

Language Translators

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* Course website: <https://www.cs.columbia.edu/~rgu/courses/4115/spring2019>

** These slides are borrowed from Prof. Edwards.

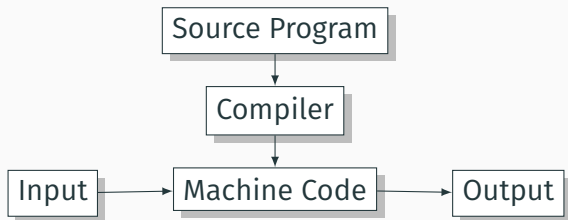
What is a Translator?

A programming language is a notation that a person and a computer can both understand.

- It allows you to express what is the **task** to compute
- It allows a computer to **execute** the computation task

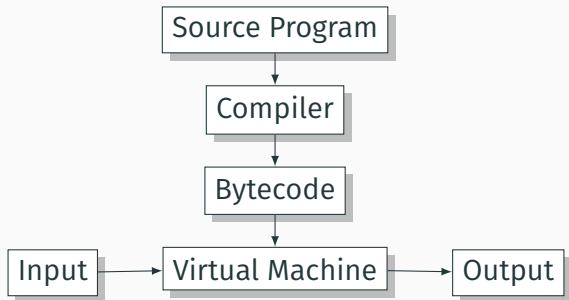
A translator translates what you express to what a computer can execute.

Compiler



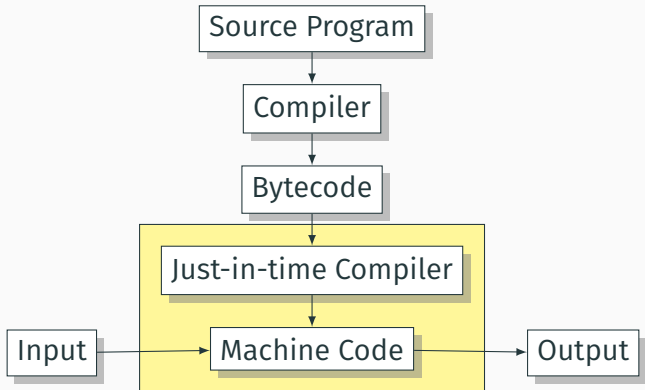
- **Pros:** translation is done once and for all; optimize code and map identifiers at compile time.
- **Cons:** long compilation time; hard to port.

Bytecode Interpreter



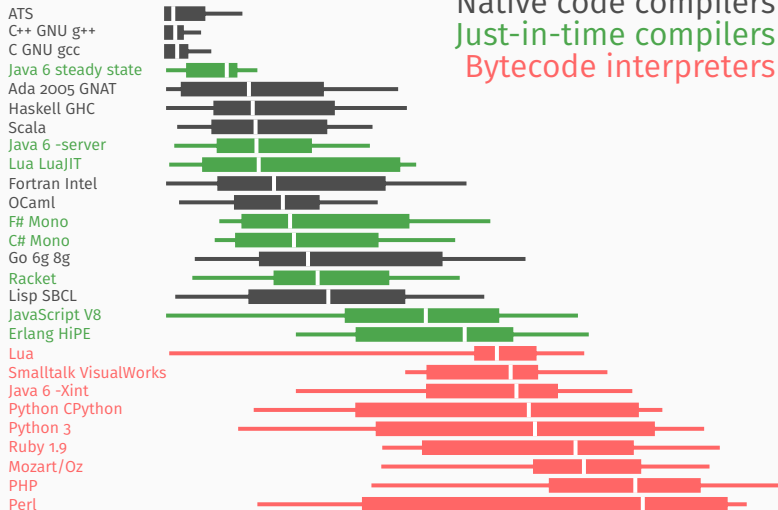
- **Pros:** bytecode is highly compressed and optimized; bytecode distribution.
- **Cons:** compilation overhead + interpreter overhead.

Just-In-Time Compiler



- **Pros:** compile and optimize many sections just before the execution; bytecode distribution.
- **Cons:** compilation overhead + warm-up overhead.

Language Speeds Compared



Source: <http://shootout.alioth.debian.org/>

Compilation Phases

Compiling a Simple Program

```
int avg(int a, int b)
{
    return (a + b) / 2;
}
```

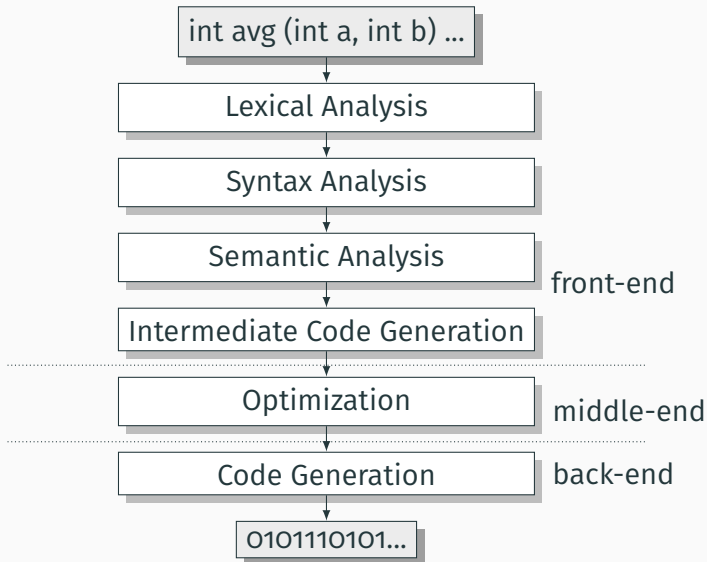
↓

Compiler

↓

0101110101...

Compilation Phases



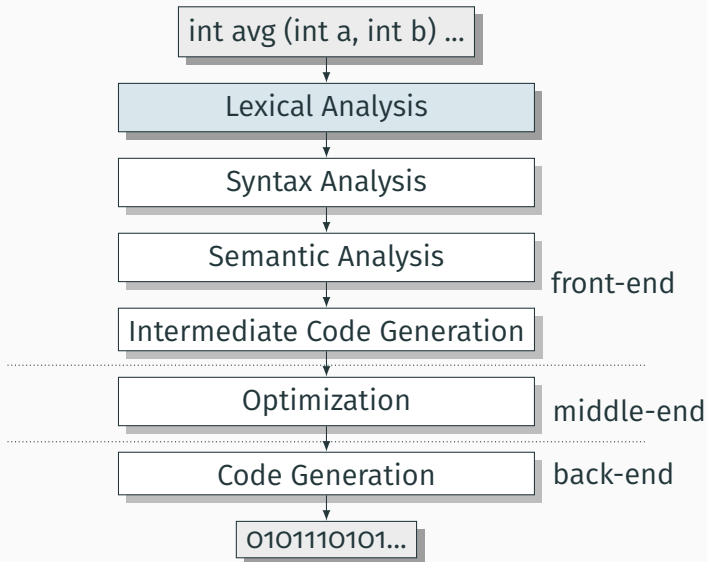
What the Compiler Sees

```
int avg(int a, int b)
{
    return (a + b) / 2;
}
```

```
i n t SP a v g ( i n t SP a , SP i n t SP b ) NL
{ NL
SP SP r e t u r n SP ( a SP + SP b ) SP / SP 2 ; NL
} NL
```

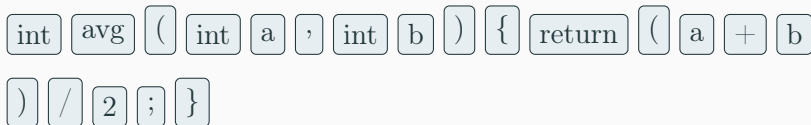
Just a sequence of characters

Lexical Analysis



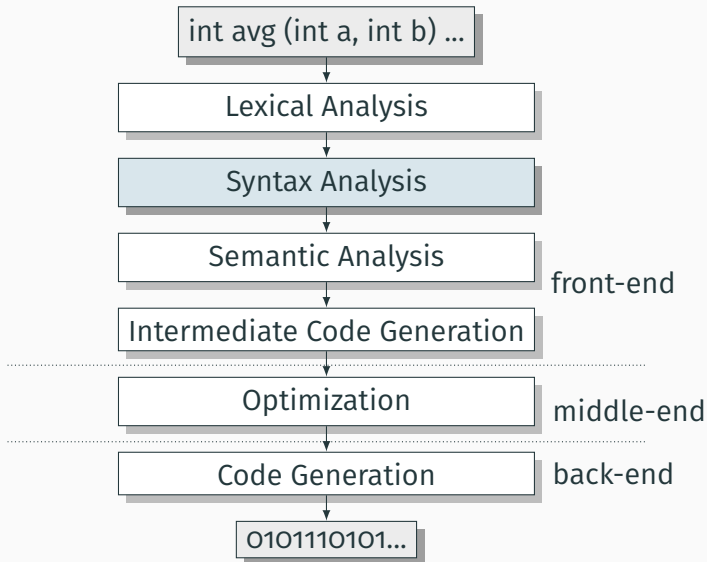
Lexical Analysis Gives Tokens

```
int avg(int a, int b)
{
    return (a + b) / 2;
}
```

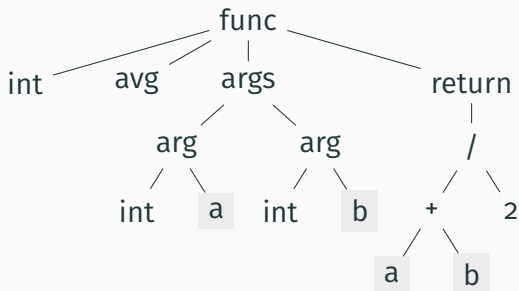


- A stream of tokens; whitespace, comments removed.
- Throw errors when failing to create tokens: malformed strings or numbers or invalid characters (such as non-ASCII characters in C).

Syntax Analysis



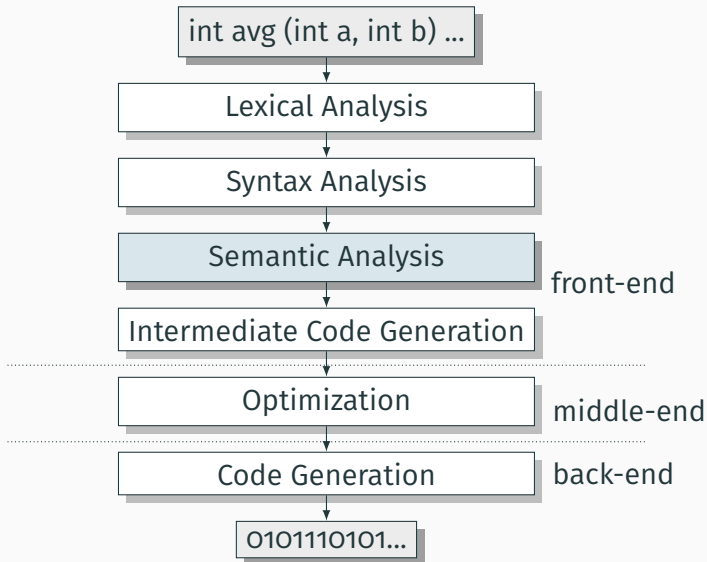
Syntax Analysis Gives an Abstract Syntax Tree



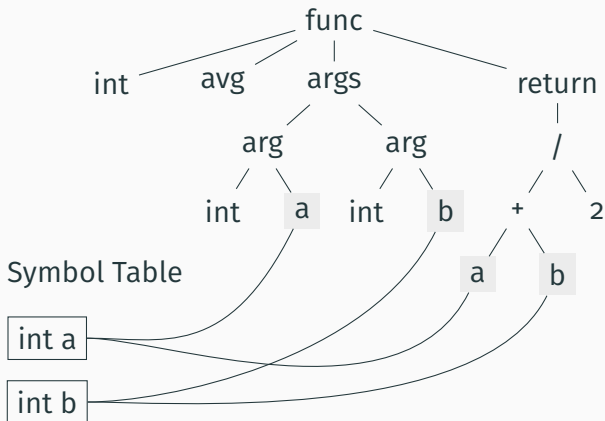
```
int avg(int a, int b)
{
    return (a + b) / 2;
}
```

- Syntax analysis will throw errors if “}” is missing. Lexical analysis will not.

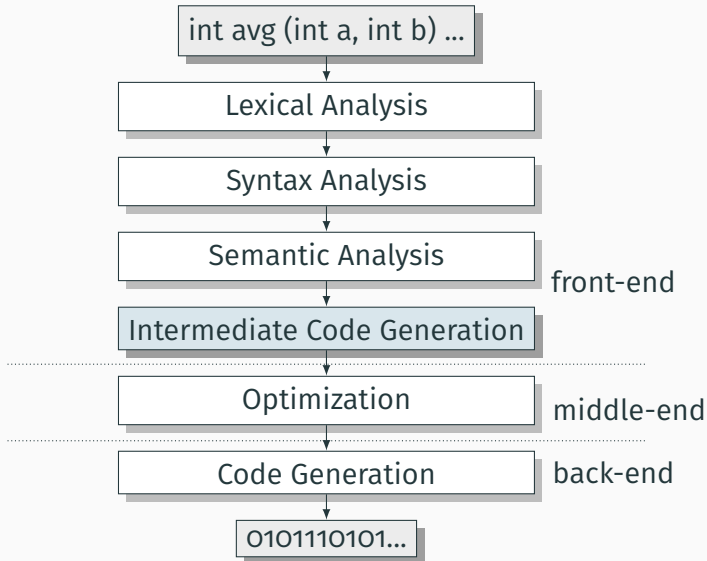
Semantic Analysis



Semantic Analysis: Resolve Symbols; Verify Types



Intermediate Code Generation



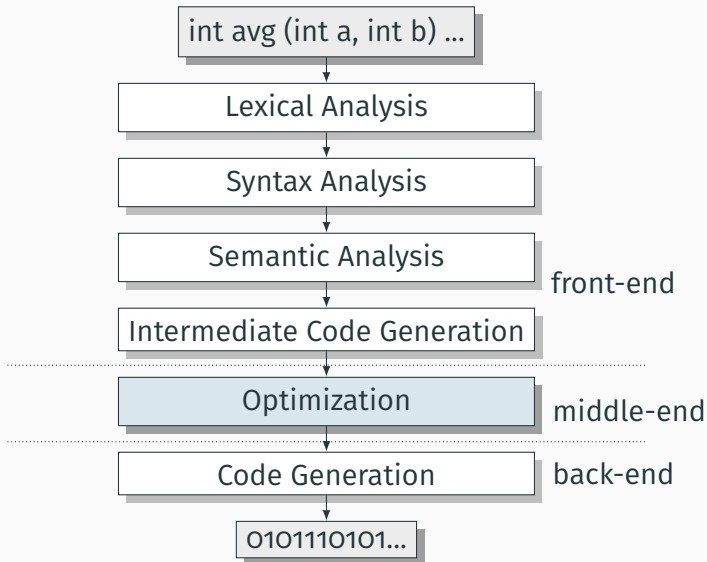
Translation into 3-Address Code

```
int avg(int a, int b)
{
    return (a + b) / 2;
}
```

Idealized assembly language w/ infinite registers

```
avg:
    t0 := a + b
    t1 := 2
    t2 := t0 / t1
    ret t2
```

Optimization



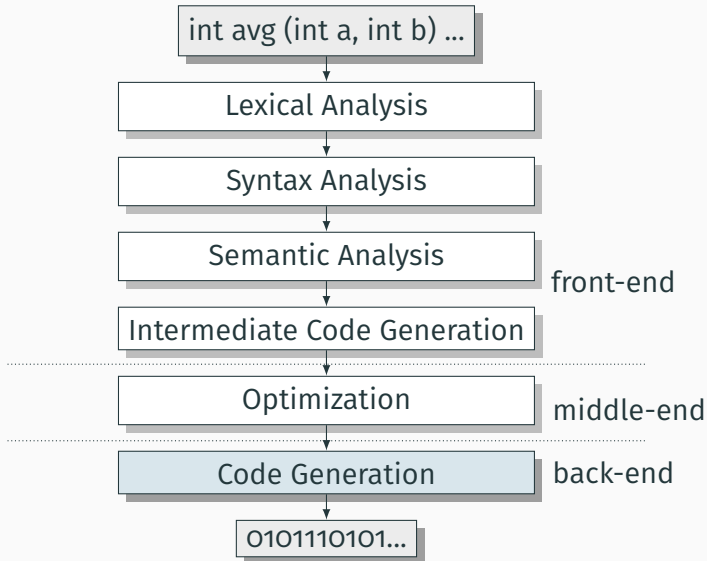
Optimization

```
avg:  
  t0 := a + b  
  t1 := 2  
  t2 := t0 / t1  
  ret t2
```

Optimization

```
avg:  
  t0 := a + b  
  t2 := t0 / 2  
  ret t2
```

Code Generation



Generation of x86 Assembly

```
avg:  
    t0 := a + b  
    t2 := t0 / 2  
    ret t2
```

Code Generation

```
avg:    pushl %ebp                # save BP  
        movl %esp,%ebp  
        movl 8(%ebp),%eax        # load a from stack  
        movl 12(%ebp),%edx       # load b from stack  
  
        addl %edx,%eax           # a += b  
        shr  $1,%eax             # a /= 2  
        ret
```