For example, suppose we want a function that computes the percent change in the price of widgets given the old and new prices as input. Our function must compute the difference between the two prices, then divide this difference by the old price to get the proportional change in price, and then multiply that by 100 to get the percent change. We can use local variables named DIFF, PROPORTION, and PERCENTAGE to hold these values.

```
(defun price-change (old new)
  (let* ((diff (- new old))
```

A common programming error is to use LET when LET* is required. Consider the following FAULTY-SIZE-RANGE function. It uses MAX and MIN to find the largest and smallest of a group of numbers. MAX and MIN are built in to Common Lisp; they both accept one or more inputs. The extra 1.0 argument to / is used to force the result to be a floating point number rather than a ratio.

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The problem is that the expression (/ BIGGEST SMALLEST 1.0) is being evaluated in a lexical context that does not include these variables. Therefore the symbol BIGGEST is interpreted as a reference to a global variable

A sequential variant of the let construct is written with keyword let*. Unlike let, which evaluates all the expressions E_1, E_2, \ldots, E_k before binding any of the variables, let* binds x_i to the value of E_i before E_{i+1} is evaluated. The syntax is

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(let ((x 2) (y x)) y); bind y before redefining x
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```
Evaluation of (h 3 4 7):

w \Rightarrow 3 \quad x \Rightarrow 4 \quad y \Rightarrow 7

the LET's local x \Rightarrow 2

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LET \Rightarrow (2+8+11)/3 = 7

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Recursive Functions

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- Our own version of f can call the supposedly already written f; but when our version is called with an argument value x, it is only allowed to call the supposedly already written f with an argument value that is <u>valid</u> for f and <u>smaller</u> in size than x.

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If $l \Rightarrow a$ proper list, then (length-of l) \Rightarrow the length of l.

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• If L ⇒ NIL, this violates the "call the supposedly already written f with an argument value that is … <u>smaller</u>" condition.

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- This definition of **length-of** is **not** circular, because when length-of calls itself it always passes an argument value that is **smaller** than the argument value it received.
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- So, for all n > 0, <u>if</u> (length-of l) returns the right result when $l \Rightarrow$ a proper list of length < n, <u>then</u> (length-of l) returns the right result when $l \Rightarrow$ a proper list of length n.

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 We've given a written explanation of a possible thought process that leads to this definition, but an experienced Lisp programmer would likely code simple definitions like this one without giving any explanation!

Example Write a function **factorial** such that: If $n \Rightarrow a$ non-negative integer, then (factorial n) $\Rightarrow n$!.

If $n \Rightarrow a$ non-negative integer, then (factorial n) $\Rightarrow n!$. Recall the following rules for writing recursive functions of 1 argument, which is a proper List or a nonnegative integer:

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If n \Rightarrow a non-negative integer, then (factorial n) \Rightarrow n!.
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Recall the following rules for writing recursive functions of 1 argument, which is a proper List or a nonnegative integer:

- When writing a recursive function f, we can first suppose a function f that correctly solves the same problem has already been written.
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• We use the fact that: For example:

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If n ⇒ 0, the above definition <u>violates</u> the "call the supposedly already written f with an argument value that is <u>valid</u> for f and <u>smaller</u> in size" condition, because (- n 1) is <u>not</u> a valid argument value for factorial if n ⇒ 0.

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- To make our function good <u>even when n ⇒ 0</u>, we add a case:
 (defun better-my-factorial (n)
 (if (zerop n)

```
1
(let ((X (factorial (- n 1))))
(* n X))))
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Example Write a function factorial such that:

If $n \Rightarrow a$ non-negative integer, then (factorial n) $\Rightarrow n!$.

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Q. How can we write factorial?
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- So, for all positive integers k, <u>if</u> (factorial i) returns the right result whenever $i \Rightarrow a$ nonnegative integer $\langle k \rangle$, <u>then</u> (factorial i) also returns the right result when $i \Rightarrow k$.

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  (defun ⇒≥≒=====factorial (n)
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• For all positive integers k, <u>if</u> (factorial i) returns
 the right result whenever i \Rightarrow a nonnegative integer < k,
 then (factorial i) also returns the right result when i \Rightarrow k.
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 We then replace the X in (* n X) with (factorial (- n 1)): (defun factorial (n) (if (zerop n); base case, where there's no recursive call 1 (* n (factorial (- n 1)))))
- As in the case of length-of, we've given a written explanation of a possible thought process that leads to this definition, but a Lisp programmer would likely code simple definitions like these without giving any explanation!

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- If the LET isn't eliminated, <u>move any case in which X needn't</u> <u>be used out of the LET</u>. If the LET <u>is</u> eliminated but <u>there's a case where the recursive call's result isn't needed, deal with such cases as base cases--i.e., without making a recursive call.</u>

Example Write a function evens such that:

If $l \Rightarrow a$ proper list of integers, then (evens $l) \Rightarrow a$ list obtained from l by omitting its odd elements.

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- If its argument value is <u>not</u> a proper list of integers, then our function *evens* may return any value whatsoever or produce an evaluation error without violating the specification!

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 - O Assignments 4 & 5 don't ask you to write such "gatekeeper" functions, but only the recursive functions themselves!

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We've seen that X is a good ____ if (oddp (car L)). To find a good ____ if (not (oddp (car L))), we try <u>another example</u>:

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- **Q.** Is there any case in which **X** is used <u>more than once</u>?
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- **Q.** Is there any case in which **X** is used <u>more than once</u>?
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- So we can <u>eliminate the LET</u> and substitute (evens (cdr L)) for each occurrence of X, to simplify the definition.

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 If l \Rightarrow a proper list of integers, then
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        (cond ((oddp (car L)) (evens (cdr L)) \cong)
                 (t (cons (car L) (evens (cdr L)) \stackrel{\succeq}{\Rightarrow}))
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- Q. Is there any case in which X is used more than once?
- A. No! X is used just once in each of the 2 cases of the cond.
- So we have <u>eliminated the LET</u> and substituted (evens (cdr L)) for each occurrence of X, to simplify the definition.

Example Write a function evens such that: If l ⇒ a proper list of integers, then (evens l) ⇒ a list obtained from l by omitting its odd elements. (defun evens (L) (if (null L) nil (cond ((oddp (car L)) (evens (cdr L)) ±) (t (cons (car L) (evens (cdr L)) ±))))

We have <u>eliminated the LET</u> and substituted (evens (cdr L))
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Example Write a function evens such that:

If L ⇒ a proper list of integers, then

(evens L) ⇒ a list obtained from L by omitting its odd elements.

(defun evens (L)

(if (null L)

nil

(cond ((oddp (car L)) (evens (cdr L)) 差)

(t (cons (car L) (evens (cdr L)) 差)))

**Touch bours of integers (cdr L))

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    We have <u>eliminated the LET</u> and substituted (evens (cdr L))

 for each occurrence of X, to simplify the definition.
• To further simplify the definition, we can replace
 (if (null L) nil (cond ... )) with (cond ((null L) nil) ... ):
    (defun evens (L)
      (cond ((null L) nil)
             ((oddp (car L)) (evens (cdr L)))
             (t (cons (car L) (evens (cdr L))))))
```