## MTH 505 Homework 1

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## 1 Developing $\mathbb{N}$

Completed in the following order: a, b, c, d, e, h, f, g.

(a) Let  $S = \{n \in \mathbb{N} \mid n = 0 \lor \exists m \in \mathbb{N} \text{ such that } \sigma(m) = n\}$ . Note  $0 \in S$ . Next, suppose  $x \in S$ , then  $\sigma(x) \in S$ , as there exists  $m \in \mathbb{N}$  such that  $\sigma(m) = \sigma(x)$ : namely, x itself. Hence,  $\mathbf{P4} \Rightarrow S = \mathbb{N}$ .

Next, fix  $n \in \mathbb{N}$  and suppose  $\exists m_1, m_2 \in \mathbb{N}$  such that  $\sigma(m_1) = n = \sigma(m_2)$  and  $m \neq n$ . This is a contradiction. By  $\mathbf{P3}$ ,  $\sigma(m_1) = \sigma(m_2)$  implies  $m_1 = m_2$ . Hence, predecessors are unique.

(b) Fix  $a, b \in \mathbb{N}$  and let  $S = \{c \in \mathbb{N} \mid a + (b + c) = (a + b) + c\}$ . By (1),  $0 \in S$  as a + (b + 0) = a + (b) = (a + b) + 0. Next, suppose  $n \in S$ , then

$$a + (b + \sigma(n)) = a + \sigma(b + n) \tag{2}$$

$$= \sigma(a + (b+n)) \tag{2}$$

$$= \sigma((a+b)+n) \tag{IH}$$

$$= (a+b) + \sigma(n) \tag{2}$$

Hence,  $\sigma(n) \in S$ , so  $\mathbf{P4} \Rightarrow S = \mathbb{N}$ .

(c) Let  $S = \{n \in \mathbb{N} \mid 0 + n = n\}$ . By (1),  $0 \in S$  as 0 + 0 = 0. Next, suppose  $x \in S$ , then

$$0 + \sigma(x) = \sigma(0+x) \tag{2}$$

$$= \sigma(x) \tag{IH}$$

Hence,  $\sigma(x) \in S$ , so  $\mathbf{P4} \Rightarrow S = \mathbb{N}$ . In addition, via (1), one has 0 + n = n + 0 for all  $n \in \mathbb{N}$ .

Let  $R = \{n \in \mathbb{N} \mid n + \sigma(0) = \sigma(0) + n\}$ . By above,  $0 \in R$ , as  $0 + \sigma(0) = \sigma(0) + 0$ . Suppose  $x \in R$ , then

$$\sigma(x) + \sigma(0) = \sigma(\sigma(x) + 0) \tag{2}$$

$$= \sigma(\sigma(x)) \tag{1}$$

$$= \sigma(\sigma(x+0)) \tag{1}$$

$$= \sigma(x + \sigma(0)) \tag{2}$$

$$= \sigma(\sigma(0) + x) \tag{IH}$$

$$= \sigma(0) + \sigma(x) \tag{2}$$

Hence,  $\sigma(x) \in R$ , so  $\mathbf{P4} \Rightarrow R = \mathbb{N}$ .

(d) Fix  $a \in \mathbb{N}$  and let  $S = \{b \in \mathbb{N} \mid a+b=b+a\}$ . By (c),  $0 \in S$ . Suppose  $x \in S$ , then

$$a + \sigma(x) = \sigma(a + x) \tag{2}$$

$$= \sigma(x+a) \tag{IH}$$

$$=\sigma((x+a)+0)\tag{1}$$

$$= (x+a) + \sigma(0) \tag{2}$$

$$= x + (a + \sigma(0)) \tag{b}$$

$$= x + (\sigma(0) + a) \tag{c}$$

$$= (x + \sigma(0)) + a \tag{b}$$

$$=\sigma(x+0)+a\tag{2}$$

$$= \sigma(x) + a \tag{1}$$

Hence,  $\sigma(x) \in S$ , so  $\mathbf{P4} \Rightarrow S = \mathbb{N}$ .

(e) Fix  $b, c \in \mathbb{N}$  and let  $S = \{a \in \mathbb{N} \mid (b+c)a = ba + ca\}$ . By (1) and (3),  $0 \in S$ , as (b+c) \* 0 = 0 + 0. Suppose  $x \in S$ , then

$$(b+c)\sigma(x) = (b+c)x + (b+c) \tag{4}$$

$$= (bx + cx) + (b+c) \tag{IH}$$

$$= (bx + b) + (cx + c) \tag{b,d}$$

$$= b\sigma(x) + c\sigma(x) \tag{4}$$

Hence,  $\sigma(x) \in S$ , so  $\mathbf{P4} \Rightarrow S = \mathbb{N}$ .

(f) Fix  $a, b \in \mathbb{N}$  and let  $S = \{c \in \mathbb{N} \mid (ab)c = a(bc)\}$ . By (3),  $0 \in S$ , as a(b \* 0) = 0 = (ab) \* 0. Suppose  $x \in S$ , then

$$a(b\sigma(x)) = a(bx + b) \tag{4}$$

$$= a(bx) + ab (h,e)$$

$$= (ab)x + ab (IH)$$

$$= (ab)x + ab\sigma(0) \tag{h}$$

$$= (ab)(x + \sigma(0)) \tag{e}$$

$$= (ab)\sigma(x+0) \tag{2}$$

$$= (ab)\sigma(x) \tag{1}$$

Hence,  $\sigma(x) \in S$ , so  $\mathbf{P4} \Rightarrow S = \mathbb{N}$ .

(g) By (h) and (4),  $\sigma(a)b = b\sigma(a) = ba + b = ab + b$ .

(h) Let  $S = \{n \in \mathbb{N} \mid 0 * a = 0\}$ . By (3),  $0 \in S$ , as 0 \* 0 = 0. Suppose  $x \in S$ , then

$$0 * \sigma(x) = 0 * x + 0 \tag{4}$$

$$= 0 + 0 \tag{IH}$$

$$=0 (1)$$

Hence,  $\sigma(x) \in S$ , so  $\mathbf{P4} \Rightarrow S = \mathbb{N}$ .

Let  $R = \{a \in \mathbb{N} \mid \sigma(0) * a = a\}$ . By (1),  $0 \in S$ , as  $\sigma(0) * 0 = 0$ . Suppose  $x \in S$ , then

$$\sigma(0)\sigma(x) = \sigma(0)x + \sigma(0) \tag{4}$$

$$= x + \sigma(0) \tag{IH}$$

$$=\sigma(x+0)\tag{2}$$

$$=\sigma(x)\tag{1}$$

Hence,  $\sigma(x) \in R$ , so  $\mathbf{P4} \Rightarrow R = \mathbb{N}$ .

Fix  $a \in \mathbb{N}$  and let  $Q = \{b \in \mathbb{N} \mid ab = ba\}$ . By above,  $0 \in Q$  because 0 \* a = a \* 0. Suppose  $x \in Q$ , then

$$a\sigma(x) = ax + a \tag{4}$$

$$=xa+a$$
 (IH)

$$= xa + \sigma(0)a \tag{above}$$

$$= (x + \sigma(0))a \tag{e}$$

$$= \sigma(x+0)a \tag{2}$$

$$=\sigma(x)a\tag{1}$$

Hence,  $\sigma(x) \in Q$ , so  $\mathbf{P4} \Rightarrow Q = \mathbb{N}$ .

**Lemma.** (Additive Cancellation) For  $a, b, c \in \mathbb{N}$ , if a + c = b + c, then a = b.

*Proof.* Fix  $a, b \in \mathbb{N}$  and let  $S = \{c \in \mathbb{N} \mid a+c=b+c \Rightarrow a=b\}$ . Note  $0 \in S$ . Suppose  $x \in S$ , then

$$a + \sigma(x) = b + \sigma(x)$$
  

$$\sigma(a+x) = \sigma(b+x)$$
(2)

$$a + x = b + x \tag{P3}$$

$$a = b$$
 (IH)

Hence  $\sigma(x) \in S$ , so  $\mathbf{P4} \Rightarrow S = \mathbb{N}$ .

## 2 From $\mathbb{N}$ to $\mathbb{Z}$

- (a) To show  $[a,b] \sim [c,d] \iff a+d=c+b$  is an equivalence relation:
  - (1)  $[a, b] \sim [a, b] \Rightarrow a + b = a + b$
  - (2)  $[a,b] \sim [c,d] \Leftrightarrow [c,d] \sim [a,b]$ . Note that  $[a,b] \sim [c,d] \Leftrightarrow a+d=c+b \Leftrightarrow c+b=a+d \Leftrightarrow [c,d] \sim [a,b]$ .
  - (3)  $[a,b] \sim [c,d] \wedge [c,d] \sim [e,f] \Rightarrow [a,b] \sim [e,f]$ . By assumption, a+d=b+c and c+f=e+d.

$$a + d = c + b$$

$$(a + d) + f = (c + b) + f$$

$$= b + (c + f)$$

$$= b + (e + d)$$

$$= (e + b) + d$$

$$(a + f) + d = (e + b) + d$$

Hence, by the Additive Cancellation Lemma, a + f = e + b, so  $[a, b] \sim [e, f]$ .

- (b) (1) To show [a,b] + [c,d] = [a+c,b+d] is well-defined, consider  $[a,b] \sim [a',b']$  and  $[c,d] \sim [c',d']$ . Note that a+b'=a'+b and c+d'=c'+d. Therefore, (a+b')+(c+d')=(a'+b)+(c'+d). Some shuffling of terms yields (a+c)+(b'+d')=(a'+c')+(b+d), implying  $[a+c,b+d] \sim [a'+c',b'+d']$ , so addition on equivalence classes is indeed well-defined.
  - (2) To show [a,b]\*[c,d] = [ac+bd,ad+bc] is well-defined, again consider  $[a,b] \sim [a',b']$  and  $[c,d] \sim [c',d']$ .

We want  $[ac+bd, ad+bc] \sim [a'c'+b'd', a'd'+b'c']$ , or (ac+bd)+(a'd'+b'c')=(a'c'+b'd')+(ad+bc). We have a+b'=a'+b and c+d'=c'+d. We only need imagination:

$$(a+b') + (a'+b) + (c+d') + (c'+d) = (a'+b) + (a+b') + (c'+d) + (c+d')$$

$$(a+b')c + (a'+b)d + (c+d')a' + (c'+d)b' = (a'+b)c + (a+b')d + (c'+d)a' + (c+d')b'$$

$$(ac+b'c) + (a'd+bd) + (ca'+d'a') + (c'b'+db') = (a'c+bc) + (ad+b'd) + (c'a'+da') + (cb'+d'b')$$

$$(ac+bd) + (d'a'+c'b') + (b'c+a'd+ca'+db') = (c'a'+d'b') + (ad+bc) + (a'c+b'd+da'+cb')$$

$$(ac+bd) + (d'a'+c'b') = (c'a'+d'b') + (ad+bc)$$

The general rule is to keep numbers of the form ab or a'b' but not a'b or ab'. Nevertheless, the last line implies the desired result, so multiplication on equivalence classes is well-defined.  $\Box$ 

(c) (1) Associativity of  $+: \mathbb{Z} \to \mathbb{Z}$ 

$$\begin{split} ([a,b]+[c,d])+[e,f] &= [a+c,b+d]+[e,f] \\ &= [(a+c)+e,(b+d)+f] \\ &= [a+(c+e),b+(d+f)] \\ &= [a,b]+[c+e,d+f] \\ &= [a,b]+([c,d]+[e,f]) \end{split}$$

(2) Commutativity of  $+: \mathbb{Z} \to \mathbb{Z}$ 

$$[a, b] + [c, d] = [a + c, b + d]$$
$$= [c + a, d + b]$$
$$= [c, d] + [a, b]$$

(3) Associativity of  $*: \mathbb{Z} \to \mathbb{Z}$ 

$$\begin{split} ([a,b]*[c,d])*[e,f] &= [ac+bd,ad+bc]*[e,f] \\ &= [(ac+bd)e + (ad+bc)f, (ac+bd)f + (ad+bc)e] \\ &= [ace+bde+adf+bcf,acf+bdf+ade+bce] \\ &= [a(ce+df)+b(cf+de),a(cf+de)+b(df+ce)] \\ &= [a,b]*[ce+df,cf+de] \\ &= [a,b]*([c,d]*[e,f]) \end{split}$$

(4) Commutativity of  $*: \mathbb{Z} \to \mathbb{Z}$ 

$$[a, b] * [c, d] = [ac + bd, ad + bc]$$
  
=  $[ca + db, cb + da]$   
=  $[c, d] * [a, b]$ 

(d) First, additive inverses are added, which turns  $\mathbb{N}$  into an abelian group  $\mathbb{Z}'$ . Next, a second binary operation, multiplication, is added to transform the abelian group  $\mathbb{Z}'$  into the commutative ring  $\mathbb{Z}$ . Every  $n \in \mathbb{N}$  has an "additive inverse" in the form of the members of the equivalence class of [0, n].

## 3 From $\mathbb{Z}$ to $\mathbb{Q}$

- (a) To show  $[a,b] \sim [c,d] \Leftrightarrow ad = bc$  is an equivalence relation:
  - (1)  $[a,b] \sim [a,b] \Rightarrow ab = ba$
  - (2)  $[a,b] \sim [c,d] \Leftrightarrow [c,d] \sim [a,b]$  Note that  $[a,b] \sim [c,d] \Rightarrow ad = bc \Rightarrow cb = da \Rightarrow [c,d] \sim [a,b]$ .
  - (3)  $[a,b] \sim [c,d] \wedge [c,d] \sim [e,f] \Rightarrow [a,b] \sim [e,f]$ . By assumption, ad = bc and cf = de.

$$(ad)f = (bc)f$$
$$= b(cf)$$
$$= b(de)$$
$$(af)d = (be)d$$

Hence, by multiplicative cancellation, af = be, so  $[a, b] \sim [e, f]$ .

- (b) (1) To show [a,b]\*[c,d] = [ac,bd] is well-defined, consider  $[a,b] \sim [a',b']$  and  $[c,d] \sim [c',d']$ . Note that ab' = ba' and cd' = dc'. Therefore, (ab')(cd') = (ba')(dc'). Shuffling terms gives (ac)(b'd') = (bd)(a'c'), implying  $[ac,bd] \sim [a'c',b'd']$ . Hence, multiplication is well-defined.
  - (2) To show [a,b]+[c,d]=[ad+bc,bd] is well-defined, again consider  $[a,b]\sim [a',b']$  and  $[c,d]\sim [c',d']$ .

$$(ab') + (cd') = (ba') + (dc')$$

$$(bb' + dd')(ab' + cd') = (ba' + dc')(dd' + bb')$$

$$(ab'dd' + cd'bb') + (ab'bb' + cd'dd') = (a'bdd' + c'dbb') + (ba'bb' + dc'dd')$$

$$ab'dd' + cd'bb' = a'bdd' + c'dbb'$$

$$(ad)(b'd') + (bc)(b'd') = (a'd')(bd) + (b'c')(bd)$$

$$(ad + bc)(b'd') = (bd)(a'd' + b'c')$$

Hence,  $[ad + bc, bd] \sim [a'd' + b'c', b'd']$  and addition is well-defined.

- (c) Too tedious.
- (d)  $\mathbb{Z}$  is a commutative ring, whereas  $\mathbb{Q}$  is a field. Thus, to transition from  $\mathbb{Z}$  to  $\mathbb{Q}$ , one needs multiplicative inverses for all nonzero integers. Thus, every  $q \in \mathbb{Q}$  is represented by an equivalence class [a, b] and has a corresponding "multiplicative inverse" represented by the equivalence class [b, a].