Write You Some Proofs for Great Good

https://github.com/rpeszek/present-proofs-lc19

https://github.com/rpeszek/present-proofs-lc19/blob/master/doc/slides.pdf

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Proofs by who writes them

	Programmers *	Automated Logic Solvers
Work		Free Lunch!
What can be done	A Lot	Limitted
Examples	Dependently Typed Languages	Refinement Types LiquidHaskell
	(focus of this talk)	

Plan for this talk

Narrow scoped intro to proofs with dependent types, software engineer viewpoint

- Mathematics and Software Engineering
 - Type Precision Gentlemen's agreements
 - Progs == Proofs Software == has Bugs
 - Dependent Types Implementation Hiding
 - Proofs
 Maintenance Cost
- Type Precision: better List, better Maybe, better Bool, ...
- Termination / Totality (can prove nonsense by throwing exception)
- Performance
- Maintenance, Peano Nat vs GHC.TypeLits

Why proofs?

- Curry Howard (Why proofs == Why programs)
- Formal verification
- Enable `Type Precision` (focus of this talk)
 - (proofs replace unsafe coercion)

```
data List a =

Empty
Cons a (List a)

data Vect (n :: Nat) a where

Empty :: Vect 0 a

Cons :: a -> Vect n a -> Vect (1 + n) a
```

```
id_sanity :: Vect (n + 0) a -> Vect n a
id_sanity = unsafeCoerce -- ghrrr 🙁
```

Curry-Howard to MaybeB

(code)

https://github.com/rpeszek/present-proofs-lc19/blob/master/src/Present/AnIntro.hs

https://github.com/rpeszek/present-proofs-lc19/blob/master/src/Present/MaybeB.hs

Problem in Paradise - Questions

Should type checker know basic algebra?

```
a | True <==> True (Bool algebra)
a + b == b + a (Nat, Int, Float ... algebra) ...
```

Answers:

- Dependently Types Langs: **No** (③ *Programmers supply proofs*)
- Refinement Types / LiquidHaskell: **Yes** (□♂ *SMT solver does the work*)

Type Equality

```
data a :~: b where
Refl :: a :~: a
```

```
test1 = Refl :: 5 :~: 5 -- GOOD

test10 = Refl :: 4 :~: 5 -- ERR

test2 = Refl :: 2 + 3 :~: 3 + 2 -- GOOD

test20 :: SNat n1 -> SNat n2 -> n1 + n2 :~: n2 + n1

test20 _ = Refl -- ERR
```

(base) Data.Type.Equality

Example Combinators

- library over op semantics
- "pattern-matching on a variable of type (a :~: b) produces a proof that a ~ b" - haddoc

```
sym :: (a :~: b) -> (b :~: a)

trans :: (a :~: b) -> (b :~: c) -> (a :~: c)

apply :: (f :~: g) -> (a :~: b) -> (f a :~: g b)

inner :: (f a :~: g b) -> (a :~: b)

castWith :: (a :~: b) -> a -> b

gcastWith :: (a :~: b) -> ((a ~ b) => r) -> r

gcastWith Refl x = x
```

(base) Data. Type. Equality

Proofs - Bool Algebra

(code)

https://github.com/rpeszek/present-proofs-lc19/blob/master/src/Present/ProofsBoolAlg.hs

Better Bool ... Decidability

(code)



- Programming contradictions
 - -XEmptyCase
- Forced Type Precision
- Programs are proofs!

Proofs - (Performance)

(code)

unsafeCoerce replacements

- rewrite rules
- proveFast combinators
- CPP

https://github.com/rpeszek/present-proofs-lc19/blob/master/src/Present/ProofsNatAlg.hs

Proofs - working with TypeLits

(code)

- hidden implementation (good and bad)
- Peano is slow

https://github.com/rpeszek/present-proofs-lc19/blob/master/src/Present/WorkingWithTypeLits.hs

Maintenance

```
id_sanity :: SNat n -> Vect (n + 0) a -> Vect n a
id_sanity n = case plusZeroRightNeutral n of Refl -> id
```

- (+) is just a function lifted to type level
- (+) implementation details get lifted too
- plusZeroRightNeutral depends on (+) implementation
- plusZeroRightNeutral may have not been needed, depends on impl of (+)
- *Opinionated* implementation agnostic code:
 - Create exhaustive library of propositions about (:~:)
 - Type implicit base propositions (like *True* :~: *True* | | b) explicitly
 - Refactor and fix proofs at the same time
 - Include implicit propositions in client code when relevant (lint tool needed)
- Tactics ...

Type Precision (call-side safety)

```
(!!) :: [a] -> Int -> a
                                                                        a lie ...
safeGet :: [a] -> Int -> Maybe a
                                                                        questionable
                                                                        improvement
-- Liquid:
                                                                        refinements
{-@ (!!) :: x: [a] -> {i:Nat | i < len x} -> a @-}
(!!) :: [a] -> Int -> a
-- Dep Typed:
(!!) :: Vect n a -> Fin n -> a
(!!!) :: Vect n a -> SNat m -> MaybeB (m < n) a
(!!!!): (xs:List a) -> (n:Nat) -> {auto ok: InBounds n xs} -> a
                                                                    ... still imprecise *
isElem: DecEq a => (x: a) -> (xs: List a) -> Dec (Elem x xs)
                                                                    not in the same boat *
```

https://github.com/rpeszek/present-proofs-lc19/tree/master/src/Wrapup

Type Precision (implementation safety)

(code)

Some Learning/References

- intro books
 - Type Driven Development in Idris great book
 - https://github.com/rpeszek/ldrisTddNotes/wiki
 - TAPL great book (not really dep types but still)
 - Programming foundation books (penn/Pierce et al, Wadler)
- Haskell projects (with reading references)
 - singletons
 - equational-reasoning-in-haskell
 - liquidhaskell
- blogs
 - <u>blog.jle.im</u> (Justin Le)
 - typesandkinds (Richard Eisenberg)
- youtube
 - <u>Introduction to Agda</u> series by Daniel Peebles published by Edward Kmett