

Intro to GPU Programming with the OpenMP API

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OpenMP Architecture Review Board

The mission of the OpenMP ARB (Architecture Review Board) is to standardize directive-based multi-language high-level parallelism that is performant, productive and portable.

The OpenMP API moves common approaches into an industry standard to simplify a developers' life.



































































Agenda

- OpenMP device and execution model
- Offload basics
- Exploit parallelism
- Asynchronous offloading
- Summary



Introduction to OpenMP Offload Features



Running Example for this Presentation: saxpy

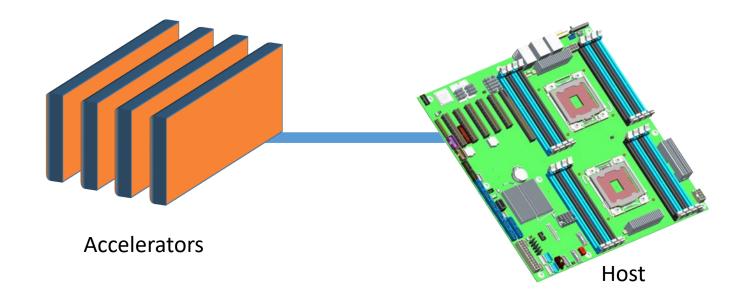
```
void saxpy() {
    float a, x[SZ], y[SZ];
    // left out initialization
    double t = 0.0;
                                                        Timing code (not needed, just to have
    double tb, te;
                                                        a bit more code to show (2)
    tb = omp_get_wtime();
#pragma omp parallel for firstprivate(a)
    for (int i = 0; i < SZ; i++) {
                                                        This is the code we want to execute on a
        y[i] = a * x[i] + y[i];
                                                        target device (i.e., GPU)
    te = omp_get_wtime();
                                                        Timing code (not needed, just to have
    t = te - tb;
                                                        a bit more code to show ③)
    printf("Time of kernel: %lf\n", t);
```

Don't do this at home!
Use a BLAS library for this!



Device Model

- As of version 4.0, the OpenMP API supports accelerators/coprocessors
- Device model:
 - One host for "traditional" multi-threading
 - Multiple accelerators/coprocessors of the same kind for offloading





OpenMP Execution Model for Devices

- Offload region and its data environment are bound to the lexical scope of the construct
 - Data environment is created at the opening curly brace
 - Data environment is automatically destroyed at the closing curly brace
 - Data transfers (if needed) are done at the curly braces, too:
 - Upload data from the host to the target device at the opening curly brace.
 - Download data from the target device at the closing curly brace.


```
!$omp target
!$omp map(alloc:A) &
!$omp map(to:A) &
!$omp map(from:A) &
    call compute(A)
!$omp end target
```

Device mem.



OpenMP for Devices - Constructs

- Transfer control and data from the host to the device
- Syntax (Fortran)

```
!$omp target [clause[[,] clause],...]
structured-block
!$omp end target
```

Clauses

```
device(scalar-integer-expression)
map([{alloc | to | from | tofrom}:] list)
if(scalar-expr)
```



```
void saxpy() {
    float a, x[SZ], y[SZ];
    double t = 0.0;
    double tb, te;
   tb = omp get_wtime();
#pragma omp target "map(tofrom:y[0:SZ])"
    for (int i = 0; i < SZ; i++) {
       y[i] = a * x[i] + y[i];
   te = omp_get_wtime();
   t = te - tb;
    printf("Time of kernel: %lf\n", t);
```

clang/LLVM: clang -fopenmp -fopenmp-targets=<target triple>

GNU: gcc -fopenmp

AMD ROCm: clang -fopenmp -offload-arch=gfx908

NVIDIA: nvcc -mp=gpu -gpu=cc80

Intel: icx -fiopenmp -fopenmp-targets=spir64
IBM XL: xlc -qsmp -qoffload -qtgtarch=sm 70

The compiler identifies variables that are used in the target region.

x[0:SZ]

y[0:SZ]

x[0:SZ]

y[0:SZ]

All accessed arrays are copied from host to device and back

Presence check: only transfer if not yet allocated on the

device.

Copying x back is not necessary: it was not changed.



IBM XL:

```
subroutine saxpy(a, x, y, n)
    use iso_fortran_env
    integer :: n, i
    real(kind=real32) :: a
    real(kind=real32), dimension(n) :: x
    real(kind=real32), dimension(n) :: y
!$omp target "map(tofrom:y(1:n))"
    do i=1,n
       y(i) = a * x(i) + y(i)
    end do
!$omp end target
end subroutine
```

```
clang/LLVM:
            flang -fopenmp -fopenmp-targets=<target triple>
GNU:
             gfortran -fopenmp
            flang -fopenmp -offload-arch=gfx908
AMD ROCm:
NVIDIA:
             nvfortran -mp=gpu -gpu=cc80
            ifx -fiopenmp -fopenmp-targets=spir64
Intel:
            xlf -qsmp -qoffload -qtgtarch=sm 70
```

The compiler identifies variables that are used in the target region.

> All accessed arrays are copied from host to device and back

> > Presence check: only transfer if not yet allocated on the device.

x(1:n)y(1:n)

x(1:n)

y(1:n)

Copying x back is not necessary: it was not changed.



```
void saxpy() {
    double a, x[SZ], y[SZ];
    double t = 0.0;
    double tb, te;
                                                      x[0:SZ]
   tb = omp_get_wtime();
                                                      y[0:SZ]
#pragma omp target map(to:x[0:SZ]) \
                   map(tofrom:y[0:SZ])
    for (int i = 0; i < SZ; i++) {
        y[i] = a * x[i] + y[i];
                                                     y[0:SZ]
   te = omp_get_wtime();
   t = te - tb;
    printf("Time of kernel: %lf\n", t);
```



```
void saxpy(float a, float* x, float* y,
           int sz) {
    double t = 0.0;
    double tb, te;
   tb = omp get wtime();
#pragma omp target map(to:x[0:sz]) \
                   map(tofrom:y[0:sz])
    for (int i = 0; i < sz; i++) {
       y[i] = a * x[i] + y[i];
   te = omp_get_wtime();
   t = te - tb;
    printf("Time of kernel: %lf\n", t);
```

The compiler cannot determine the size of memory behind the pointer.

Observation when running this: the loop is a sequential loop, and the capabilities of the GPU are not really used! \odot

y[0:sz]

Programmers have to help the compiler with the size of the data transfer needed.



Commercial Break... Community Interaction











Check out openmp.org/news/events-calendar/



Exploiting (Multilevel) Parallelism



Creating Parallelism on the Target Device

- ■The target construct transfers the control flow to the target device
 - Transfer of control is sequential and synchronous
 - This is intentional!

- OpenMP separates offload and parallelism
 - Programmers need to explicitly create parallel regions on the target device
 - In theory, this can be combined with any OpenMP construct
 - In practice, there is only a useful subset of OpenMP features for a target device such as a GPU, e.g., no I/O, limited use of base language features.



GPUs are multi-level devices: SIMD, threads, thread blocks

Create a team of threads to execute the loop in parallel using SIMD instructions.



Multi-level Parallel saxpy

Manual code transformation

- Tile the loop into an outer loop and an inner loop.
- Assign the outer loop to "teams".
- Assign the inner loop to the "threads".
- (Assign the inner loop to SIMD units.)



Multi-level Parallel saxpy

■ For convenience, OpenMP defines composite constructs to implement the required code transformations

```
subroutine saxpy(a, x, y, n)
   ! Declarations omitted
!$omp omp target teams distribute parallel do simd &
!$omp& num_teams(num_blocks) map(to:x) map(tofrom:y)
   do i=1,n
        y(i) = a * x(i) + y(i)
   end do
!$omp end target teams distribute parallel do simd
end subroutine
```



teams Construct

- Support multi-level parallel devices
- Syntax (C/C++):

```
#pragma omp teams [clause[[,] clause],...]
structured-block
```

■Syntax (Fortran):

```
!$omp teams [clause[[,] clause],...]
structured-block
```

■ Clauses

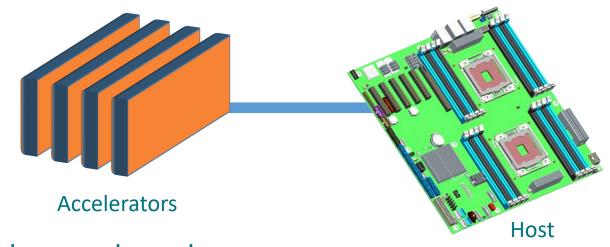
```
num_teams(integer-expression), thread_limit(integer-expression)
default(shared | firstprivate | private none)
private(list), firstprivate(list), shared(list), reduction(operator:list)
```



Optimizing Data Transfers



Optimizing Data Transfers is Key to Performance



 Connections between host and accelerator are typically lower-bandwidth, higher-latency interconnects

Bandwidth host memory: hundreds of GB/sec

Bandwidth accelerator memory: TB/sec

■ PCle Gen 4 bandwidth (16x): tens of GB/sec

- Unnecessary data transfers must be avoided, by
 - only transferring what is actually needed for the computation, and
 - making the lifetime of the data on the target device as long as possible.



Role of the Presence Check

■If map clauses are not added to target constructs, presence checks determine if data is already available in the device data environment:

```
subroutine saxpy(a, x, y, n)
    use iso_fortran_env
    integer :: n, i
    real(kind=real32) :: a
    real(kind=real32), dimension(n) :: x
    real(kind=real32), dimension(n) :: y
!$omp target "present?(y)" "present?(x)"
    do i=1,n
       y(i) = a * x(i) + y(i)
    end do
!$omp end target
end subroutine
```

- OpenMP maintains a mapping table that records what memory pointers have been mapped.
- That table also maintains the translation between host memory and device memory.
- Constructs with no map clause for a data item then determine if data has been mapped and if not, a map(tofrom:...) is added for that data item.



Optimize Data Transfers

- Reduce the amount of time spent transferring data:
 - Use map clauses to enforce direction of data transfer.
 - Use target data, target enter data, target exit data constructs to keep data environment on the target device.

```
subroutine saxpy(a, x, y, n)
  ! Declarations omitted

!$omp target "present?(y)" "present?(x)"
  do i=1,n
      y(i) = a * x(i) + y(i)
  end do
!$omp end target
end subroutine
```



Optimize Data Transfers

- Reduce the amount of time spent transferring data:
 - Use map clauses to enforce direction of data transfer.
 - Use target data, target enter data, target exit data constructs to keep data environment on the target device.

```
void zeros(float* a, int n) {
#pragma omp target teams distribute parallel for
    for (int i = 0; i < n; i++)
        a[i] = 0.0f;
}</pre>
```

```
void saxpy(float a, float* y, float* x, int n) {
#pragma omp target teams distribute parallel for
   for (int i = 0; i < n; i++)
      y[i] = a * x[i] + y[i];
}</pre>
```



Example: target data and target update

```
#pragma omp target data device(0) map(alloc:tmp[:N]) map(to:input[:N)) map(from:res)
#pragma omp target device(0)
#pragma omp parallel for
    for (i=0; i<N; i++)
      tmp[i] = some computation(input[i], i);
    update_input_array_on_the_host(input);
#pragma omp target update device(0) to(input[:N])
#pragma omp target device(0)
#pragma omp parallel for reduction(+:res)
    for (i=0; i<N; i++)
      res += final_computation(input[i], tmp[i], i)
```



target data Construct Syntax

- Create scoped data environment and transfer data from the host to the device and back
- Syntax (Fortran)

```
!$omp target data [clause[[,] clause],...]
structured-block
!$omp end target data
```

Clauses

```
device(scalar-integer-expression)
map([{alloc | to | from | tofrom | release | delete}:] list)
if(scalar-expr)
```



target update Construct Syntax

- Issue data transfers to or from existing data device environment
- Syntax (C/C++)

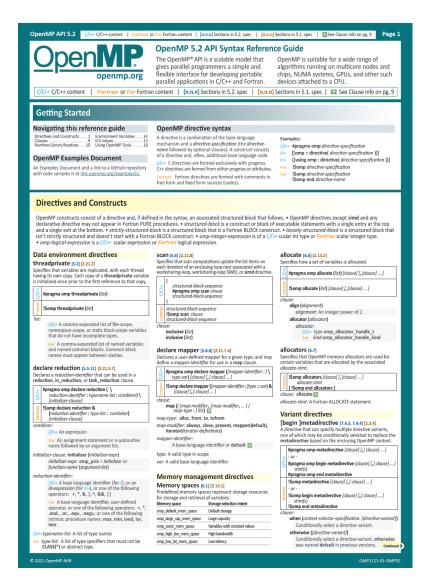
 #pragma omp target update [clause[[,] clause],...]
- Syntax (Fortran)
 !\$omp target update [clause[[,] clause],...]

Clauses

```
device(scalar-integer-expression)
to(list)
from(list)
if(scalar-expr)
```



Commerical Break...







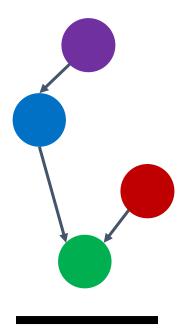
Asynchronous Offloading



Asynchronous Offloads

- OpenMP target constructs are synchronous by default
 - The encountering host thread awaits the end of the target region before continuing
 - The nowait clause makes the target constructs asynchronous (in OpenMP speak: they become an OpenMP task)

```
depend(out:a)
#pragma omp task
    init data(a);
                                                               depend(in:a) depend(out:x)
#pragma omp target map(to:a[:N]) map(from:x[:N])
                                                  nowait
    compute 1(a, x, N);
#pragma omp target map(to:b[:N]) map(from:z[:N])
                                                               depend(out:y)
                                                  nowait
    compute 3(b, z, N);
#pragma omp target map(to:y[:N]) map(to:z[:N])
                                                               depend(in:x) depend(in:y)
                                                   nowait
    compute 4(z, x, y, N);
#pragma omp taskwait
```





Summary

- OpenMP API is ready to use GPUs for offloading computations
 - Mature offload model w/ support for asynchronous offload/transfer
 - Tightly integrates with OpenMP multi-threading on the host
- More, advanced features (not covered here)
 - Memory management API
 - Interoperability with native data management
 - Interoperability with native streaming interfaces
 - Unified shared memory support

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Visit www.openmp.org for more information