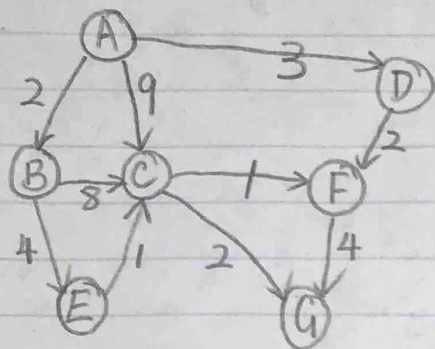


1. (C, G) (F, G) (B, E) (E, C) (C, F) (B, C) (D, F) (A, D) (A, C) (A, B)



Iteration 1: $\begin{array}{c|c|c|c|c|c|c|c} 0 & \infty & \infty & \infty & \infty & \infty & \infty & \infty \\ \hline A & B & C & D & E & F & G & \end{array}$

\Downarrow

$\begin{array}{c|c|c|c|c|c|c|c} 0 & 2 & 9 & 3 & \infty & \infty & \infty & \infty \\ \hline A & B & C & D & E & F & G & \end{array}$

Iteration 2: $\begin{array}{c|c|c|c|c|c|c|c} 0 & 2 & 7 & 3 & \infty & \infty & \infty & \infty \\ \hline A & B & C & D & E & F & G & \end{array} \Rightarrow \begin{array}{c|c|c|c|c|c|c|c} 0 & 2 & 7 & 3 & 6 & 5 & 11 & \infty \\ \hline A & B & C & D & E & F & G & \end{array}$

Iteration 3: $\begin{array}{c|c|c|c|c|c|c|c} 0 & 2 & 7 & 3 & 6 & 5 & 11 & \infty \\ \hline A & B & C & D & E & F & G & \end{array} \Rightarrow \begin{array}{c|c|c|c|c|c|c|c} 0 & 2 & 7 & 3 & 6 & 5 & 9 & \infty \\ \hline A & B & C & D & E & F & G & \end{array}$

Iteration 4: $\begin{array}{c|c|c|c|c|c|c|c} 0 & 2 & 7 & 3 & 6 & 5 & 9 & \infty \\ \hline A & B & C & D & E & F & G & \end{array} \Rightarrow \begin{array}{c|c|c|c|c|c|c|c} 0 & 2 & 7 & 3 & 6 & 5 & 9 & \infty \\ \hline A & B & C & D & E & F & G & \end{array}$

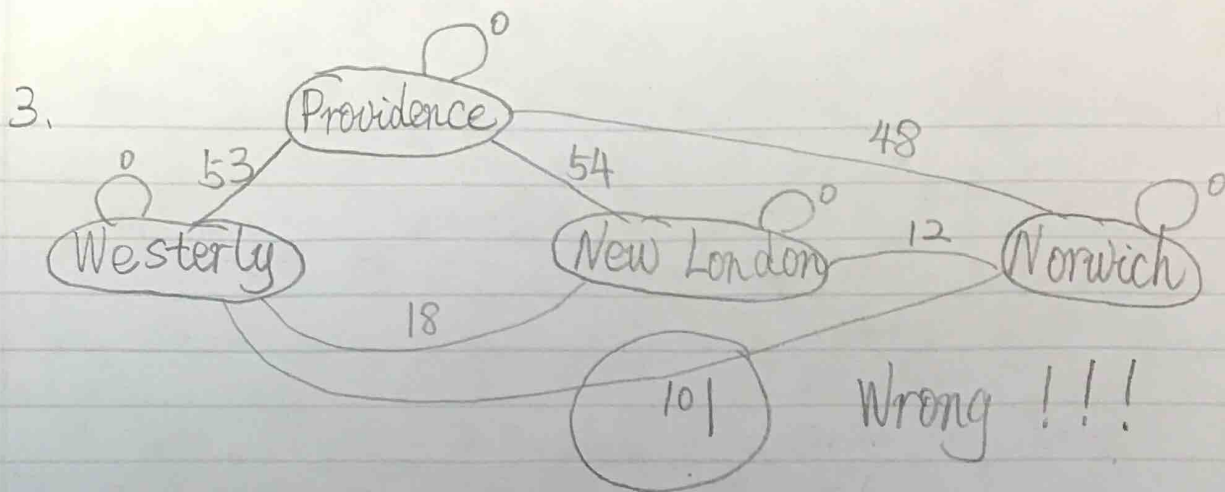
Iteration 5: $\begin{array}{c|c|c|c|c|c|c|c} 0 & 2 & 7 & 3 & 6 & 5 & 9 & \infty \\ \hline A & B & C & D & E & F & G & \end{array} \Rightarrow \begin{array}{c|c|c|c|c|c|c|c} 0 & 2 & 7 & 3 & 6 & 5 & 9 & \infty \\ \hline A & B & C & D & E & F & G & \end{array}$

Iteration 6: $\begin{array}{c|c|c|c|c|c|c|c} 0 & 2 & 7 & 3 & 6 & 5 & 9 & \infty \\ \hline A & B & C & D & E & F & G & \end{array} \Rightarrow \begin{array}{c|c|c|c|c|c|c|c} 0 & 2 & 7 & 3 & 6 & 5 & 9 & \infty \\ \hline A & B & C & D & E & F & G & \end{array}$

2. ① Initialize distance $[] = \{-\infty, -\infty, \dots\}$ and $\text{dist}[s] = 0$
 ② Create a topological order for all vertices
 ③ Do following for every vertex u in topological order:
 Do following for every adjacent vertex v to u :
 if ($\text{distance}[v] < \text{distance}[u] + \text{weight}(u, v)$)
 $\text{distance}[v] = \text{distance}[u] + \text{weight}(u, v)$

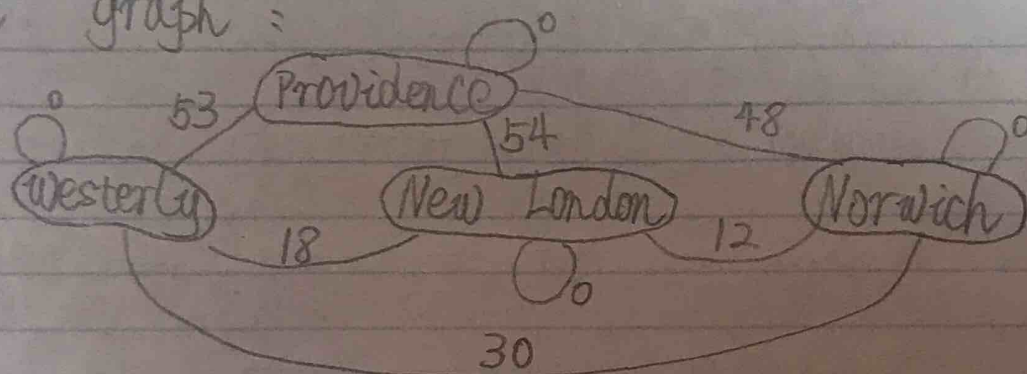
Auxiliary data structure used = topological order.

Running Time = • topological sort for vertex V and edges E is $O(V+E) \iff$ step ②
 • For every vertices $O(V)$, it checks every adjacent vertices. Total adjacent vertices is $O(E)$, so it runs $O(V+E)$ times. \iff step ③
 \therefore the running time is $O(V+E)$



From Westerly to Norwich shortest distance shouldn't be 101, we can go westerly - New London - Norwich, so the shortest distance is $18 + 12 = 30$

\therefore correct graph :



4. (a) Boolean or : if there is 1 result is 1

Boolean and : if there is 0 result is 0

Because $M^2(i, j)$ is first take boolean and, then take boolean or, so if $M^2(i, j) = 1$ means there must be a path between vertex i and j and between i and j we can take one intermediate vertex which can be between 1 to n . If $M^2(i, j) = 0$ means there is no paths between vertex i and j .

$$4. (b) \quad M^2(i, j) = (M(i, 1) \cdot (1, j)) + \dots + (M(i, n) \cdot M(n, j))$$

$$M^2(i, j) = (M(i, 1_2) \cdot (1_2, j)) + \dots + (M(i, n_2) \cdot M(n_2, j))$$

$$M^4(i, j) = (M(i, 1) \cdot M(1, 1_2) \cdot M(1_2, 1_3) \cdot M(1_3, j)) + \dots + (M(i, n) \cdot M(n, n_2) \cdot M(n_2, n_3) \cdot M(n_3, j))$$

So $M^4(i, j) = 1$ means there must be a path between vertex i and j , and between i and j we can take three intermediate vertices in set $V(1, \dots, n)$. $M^4(i, j) = 0$ means there is no paths between i and j .

4. (b) $M^k(i, j) = 1$ means there must be a path between i and j , and between i and j we can take $k-1$ intermediate vertices in set $V(1 \dots n)$. $M^k(i, j) = 0$ means there is no paths between i and j .

5. $M(i, j) = 0$, $i = j$
 $= \text{weight}(i, j)$, if there is an edge between i, j
 $= \infty$, otherwise

According to question 4, we can see $M^2(i, j) = d$ means it takes the min. of all paths between vertex i and j , and taking one intermediate vertex which between 1 to n . So if $M^2(i, j) = d$, we can find the minimum path $= d$ between vertex i and j , and taking one intermediate vertex which between 1 to n .

According to question 4, we can see $M^k(i, j) = d$ means it takes the min. of all paths between vertex i and j , and taking $k-1$ intermediate vertices in set $V(1 \dots n)$. So if $M^k(i, j) = d$, we can find the minimum path $= d$ between vertex i and j , and taking $k-1$ intermediate vertices in set $V(1 \dots n)$.