MFAI

- Rules CP or Deduction theorem ;

$(P \wedge R) \rightarrow S \Leftrightarrow P \rightarrow (R \rightarrow S)$

- $(H_1 \wedge H_2 \wedge \dots \wedge H_n \wedge P) \Longrightarrow Q$

$(H_1 \wedge H_2 \wedge \dots \wedge H_n) \Longrightarrow (P \rightarrow Q)$

Instead of this

u can do this

? $\{P \vee Q, P \rightarrow R, Q \rightarrow S\}$ $S \vee R$

$P \vee Q$
$\neg P \rightarrow Q$
$\neg R \rightarrow \neg P$
$\neg R \rightarrow Q$

$Q \rightarrow S$
$\neg R \rightarrow S$
$R \vee S$
$S \vee R$ ⊤

? $\{P \vee Q, Q \rightarrow R, P \rightarrow M, \neg M\}$

Conc $R \wedge (P \vee Q)$

$P \rightarrow M$
$\neg M$
$\neg M \rightarrow \neg P$
$\neg P$
$P \vee Q$
$\neg P \rightarrow Q$
Q has to be true

TO DO LIST;

- Take keyboard for tablet
- Install termux in tab
- Modify termux for colorscheme
- install python in it
  ↳ check if python in full power accessable in it
-

$\{P, Q \rightarrow R$
R has to be true

? $\phi \vdash P \rightarrow (Q \rightarrow R)$

$\phi \cup \{P\} \vdash Q \rightarrow R$

$\phi \cup \{P, Q\} \vdash R$

↳ where

in the place of ϕ we can use

$\{H_1, H_2, H_3 \text{ ---- } H_n\}$

$\{H_1, H_2, \text{ ---- } H_n\} \vdash C$

$H_1 \wedge H_2 \text{ - - - } \wedge H_n \Rightarrow C$

inconsistent

For all has to be false

One should be false

To make C false

In other way;

$H_1 \wedge H_2 \wedge \dots \wedge H_n \Rightarrow C \wedge \neg C$

this gives the birth of contradiction.

$H_1 \wedge H_2 \wedge \dots \wedge H_n \Rightarrow C \to C$

$H_1 \wedge H_2 \wedge \dots \wedge H_n \wedge C \Rightarrow C$

[$\because$ deduction theorem]

in consistent