

Artificial Intelligence



8 First Order Logic

Russell & Norvig, AI: A Modern Approach, 3rd Ed

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First Order Logic

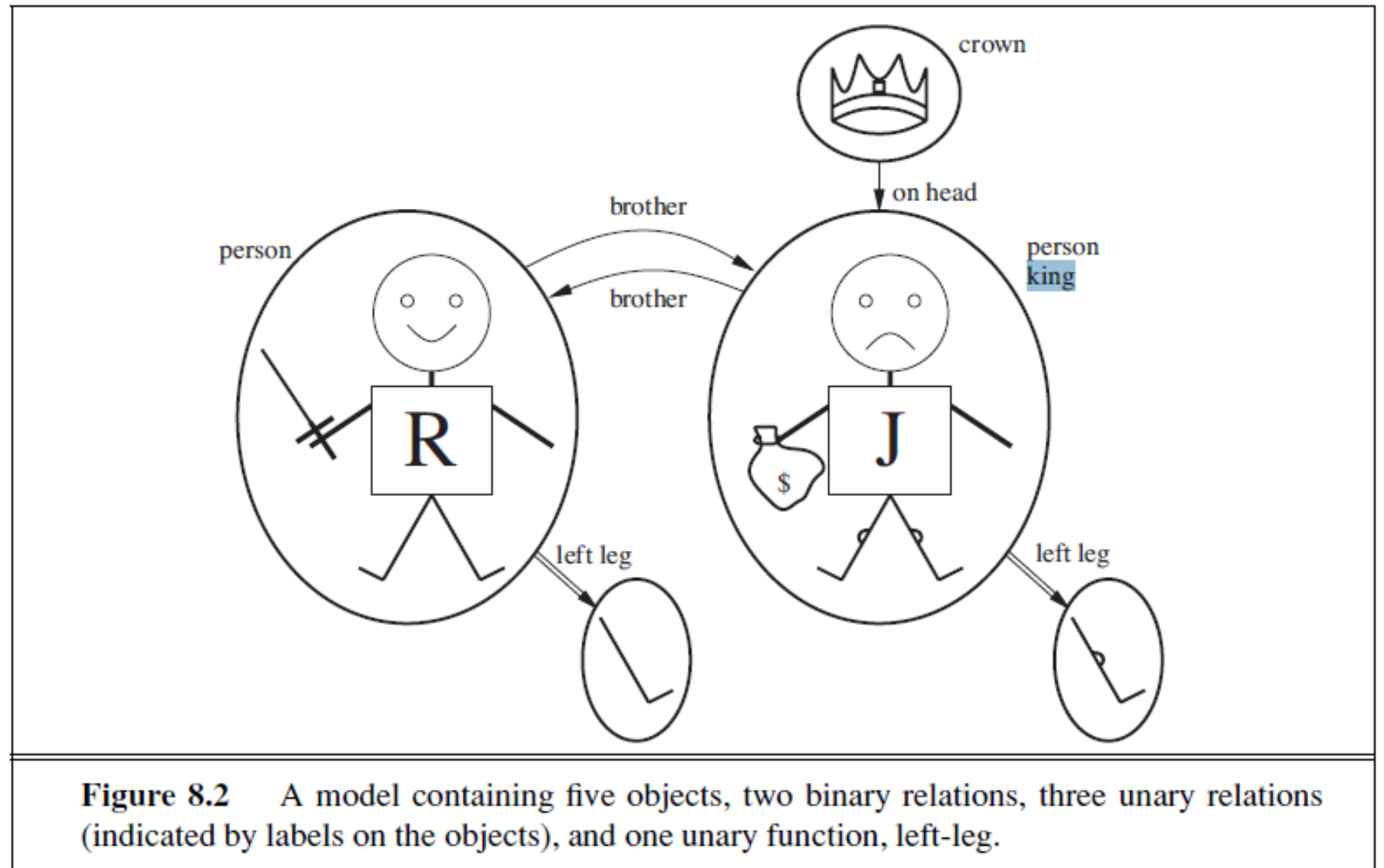
Used for representing a complex environments.

❖ Difference with propositional logic:

- ✓ Existence of predicates
- ✓ And Quantifiers

First Order Logic - Model

- ❑ Object
- ❑ Relation
- ❑ Function



Syntax

- Constant
- Variable – Small Letter
- Predicate – Ret T/F
- Function – Ret any value

$$\begin{aligned} \text{Sentence} &\rightarrow \text{AtomicSentence} \mid \text{ComplexSentence} \\ \text{AtomicSentence} &\rightarrow \text{Predicate} \mid \text{Predicate}(\text{Term}, \dots) \mid \text{Term} = \text{Term} \\ \text{ComplexSentence} &\rightarrow (\text{Sentence}) \mid [\text{Sentence}] \\ &\mid \neg \text{Sentence} \\ &\mid \text{Sentence} \wedge \text{Sentence} \\ &\mid \text{Sentence} \vee \text{Sentence} \\ &\mid \text{Sentence} \Rightarrow \text{Sentence} \\ &\mid \text{Sentence} \Leftrightarrow \text{Sentence} \\ &\mid \text{Quantifier Variable}, \dots \text{Sentence} \end{aligned}$$
$$\begin{aligned} \text{Term} &\rightarrow \text{Function}(\text{Term}, \dots) \\ &\mid \text{Constant} \\ &\mid \text{Variable} \end{aligned}$$
$$\begin{aligned} \text{Quantifier} &\rightarrow \forall \mid \exists \\ \text{Constant} &\rightarrow A \mid X_1 \mid \text{John} \mid \dots \\ \text{Variable} &\rightarrow a \mid x \mid s \mid \dots \\ \text{Predicate} &\rightarrow \text{True} \mid \text{False} \mid \text{After} \mid \text{Loves} \mid \text{Raining} \mid \dots \\ \text{Function} &\rightarrow \text{Mother} \mid \text{LeftLeg} \mid \dots \end{aligned}$$

OPERATOR PRECEDENCE : $\neg, =, \wedge, \vee, \Rightarrow, \Leftrightarrow$

Syntax

- ❖ Two types of quantifiers
 - ✓ Universal quantification, \forall : for all
 - ✓ Existential quantification, \exists : for some

First Order Logic – Example 1

Sentences: *Sakib is a cricketer.*

Atomic Sentences: is formed from a predicate symbol optionally followed by a parenthesized list of terms.

1. x is a cricketer \equiv Cricketer(x)

FOL: *Crickter(Sakib)*

First Order Logic – Example 2

Sentences: *Romeo loves Juliet.*

Atomic Sentences: is formed from a predicate symbol optionally followed by a parenthesized list of terms.

1. $x \text{ loves } y \equiv \text{Loves}(x, y)$

FOL: *Loves(Romeo, Juliet)*

First Order Logic – Example 3

Sentences: *Mary loves everyone.*

Atomic Sentences: is formed from a predicate symbol optionally followed by a parenthesized list of terms.

1. $x \text{ loves } y \equiv \text{Loves}(x, y)$

FOL: $\forall x \text{ Loves}(\text{Mary}, x)$

First Order Logic – Example 4

Sentences: *There is someone who loves Mary.*

Atomic Sentences: is formed from a predicate symbol optionally followed by a parenthesized list of terms.

1. $x \text{ loves } y \equiv \text{Loves}(x, y)$

FOL: $\exists x \text{ Loves}(x, \text{Mary})$

First Order Logic – Example 5

Sentences: *Everyone loves it's mother.*

Way 1:

Atomic Sentences: is formed from a predicate symbol optionally followed by a parenthesized list of terms.

1. x loves y \equiv Loves (x , y)
2. Return mother of x \equiv Mother (x)

FOL: $\forall_x \text{Loves } (x, \text{Mother } (x))$

First Order Logic – Example 5 . . .

Sentences: *Everyone loves it's mother.*

Way 2:

Atomic Sentences: is formed from a predicate symbol optionally followed by a parenthesized list of terms.

1. x's mother is y \equiv Mother (x, y)
2. x loves y \equiv Loves (x, y)

FOL: $\forall_x \exists_y \text{Mother (x, y) } \wedge \text{Loves (x, y)}$

First Order Logic – Example 6

Sentences: *No person likes a professor unless the professor is smart.*

Atomic Sentences:

- | | | |
|------------------|----------|----------------------------|
| 1. Person(x) | \equiv | x is a person. |
| 2. Professor (x) | \equiv | x is a professor. |
| 3. Smart(x) | \equiv | x is a smart. |
| 4. Likes(x, y) | \equiv | subject x likes subject y. |

FOL: $\forall x \forall y \text{ Person}(x) \wedge \text{Professor}(y) \wedge \neg \text{Smart}(y) \Rightarrow \neg \text{likes}(x,y)$

Example

Sentence	First Order Logic Representation
Lucy is a professor	is-prof(lucy)
All professors are people.	$\forall x (\text{is-prof}(x) \Rightarrow \text{is-person}(x))$
Fuchs is the dean.	is-dean(fuchs)
All Deans are professors.	$\forall x (\text{is-dean}(x) \Rightarrow \text{is-prof}(x))$
Everyone is a friend of someone.	$\forall x (\exists y (\text{is-friend-of}(y, x)))$
Lucy criticized Fuchs.	criticize(lucy,fuchs)

Nested quantifiers

- Consecutive quantifiers of the same type can be written as one quantifier with several variables.
 - Siblinghood is a symmetric relationship
 - $\forall x, y \text{ Sibling}(x, y) \Leftrightarrow \text{Sibling}(y, x)$
- The order of quantification is therefore very important. It becomes clearer if we insert parentheses.
 - Everybody loves somebody
 - $\forall x \exists y \text{ Loves}(x, y)$
 - $\exists y \forall x \text{ Loves}(x, y)$

Connections between \forall and \exists

- The two quantifiers are actually intimately connected with each other, through negation
- \forall is a conjunction over the universe of objects and \exists is a disjunction
- They obey De Morgan's rules
 - $\forall x \neg P \equiv \neg \exists x P$
 - $\neg \forall x P \equiv \exists x \neg P$
 - $\forall x P \equiv \neg \exists x \neg P$
 - $\exists x P \equiv \neg \forall x \neg P$
- Everyone likes ice cream : $\forall x \text{ Likes}(x, \text{IceCream})$
- There is no one who does not like ice cream:
 $\neg \exists x \neg \text{Likes}(x, \text{IceCream})$

Equality

- **Equality symbol** can be used to signify that two terms refer to the same object.
 - $\text{Father}(\text{John}) = \text{Henry}$
- Richard has at least two brothers.
 - $\exists x, y \text{ Brother}(x, \text{Richard}) \wedge \text{Brother}(y, \text{Richard})$
 - $\exists x, y \text{ Brother}(x, \text{Richard}) \wedge \text{Brother}(y, \text{Richard}) \wedge \neg(x=y)$

More Examples – on Kinship

- One's mother is one's female parent:
- $\forall m, c \text{ Mother}(c)=m \Leftrightarrow \text{Female}(m) \wedge \text{Parent}(m, c) .$
- One's husband is one's male spouse:
- $\forall w, h \text{ Husband}(h,w) \Leftrightarrow \text{Male}(h) \wedge \text{Spouse}(h,w) .$
- Male and female are disjoint categories:
- $\forall x \text{ Male}(x) \Leftrightarrow \neg \text{Female}(x) .$
- Parent and child are inverse relations:
- $\forall p, c \text{ Parent}(p, c) \Leftrightarrow \text{Child}(c, p) .$
- A grandparent is a parent of one's parent:
- $\forall g, c \text{ Grandparent}(g, c) \Leftrightarrow \exists p \text{ Parent}(g, p) \wedge \text{Parent}(p, c) .$
- A sibling is another child of one's parents:
- $\forall x, y \text{ Sibling}(x, y) \Leftrightarrow x \neq y \wedge \exists p \text{ Parent}(p, x) \wedge \text{Parent}(p, y) .$

Try yourself:

1. All professors consider the dean a friend or don't know him.
2. People only criticize people that are not their friends.
3. Only one student failed in History.
4. No person likes a leader unless the leader is smart
5. Everyone who loves all animals is loved by someone
6. All kids are short.
7. Certain kids own shoes.
8. If someone is a kid and a boy, he likes cars.
9. All kids that own shoes, wear them.
10. All kids have a mother.
11. If a woman is a mother, she has a kid.
12. No kid likes a vegetable if it's overcooked.
13. No mother likes a teacher if he punishes her kid.

The knowledge-engineering process

The general process of knowledge-base construction:

1. Identify the task. – PEAS
2. Assemble the relevant knowledge. – knowledge acquisition
3. Decide on a vocabulary of predicates, functions, and constants.
4. Encode general knowledge about the domain.
5. Encode a description of the specific problem instance.
6. Pose queries to the inference procedure and get answers.
7. Debug the knowledge base.

Thank you 😊