

Recursive Data Structures

FOCS, Fall 2020

Recursive data structures

A recursive data structure is a data structure made up of pieces that are instances of the data structure

Simplest example: a linked list is one of:

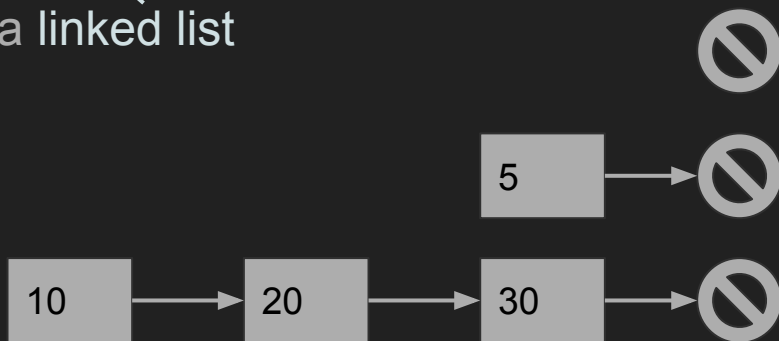
- an empty list, or
- a pair of an item and a linked list

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Simplest example: a **linked list** is one of:

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Recursive functions over recursive data structures

Most functions that works recursive data structures tend to be recursive functions

- Functions usually follow the shape of the data they work on

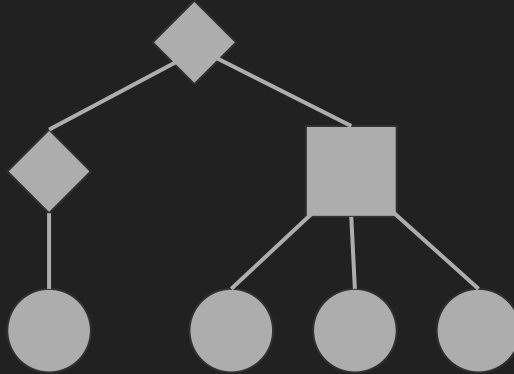
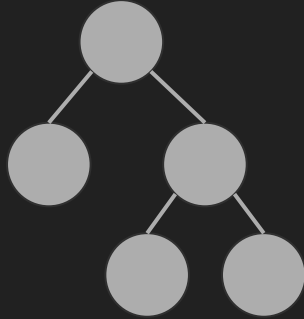
To **sum all the elements** of a list:

- if the list is the empty list: the sum is 0
- if the list if an integer N followed by a list L :
recursively sum all the elements in L , and add N to it

We know all of this already for lists

Tree data structures

The general form of a recursive data structure is a tree:



- where do we store values?
- kind of nodes?
- number of children?

Abstract syntax trees: arithmetic expressions

n

An **arithmetic expression** is one of:

- a number n
- $-(e_1)$ where e_1 is an arithmetic expression
- $(e_1 + e_2)$ where e_1 and e_2 are arithmetic expressions
- $(e_1 \times e_2)$ where e_1 and e_2 are arithmetic expressions

-

e_1

+

e_1

e_2

x

e_1

e_2

Abstract syntax trees: arithmetic expressions

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An arithmetic expression is one of:

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- $-(e_1)$ where e_1 is an arithmetic expression
- $(e_1 + e_2)$ where e_1 and e_2 are arithmetic expressions
- $(e_1 \times e_2)$ where e_1 and e_2 are arithmetic expressions

Often written as a BNF grammar:

```
e ::=  num
      -(e1)
      (e1 + e2)
      (e1 x e2)
```

-

e_1

+

e_1

e_2

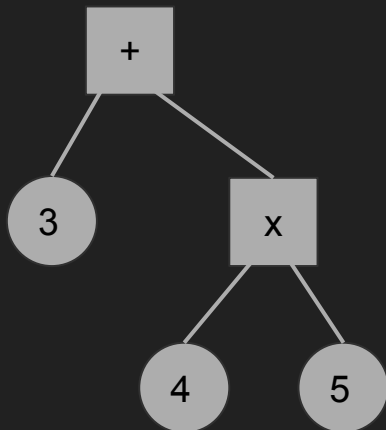
x

e_1

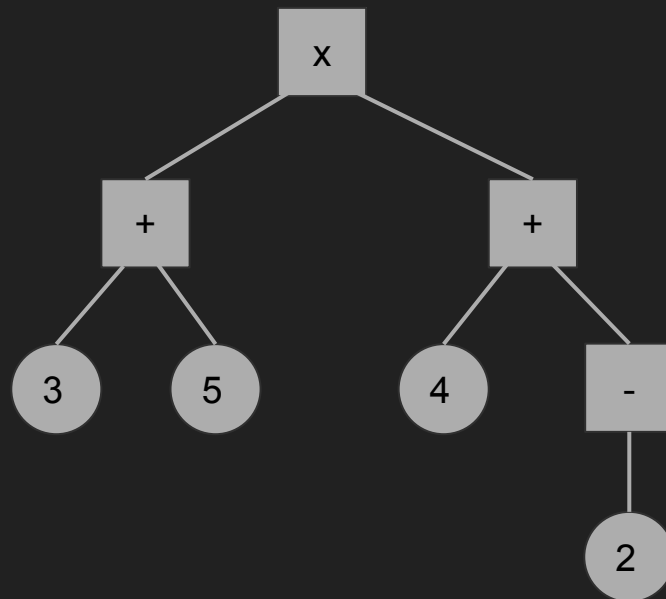
e_2

Examples


$$3 + (4 \times 5)$$





$$(3 + 5) \times (4 + (-2))$$



Evaluation

eval  = n

eval  = - eval e₁

eval  = eval e₁ + eval e₂

eval  = eval e₁ * eval e₂

Abstract syntax trees: adding identifiers

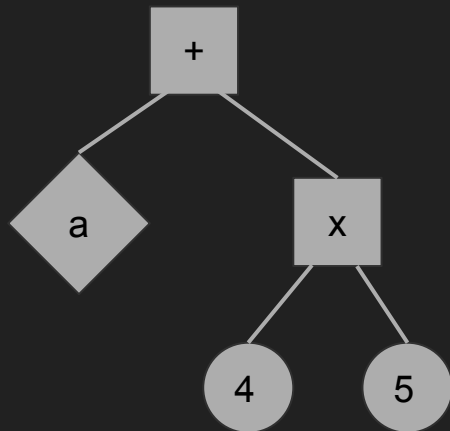
An **arithmetic expression** is one of:

- a number n
- an identifier id
- $-(e_1)$ where e_1 is an arithmetic expression
- $(e_1 + e_2)$ where e_1 and e_2 are arithmetic expressions
- $(e_1 \times e_2)$ where e_1 and e_2 are arithmetic expressions

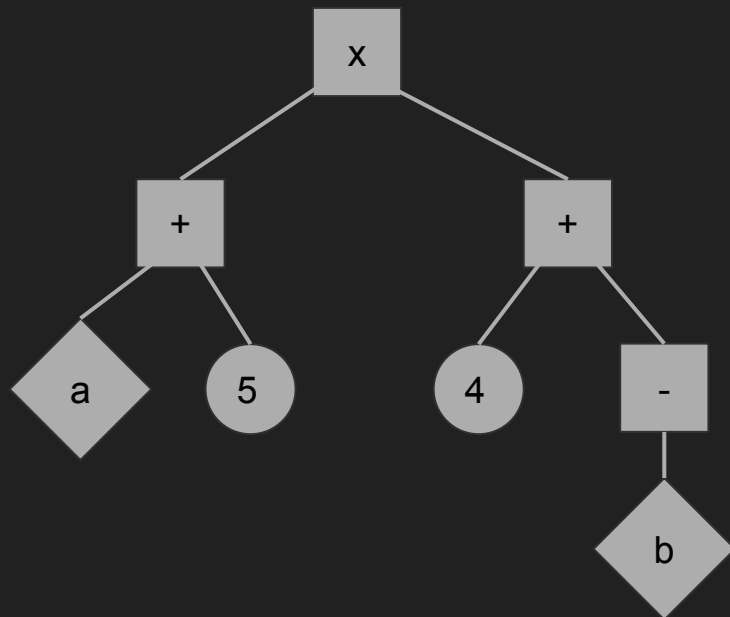


Examples

$$a + (4 \times 5)$$



$$(a + 5) \times (4 + (-b))$$



Evaluation with environments

$$\text{eval } E \quad \text{○ } n \quad = n$$

$$\text{eval } E \quad \begin{array}{c} \text{□ } - \\ | \\ e_1 \end{array} \quad = - \text{eval } E \ e_1$$

$$\text{eval } E \quad \begin{array}{c} \text{□ } + \\ / \quad \backslash \\ e_1 \quad e_2 \end{array} \quad = \text{eval } E \ e_1 + \text{eval } E \ e_2$$

$$\text{eval } E \quad \begin{array}{c} \text{□ } x \\ / \quad \backslash \\ e_1 \quad e_2 \end{array} \quad = \text{eval } E \ e_1 * \text{eval } E \ e_2$$

$$\text{eval } E \quad \text{◇ } \text{id} \quad = E(\text{id})$$

Implementation using algebraic datatypes

```
type aexp =  
  | Num of int  
  | Neg of aexp  
  | Plus of aexp * aexp  
  | Times of aexp * aexp
```

```
let rec eval aexp =  
  match aexp with  
  | Num n -> n  
  | Neg e1 -> - (eval e1)  
  | Plus (e1, e2) -> (eval e1) + (eval e2)  
  | Times (e1, e2) -> (eval e1) * (eval e2)
```

Example: `Plus (Num 3, Times (Num 4, Num 5))`

Implementation using algebraic datatypes

```
type aexp =  
  | Num of int  
  | Neg of aexp  
  | Plus of aexp * aexp  
  | Times of aexp * aexp  
  | Ident of string
```

```
let rec eval env aexp =  
  match aexp with  
  | Num n -> n  
  | Neg e1 -> - (eval env e1)  
  | Plus (e1, e2) -> (eval env e1) + (eval env e2)  
  | Times (e1, e2) -> (eval env e1) * (eval env e2)  
  | Ident s -> lookup env s
```

Example: `Plus (Ident "a", Times (Num 4, Num 5))`

Implementation using classes

```
abstract class AExp {  
    abstract int eval ();  
}
```

```
class Num extends AExp {  
    int n;  
  
    Num (int n) {  
        this.n = n;  
    }  
  
    int eval () {  
        return this.n;  
    }  
}
```

Implementation using classes

```
abstract class AExp {  
    abstract int eval ();  
}
```

```
class Neg extends AExp {  
    AExp e1;  
  
    Neg (AExp e1) {  
        this.e1 = e1;  
    }  
  
    int eval () {  
        return -(this.e1.eval());  
    }  
}
```


Implementation using classes

```
abstract class AExp {  
    abstract int eval ();  
}
```

```
class Plus extends AExp {  
    AExp e1, e2;  
  
    Plus (AExp e1, AExp e2) {  
        this.e1 = e1;  
        this.e2 = e2;  
    }  
  
    int eval () {  
        return this.e1.eval() + this.e2.eval();  
    }  
}
```

Implementation using classes

```
abstract class AExp {  
    abstract int eval ();  
}
```

```
class Times extends AExp {  
    AExp e1, e2;  
  
    Times (AExp e1, AExp e2) {  
        this.e1 = e1;  
        this.e2 = e2;  
    }  
  
    int eval () {  
        return this.e1.eval() * this.e2.eval();  
    }  
}
```

Implementation using classes

```
abstract class AExp {  
    abstract int eval ();  
}
```

Example:

```
AExp e = new Plus(new Num(3),  
                  new Times(new Num(4),  
                             new Num(5)));
```

JSON in OCaml

A **JSON item** is one of:

- an integer
- a string
- a list of JSON items
- an object, mapping string keys to JSON items

Standard notation: {
 "a" : [1, 2, 3],
 "b" : { "c" : 1, "d" : "two" }
 }

JSON in OCaml

```
type json =  
  | JInteger of int  
  | JString of string  
  | JList of json list  
  | JObject of (string * json) list
```

Recursive functions to

- dump a JSON object to a file or a string
- transform (map) a JSON object
- traverse (fold) a JSON object