

Graph Algorithms

DSA, Fall 2022

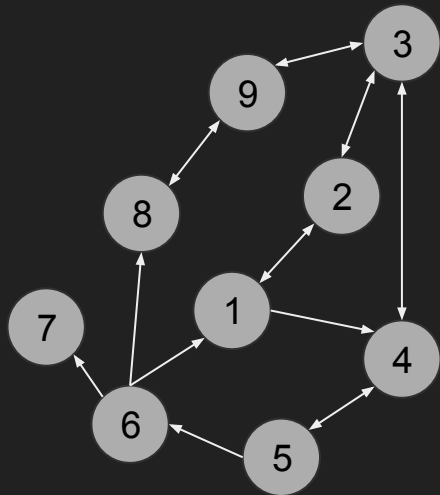
Graphs

A **graph** is a set of things (vertices) connected together (edges)

- if connections have a direction — directed graph
- if connections do not have a direction — undirected graph

Can be used to represent

- networks of computers with communication links
- networks of friends with "friend of" or "follows" connectivity
- cities and flights between them
- etc...



Definitions

A **directed graph** is a pair (V, E) consisting of:

- A finite set V of vertices
- A finite set E of edges of the form (u, v) connecting vertex u to vertex v

Edge (u, v) — u is the source, v is the target — written $u \rightarrow v$

For undirected graphs, take edges to be pairs $\{u, v\}$

v is **reachable** from u if there exists vertices v_1, \dots, v_k such that
 $u \rightarrow v_1, v_1 \rightarrow v_2, \dots, v_{k-1} \rightarrow v_k$, and $v_k \rightarrow v$

Graph Representations

Need to represent both vertices and edges

- For simplicity, assume vertices are $0, 1, \dots, n$

Most algorithms involve querying the graph more than updating the graph

Adjacency list

attach to each vertex a list of connected vertices

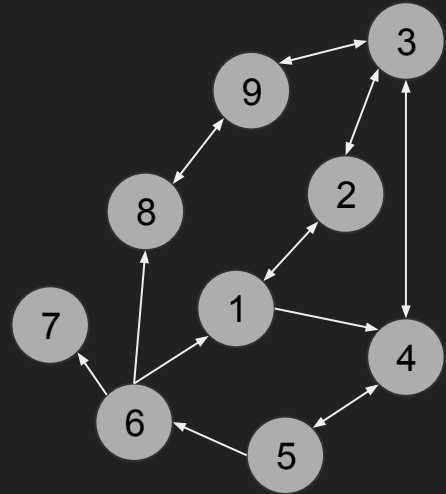
Adjacency matrix

matrix indexed by vertices with 1 at (i, j) when there's an edge from i to j

Adjacency List Representation

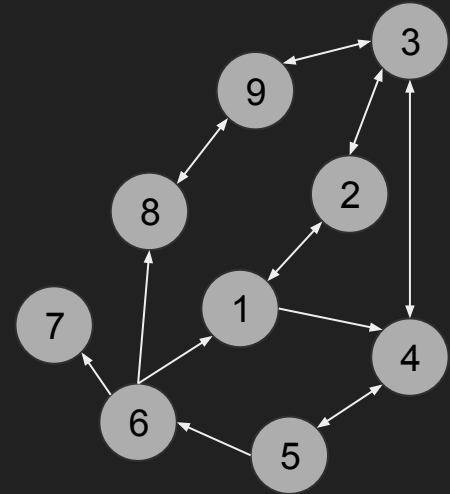
```
type Graph struct {  
    vertices int  
    edges []*Edge  
}
```

```
type Edge struct {  
    target int  
    next *Edge  
}
```



Adjacency Matrix Representation

```
type Graph struct {  
    vertices int  
    edges [][]int  
}
```



Choice of Representation

The sparser a graph (small number of edges vs vertices), the more benefit from the adjacency list representation

- maximum number of edges in a (directed) graph is $|V|^2$
- a graph is **sparse** if $|E| \ll |V|^2$
- most practical graphs are sparse

Space for adjacency list representation $O(|V| + |E|)$

Space for adjacency matrix representation $O(|V|^2)$

Both pretty straightforwardly generalize to weighted graphs representations

Core Algorithms — Search

Given a starting vertex u , search for all reachable vertices from u

- construct a path from u to each reachable vertex
- it's basically a tree rooted at u with reachable vertices as nodes

Two basic approaches

- systematically look at all vertices distance 1, 2, 3, ... from starting vertex
- blaze through from starting vertex, backtracking when you get stuck

Search Skeleton

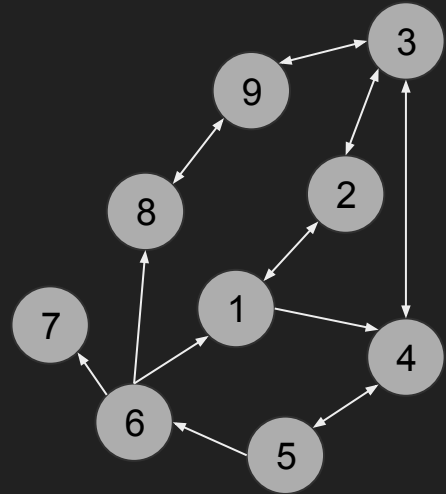
```
color every vertex gray
color[start] ← red
while there is a red vertex:
    pick a red vertex v
    color[v] ← green
    for every adjacent vertex u of v:
        if color[u] = gray:
            color[u] ← red
```

How to pick next red vertex to look at?

Breadth-First Search

Add red vertices into a QUEUE

```
color every vertex gray
color[start] ← red
ENQUEUE(Q, start)
while not ISEMPTY(Q):
  v ← DEQUEUE(Q)
  color[v] ← green
  for every adjacent vertex u of v:
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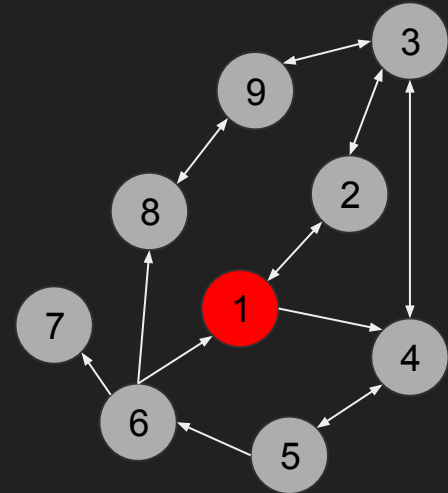


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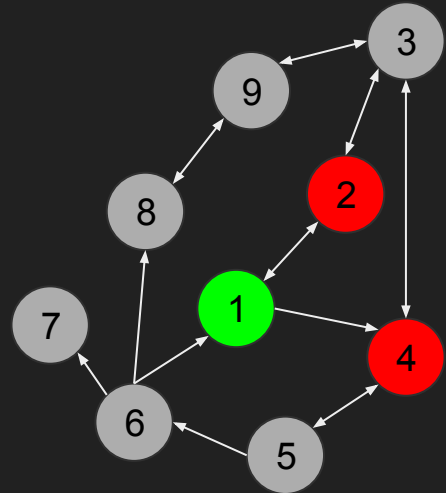


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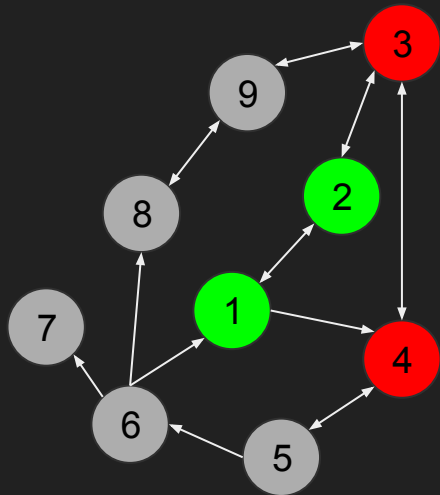


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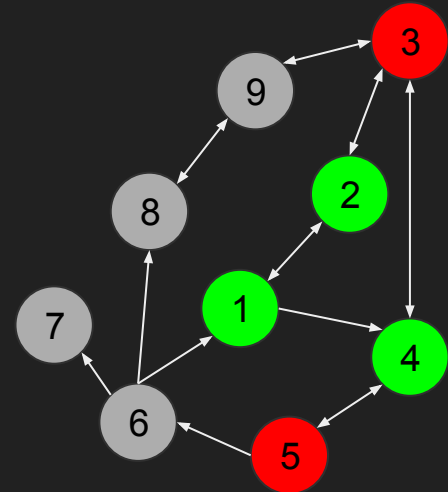


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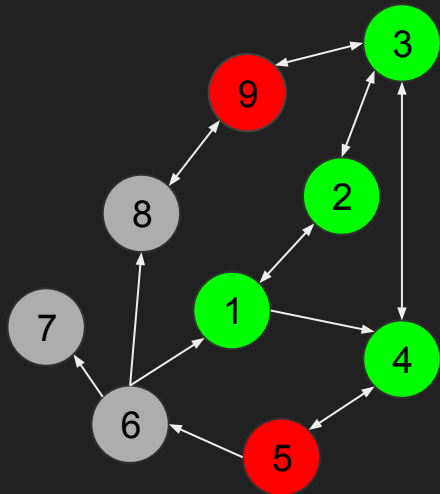


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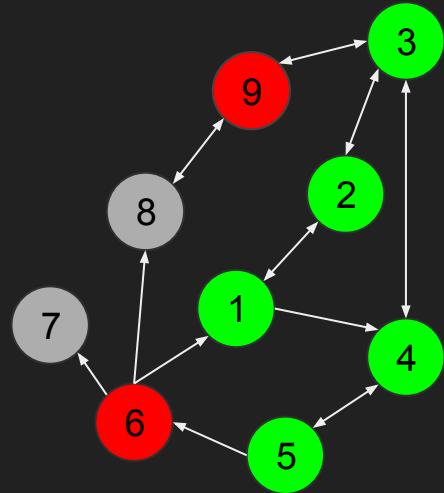


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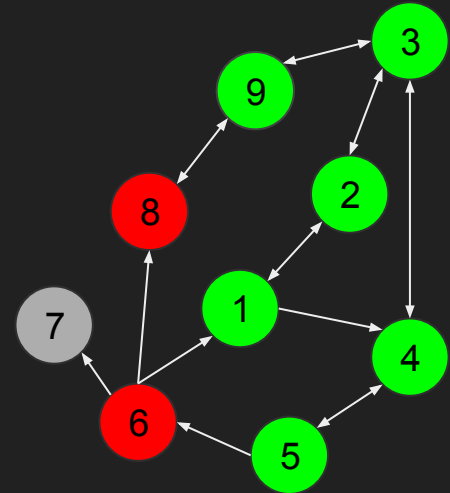
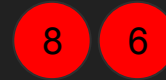


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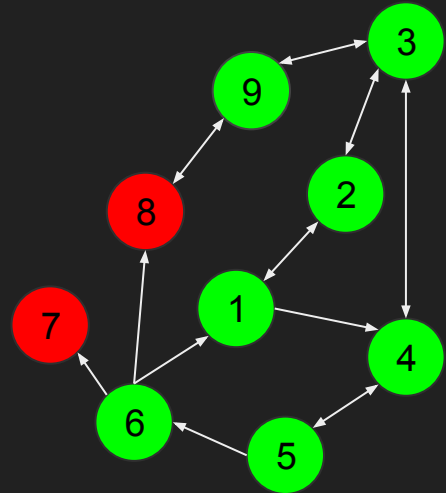


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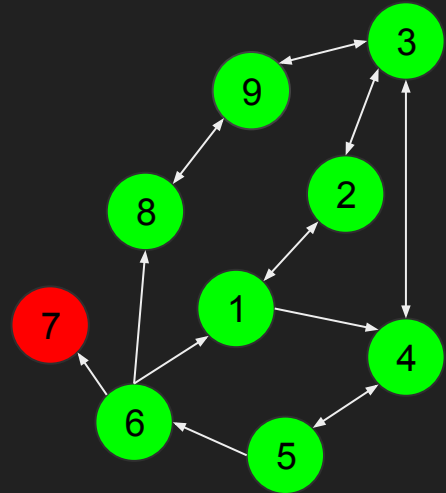
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7

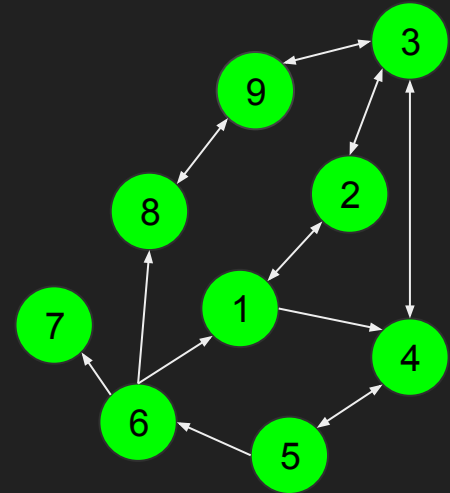


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```

QUEUE:



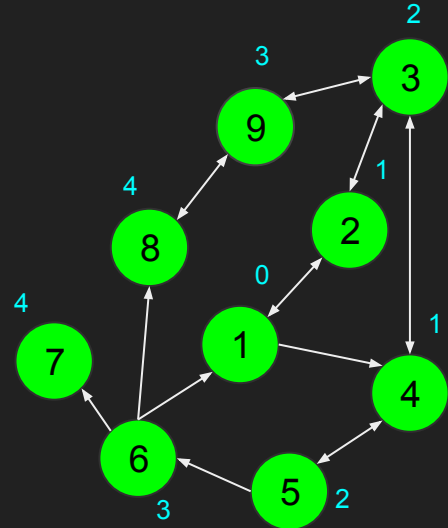
Application — Shortest Paths

A feature of BFS is that it searches vertices in order of "closeness" from the start vertex

- "closeness" defined in terms of number of edges to follow to reach vertex

Application — Shortest Paths

```
for every vertex u: color[u] ← gray
for every vertex u: dist[u] ← ∞
color[start] ← red
ENQUEUE(Q, start)
dist[start] ← 0
while not ISEMPTY(Q):
    v ← DEQUEUE(Q)
    color[v] ← green
    for every adjacent vertex u of v:
        if color[u] = gray:
            color[u] ← red
            dist[u] ← dist[v] + 1
            ENQUEUE(Q, u)
```



Application — Shortest Paths

A feature of BFS is that it searches vertices in order of "closeness" from the start vertex

- "closeness" defined in terms of number of edges to follow to reach vertex

Can also construct the shortest path from start vertex for every reachable vertex

- shortest path is not unique

Application — Shortest Paths

```
for every vertex u: color[u] ← gray
for every vertex u: dist[u] ← ∞  parent[u] ← nil
color[start] ← red
ENQUEUE(Q, start)
dist[start] ← 0
while not ISEMPTY(Q):
    v ← DEQUEUE(Q)
    color[v] ← green
    for every adjacent vertex u of v:
        if color[u] = gray:
            color[u] ← red
            dist[u] ← dist[v] + 1  parent[u] ← v
            ENQUEUE(Q, u)
```

When $\text{dist}[u] < \infty$
following $\text{parent}[u]$ up to
start yields a path from
start to u

Depth-First Search

Add red vertices into a STACK

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color every vertex gray
color[start] ← red
ENQUEUE(Q, start)
while not ISEMPTY(Q):
    v ← DEQUEUE(Q)
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    for every adjacent vertex u of v:
        if color[u] = gray:
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```

Depth-First Search

Add red vertices into a STACK

color every vertex gray

color[start] \leftarrow red

PUSH(T, start)

while not ISEMPTY(T):

 v \leftarrow POP(T)

 if color[v] \neq green:

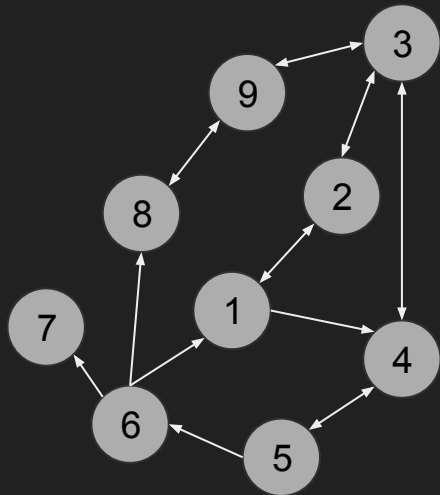
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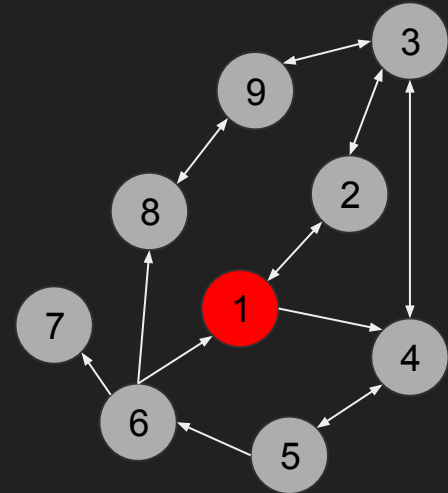


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while not ISEMPTY(T):
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  if color[v] != green:
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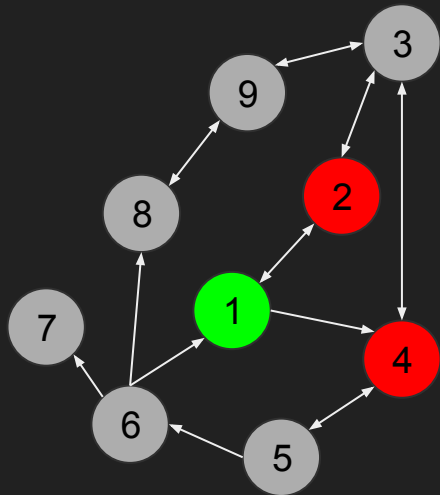
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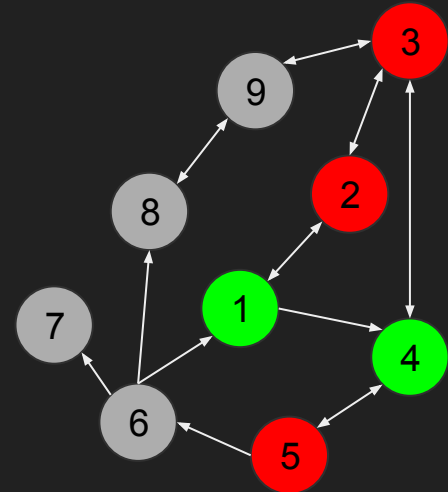
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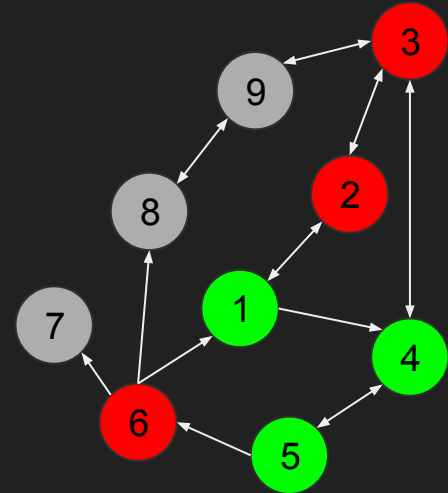
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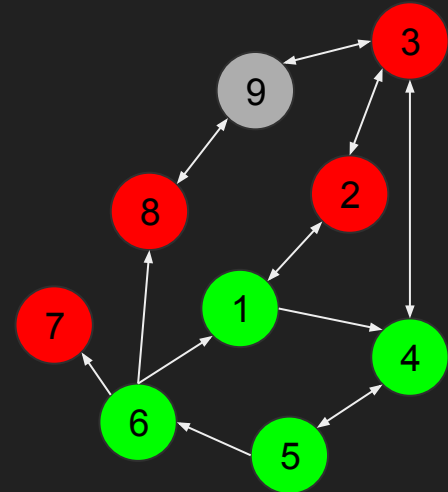
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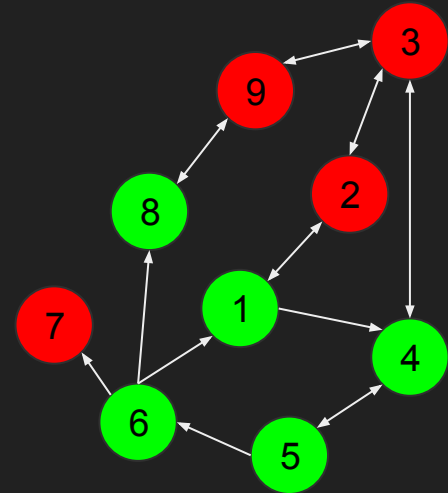
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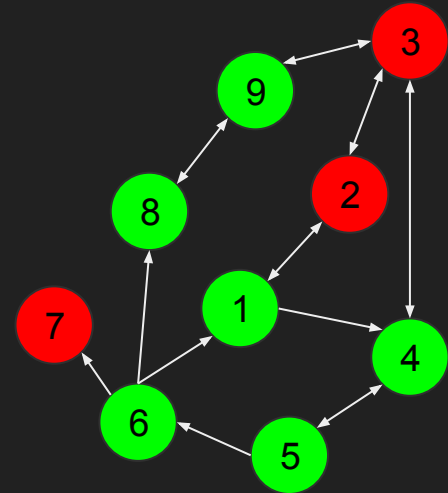
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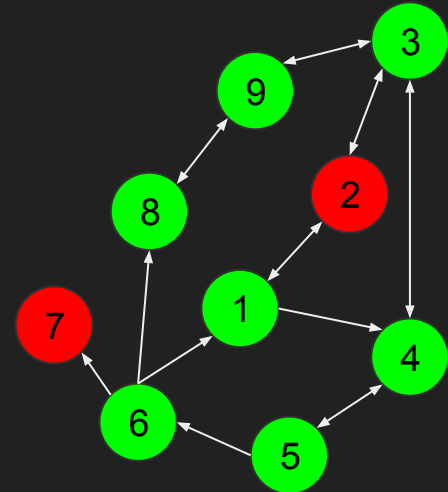
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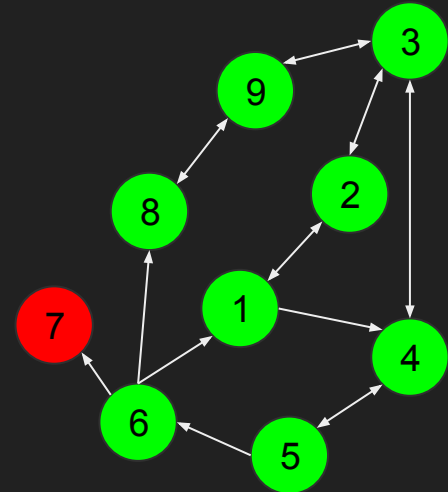
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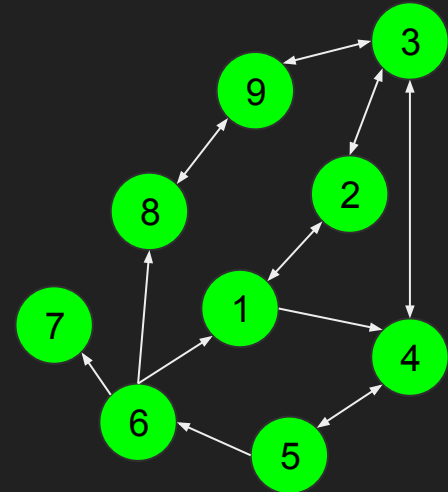
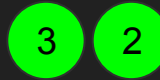
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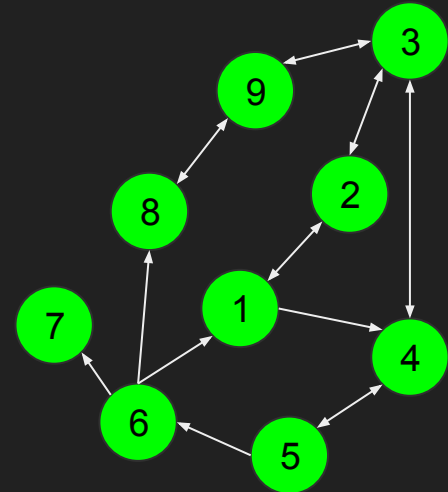
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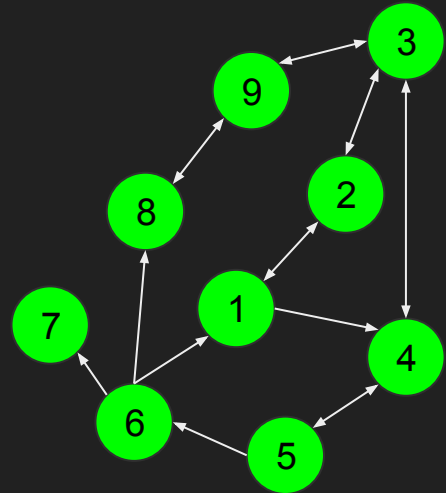
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STACK:



Application — Topological Sort

An acyclic directed graph is a directed graph with no cycles

- no path from a vertex to itself

Given an acyclic directed graph G , the **topological sort** of G is an ordering of the vertices of G such that when there's an edge $u \rightarrow v$ in G , then u comes before v in the ordering

Classic: vertices of G are tasks, $u \rightarrow v$ means u must be done before v , then a topological sort is a way to schedule tasks so that required tasks are done first

Practical: vertices are modules, $u \rightarrow v$ means v depends on u , a topological sort is an order of loading modules so you don't get errors

Application — Topological Sort

Starting point: recursive version of DFS

Blue vertices represent vertices being worked on

color every vertex gray

DFS(G, start)

DFS(G, v) :=

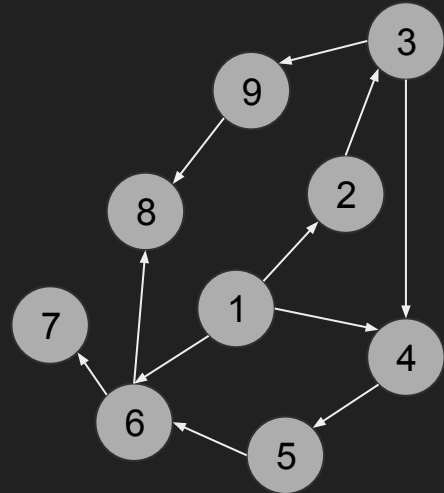
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 for every adjacent vertex u of v:

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 color[v] ← green

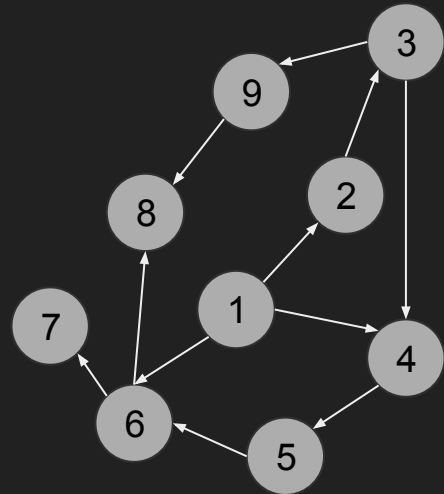


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

```
DFS( $G, v$ ) :=  
  color[ $v$ ] ← blue  
  for every adjacent vertex  $u$  of  $v$ :  
    if color[ $u$ ] = blue:  
      error "cycle in graph"  
    if color[ $u$ ] = gray:  
      DFS( $G, u$ )  
  color[ $v$ ] ← green  
  output( $v$ )
```

OUTPUT:

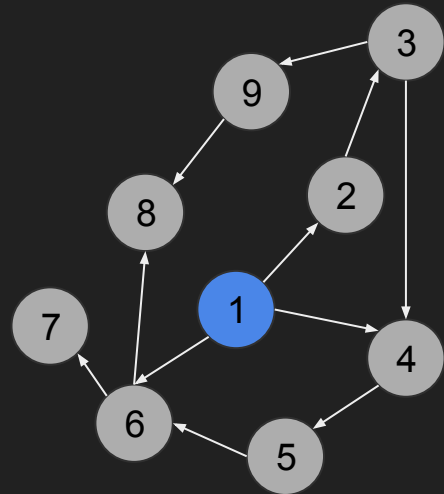


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 output(v)

OUTPUT:

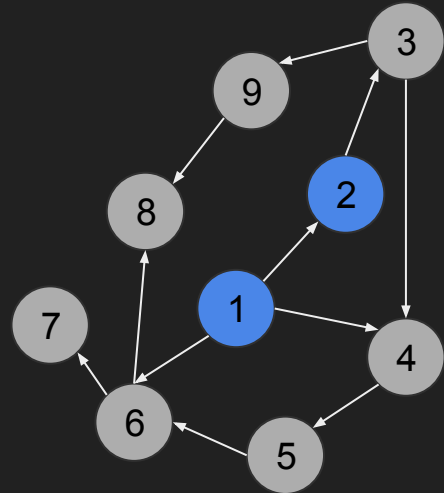


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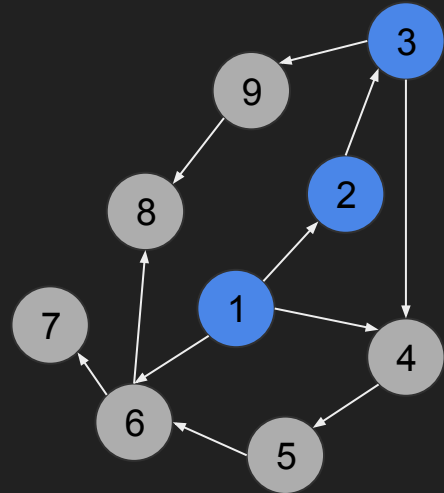


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      DFS( $G, u$ )  
  color[ $v$ ] ← green  
  output( $v$ )
```

OUTPUT:

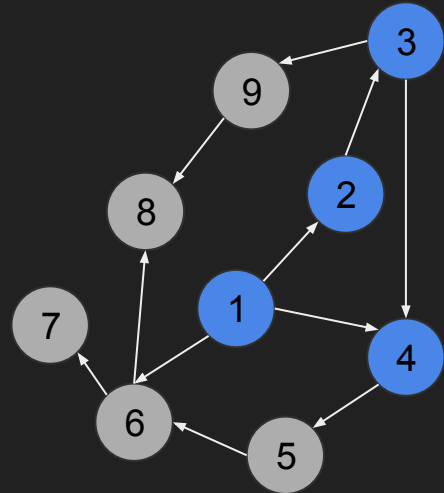


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

```
DFS( $G, v$ ) :=  
  color[ $v$ ] ← blue  
  for every adjacent vertex  $u$  of  $v$ :  
    if color[ $u$ ] = blue:  
      error "cycle in graph"  
    if color[ $u$ ] = gray:  
      DFS( $G, u$ )  
  color[ $v$ ] ← green  
  output( $v$ )
```

OUTPUT:

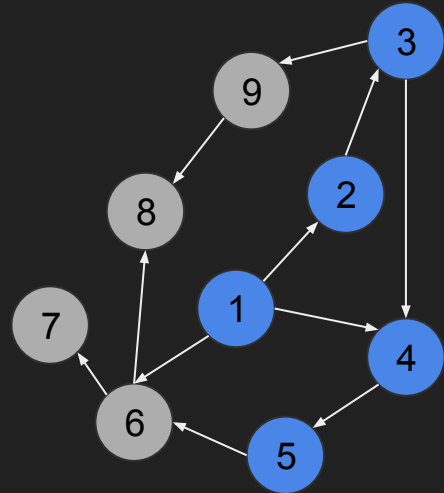


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

```
DFS( $G, v$ ) :=  
  color[ $v$ ] ← blue  
  for every adjacent vertex  $u$  of  $v$ :  
    if color[ $u$ ] = blue:  
      error "cycle in graph"  
    if color[ $u$ ] = gray:  
      DFS( $G, u$ )  
  color[ $v$ ] ← green  
  output( $v$ )
```

OUTPUT:

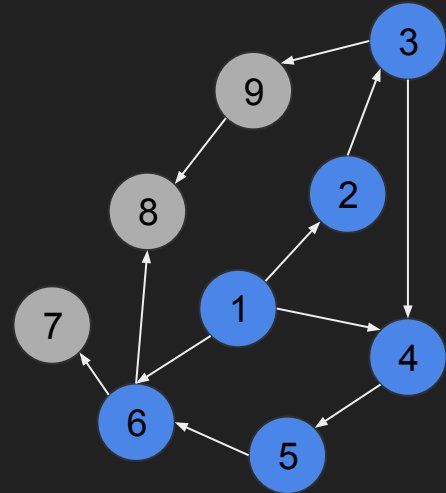


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

```
DFS( $G, v$ ) :=  
  color[ $v$ ] ← blue  
  for every adjacent vertex  $u$  of  $v$ :  
    if color[ $u$ ] = blue:  
      error "cycle in graph"  
    if color[ $u$ ] = gray:  
      DFS( $G, u$ )  
  color[ $v$ ] ← green  
  output( $v$ )
```

OUTPUT:

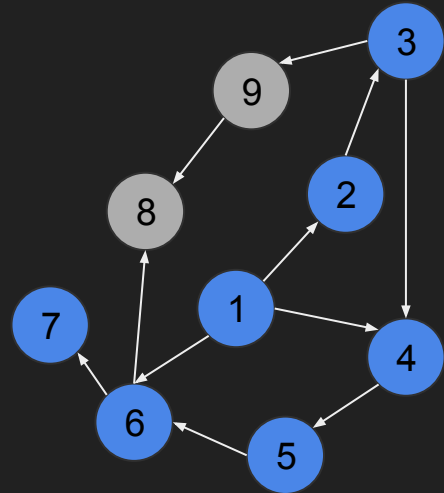


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

```
DFS( $G, v$ ) :=  
  color[ $v$ ] ← blue  
  for every adjacent vertex  $u$  of  $v$ :  
    if color[ $u$ ] = blue:  
      error "cycle in graph"  
    if color[ $u$ ] = gray:  
      DFS( $G, u$ )  
  color[ $v$ ] ← green  
  output( $v$ )
```

OUTPUT:

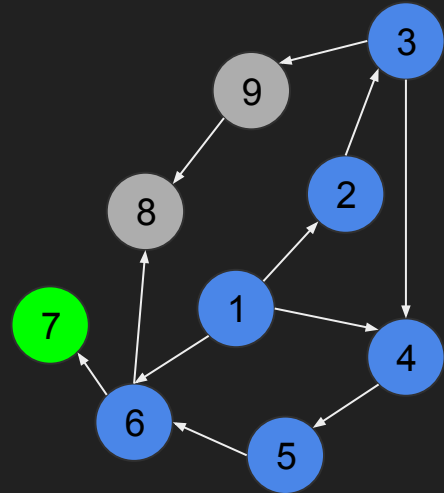


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

DFS(G, v) :=
 color[v] \leftarrow blue
 for every adjacent vertex u of v :
 if color[u] = blue:
 error "cycle in graph"
 if color[u] = gray:
 DFS(G, u)
 color[v] \leftarrow green
 output(v)

OUTPUT: 7

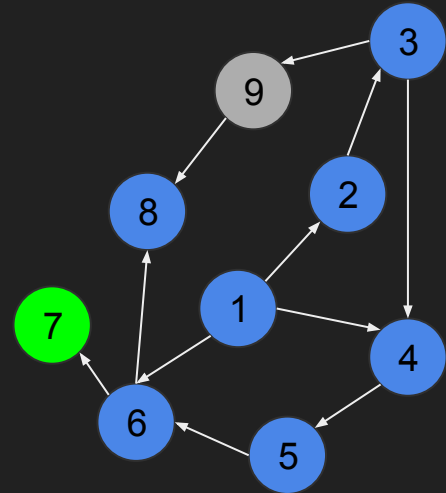


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

```
DFS( $G, v$ ) :=  
  color[ $v$ ] ← blue  
  for every adjacent vertex  $u$  of  $v$ :  
    if color[ $u$ ] = blue:  
      error "cycle in graph"  
    if color[ $u$ ] = gray:  
      DFS( $G, u$ )  
  color[ $v$ ] ← green  
  output( $v$ )
```

OUTPUT: 7

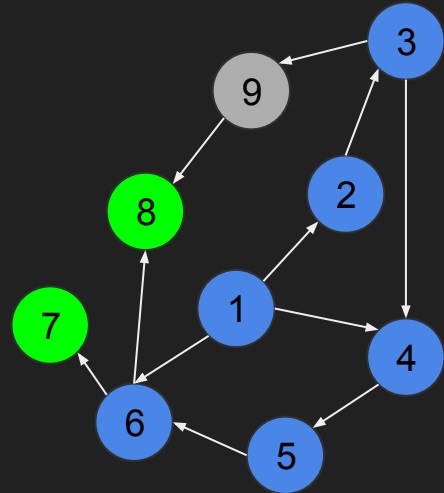


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

DFS(G, v) :=
 color[v] \leftarrow blue
 for every adjacent vertex u of v :
 if color[u] = blue:
 error "cycle in graph"
 if color[u] = gray:
 DFS(G, u)
 color[v] \leftarrow green
 output(v)

OUTPUT: 7 8

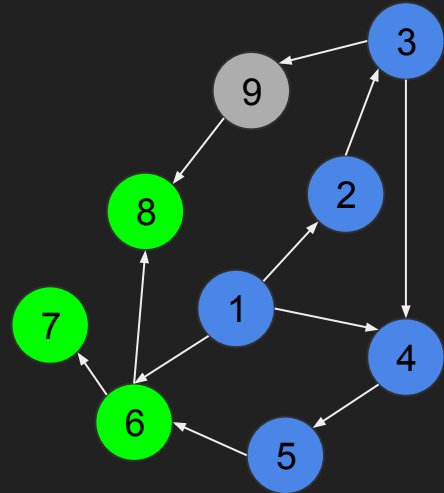


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

DFS(G, v) :=
 color[v] \leftarrow blue
 for every adjacent vertex u of v :
 if color[u] = blue:
 error "cycle in graph"
 if color[u] = gray:
 DFS(G, u)
 color[v] \leftarrow green
 output(v)

OUTPUT: 7 8 6

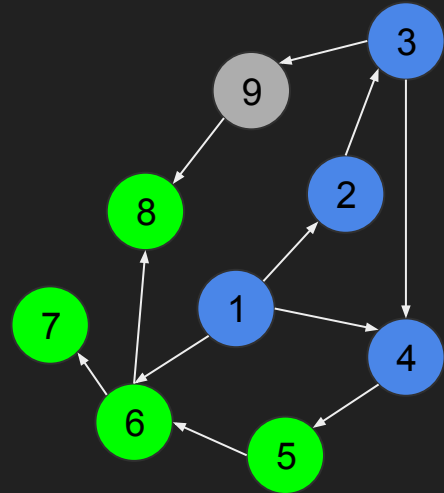


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

DFS(G, v) :=
 color[v] \leftarrow blue
 for every adjacent vertex u of v :
 if color[u] = blue:
 error "cycle in graph"
 if color[u] = gray:
 DFS(G, u)
 color[v] \leftarrow green
 output(v)

OUTPUT: 7 8 6 5

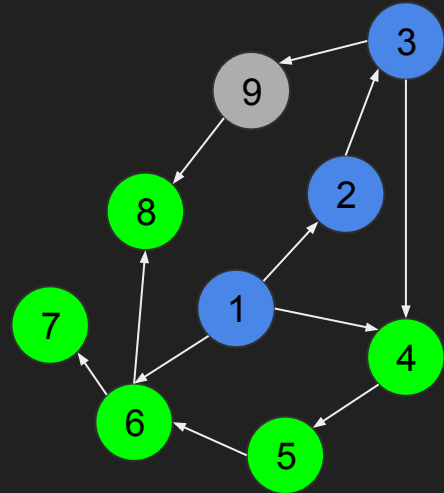


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

DFS(G, v) :=
 color[v] \leftarrow blue
 for every adjacent vertex u of v :
 if color[u] = blue:
 error "cycle in graph"
 if color[u] = gray:
 DFS(G, u)
 color[v] \leftarrow green
 output(v)

OUTPUT: 7 8 6 5 4

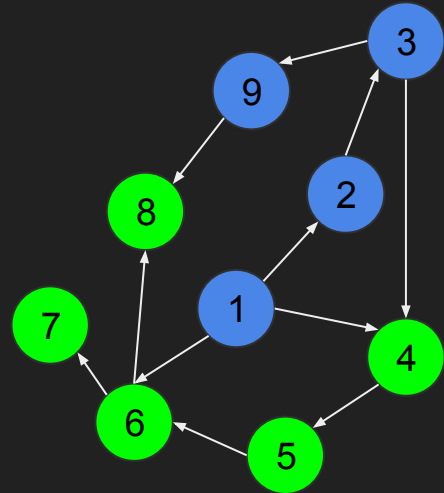


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

DFS(G, v) :=
 color[v] \leftarrow blue
 for every adjacent vertex u of v :
 if color[u] = blue:
 error "cycle in graph"
 if color[u] = gray:
 DFS(G, u)
 color[v] \leftarrow green
 output(v)

OUTPUT: 7 8 6 5 4

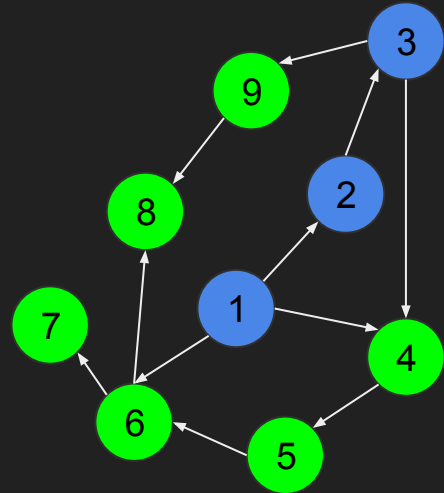


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

DFS(G, v) :=
 color[v] \leftarrow blue
 for every adjacent vertex u of v :
 if color[u] = blue:
 error "cycle in graph"
 if color[u] = gray:
 DFS(G, u)
 color[v] \leftarrow green
 output(v)

OUTPUT: 7 8 6 5 4 9

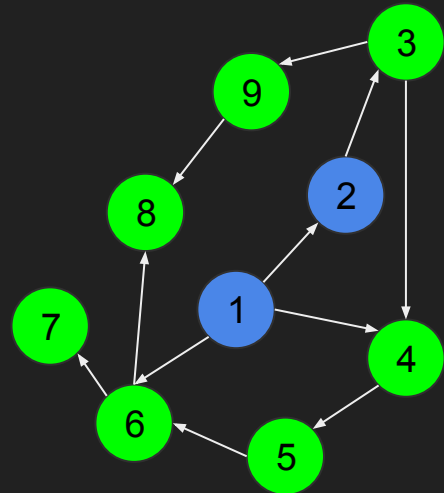


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

DFS(G, v) :=
 color[v] \leftarrow blue
 for every adjacent vertex u of v :
 if color[u] = blue:
 error "cycle in graph"
 if color[u] = gray:
 DFS(G, u)
 color[v] \leftarrow green
 output(v)

OUTPUT: 7 8 6 5 4 9 3

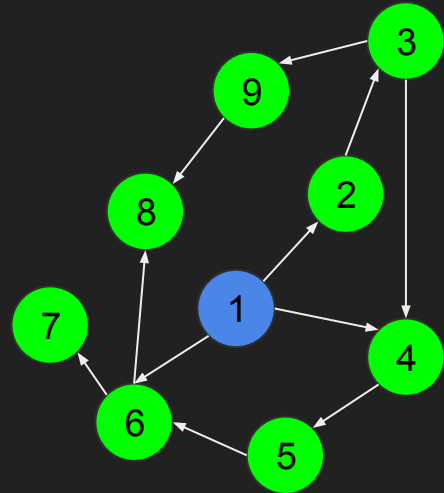


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

DFS(G, v) :=
 color[v] \leftarrow blue
 for every adjacent vertex u of v :
 if color[u] = blue:
 error "cycle in graph"
 if color[u] = gray:
 DFS(G, u)
 color[v] \leftarrow green
 output(v)

OUTPUT: 7 8 6 5 4 9 3 2

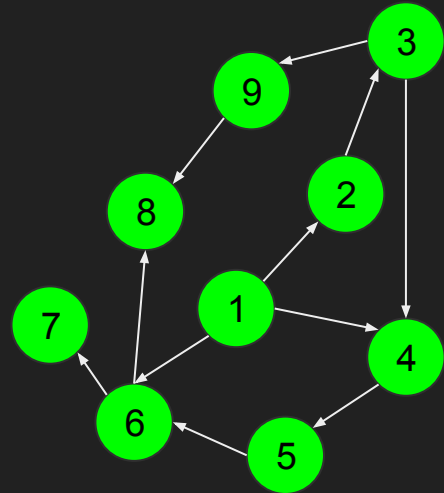


Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

DFS(G, v) :=
 color[v] \leftarrow blue
 for every adjacent vertex u of v :
 if color[u] = blue:
 error "cycle in graph"
 if color[u] = gray:
 DFS(G, u)
 color[v] \leftarrow green
 output(v)

OUTPUT: 7 8 6 5 4 9 3 2 1



Application — Topological Sort

color every vertex gray
for every vertex u with no edge into it:
 DFS(G, u)

DFS(G, v) :=
 color[v] \leftarrow blue
 for every adjacent vertex u of v :
 if color[u] = blue:
 error "cycle in graph"
 if color[u] = gray:
 DFS(G, u)
 color[v] \leftarrow green
 output(v)

OUTPUT: 7 8 6 5 4 9 3 2 1

Reverse to get a topological sort

1 2 3 9 4 5 6 8 7

