

Simulation and Optimization

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Welcome

This is the website for the **Simulation and Optimization** book that will teach you the basics of simulation approaches and optimization techniques in the context of modern AI systems. The source code of the book and the examples are provided as open-source and free to use [Creative Commons Attribution-NonCommercial-NoDerivs 4.0](#) license.

Preface

This book was created as companion material for a semester graduate course on simulation and optimization. It is the author's opinion that in an age of rapid advances in the field of artificial intelligence, it is of utmost importance to focus not only on machine learning, but to study in detail the techniques that make current advances in AI possible. From those, the areas of simulation and optimization have the highest potential to reveal how intricate current AI is intertwined with other areas of mathematics, statistics and computer science.

Simulation techniques are widely used in many scientific disciplines, ranging from climate models, epidemiology, and engineering to finance and logistics. These methods allow researchers and practitioners to analyze complex systems, evaluate scenarios, and make informed decisions when analytical solutions are infeasible or unavailable. Throughout this book, we will explore foundational concepts and practical approaches to simulation and optimization, providing both theoretical background and hands-on examples. In the context of AI, simulation approaches can be used to produce synthetic data for training in situations where these data are scarce, expensive, or simply impossible to collect. Another uses of simulation approaches include stress-testing algorithms, validating models under various hypothetical scenarios, and supporting decision-making in uncertain environments. By leveraging simulation, practitioners can gain insights into system behavior, identify potential risks, and optimize performance before deploying solutions in real-world settings.

Optimization approaches lie at the core of how machine learning is used in modern AI systems. Foundational algorithms like stochastic gradient descent make it possible to find optimal parameters for machine learning models using training datasets composed of millions of data points. Additionally, optimization algorithms are used for hyperparameter tuning and can be found at the heart of classical approaches like support vector machines and logistic regression. In this context, both classical and metaheuristic approaches play a pivotal role in finding optimal or near-optimal solutions which are used in the broader context of specific applications in practice.

Throughout the book, we will assume that the reader has familiarity with linear algebra and calculus and possesses a good command of statistics and the basics of machine learning. Additionally, good background knowledge of the Python programming language is advised for the practical part of this book.

1 Introduction

Simulation and optimization approaches are present in our everyday lives, albeit most of the time operating in a background plane. For example, when navigating with a GPS, the system simulates different routes and optimizes for the shortest or fastest path. Similarly, supply chains use optimization algorithms to minimize costs and maximize efficiency, while simulations help predict demand and manage inventory. These techniques are fundamental tools in decision-making processes across various industries, from transportation and logistics to finance and healthcare.

But what do simulation and optimization approaches have in common, apart from being complementary tools? The answer lies in the concept of a *model*. In the context of machine learning, we normally refer to a model as a mathematical or computational representation that captures the relationships between input data and output predictions. In simulation and optimization, a model similarly serves as an abstraction of a real-world system or process, allowing us to analyze, predict, and improve its behavior through experimentation and algorithmic techniques.

In the following, we will delve deeper into the concept of a model and how models are used in simulation and optimization contexts using some practical examples.

1.1 Example 1: Simulating supermarket dynamics

Imagine you are in your favourite grocery store waiting at the checkout queue. For simplicity, let's assume there is only one open counter. When you arrive at the queue, there might be other customers already waiting, while the first customer at the queue is currently being served. Shortly after you, a new customer arrives, taking the next free spot right behind you. And then another customer arrives, and another one, and another one...

Let's try to break down how this system behaves and what are the most important interactions between the parts of the system. In general, we will distinguish between *components*, *states*, *events*, *inputs* and *metrics*.

- **Components:** These are the entities that interact with each other. In our example, we have customers, cashiers and the queue itself.

- **States:** The configurations of the system that represent valid combinations of specific properties of the components at a given moment of time. For instance, at each time the queue has a specific length: zero if it's empty, one customer, two customers, etc. Additionally, the cashier can be busy or idle. We can also count the number of customers currently present in the supermarket which have not yet arrive at the checkout queue.
- **Events:** The interactions themselves, like a new customer arriving at the queue, checkout start or checkout completion.
- **Inputs:** Whatever information is fed into the system, e.g. arrival times, service times, etc. These inputs can contain statistical assumptions, like the distribution of arrival times.
- **Metrics:** How we evaluate the system as a whole in a given time step. For instance, what is the average waiting time? How much time are the cashiers busy? How is the queue length distributed?

The system could be represented by the following Python code as a minimal variant.

```
import heapq, random

# event = (time, type, customer_id)
event_list = []
heapq.heappush(event_list, (first_arrival_time, 'arrival', 1))

while event_list and time < sim_end:
    time, ev_type, cid = heapq.heappop(event_list)
    if ev_type == 'arrival':
        if any_cashier_free():
            start_checkout(cid, time)
            heapq.heappush(event_list, (time + service_time(cid),
                                         'departure', cid))
        else:
            enqueue(cid, time)
            heapq.heappush(event_list, (time + next_interarrival(),
                                         'arrival', next_id()))
    elif ev_type == 'departure':
        finish_service(cid, time)
        if queue_not_empty():
            next_cid = dequeue()
            start_service(next_cid, time)
            heapq.heappush(event_list, (time + service_time(next_cid),
                                         'departure', next_cid))
```

This code assumes that customers arrive at regular subsequent intervals after each arrival event. The parameter `sim_end` defines how long (how many steps) we want to simulate in this case. The function `service_time` returns the time the cashier needs for checking out customer `cid`.

The next customer will arrive after a time given by the function `next_interarrival`, which can implement different stochastic behaviours.

We can represent this system graphically as shown in the following illustration:

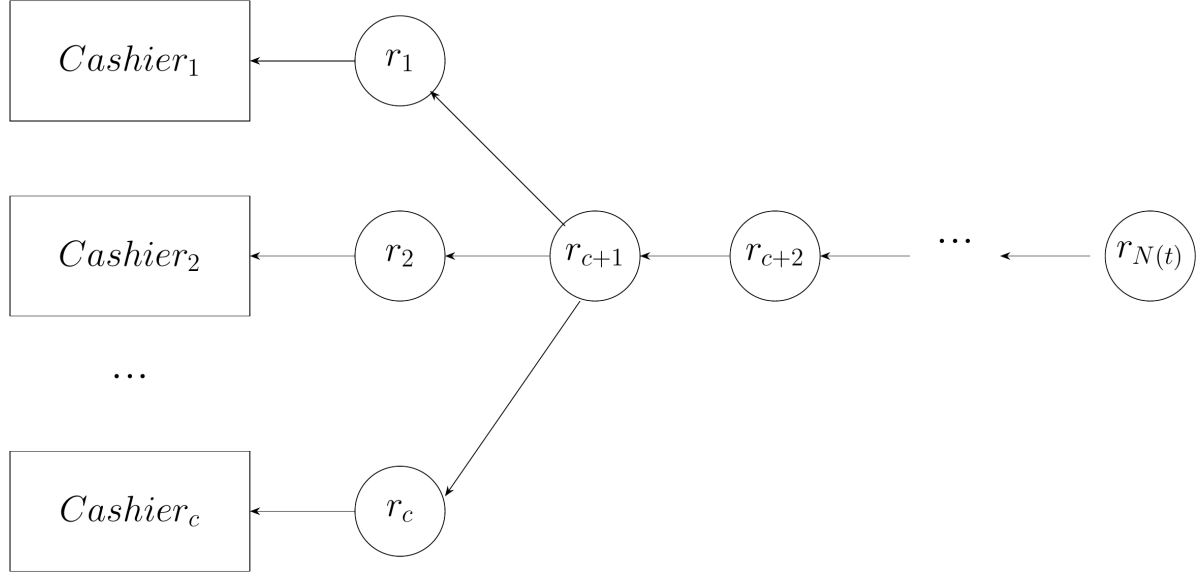


Figure 1.1: An illustration of the supermarket dynamics.

In this figure, customers are denoted by r_i , the amount of cashiers is c and the total number of customers in the supermarket at time t is denoted by $N(t)$.

Now let's try to refine the dynamics of this system. We will now write some equations to describe the system's dynamics according to the **infinite waiting room** $M/M/c$ model. Let's make the following assumptions:

- Arrivals follow a Poisson distribution with mean λ (arrivals per second), which for this case will be assumed to be stationary.
- The service times are assumed to be exponentially distributed with mean $1/\mu$.

What would be now the *traffic intensity per cashier*? That is, what is the mean customer flow that each cashier experiences from their own point of view? Let's call this number ρ and calculate it as follows:

$$\rho = \frac{\lambda}{c\mu}$$

In words, if customers arrive at a rate of $\lambda = 10$ customers/s and each cashier serves 2 customers/s (yes, it's a fast supermarket). With 5 cashiers, that means that $\rho = 10/5 \times 2 = 1$. This means that each cashier has quite a lot to do right now.

We are now interested in the probabilities of the states in this systems. In this case, we define a system by the number of customers currently present in the supermarket. So we can have $N = 1$ if there is currently 1 customer present, or any other number of customers (we assume the supermarket is so large, we can accomodate any number of them). Let's denote these probabilities by $p_n = \Pr\{N = n\}$. We have:

$$p_n = \lim_{t \rightarrow \infty} \int_0^t \mathbb{1}_{\{N(s)=n\}} ds$$

Intuitively, p_n represents the fraction of time where the supermarket has exactly n customers. As mentioned earlier, we will assume that arrivals do not depend of the current state n , so we write $\lambda_n = \lambda$ for all $n \geq 0$. However, note that the completion rates μ_n do *indeed* depend of the current state. To see this, imagine that there is only one customer in the supermarket ($n = 1$). The completion rate is then $\mu_1 = \mu$ since the only one cashier is needed to perform checkout. However, if there are $n = 2$ customers in the supermarket, two cashiers can serve those two customers in parallel, increasing the completion rate to $\mu_2 = 2\mu$. The same reasoning applies until $n = c$, the total number of cashiers. In this case, $\mu_c = c\mu$ and the next customer will have to wait in the queue. So we have:

$$\begin{aligned} \lambda_n &= \lambda \text{ for all } n \geq 0 \\ \mu_n &= \min(n, c)\mu \end{aligned}$$

Now we are going to state our **main modeling assumption**. Consider how we transition *between states*. Specifically, we transition from state n to state $n + 1$ when a new customer enters the supermarket, and there were already n customers in it. Similarly, we transition from state n to state $n - 1$ when a customer leaves the supermarket (in this case, all customers are served by the cashiers, so there is no way you leave the supermarket without paying first). Remember that the rate of customers arriving at the supermarket is always λ , and the rate of customers being served (i.e. leaving) when at state n is μ_n . In general, for each state we can define an *incoming* and an *outgoing flow*. This quantifies the transitions in resp. out of a given state.

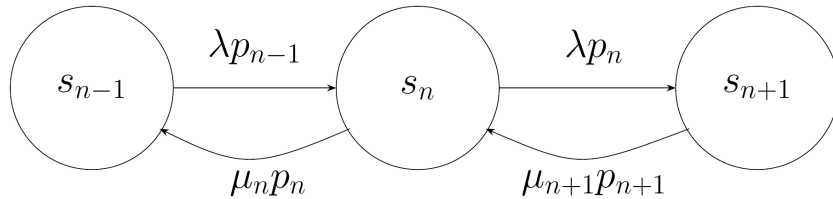


Figure 1.2: Transition dynamics between neighboring states

As can be seen in the previous figure, transitions flow away from state n in two ways: first, to state $n - 1$ when a customer is served with a rate $\mu_n p_n$ and to state $n + 1$ when a new customer

arrives with a rate λp_n . Similarly, one can transition from the other states to state n either by having a customer served in state $n + 1$ with a rate $\mu_{n+1} p_{n+1}$ or when being in state $n - 1$ and a new customer arrives with a rate λp_{n-1} . Our modeling assumption now is that, for each state n , the flow outwards balances out with the flow inwards (global balance):

$$\lambda p_{n-1} + \mu_{n+1} p_{n+1} = \lambda p_n + \mu_n p_n, \text{ for } n \geq 1$$

We can also write that, in the long term, the rate of transitions from n to $n + 1$ equals the transitions from $n + 1$ to n , which results in the more simple form (local balance)

$$\lambda p_n = \mu_{n+1} p_{n+1}, \text{ for } n \geq 1$$

This form follows from the global balance condition when only neighboring states are connected. Now, using this short form, we can provide closed-form expressions for different probabilities, using the following recursion (which directly follows from the above)

$$p_{n+1} = \frac{\lambda}{\mu_{n+1}} p_n$$

For instance, one can calculate that the probability of a customer having to wait (because all cashiers are busy at the moment) is

$$P_W = p_0 \frac{(\lambda/\mu)^c}{c!} \frac{1}{1 - \rho}$$

This is also called the *Erlang-C* probability and we will delve deeper into the details in the coming chapters.

We can now write computer code that performs a step-by-step simulation of the system (i.e. in discrete time steps). This is specially useful if we are in a situation where there is no closed-form analytical solution, or the analytical solution is too complex to calculate. For instance, we can run the above simulator for a large number of steps (say $T = 10^6$) and then calculate specific metrics like:

- Mean queue length.
- Mean waiting time.
- Mean time in the system.
- Total fraction of time that the cashiers were busy.
- Overall utilization.

We will see examples of such simulators in the first part of the book.

1.2 Example 2: The traveling salesman

Let's not turn our attention to the other type of problems which are central to this book: *optimization problems*. Imagine you are a sales representative for a vacuum cleaner manufacturer. Your task is to visit potential customers in cities across your area and, at the end of the day, return to where you started your journey. As an environmental conscious employer and in order to save transport costs, your company introduces the restriction that each customer in the route has to be visited *exactly once*. So you need to think carefully before getting into your car and starting your route.

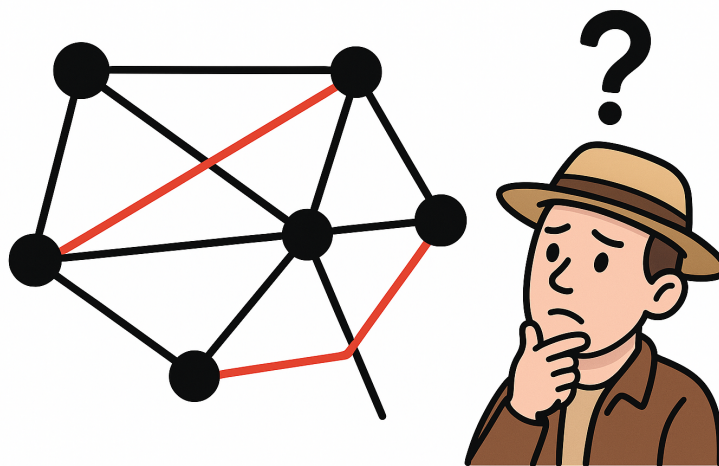


Figure 1.3: The traveling salesman has to find a good and practical solution

Now, the traveling salesman needs to consider two things:

- A method for *constructing valid* tours that start and end in the same location.
- A method for *evaluating* those tours so that we can quantitatively decide if a tour is better than another one.

Let's address each of these considerations in detail. Assume that $N = \{1, 2, \dots, n\}$ is our set of possible locations. Let's define the following variables:

$$x_{ij} = \begin{cases} 1 & \text{if the tour goes directly to location } i \text{ to location } j \\ 0 & \text{otherwise} \end{cases}$$

where $i, j \in N$. We call x_{ij} our *decision variables*. So if the traveling salesman specifies the value of each x_{ij} , we have a candidate route to consider. However, not every assignment of

the x_{ij} variables to $\{0, 1\}$ will make sense for the traveling salesman. For instance, imagine that we have in one assignment both $x_{23} = 1$ and $x_{43} = 1$. That would mean that location 3 is visited twice, once from location 2 and another time from location 4. That violates the requirement that each location is visited exactly once.

To model this situation, we need to introduce *constraints*. In optimization problems, constraints take usually the form of equalities or inequalities as functions of the decision variables. In our case, the requirement that each location is visited only once can be expressed by the following (linear) equalities:

$$\begin{aligned}\sum_{j=1}^n x_{ij} &= 1 \text{ for all } i \in N \\ \sum_{i=1}^n x_{ij} &= 1 \text{ for all } j \in N\end{aligned}$$

The first equality means that, fixed a location i , the sum of all *outgoing* edges is exactly one. Conversely, the second states that for a fixed location j , the sum of all *incoming* edges is also exactly one. This ensures that each location is visited exactly once.

We need another technical condition to guarantee that the tour is a single one and not composed of multiple sub-tours. There should be no subset $S \subset N$ such that is self contained, the number of visited cities equals exactly its size. Mathematically:

$$\sum_{i \in N} \sum_{j \in N} x_{ij} \leq |S| - 1 \text{ for all subsets } S \subset N$$

To sum up, we now have modeled how valid tours should look like. If we find a tour $T = \{x_{ij}\}$ that satisfies the constraints outlined before, we can be sure it is a valid tour.

But surely there are some tours that are better than others? This is where the second issue becomes important: we need an evaluation method to distinguish between good and bad solutions. In the optimization literature, we normally talk about **objective functions**. In our case, the traveling salesman would like the total distance to be *minimized*, meaning the sum of all distances between locations of the tour. Assume that $c_{ij} > 0$ is the distance between location i and j (for consistency assume $c_{ii} = 0$). Now we want to minimize the distance. For this we write:

$$\min \sum_{i \in N} \sum_{j \in N} c_{ij} x_{ij}$$

1.3 Exercises

1. Prove that in the supermarket example the local balance condition follows from the global balance condition (Hint: use induction).

Part I

PART I: SIMULATION

2 Simulation basics

3 Monte Carlo

4 Discrete events

5 Agent-based simulation

Part II

PART II: OPTIMIZATION

6 Optimization basics

7 Gradient descent

8 Evolutionary algorithms

9 Summary

In summary, this book has no content whatsoever.

References