## Homework 1

## Salim M'jahad msm2243

1 - list of functions from the lowest growth rate to the highest:

- 2/N
- 37
- $\sqrt{N}$
- N
- N\*log(log(N))
- $N * log(N) \equiv N * log(N^2)$
- $N * log^2(N)$
- $N^{1.5}$
- N<sup>2</sup>
- $N^2 * log(N)$
- $\bullet$   $N^3$
- 2 <sup>N/2</sup>
- 2 <sup>N</sup>

**2 - a -** the sequence in dollars for the fine on day N is  $v_N = 2^{2^{N-1}}$ 

**b** - to reach D dollar we do the inverse which is log base 2 of log base 2  $N = log_2(log_2(D))$ 

3 - a - for this program fraction an addition and an assignment of value will be done for n times, and all that will be done for 23 times. Thus, it will take 2\*N\*23 TU (Time Units). O = N

**b** - for this program fraction an addition and an assignment of value will be done for -

 $\sum_{i=1}^{N} i \approx N^2/2$  times. As seen in the book, we consider the worst case for the Big-Oh. O =  $N^2$ 

**c** - this program fraction uses recurrence. Every time the method is called the number is split into k parts and the program returns 1 when  $n \le k$ . The method is called around  $log_k N + constant$  thus O = log(N)

- **4 a.** 2.5 ms
  - **b.**  $2.5*\log(500)/\log(100) = 3.37$
  - **c.** 12.5 ms
  - **d.** 62.5 ms

**5** - since the  $A_i$  are ordered, the array is sorted. We can use binary search to find i (which will be the position of  $A_i+1$ ). The running time of the algorithm is  $O(\log N)$