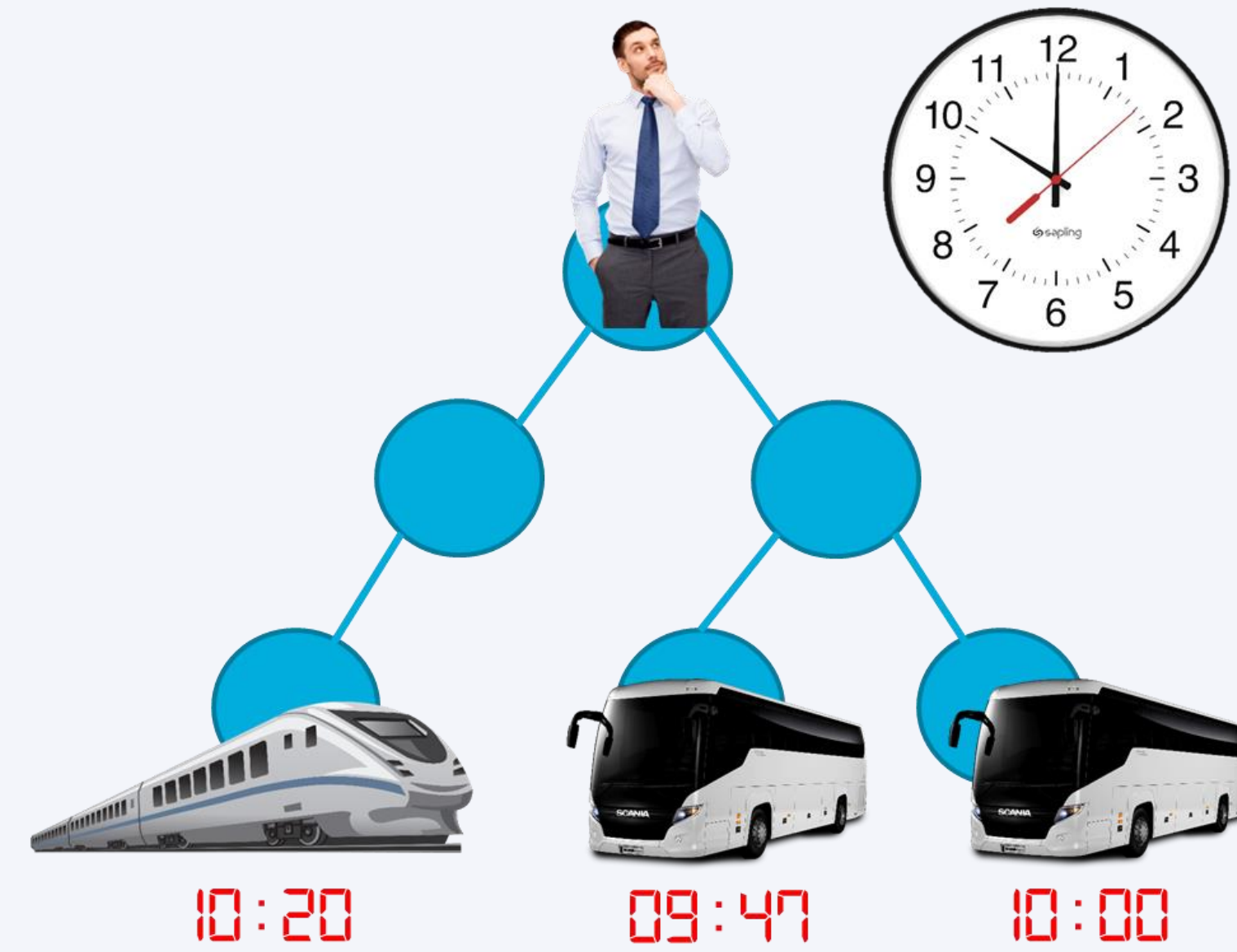


# Situated Temporal Planning Using Deadline-aware Metareasoning

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## 1 Situated Temporal Planning [Cashmore et al. 2018]



**Example: planning a route:**

- 'take 10:00 bus' action expires at 10:00 subtree of plans becomes invalid; consider only if sufficient time to complete plan
- exploring 'take 9:47 bus' action can invalidate 10:00 action; searching under multiple nodes means less time for each
- a plan expiration time and cost are uncertain until the plan is complete but completion effort also uncertain

**which plans to explore?**

## 2 Allocating Effort when Actions Expire (AE2) [Shperberg et al. 2019]

$n$  partial plans/nodes/processes to share CPU time. Given for each process  $i$ :

- **effort CDF:**  $M_i(t)$  = probability that  $i$  requires CPU time  $\leq t$
- **success probability:**  $P_i$  = probability that  $i$  finds a solution (without considering time found)
- **deadline CDF:**  $D_i(t)$  = probability that  $i$  expires before  $t$  (clock wall time). Not certain until solution is complete

Find a **schedule** for the processes that **maximizes probability** of finding a solution that is **still valid** when found.

**AE2 is NP-Hard.** The (AE)<sup>2</sup> can be modeled as an MDP, but the **state space** is exponential in  $n$ .

**Known deadlines:**

Linear contiguous policies (LCP):  
(1, 1, 1, 3, 3, 2, 2, 2, ...)  
No feedback and allocation for all process are contiguously

With known deadlines, there exists an LCP that is an optimal solution.

## 3 Known deadlines: pseudo-polynomial algorithm

- $LPF_i(t_0, t_u)$ : the log probability of failure for the process to find a timely solution within  $t_u$  processing time units, starting at time  $t_0$
- -LPF is equivalent to utility

**DP algorithm:**

1. Sort processes in a non-decreasing order of deadlines
2. Compute the optimal utility of scheduling processes **I** thought **n** starting from time **t**:

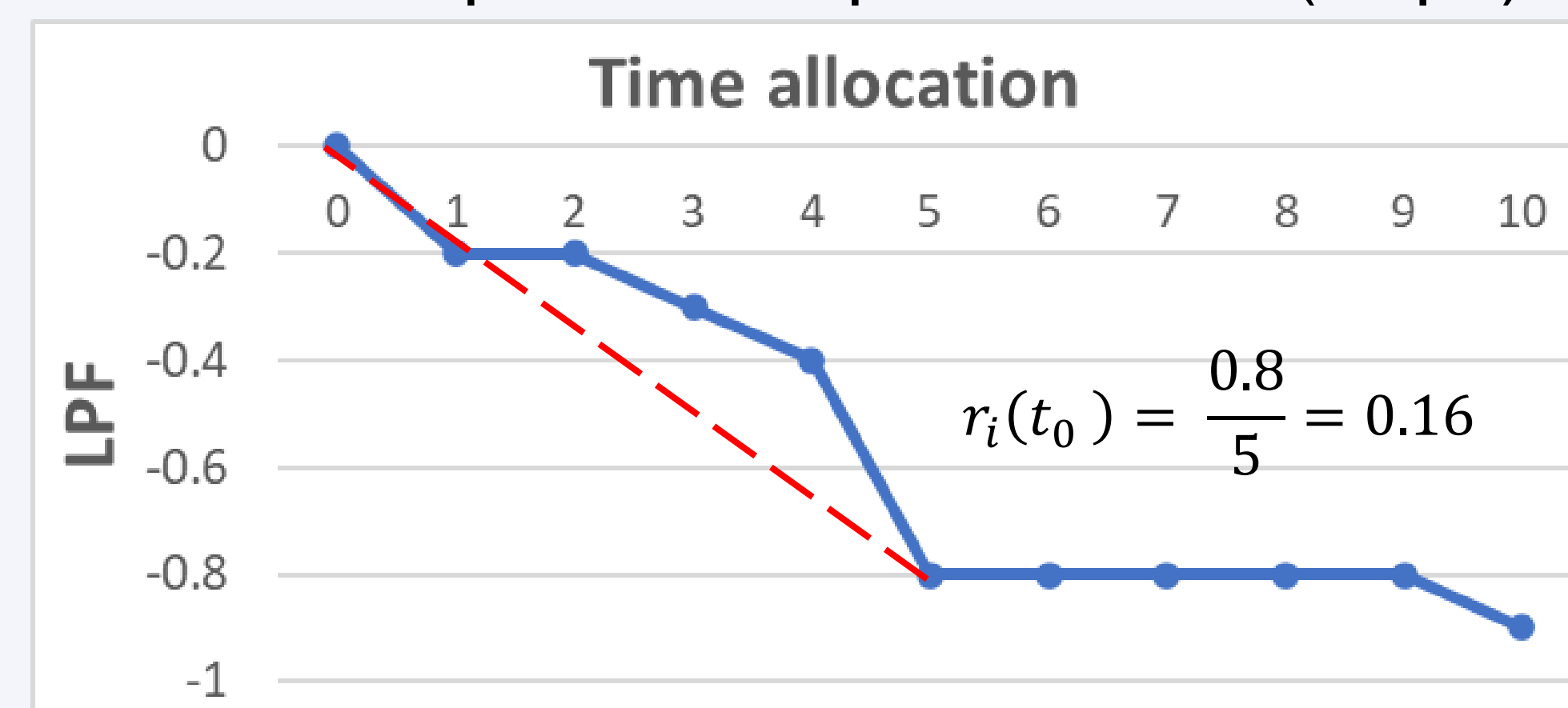
$$OPT(t, l) = \max_{0 \leq j \leq d_i - t} (OPT(t + j, l + 1) - LPF_i(t, j))$$

3. Return  $OPT(0, 1)$  as the utility of the optimal policy.

- runs in time polynomial in  $n$  and  $\max_i d_i$
- Usable when a solution needs to be found before a known (common to all processes) timeout, e.g., Algorithm Portfolio

## 4 Unknown deadlines: greedy schemes

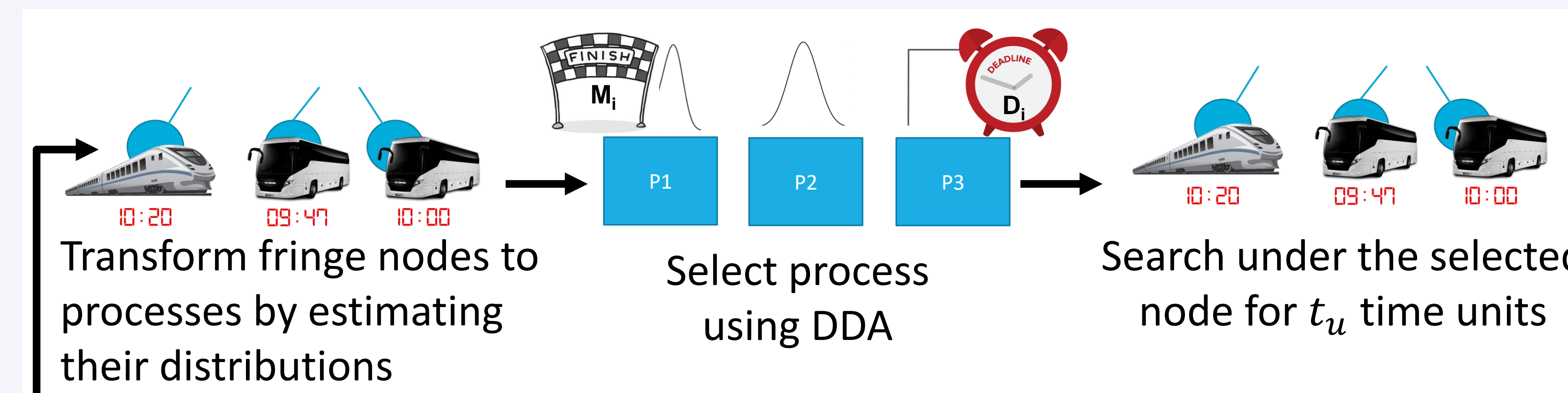
- $r_i(t_0)$ : the best rate of improvement per time unit (slope) for process  $i$ :



- **P-Greedy** [Shperberg et al. 2019]: Allocate  $t_u$  time units to the process that maximizes:  $Q_i(t_0) = r_i(t_0) + \frac{\alpha}{E(D_i)}$
- **Delay-Damage Aware (DDA)**: Allocate  $t_u$  time units to the process that maximizes:  $Q_i(t_0) = r_i(t_0) - \gamma \cdot r_i(t_0 + t_u)$

DDA uses a more methodological way to decide which process is damaged the most by delay

## 5 Search -> Metareasoning -> Search



**Estimating the distributions:**

- OPTIC uses a temporal relaxed planning graph (TRPG) (Coles et al. 2010) to estimate **distance-to-go**  $d(s)$  and **latest start time**  $lst(s)$  for every state.
- $d(s)$  is used for estimating distribution of remaining search time ( $M_i$ )
- $lst(s)$  is used for estimating distribution for deadlines ( $D_i$ )

## 6 Empirical Evaluation - results

Domain	baseline		DDA		DDA (dom tuned)	
airport	19.0	(19–19)	20.0	(20–20)	20.5	(19–21)
pw-nt	4.0	(3–4)	4.0	(3–5)	3.9	(3–5)
rc1 1	37.7	(37–40)	73.7	(53–92)	83.9	(59–99)
rc1 2	1.0	(1–1)	4.0	(2–23)	2.7	(0–13)
sat cmplx	5.0	(5–5)	5.0	(5–5)	3.8	(2–5)
sat tw	5.0	(5–5)	5.0	(5–5)	3.6	(3–5)
trucks	6.0	(6–6)	6.9	(6–9)	5.7	(5–8)
turtlebot	14.0	(14–14)	12.5	(10–13)	13.0	(13–13)
umts-flaw	4.1	(4–5)	5.1	(5–6)	5.0	(5–5)
umts	48.0	(48–48)	45.5	(42–49)	45.7	(44–49)
TOTAL	143.8	(142–147)	181.7	(151–227)	187.7	(153–223)

## 7 Conclusion

Considering planning time is **hard**, but ignoring it is not the solution!

