Q1

```
for (i=0 to n)
{ if (arr[i]== key)
    break;
}
```

Q2)

```
void insertion sort (int arr[], int n)
{ for (int i= 1 ; i<n ; i++)
  {
     j= i-1
     x = arr[i]
     while (j>-1 && arr[j]>x)
      { arr[j+1]= arr[j];
        j--;
      }
     arr[j+1] =x;
  }
}

void insertion (int arr[], int n)
{
   if (n<=1) return
   insertion sort (arr, n-1);
   int last = arr[n-1];
   int j =n-2
   while (j>=0 && arr[j]> last)
    { arr[j+1] =arr[j];
      j--;
    }
   arr[j+1]= last;
}
```

Insertion sort doesn't need to know about what value it will sort during running & hence called online sort.

Other sorting Algos

1) Bubble Sort
2) Quick Sort
3) Merge Sort
4) Selection Sort
5) Heap Sort

Time complexity.

Name	Best	Worst	Avg.
Selection	$O(n^2)$	$O(n^2)$	$O(n^2)$
Bubble	$O(n)$	$O(n^2)$	$O(n^2)$
Insertion	$O(n)$	$O(n^2)$	$O(n^2)$
Heap	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Quick	$O(n \log n)$	$O(n^2)$	$O(n \log n)$
Merge	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$

Inplace	stable	Online	← Sorting Type
Bubble	Merge	Insertion	
Selection	Bubble		
Insertion	Insertion		
Quick	~~Heap~~		
Heap	Count Sort		

Iterative

```
int bsearch (int a[], int l, int r, int key)
{
  while (l<=r)
  int m = (l+r)/2;
  if (a[m] == key)
    return m
    else if (key < a[m])
     r = m-1;
  else l = m+1
  }
 return -1;
}
```

Recursive

```
int bs (int ar, int l, int r, int
{ while (l<=r)
  { m = (l+r)/2
    if (k == a[m]
    return m;
    else if (kg < a[m])
     return bs (ar, l, mid-1, k
    else
    return bs (a, mid+1, r,
  }
return -1
```

Linear search - O(n)
Binary = O(log n)

Qu1

$T(n) = T(n/2) + 1$ - ①
$T(n/2) = T(n/4) + 1$ - ②
$T(n/4) = T(n/8) + 1$ - ③

$T(n) = T(n/2) + 1$
$T(n/4) + 1 + 1$
$T(n/8) + 1 + 1 + 1$
;
$T(n/2^n) + 1$ (k times)

Let $2^k = n$
$k = \log n$
$T(n) = T(n/n) + \log n$
$T(n) = T(1) + \log(n)$
$T(n) = O(\log n)$ →

Quick sort is the fastest general sort.
It is stable & has the avg and best running time of $O(n \log n)$

Ques 31

(10) Quick sort gives the worst time complexity in

1) The array is sorted and either the first or the last element is selected as a pivot.

2) ~~Worst~~ Best case when the pivot is a mean element.

(11)

Merge sort →

Best case → $T(n) = 2T(n/2) + O(n)$

Worst case → $T(n) = 2T(n/2) + O(n)$

Quick sort

Best case → $T(n) = 2T(n/2) + O(n) \rightarrow O(n \log n)$

Worst case → $T(n) = T(n-1) + O(n) \rightarrow O(n^2)$

To prevent bubble sort from scanning the whole array if it is sorted already then we can use a counter to check if any exchanges were made. If not then we break the loop and conclude that the array is sorted

```
void bubb(int ar, int n)
{   int cnt=0;
    for(int i=0; i<n; ++j)
    {
        for(int j=0; j<n; ++j)
        {
            if(ar[j] > ar[j+1])
            { swap(ar[j], ar[j+1]) cnt++;}
        }
        if(!cnt) break;
}
```