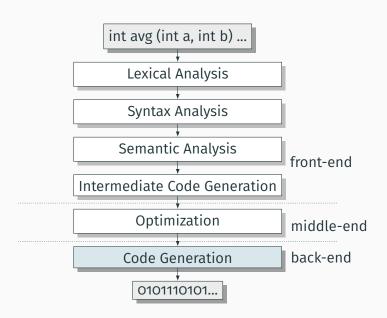
# **Optimization**



# **Optimization**



#### **Code Generation**



## **Generation of x86 Assembly**

```
t0 := a + b
                t2 := t0 / 2
                ret t2
               Code Generation
avg: pushl %ebp
                # save BP
     movl %esp,%ebp
     movl 8(%ebp), %eax # load a from stack
     movl 12(%ebp),%edx # load b from stack
     addl \%edx,\%eax \# a += b
     shr $1,\%eax # a = 2
     ret
```

avg: