

Basic Electrical and Electronics Engineering

Module 4

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LECTURE 5

Module 4
Digital Systems
Lecture 4
Topics to be covered

- **HALF ADDER**
- **FULL ADDER**

Half Adder

Half Adder: is a combinational circuit that performs the addition of two bits, this circuit needs two binary inputs and two binary outputs.

Inputs		Outputs	
<i>X</i>	<i>Y</i>	<i>C</i>	<i>S</i>
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0
Truth table			

The simplified Boolean function from the truth table:

$$\left\{ \begin{array}{l} S = \bar{X}Y + X\bar{Y} \\ C = XY \end{array} \right. \quad \mathbf{1} \quad \text{(Using sum of product form)}$$

Where **S** is the sum and **C** is the carry.

$$\left\{ \begin{array}{l} S = X \oplus Y \\ C = XY \end{array} \right. \quad \mathbf{2} \quad \text{(Using XOR and AND Gates)}$$

Half Adder Truth Table:

Inputs		Outputs	
X	Y	S	C-out
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

$$S(X,Y) = \Sigma (1,2)$$

$$S = X'Y + XY'$$

$$S = X \oplus Y$$

$$C\text{-out}(x, y, C\text{-in}) = \Sigma (3)$$

$$C\text{-out} = XY$$

For Carry

		A	B
		0	0
		1	0
		0	1
		1	1

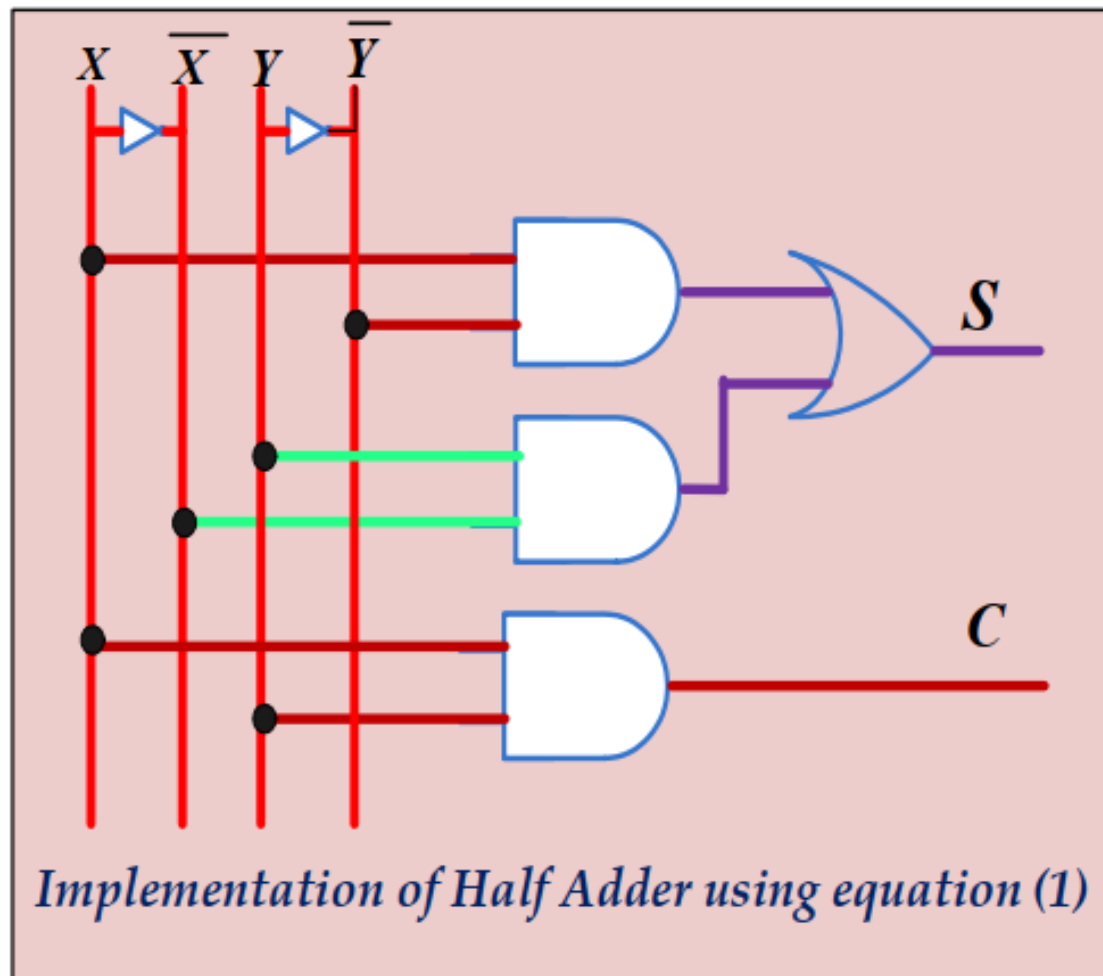
$$\text{Carry} = AB$$

For Sum

		A	B
		0	0
		1	0
		0	1
		1	1

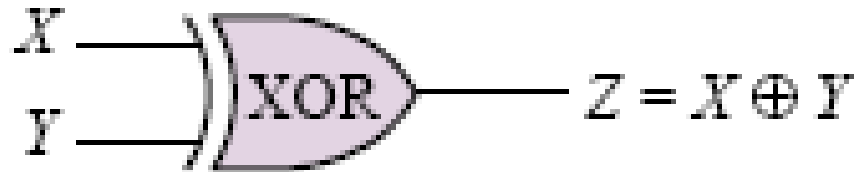
$$\begin{aligned} \text{Sum} &= \overline{A}\overline{B} + \overline{A}B \\ &= A \oplus B \end{aligned}$$

$$\left. \begin{aligned} S &= \bar{X}Y + X\bar{Y} \\ C &= XY \end{aligned} \right\} \text{ (Using sum of product form)}$$



XOR (Exclusive OR) Gate

$$Z = X'Y + XY'$$

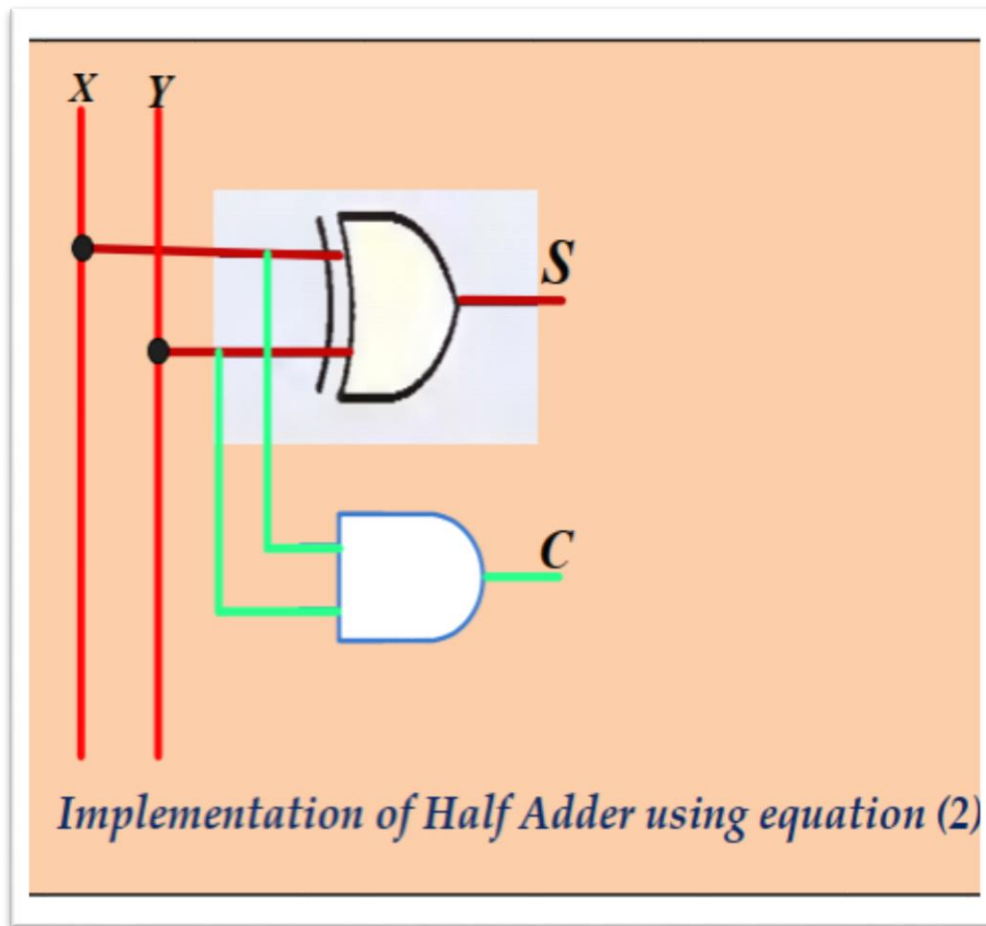


$$Z = X \oplus Y$$

X	Y	Z
0	0	0
0	1	1
1	0	1
1	1	0

Truth table

$$\left. \begin{array}{l} S = X \oplus Y \\ C = XY \end{array} \right\} \text{ (Using XOR and AND Gates)}$$



Full Adder

Full Adder is a combinational circuit that performs the addition of three bits (two significant bits and previous carry).

- It consists of *three inputs and two outputs*, two inputs are the bits to be added, the third input represents the carry from the previous position.
- The full adder is usually a component in a cascade of adders, which add 8, 16, etc, binary numbers.

Full Adder Truth Table

Inputs			Outputs	
<i>X</i>	<i>Y</i>	<i>C_{in}</i>	<i>S</i>	<i>C_{out}</i>
<i>0</i>	<i>0</i>	<i>0</i>	<i>0</i>	<i>0</i>
<i>0</i>	<i>0</i>	<i>1</i>	<i>1</i>	<i>0</i>
<i>0</i>	<i>1</i>	<i>0</i>	<i>1</i>	<i>0</i>
<i>0</i>	<i>1</i>	<i>1</i>	<i>0</i>	<i>1</i>
<i>1</i>	<i>0</i>	<i>0</i>	<i>1</i>	<i>0</i>
<i>1</i>	<i>0</i>	<i>1</i>	<i>0</i>	<i>1</i>
<i>1</i>	<i>1</i>	<i>0</i>	<i>0</i>	<i>1</i>
<i>1</i>	<i>1</i>	<i>1</i>	<i>1</i>	<i>1</i>
<i>Truth table for the full adder</i>				

➤ The ***S*** output is equal to ***1*** when only one input is equal to ***1*** or when all three inputs are equal to ***1***.

➤ The ***C_{out}*** output has a carry ***1*** if two or three inputs are equal to ***1***.

➤ The Karnaugh maps and the simplified expression are shown in the following figures:

Full Adder Truth Table

Inputs			Outputs	
X	Y	C-in	S	C-out
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

$$S(X, Y, C\text{-in}) = \Sigma (1, 2, 4, 7)$$

$$C\text{-out}(x, y, C\text{-in}) = \Sigma (3, 5, 6, 7)$$

Sum S

C-in \ XY		X			
		00	01	11	10
0	0	0	2 1	6	4 1
1	1	1 1	3	7 1	5

Y

$$S = X'Y'(C\text{-in}) + X'Y(C\text{-in})' + XY'(C\text{-in})' + XY(C\text{-in})$$

$$S = X \oplus Y \oplus (C\text{-in})$$

Carry C-out

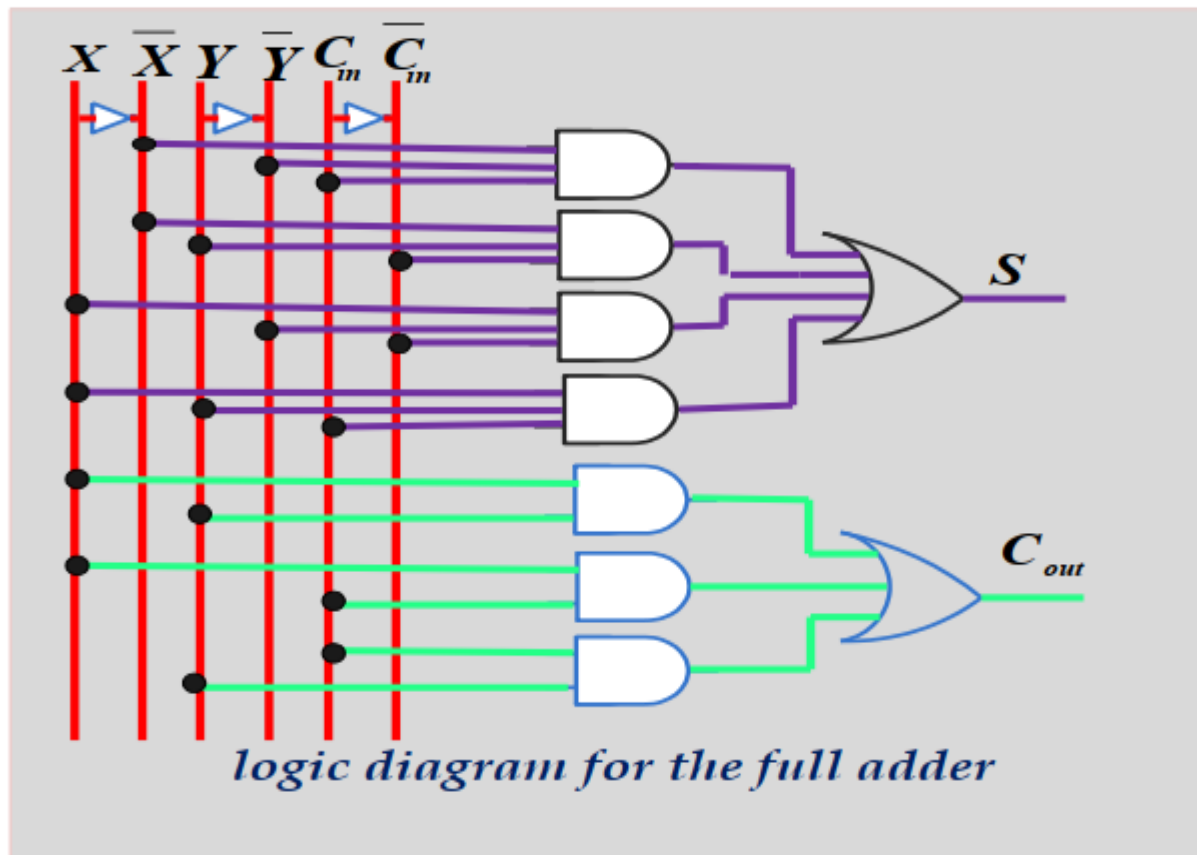
C-in \ XY		X			
		00	01	11	10
0	0	0	2	6 1	4
1	1	1	3 1	7 1	5 1

Y

$$C\text{-out} = XY + X(C\text{-in}) + Y(C\text{-in})$$

$$\left\{ \begin{array}{l} S = \bar{X} \bar{Y} C_{in} + \bar{X} Y \bar{C}_{in} + X \bar{Y} \bar{C}_{in} + X Y C_{in} \\ C_{out} = X Y + X C_{in} + Y C_{in} \end{array} \right\} \text{ (Sum of products)}$$

- The *logic diagrams* for the full adder implemented in *sum-of-products* form are the following:



- It can also be implemented using *two half adders* and *one OR gate* (using **XOR** gates).

$$\begin{cases} S = C_{in} \oplus (X \oplus Y) \\ C_{out} = C_{in} \cdot (X \oplus Y) + XY \end{cases}$$

Proof:

The sum:

$$\begin{aligned} S &= \bar{X}\bar{Y}C_{in} + \bar{X}Y\bar{C}_{in} + X\bar{Y}\bar{C}_{in} + XYC_{in} \\ &= \bar{C}_{in}(\bar{X}Y + X\bar{Y}) + C_{in}(\bar{X}\bar{Y} + XY) \\ &= \bar{C}_{in}(\bar{X}Y + X\bar{Y}) + C_{in}(\overline{\bar{X}Y + X\bar{Y}}) \\ S &= C_{in} \oplus (X \oplus Y) \end{aligned}$$

The carry output:

$$\begin{aligned} C_{out} &= \bar{X}YC_{in} + X\bar{Y}C_{in} + XYC_{in} + XY\bar{C}_{in} \\ &= C_{in}(\bar{X}Y + X\bar{Y}) + XY(C_{in} + \bar{C}_{in}) \\ C_{out} &= C_{in} \cdot (X \oplus Y) + XY \end{aligned}$$

