

Write You Some Proofs for Great Good

<https://github.com/rpeszek/present-proofs-lc19>

<https://github.com/rpeszek/present-proofs-lc19/blob/master/doc/slides.pdf>

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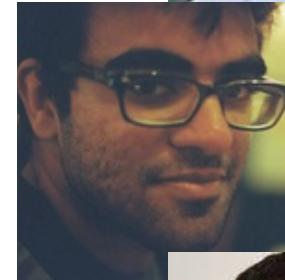
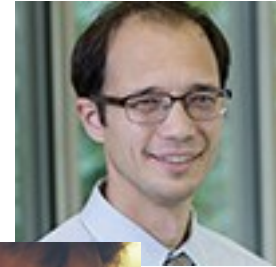
FP @ Innovation Lab, Holland & Hart

<https://www.meetup.com/Boulder-Haskell-Programmers/>

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Why proofs?

- Curry Howard (Why proofs == Why programs)*
- Formal verification
- Enable `Type Precision` (*focus of this talk*)

```
data List a =  
    Empty  
    | Cons a (List a)
```





v
s

```
data Vect (n :: Nat) a where  
    Empty :: Vect 'Z a  
    Cons :: a -> Vect n a -> Vect ('S n) a
```

Plan for this talk

- Narrow scoped, pragmatic intro to proofs with dependent types
- Focus on what can be done vs how
- Anything wrong with List?
- Anything wrong with Maybe?
- Anything wrong with Bool?
- Q: Proofs and software engineering - together?
- Concerns
 - Performance; Termination / Totality; Maintenance; ...

Proofs by who writes them

	Programmers *	Automated Logic Solvers
Work		Free Lunch! 
What can be done	A Lot 	Limited 
Examples	Dependently Typed Languages	Refinement Types LiquidHaskell
	<i>(focus of this talk)</i>	

Motivation (*call-side safety*)

```
(!!) :: [a] -> Int -> a
```

```
safeGet :: [a] -> Int -> Maybe a
```

```
-- Liquid:
```

```
{-@ (!! ) :: x: [a] -> {i:Nat | i < len x} -> a @-}
```

```
(!!) :: [a] -> Int -> a
```

```
-- Dep Typed:
```

```
(!!) :: Vect n a -> Fin n -> a
```

```
(!!!) :: Vect n a -> SNat m -> MaybeB (m < n) a
```

```
(!!!!) : (n:Nat) -> (xs:List a) -> {auto ok: InBounds n xs} -> a
```

- Questionable improvement
- Refinements
- Precise types (*talk focus*)

... and much more

<https://github.com/rpeszek/present-proofs-lc19/tree/master/src/Motivation>

Motivation *(implementation safety)*

(code)

<https://github.com/rpeszek/present-proofs-lc19/blob/master/src/Motivation/Exfer.idr>

Intro

(code) Curry-Howard vs Imperative

<https://github.com/rpeszek/present-proofs-lc19/blob/master/src/Present/AnIntro.hs>

<https://github.com/rpeszek/present-proofs-lc19/blob/master/src/Present/MaybeB.hs>

Problem in Paradise - Questions

Should type checker know basic algebra?

`a || True <==> True (Bool algebra)`

`a + b == b + a (Nat, Int, Float ... algebra) ...`

Answers:

- Dependently Types Langs: **No** (*Programmers supply proofs*)
- Refinement Types / LiquidHaskell: **Yes** (*SMT solver does the work*)

Type Equality

```
data a ~: b where  
  Refl :: a ~: a
```

typecheck **Refl** only if same types

```
test1  = Refl :: 5 ~: 5      -- GOOD  
test10 = Refl :: 4 ~: 5      -- ERR  
  
test2  = Refl :: 2 + 3 ~: 3 + 2 -- GOOD  
test20 :: SNat n1 -> SNat n2 -> n1 + n2 ~: n2 + n1  
test20 _ _ = Refl           -- ERR
```

[\(base\) Data.Type.Equality](#)

Example Combinators

- library *over* op semantics
- "pattern-matching on a variable of type $(a \sim b)$ produces a proof that $a \sim b$ " - *haddock*

```
sym :: (a ~ b) -> (b ~ a)
```

```
trans :: (a ~ b) -> (b ~ c) -> (a ~ c)
```

```
apply :: (f ~ g) -> (a ~ b) -> (f a ~ g b)
```

```
inner :: (f a ~ g b) -> (a ~ b)
```

```
castWith :: (a ~ b) -> a -> b
```

```
gcastWith :: (a ~ b) -> ((a ~ b) => r) -> r
```

```
gcastWith Refl x = x
```

math useful and tedious

a bit unexpected ?

fullfil ~ constraint using ~

[\(base\) Data.Type.Equality](#)

Proofs - Bool Algebra

(code)

<https://github.com/rpeszek/present-proofs-lc19/blob/master/src/Present/ProofsBoolAlg.hs>

Better Bool ... Decidability

```
import Data.Void

data Dec prop = Yes prop |
               No  (prop -> Void)
```

<https://github.com/rpeszek/present-proofs-lc19/blob/master/src/Present/ProofsDecidable.hs>

Proofs - Nat (Performance)

```
reverse :: Vect n a -> Vect n a  
reverse = ...
```

```
-- recursive calls
```

```
plusCommutative :: SNat left -> SNat right -> ((left + right) :~: (right + left))  
plusCommutative left right = case left of  
  SZ -> lemma1 right  
  (SS k) -> case plusCommutative k right of Refl -> sym (lemma2 right k)
```

unsafeCoerce replacements

[https://github.com/rpeszek/present-proofs-lc19/blob/master/src/Present/ProofsNatAlg.
hs](https://github.com/rpeszek/present-proofs-lc19/blob/master/src/Present/ProofsNatAlg.hs)

Some Learning/References

- intro books
 - [Type Driven Development in Idris](#) great book
 - <https://github.com/rpeszek/IdrisTddNotes/wiki>
 - [TAPL](#) great book (*not really dep types but still*)
 - Programming foundation books ([penn/Pierce et al](#), [Wadler](#))
- Haskell projects (with reading references)
 - [singletons](#)
 - [equational-reasoning-in-haskell](#)
 - [liquidhaskell](#)
- blogs
 - [blog.jle.im](#) (Justin Le)
 - [typesandkinds](#) (Richard Eisenberg)
- youtube
 - [Introduction to Agda](#) series by Daniel Peebles published by Edward Kmett