

Streaming Models

FOCS, Spring 2022

1. Dataflow Networks

Streaming models

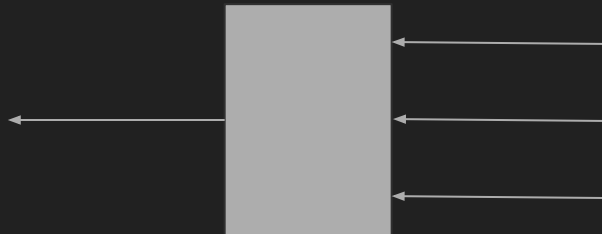
Working with infinitely streaming data

- multiple input streams
- single output stream

Process and create output stream as input comes in

- Ideally don't buffer

What goes in the box?



Dataflow networks

Dataflow networks take streams of values as inputs and produce streams of values as outputs

- type of values depends on the kind of network developed
 - floating point for approximation algorithms
 - images for streaming movies
- sequential components connected by buffered communication channels
- model assumes an underlying sequential language

Primitive components

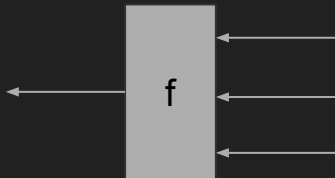
Constant k : produces an infinite stream of k



Primitive components

map f: transforms one or more streams by applying f to the inputs

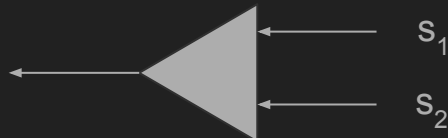
- blocks until all input streams have at least one value)
- transformation f written in underlying sequential language
- transformation f holds no state



Primitive components

followed by : produces a stream from the first element of s_1 followed by everything from s_2

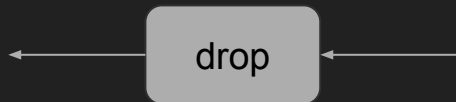
- blocks until an element of s_1 arrives
- then simply forwards values that arrive on s_2



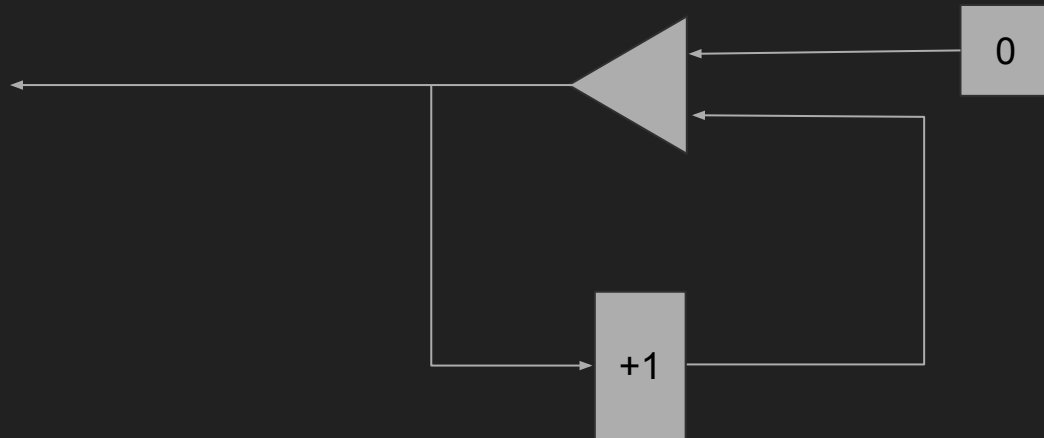
Primitive components

drop : produces a stream from the input stream by "dropping" the first element of the stream

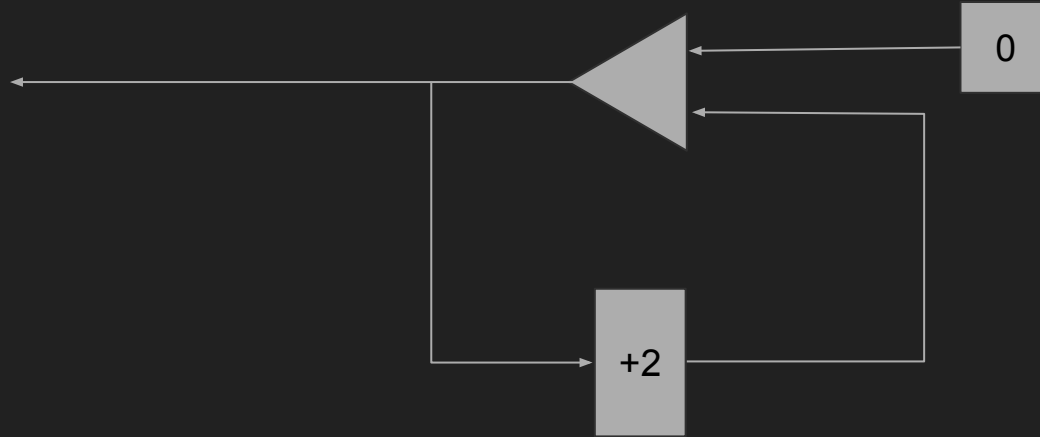
- input *a b c d e f ...* output *b c d e f ...*
- discards the first element that arrives (produce no output)
- then simply forwards everything that arrives to its output



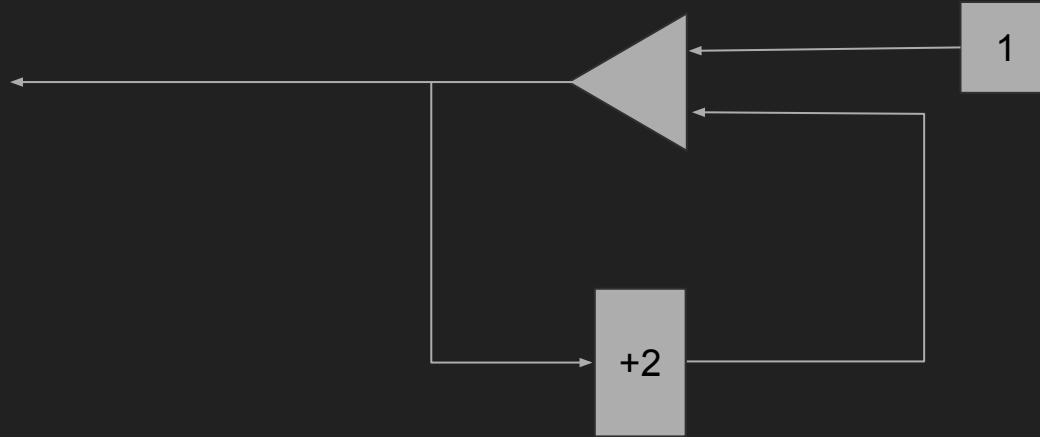
Sequences: nats



Sequences: evens



Sequences: odds



Sequences: odds



Sequences: triangular numbers

Want to create 0, 1, 3, 6, 10, 15, 21, ...

Observe:

$$0 + 1 = 1$$

$$1 + 2 = 3$$

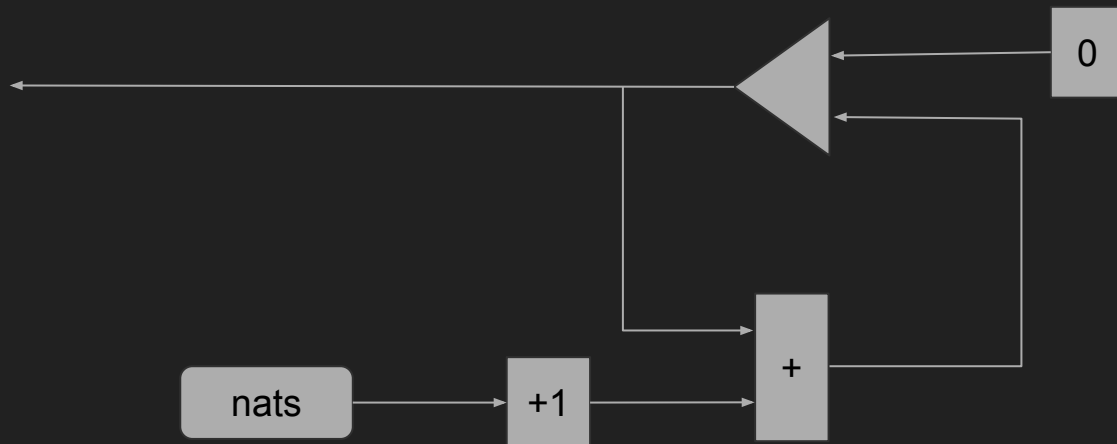
$$3 + 3 = 6$$

$$6 + 4 = 10$$

$$10 + 5 = 15$$

$$15 + 6 = 21$$

...



Sequences: square numbers

Want to create 0, 1, 4, 9, 16, 25, 36, ...

Observe:

$$0 + 1 = 1$$

$$1 + 3 = 4$$

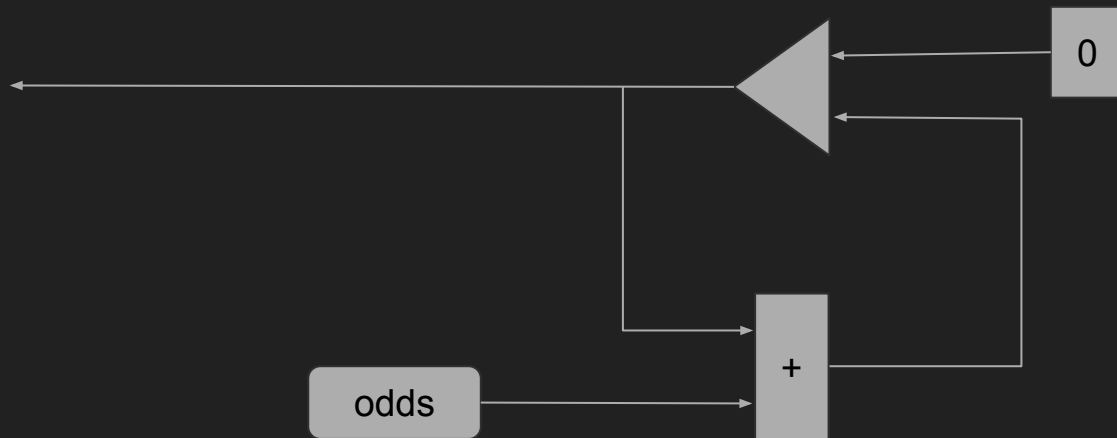
$$4 + 5 = 9$$

$$9 + 7 = 16$$

$$16 + 9 = 25$$

$$25 + 11 = 36$$

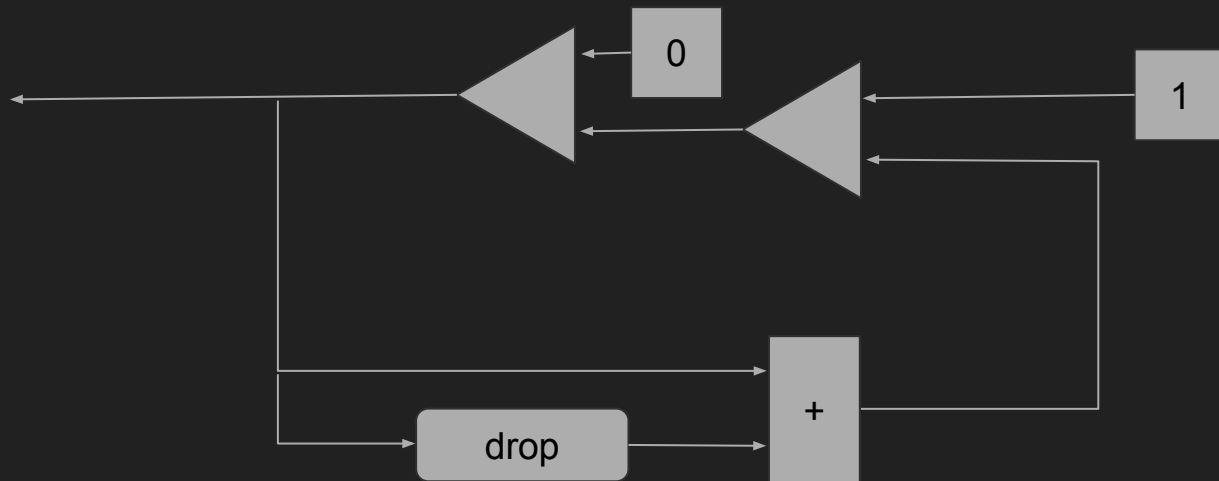
...



Sequences: Fibonacci numbers

Want to create 0, 1, 1, 2, 3, 5, 8, 13, 21, ...

Each number in the sequence is the sum of the previous two

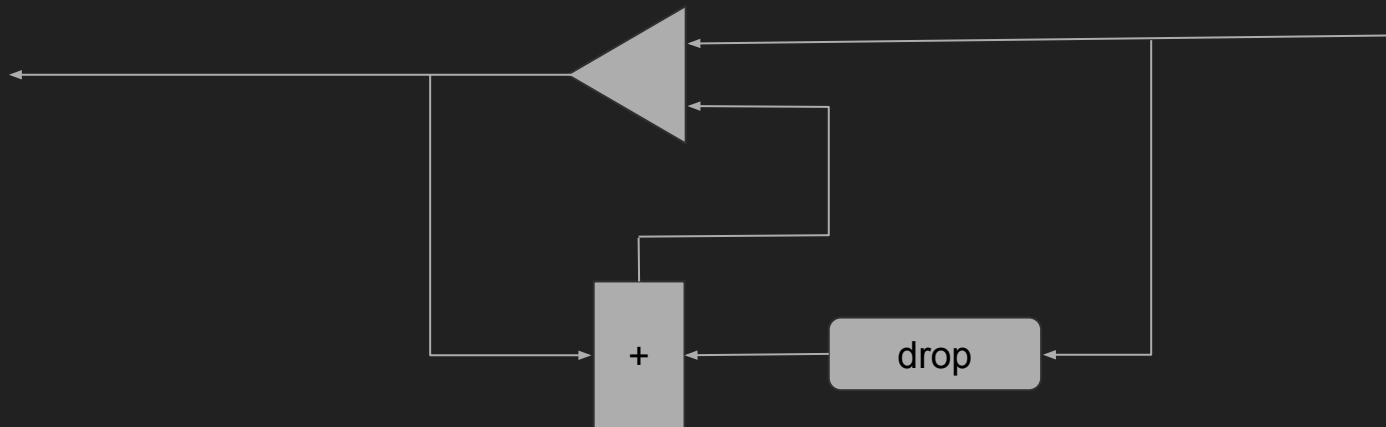


Transformation: partial sums

Input:	a	b	c	d	e	f	...
Output:	a	a+b	a+b+c	a+b+c+d	a+b+c+d+e	a+b+c+d+e+f	...

Transformation: partial sums

Input:	a	b	c	d	e	f	...
Output:	a	a+b	a+b+c	a+b+c+d	a+b+c+d+e	a+b+c+d+e+f	...



Transformations: running averages

Input:	a	b	c	d	e	...
Output:	$a/1$	$(a+b)/2$	$(a+b+c)/3$	$(a+b+c+d)/4$	$(a+b+c+d+e)/5$...



Definitions

A **dataflow network** with inputs I and outputs O is a finite network of components where:

1. every component is either a primitive component or an already defined dataflow network
2. every component's input is either in I or connected to exactly one output
3. every component's output can be connected to zero or more inputs and can also appear in O

Main theorem

A **cycle** in a dataflow network is a path from the output of some component back to an input of the same component by following links in the network

Theorem: *If every cycle in a dataflow network goes through the lower input of at least one "followed by" primitive component, then the dataflow network computes a function from its input streams to its output stream*

2. Stream Programming

Dataflow networks

Dataflow networks are fundamentally a graphical model

We'll see how to program graphical models next homework

Another way to program over streaming data is to consider a stream to be an **infinite list** and write list processing functions as usual

Infinite lists

A list in OCaml or most programming languages is a finite structure

You cannot construct an infinite list in Ocaml

- You'd spend all your time building it

Yet, some languages lets you define infinite lists, the equivalent of writing:

```
let rec from k = k :: (from (k + 1))
```

where `from 10` would be the list `[10, 11, 12, 13, 14, 15, 16, ...]`

Lazy evaluation

One way to do it is to use **lazy evaluation**

- only evaluate an expression if you need its value

Consider the function

```
let test a b = a
```

It does not use the second argument at all

In a language with lazy evaluation, `test 10 (g 20)` returns `10` *even if* `(g 20)` *is an infinite loop*

Haskell

A functional programming language like OCaml but which uses lazy evaluation

- An expression is evaluated only if its value is needed

```
square x = x * x
```

```
length [] = 0
```

```
length (h : t) = 1 + length t
```

```
squares lst = map (\a -> a * a) lst
```

Lists in Haskell

Like Ocaml, if L is a list, $a : L$ is a list with first element a and rest of the list L

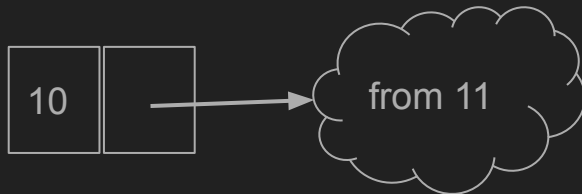
$10 : [1, 2, 3] \rightarrow [10, 1, 2, 3]$

But $a : L$ is **lazy** in its second argument

- it does not evaluate L until L is needed
(e.g., somebody trying to access the elements of L)

With $\text{from } k = k : (\text{from } (k + 1))$

$\text{from } 10$ is perfectly well defined



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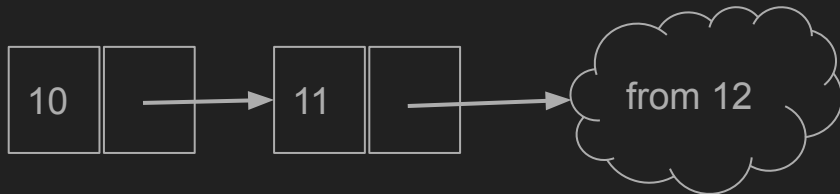
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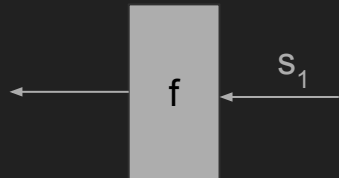
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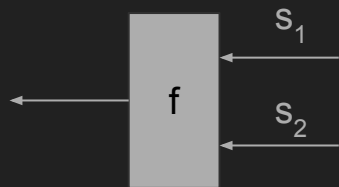
Primitive components



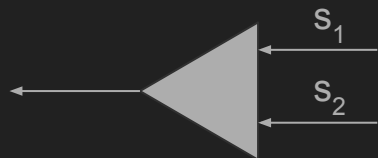
$\text{cst } k$



$\text{map } f \ s_1$



$\text{map2 } f \ s_1 \ s_2$



$\text{fby } s_1 \ s_2$

$\text{cst } a = a : (\text{cst } a)$

$\text{map } f \ (h : t) = (f \ h) : (\text{map } f \ t)$

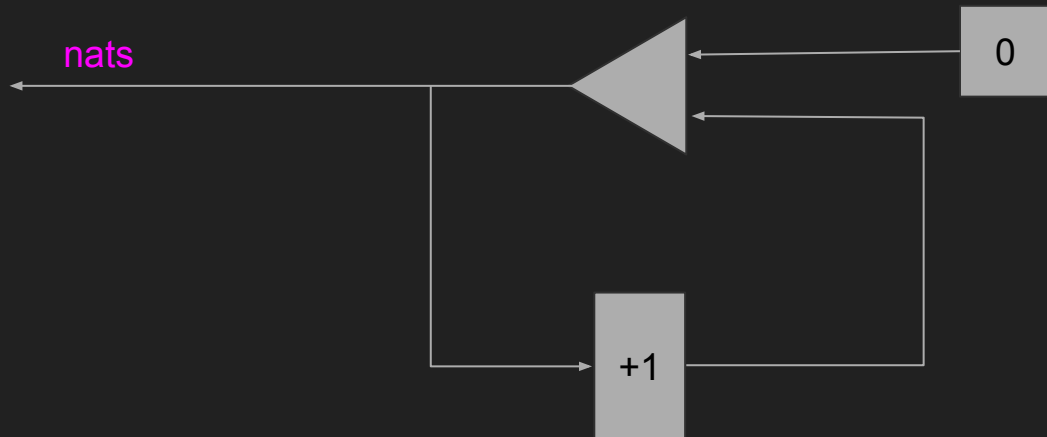
$\text{map2 } f \ (h_1 : t_1) \ (h_2 : t_2) = (f \ h_1 \ h_2) : (\text{map2 } f \ t_1 \ t_2)$

$\text{fby } (h : t) \ s = h : s$

$\text{tail } (h : t) = t$



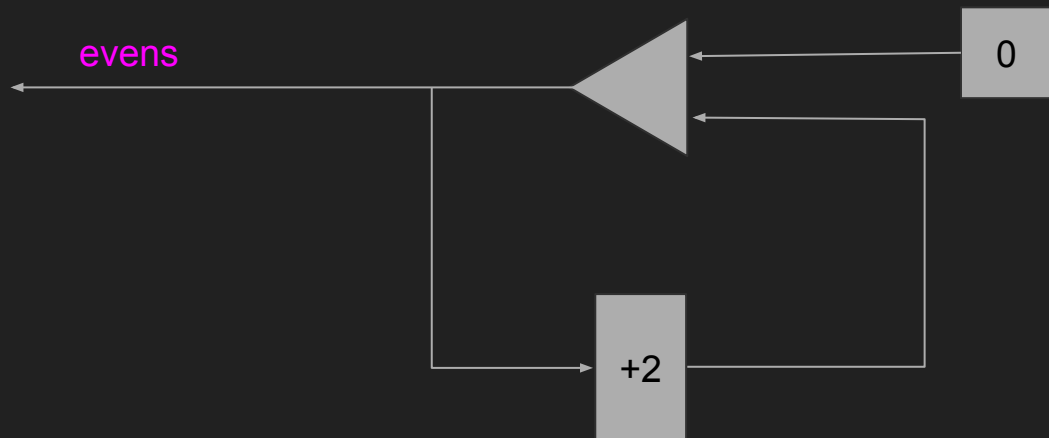
Example: nats



`plus1 s = map (\a -> a + 1) s`

`nats = fby (cst 0) (plus1 nats)`

Example: evens



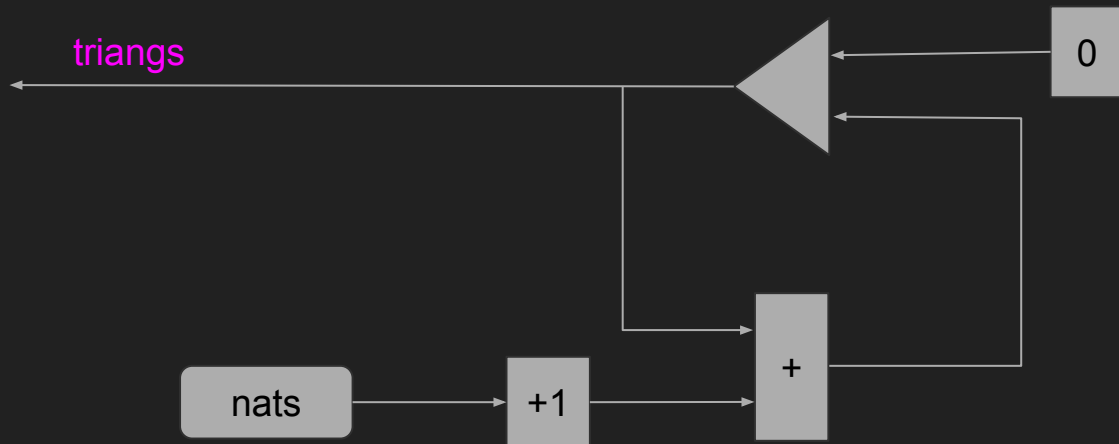
`evens = fby (cst 0) (plus1 (plus 1 evens))`

Example: odds



odds = plus1 evens

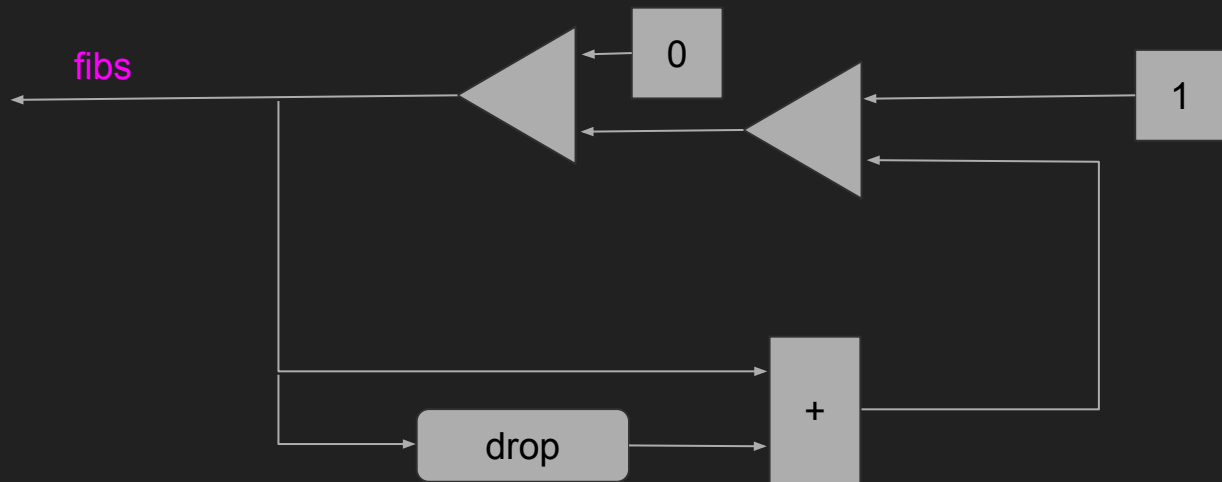
Example: triangular numbers



$\text{plus } s_1 \ s_2 = \text{map2 } (\backslash a \rightarrow \backslash b \rightarrow a + b) \ s_1 \ s_2$

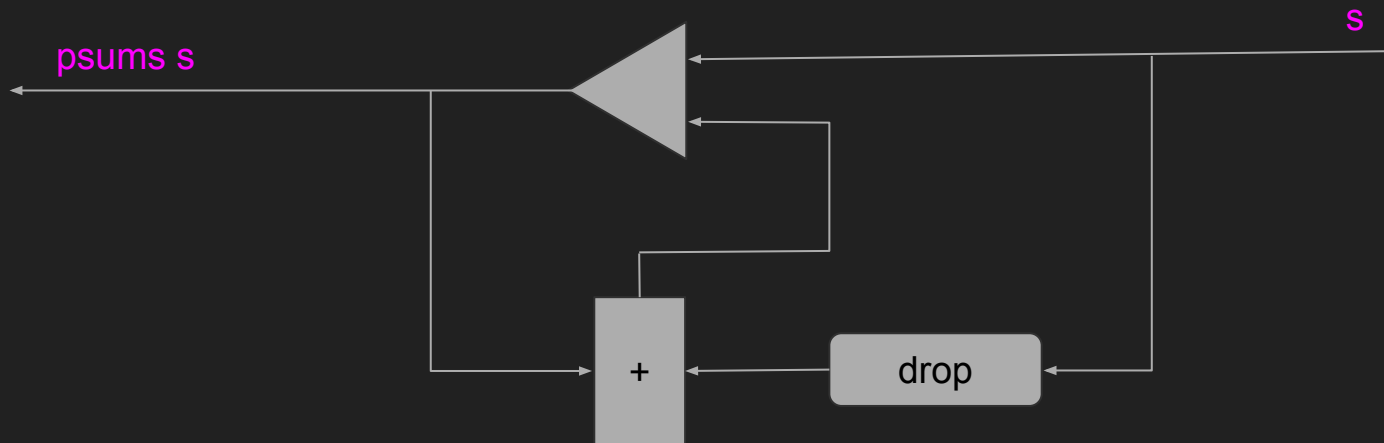
$\text{triangs} = \text{fby } (\text{cst } 0) \ (\text{plus triangs } (\text{plus1 nats}))$

Example: Fibonacci numbers



`fibs = fby (cst 0) (fby (cst 1) (plus fibs (tail fibs)))`

Example: partial sums



$\text{psums } s = \text{fby } s \text{ (plus (psums } s) \text{ (tail } s))}$

Recursive dataflow networks

A dataflow network that uses a copy of itself as a component

Sieve of Eratosthenes - how to compute the stream of prime numbers:

- sieving a stream: take a stream of values, keep the first value, and sieve the rest of the stream *after* removing all multiples of the first value
- sieving the natural numbers starting from 2 yields the prime numbers

```
divides c x = (mod x c == 0)
```

```
sieve s = fby s (sieve (filter (\a -> not (divides (head s) a)) (tail s)))
```

```
primes = sieve (plus1 (plus1 nats))
```

Infinite lists without lazy evaluation

Even if you do not have lazy evaluation, you can still program with infinite streams

- OCaml: a stream is a pair of an element and a **thunk**
 - a thunk is a function of no arguments that returns a stream
 - until the thunk is called, it's like an "unevaluated stream"

```
type 'a stream = St of 'a * (unit -> 'a stream)
```

```
let map (f: 'a -> 'b) (s: 'a stream): 'b stream =  
  match s with (h, th) -> St (f h, fun () -> map f (th ()))
```

```
let nats (): int stream =  
  St (0, fun () -> map (fun x -> x + 1) (nats ()))
```