

Common Lisp S-Expressions: Atoms and Lists

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If a function you write for this course returns a dotted list, then either your code has a bug or an inappropriate argument value was passed to the function.

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For example, the (proper) list

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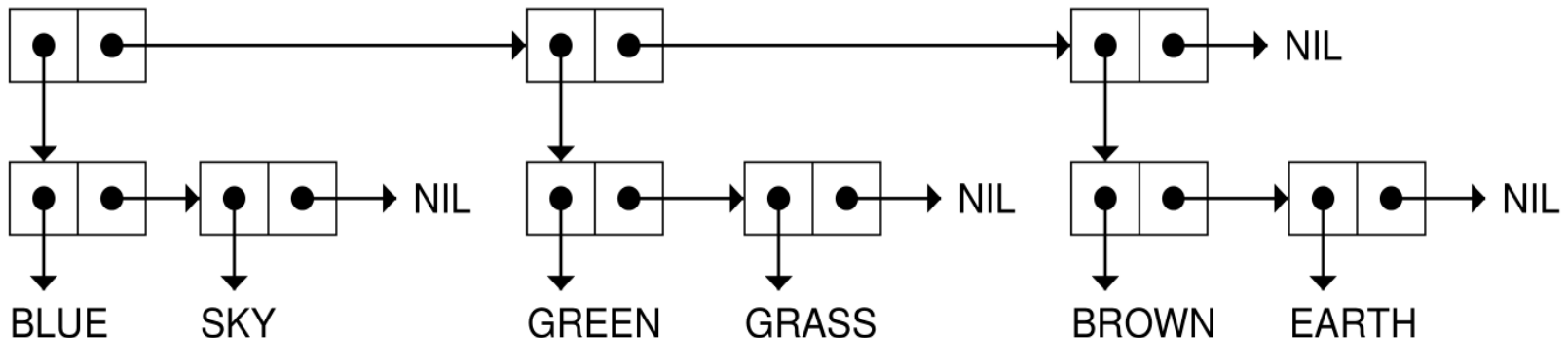
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- If a Lisp variable has a value, then the value will usually be an S-expression.
- More generally, if a Lisp expression can be evaluated, then its value (i.e., the result of its evaluation) will usually be an S-expression.

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- Lisp *macros*, which are used to extend the syntax of Lisp (i.e., "define new keywords"), depend on this property.
 - Ch. 14 of Touretzky explains how to write macros, but you will not be expected to do that: All the macros you need to use will be predefined (i.e., built-in) macros.

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From the Common Lisp Hyperspec

(<http://www.lispworks.com/documentation/common-lisp.html>):

The set of *special operator names* is fixed in Common Lisp; no way is provided for the user to define a *special operator*. The next figure lists all of the Common Lisp *symbols* that have definitions as *special operators*.

<u>block</u>	<u>let*</u>	<u>return-from</u>
<u>catch</u>	<u>load-time-value</u>	<u>setq</u>
<u>eval-when</u>	<u>locally</u>	<u>symbol-macrolet</u>
<u>flet</u>	<u>macrolet</u>	<u>tagbody</u>
<u>function</u>	<u>multiple-value-call</u>	<u>the</u>
<u>go</u>	<u>multiple-value-prog1</u>	<u>throw</u>
<u>if</u>	<u>progn</u>	<u>unwind-protect</u>
<u>labels</u>	<u>prog1</u>	
<u>let</u>	<u>quote</u>	

Just a few of these
25 special operators
(probably just QUOTE,
IF, LET, LET*, and
LABELS) will be covered
later. You will **not** be
expected to know the
rest!

Figure 3-2. Common Lisp Special Operators

More on Numbers, Symbols, and Lists in Common Lisp

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 - Note that the default floating point type in Java/C++ is ***double*** (64 bits) rather than ***float*** (32 bits).

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- Numbers are one of the two kinds of Common Lisp atom that will be used extensively in this course.
- Symbols**, which we consider next, are the other kind of atom we will use extensively.

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Page 7 of Touretzky defines symbols as follows:

symbol	Any sequence of letters, digits, and permissible special characters that is not a number.
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So FOUR is a symbol, 4 is an integer, +4 is an integer, but + is a symbol. And 7-11 is also a symbol.

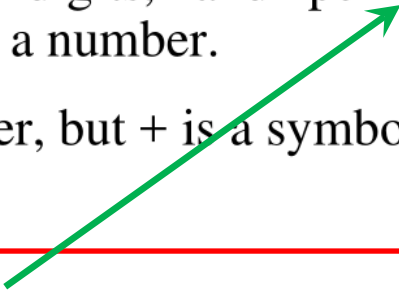
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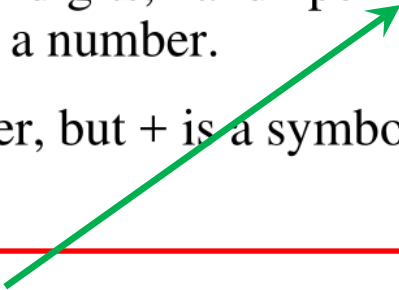
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Q. Which special characters are "permissible"?

A. Any character that is not a whitespace character and also is not one of () ' ` " , ; | \ : is permissible, but # can't be the *first* character of a symbol and a symbol can't consist only of . characters.

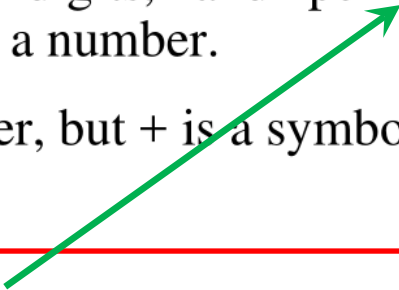
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So FOUR is a symbol, 4 is an integer, +4 is an integer, but + is a symbol. And 7-11 is also a symbol.



Q. Which special characters are "permissible"?

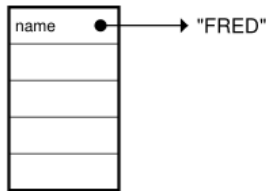
A. Any character that is not a whitespace character and also is not one of () ' ` " , ; | \ : is permissible, but # can't be the *first* character of a symbol and a symbol can't consist only of . characters.

Symbols as defined here are also called **symbol names**.

- The data object represented by a symbol name (see pp. 105-6) is called a symbol too, so use of the term *symbol name* may avoid confusion of the two concepts.

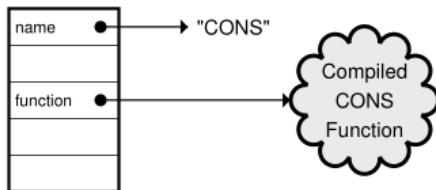
3.18 INTERNAL STRUCTURE OF SYMBOLS

So far in this book we have been drawing symbols by writing their names. But symbols in Common Lisp are actually composite objects, meaning they have several parts to them. Conceptually, a symbol is a block of five pointers, one of which points to the representation of the symbol's name. The others will be defined later. The internal structure of the symbol FRED looks like this:

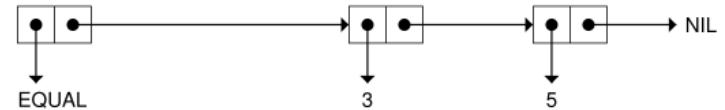


The "FRED" appearing above in quotation marks is called a **string**. Strings are sequences of characters; they will be covered more fully in Chapter 9. For now it suffices to note that strings are used to store the names of symbols; a symbol and its name are actually two different things.

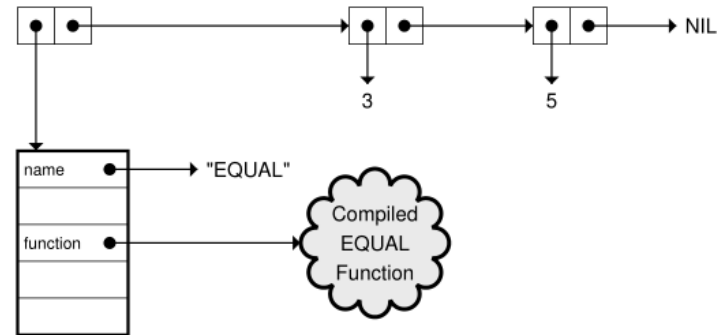
Some symbols, like CONS or +, are used to name built-in Lisp functions. The symbol CONS has a pointer in its **function cell** to a "compiled code object" that represents the machine language instructions for creating new cons cells.



When we draw Lisp expressions such as (EQUAL 3 5) as cons cell chains, we usually write just the name of the symbol instead of showing its internal structure:



But if we choose we can show more detail, in which case the expression (EQUAL 3 5) looks like this:



We can extract the various components of a symbol using built-in Common Lisp functions like SYMBOL-NAME and SYMBOL-FUNCTION. The following dialog illustrates this; you'll see something slightly different if you try it on your computer, but the basic idea is the same.

```
> (symbol-name 'equal)
"EQUAL"

> (symbol-function 'equal)
#<Compiled EQUAL function {60463B0}>
```

Note: Symbols are memory unique; see sec. 6.13.

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- Symbol names are very often used like Java/C++ **enum** constants, but they don't need to be declared.
- The value of a variable/parameter or of any other Lisp expression may well be a symbol.

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Comment: Escape characters can also be used to create symbols with names that would otherwise not be allowed. But we will not use such symbols in this course.

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- NIL is also used to mean **false** in Common Lisp.
- **T** is a constant that is the *usual* way to represent **true**, though any S-expression other than NIL can also be used to represent **true**.
- T and NIL evaluate to themselves:
The value of T is always T itself;
the value of NIL is always NIL itself.
The values of T and NIL can't be changed, and you **can't** use T or NIL as a variable / formal parameter!

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(1) proper lists

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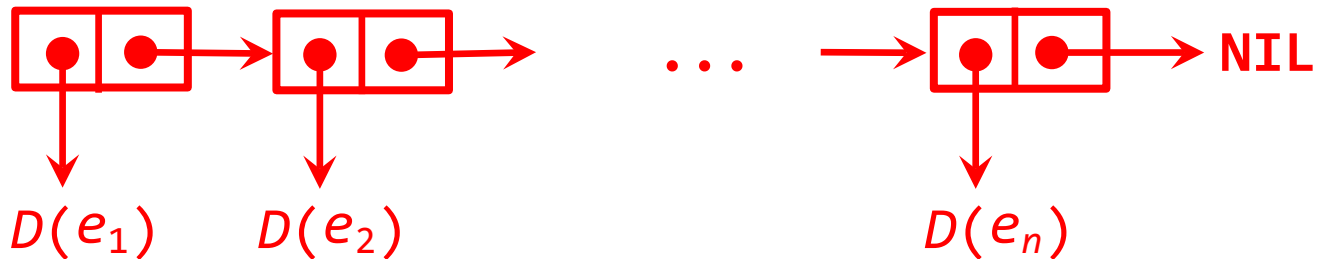
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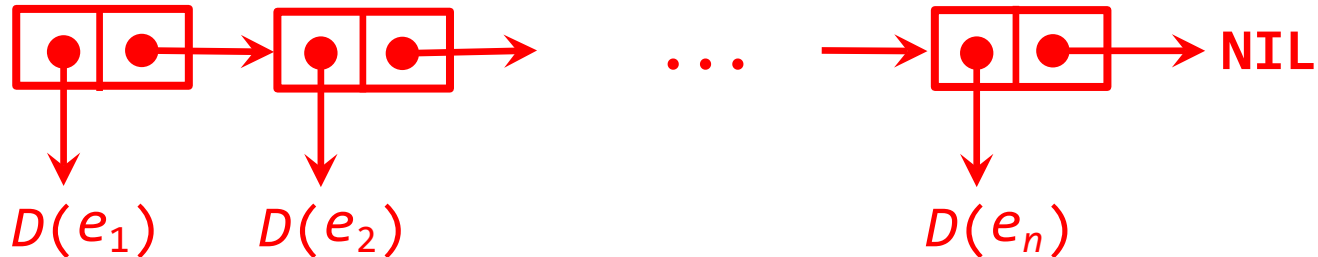
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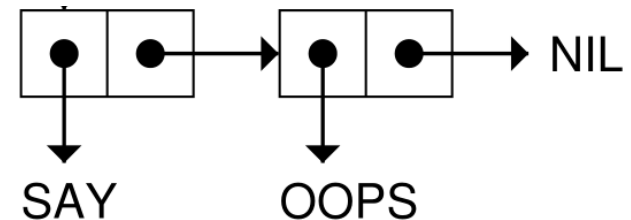
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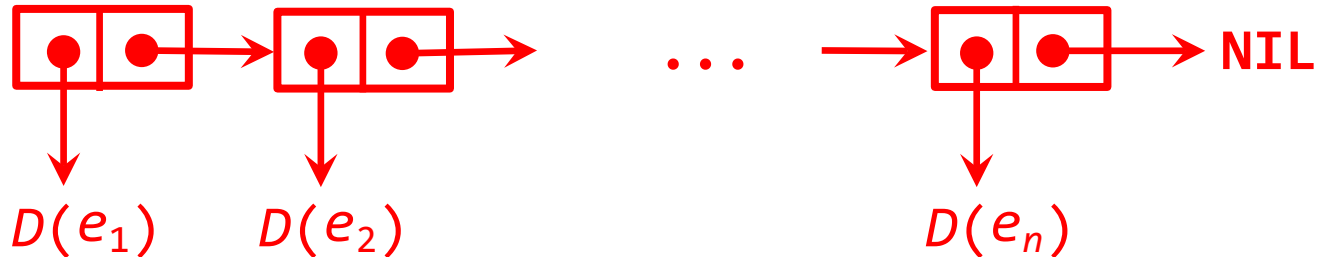
Example from p. 34 of Touretzky.

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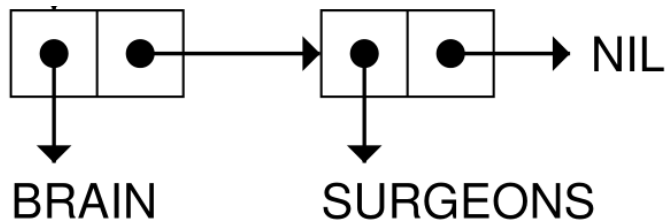


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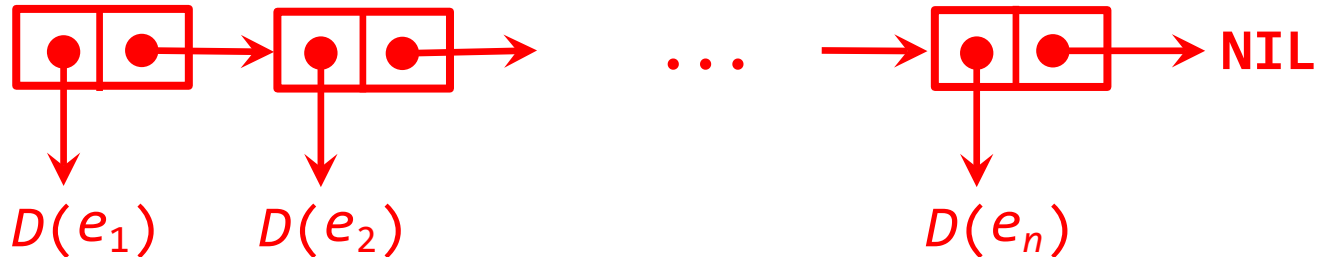
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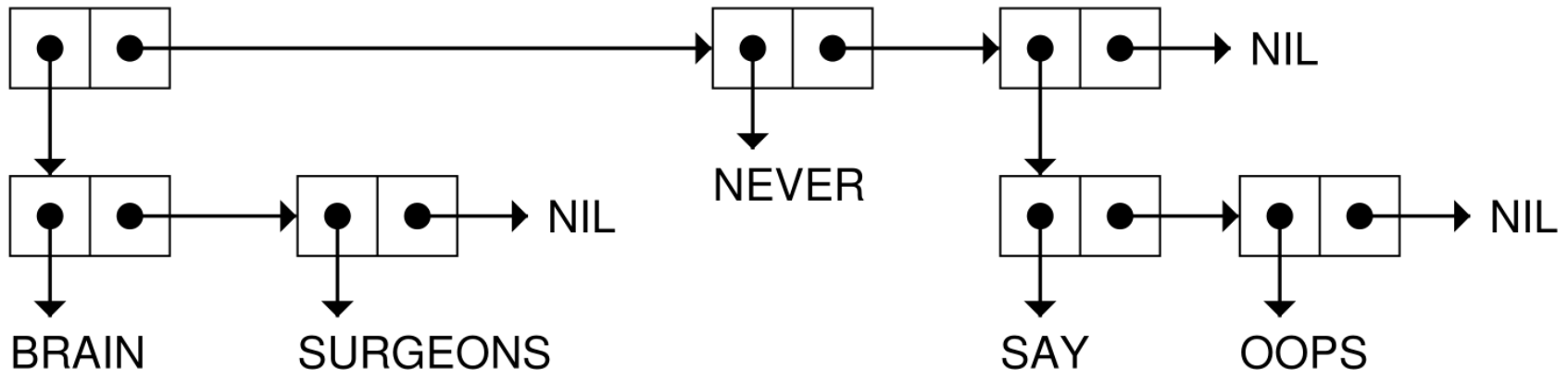
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Examples from pp. 72-3 of Touretzky

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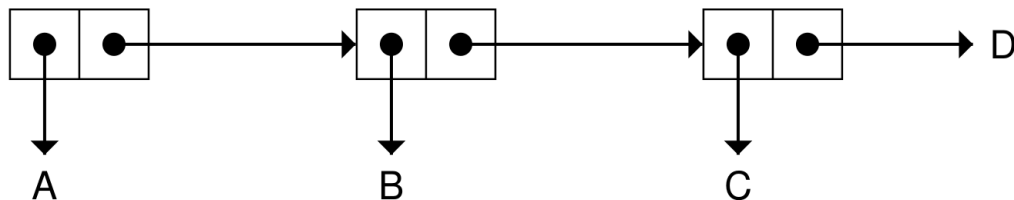
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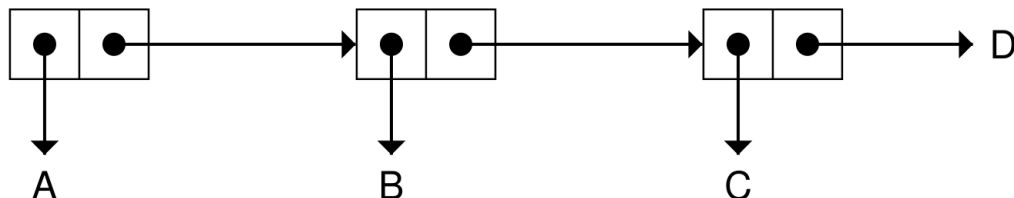
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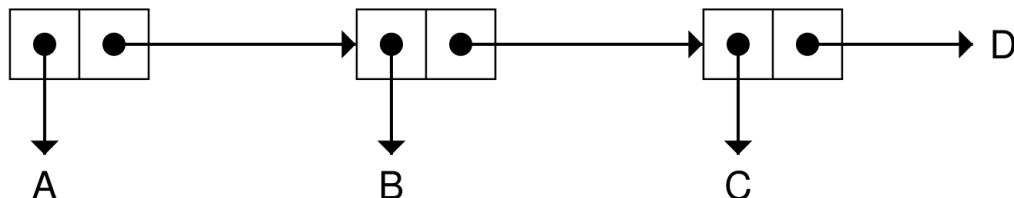
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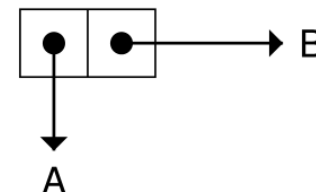
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$$(e_1 \dots e_n . \text{NIL}) =$$

$$(e_1 \dots e_n . (e_{n+1})) =$$

$$(e_1 \dots e_n . (e_{n+1} \dots e_{n+k})) =$$

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Lists

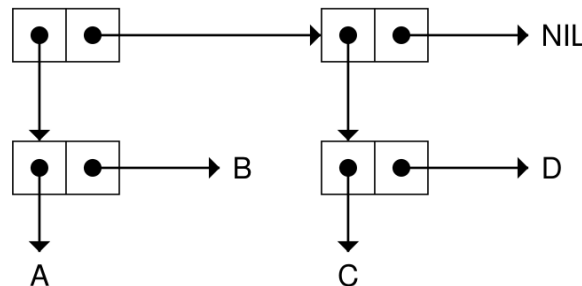
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This example is from
p. 73 and p. C-18 of Touretzky.

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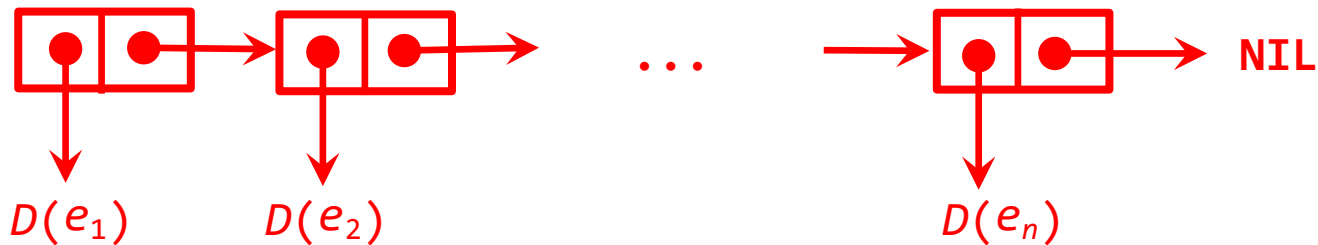


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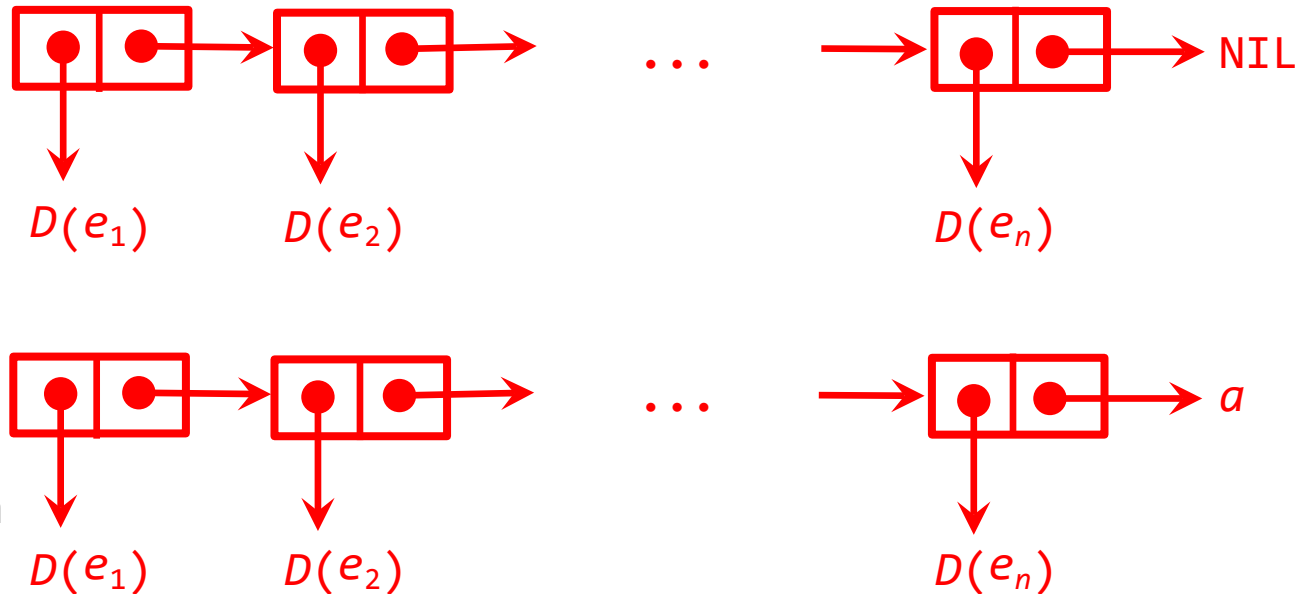
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- Dotted lists are used much less than proper lists.
- The term **List** is often used to mean **proper List**!
- If a function you write for this course returns a dotted list, then either your code has a bug or an inappropriate argument value was passed to the function!

Tree Representation of S-Expressions

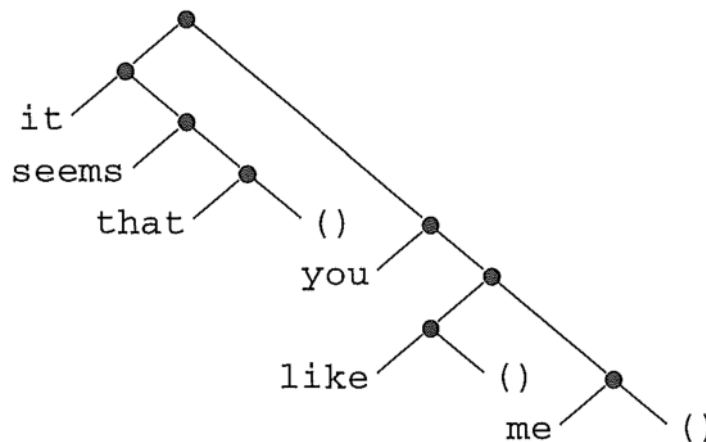
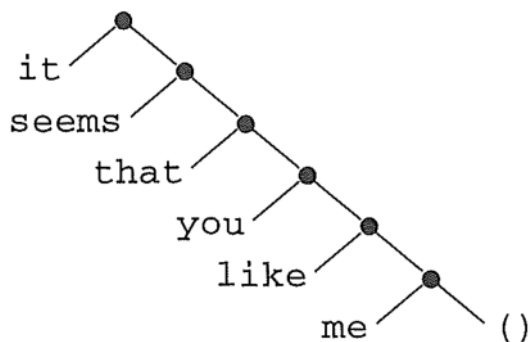
These "box and arrow" drawings are quite often simplified to binary trees in which the two children of any internal node v are *either* south and east of v *or* southwest and southeast of v :



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These two examples are from p. 393 of Sethi's book.

(it seems that you like me)

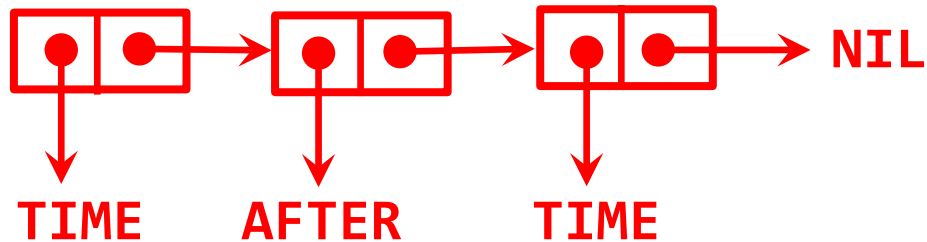
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Memory Uniqueness of Symbols

(TIME AFTER TIME) represents:

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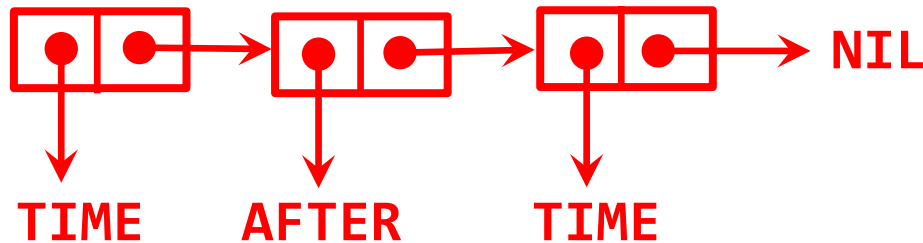
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Touretzky explains this as follows on p. 195:

In Lisp, symbols are unique, meaning there can be only one symbol in the computer's memory with a given name.^{**} Every object in the memory has a numbered location, called its **address**. Since a symbol exists in only one place in memory, symbols have unique addresses. So in the list (TIME AFTER TIME), the two occurrences of the symbol TIME must refer to the same address. There cannot be two separate symbols named TIME.

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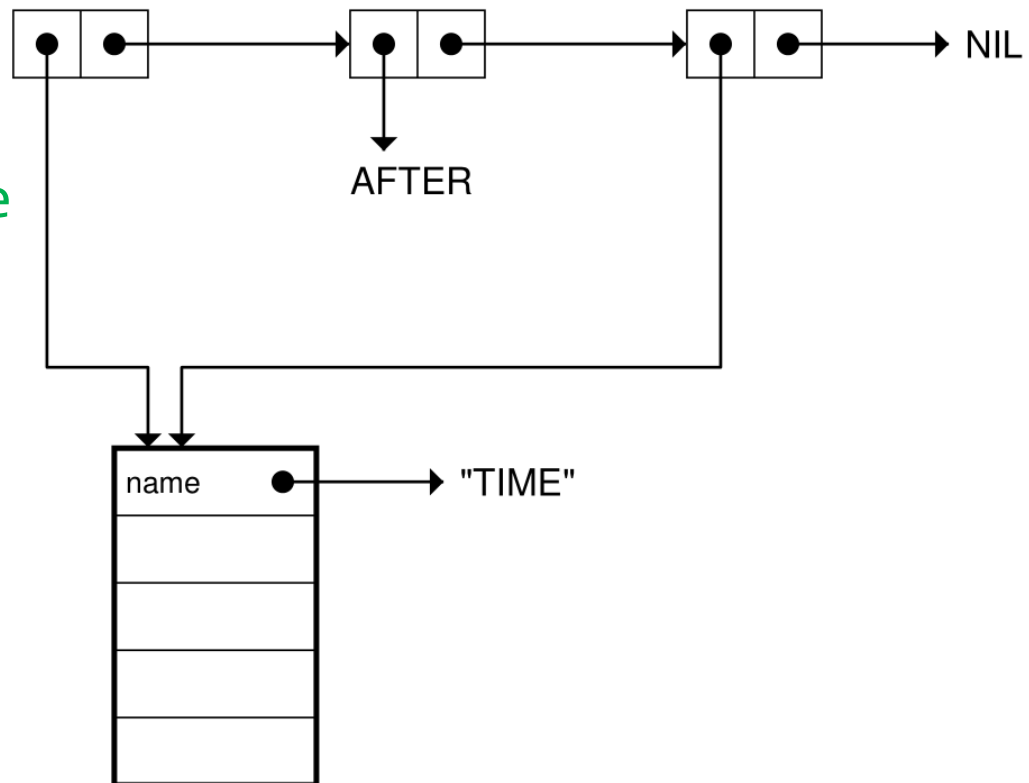
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There is *just one*
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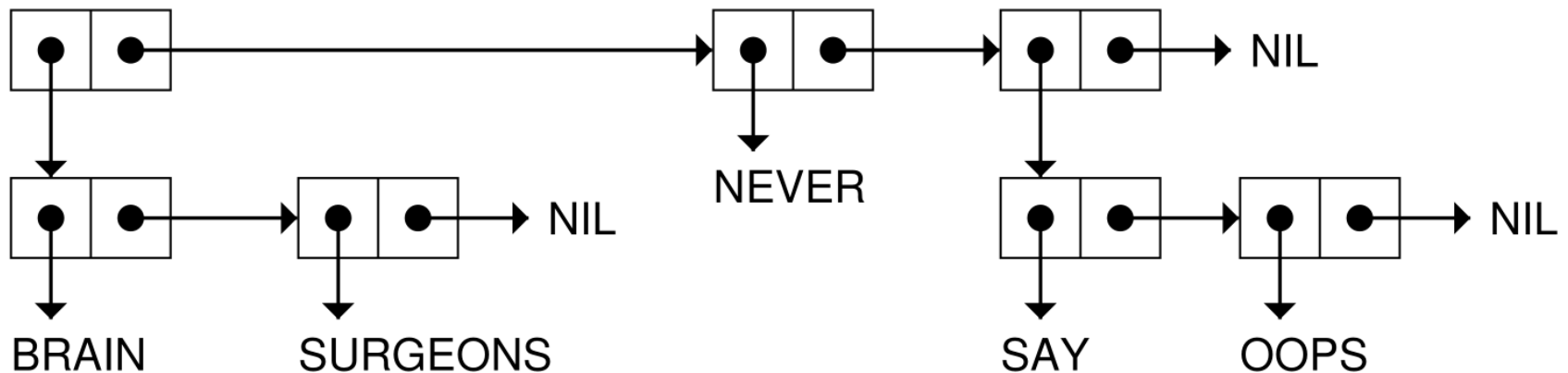
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((BRAIN SURGEONS) NEVER (SAY OOPS)) represents:



Evaluation of S-Expressions

- To evaluate an S-expression is to compute its value; evaluation is said to return the value, and the S-expression is said to evaluate to that value.

- **Example:**

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 - $(e_1 \dots e_n)$ is evaluated as a special/macro form in a way that depends on the special operator/macro.
 - **Example:** $(SETF X 9)$ is evaluated as a macro form.

How S-Expressions are Evaluated

3. A **nonempty proper list** $(e_1 \dots e_n)$ may be evaluated as a **function call** **or** as a **special/macro form** depending on e_1 , the *first element* of the list:
- If e_1 is *the name of a function* or e_1 is a *lambda expression* (a kind of list we will study later), then:
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Note: The result of evaluating a *dotted* list is undefined!