Linking II: Static and Dynamic Linking

COMP400727: Introduction to Computer Systems

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Today

- Libraries and Static Linking
- Dynamic Linking
- Case Study: Library Interpositioning

Libraries: Packaging a Set of Functions

- How to package functions commonly used by programmers?
 - Math, I/O, memory management, string manipulation, etc.

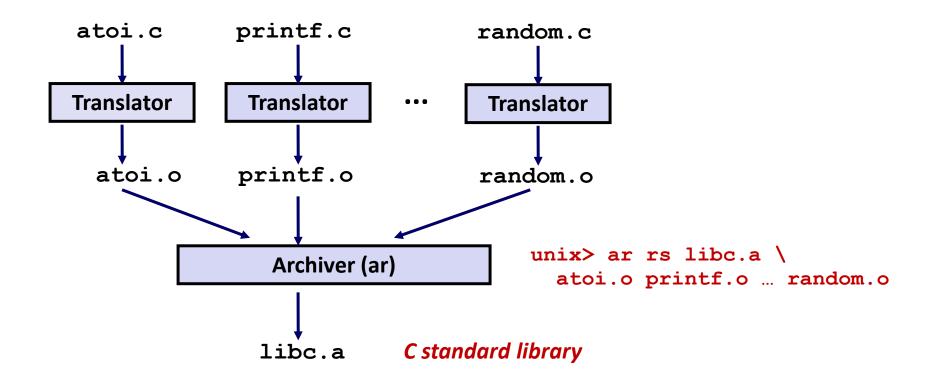
Awkward, given the linker framework so far:

- Option 1: Put all functions into a single source file
 - Programmers link big object file into their programs
 - Space and time inefficient
- Option 2: Put each function in a separate source file
 - Programmers explicitly link appropriate binaries into their programs
 - More efficient, but burdensome on the programmer

Old-fashioned Solution: Static Libraries

- Static libraries (.a archive files)
 - Concatenate related relocatable object files into a single file with an index (called an archive).
 - Enhance linker so that it tries to resolve unresolved external references by looking for the symbols in one or more archives.
 - If an archive member file resolves reference, link it into the executable.

Creating Static Libraries



- Archiver allows incremental updates
- Recompile function that changes and replace .o file in archive.

Commonly Used Libraries

libc.a (the C standard library)

- 4.6 MB archive of 1496 object files.
- I/O, memory allocation, signal handling, string handling, data and time, random numbers, integer math

libm. a (the C math library)

- 2 MB archive of 444 object files.
- floating point math (sin, cos, tan, log, exp, sqrt, ...)

```
% ar -t /usr/lib/libc.a | sort
...
fork.o
...
fprintf.o
fpu_control.o
fputc.o
freopen.o
fscanf.o
fscanf.o
fseek.o
fstab.o
...
```

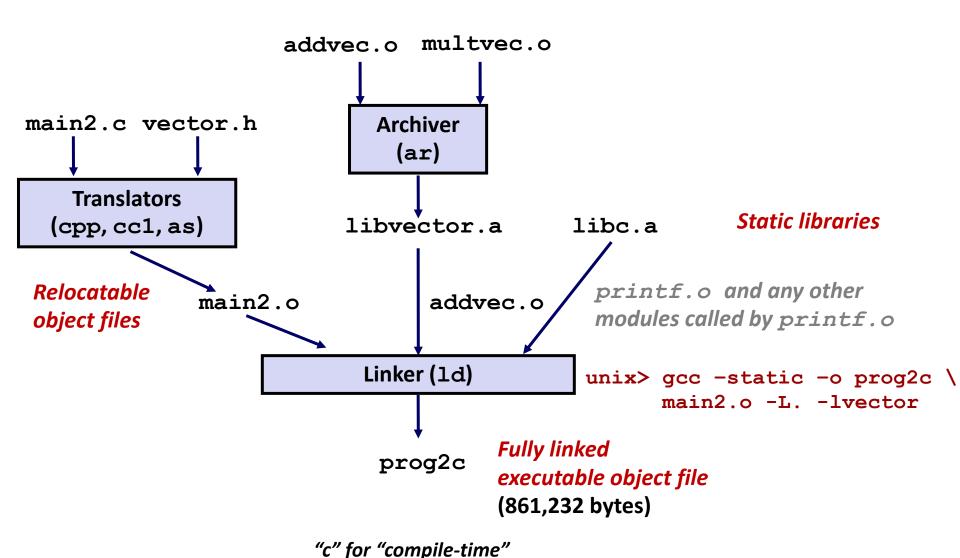
```
% ar -t /usr/lib/libm.a | sort
...
e_acos.o
e_acosf.o
e_acosh.o
e_acoshf.o
e_acoshl.o
e_acosl.o
e_asin.o
e_asinf.o
e_asinf.o
e_asinf.o
```

Linking with Static Libraries

```
#include <stdio.h>
#include "vector.h"
int x[2] = \{1, 2\};
int y[2] = \{3, 4\};
int z[2];
int main(int argc, char**
argv)
{
    addvec(x, y, z, 2);
   printf("z = [%d %d] \n",
           z[0], z[1]);
    return 0;
                   main2.c
```

libvector.a

Linking with Static Libraries



Using Static Libraries

Linker's algorithm for resolving external references:

- Scan .o files and .a files in the command line order.
- During the scan, keep a list of the current unresolved references.
- As each new .o or .a file, obj, is encountered, try to resolve each unresolved reference in the list against the symbols defined in obj.
- If any entries in the unresolved list at end of scan, then error.

Problem:

- Command line order matters!
- Moral: put libraries at the end of the command line.

```
unix> gcc -static -o prog2c -L. -lvector main2.o
main2.o: In function `main':
main2.c:(.text+0x19): undefined reference to `addvec'
collect2: error: ld returned 1 exit status
```

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Modern Solution: Shared Libraries

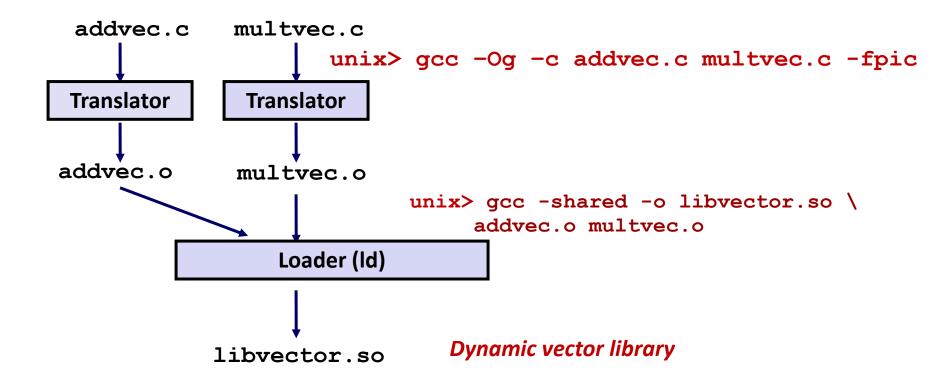
Static libraries have the following disadvantages:

- Duplication in the stored executables (every function needs libc)
- Duplication in the running executables
- Minor bug fixes of system libraries require each application to explicitly relink
 - Rebuild everything with glibc?
 - https://security.googleblog.com/2016/02/cve-2015-7547-glibc-getaddrinfo-stack.html

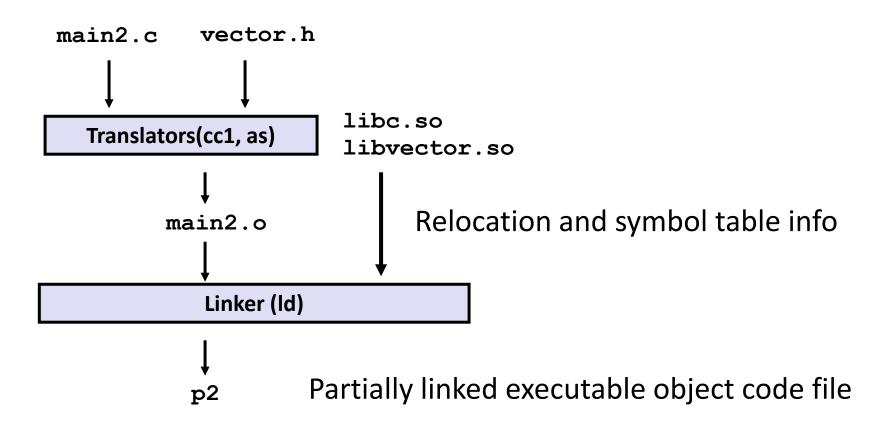
Modern solution: Shared Libraries

- Object files that contain code and data that are loaded and linked into an application dynamically, at either load-time or run-time
- Also called: dynamic link libraries, DLLs, .so files

Shared Library Example



Partially Linking with Shared Libraries



Unix>gcc -o p2 main2.c ./libvector.so

Partially Linking with Shared Libraries

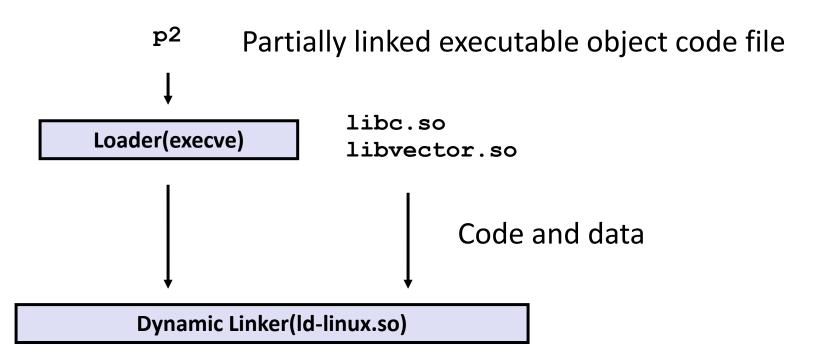
Which parts in libvector.so are copied into p2

The code and data sections

No

Relocation and symbol table information Some

Dynamic Linking at Load-time



Fully linked executable in memory

What have done by dynamic linker?

- Done by execve() & Id-linux.so
 - Copy code and data of libc.so and libvector.so into some memory segments
 - Relocate any references in p2 to symbols defined by libc.so and libvector.so
- After linking, the locations of the shared libraries are fixed and do not change during the execution time

What dynamic libraries are required?

.interp section

Specifies the dynamic linker to use (i.e., ld-linux.so)

.dynamic section

- Specifies the names, etc of the dynamic libraries to use
- Follow an example of prog

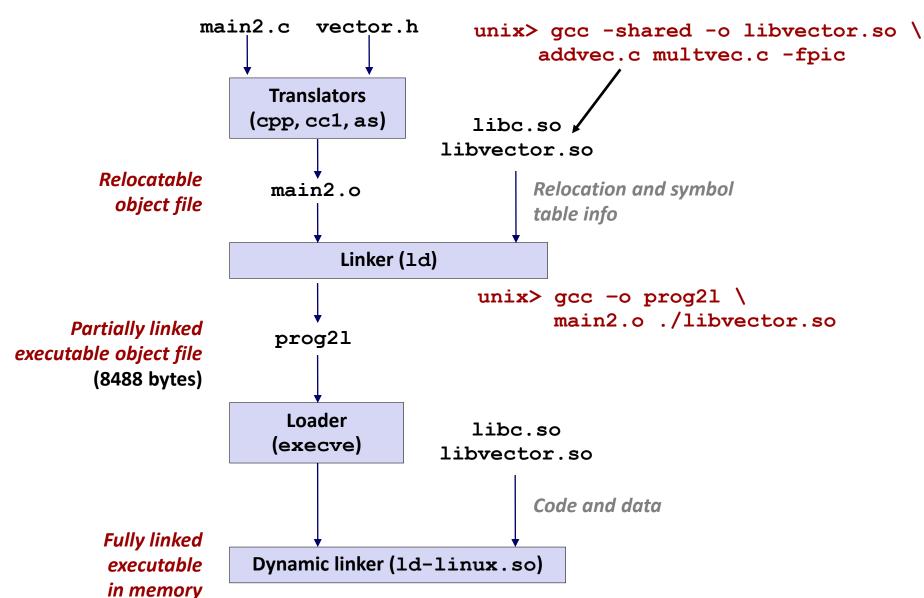
```
(NEEDED) Shared library: [libm.so.6]
```

Where are the libraries found?

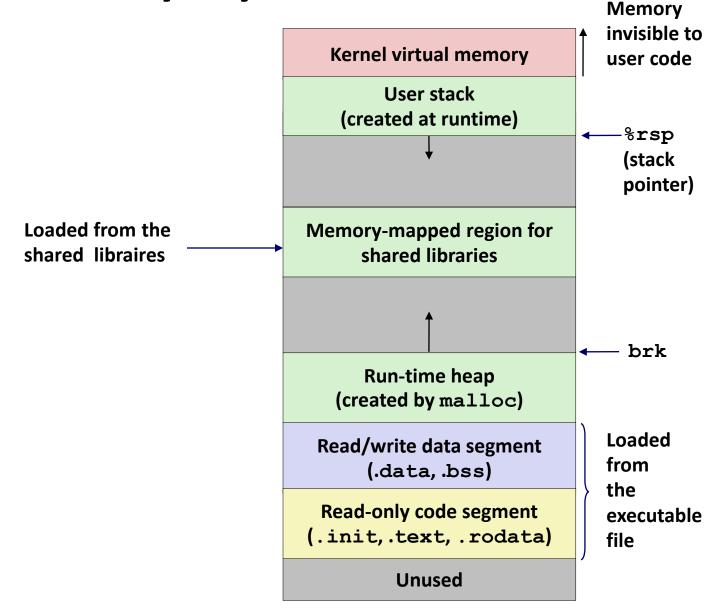
Use "ldd" to find out:

```
unix> ldd prog
linux-vdso.so.1 => (0x00007ffcf2998000)
libc.so.6 => /lib/x86_64-linux-gnu/libc.so.6 (0x00007f99ad927000)
/lib64/ld-linux-x86-64.so.2 (0x00007f99adcef000)
```

Dynamic Linking at Load-time (Complete)



Memory Layout for Shared Libraries



Dynamic Linking at Runtime

- Dynamic linking can occur when executable is first loaded and run (load-time linking).
 - Common case for Linux, handled automatically by the dynamic linker (ld-linux.so).
 - Standard C library (libc.so) usually dynamically linked.
- Dynamic linking can also occur after program has begun (run-time linking).
 - In Linux, this is done by calls to the dlopen() interface.
 - Distributing software.
 - High-performance web servers.
 - Runtime library interpositioning.

Why Linking at Run-time?

Distributing software

- Developers of Microsoft Windows applications frequently use shared libraries to distribute software updates.
- They generate a new copy of a shared library, which users can then download and use as a replacement for the current version.
- The next time they run their application, it will automatically link and load the new shared library.

Why Linking at Run-time?

Building high-performance Web servers

- Modern high-performance Web servers can generate dynamic content using a more efficient and sophisticated approach based on dynamic linking.
- package each function that generates dynamic content in a shared library.
- When a request arrives from a Web browser, the server dynamically loads and links the appropriate function and then calls it directly.

Why Linking at Run-time?

Building high-performance Web servers

- The function remains cached in the server's address space, so subsequent requests can be handled at the cost of a simple function call.
- This can have a significant impact on the throughput of a busy site.
- Further, existing functions can be updated, and new functions can be added at run time, without stopping the server.

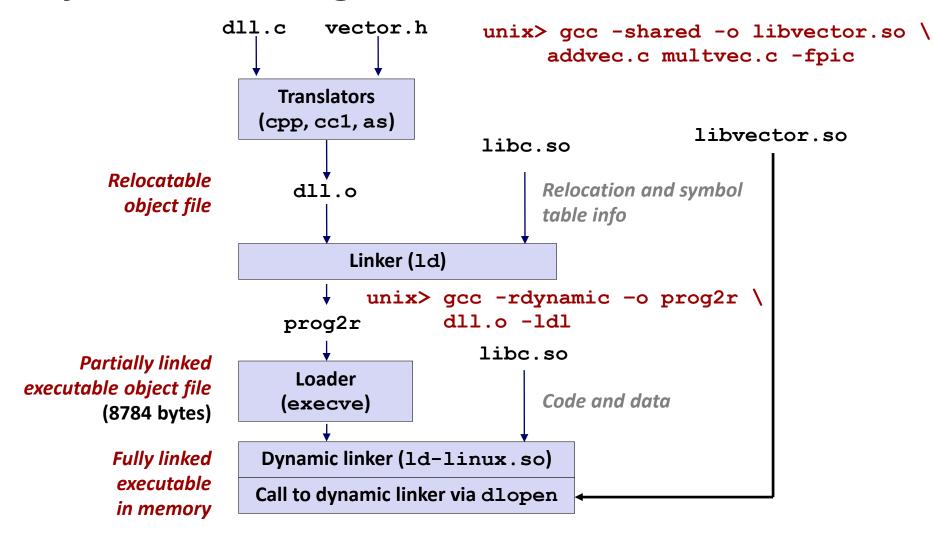
Dynamic Linking at Run-time

```
#include <stdio.h>
#include <stdlib.h>
#include <dlfcn.h>
int x[2] = \{1, 2\};
int y[2] = \{3, 4\};
int z[2];
int main(int argc, char** argv)
{
   void *handle;
   void (*addvec)(int *, int *, int *, int);
    char *error;
    /* Dynamically load the shared library that contains addvec() */
    handle = dlopen("./libvector.so", RTLD LAZY);
    if (!handle) {
        fprintf(stderr, "%s\n", dlerror());
       exit(1);
                                                                dll.c
```

Dynamic Linking at Run-time (cont)

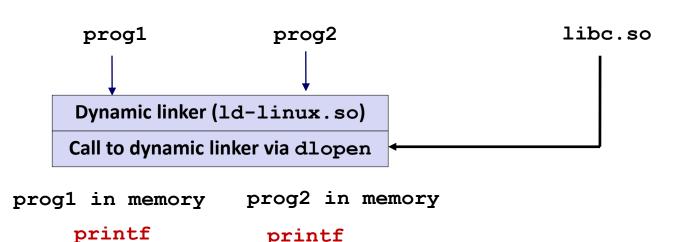
```
/* Get a pointer to the addvec() function we just loaded */
addvec = dlsym(handle, "addvec");
if ((error = dlerror()) != NULL) {
    fprintf(stderr, "%s\n", error);
    exit(1);
/* Now we can call addvec() just like any other function */
addvec(x, y, z, 2);
printf("z = [%d %d]\n", z[0], z[1]);
/* Unload the shared library */
if (dlclose(handle) < 0) {</pre>
    fprintf(stderr, "%s\n", dlerror());
    exit(1);
return 0;
                                                        d11.c
```

Dynamic Linking at Run-time



Share Libraires across Executables

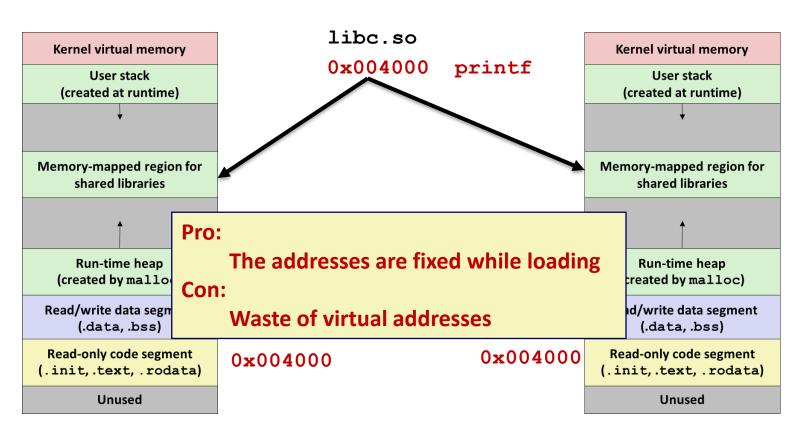
- Space: Libraries. How do libraries save space?
 - Option 2: Dynamic linking
 - Executable files contain no library code
 - During execution, single copy of library code can be shared across all executing processes



Fully linked executable in memory

How to know the address of printf?

- Naïve Solution: Fixed address
 - libc.so fixes the address of each function
 - Process reserves those addresses while loading



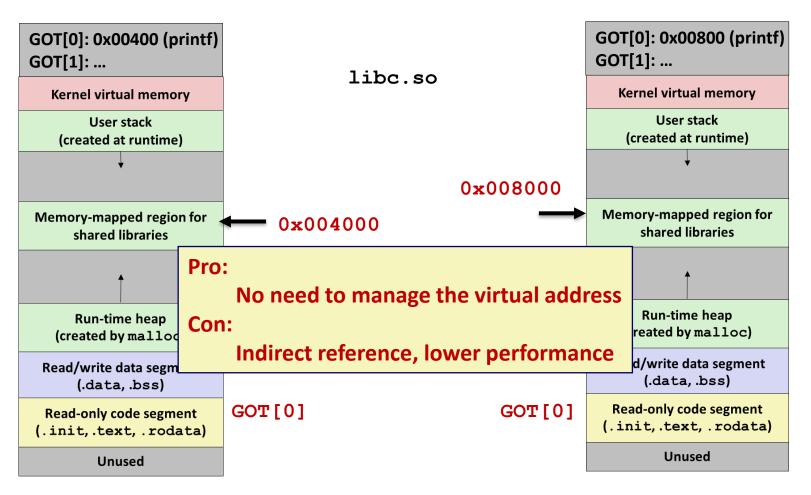
Position Independent Code (PIC)

- Code that can be execute from any address
- Internally-defined procedures
 - PC-relative reference

```
00000000004004d0 <main>:
 4004d0:
              48 83 ec 08
                                sub
                                      $0x8,%rsp
 4004d4:
              be 02 00 00 00
                                      $0x2,%esi
                                mov
 4004d9:
              bf 18 10 60 00
                                      $0x601018, %edi # %edi = &array
                                mov
                                      4004e8 <sum>
 4004de:
              e8 05 00 00 00
                                callq
                                                     # sum()
 4004e3:
              48 83 c4 08
                                add
                                      $0x8,%rsp
 4004e7:
               c3
                                retq
00000000004004e8 <sum>:
 4004e8:
              b8 00 00 00 00
                                            $0x0,%eax
                                     mov
 4004ed:
              ba 00 00 00 00
                                            $0x0, %edx
                                     mov
 4004f2:
              eb 09
                                     jmр
                                            4004fd < sum + 0x15 >
 4004f4:
              48 63 ca
                                     movslq %edx,%rcx
 4004f7:
              03 04 8f
                                     add (%rdi,%rcx,4),%eax
              83 c2 01
 4004fa:
                                     add
                                            $0x1,%edx
              39 f2
 4004fd:
                                            %esi,%edx
                                     cmp
 4004ff:
              7c f3
                                            4004f4 < sum + 0xc >
                                     il
 400501:
               f3 c3
                                     repz retq
```

Position Independent Code (PIC)

- Externally-defined procedures and global variables
 - Global offset table (GOT)



Position-Independent Code (PIC)

PIC Data References

```
Data segment
Global offset table (GOT)
GOT[0]:...
GOT[1]:...
GOT[2]:...
GOT[3]: &addcnt

Code segment
addvec:
mov 0x2008b9(%rip), %rax #%rax=*GOT[3]=&addcnt
→addl $0x1, (%rax) # addcnt++
```

Linking Summary

- Linking is a technique that allows programs to be constructed from multiple object files.
- Linking can happen at different times in a program's lifetime:
 - Compile time (when a program is compiled)
 - Load time (when a program is loaded into memory)
 - Run time (while a program is executing)
- Understanding linking can help you avoid nasty errors and make you a better programmer.

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Case Study: Library Interpositioning

- Library interpositioning: powerful linking technique that allows programmers to intercept calls to arbitrary functions
- Interpositioning can occur at:
 - Compile time: When the source code is compiled
 - Link time: When the relocatable object files are statically linked to form an executable object file
 - Load/run time: When an executable object file is loaded into memory, dynamically linked, and then executed.

Some Interpositioning Applications

Security

- Confinement (sandboxing)
- Behind the scenes encryption

Debugging

- In 2014, two Facebook engineers debugged a treacherous 1-year old bug in their iPhone app using interpositioning
- Code in the SPDY networking stack was writing to the wrong location
- Solved by intercepting calls to Posix write functions (write, writev, pwrite)

Source: Facebook engineering blog post at:

https://code.facebook.com/posts/313033472212144/debugging-file-corruption-on-ios/

Some Interpositioning Applications (cont)

Monitoring and Profiling

- Count number of calls to functions
- Characterize call sites and arguments to functions
- Malloc tracing
 - Detecting memory leaks
 - Generating address traces

Error Checking

- C Programming Lab used customized versions of malloc/free to do careful error checking
- Other labs (malloc, shell, proxy) also use interpositioning to enhance checking capabilities

Example program

```
#include <stdio.h>
#include <malloc.h>
#include <stdlib.h>
int main(int argc, char *argv[])
{
  int i;
  for (i = 1; i < argc; i++) {
    void *p =
          malloc(atoi(argv[i]));
    free(p);
  }
  return(0);
                             int.c
```

Goal: trace the addresses and sizes of the allocated and freed blocks, without breaking the program, and without modifying the source code.

Three solutions: interpose on the library malloc and free functions at compile time, link time, and load/run time.

Compile-time Interpositioning

```
#ifdef COMPILETIME
#include <stdio.h>
#include <malloc.h>
/* malloc wrapper function */
void *mymalloc(size t size)
    void *ptr = malloc(size);
    printf("malloc(%d)=%p\n", (int)size, ptr);
    return ptr;
/* free wrapper function */
void myfree(void *ptr)
    free (ptr) ;
    printf("free(%p)\n", ptr);
#endif
                                                     mymalloc.c
```

Compile-time Interpositioning

```
#define malloc(size) mymalloc(size)
#define free(ptr) myfree(ptr)
void *mymalloc(size t size);
void myfree(void *ptr);
                                                            malloc.h
linux> make intc
qcc -Wall -DCOMPILETIME -c mymalloc.c
qcc -Wall -I. -o intc int.c mymalloc.o
linux> make runc
./intc 10 100 1000
                               Search for <malloc.h> leads to
malloc(10) = 0 \times 1 ba 70 \sqrt{0}
                               /usr/include/malloc.h
free (0x1ba7010)
malloc(100) = 0 \times 1 ba7030
free (0x1ba7030)
malloc(1000) = 0x1ba70a0
                             Search for <malloc.h> leads to
free (0x1ba70a0)
linux>
```

Link-time Interpositioning

```
#ifdef LINKTIME
#include <stdio.h>
void * real malloc(size t size);
void real free(void *ptr);
/* malloc wrapper function */
void * wrap malloc(size t size)
   void *ptr = real malloc(size); /* Call libc malloc */
   printf("malloc(%d) = p\n", (int)size, ptr);
   return ptr;
/* free wrapper function */
void wrap free(void *ptr)
    real free (ptr); /* Call libc free */
   printf("free(%p)\n", ptr);
#endif
```

Link-time Interpositioning

```
linux> make intl
gcc -Wall -DLINKTIME -c mymalloc.c
gcc -Wall -c int.c
gcc -Wall -Wl,--wrap, malloc -Wl,--wrap, free -o intl \
   int.o mymalloc.o
linux> make runl
./intl 10 100 1000
malloc(10) = 0x91a010
free(0x91a010)
. . .
```

- The "-W1" flag passes argument to linker, replacing each comma with a space.
- The "--wrap, malloc" arg instructs linker to resolve references in a special way:
 - Refs to malloc should be resolved as __wrap_malloc
 - Refs to __real_malloc should be resolved as malloc

Load/Run-time Interpositioning

```
#ifdef RUNTIME
                                           Interpositioning
#define GNU SOURCE
#include <stdio.h>
#include <stdlib.h>
                            Observe that DON'T have
#include <dlfcn.h>
                            #include <malloc.h>
/* malloc wrapper function */
void *malloc(size t size)
   void *(*mallocp) (size t size);
    char *error;
   mallocp = dlsym(RTLD NEXT, "malloc"); /* Get addr of libc malloc */
    if ((error = dlerror()) != NULL) {
        fputs(error, stderr);
       exit(1);
    char *ptr = mallocp(size); /* Call libc malloc */
   printf("malloc(%d) = %p\n", (int)size, ptr);
    return ptr;
                                                            mymalloc.c
```

Load/Run-time Interpositioning

```
/* free wrapper function */
void free(void *ptr)
   void (*freep) (void *) = NULL;
    char *error;
    if (!ptr)
        return;
    freep = dlsym(RTLD NEXT, "free"); /* Get address of libc free */
    if ((error = dlerror()) != NULL) {
        fputs(error, stderr);
        exit(1);
    freep(ptr); /* Call libc free */
   printf("free(%p)\n", ptr);
#endif
```

mymalloc.c

Load/Run-time Interpositioning

```
linux> make intr
gcc -Wall -DRUNTIME -shared -fpic -o mymalloc.so mymalloc.c -ldl
gcc -Wall -o intr int.c
linux> make runr
(LD_PRELOAD="./mymalloc.so" ./intr 10 100 1000)
malloc(10) = 0x91a010
free(0x91a010)
. . .
linux>
Search for <malloc.h> leads to
/usr/include/malloc.h
linux>
```

- The LD_PRELOAD environment variable tells the dynamic linker to resolve unresolved refs (e.g., to malloc) by looking in mymalloc.so first.
- Type into (some) shells as:

```
env LD_PRELOAD=./mymalloc.so ./intr 10 100 1000)
```

Interpositioning Recap

Compile Time

- Apparent calls to malloc/free get macro-expanded into calls to mymalloc/myfree
- Simple approach. Must have access to source & recompile

Link Time

- Use linker trick to have special name resolutions
 - malloc → wrap malloc
 - __real_malloc → malloc

Load/Run Time

- Implement custom version of malloc/free that use dynamic linking to load library malloc/free under different names
- Can use with ANY dynamically linked binary

```
env LD_PRELOAD=./mymalloc.so gcc -c int.c)
```

Linking Recap

- Usually: Just happens, no big deal
- Sometimes: Strange errors
 - Bad symbol resolution
 - Ordering dependence of linked .o, .a, and .so files
- For power users:
 - Interpositioning to trace programs with & without source