CS/ECE 473 Fall 2020 Homework 6 Problem 1 Ryan Prendergast (ryanp4) Noah Watson (nwatson3) Lawson Probasco (lawsonp2)

**Solution:** In order to prove there is a |f'| - |f| valued (s, t)-flow in  $G_f$ , it is first useful to bound these quantities. Starting with a bound on the flow of  $G_f$ , which will be referred to abstractly as  $f_{G_f}$ , the value of this flow will be the sum of the outbound flow from s in  $G_f$  minus the sum of the inbound flow to s, e.g.  $|f_{G_f}| = \Sigma_w f(s \to w) - \Sigma_u f(u \to s)$ .

$$|f_{G_f}| = \Sigma_w f(s \to w) - \Sigma_u f(u \to s) \tag{1}$$

$$\Sigma_{w} f(s \to w) \le \Sigma_{w} c(s \to w) - |f| \tag{2}$$

$$\Sigma_{u} f(u \to s) \le -|f| \tag{3}$$

$$-|f| \le |f_{G_f}| \le \sum_{w} c(s \to w) - |f| \tag{4}$$

The justification for (2) is by the definition of residual graphs - namely  $c_f(s \to w) = c(s \to w) - f(s \to w)$ . The justification for (3) is also by the definition of residual graphs, where all  $u \to s \in E_f$  are backtrack edges such that  $c_f(u \to s) = f(s \to u)$ . Note for this equation I dropped the value of outbound edges, meaning I minimized full quantity in the case where none of the residual edges have a flow along them. Now we proceed to the bounds on the quantity |f'| - |f| below.

$$0 \le |f'| \le \Sigma_w c(s \to w) \tag{5}$$

$$-|f| \le |f'| - |f| \le \Sigma_w c(s \to w) \tag{6}$$

The justification for (1) above is just satisfying the capacity constraint, and (2) just plugs in the bounds of |f'| into |f'| - |f|. Seeing clearly now that the bounds of the quantities |f'| - |f| and  $f_{G_f}$  are the same, it is sufficient to prove that the upper and lower bound values for  $f_{G_f}$  are attainable for all the values in the range to be attainable. Starting with the easy case, the quantity -|f| is just flow in  $G_f$  that saturates all backtracking edges  $u \to s$ . This is clearly attainable because for any edge  $u \to s$  with capacity  $c_f$ , there is a set of vertices s.t.  $\Sigma_a a \to u = c_f$  since the original flow f satisfied the conservation constraint. This identity can be applied outwards to those vertices a and connected edges all the way to b. For the upper bound b0 constraint can be kept to reach this quantity. Chiefly, the flow that saturates all edges b1 w in b2 produces exactly this sum. It should be noted that the equality can most certainly be satisfied up to the sum |f'|, because if there was a bottleneck in some edge such that its flow in b2 and its value in b3 could not sum to its value b4, b5 could not exist because it would be impossible that b6 doesn't run into a bottleneck at the same edge.

Using the above knowledge, it can be said that f is a maximum flow iff there is no s-t path in  $G_f$  in the following way. Consider the case where f is a maximum flow but there IS an s-t path in  $G_f$ . In thie case, this s-t path means there exist  $|f_{G_f}|$  with positive flow, or in other words, some f' such that |f'|-|f|>0. But if that's the case, the |f'|>|f|, meaning f was not the maximum flow and we have a contradiction. Similarly, assuming there is some f that is not the maximum flow for which  $G_f$  has no s-t path, it is clear that there is no such positive flow in  $G_f$ , or in other words no f' such that |f'|-|f|>0. However if this is the case, then clearly |f| is the maximum flow and we have another contradiction.