

All about CP

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1 Development

This book is publicly developed on [GitHub](#). If you find anything confusing, or you think that there is a better way to express the idea, please make a pull request.

Contribution

- Avoid comments: Name your variables and functions elaborately, instead
- Set indentation to two spaces
- Keep it concise

Specific to C++:

- Make sure your code supports C++17 standard
- Use `ll` instead of `long long`
- Assume `bits/stdc++.h` is included

2 Fast I/O

2.1 C++

This will save you from TLEs in many problems

```
cin.tie(0) -> sync_with_stdio(0);
```

Caution: Either use printf/scanf, or use cin/cout with fast I/O. But don't use both

2.2 Python

```
import sys
input=lambda:sys.stdin.readline().strip()
```

3 Language

3.1 Return by reference

Lets define an array for demonstration:

```
int vals[] = {10, 12, 83, 122, 5, 34};
```

The following function returns the value of the i-th element

```
int getVal(int i) {  
    return vals[i];  
}
```

The following function returns a reference to the i-th element

```
int& getRef(int i) {  
    return vals[i];  
}
```

Demonstration:

```
int x = getVal(2);  
x++; // Doesn't change array  
cout << getVal(2) << '\n'; // Output: 83  
int y = getRef(2);  
y++; // Changes array  
cout << getVal(2) << '\n'; // Output: 84
```

Practice:

- [This code](#) causes TLE, but [this](#) gets AC

3.2 Inner Functions

4 Maths

4.1 Summation

4.1.1 Identities

$$\sum c \times f(n) = c \times \sum f(n), c \text{ is constant}$$

$$\sum (f(n) \pm g(n)) = \sum f(n) \pm \sum g(n)$$

$$\sum_{i=1}^n \sum_{j=1}^n a_i b_j = (\sum_{i=1}^n a_i) (\sum_{j=1}^n b_j)$$

Practice:

- [AtCoder ARC A - Simple Math](#)

5 String

5.1 Big Integers

1. Take input as string
2. Reverse the string

6 Constructive Algorithm

6.1 From A to B

Think out of the box: Convert from B to A instead

Practice:

- [CF1810 - Candies](#)

7 Searching

7.1 Linear Search

Linear search sequentially checks each element of a range¹ until it finds the targeted value

```
int linearSearch(int a[], int n, int target) {
    for(int i=0; i<n; i++)
        if(a[i]==target)
            return i;
    return -1; // alternatively: return n;
}
```

Worst Case Time Complexity: $O(n)$

7.1.1 STL Algorithms with Linear Search

Let's define an integer vector for demonstration:

```
vector<int> v {3, 2, 1, 5, 3, 2, 6};
```

find

Returns the first occurrence of the target in range

```
int i = find(v.begin(), v.end(), 2) - v.begin();
cout << "First occurrence: " << i << '\n';

int j = find(v.rbegin(), v.rend(), 2) - v.begin();
cout << "Last occurrence: " << n-j-1 << '\n';
```

count

Returns the number of elements that are equal to the target “cpp int cnt = count(v.begin(), v.end(), 3); // 2

Binary Search

Binary search compares the target value with the middle value of *sorted* range

Hash Table

\pagebreak

Bit Manipulation

Non-decimal Literal in C++

To express integers in number systems other than decimal, use the following prefixes

| Base | Prefix |

¹A range is any sequence of objects that can be accessed through iterators or pointers, such as an array or an instance of some of the STL containers.

Source: <https://cplusplus.com/reference/algorithm/>

	----		-----	
	bin		0b	
	hex		0x	
	oct		0	

```
```cpp
assert(13 == 0b1101);
assert(13 == 0xd);
assert(13 == 015);
```

## 7.2 How integers are stored

Integers are stored as blocks of bytes

Data Type	No. of Bytes
char	1
short	2
int	4
long long	8

## 7.3 I/O with Non-Decimal Numbers

```
int x;
cin >> hex >> x; // takes input in hex
cout << hex << x; // prints output in hex
cin >> oct >> x; // takes input in oct
cout << oct << x; // prints output in oct
cin >> dec >> x; // takes input in dec
cout << dec << x; // prints output in dec

bitset<32> b;
cin >> b; // takes 32 bit input in bin
cout << b; // prints 32 bit output in bin
cout << b.to_ulong();
```

## 7.4 Signed and Unsigned Integers

Positive integers (both signed and unsigned) are just represented with their binary digits, and negative signed numbers (which can be positive and negative) are usually represented with the Two's complement<sup>2</sup>

```
cout << bitset<3>(5) << "\n";
// Output: 101
```

<sup>2</sup>[cp-algorithms - bit manipulation](#)

```
cout << bitset<32>(-1) << "\n";
// Output: 11111111111111111111111111111111
```

## 7.5 Index of Bit

Bits in a bit string are indexed from right to left, starting with 0.

```
bit string: 1 0 1 1 0 1
index: ... 3 2 1 0
```

In this text, i-th bit means bit with index i.

## 7.6 Terminologies

Terms	Meaning
<b>Set bit</b>	Make the bit 1
<b>Unset/Clear bit</b>	Make the bit 0
<b>Flip bit</b>	Make the bit opposite
<b>Lower bit/ Higher bit</b>	i-th bit is lower than j-th bit if $i < j$ MSB is highest bit, LSB is the lowest bit

## 7.7 Thinking in Binary

**Fun fact:** Bit stands for “Binary Digit”

### 7.7.1 Position of bits

Every position in a binary number has an index as mentioned [here](#).

Each position also has a positional value:  $2^{index}$

Here is an example:

```
bit string: 0 1 1 0 1
index: 4 3 2 1 0
value: 16 8 4 2 1
```

### 7.7.2 Converting to Decimal

To find decimal representation, we have to add up the positional values of set bits

As an example, let's convert 0b1101 to decimal:

```
bit string: 1 1 0 1
index: 3 2 1 0
value: 8 4 2 1
```

Therefore, 0b1101 in decimal is -  
 $8 \times 1 + 4 \times 1 + 2 \times 0 + 1 \times 1$   
 $= 8 + 4 + 1$   
 $= 13$

### 7.7.3 Maximizing/Minimizing bitstrings

Bit string, A is greater than B if: - Length of binary representation of A is greater than that of B (ignoring leading zeros) - If their lengths are equal, then in the first position where they differ, A has 1 and B has 0.

This idea is important to solve problems related to maximizing/minimizing number with binary operations

These problems can be reduced to the following pattern: You are given a number n. Do some bitwise operations on n such that n is maximized.

The idea behind this pattern of problems is that you have to maximize the length of the bitstring (without leading zeros) and set more significant bits that are unset

## 7.8 NOT

Operation	Meaning
$\sim x$	1's complement of x
$\sim x + 1$	2's complement of x

### 7.8.1 $\sim x$ in terms of $\sim x$

```
assert($\sim x == \sim x + 1$);
```

## 7.9 XOR

### 7.9.1 How to Visualize XOR

#### Method-1: Non-equivalence Operator

$A \wedge B$  is true if truth value of A and B are different.  
The following function uses this algorithm:

```
bool xor(bool a, bool b) {
 if (a!=b)
 return true;
 return false;
}
```

## Method-2: Programmable Inverter

Think of XOR as a machine with an on/off button, that takes one bit as input and one bit as output. If the machine is on, the output bit will be inverse of input bit, otherwise, the output bit will be the same as input bit.

In  $A \oplus B$ , one bit decides if the other should be flipped.

The following function uses this algorithm:

```
bool xor(bool a, bool b) {
 if (a)
 return !b;
 return b;
}
```

## 7.10 XOR with AND OR

$A$	$B$	$A \vee B$	$A \wedge B$	$A \oplus B$
0	0	0	0	0
0	1	1	0	1
1	0	1	0	1
1	1	1	1	0

From the truth table it is clear that:

$$A \oplus B = (A \vee B) - (A \wedge B)$$

$$A \oplus B = (A \vee B) \wedge \neg(A \wedge B)$$

This is also valid for bitwise operations:

$$a \oplus b = (a | b) - (a \& b)$$

$$a \oplus b = (a | b) \& \sim(a \& b)$$

## 7.11 Thinking of Bitwise Operators: Fixing one operand

### 7.11.1 AND

Operation	Meaning
$\& 1$	same
$\& 0$	0

### 7.11.2 OR

Operation	Meaning
$  1$	1

Operation	Meaning
0	same

### 7.11.3 XOR

Operation	Meaning
^ 1	flip
^ 0	same

## 7.12 Visualizing n-1

When we subtract 1 from a number, the rightmost set bit becomes unset and all the bits to its right become set

$n = \text{xxxx}10000$

$n-1 = \text{xxxx}01111$

There for value of a binary number with all 1s of length  $n$  is  $2^n - 1$

## 7.13 Common Bit Operations and Checks

### 7.13.1 Parity Check

If  $n$  is an integer(positive or negative) then  $n\&1$  represent parity of  $n$ . It is similar to  $n\%2$  but better, because unlike  $n\%2$ ,  $n\&1$  works for both positive and negative numbers.

```
int n;
n = 5;
assert((n & 1) == 1);
assert((n % 2) == 1);
n = -5;
assert((n & 1) == 1);
assert((n % 2) == -1);
```

In general,  $x \% (1\<<k)$  is equivalent to  $x \& (1\<<k)-1$

### 7.13.2 Left Shift as $\times 2$

$x\<<y$  is equivalent to  $x \times 2^y$

```
assert((5<<2) == 20);
```

### 7.13.3 Right Shift as $\div 2$

$x \gg y$  is equivalent to  $\lfloor \frac{x}{2^y} \rfloor$

```
assert((11>>2) == 2);
```

### 7.13.4 Set i-th Bit

OR with a number with only i-th bit set

```
n = n | (1<<i)
```

### 7.13.5 Clear i-th Bit

AND with a number with only i-th bit unset

```
n = n & ~(1<<i)
```

### 7.13.6 Flip i-th Bit

We already discussed that, XOR works as a programmable inverter.

```
n = n ^ (1<<i)
```

### 7.13.7 Check i-th Bit

AND with a number with only i-th bit set

```
n & (1<<i)
```

There is a chance of making overflow of bits, if n is of type long long. In general, it is better to write it this way:

```
(n>>i) & 1
```

### 7.13.8 Unset Rightmost Set Bit

See [Visualizing n-1](#)

```
n & (n-1)
```

### 7.13.9 Check if Power of Two

n is power of two, if there is only one set bit. To check if there is only one set bit, unset the last set bit, and check if it becomes zero or greater. If it becomes zero it's a power of two

```
n!=0 && (n&(n-1))>0;
```

**Edge Case: 0**

### 7.13.10 Count Set Bits: Brian Kernighan's Algorithm

The idea is to count how many times we can unset the rightmost set bit, until we reach 0

```
int countSetBits(int n) {
 int cnt=0;
 while(n)
 n = n & (n-1), cnt++;
 return cnt;
}
```

### 7.13.11 Set Range of bits

The concept is similar to setting i-th bit, except we will use a different bit mask.

$1 \ll i$  is i 0s after one 1

$(1 \ll i) - 1$  is i 1s at the end, and rest are 0s

Now we can leftshift these 1s to fit into the range

```
// Set bits in the range [start, stop)
int setRange(int n, int start, int stop) {
 int length = stop - start;
 int mask = (((1<<length) - 1) << start);
 return n | mask;
}
```

### 7.13.12 Clear Range of bits

```
int clearRange(int n, int start, int stop) {
 int length = stop - start;
 int mask = ~(((1<<length) - 1) << start);
 return n & mask;
}
```

### 7.13.13 Flip Range of bits

```
int flipRange(int n, int start, int stop) {
 int length = stop-start;
 int mask = (((1<<length) - 1) << start);
 return n ^ mask;
}
```



## 7.14 Builtin Functions

The g++ compiler provides the following functions for counting bits:

- `__builtin_clz(x)`: the number of zeros at the beginning of the number
- `__builtin_ctz(x)`: the number of zeros at the end of the number
- `__builtin_popcount(x)`: the number of ones in the number
- `__builtin_parity(x)`: the parity (even or odd) of the number of ones

The functions can be used as follows:

```
int x = 5328; // 000000000000000000001010011010000
cout << __builtin_clz(x) << '\n'; // 19
cout << __builtin_ctz(x) << '\n'; // 4
cout << __builtin_popcount(x) << '\n'; // 5
cout << __builtin_parity(x) << '\n'; // 1
```

While the above functions only support int numbers, there are also long versions of the functions available with the suffix ll.

Source: [CSES Book](#)

## 7.15 Sum with bit operations

$a \& b$  represents the carry bit and  $a \wedge b$  represents the last digit of  $a+b$ . Therefore the following formula holds:

$$a + b = ((a \& b) \ll 1) + (a \wedge b)$$

Since,  $a \wedge b$  can be written in terms of AND, OR the sum can also be written in terms of AND, OR

$$\begin{aligned} a + b &= ((a \& b) \ll 1) + (a \wedge b) \\ &= 2 * (a \& b) + (a \mid b) - (a \& b) \\ &= (a \& b) + (a \mid b) \end{aligned}$$

Practice: [CF](#)

## 8 Ranges

### 8.1 Multiple Ranges

#### 8.1.1 Intersections of ranges

There are  $n$  ranges  $[l_1, r_1], [l_2, r_2], [l_3, r_3], \dots [l_n, r_n]$   
Now, the intersections of these ranges is  $[L, R]$ , where  
$$L = \max_{i=1}^n l_i$$
$$R = \min_{i=1}^n r_i$$

**Edge Case:**

If  $R < L$ , then the intersection is an empty range.

**Practice:**

- [CF1282A - Temporarily unavailable](#)
- [CF124A - The number of positions](#)
- [CF714A - Meeting of Old Friends](#)

## 9 Interactive Problem

Make a function called `ask` that takes the parameters of the question as parameter and returns the return value of the question

Example:

```
int ask(string s, int a, int b) {
 cout << s << ' ' << a << ' ' << b << '\n';
 cout << flush;
 int res;
 cin >> res;
 return res;
}
```

## 10 Big Integers

If  $N$  is a very large number (containing more than **40** digits), then  $N$  **should be read as string**.

### 10.1 Divisibility of Integer $N$

$N$  is

1. Always divisible by 1.
2. Divisible by 2 if the last digit of  $N$  is divisible by 2 i.e., last digit is even.
3. Divisible by 3 if sum of digits is divisible by 3.
4. Divisible by 4 if the number containing only the last two digits of  $N$  is divisible by 4.
5. Divisible by 5 if last digit is 0 or 5.
6. Divisible by 6 if it is divisible by both 2 and three i.e., last digit is even and the sum of all digits is divisible by 3.

## 11 Number Theory

### 11.1 Binary Exponentiation

$$x^n = \begin{cases} 1, & n = 0 \\ x^{(n/2)} \cdot x^{(n/2)}, & n \bmod 2 = 0 \\ x^{(n-1)} \cdot x, & n \bmod 2 = 1 \end{cases}$$

**C++ Code:**

```
ll binPow(ll x, ll n) {
 if (n == 0) return 1;
 ll z = modpow(x, n/2);
 z *= z;
 if (n & 1) return x * z;
 return z;
}
```

**Complexity:**  $O(\log n)$

### 11.2 Modular Arithmetic

#### 11.2.1 Basic Modular Operations

**Modular Addition:**

$$(a + b) \bmod m = ((a \bmod m) + (b \bmod m)) \bmod m$$

**Modular Multiplication**

$$(a \times b) \bmod m = ((a \bmod m) \times (b \bmod m)) \bmod m$$

#### 11.2.2 Modular Exponentiation

**Calculate  $x^n \% m$  :**

$$x^n \bmod m = \begin{cases} 1 \bmod m, & n = 0 \\ ((x^{n/2} \bmod m) \cdot (x^{n/2} \bmod m)) \bmod m, & n \bmod 2 = 0 \\ ((x^{n/2} \bmod m) \cdot (x \bmod m)) \bmod m, & n \bmod 2 = 1 \end{cases}$$

```
ll modPow(ll x, ll n, ll m) {
 if(n == 0) return 1 % m;
 ll z = modpow(x, n/2, m) % m;
 z = (z * z) % m;
 if (n & 1) return ((x % m) * z) % m;
 return z;
}
```

**Complexity:**  $O(\log n)$

## 12 Regular Expression

### 12.1 Matching Substring

Calculate how many times a certain pattern appears in a string(with duplicates).

#### Problem-1:

The pattern starts and ends with 1, and there are one or more 0s in-between.

**C++ code:**

```
#include <bits/stdc++.h>
#include <regex>
using namespace std;
int countSubstrings(const string &s) {
 regex pattern("(?=(10+1))");
 sregex_iterator iter(s.begin(), s.end(), pattern);
 sregex_iterator end;
 int count = 0;
 while (iter != end) {
 ++count;
 ++iter;
 }
 return count;
}
```

**Python Code:**

```
def countSubstrings(s):
 pattern = r'(?=(10+1))'
 matches = finditer(pattern, s)
 count = 0
 for match in matches: count+=1
 return count
```

Now, if we want only the unique substrings we can store those substrings in a set.

**C++ Code:**

```
set<string> uniqueSubstrings;
while (iter != end){
 uniqueSubstrings.insert((*iter)[1]);
 ++iter;
}
```

**Python Code:**

```
uniqueSubstrings = set()
for match in matches:
 uniqueSubstrings.add(match.group(1))
```

*The size of the set **uniqueSubstrings** is the number of unique substrings that matches the pattern.*