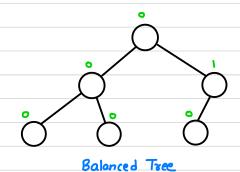
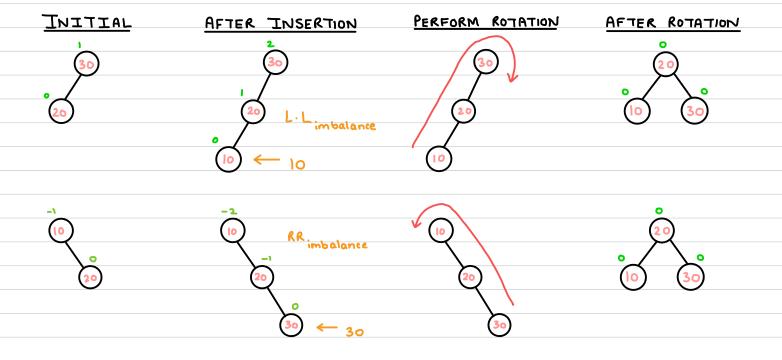


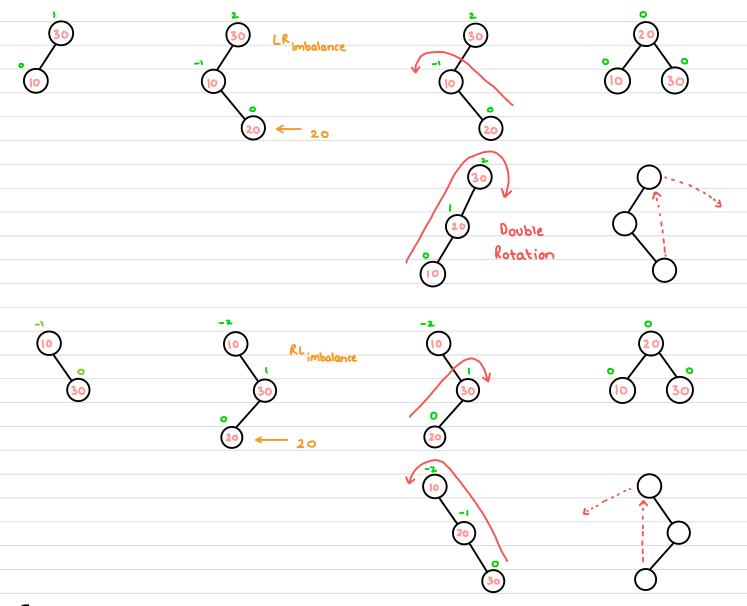
No of edges No of edges Height Balanced Binary Search Trees Balance factor = Height of left subtree - Height of right subtree \[\{-1,0,1\} \rightarrow \text{Balanced Tree} \]



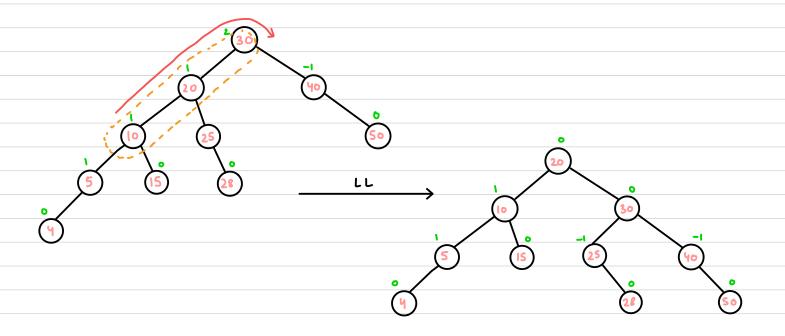
Rotations for insertion in AVL trees

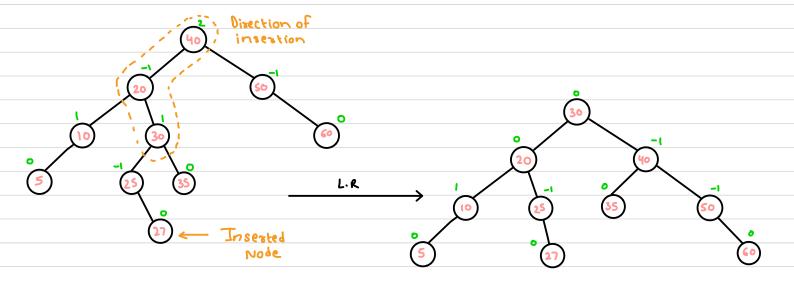
- (1) LL Rotation 7 Single Rotation
- (2) RR Rotation
- (3) LR Rotation 7 Double Rotation
- (4) RL Rotation





FORMULA OF ROTATION FOR INSERTION





PROGRAM FOR LL ROTATION

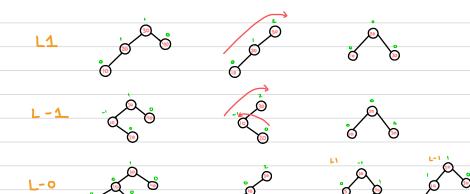
```
Struct Node
      Struct Node "Ichild;
      int data;
      int height; // We will set height for
      Struct Node * rihild; each and every Node
} * root = NULL;
Struct Node " Insert (Node "p, int key)
       if (p = = NULL)
             t = new Node;
              t → data = key;
              t \rightarrow height = 1;
              t \rightarrow l child = t \rightarrow r child = NULL;
              return t;
      if ( key  data)
              p \rightarrow l child = insert (p \rightarrow l child, key);
      else if ( key > p -> data)
              p → rchild = insert (p → rchild, key);
      p → height = Node Height (p); // update height of each
                                  node at returning time
```

```
if ( balance Factor (p) == 2 dd Balance Factor (p → lchild) = = 1)
               seturn LL Rotation (p);
      else if (Balance Factor (p) == 2 dd Balance Factor (p → schild) = = -1)
               return LRRotation (P);
      else if L Balance Factor (p) == -2 dd Balance Factor (x → child) = = -1)
               return RRRotation (p);
      else if ( Balance Factor (p) = = -2 dd Balance Factor (r → child) = = 1)
               return RL Rotation (p);
      zetuen p;
int NodeHeight (Struct Node *p)
     int hl, hr; // keight of left subtree (HL), height of night subtree (HR)
   - hl = \rho dd \rho \rightarrow lchild ? \rho \rightarrow lchild \rightarrow height : 0;
   - hr = ρ dd ρ → rohild? ρ → rohild → height: 0;
    → NOT NULL
     return he + thro? he + 1;
int Balance Factor ( struct Node *p)
     int hl, hr;
     hl = p dd p → lchild? p → lchild → height : 0;
     hr = p dd p → rohild? p → rohild → height: 0;
     return hl-hr;
3
Struct Node "LL Rotation (Struct Node "P)
     Struct Node "pl = p → Ichild;
     Shuct Node *plr=pl → lchild;
     pl -> rchild = p;
     p -> lchild = par;
     p → height = Node Keight (p);
     pl → height = Node Height (pl);
```

```
if (soot == p) // I sotation was performed on
          root = pl; root node, root needs to be
                       updated.
    return pl;
Struct Node " LRRotation ( struct Node "p)
     Struct Node *pl = p → lchild;
     struct Node * plr = pl -> rchild;
     pl → rehild = plr - lehild;
     p → I child = plr -> r child;
     plr -> luhild = pl;
     plo -> whild = p;
     pl → height = Node Height (pl);
     \rho \rightarrow height = Node Height (p);
     plr > height = Node Height (plr);
    if (300t == p)
       i ald = tool
    return plr; // New root;
```

DELETION FROM AVL TREES WITH ROTATION

1. L 1 Rotation 4. R 1 Rotation
2. L -1 Rotation 5. R -1 Rotation
3. L O Rotation 6. R - 0 Rotation

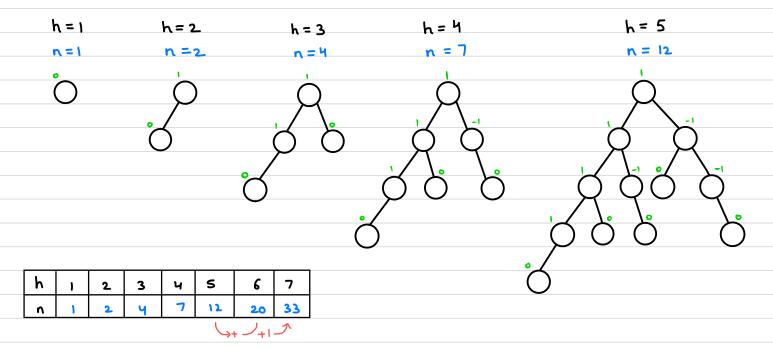


// Other three will be mirror images

HEIGHT VS NODES OF AVL TREES

If height is given:

- Max nodes $n = 2^h 1$ // Not $2^{h+1} 1$ because height is starting from 1.
- · Min nodes n = Look in table



$$N(h) = \begin{cases} 0 & 0 \\ 1 & 1 \end{cases}$$

$$N(h-2) + N(h-1) + 1 \qquad \text{Otherwise}$$

$$// \text{formula Same as Fibonacci series}$$

$$\Rightarrow \text{balanced series}$$

I 'N' Nodes are given find:

- · Min height = log_ (n+1)
- · Max height = Look in table

Foreg for 13 nodes h = 5 // Look from node towards height