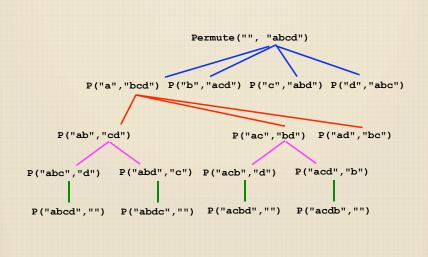
Admin

- ♦ Today's topics
 - Recursive backtracking
- Reading
 - Reader ch. 6 (today)
 - Next: pointers! 2.2-2.3, linked lists 9.5(sort of)
- ♦ Terman cafe after class

Lecture #10

Refresh: permute code

Tree of recursive calls



Subsets

- Enumerate all subsets of input
 - "abc" has subsets "a", "b", "ab", "ac", ...
 - Order doesn't matter, "ab" is same as "ba"
- Solving recursively
 - Separate one element from input
 - Can either include in current subset or not
 - Recursively form subsets including it
 - Recursively form subsets not including it
 - What is the base case?
- ◇ Remind you of any other problem you've seen?
 - Same patterns often resurface!

Subset strategy

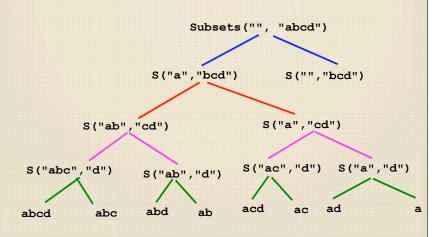
- ♦ Result is empty, starting input is "abcd"
- ♦ Consider first element: "a"
- ♦ Add to subset, remaining input is "bcd"
- Recursively find all subsets from here
- ♦ Repeat recursion without including "a"

Subsets code

```
void RecSubsets(string soFar, string rest)
{
   if (rest == "")
        cout << soFar << endl;
   else {
        // add to subset, remove from rest, recur
        RecSubsets(soFar + rest[0], rest.substr(1));
        // don't add to subset, remove from rest,recur
        RecSubsets(soFar, rest.substr(1));
   }
}

void ListSubsets(string str)
{
   RecSubsets("", str);
}</pre>
```

Tree of recursive calls



Exhaustive recursion

- ♦ Permutations/subsets are about choice
 - Both have deep/wide tree of recursive calls
 - Depth represents total number of decisions made
 - Width of branching represents number of available options per decision
- Exhaustive recursion is, well, exhaustive
 - Explores every possible option at every decision point
 - Typically very expensive
 - N! permutations, 2^N subsets
 - (Recursion isn't the problem, there just is a huge space to explore)
- Consider partial exploration of exhaustive space
 - Similar exhaustive structure, but stop at first "satisfactory" outcome

Recursive backtracking

- Cast problem in terms of decision points
 - Identify what decisions need to be made
 - Identify what options are available for each decision
 - A recursive call makes one decision, and recurs on remaining decisions
- Backtracking approach
 - Design recursion function to return success/failure
 - At each call, choose one option and go with it
 - Recursively proceed and see what happens
 - If it works out, great, otherwise unmake choice and try again
 - If no option worked, return fail result which triggers backtracking (i.e. un-making earlier decisions)
- Heuristics may help efficiency
 - Eliminate dead ends early by pruning
 - Pursue most likely choice(s) first

Backtracking pseudocode

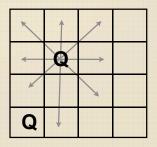
Permute -> anagram finder

```
void RecPermute(string soFar, string rest)
{
   if (rest == "") {
      cout << soFar << endl;
} else {
      for (int i = 0; i < rest.length(); i++) {
            RecPermute(soFar+rest[i], rest.substr(0,i)+rest.substr(i+1));
      }
}

bool IsAnagram(string soFar, string rest, Lexicon &lex)
{
   if (rest == "") {
      return lex.containsWord(soFar);
} else {
      for (int i = 0; i < rest.length(); i++) {
        if (IsAnagram(soFar+rest[i], rest.substr(0,i)+rest.substr(i+1), lex))
            return true;
      }
   }
   return false;
}</pre>
```

8 Queens

- ♦ Goal: place N queens on board so none threatened
 - Queen can attack in any straight line (horizontally, vertically, diagonally)
- ♦ Cast as in terms of decision
 - · Each call will make one decision and recur on rest
 - How many decisions do you have to make?
 - What options do you have for each?



N queens code

```
bool Solve(Grid<bool> &board, int col)
{
   if (col >= board.numCols()) return true; // base case

   for (int rowToTry=0; rowToTry<board.numRows(); rowToTry++) {
      if (IsSafe(board,rowToTry,col)) {
            PlaceQueen(board, rowToTry, col);
            if (Solve(board, col + 1)) return true;
            RemoveQueen(board, rowToTry, col);
      }
   }
   return false;
}</pre>
```