## Conjugate Gradient

Suppose we want to solve the system of linear equations

$$Ax = b$$

for the vector  $\mathbf{x}$ , where the known  $n \times n$  matrix  $\mathbf{A}$  is symmetric (i.e.,  $\mathbf{A}^T = \mathbf{A}$ ), positive-definite (i.e.  $\mathbf{x}^T \mathbf{A} \mathbf{x} > 0$  for all non-zero vectors  $\mathbf{x} \in \mathbb{R}^n$ ), and real, and  $\mathbf{b}$  is known as well. We denote the unique solution of this system by  $\mathbf{x}^*$ .

Solution of the above system of equations is the same as  $\operatorname{argmin}_{\mathbf{x} \in \mathbb{R}^n} \frac{1}{2} \mathbf{x}^T \mathbf{A} \mathbf{x} - \mathbf{b}^T \mathbf{x}$ .

**Definition 1** (**A**-conjugate directions). Let **A**  $(n \times n)$  be a symmetric matrix. The vectors  $\{\mathbf{d}_1, \ldots, \mathbf{d}_k\}$  (all  $\in \mathbb{R}^n$ ) are called conjugate (or orthogonal) with respect to **A** if  $\mathbf{d}_i^T \mathbf{A} \mathbf{d}_j = 0$  for all  $i \neq j$ .

<u>Lemma</u> 1 (Linear independence). Let **A**  $(n \times n)$  be positive definite. If the vectors  $\{\mathbf{d}_1, \ldots, \mathbf{d}_k\}$  (all  $\in \mathbb{R}^n$ ) are conjugate (orthogonal) with respect to **A**, then they are linearly independent.

*Proof.* Let  $\{\mathbf{d}_1, \dots, \mathbf{d}_k\}$  are linearly dependent and let  $\mathbf{d}_k = \sum_{j=1}^{k-1} \alpha_j \mathbf{d}_j$ .

By definition of conjugacy, we have  $\mathbf{d}_k^T \mathbf{A} \mathbf{d}_j = 0$  for all  $j \neq k$ . Since,  $\mathbf{A}$  is pd, we have  $\mathbf{d}_k^T \mathbf{A} \mathbf{d}_k > 0$ . On the other hand,  $\mathbf{d}_k^T \mathbf{A} \mathbf{d}_k = \sum_{j=1}^{k-1} \alpha_j \mathbf{d}_k^T \mathbf{A} \mathbf{d}_j = 0$  (Contradiction!)

If  $\mathbf{x}^*$  is a solution of the above system of equation, we can represent it as a linear combination of  $\{\mathbf{d}_1, \dots, \mathbf{d}_n\}$ , which are conjugate vectors with respect to  $\mathbf{A}$ .

Let  $\mathbf{x}^* = \sum_{j=1}^n \alpha_j \mathbf{d}_j$ . It is easy to check that  $\mathbf{d}_j^T \mathbf{A} \mathbf{x}^* = \alpha_j \mathbf{d}_j^T \mathbf{A} \mathbf{d}_j$ . Thus  $\alpha_j = \frac{\mathbf{d}_j^T \mathbf{A} \mathbf{x}^*}{\mathbf{d}_j^T \mathbf{A} \mathbf{d}_j}$ .

But by assumption,  $\mathbf{A}\mathbf{x}^* = \mathbf{b}$ . Thus  $\alpha_j = \frac{\mathbf{d}_j^T \mathbf{b}}{\mathbf{d}_j^T \mathbf{A} \mathbf{d}_j}$ . Hence, to know  $\alpha_j$ , we do not need to know  $\mathbf{x}^*$ .

Hence  $\mathbf{x}^* = \sum_{j=1}^n \frac{\mathbf{d}_j^T \mathbf{b}}{\mathbf{d}_i^T \mathbf{A} \mathbf{d}_j} \mathbf{d}_j$ , no matrix inversion or nothing.

In practice, we do not know the conjugate vectors. We thus may follow Gram–Schmidt type algorithm.

## Algorithm 1: Conjugate gradient

- (i) Start setting  $\mathbf{r}_1 = \mathbf{b} \mathbf{A}\mathbf{x}_1$  with some starting value  $\mathbf{x}_1$ .
- (ii) If  $\mathbf{r}_0$  is small enough, return  $\mathbf{x} = \mathbf{x}_1$  as solution.
- (iii) Otherwise, set  $\mathbf{d}_1 = \mathbf{r}_1$  and k = 1.
- (iv) Start loop

$$\bullet \ \alpha_k = \frac{\mathbf{r}_k^T \mathbf{r}_k}{\mathbf{d}_k^T \mathbf{A} \mathbf{d}_k}$$

$$\bullet \ \mathbf{x}_{k+1} = \mathbf{x}_k + \alpha_k \mathbf{d}_k$$

• 
$$\mathbf{r}_{k+1} = \mathbf{r}_k - \alpha_k \mathbf{A} \mathbf{d}_k$$

- If  $\mathbf{r}_{k+1}$  is small enough, return  $\mathbf{x} = \mathbf{x}_{k+1}$  as solution.
- Otherwise,  $\beta_k = \frac{\mathbf{r}_{k+1}^T \mathbf{r}_{k+1}}{\mathbf{r}_k^T \mathbf{r}_k}$ ,  $\mathbf{d}_{k+1} = \mathbf{r}_{k+1} + \beta_k \mathbf{d}_k$ , move k = k+1 and Repeat

In the above algorithm, how the term 'gradient' fits in? Say  $f(\mathbf{x}) = \frac{1}{2}\mathbf{x}^T\mathbf{A}\mathbf{x} - \mathbf{b}^T\mathbf{x}$  and we are trying to solve  $\min_{\mathbf{x}} f(\mathbf{x})$ . Following the steps in gradient descent, the updates should look like  $\mathbf{x}_{k+1} = \mathbf{x}_k - t_k \nabla f(\mathbf{x}_k)$ . For our function, we have  $\nabla f(\mathbf{x}) = \mathbf{A}\mathbf{x} - \mathbf{b}$ . Thus, we need to move towards  $\mathbf{b} - \mathbf{A}\mathbf{x}$ . The other vectors in the basis will be made conjugate to this gradient vector. Thus, the term 'conjugate gradient' comes in.

To see above, note that 
$$\beta_k = \frac{\mathbf{r}_{k+1}^T \mathbf{r}_{k+1}}{\mathbf{r}_k^T \mathbf{r}_k} = -\frac{\mathbf{r}_{k+1} \mathbf{A} \mathbf{d}_k}{\mathbf{d}_k^T \mathbf{A} \mathbf{d}_k}$$
. Then,  $\mathbf{d}_{k+1}^T \mathbf{A} \mathbf{d}_k = 0$ 

If the condition number of  $\mathbf{A}$  is large, the convergence is slow as it leads to slower improvement. Hence, to ensure faster convergence in this case, **Preconditioning conjugate** gradient is used. [Condition number of a matrix  $\mathbf{A}$  is  $\kappa(\mathbf{A}) = \frac{\sigma_{\max}(\mathbf{A})}{\sigma_{\min}(\mathbf{A})}$ , where  $\sigma_{\max}(\mathbf{A})$  and  $\sigma_{\min}(\mathbf{A})$  are maximal and minimal singular values of  $\mathbf{A}$  respectively. In our case, the singular values and eigenvalues are the same as  $\mathbf{A}$  is symmetric]

## Algorithm 2: Pre-conditioned conjugate gradient

$$\mathbf{r}_0 = \mathbf{b} - \mathbf{A}\mathbf{x}_0$$
$$\mathbf{z}_0 = \mathbf{M}^{-1}\mathbf{r}_0$$
$$\mathbf{p}_0 = \mathbf{z}_0$$
$$k = 0$$

repeat

$$\alpha_k = \frac{\mathbf{r}_k^\mathsf{T} \mathbf{z}_k}{\mathbf{p}_k^\mathsf{T} \mathbf{A} \mathbf{p}_k}$$
$$\mathbf{x}_{k+1} = \mathbf{x}_k + \alpha_k \mathbf{p}_k$$
$$\mathbf{r}_{k+1} = \mathbf{r}_k - \alpha_k \mathbf{A} \mathbf{p}_k$$

if  $r_{k+1}$  is sufficiently small, then exit loop end if

$$\mathbf{z}_{k+1} = \mathbf{M}^{-1} \mathbf{r}_{k+1}$$
$$\beta_k = \frac{\mathbf{r}_{k+1}^T \mathbf{z}_{k+1}}{\mathbf{r}_k^T \mathbf{z}_k}$$
$$\mathbf{p}_{k+1} = \mathbf{z}_{k+1} + \beta_k \mathbf{p}_k$$
$$k = k+1,$$

end repeat; The result is  $x_{k+1}$ 

In the above algorithm  $\mathbf{M}$  is fixed. We often set  $\mathbf{M}$  as  $\mathbf{L}\mathbf{L}^T$  where  $\mathbf{L}$  is an incomplete Cholesky decomposition of  $\mathbf{A}$  or sometimes the diagonal entries in  $\mathbf{A}$  (Jacobi preconditioner). Or, we often set  $\mathbf{M}^{-1}$ = some approximate inverse of  $\mathbf{A}$ . The approximate inverse computation would depend on the shape and structure of  $\mathbf{A}$ .