

# 1. Dataflow Networks

# Streaming models

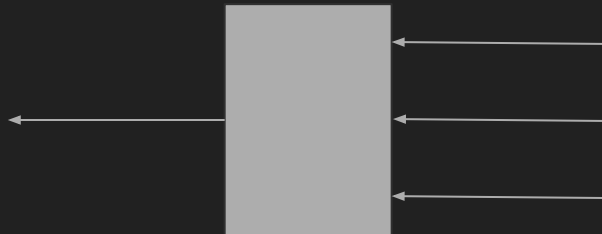
Working with infinitely streaming data

- multiple input streams
- single output stream

Process and create output stream as input comes in

- Ideally don't buffer

What goes in the box?



# Dataflow networks

Dataflow networks take streams of values as inputs and produce streams of values as outputs

- type of values depends on the kind of network developed
  - floating point for approximation algorithms
  - images for streaming movies
- sequential components connected by buffered communication channels
- model assumes an underlying sequential language

# Primitive components

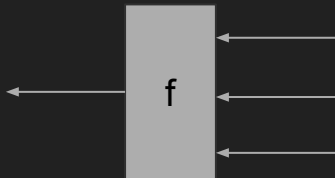
Constant  $k$ : produces an infinite stream of  $k$



# Primitive components

**map f:** transforms one or more streams by applying  $f$  to the inputs

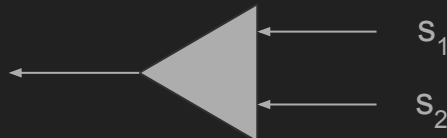
- blocks until all input streams have at least one value)
- transformation  $f$  written in underlying sequential language
- transformation  $f$  holds no state



# Primitive components

**followed by** : produces a stream from the first element of  $s_1$  followed by everything from  $s_2$

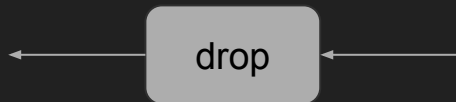
- blocks until an element of  $s_1$  arrives
- then simply forwards values that arrive on  $s_2$



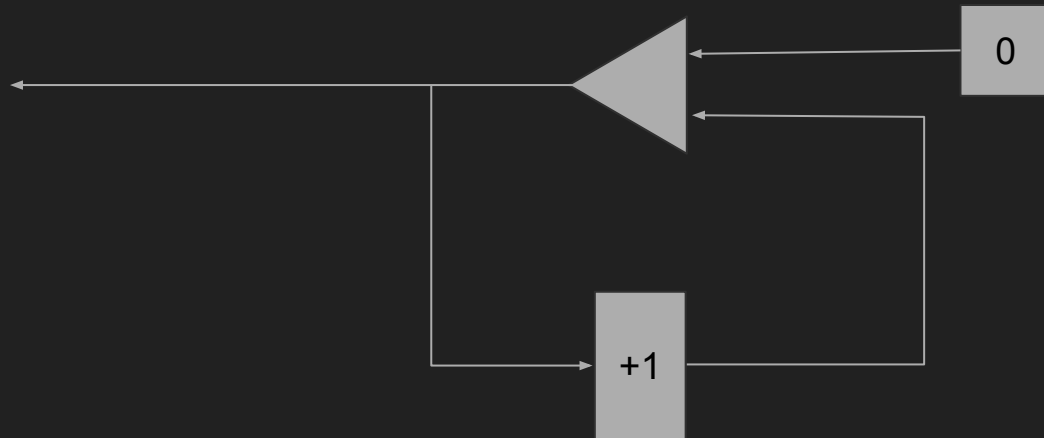
# Primitive components

**drop** : produces a stream from the input stream by "dropping" the first element of the stream

- input *a b c d e f ...* output *b c d e f ...*
- discards the first element that arrives (produce no output)
- then simply forwards everything that arrives to its output

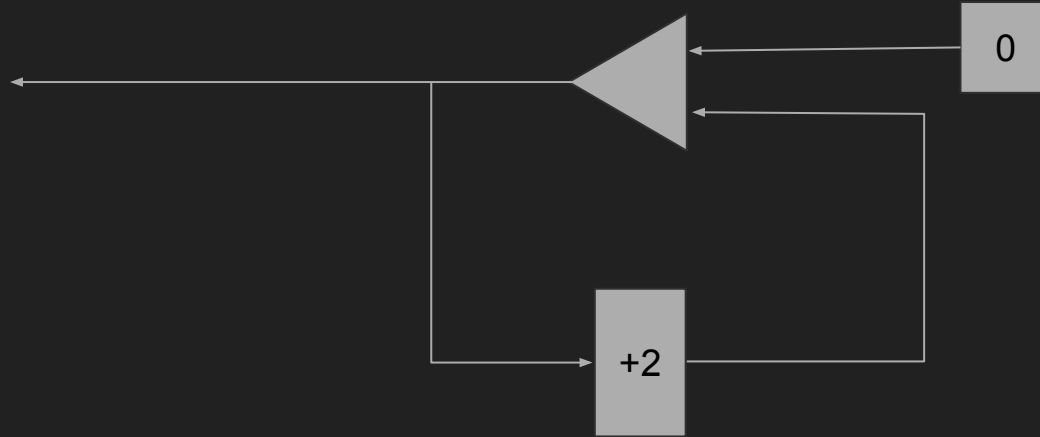


# Sequences: nats

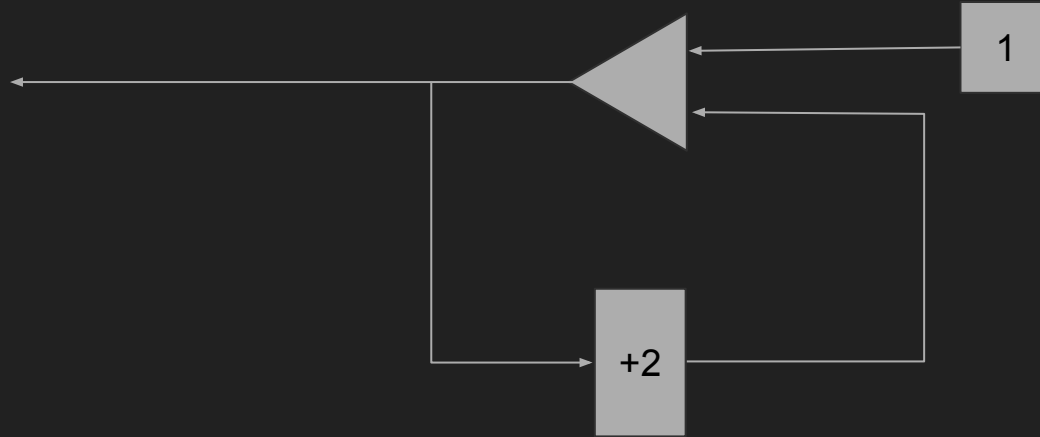




# Sequences: evens



# Sequences: odds



# Sequences: odds



# Sequences: triangular numbers

Want to create 0, 1, 3, 6, 10, 15, 21, ...

Observe:

$$0 + 1 = 1$$

$$1 + 2 = 3$$

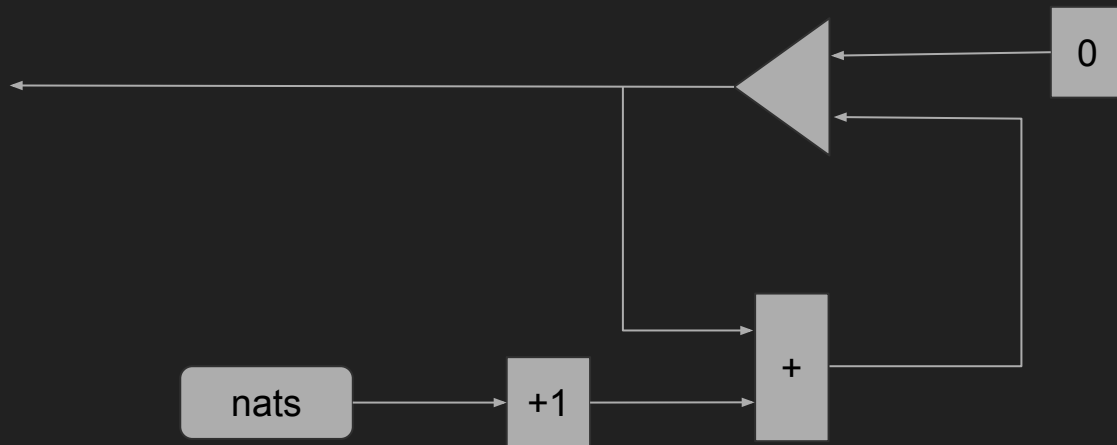
$$3 + 3 = 6$$

$$6 + 4 = 10$$

$$10 + 5 = 15$$

$$15 + 6 = 21$$

...



# Sequences: square numbers

Want to create 0, 1, 4, 9, 16, 25, 36, ...

Observe:

$$0 + 1 = 1$$

$$1 + 3 = 4$$

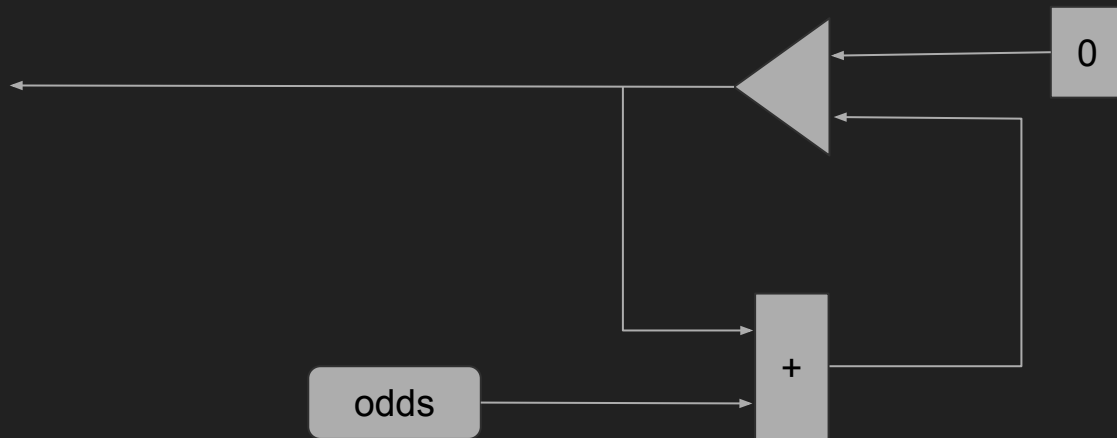
$$4 + 5 = 9$$

$$9 + 7 = 16$$

$$16 + 9 = 25$$

$$25 + 11 = 36$$

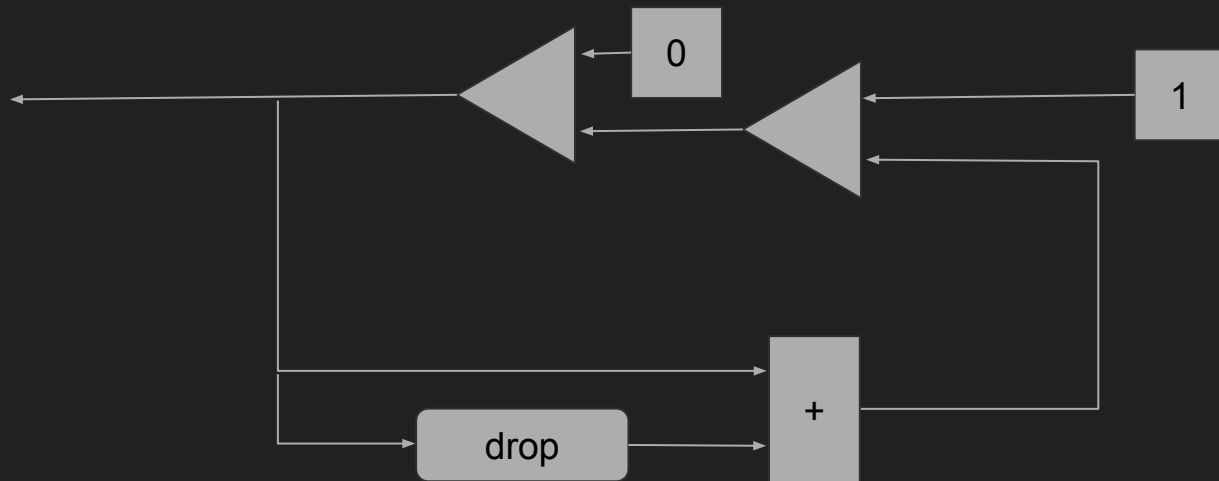
...



# Sequences: Fibonacci numbers

Want to create 0, 1, 1, 2, 3, 5, 8, 13, 21, ...

Each number in the sequence is the sum of the previous two

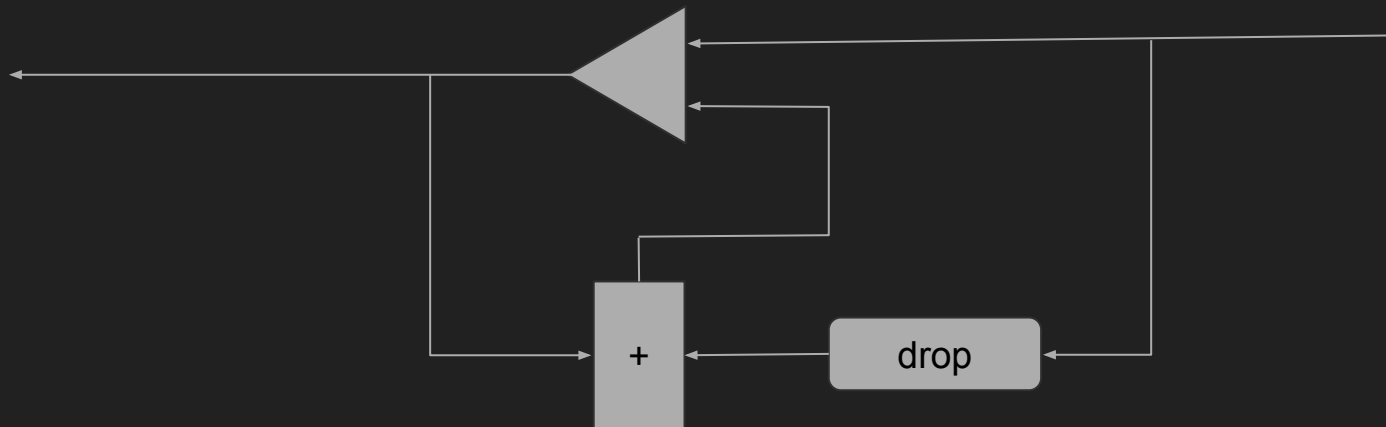


# Transformation: partial sums

Input:	a	b	c	d	e	f	...
Output:	a	a+b	a+b+c	a+b+c+d	a+b+c+d+e	a+b+c+d+e+f	...

# Transformation: partial sums

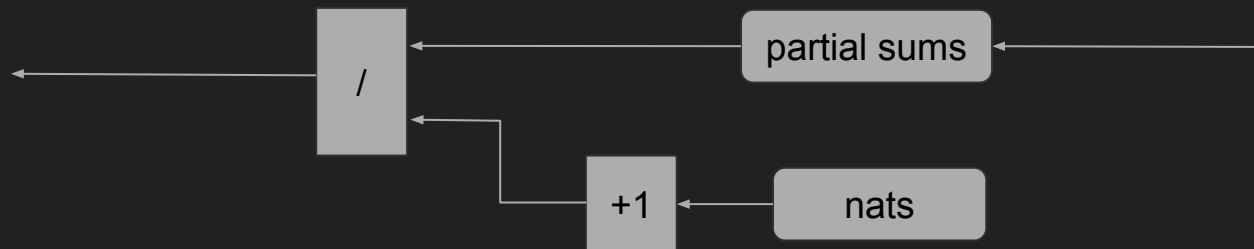
Input:	a	b	c	d	e	f	...
Output:	a	a+b	a+b+c	a+b+c+d	a+b+c+d+e	a+b+c+d+e+f	...





# Transformations: running averages

Input:	a	b	c	d	e	...
Output:	$a/1$	$(a+b)/2$	$(a+b+c)/3$	$(a+b+c+d)/4$	$(a+b+c+d+e)/5$	...



# Definitions

A **dataflow network** with inputs  $I$  and outputs  $O$  is a finite network of components where:

1. every component is either a primitive component or an already defined dataflow network
2. every component's input is either in  $I$  or connected to exactly one output
3. every component's output can be connected to zero or more inputs and can also appear in  $O$

# Main theorem

A **cycle** in a dataflow network is a path from the output of some component back to an input of the same component by following links in the network

**Theorem:** *If every cycle in a dataflow network goes through the lower input of at least one "followed by" primitive component, then the dataflow network computes a function from its input streams to its output stream*

## 2. Stream Programming

# From dataflow networks to streams programming

Dataflow networks are fundamentally a graphical model

We'll see how to program graphical models next homework

Another way to program over streaming data is to consider a stream to be an **infinite list** and write list processing functions as usual

# Infinite lists

Most languages do not allowing you to define infinite lists:

- You'd spend all your time building it

Yet, Haskell (and some other languages) lets you define infinite lists:

```
from k = k : (from (k + 1))
```

where `from 10` would be the list `[10, 11, 12, 13, 14, 15, 16, ...]`

It does so using *lazy evaluation*

# Lazy evaluation

Haskell uses lazy evaluation everywhere (but relevant mostly at function calls)

It only evaluates an expression if it *need* its value

Consider the function

```
test a b = a
```

The body does not use the second argument at all

In Haskell, `test 10 (loop 0)` returns `10` even if *(Loop 0)* is an infinite loop

# Lists in Haskell

If  $L$  is a list,  $a : L$  is a list with first element  $a$  and rest of the list  $L$

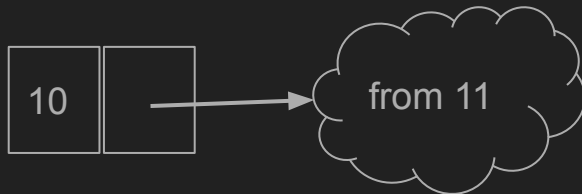
$10 : [1, 2, 3] \rightarrow [10, 1, 2, 3]$

But  $a : L$  is **lazy** in its second argument

- it does not evaluate  $L$  until  $L$  is needed  
(e.g., somebody trying to access the elements of  $L$ )

With  $\text{from } k = k : (\text{from } (k + 1))$

$\text{from } 10$  is perfectly well defined





# Lists in Haskell

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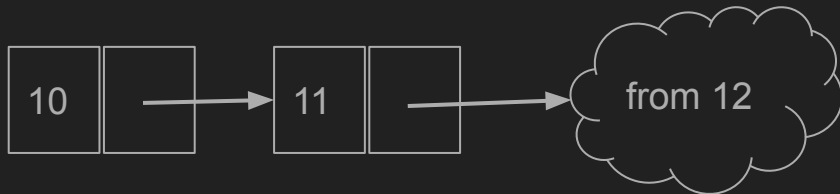
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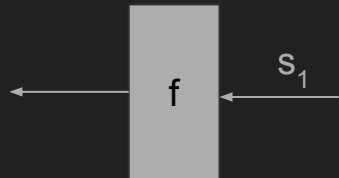
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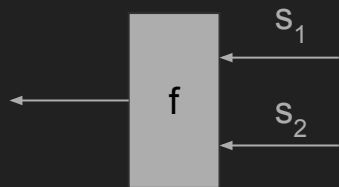
# Primitive components



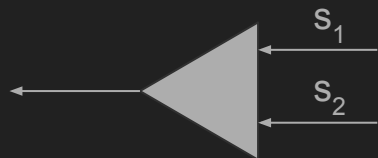
$\text{cst } k$



$\text{map } f \ s_1$



$\text{map2 } f \ s_1 \ s_2$



$\text{fby } s_1 \ s_2$

$\text{cst } a = a : (\text{cst } a)$

$\text{map } f \ (h : t) = (f \ h) : (\text{map } f \ t)$

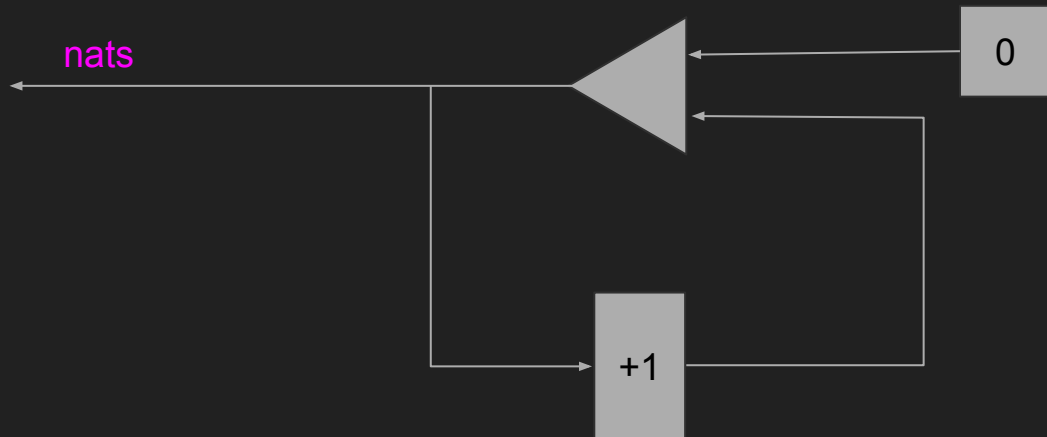
$\text{map2 } f \ (h_1 : t_1) \ (h_2 : t_2) = (f \ h_1 \ h_2) : (\text{map2 } f \ t_1 \ t_2)$

$\text{fby } (h : t) \ s = h : s$

$\text{tail } (h : t) = t$



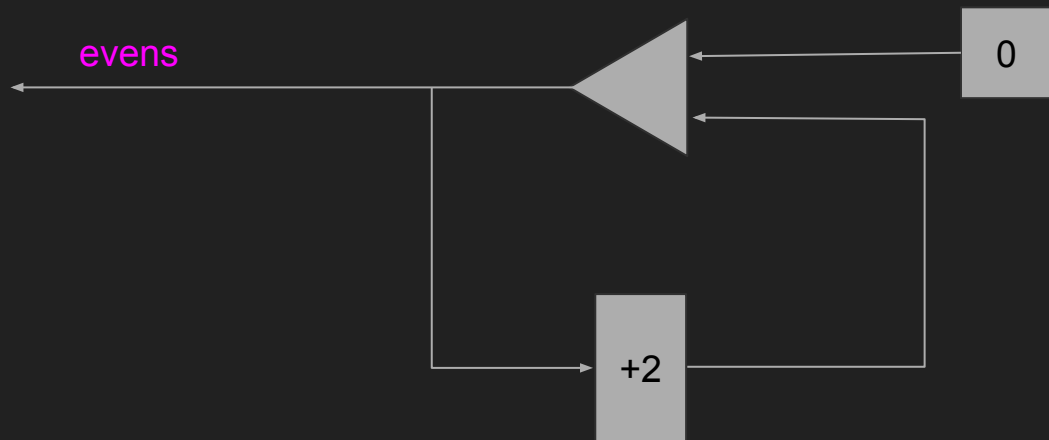
# Example: nats



`plus1 s = map (\a -> a + 1) s`

`nats = fby (cst 0) (plus1 nats)`

# Example: evens



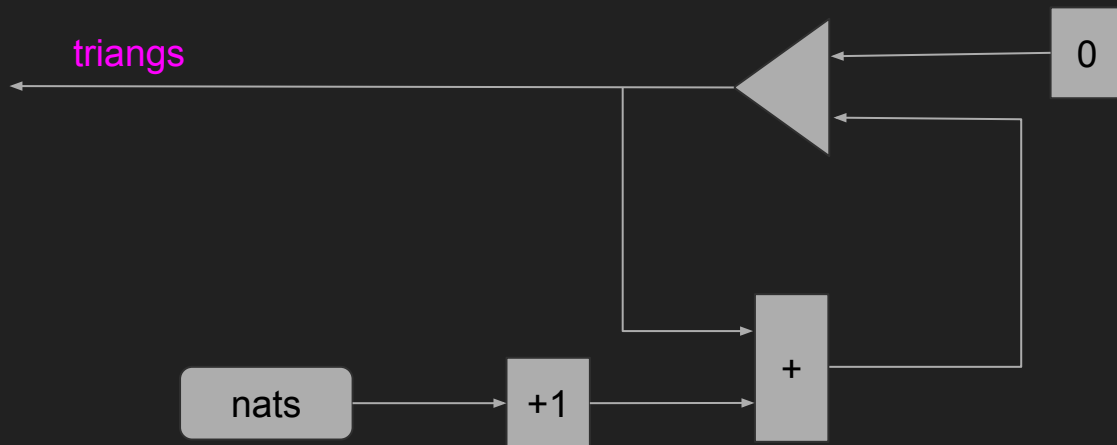
`evens = fby (cst 0) (plus1 (plus 1 evens))`

# Example: odds



odds = plus1 evens

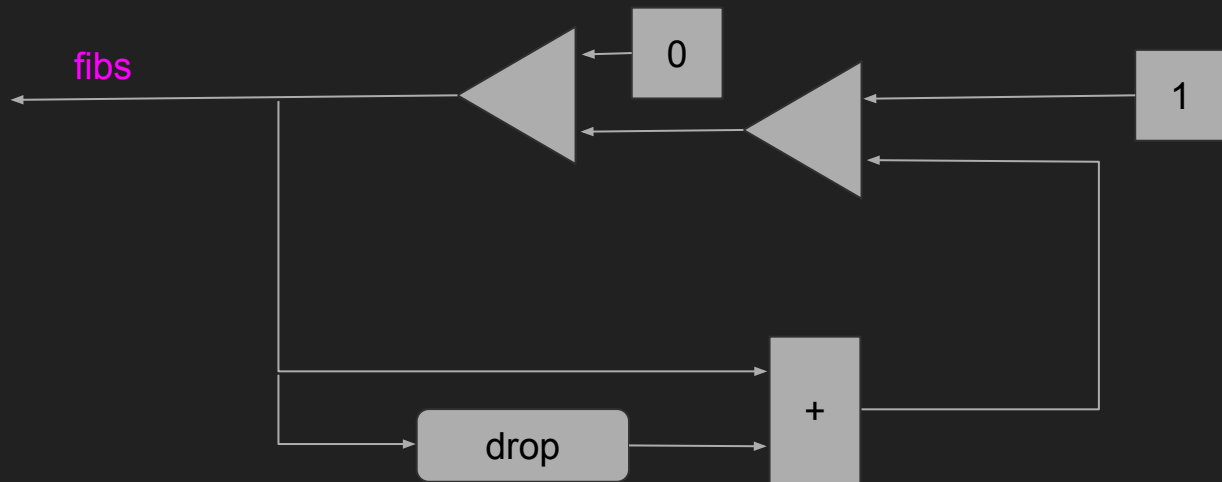
# Example: triangular numbers



$\text{plus } s_1 \ s_2 = \text{map2 } (\backslash a \rightarrow \backslash b \rightarrow a + b) \ s_1 \ s_2$

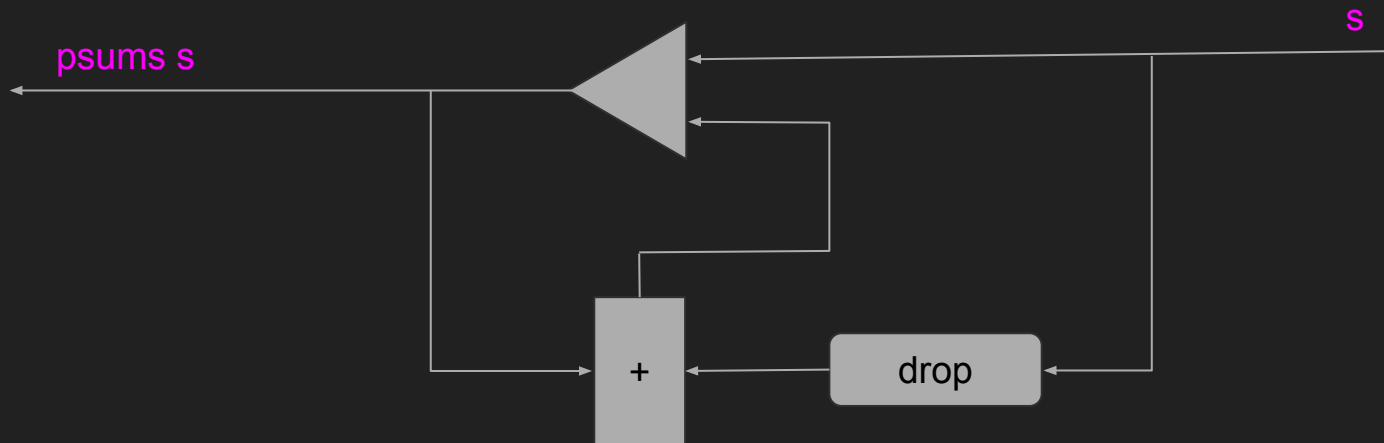
$\text{triangs} = \text{fby } (\text{cst } 0) \ (\text{plus triangs } (\text{plus1 nats}))$

# Example: Fibonacci numbers



`fibs = fby (cst 0) (fby (cst 1) (plus fibs (tail fibs)))`

# Example: partial sums



$\text{psums } s = \text{fby } s \text{ (plus (psums } s) \text{ (tail } s))}$



# Recursive stream programs

Sieve of Eratosthenes - how to compute the stream of prime numbers:

- sieving a stream: take a stream of values, keep the first value, and sieve the rest of the stream *after* removing all multiples of the first value
- sieving the natural numbers starting from 2 yields the prime numbers

```
divides c x = (mod x c == 0)
sieve s = fby s (sieve (filter (\a -> not (divides (head s) a)) (tail s)))

primes = sieve (plus1 (plus1 nats))
```