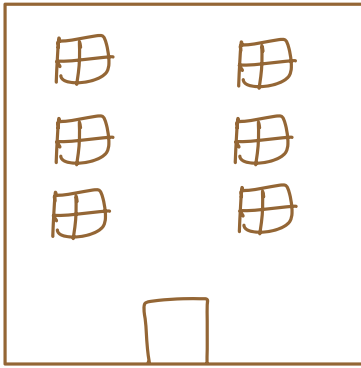


Hashing 3 : Internal Implementation & Problems

mohan



Check if any particular room is available?

Room no.	Free
1	✓
2	✗
3	✓
⋮	⋮

Key (unique) →

→ Value (multiple)

capacity, AC, balcony etc. →

Ques → Given an integer array & multiple queries.

for every query → check if element x is present in array?

$A = [\overset{0}{2} \overset{1}{4} \overset{2}{11} \overset{3}{6} \overset{4}{8} \overset{5}{9} \overset{6}{1}]$

Query

$x = 10$

ans = false

2

ans = true

Bruteforce \rightarrow \forall query, travel the array & check.

$$TC = O(Q * N)$$

$$SC = O(1)$$

Sol 2 \rightarrow $\forall A[i]$, mark them visited in a separate array.

Direct
Access
Table (DAT)

$data[i] \rightarrow$ true \Rightarrow element 'i' is present in A
 \rightarrow false \Rightarrow element 'i' is not present in A

$A = [\overset{0}{2} \ \overset{1}{4} \ \overset{2}{11} \ \overset{3}{6} \ \overset{4}{8} \ \overset{5}{9} \ \overset{6}{1}]$
 \rightarrow $\checkmark \ \checkmark \ \checkmark \ \checkmark \ \checkmark \ \checkmark \ \checkmark$

$data = [\overset{0}{f} \ \overset{1}{T} \ \overset{2}{T} \ \overset{3}{T} \ \overset{4}{T} \ \overset{5}{f} \ \overset{6}{T} \ \overset{7}{f} \ \overset{8}{T} \ \overset{9}{T} \ \overset{10}{f} \ \overset{11}{T}]$

$\forall i, data[i] = false$

for $(i = 0 \text{ to } N-1) \{$

$data[A[i]] = true$

$\}$

$$TC = O(N)$$

Query $x \rightarrow ans = data[x]$

$\hookrightarrow TC = O(1)$ per query

$$\text{total } TC = O(N + Q)$$

$$SC = O(\max(A[i]))$$

Max array size allowed = $10^6 - 10^7$

$10^7 \rightarrow \text{int}$

$10^7 \times 4B$

$$= \frac{4 \times 10^3 \times 10^3 \times 10^1}{\text{MB}} = 40 \text{ MB}$$

if $0 \leq A[i] \leq 10^9 \rightarrow \text{MLE error}$

Ques \rightarrow If memory limit is M , store $A[i]$ that it is present in available memory limit.

$M=10$

0	1	2	3	4	5	6	7	8	9
		T					T	T	

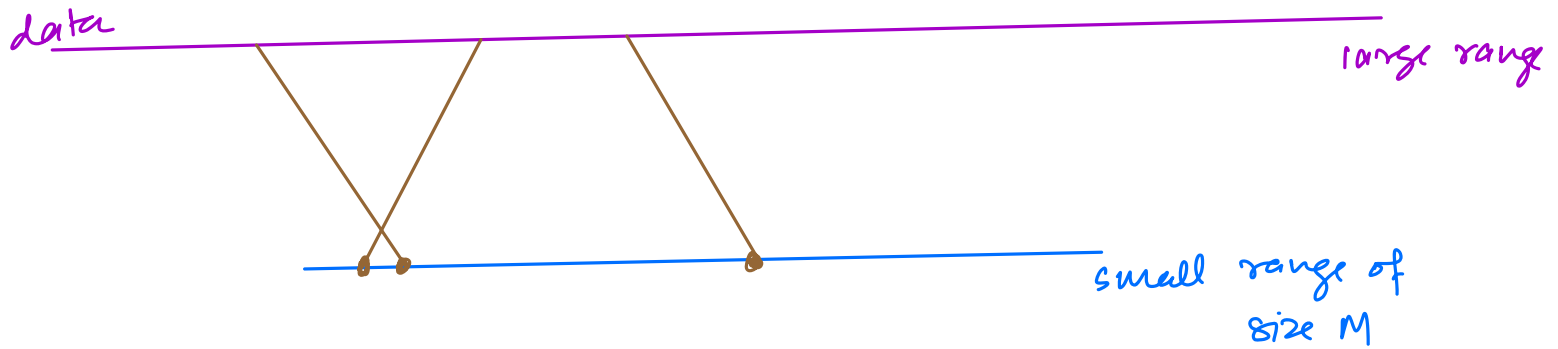
$$0 \leq h(A[i]) \leq 9$$

hash function $\rightarrow A[i] \div 10$
 \uparrow
 M

$A[i]$	$h(A[i])$
27	$27 \div 10 = 7$
18	$18 \div 10 = 8$
32	$32 \div 10 = 2$
37	$37 \div 10 = 7$

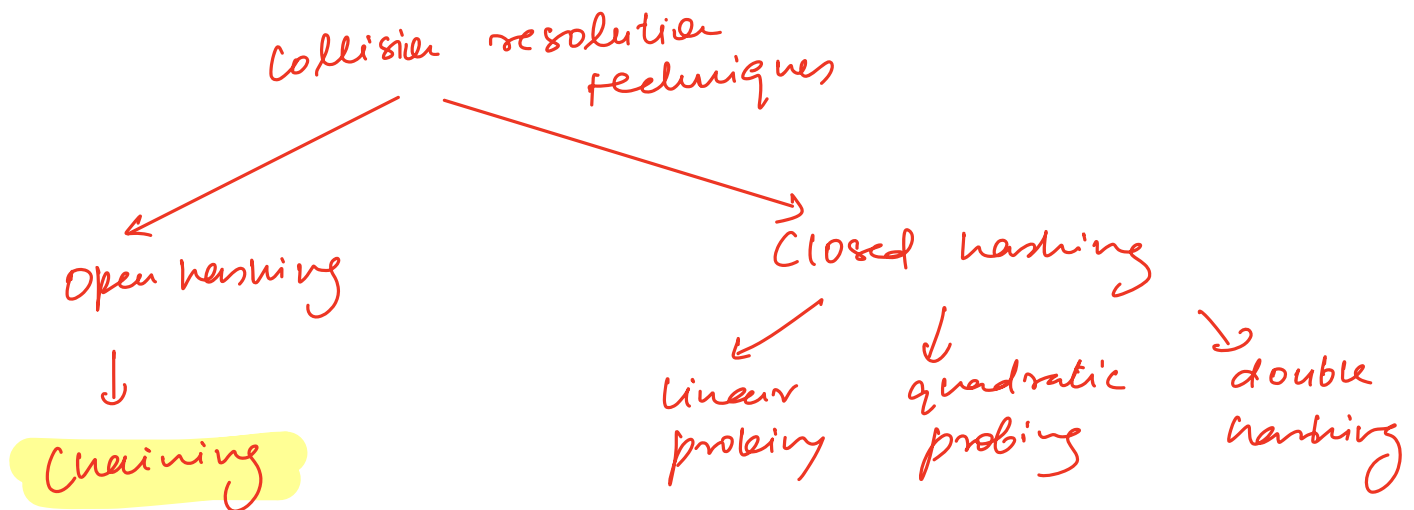
collision

Can we **avoid** collision? \rightarrow NO

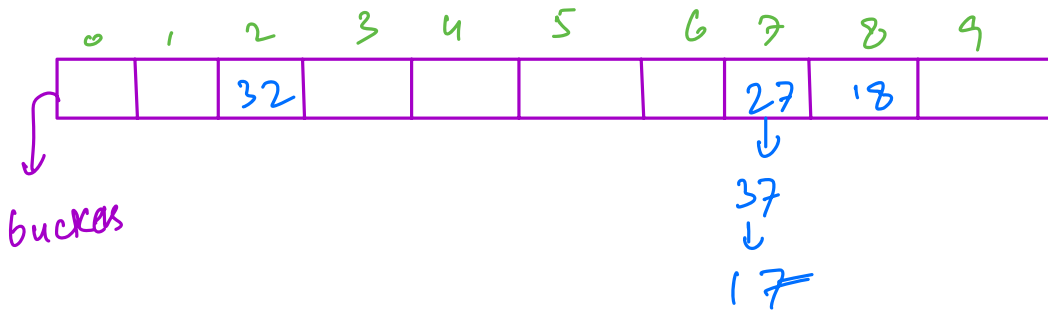


Pigeon Hole Principle \rightarrow If there are N pigeons & $(N-1)$ holes then there will be at least 1 hole with more than 1 pigeon.

Can we **handle** collision? \rightarrow YES



Chaining



TC per query $> O(1)$

(worst case TC = $O(N)$)

TC on average $<= \lambda$ (lambda)



(load factor)

elements inserted $\rightarrow 8$

Size of data array (M) $\rightarrow 10$

$$\lambda = \frac{8}{10} = 0.8$$

$$20/10 = 2$$

$$200/10 = 20$$

AU)

$u(AU)$

27

$$27 \times 10 = 7$$

18

$$18 \times 10 = 8$$

32

$$32 \times 10 = 2$$

37

$$37 \times 10 = 7$$

17

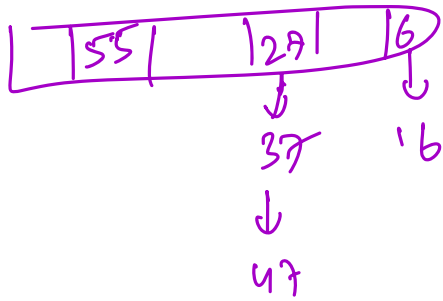
$$17 \times 10 = 7$$

There is a predefined threshold for $\lambda = 0.7$

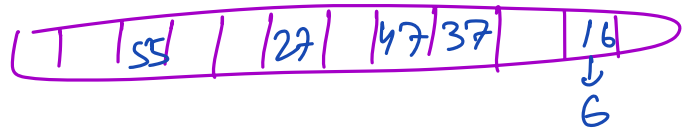
If λ becomes greater than threshold \Rightarrow

$\rightarrow TC = O(N)$

Redistribute the existing elements on a new data array with double size. $\lambda \rightarrow \text{half}$



⇒



Code Implementation

```
class HashMap <K, V> {
```

```
    private class HMNode {
```

```
        K key;
```

```
        V value;
```

```
        public HMNode (key, val) {
```

```
            this.key = key
```

```
            this.value = val
```

```
        }
```

```
    }
```

```
    private ArrayList <HMNode> () buckets;
```

```
    private int size; // number of key-value pairs
```

```
    public HashMap () {
```

```
        initBuckets();
```

```
        size = 0;
```

```
    }
```

```
private void initbuckets() {
```

```
    buckets = new ArrayList[4]; ← initial size of hashmap
```

```
    for (i=0 to 3)
```

```
        buckets[i] = new ArrayList<>();
```

```
    }
```

```
}
```

put(2,3)

⇒

2 → 3

put(2,4)

⇒

2 → 4

Insertion

```
void put ( K key, V value ) {
```

```
    int bi = hash(key); → index of bucket
```

```
    int di = getIndexWithinBucket(key, bi);
```

```
    if ( di != -1 ) { → key is present
```

```
        buckets[bi].get(di).value = value
```

```
    }
```

```
    else {
```

```
        HMNode temp = new HMNode(key, value)
```

```
        buckets[bi].add(temp);
```

```
        size++;
```

```
        lambda = size * 1.0 / buckets.length
```

↑ for decimal

```
        if ( lambda > 0.7 ) // rehashing
```

```
            rehash();
```

```
    }
```

```

3
3
private int hash( K key ) {
    int hc = key.hashCode()
    return hc % buckets.length
}

```

```

3
private int Get Index With Bucket( K key, int bi ) {
    int di = 0
    for ( HMNode node : buckets[bi] ) {
        if ( node.Key.equals(key) )
            return di;
        di++;
    }
    return -1 // key not found
}

```

→ TC = O(len of array-list)

```

private void rehash() {
    ArrayList<HMNode>[] oldBuckets = buckets;
    buckets = new ArrayList[oldBuckets.length * 2]
    // copy old bucket to new bucket
}

```

```
public V get(K key) {
```

```
    bi = hash(key)
```

```
    di = getIndexWithinBucket(key, bc);
```

```
    if (di == -1) {  
        return null
```

```
    }
```

```
    else {
```

```
        return buckets[bi].get(di).value
```

```
    }
```

```
}
```

```
public boolean containsKey()  $\Rightarrow$  TODO
```

```
public V remove(K key) {
```

```
    bi = hash(key)
```

```
    di = getIndexWithinBucket(key, bc);
```

```
    if (di == -1) {  
        return null
```

```
    }
```

```
    else {
```

```
        size --
```

```
return buckets(bi).remove(di).value;
```

```
}  
}
```

```
public int size() {
```

```
    return size
```

```
}
```

```
public ArrayList<K> keySet() {
```

```
    ArrayList<K> keys = new ArrayList<>();
```

```
    for (ArrayList<umNode> bucket : buckets) {
```

```
        for (umNode node : bucket) {
```

```
            keys.add(node.key)
```

```
        }
```

```
    }
```

```
    return keys
```

```
}
```