Reaction Network Modeling

Molecular Programming tutorial ASE 2015

A mathematical model.

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A programming language for molecular programming in a well-mixed solution.

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"Computing in soup"

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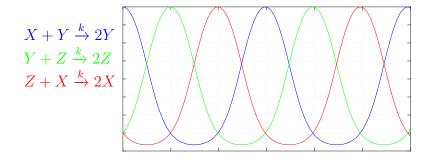
Essentially equivalent to

Population protocols

Petri nets

Vector addition systems

Example: A three-phase oscillator (\approx Lotka-Volterra)



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1. Descriptive (20th century): Use CRNs to describe/explain chemical/biochemical processes.

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- 2. Mathematical (20th century): Use algebra, analysis, and probability to analyze structure and dynamics of CRNs.

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 - (a) Write down an (abstract) CRN that does something cool.

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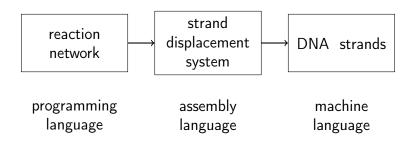
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- 4. Systematic prescriptive (21st century):

Replace (c) with uniform translation into DNA strand displacement systems.

Soloveichik, Seelig, and Winfree 2010



Chen, Dalchau, Srinivas, Phillips, Cardelli, Soloveichik, and Seelig 2013.

"Real programmers code in chemistry!"

Soloveichik, 2011

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A reaction over a finite set $S \subseteq \mathbf{S}$ is a triple

$$\rho = (\mathbf{r}, \mathbf{p}, k) \in \mathbb{N}^S \times \mathbb{N}^S \times (0, \infty)$$

with $\mathbf{r} \neq \mathbf{p}$.

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with $\mathbf{r} \neq \mathbf{p}$.

$$\mathbf{r} = \mathbf{r}(\rho) = \text{ reactant vector of } \rho$$

 $\mathbf{p} = \mathbf{p}(\rho) = \text{ product vector of } \rho$
 $k = k(\rho) = \text{ rate constant of } \rho$

Example: Let $S = \{X, Y, Z\}$. The intuitive notation

$$X + Z \xrightarrow{k} 2Y + Z$$

refers to the formal reaction $(\mathbf{r}, \mathbf{p}, k)$, where

$$\mathbf{r}(X) = \mathbf{r}(Z) = 1, \ \mathbf{r}(Y) = 0,$$

 $\mathbf{p}(X) = 0, \ \mathbf{p}(Y) = 2, \ \mathbf{p}(Z) = 1.$

The two most widely used semantics (operational meanings) ascribed to a CRN N=(S,R).

1. Deterministic mass action

A state is a vector $\mathbf{x} \in [0, \infty)^S$.

 $\mathbf{x}(Y) = \text{the concentration of } Y \text{ in state } \mathbf{x}.$

Behavior of N governed by autonomous polynomial ODEs.

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Behavior of N governed by autonomous polynomial ODEs.

Used to model systems with many molecules of each species.

Law of Mass Action (deterministic)

The rate of a reaction depends on its rate constant and the state!

Example. In state x, the reactions
$$\rho_1 \equiv X + Z \xrightarrow{a} 2Y + Z$$
 and
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Example (continued). The CRN

$$X + Z \xrightarrow{a} 2Y + Z$$
$$2Y + Z \xrightarrow{b} X + Z$$

is governed by the system

$$\frac{dx}{dt} = -axz + by^2z$$

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is governed by the system

$$\begin{split} \frac{dx}{dt} &= -axz + by^2z\\ \frac{dy}{dt} &= 2axz - 2by^2z\\ \frac{dz}{dt} &= 0 \quad (Z \text{ is a catalyst.}) \end{split}$$

Open question: Are deterministic mass action CRNs Turing universal?

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What is robustness here?

Klinge, Lathrop, and Lutz 2015: Uniform translation of a nondeterministic finite automaton M into a deterministic mass action CRN N.

```
N simulates M size(N) linear in size(M) N robust with respect to input signal initial concentrations output rate constants of reactions Proof, not simulation.
```

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2. Stochastic mass action

A state is a vector $\mathbf{x} \in \mathbb{N}^S$.

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Behavior is governed by continuous time Markov chains.

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Used to model systems in which small counts of molecules are significant.

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Theorem (Kurtz 1972). Roughly stated, deterministic mass action is the limiting case of stochastic mass action as populations and volumes go to infinity.

The computational power of stochastic mass action CRNs.

Angluin, Aspnes, and Eisenstat 2006;
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Stable computation (correctness on all paths) can only compute semilinear predicates and functions.

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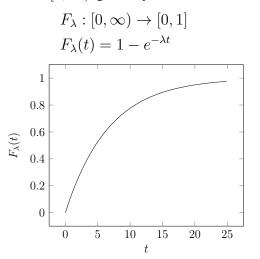
- Soloveichik, Cook, Winfree, and Bruck 2008: Allowing any error $\epsilon > 0$ makes the model Turing universal.
- The known simulation is not robust.

Recall For $\lambda \in (0, \infty)$, the exponential distribution with rate λ is the probability measure on $[0, \infty)$ given by the c.d.f.

$$F_{\lambda}: [0, \infty) \to [0, 1]$$

 $F_{\lambda}(t) = 1 - e^{-\lambda t}$

Recall For $\lambda \in (0, \infty)$, the exponential distribution with rate λ is the probability measure on $[0, \infty)$ given by the c.d.f.



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 $\Delta(Q) = \{ \text{probability measures on } Q \}.$

Semantics of a CTMC $C = (Q, \lambda, \pi)$.

• At each time $t \in [0, \infty)$, C is probabilistically in some state state(t).

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- At each time $t \in [0, \infty)$, C is probabilistically in some state state(t).
- state(0) is chosen according to π .
- For each state $q \in Q$,

$$\lambda_q = \sum_{r \in Q} \lambda(q, r)$$

is the rate out of q.

• If $\lambda_q=0$, then state q is terminal.

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- if $\lambda_q > 0$, then state q is nonterminal If C enters state q at time t, then the sojourn time for which C remains in state q before moving to a new state has exponential distribution with rate λ_q .

- If $\lambda_a = 0$, then state q is terminal.
- if $\lambda_q > 0$, then state q is nonterminal If C enters state q at time t, then the sojourn time for which C remains in state q before moving to a new state has exponential distribution with rate λ_q .

When ${\cal C}$ does move to a new state, this state is r with probability

$$p(q,r) = \frac{\lambda(q,r)}{\lambda_a}.$$

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Stochastic mass action semantics of a CRN ${\cal N}=(S,R)$

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- $Q = \mathbb{N}^S$
- $\pi \in \Delta(Q)$ is how you want to initialize the species counts in N.
- Still need $\lambda: \mathbb{N}^S \times \mathbb{N}^S \to [0, \infty)$.

Law of Mass Action (stochastic)

Intuition: The rate of a unary reaction

$$\rho \equiv X \xrightarrow{k} < \text{anything} >$$

in a state $\mathbf{x} \in \mathbb{N}^S$ is

$$rate_{\mathbf{x}}(\rho) = k\mathbf{x}(X).$$

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The rate of a binary reaction

$$\rho \equiv X + Y \xrightarrow{k} < anything >$$

in a state $\mathbf{x} \in \mathbb{N}^S$ and volume V is

$$rate_{\mathbf{x}}(\rho) = \begin{cases} k\mathbf{x}(X)\mathbf{x}(Y)/V & \text{if } X \neq Y \\ k\mathbf{x}(X)(\mathbf{x}(X) - 1)/V & \text{if } X = Y \end{cases}$$

The rate of a reaction $\rho = (\mathbf{r}, \mathbf{p}, k)$ in a state $\mathbf{x} \in \mathbb{N}^S$ is

$$\mathsf{rate}_{\mathbf{x}}(\rho) = \frac{k}{V^{\|\mathbf{r}\|-1}} \prod_{Y \in S} \frac{\mathbf{x}(Y)!}{(\mathbf{x}(Y) - \mathbf{r}(Y))!},$$

where V is the volume of the solution and

$$\|\mathbf{r}\| = \sum_{Y \in \mathcal{C}} \mathbf{r}(Y).$$

Gillespie, 1977

In the CTMC $C=(\mathbb{N}^S,\lambda,\pi)$ for a CRN N=(S,R) we define the rate matrix $\lambda:Q\times Q\to [0,\infty)$ by

$$\lambda(\mathbf{x}, \mathbf{y}) = \sum \left\{ rate_{\mathbf{x}}(\rho) \mid \rho \in R \ takes \ \mathbf{x} \ to \ \mathbf{y} \right\}.$$

We have now completely specified the most often-used semantics (kinetics) for CRNs, namely,
deterministic mass action
and

stochastic mass action.

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Caution: These are sometimes called

stochastic, respectively.

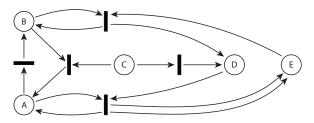
mass action

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1. Petri nets (Petri 1962).

A Petri net is a directed bipartite graph whose vertices are places (circles) and transitions (bars).

Example (CSWB 2009)



A state of a Petri net is a placement of zero or more tokens, all of which are identical, on its places.

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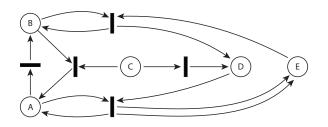
A transition is enabled if there is a token available for each of its incoming edges.

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In this case, the transition may fire, causing these tokens to be removed and tokens to be placed at the targets of its outgoing edges.

It is routine to check that the Petri net



has the same behavior as the CRN

$$A \rightarrow B$$
 $A + D \rightarrow A + 2E$ $C \rightarrow D$ $B + C \rightarrow A$

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2. Vector addition systems (Karp and Miller 1969)

Fix $n \in \mathbb{Z}^+$ and a finite set V of vectors $\mathbf{v} \in \mathbb{Z}^n$. Consider walks in \mathbb{N}^n that start at a location $\mathbf{x}_0 \in \mathbb{N}^n$ and proceed in steps, each of which adds a vector in V to the current location, leading to a new location in \mathbb{N}^n .

Example (CSWB 2009) The CRN

$$A \rightarrow B$$
 $A + D \rightarrow A + 2E$ $C \rightarrow D$ $B + C \rightarrow A$

is equivalent to the vector addition system

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3. Population protocols (Angluin, Aspnes, and Eisenstat 2008)

These are essentially CRNs in which every reaction has exactly two reactants and two products.

Note that this fixes the total population.

Example (AAE 2008) The population protocol

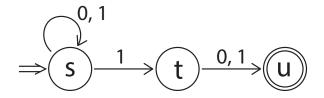
$$X + Y \rightarrow X + B$$
 $X + B \rightarrow X + X$
 $X + Y \rightarrow Y + B$ $Y + B \rightarrow Y + Y$

computes approximate majority.

A CRN That Simulates an NFA

with Titus Klinge and Jim Lathrop

Let M be the NFA



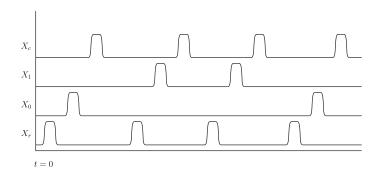
Objective: Design a CRN ${\cal N}$ that simulates ${\cal M}$ in deterministic mass action.

A CRN That Simulates an NFA

M gets its input $w \in \{0,1\}^*$ sequentially, so N should get its input sequentially.

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The user will provide an input $w = b_0 b_1 \cdots b_{n-1} \in \{0, 1\}^n$ approximately in the square-wave form



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We thus have input species X_0, X_1, X_c, X_r .

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- ullet state species Y_s, Y_t, Y_u ,
- portal species Z_s, Z_t, Z_u , and
- dual species $\overline{Y_s}, \overline{Y_t}, \overline{Y_u}, \overline{Z_s}, \overline{Z_t}, \overline{Z_u}$.

• For each state q,

$$X_r + Z_q \to X_r + \overline{Z_q}$$
.

reset

For each state q,

$$X_r + Z_q \to X_r + \overline{Z_q}$$
. reset

• For each transition $q \xrightarrow{b} r$

$$X_b + Y_q + \overline{Z_r} \to X_b + Y_q + Z_q.$$

compute

For each state q,

$$X_r + Z_q \to X_r + \overline{Z_q}.$$
 reset

• For each transition $q \xrightarrow{b} r$

$$X_b + Y_q + \overline{Z_r} \to X_b + Y_q + Z_q.$$
 compute

• For each state q

$$X_c + Z_q + \overline{Y_q} \to X_c + Z_q + Y_q$$
 copy $X_c + \overline{Z_q} + Y_q \to X_c + \overline{Z_q} + \overline{Y_q}$

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• For each state q,

$$\begin{split} 2Y_q + \overline{Y_q} &\to 3Y_q \\ 2\overline{Y_q} + Y_q &\to 3\overline{Y_q} \end{split}$$

error correction

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Thank you!