

Question 1

I used these tools from <http://pythonfiddle.com/binary-finite-field-multiplication/> to calculate GF products

- a) The product of $x^7 + x^6 + x^4 + x^3 + x^2 + 1$ and $x^7 + x^5 + x^3 + x^2 + x$ gives $x^6 + x^5 + 1$
- b) The product of the first and second equations should be 1 for it to be the inverse.

Question 2

By following the slides I implemented a way to calculate the passwords (details on code) ['RAF!MX', 'IZSROR', '.I,BA,', 'AC.B,A', 'ZHKVHG', 'CMDLNJ', '?G.D,N', 'KOGSLE', 'AM!?OU', 'TKCQZA']

Question 3

- a) So if we know cp or cq we can take the gcd of one of them and n to find p or q . This is because when we multiply both equations we get $cp * cq = k^{(2e)} * (pq)^e \pmod{n}$ which is 0 since $pq = n$ $cp * cq$ is a multiple of n hence taking the gcd of one of them and n will give us a factor of n . Once we found p or q we can find q or p , hence finding the factors of n . Then we are able to decrypt the message.
- b) Insanity is doing the same thing, over and over again, but expecting different results.

Question 4

- a) The attacker can choose an integer x which is relatively prime to n . Then they can generate a ciphertext c_x by taking the modulus of x to the power e . ($c_x = x^e \pmod{n}$). Next, they can multiply c_x with the given ciphertext and give it to the oracle. The oracle will return us with a message m times x to the power of $e*d \pmod{N}$. ($c_x * c = (m*x)^{e*d} \pmod{n}$). We know that d is the modular inverse of e so $e*d = 1$ and since x is relatively prime with n we can multiply each side to get our message. $(c_x * c)/x = m \pmod{n}$
- b) You discovered my verry secret message:) Bravo