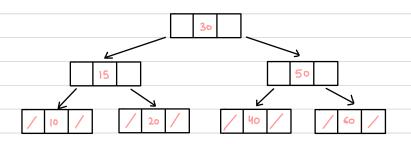


# 30 15 50 10 60



- · Left subtree is smaller
- · Right subtree is greater
- · Useful for searching (in less number of comparisons)
- · No duplicates
- · Inorder gives sorted order
- · For n number of nodes, catalan number of trees can be generated

$$T(n) = \frac{2n(n)}{n+1}$$

- \* BST is represented using linked representation.
- · Can also be represented using arrays.

O(P)

logn = h = n

SEARCHING IN A BINARY SEARCH TREE ( searching takes maximum time same as height of type) O(logn)

BINARY SEARCH TREE

```
(1) RECURSIVE SEARCH (tail recursion)
```

```
Node * Rsearch (Node *t, int key)

if (t == NULL)  // If element is not

return NULL; found

if (key == t -> data)

return t;

else if (key < t -> data)

return Rsearch (t -> Rchild, key);

else

return Rsearch (t -> rchild, key);
```

(2) ITERATIVE VERSION 0 (log n)

else if (key < 
$$t \rightarrow data$$
)
$$t = t \rightarrow lchild;$$
else
$$t = t \rightarrow rchild;$$
return NULL;

#### INSERTING IN BINARY SEARCH TREE

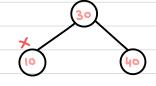
```
void Insert (Node *t, int key)
         Node " = NULL, " p;
         While (t! = NULL)
                if ( key == t → data)
                        return; // To avoid duplication
                else if (key < t → data)
                      t=t\rightarrow lchild;
                        t= t -> rchild;
         3
         p=new Node;
         p → data = key;
         \rho \rightarrow l child = \rho \rightarrow \pi child = NULL;
        if (p → data < r → data)
                > -> lchild = p;
        else
                r → rchild = p;
3
```

#### RECURSIVE INSERT IN BST

```
Node " Insert (Node "p, int key)
                                                           void main()
                                                                  Node "TOOL = NULL;
      if (p = = NULL)
                                                                  root = insert (root, 30);
             t = new Node;
                                                                  insert troot, 20);
             t → data = key;
                                                                  insert (root, 25);
             t → lchild = t → rchild = NULL;
                                                           3
      3
      if ( key 
             p \rightarrow lchild = insert (p \rightarrow lchild, key);
      else if ( key > p -> data)
              p → rchild = insert (p → rchild, key))
      zetuzn p;
3
```

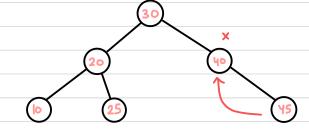
#### DELETING FROM BST

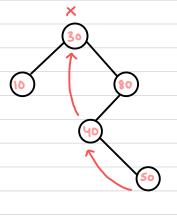
#### CASE 1:



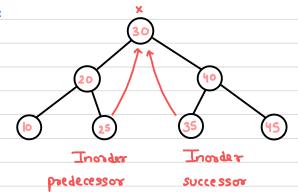
## CASE 4 :

# CASE 2:





#### CASE 3:



// Multiple changes to be

```
int Height (struct Node *p)

{

int x, y;

if (p == NULL)

return 0;

x = Height (p → 1child);

y = Height (p → rchild);

return x>y? x+1: y+1;

}
```

```
Struct Node * Infre(struct Node *p)

{

while p \neq p \Rightarrow \pi \text{ child} := null)

p = p \Rightarrow \pi \text{ child};

return p;
}
```

```
Struct Node *In Succ (Struct Node *p)

{

while l p dd p \rightarrow lchild! = NULL)

p = p \rightarrow lchild;

return p;
}
```

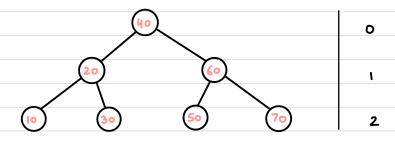
```
Struct Node " Delete ( Struct Node "p, int key)
E
        Struct Node *9;
        if (p = = NULL)
                return NULL;
         if (p → lchild = NULL of p → rchild = NULL) // Leaf Node
                  if lp = = root
                           root = NULL;
                  free (p);
                  return NULL;
         3
         if ( key < p → data)
                  p → lchild = Delete (p → lchild, key);
         else if (key > p → data)
                  p → rchild = Delete Lp → rchild, key);
         else
         ٤
                  if (Height (p → Ichild) > Height (p → rchild))
                           q = InPre(ρ→ lchild);
                           p \rightarrow data = q \rightarrow data;
                           \rho \rightarrow l child = Delete (\rho \rightarrow l child, q \rightarrow data);
                  3
                  else
                           q = \text{InSucc}(\rho \rightarrow \text{schild});
                            p \rightarrow data = q \rightarrow data;
                            \rho \rightarrow schild = Delete(\rho \rightarrow schild, q \rightarrow data);
                   }
          return p;
3
```

#### GENERATING BINARY SEARCH TREE

```
Void Createpre (int pre[], int n) O(n)
                                          BST can be generated using
        Stack St;
                                               · Preorder + Inorder
        Node "t;
                                               · Postorder + Inorder
        int i=0;
                                          Souted preorder of BST gives its Inorder
        root = new Node;
                                                 postorder
        root -> data = pre[i++];
        root -> lchild = root -> rchild = NULL;
        p = voot;
        while (i<n)
        ٤
               if [pre[i] < p → data]
               ٢
                       t = new Node;
                       t -> data = pre[i++];
                       t → Schild = t → rchild = NULL;
                        p \rightarrow lchild = t;
                       push (Jstk, p);
                       p = t;
              3
              else
                       if (pre[i] > p → data dd pre[i] < Stacktop(Stk) → data)
                               t = new Node;
                               t → data = pre [i++];
                               t → 1child = t → rchild = NULL;
                                p → rchild = t;
                               p = t;
                       3
                       else
                              p = pop(dstk);
            1
3
```

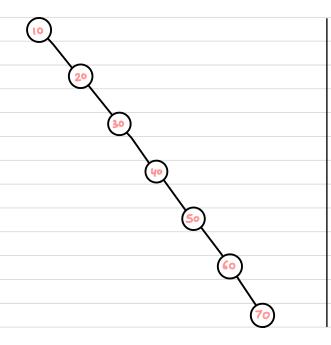
## DRAWBACK OF BINARY SEARCH TREE

KEYS: 40, 20, 30, 60, 50, 10, 70



 $h = \log_2(n+1) - 1$   $O(\log n)$ 

KEYS: 10,20,30,40,50,60,70



h = n - 1 O(n)

0

2

3

4

5

6

There is no control over height of binary search tree.

It only depends on user's input or sequence of insertion