

CT-1 Question 3(i)

Tokenization of the sample code :

< NumType, %d > < keyword, main > < Bracket, (> < Bracket,) >
< Bracket, { >

< NumType, %lf > < Id, 1 > < Operator, = > < Number, 127 >
< Semicolon, ; >

< NumType, %d > < ~~se~~ Semicolon >

< keyword, printf > < Bracket, (> < Literal, 1000 >
< Bracket,) > < Semicolon >

< keyword, if > < Bracket, (> < Id, 1 > < Operator, < >
< Number, 0.0 > < Bracket,) >
< Bracket, { >

< keyword, printf > < Bracket, (> < Literal, 1001 >
< Bracket,) > < Semicolon >

< Bracket, } >

< keyword, else >
< Bracket, { >

< keyword, printf > < Bracket, (> < Literal, 1002 >
< Bracket,) > < Semicolon >

< Bracket, } >

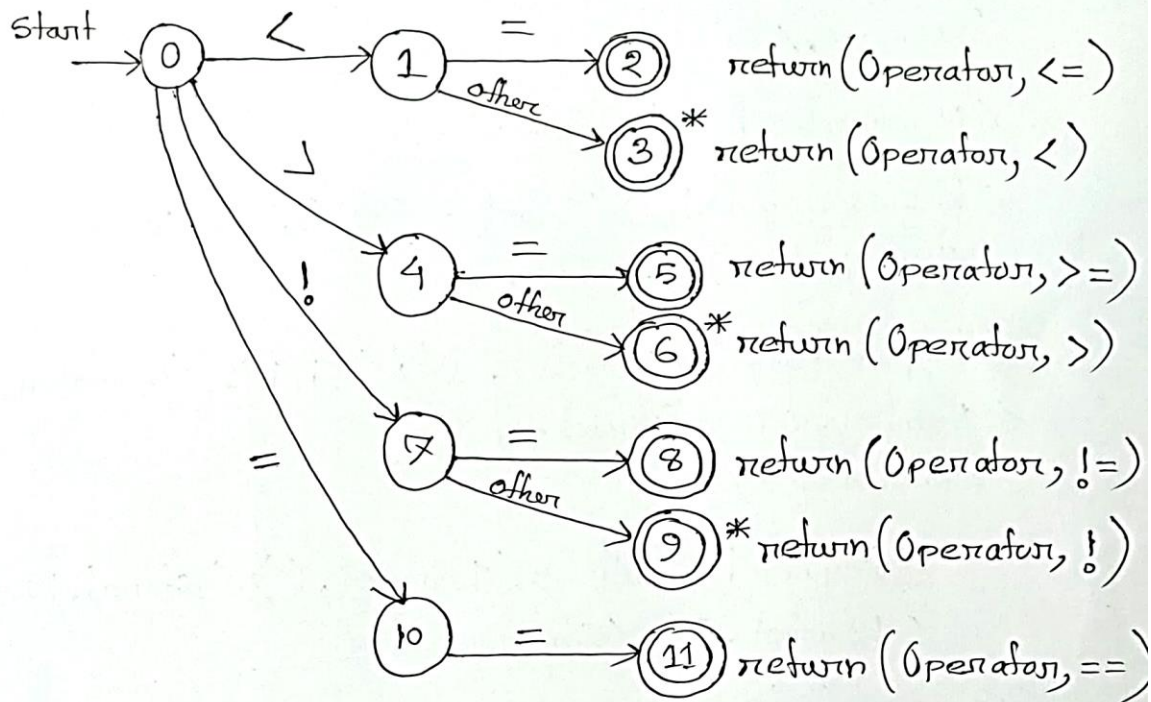
< Number, 0 > < Semicolon >

< Bracket, } >

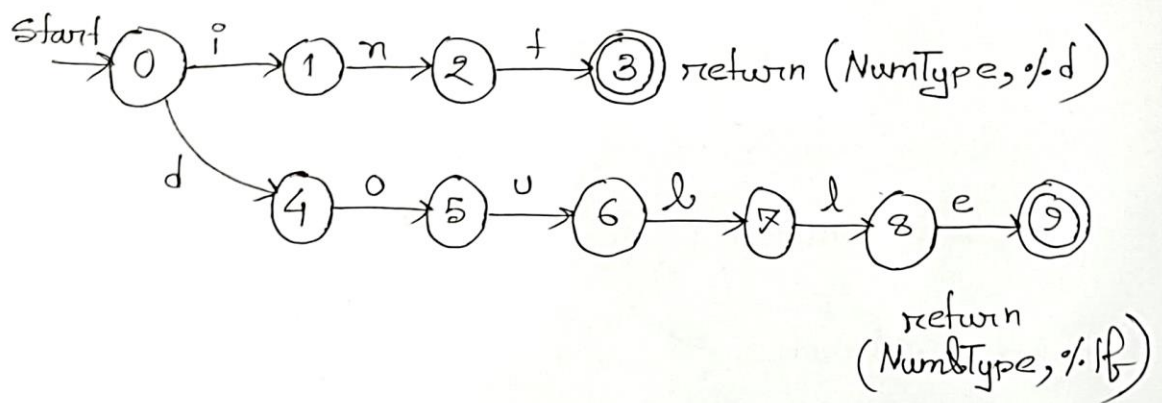
CT-01 Question 3(ii)

Number
NumType : (Same as the transition diagram of number given in your slide)

Operator :



NumType :



Mid-term Question 01

$$S \rightarrow ABC$$

$$A \rightarrow a | c | b | \epsilon$$

$$B \rightarrow c | d | A | \epsilon$$

$$C \rightarrow e | f$$

FIRST Sets :

$$\begin{aligned} \text{FIRST}(S) &= \{\text{FIRST}(A)\} = \{a, e, b, (\epsilon)\} \\ &= \{a, e, b, \text{FIRST}(B)\} \\ &= \{a, e, b, c, d, (\epsilon)\} \\ &= \{a, e, b, c, d, \text{FIRST}(C)\} \\ &= \{a, e, b, c, d\} \end{aligned}$$

$$\text{FIRST}(A) = \{a, \text{FIRST}(C), \epsilon\} = \{a, e, b, \epsilon\}$$

$$\text{FIRST}(B) = \{c, d, \epsilon\}$$

$$\text{FIRST}(C) = \{e, f\}$$

FOLLOW SETS :

	①	②	③	③
FOLLOW(S)	{ \$ }	{ \$ }	{ \$ }	{ \$ }
FOLLOW(A)		{ c, d, e, f }	{ c, d, e, f }	{ c, d, e, f }
FOLLOW(B)		{ e, f }	{ e, f }	{ e, f }
FOLLOW(C)		{ b }	{ \$, b }	{ b, \$ }

$\boxed{\text{①} \rightarrow \text{Rule 1}, \text{②} \rightarrow \text{Rule 2. } (A \rightarrow \alpha B \beta), \text{③} \rightarrow \text{Rule 3 } (A \rightarrow \alpha B)}$
 $S \rightarrow \underline{ABC} \quad \therefore \text{When } B=A, \beta=BC, \therefore \text{FOLLOW}(A) = \overline{\text{FIRST}(BC)}$
 $\quad \alpha B \beta \quad \quad \quad = \{c, d, e, f\}$

Left Factoring

Mid-term Ques-02

$$\begin{aligned} * A &\rightarrow abAB \mid abeBe \mid abeA \\ B &\rightarrow b \mid d \mid \epsilon \end{aligned}$$

$$\text{Soln : } A \rightarrow \underline{ab}AB \mid \underline{abe}Be \mid \underline{abe}A$$

$$\begin{aligned} A &\rightarrow abA' \\ A' &\rightarrow AB \mid \underline{e}Be \mid \underline{e}A \end{aligned}$$

$$\begin{aligned} A' &\rightarrow AB \mid eA'' \\ A'' &\rightarrow Be \mid A \end{aligned}$$

$$\begin{aligned} \therefore \text{Finally, } A &\rightarrow abA' \\ A' &\rightarrow AB \mid eA'' \\ A'' &\rightarrow Be \mid A \\ B &\rightarrow b \mid d \mid \epsilon \end{aligned}$$