

Virtual Memory: Concepts

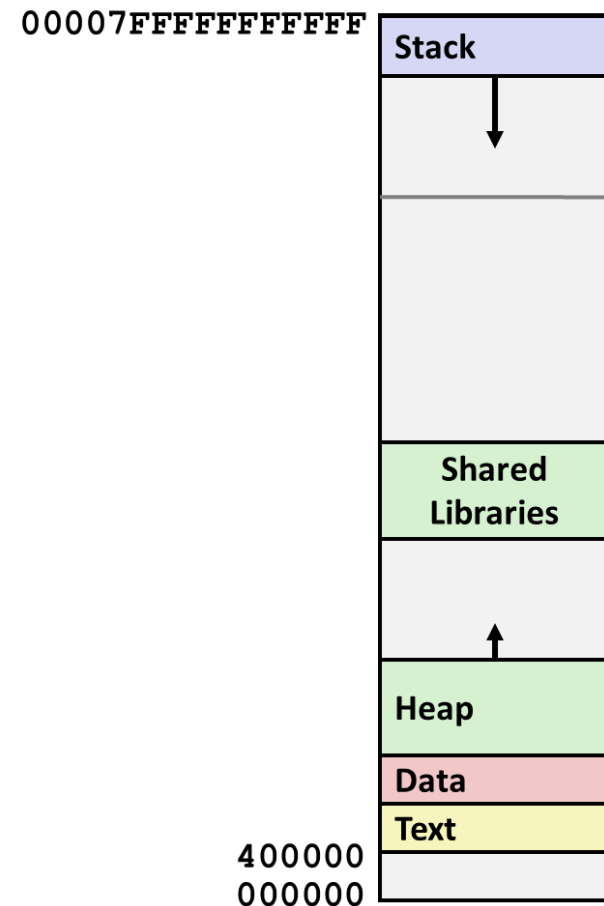
COMP400727: Introduction to Computer Systems

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This Picture is a Lie

- This is RAM, we said...
- But the computer can run more than one program at a time!
- Where are all the other programs?
- Let's *investigate*.



Processes (Recap)

■ **Definition:** A *process* is an instance of a running program.

- One of the most profound ideas in computer science
- Not the same as “program” or “processor”

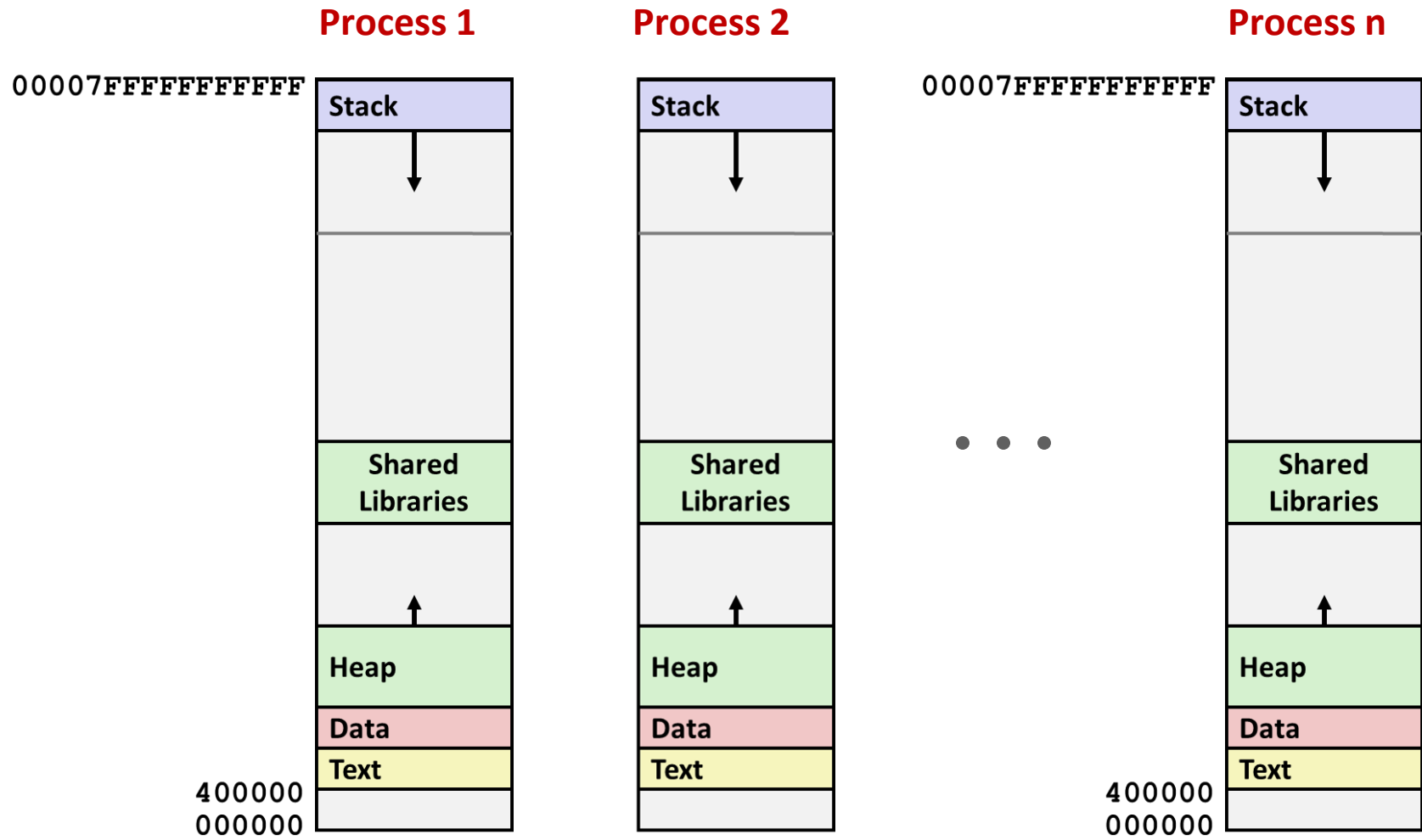
■ **Unix:** A *parent process* creates a new *child process* by calling `fork`

- Child is (sort of) a copy of the parent
- `fork` returns *twice*—once in each process
 - Different return value in each

■ **Parent can wait for child to finish by calling `waitpid`**

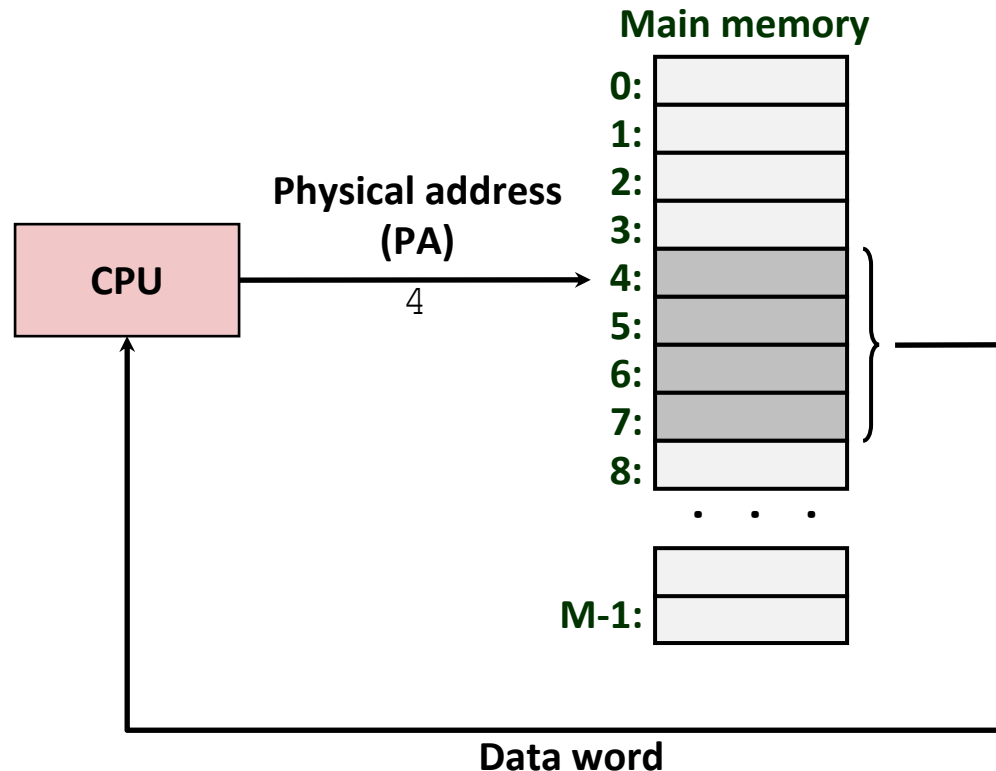
- For now, think of this as “what `main` returns to”

Hmmm, How Does This Work?!



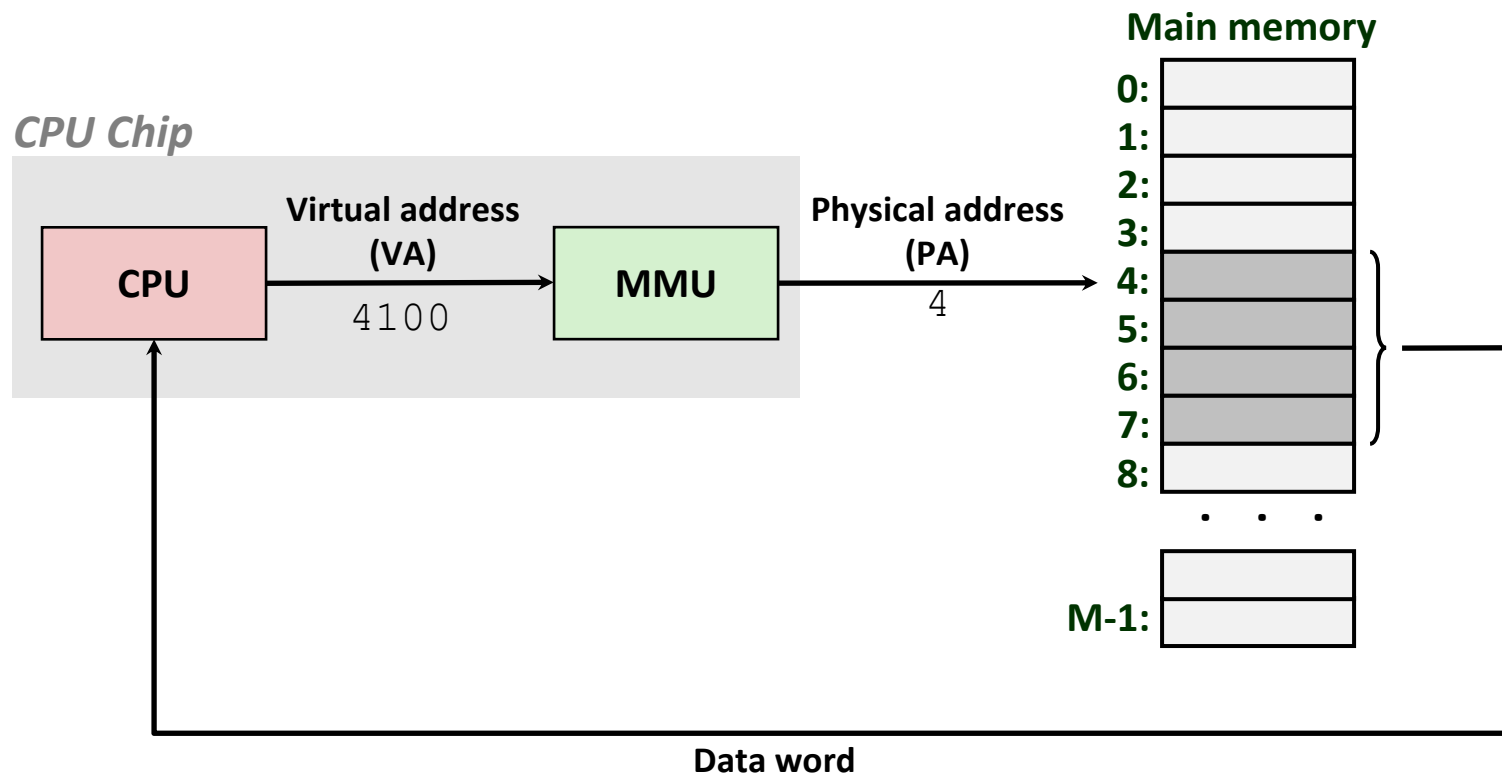
Solution: Virtual Memory (today and next lecture)

A System Using Physical Addressing



- Used in “simple” systems like embedded microcontrollers in devices like cars, elevators, and digital picture frames

A System Using Virtual Addressing

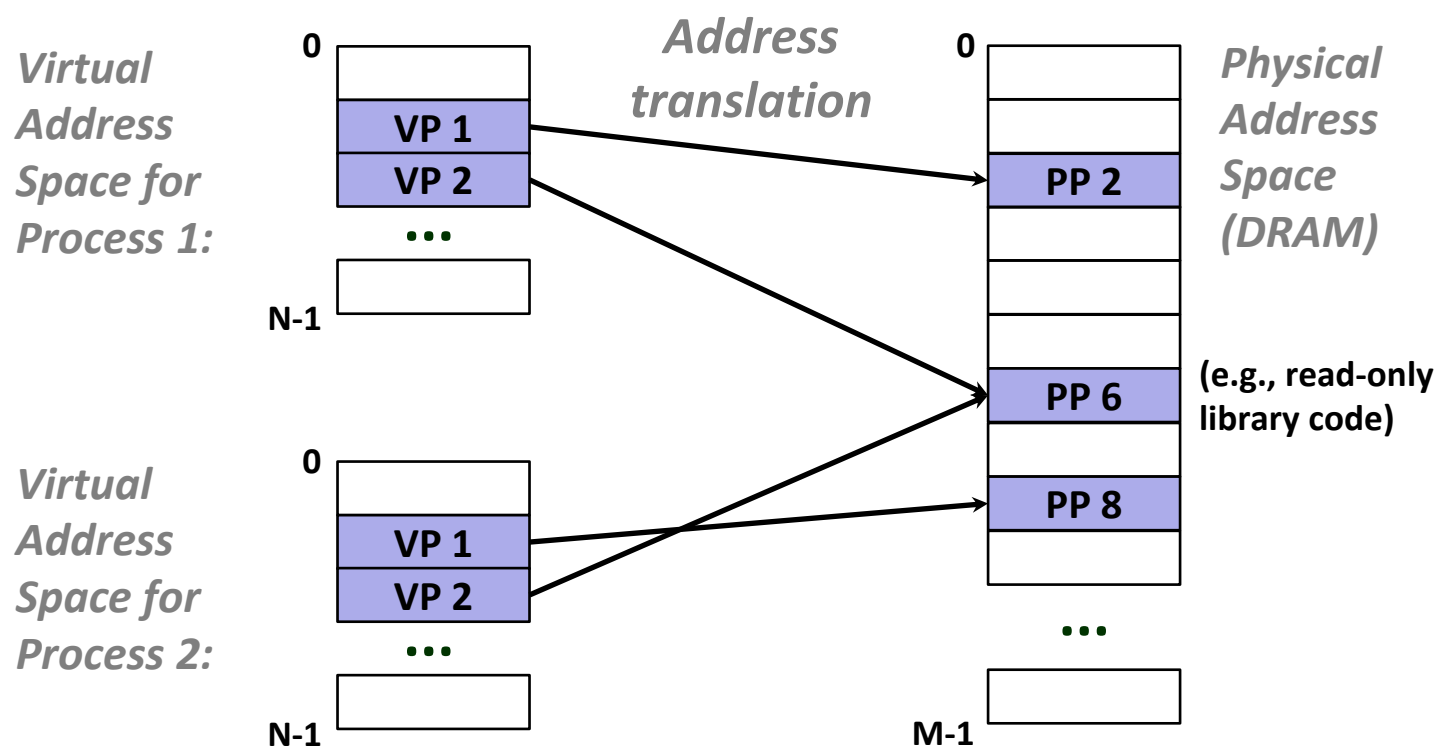


- Used in all modern servers, laptops, and smart phones
- One of the great ideas in computer science

VM as a Tool for Memory Management

■ Key idea: each process has its own virtual address space

- It can view memory as a simple linear array
- Mapping function scatters addresses through physical memory
 - Well-chosen mappings can improve locality



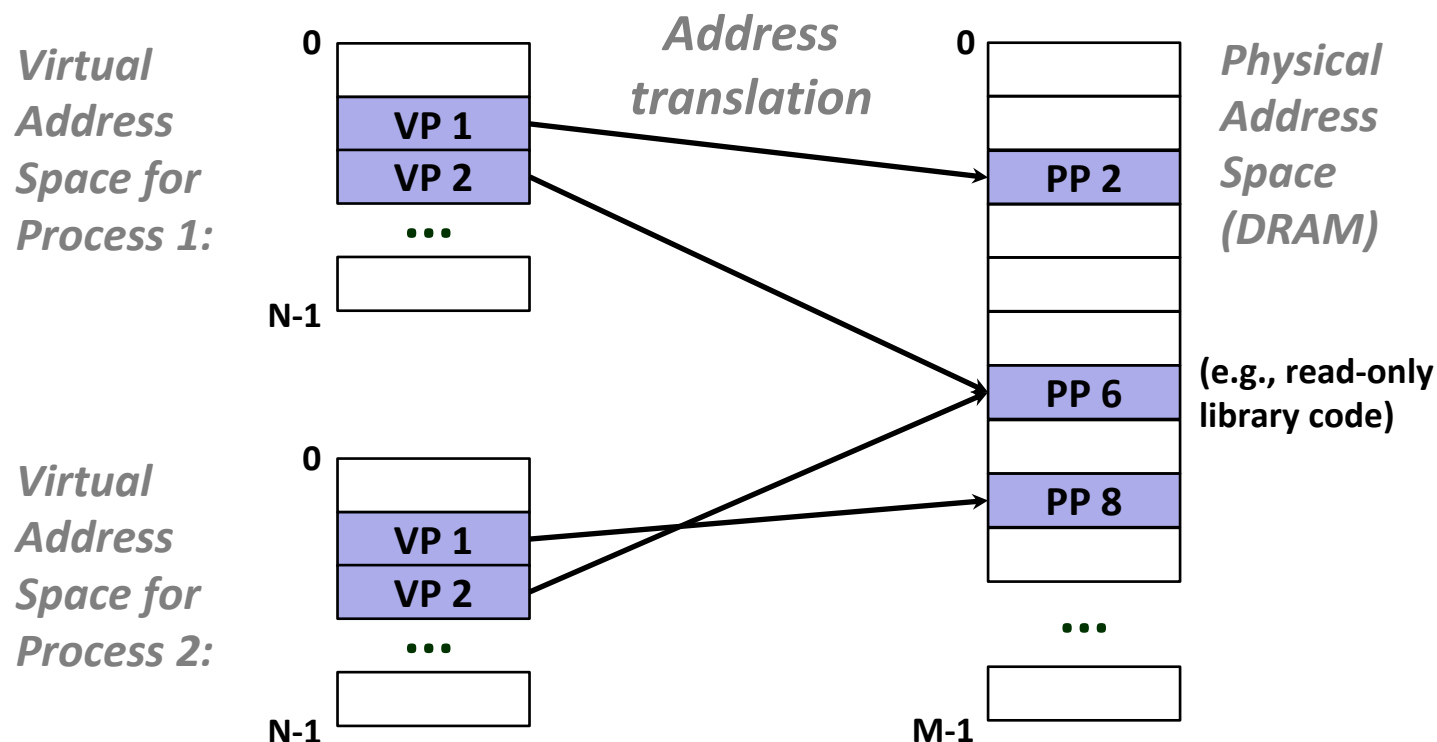
VM as a Tool for Memory Management

■ Simplifying memory allocation

- Each virtual page can be mapped to any physical page
- A virtual page can be stored in different physical pages at different times

■ Sharing code and data among processes

- Map virtual pages to the same physical page (here: PP 6)



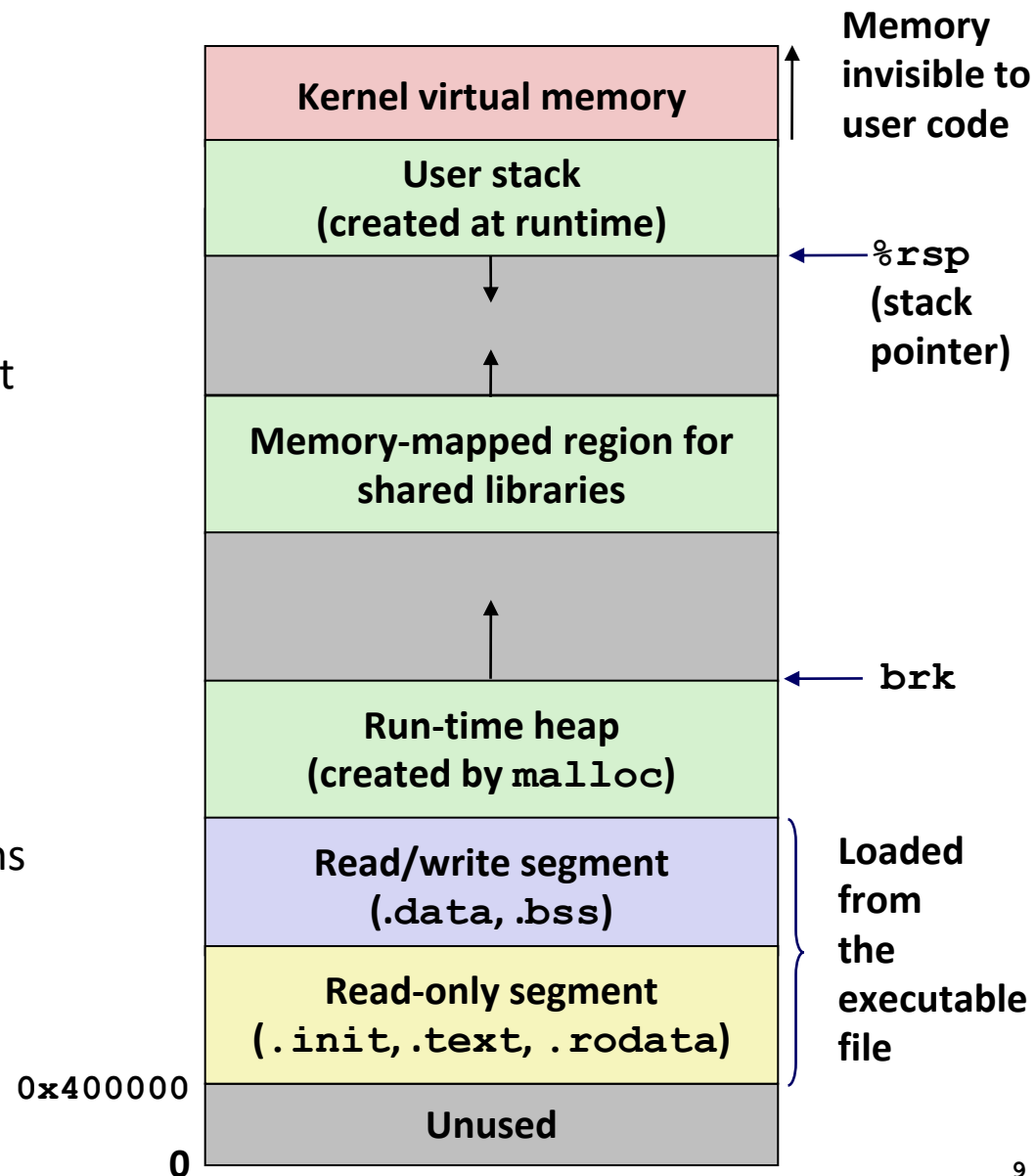
Simplifying Linking and Loading

■ Linking

- Each program has similar virtual address space
- Code, data, and heap always start at the same addresses.

■ Loading

- **execve** allocates virtual pages for `.text` and `.data` sections & creates PTEs marked as invalid
- The `.text` and `.data` sections are copied, page by page, on demand by the virtual memory system



Address Spaces

- **Linear address space:** Ordered set of contiguous non-negative integer addresses:

$\{0, 1, 2, 3 \dots \}$

- **Virtual address space:** Set of $N = 2^n$ virtual addresses

$\{0, 1, 2, 3, \dots, N-1\}$

- **Physical address space:** Set of $M = 2^m$ physical addresses

$\{0, 1, 2, 3, \dots, M-1\}$

Why Virtual Memory (VM)?

■ Uses main memory efficiently

- Use DRAM as a cache for parts of a virtual address space

■ Simplifies memory management

- Each process gets the same uniform linear address space

■ Isolates address spaces

- One process can't interfere with another's memory
- User program cannot access privileged kernel information and code

VM Address Translation

■ Virtual Address Space

- $V = \{0, 1, \dots, N-1\}$

■ Physical Address Space

- $P = \{0, 1, \dots, M-1\}$

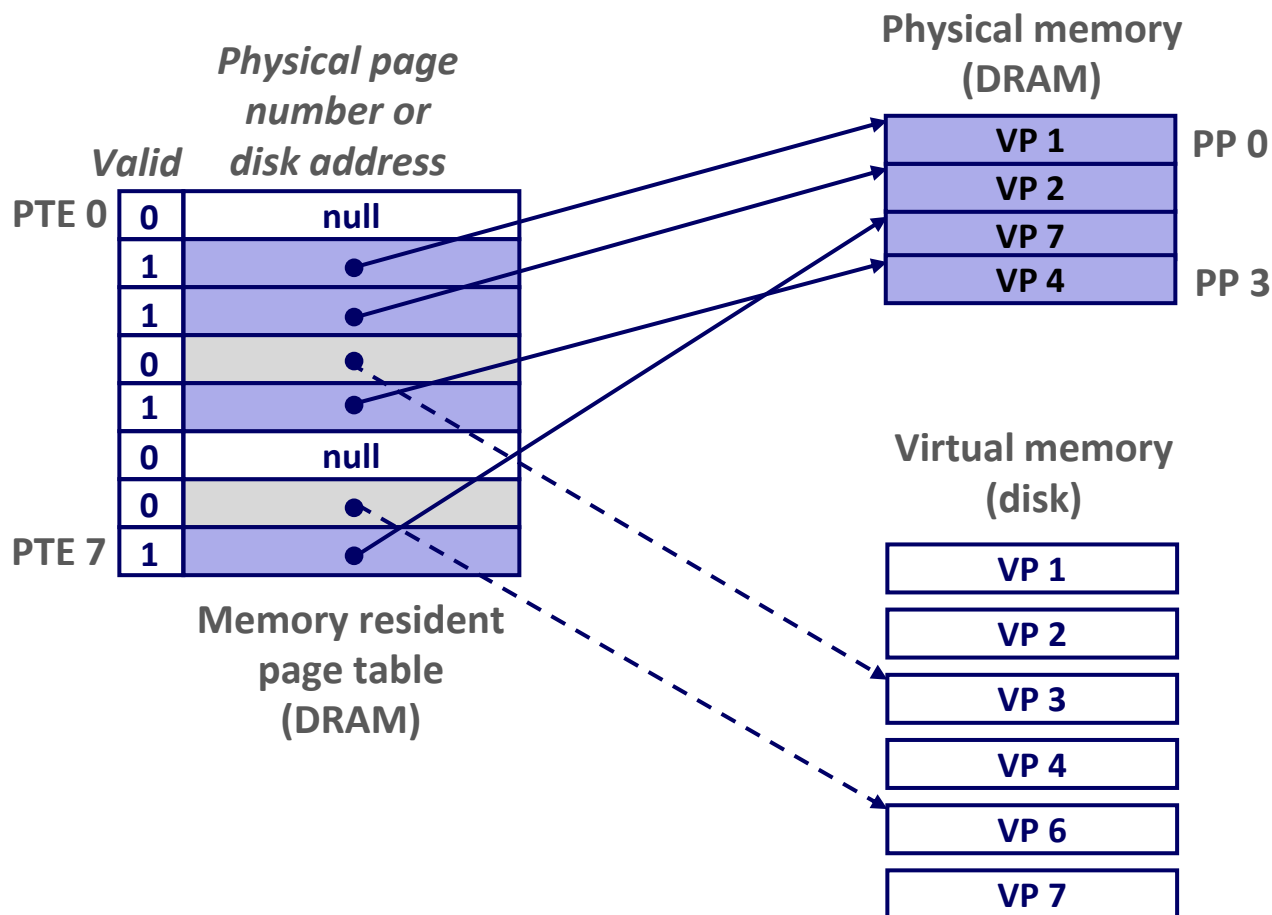
■ Address Translation

- $MAP: V \rightarrow P \cup \{\emptyset\}$
- For virtual address a :
 - $MAP(a) = a'$ if data at virtual address a is at physical address a' in P
 - $MAP(a) = \emptyset$ if data at virtual address a is not in physical memory
 - Either invalid or stored on disk

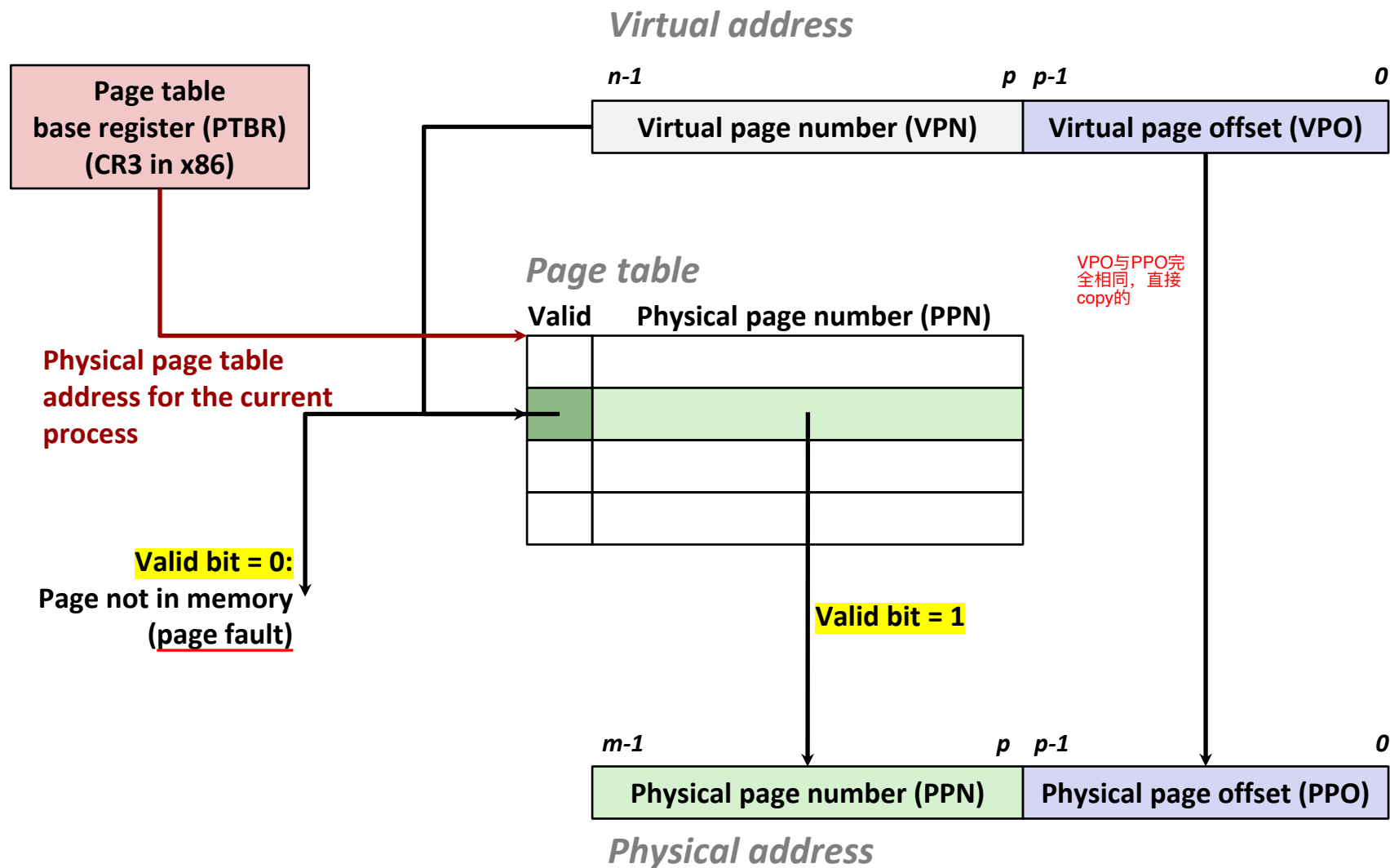
Enabling Data Structure: Page Table

■ A **page table** is an array of page table entries (PTEs) that maps virtual pages to physical pages.

- Per-process kernel data structure in DRAM

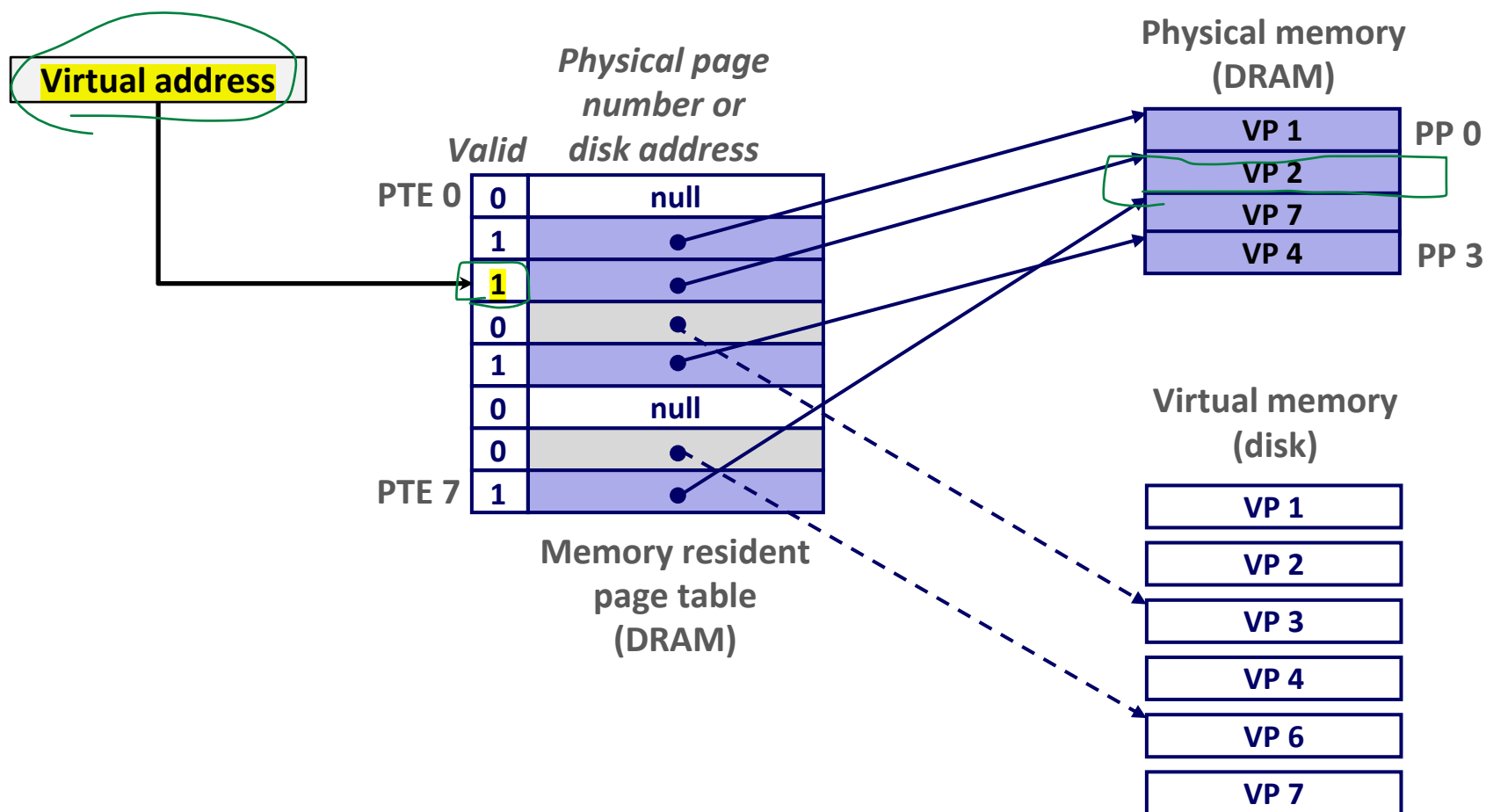


Address Translation With a Page Table

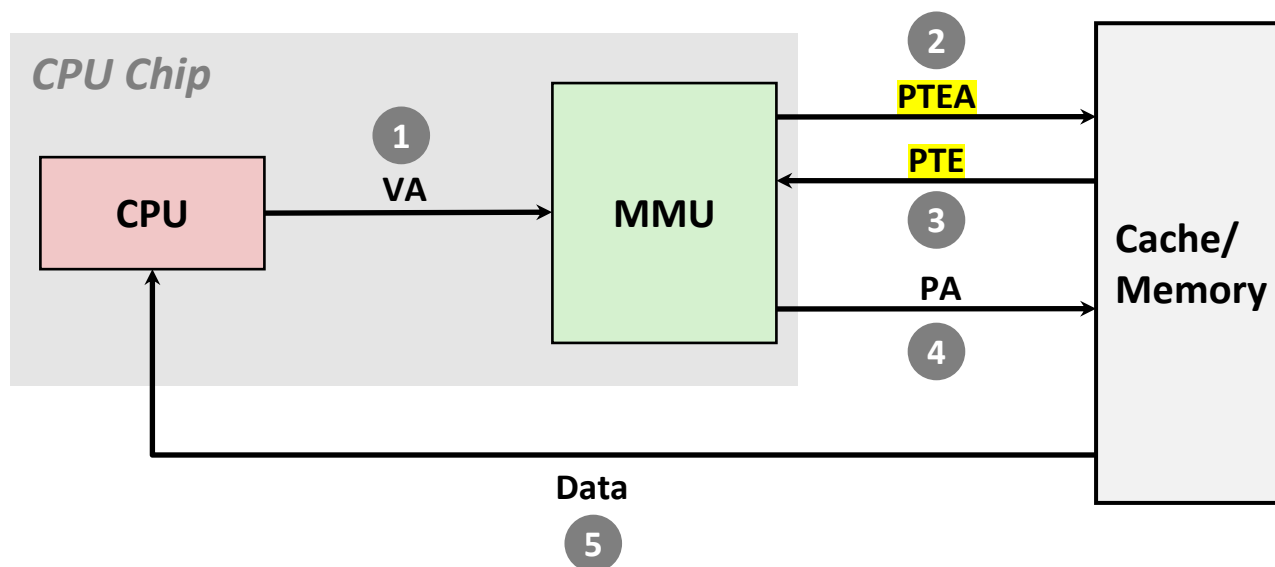


Page Hit

- **Page hit:** reference to VM word that is in physical memory (DRAM cache hit)



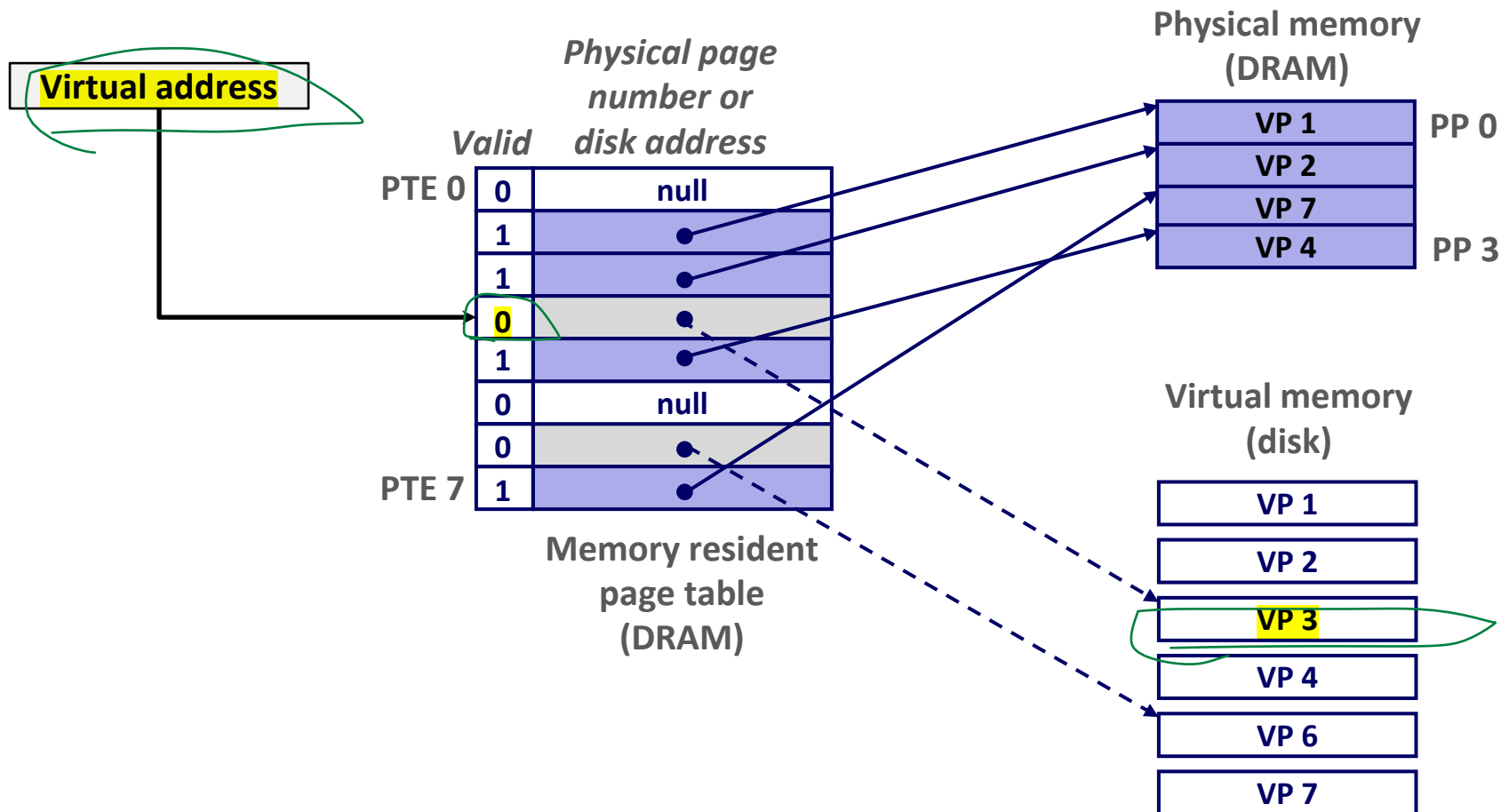
Address Translation: Page Hit



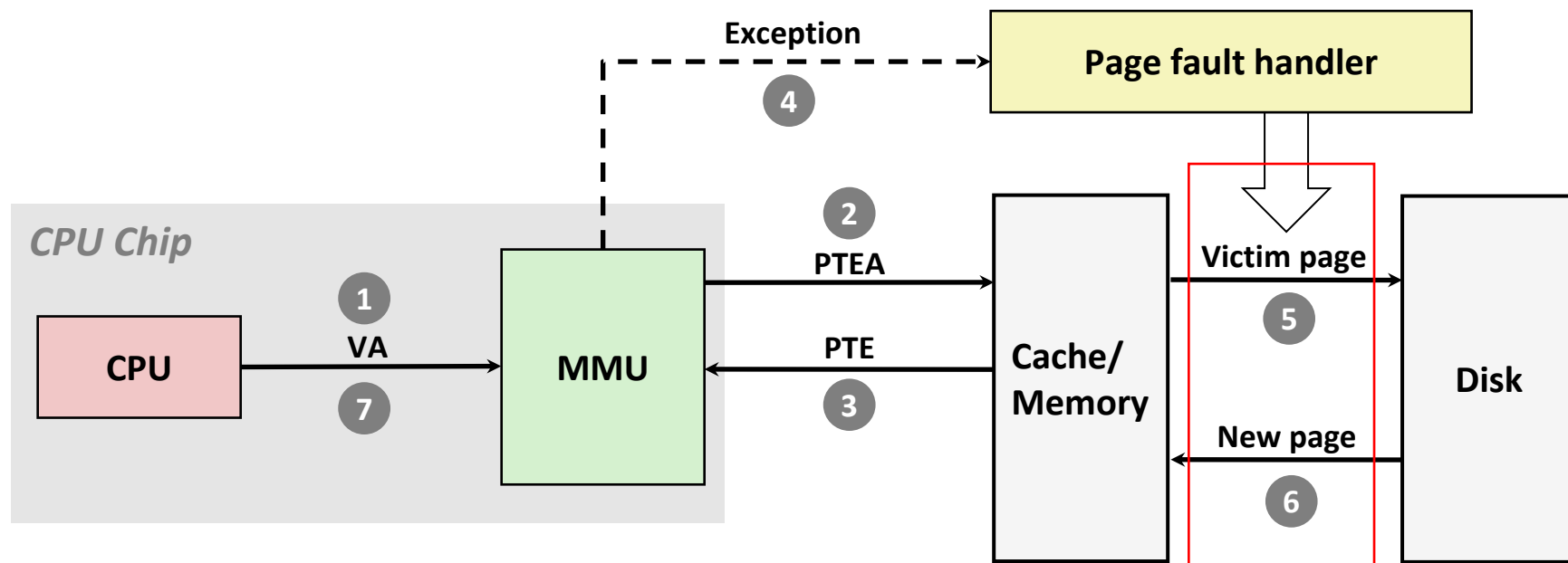
- 1) Processor sends virtual address to MMU
- 2-3) MMU fetches PTE from page table in memory
- 4) MMU sends physical address to cache/memory
Get the correct content in corresponding PA
- 5) Cache/memory sends data word to processor

Page Fault

■ **Page fault:** reference to VM word that is not in physical memory (DRAM cache miss)



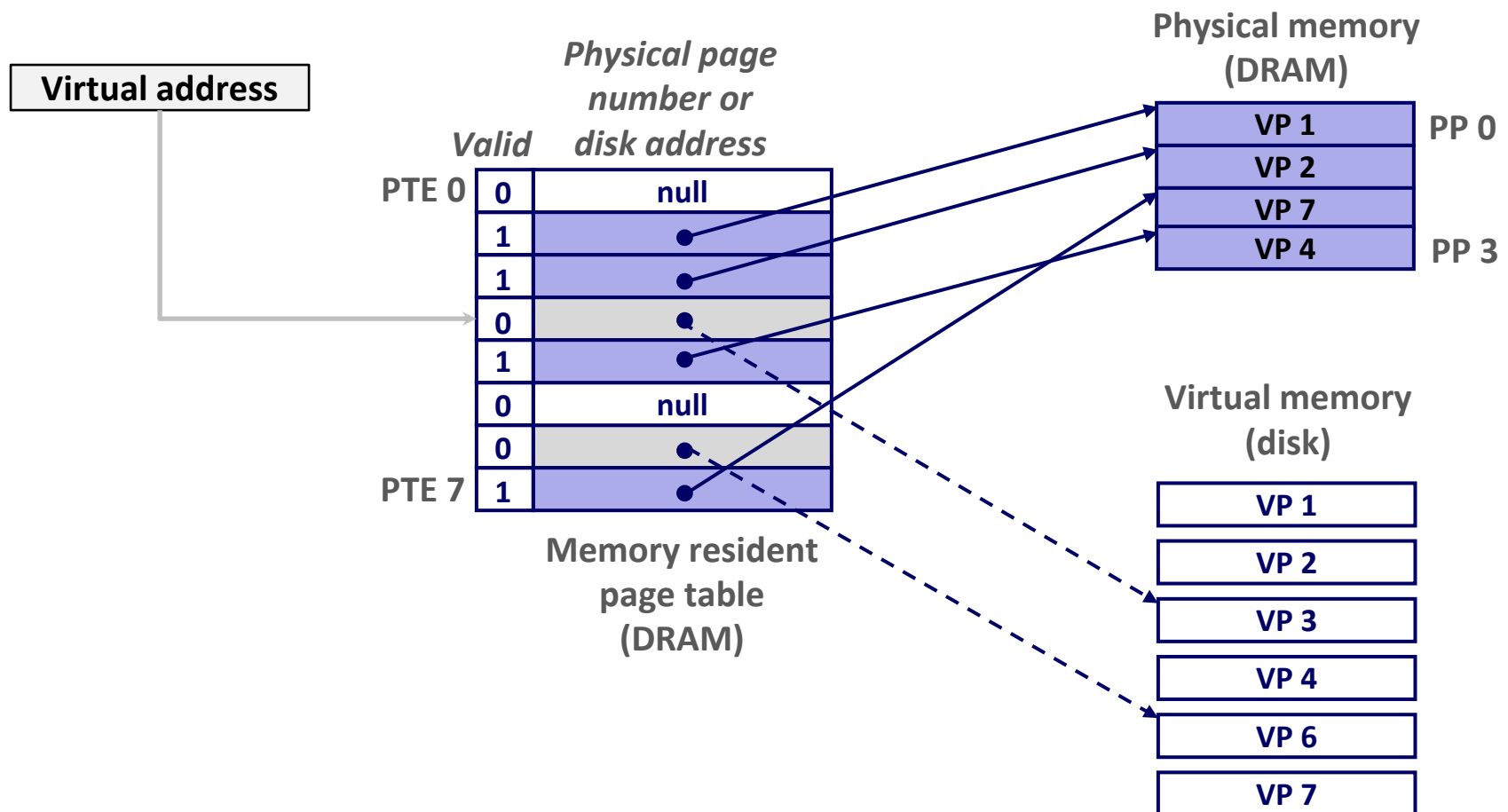
Address Translation: Page Fault



- 1) Processor sends virtual address to MMU
- 2-3) MMU fetches PTE from page table in memory
- 4) **Valid bit is zero**, so **MMU triggers page fault exception**
- 5) Handler identifies victim (and, if dirty, pages it out to disk)
- 6) Handler pages in new page and **updates PTE in memory**
- 7) Handler **returns to original process, restarting faulting instruction**

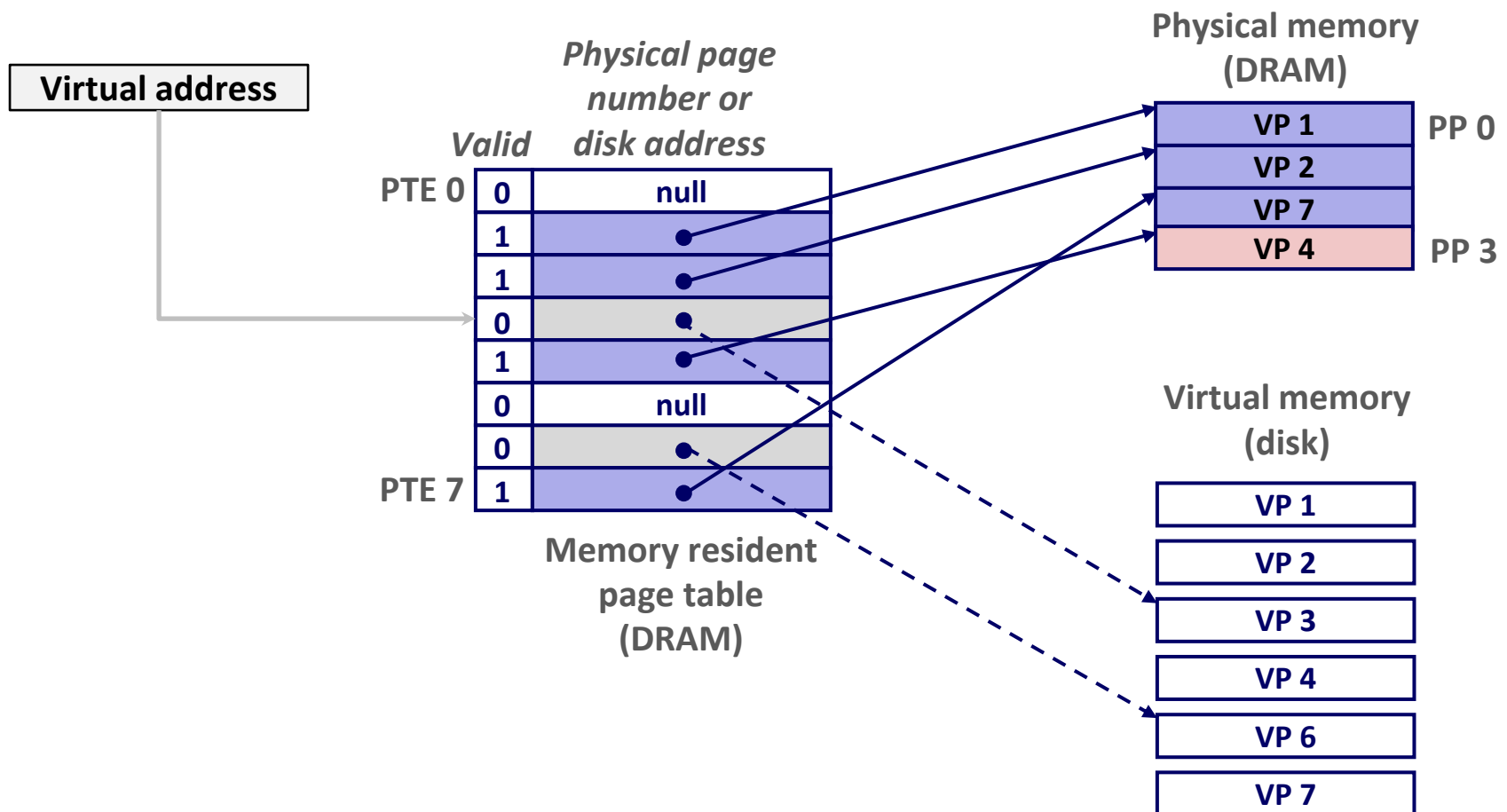
Handling Page Fault

- Page miss causes page fault (an exception)



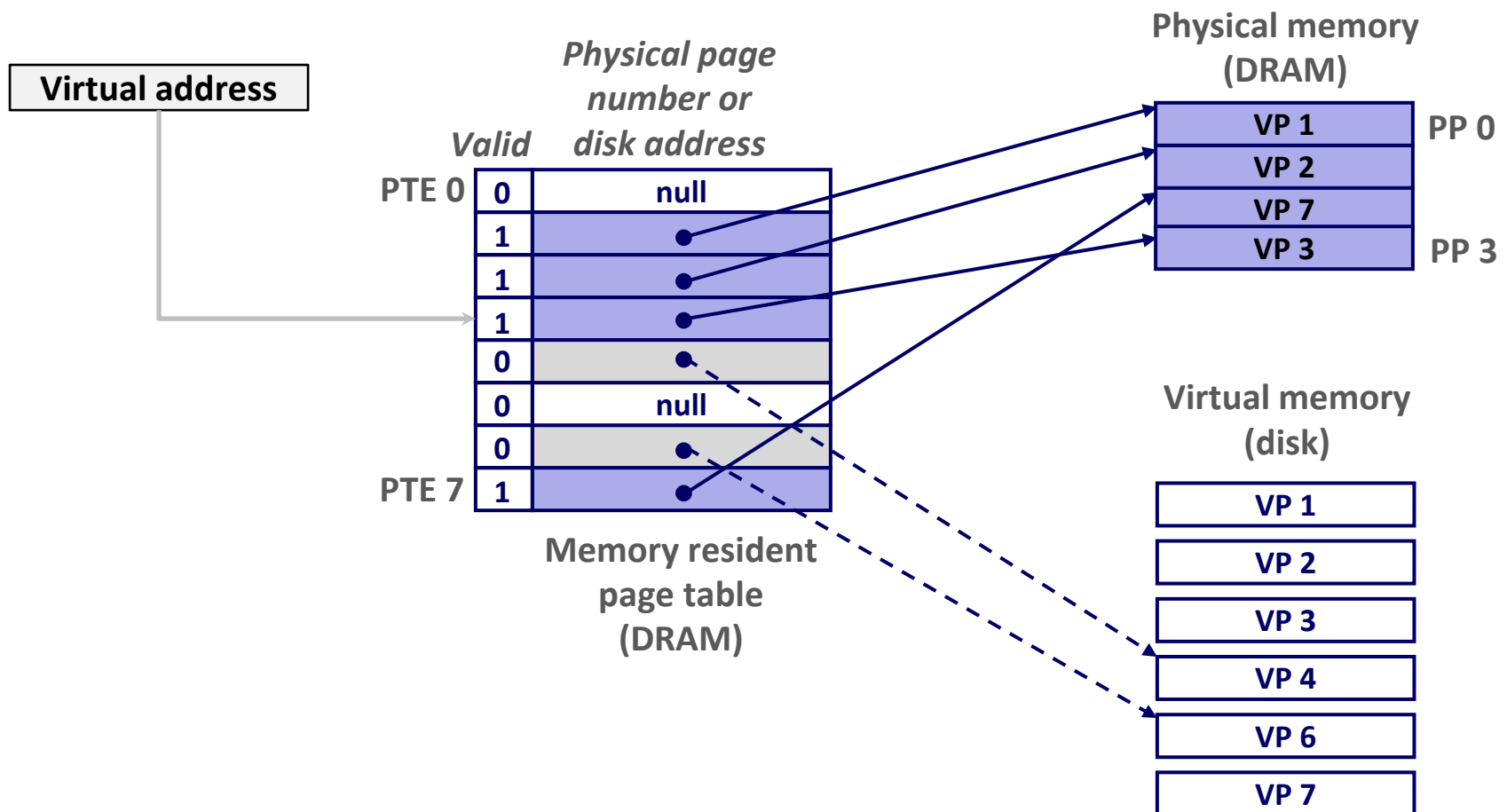
Handling Page Fault

- Page miss causes page fault (an exception)
- Page fault handler selects a victim to be evicted (here VP 4)



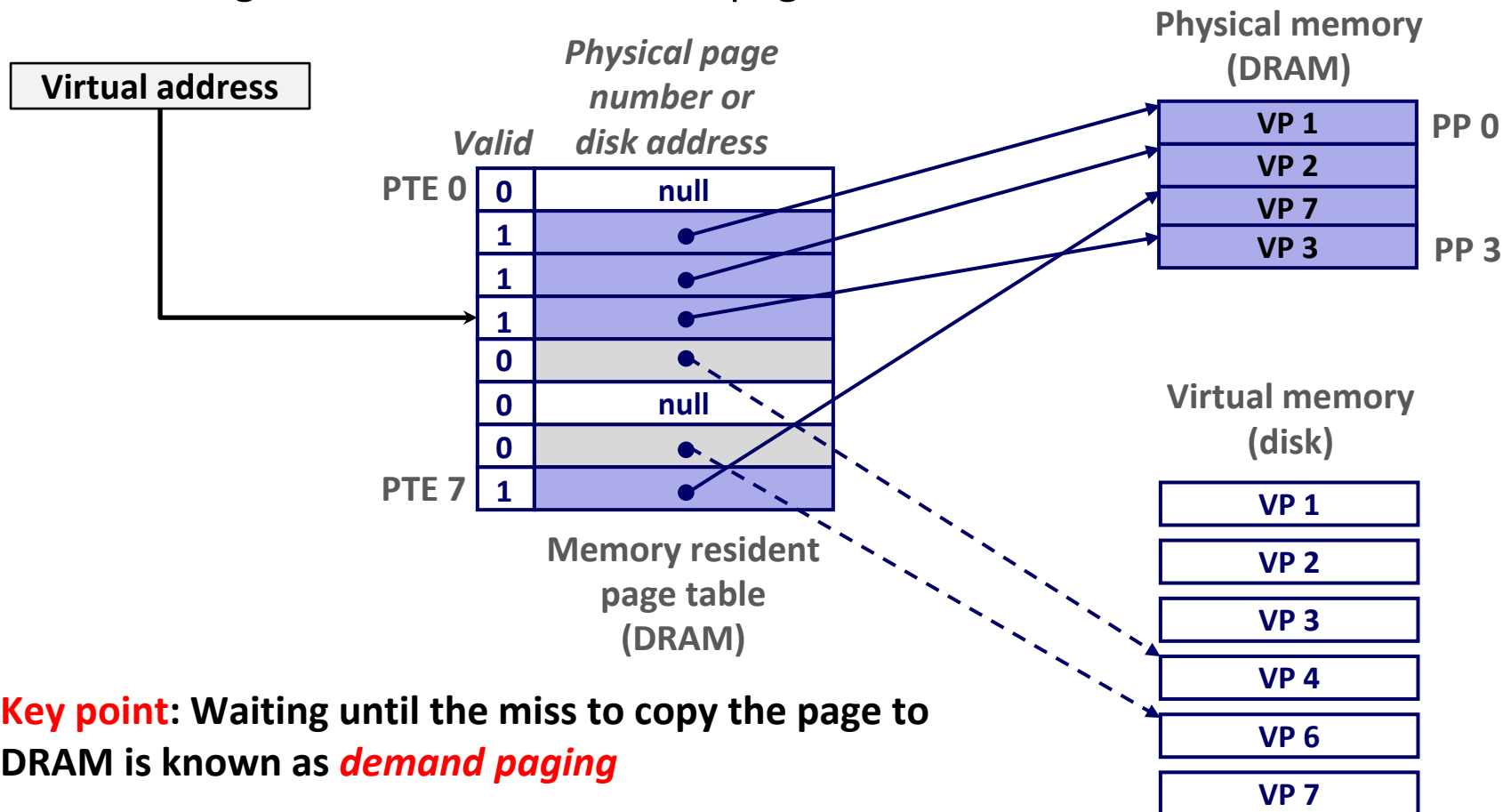
Handling Page Fault

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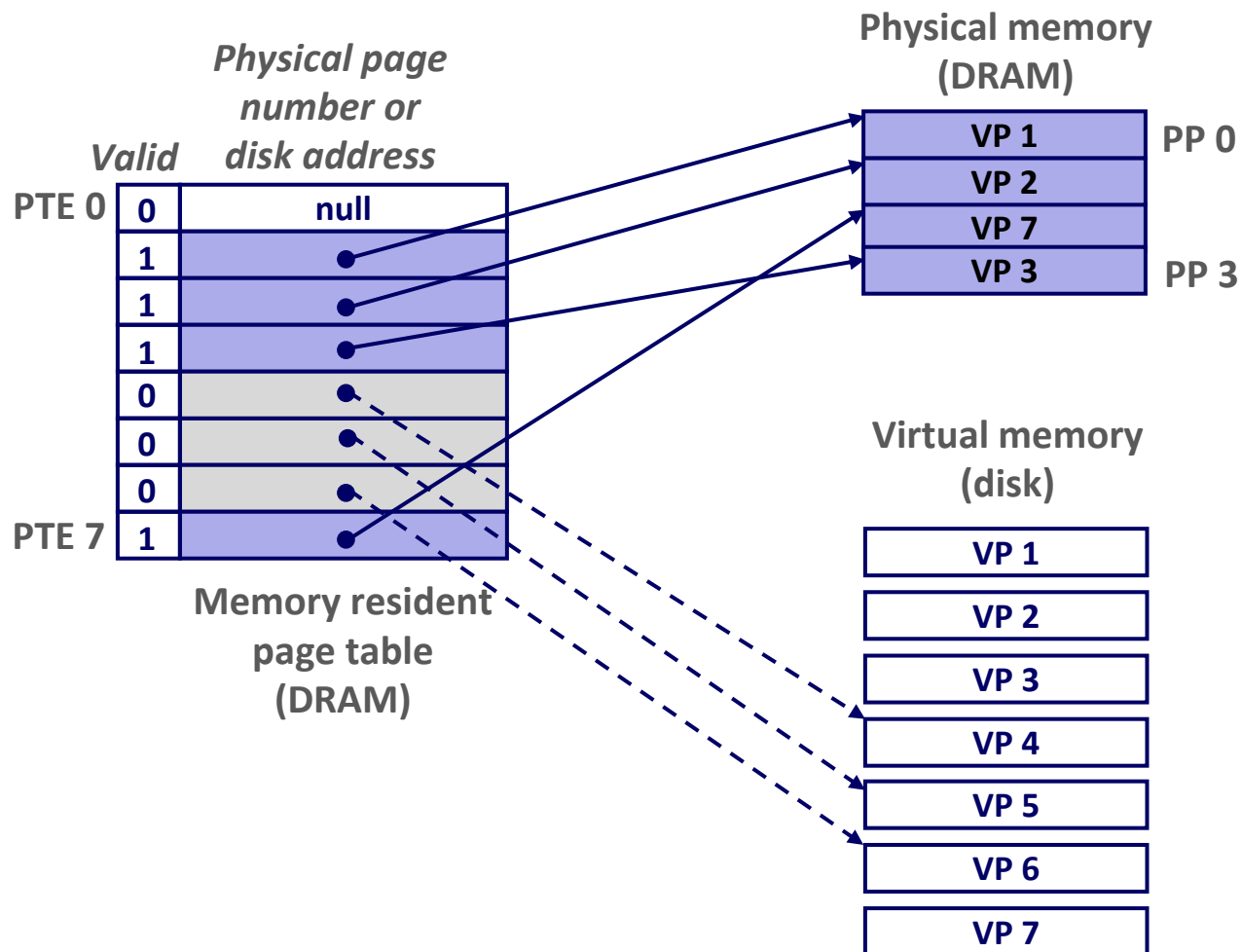
Handling Page Fault

- Page miss causes page fault (an exception)
- Page fault handler selects a victim to be evicted (here VP 4)
- Offending instruction is restarted: page hit!



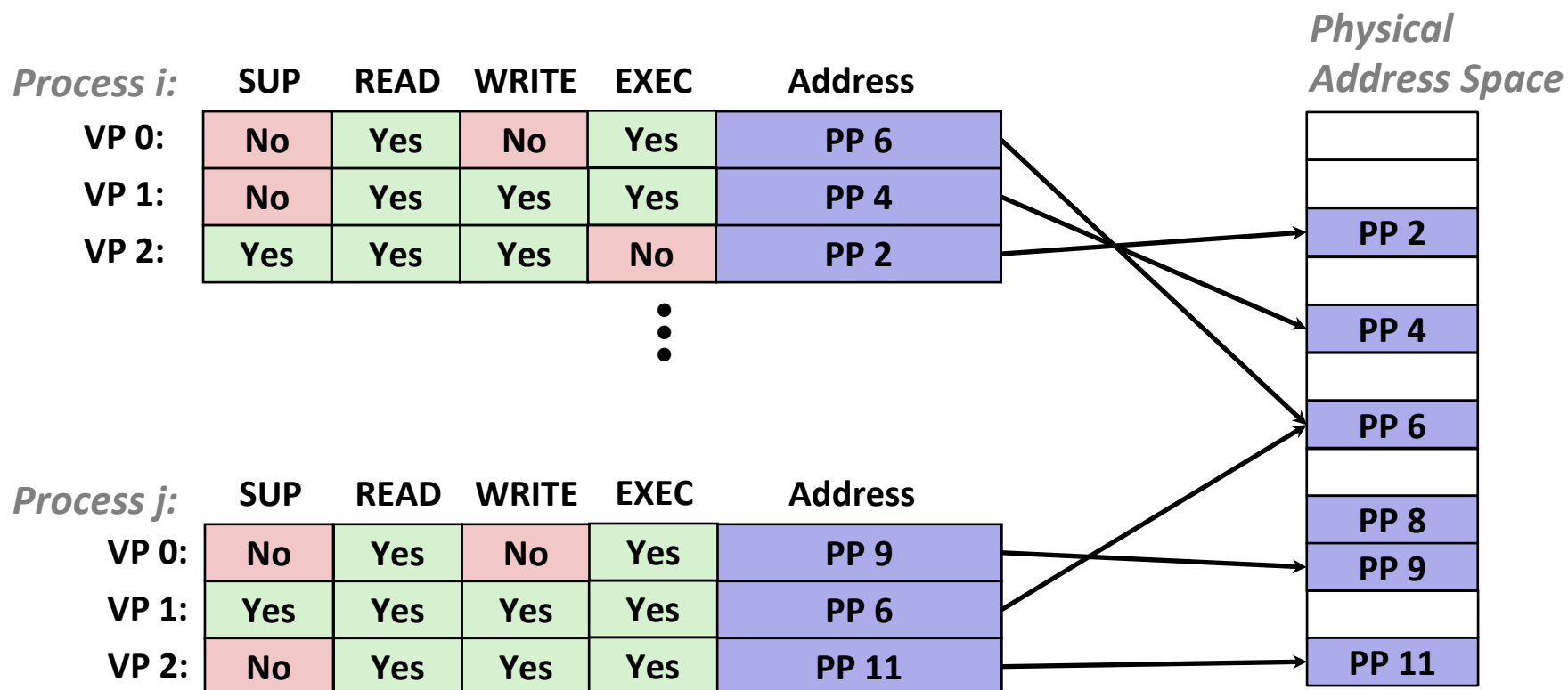
Allocating Pages

- Allocating a new page (VP 5) of virtual memory.



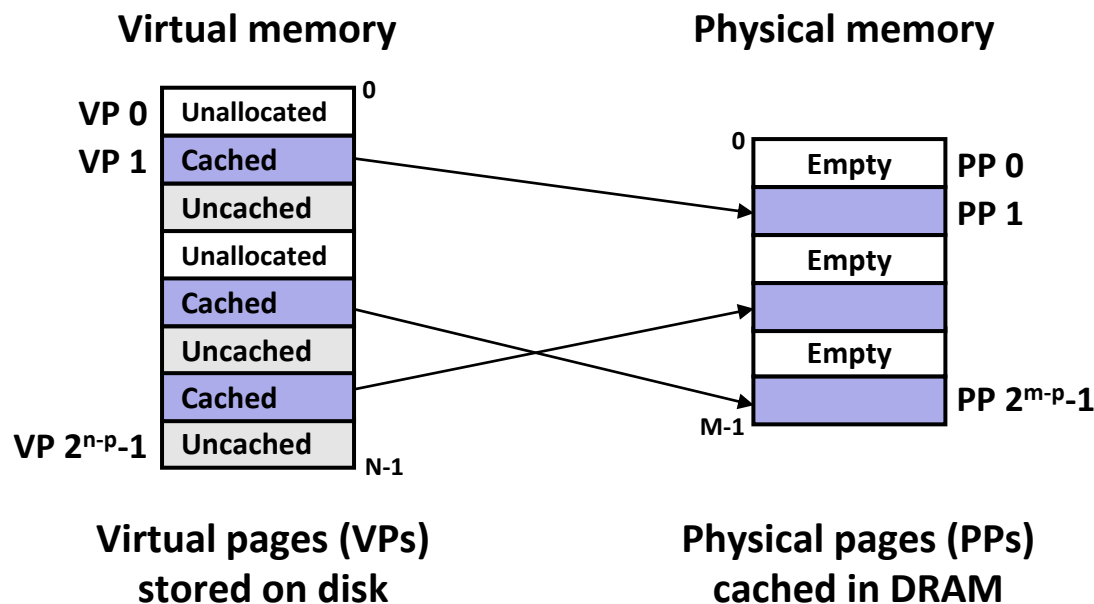
VM as a Tool for Memory Protection

- Extend PTEs with permission bits
- MMU checks these bits on each access



VM as a Tool for Caching

- Conceptually, **virtual memory** is an array of N contiguous bytes stored on disk.
- The contents of the array on disk are cached in **physical memory (DRAM cache)**
 - These cache blocks are called *pages* (size is $P = 2^p$ bytes)



Remember: Set Associative Cache

$E = 2$: Two lines per set

Assume: cache block size 8 bytes

Address :

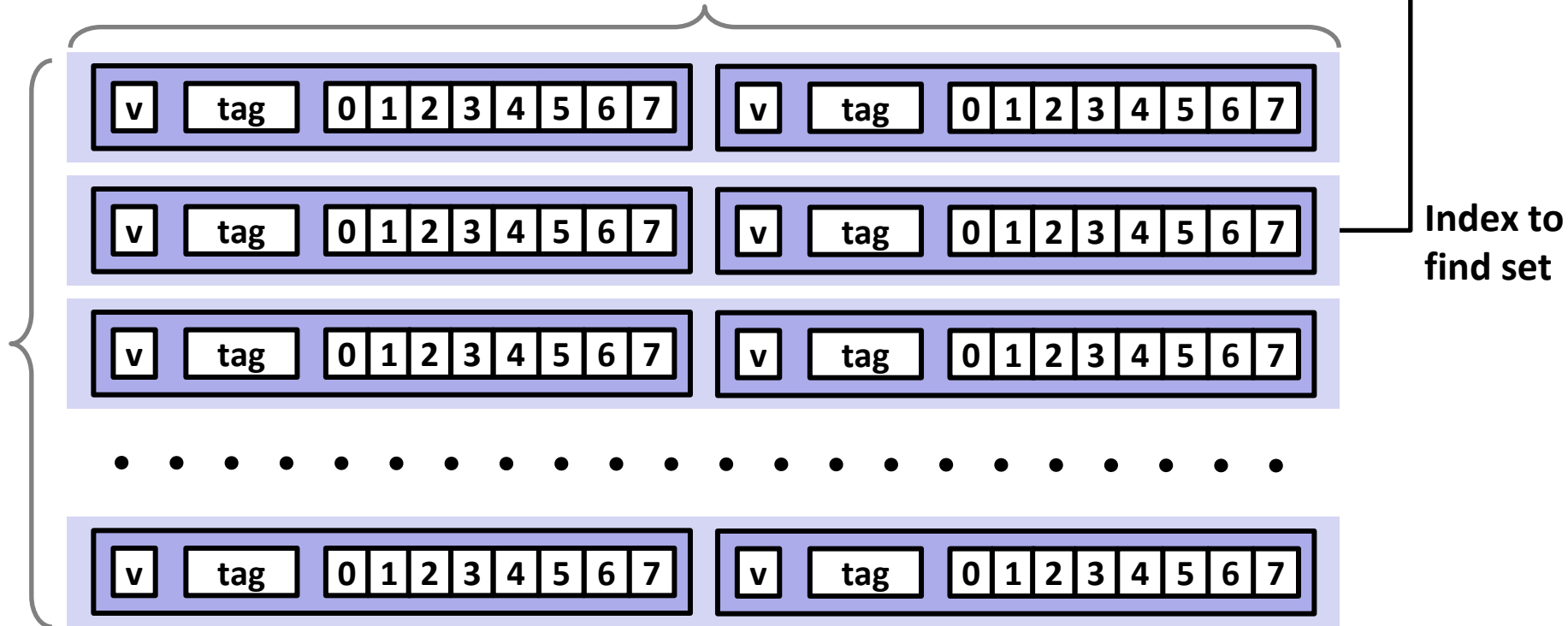
Block
offset

2 lines per set

t bits

0...01

100



S sets

DRAM Cache Organization

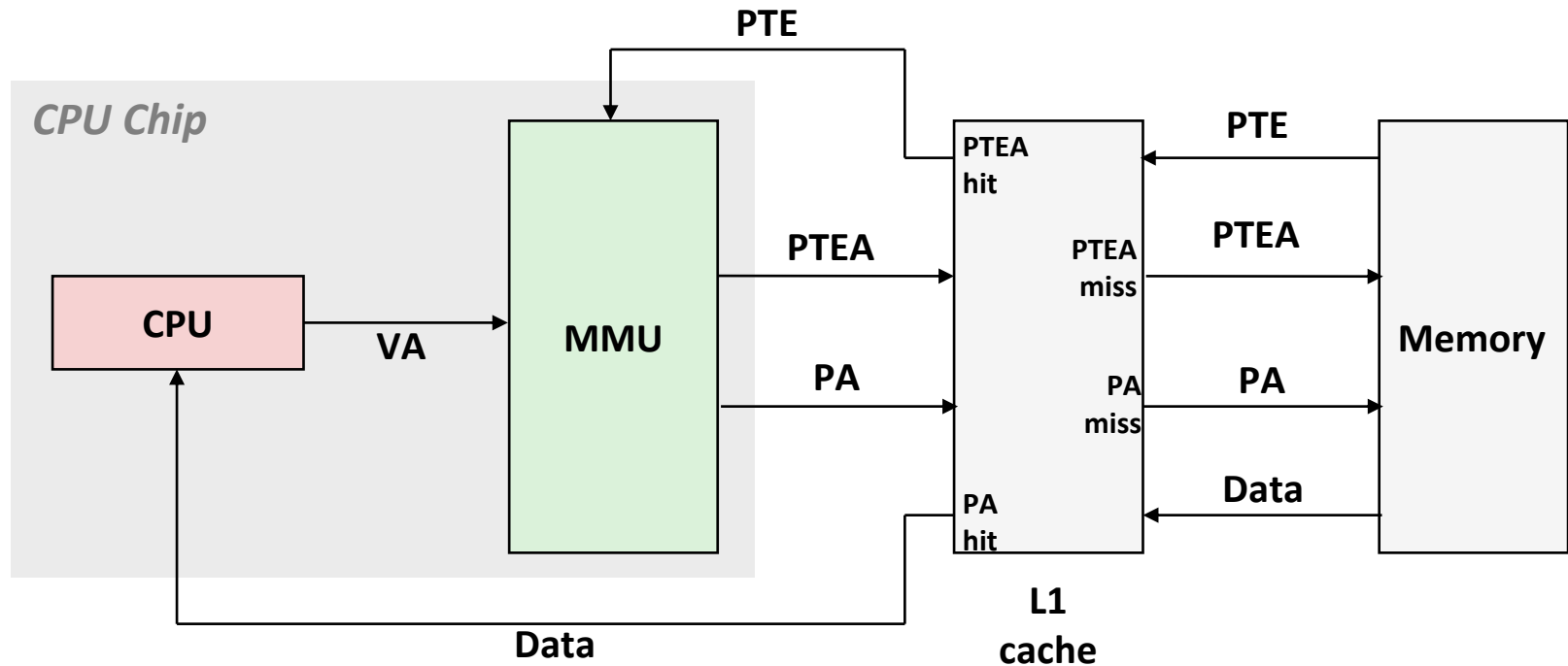
■ DRAM cache organization driven by the enormous miss penalty

- DRAM is about **10x** slower than SRAM
- Disk is about **10,000x** slower than DRAM

■ Consequences

- Large page (block) size: typically 4 KB, sometimes 4 MB
- Fully associative
 - Any VP can be placed in any PP
 - Requires a “large” mapping function – different from cache memories
- Highly sophisticated, expensive replacement algorithms
 - Too complicated and open-ended to be implemented in hardware
- Write-back rather than write-through

Integrating VM and Cache



VA: virtual address, PA: physical address, PTE: page table entry, PTEA = PTE address

Locality to the Rescue Again!

- Virtual memory seems terribly inefficient, but it works because of locality.
- At any point in time, programs tend to access a set of active virtual pages called the *working set*
 - Programs with better temporal locality will have smaller working sets
- If (working set size < main memory size)
 - Good performance for one process after compulsory misses
- If (SUM(working set sizes) > main memory size)
 - *Thrashing*: Performance meltdown where pages are swapped (copied) in and out continuously

Speeding up Translation with a TLB

■ Page table entries (PTEs) are cached in L1 like any other memory word

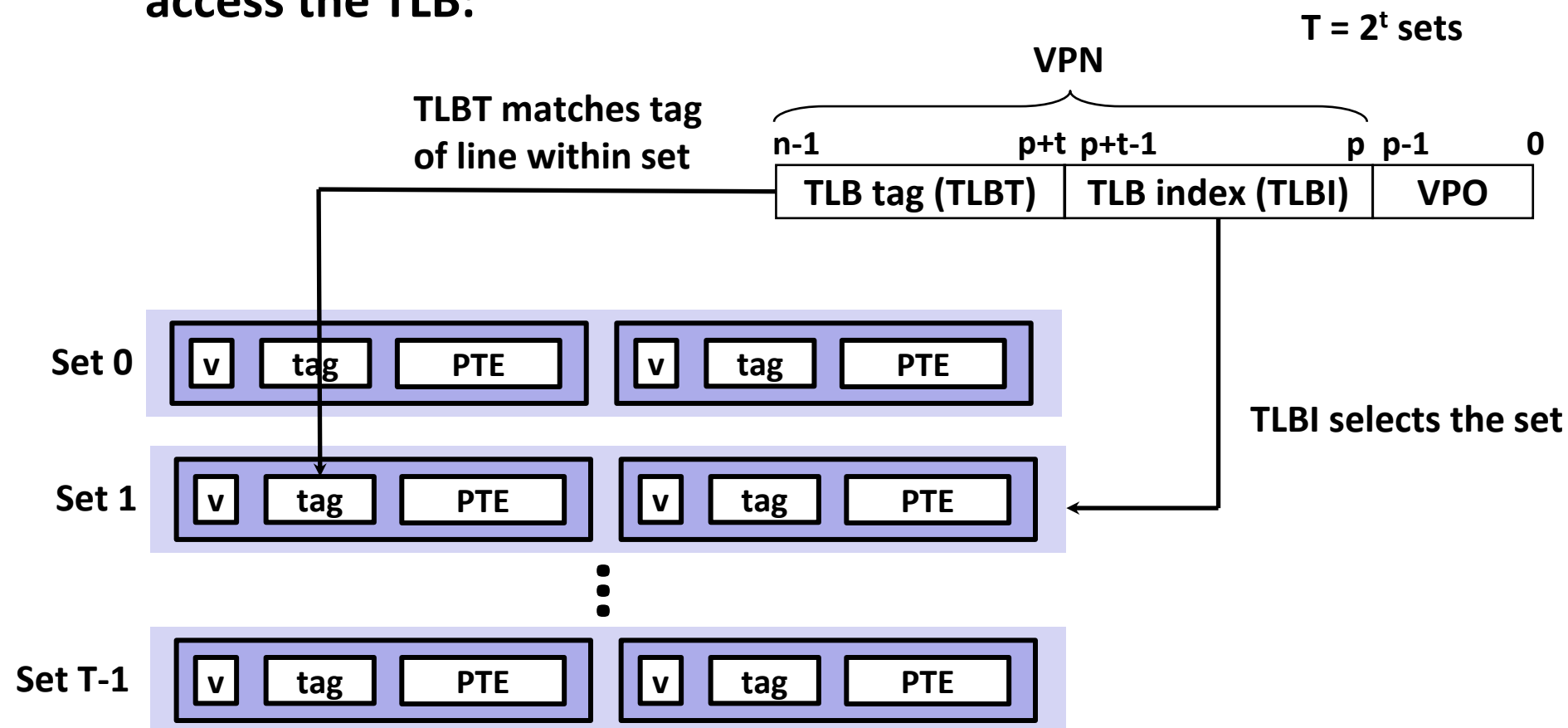
- PTEs may be evicted by other data references
- PTE hit still requires a small L1 delay

■ Solution: *Translation Lookaside Buffer* (TLB)

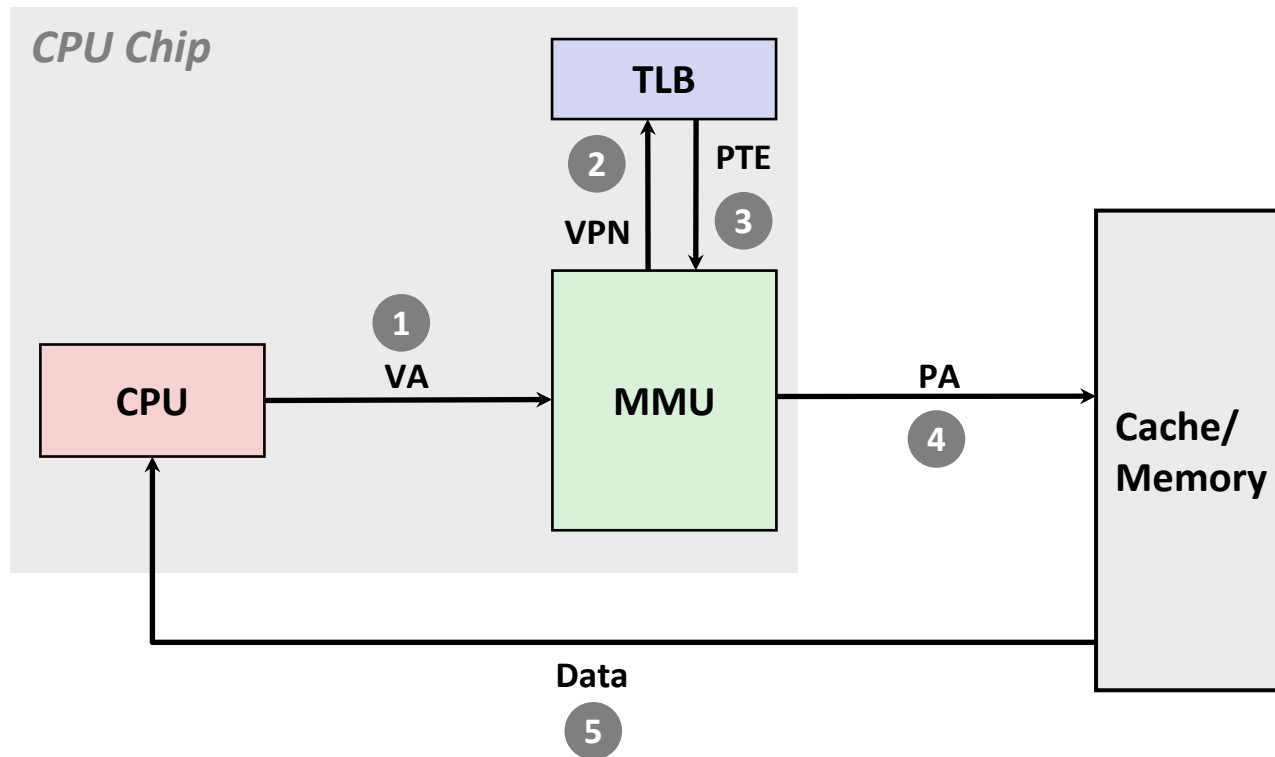
- Small set-associative hardware cache in MMU
- Maps virtual page numbers to physical page numbers
- Contains complete page table entries for small number of pages

Accessing the TLB

- MMU uses the VPN portion of the virtual address to access the TLB:

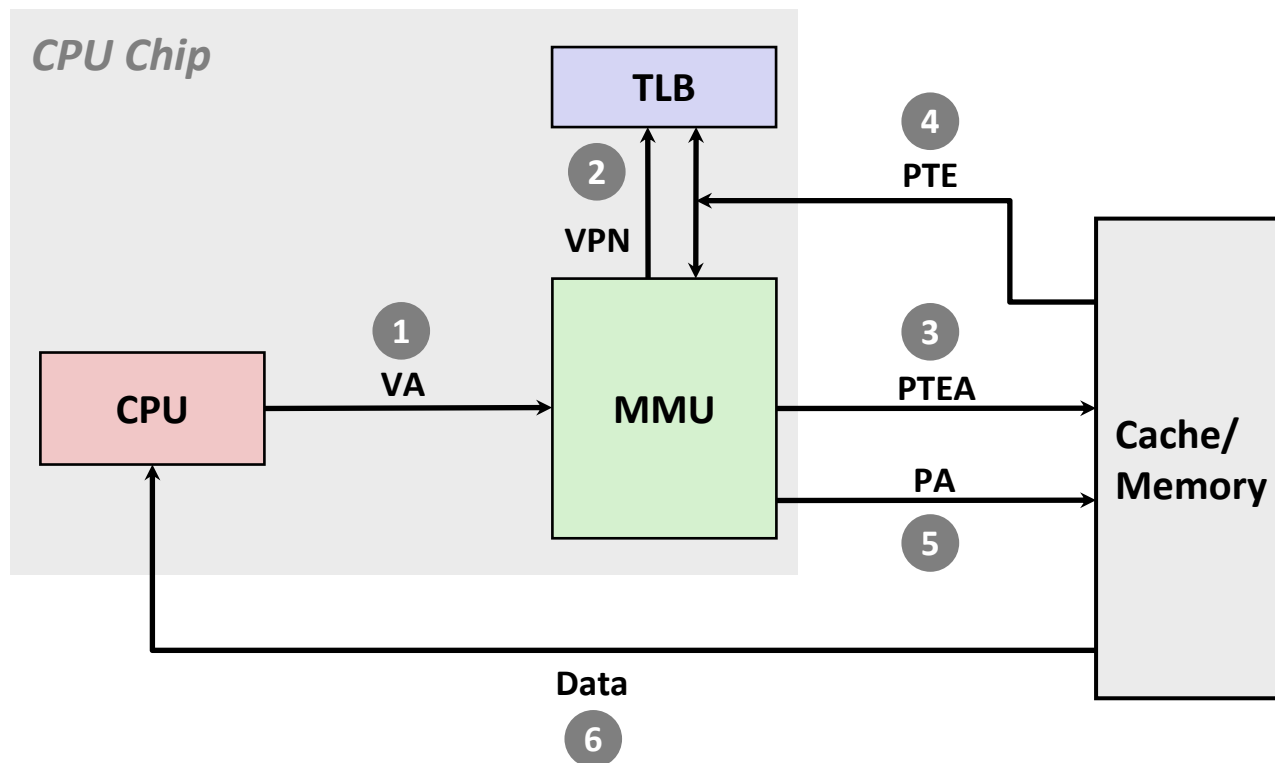


TLB Hit



A TLB hit eliminates a memory access

TLB Miss



A TLB miss incurs an additional memory access (the PTE)

Fortunately, TLB misses are rare. Why?

Summary of Address Translation Symbols

■ Basic Parameters

- $N = 2^n$: Number of addresses in virtual address space
- $M = 2^m$: Number of addresses in physical address space
- $P = 2^p$: Page size (bytes)

■ Components of the virtual address (VA)

- **TLBI**: TLB index
- **TLBT**: TLB tag
- **VPO**: Virtual page offset
- **VPN**: Virtual page number

■ Components of the physical address (PA)

- **PPO**: Physical page offset (same as VPO)
- **PPN**: Physical page number

Multi-Level Page Tables

■ Suppose:

- 4KB (2^{12}) page size, 48-bit address space, 8-byte PTE

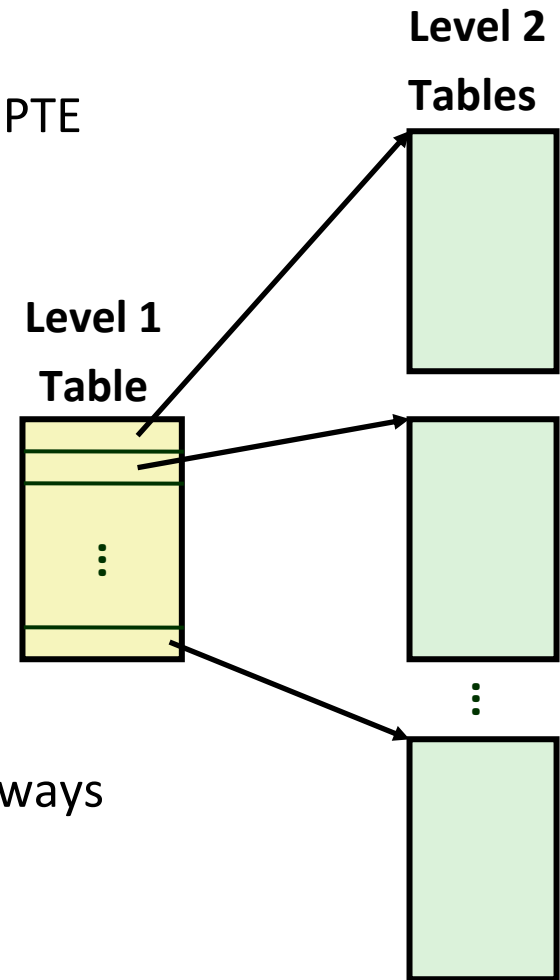
■ Problem:

- Would need a 512 GB page table!
 - $2^{48} * 2^{-12} * 2^3 = 2^{39}$ bytes

■ Common solution: Multi-level page table

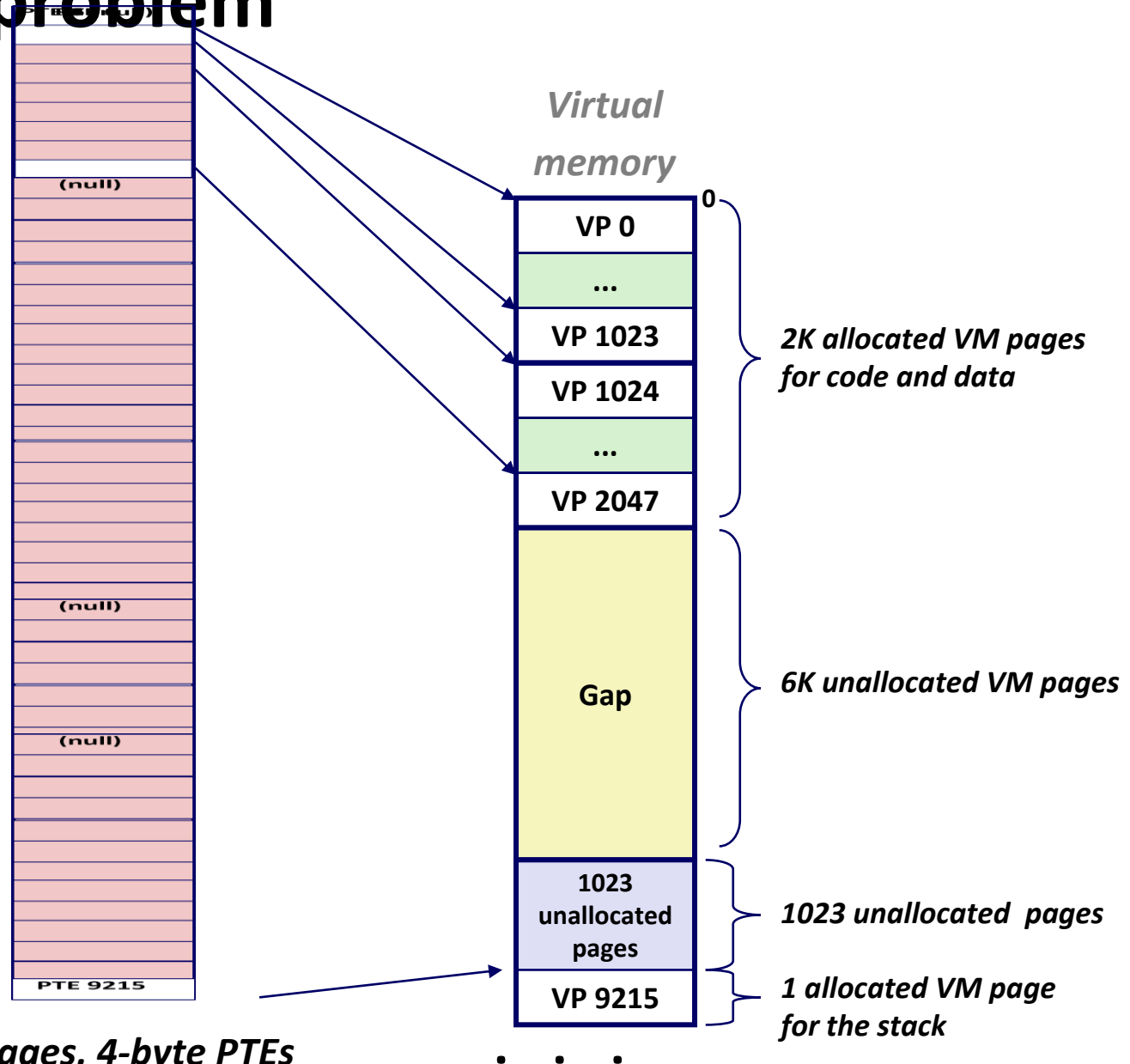
■ Example: 2-level page table

- Level 1 table: each PTE points to a page table (always memory resident)
- Level 2 table: each PTE points to a page (paged in and out like any other data)



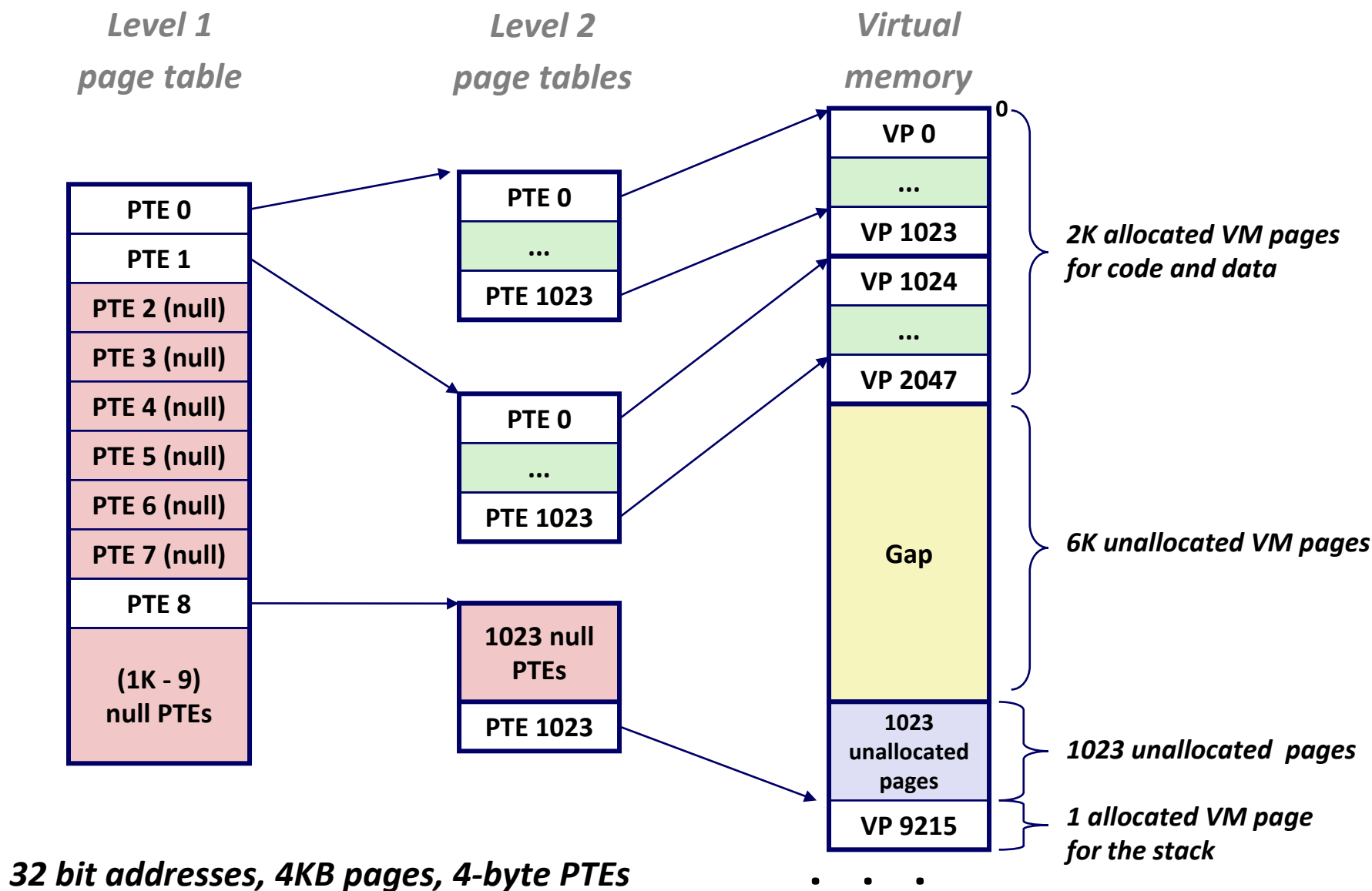
We have a problem

2^{20} Entries of
4 bytes each

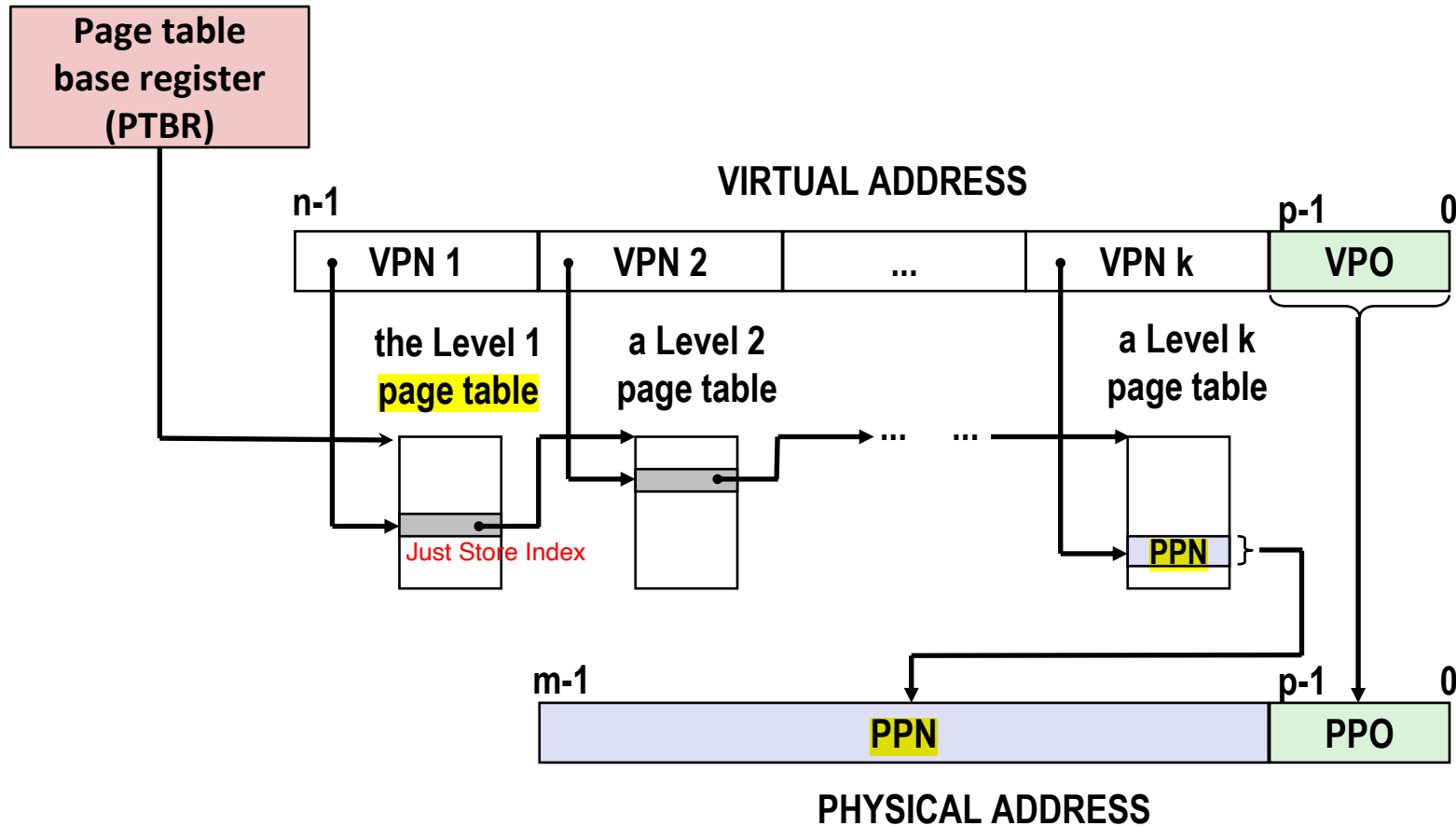


32 bit addresses, 4KB pages, 4-byte PTEs

A Two-Level Page Table Hierarchy



Translating with a k-level Page Table



Summary

■ Programmer's view of virtual memory

- Each process has its own private linear address space
- Cannot be corrupted by other processes

■ System view of virtual memory

- Uses memory efficiently by caching virtual memory pages
 - Efficient only because of locality
- Simplifies memory management and programming
- Simplifies protection by providing a convenient interpositioning point to check permissions