

Question 1)

- a) There are 16 generators in  $Z^*_{61}$ . 2 and 6 are some of the generators.
- b)  $\langle 9 \rangle$  is a subgroup of  $Z^*_{61}$  with order 5 its generators are [9, 20, 58, 34, 1].

Question 2)

Since  $p$  and  $q$  are prime  $\phi(n) = \phi(p) * \phi(q) = p-1 * q-1$

By using the right to left algorithm to calculate the modular of a power number we calculate

$$m =$$

30256242323116471143377579036851734956599566784957808777900084159344311403348  
07724330546889390094435679786549374287072364078639802603092317776014730871103  
74524033894222588248814856928431210421802802119552503278962630326720552941119  
00141070981523217773015658586930923177732713966172972763510074059270940820416

Question 3)

- Since the gcd of  $a$  and  $n$  is 2 and  $b$  is divisible by 2 then there are 2 incongruent solutions to this congruence.  
 $x = [1632949440937904653800123918, 168551507567775226236841370368]$
- The congruence has no solutions because gcd of  $a$  and  $n$  does not divide  $b$ , so there is no solutions for this congruence.
- Since the gcd of  $a$  and  $n$  is 1 then  $a$  and  $n$  are relatively prime. Hence there is only one incongruent solution for this congruence.  
 $x = 327252728639173874206458501252$

Question 4)

For  $p_1(x)$  it generates the maximum period sequences when the length is greater than 61.  $P_1(x)$  maximum period is  $2^L - 1$  which is 31.

For  $p_2(x)$  the maximum period is  $2^L - 1$  however given the initial state it does not generate the maximum period, it only generates values smaller than 31. It might be because it is not a prime polynomial.

Question 5)

[illegible]

#### Question 6)

Since we know that the message starts with 'Dear Student', we know the first 84 bits of the plain text and when we XOR it with the cipher text we should be able to find the first 84 keys.

Now we can use Berlekamp-Massey Algorithm to find the length of the polynomial for the LFSR and the first 84 keys.

Let the length of the polynomial coefficients (not including the constant) be  $L$ .

To calculate the  $i$ th key we take the last  $L$  known keys (from  $i-L-1$  to  $i-1$ ) and we add them and take the modulus of 2.

OR we can simply take the dot product of the last  $L$  known keys with the polynomial coefficients and take the modulus of 2.

after that we will have all the keys, we can xor the keys with the cipher text to get our plain text in binary format. Then convert it into ascii to get our message:

Dear Student,

You have worked hard and that paid off:)

You have just earned 20 bonus points. Congrats!

Best,

Erkay Savas