

**Bayesuvius,**  
a small visual dictionary of Bayesian Networks

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Figure 1: View of Mount Vesuvius from Pompeii



Figure 2: Mount Vesuvius and Bay of Naples

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## 0.1 Foreword

Welcome to Bayesuvius! a proto-book uploaded to github.

A different Bayesian network is discussed in each chapter. Each chapter title is the name of a B net. Chapter titles are in alphabetical order.

This is a volcano in its early stages. First version uploaded to a github repo called Bayesuvius on June 24, 2020. First version only covers 2 B nets (Linear Regression and GAN). I will add more chapters periodically. Remember, this is a moonlighting effort so I can't do it all at once.

For any questions about notation, please go to Notational Conventions section.

Requests and advice are welcomed.

Thanks for reading this.

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## 0.2 Notational Conventions

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bnet=Bayesian Network

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Define  $\mathbb{Z}, \mathbb{R}, \mathbb{C}$  to be the integers, real numbers and complex numbers, respectively.

For  $a < b$ , define  $\mathbb{Z}_I$  to be the integers in the interval  $I$ , where  $I = [a, b], [a, b), (a, b], (a, b)$  (i.e.,  $I$  can be closed or open on either side).

$A_{>0} = \{k \in A : k > 0\}$  for  $A = \mathbb{Z}, \mathbb{R}$ .

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Random Variables will be indicated by underlined letters and their values by non-underlined letters. Each node of a bnet will be labelled by a random variable. Thus,  $\underline{x} = x$  means that node  $\underline{x}$  is in state  $x$ .

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$P_{\underline{x}}(x) = P(\underline{x} = x) = P(x)$  is the probability that random variable  $\underline{x}$  equals  $x \in S_{\underline{x}}$ .  $S_{\underline{x}}$  is the set of states (i.e., values) that  $\underline{x}$  can assume and  $n_{\underline{x}} = |S_{\underline{x}}|$  is the size (aka cardinality) of that set. Hence,

$$\sum_{x \in S_{\underline{x}}} P_{\underline{x}}(x) = 1 \quad (1)$$


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$$P_{\underline{x}, \underline{y}}(x, y) = P(\underline{x} = x, \underline{y} = y) = P(x, y) \quad (2)$$


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$$P_{\underline{x}|\underline{y}}(x|y) = P(\underline{x} = x | \underline{y} = y) = P(x|y) = \frac{P(x, y)}{P(y)} \quad (3)$$


---

Kronecker delta function: For  $x, y$  in discrete set  $S$ ,

$$\delta(x, y) = \begin{cases} 1 & \text{if } x = y \\ 0 & \text{if } x \neq y \end{cases} \quad (4)$$


---

Dirac delta function: For  $x, y \in \mathbb{R}$ ,

$$\int_{-\infty}^{+\infty} dx \delta(x - y) f(x) = f(y) \quad (5)$$


---

Transition probability matrix of a node of a bnet can be either a discrete or a continuous probability distribution. To go from continuous to discrete, one replaces integrals over states of node by sums over new states, and Dirac delta functions by Kronecker delta functions. More precisely, consider a function  $f : S \rightarrow \mathbb{R}$ . Let  $S_{\underline{x}} \subset S$  and  $S \rightarrow S_{\underline{x}}$  upon discretization (binning). Then

$$\int_S dx P_{\underline{x}}(x) f(x) \rightarrow \frac{1}{n_{\underline{x}}} \sum_{x \in S_{\underline{x}}} f(x) . \quad (6)$$

Both sides of last equation are 1 when  $f(x) = 1$ . Furthermore, if  $y \in S_{\underline{x}}$ , then

$$\int_S dx \delta(x - y) f(x) = f(y) \rightarrow \sum_{x \in S_{\underline{x}}} \delta(x, y) f(x) = f(y) . \quad (7)$$


---

Indicator function (aka Truth function):

$$\mathbb{1}(\mathcal{S}) = \begin{cases} 1 & \text{if } \mathcal{S} \text{ is true} \\ 0 & \text{if } \mathcal{S} \text{ is false} \end{cases} \quad (8)$$

For example,  $\delta(x, y) = \mathbb{1}(x = y)$ .

---


$$\vec{x} = (x[0], x[1], x[2] \dots, x[nsam(\vec{x}) - 1]) = x[:] \quad (9)$$

$nsam(\vec{x})$  is the number of samples of  $\vec{x}$ .  $x[i]$  are i.i.d. (independent identically distributed) samples with

$$x[i] \sim P_{\underline{x}} \text{ (i.e. } P_{x[i]} = P_{\underline{x}}) \quad (10)$$

$$P(\underline{x} = x) = \frac{1}{nsam(\vec{x})} \sum_i \mathbb{1}(x[i] = x) \quad (11)$$

If we use two sampled variables, say  $\vec{x}$  and  $\vec{y}$ , in a given bnnet, their number of samples  $nsam(\vec{x})$  and  $nsam(\vec{y})$  need not be equal.

---


$$P(\vec{x}) = \prod_i P(x[i]) \quad (12)$$

$$\sum_{\vec{x}} = \prod_i \sum_{x[i]} \quad (13)$$

$$\partial_{\vec{x}} = [\partial_{x[0]}, \partial_{x[1]}, \partial_{x[2]}, \dots, \partial_{x[nsam(\vec{x})-1]}] \quad (14)$$

---


$$P(\vec{x}) \approx \left[ \prod_x P(x)^{P(x)} \right]^{nsam(\vec{x})} \quad (15)$$

$$= e^{nsam(\vec{x}) \sum_x P(x) \ln P(x)} \quad (16)$$

$$= e^{-nsam(\vec{x}) H(P_{\underline{x}})} \quad (17)$$

---


$$f^{[1, \partial_x, \partial_y]}(x, y) = [f, \partial_x f, \partial_y f] \quad (18)$$

$$f^+ = f^{[1, \partial_x, \partial_y]} \quad (19)$$

---

For probabilty distributions  $p(x), q(x)$  of  $x \in S_{\underline{x}}$

- Entropy:

$$H(p) = - \sum_x p(x) \ln p(x) \geq 0 \quad (20)$$

- Kullback-Liebler divergence:

$$D_{KL}(p \parallel q) = \sum_x p(x) \ln \frac{p(x)}{q(x)} \geq 0 \quad (21)$$

- Cross entropy:

$$CE(p \rightarrow q) = - \sum_x p(x) \ln q(x) \quad (22)$$

$$= H(p) + D_{KL}(p \parallel q) \quad (23)$$

---

Normal Distribution:  $x, \mu, \sigma \in \mathbb{R}, \sigma > 0$

$$\mathcal{N}(\mu, \sigma^2)(x) = \frac{1}{\sigma\sqrt{2\pi}} e^{-\frac{1}{2}\left(\frac{x-\mu}{\sigma}\right)^2} \quad (24)$$


---

Uniform Distribution:  $a < b, x \in [a, b]$

$$\mathcal{U}(a, b)(x) = \frac{1}{b-a} \quad (25)$$


---

Expected Value

Given a random variable  $\underline{x}$  with states  $S_{\underline{x}}$  and a function  $f : S_{\underline{x}} \rightarrow \mathbb{R}$ , define

$$E_{\underline{x}}[f(\underline{x})] = E_{x \sim P(x)}[f(x)] = \sum_x P(x)f(x) \quad (26)$$


---

Conditional Expected Value

Given a random variable  $\underline{x}$  with states  $S_{\underline{x}}$ , a random variable  $\underline{y}$  with states  $S_{\underline{y}}$ , and a function  $f : S_{\underline{x}} \times S_{\underline{y}} \rightarrow \mathbb{R}$ , define

$$E_{\underline{x}|\underline{y}}[f(\underline{x}, \underline{y})] = \sum_x P(x|\underline{y})f(x, \underline{y}) , \quad (27)$$

$$E_{\underline{x}|\underline{y}=\underline{y}}[f(\underline{x}, \underline{y})] = E_{\underline{x}|\underline{y}}[f(\underline{x}, \underline{y})] = \sum_x P(x|\underline{y})f(x, \underline{y}) . \quad (28)$$

Note that

$$E_{\underline{y}}[E_{\underline{x}|\underline{y}}[f(\underline{x}, \underline{y})]] = \sum_{x,y} P(x|\underline{y})P(\underline{y})f(x, \underline{y}) \quad (29)$$

$$= \sum_{x,y} P(x, \underline{y})f(x, \underline{y}) \quad (30)$$

$$= E_{\underline{x}, \underline{y}}[f(\underline{x}, \underline{y})] . \quad (31)$$


---

Sigmoid function: For  $x \in \mathbb{R}$ ,

$$\text{sig}(x) = \frac{1}{1 + e^{-x}} \quad (32)$$


---

$\mathcal{N}(!a)$  will denote a normalization constant that does not depend on  $a$ . For example,  $P(x) = \mathcal{N}(!x)e^{-x}$  where  $\int_0^\infty dx P(x) = 1$ .

A **one hot** vector of zeros and ones is a vector with all entries zero with the exception of a single entry which is one. A **one cold** vector has all entries equal to one with the exception of a single entry which is zero. For example, if  $x^n = (x_0, x_1, \dots, x_{n-1})$  and  $x_i = \delta(i, 0)$  then  $x^n$  is one hot.

Short Summary of Boolean Algebra. See Ref.[1] for more info about this topic.

---

Suppose  $x, y, z \in \{0, 1\}$ . Define

$$x \text{ or } y = x \vee y = x + y - xy , \quad (33)$$



$$x \text{ and } y = x \wedge y = xy , \quad (34)$$

and

$$\text{not } x = \bar{x} = 1 - x , \quad (35)$$

where we are using normal addition and multiplication on right hand sides.

Associativity	$x \vee (y \vee z) = (x \vee y) \vee z$ $x \wedge (y \wedge z) = (x \wedge y) \wedge z$
Commutativity	$x \vee y = y \vee x$ $x \wedge y = y \wedge x$
Distributivity	$x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z)$ $x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z)$
Identity	$x \vee 0 = x$ $x \wedge 1 = x$
Annihilator	$x \wedge 0 = 0$ $x \vee 1 = 1$
Idempotence	$x \vee x = x$ $x \wedge x = x$
Absorption	$x \wedge (x \vee y) = x$ $x \vee (x \wedge y) = x$
Complementation	$x \wedge \bar{x} = 0$ $x \vee \bar{x} = 1$
Double negation	$\overline{(\bar{x})} = x$
De Morgan Laws	$\bar{x} \wedge \bar{y} = \overline{(x \vee y)}$ $\bar{x} \vee \bar{y} = \overline{(x \wedge y)}$

Table 1: Boolean Algebra Identities

Actually, since  $x \wedge y = xy$ , we can omit writing the symbol  $\wedge$ . The symbol  $\wedge$  is useful to exhibit the symmetry of the identities, and to remark about the analogous identities for sets, where  $\wedge$  becomes intersection  $\cap$  and  $\vee$  becomes union  $\cup$ . However, for practical calculations,  $\wedge$  is an unnecessary nuisance.

Since  $x \in \{0, 1\}$ ,

$$P(\bar{x}) = 1 - P(x) . \quad (36)$$

Clearly, from analyzing the simple event space  $(x, y) \in \{0, 1\}^2$ ,

$$P(x \vee y) = P(x) + P(y) - P(x \wedge y) . \quad (37)$$

# Chapter 1

## Back Propagation (Auto Differentiation): COMING SOON

## Chapter 2

# Basic Curve Fitting Using Gradient Descent



Figure 2.1: Basic curve fitting bnet.

Samples  $(x[i], y[i]) \in S_{\underline{x}} \times S_{\underline{y}}$  are given.  $nsam(\vec{x}) = nsam(\vec{y})$ .

Estimator function  $\hat{y}(x; \phi)$  for  $x \in S_{\underline{x}}$  and  $\phi \in \mathbb{R}$  is given.

Let

$$P_{\underline{x}, \underline{y}}(x, y) = \frac{1}{nsam(\vec{x})} \sum_i \mathbb{1}(x = x[i], y = y[i]) . \quad (2.1)$$

Let

$$\mathcal{E}(\vec{x}, \vec{y}, \phi) = \frac{1}{nsam(\vec{y})} \sum_i |y[i] - \hat{y}(x[i]; \phi)|^2 \quad (2.2)$$

$\mathcal{E}$  is called the mean square error.

Best fit is parameters  $\phi^*$  such that

$$\phi^* = \underset{\phi}{\operatorname{argmin}} \mathcal{E}(\vec{x}, \vec{y}, \phi) . \quad (2.3)$$

The node transition matrices for the basic curve fitting bnet Fig.2.1 are printed below in blue.

$$P(\phi) = \text{given} . \quad (2.4)$$

The first time it is used,  $\phi$  is arbitrary. After the first time, it is determined by previous stage.

$$P(\vec{x}) = \prod_i P_{\underline{x}}(x[i]) \quad (2.5)$$

$$P(\vec{y}|\vec{x}) = \prod_i P_{\underline{y}|\underline{x}}(y[i] | x[i]) \quad (2.6)$$

$$P(\hat{y}[i]|\phi, \vec{x}) = \delta(\hat{y}[i], \hat{y}(x[i]; \phi)) \quad (2.7)$$

$$P(\mathcal{E}|\vec{\hat{y}}, \vec{y}) = \delta(\mathcal{E}, \frac{1}{nsam(\vec{x})} \sum_i |y[i] - \hat{y}[i]|^2) . \quad (2.8)$$

$$P(\phi'|\phi, \mathcal{E}) = \delta(\phi', \phi - \eta \partial_{\phi} \mathcal{E}) \quad (2.9)$$

$\eta > 0$  is the descent rate. If  $\Delta\phi = \phi' - \phi = -\eta \frac{\partial \mathcal{E}}{\partial \phi}$ , then  $\Delta\mathcal{E} = \frac{-1}{\eta} (\Delta\phi)^2 < 0$  so this will minimize the error  $\mathcal{E}$ . This is called “gradient descent”.

## Chapter 3

# Bell and Clauser-Horne Inequalities in Quantum Mechanics



Figure 3.1: bnet used to discuss Bell and Clauser-Horne inequalities in Quantum Mechanics.

I wrote an article about this in 2008 for my blog “Quantum Bayesian Networks”. See Ref.[4].

# Chapter 4

## Binary Decision Diagrams

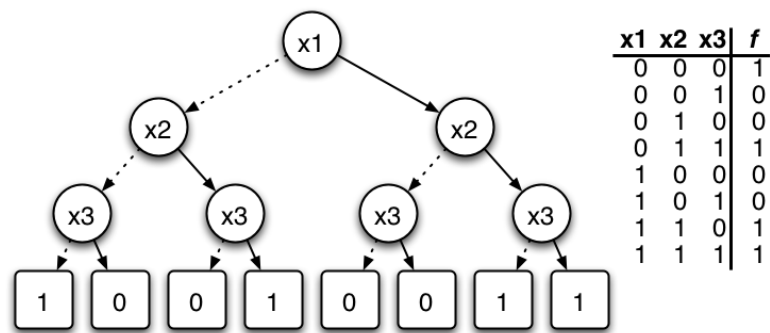


Figure 4.1: Binary decision tree and truth table for the function  $f(x_1, x_2, x_3) = \bar{x}_1(x_2 + \bar{x}_3) + x_1x_2$

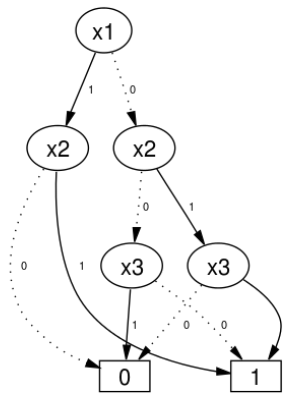


Figure 4.2: BDD for the function  $f$  of Fig.4.1.

This chapter is based on Wikipedia article Ref.[5].

Binary Decision Diagrams (BDDs) can be understood as a special case of Decision Trees (dtrees). We will assume that the reader has read Chapter 5 on dtrees before reading this chapter.

Both Figs.4.1 and 4.2 were taken from the aforementioned Wikipedia article. They give a simple example of a function  $f : \{0, 1\}^3 \rightarrow \{0, 1\}$  represented in Fig.4.1 as a **binary decision tree** and in Fig.4.2 as a **binary decision diagram (BDD)**. The goal of this chapter is to find for each of those figures a bnet with the same graph structure.

We begin by noting that the function  $f : \{0, 1\}^3 \rightarrow \{0, 1\}$  is a special case of a probability distribution  $P : \{0, 1\}^3 \rightarrow [0, 1]$ . In fact, if we restrict  $P$  to be deterministic, then  $P_{det} : \{0, 1\}^3 \rightarrow \{0, 1\}$  has the same domain and range as  $f$ . Henceforth, we will refer to  $f(x_1, x_2, x_3)$  as  $P(x_1, x_2, x_3)$ , keeping in mind that we are restricting our attention to deterministic probability distributions.

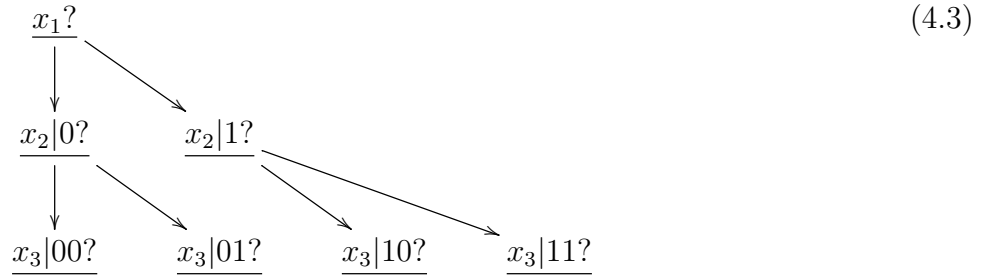
If we apply the chain rule for conditional probabilities to  $P(x_1, x_2, x_3)$ , we get

$$P(x_1, x_2, x_3) = P(x_3|x_1, x_2)P(x_2|x_1)P(x_1), \quad (4.1)$$

which can be represented by the bnet:



But in Chapter 5, we learned how to represent the bnet of Eq.(4.2) as the bnet tree Eq.(4.3). In that tree, the nodes pose questions with 3 possible answers 0, 1, *null*. In Eq.(4.3),  $x_2|a?$  stands for “what is  $x_2$  if  $x_1 = a$ ?” and  $x_3|a, b?$  stands for “what is  $x_3$  if  $x_1 = a, x_2 = b$ ?”.



The node transition probability matrices, printed in blue, for the bnet of Eq.(4.3) are as follows. If  $x_1, x_2, x_3 \in \{0, 1, null\}$  and  $a, b \in \{0, 1\}$ , then

$$P(\underline{x_1?} = x_1) = \begin{cases} P_{x_1}(x_1) & \text{if } x_1 \in \{0, 1\} \\ 0 & \text{if } x_1 = null \end{cases} \quad (4.4)$$

$$P(\underline{x_2|a?} = x_2 \mid \underline{x_1?} = x_1) = \begin{cases} P_{x_2|x_1}(x_2|a) & \text{if } x_1 = a \\ \mathbb{1}(x_2 = null) & \text{otherwise} \end{cases} \quad (4.5)$$

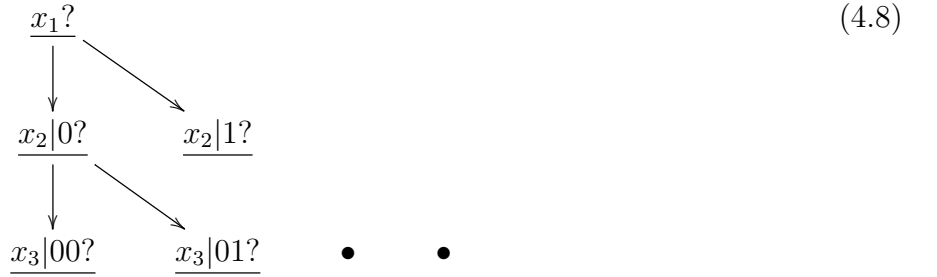
$$P(\underline{x_3|a, b? = x_3} \mid \underline{x_2|b? = x_2}) = \begin{cases} P_{\underline{x_3|\underline{x_1}, \underline{x_2}}}(x_3|a, b) & \text{if } (x_1, x_2) = (a, b) \\ \mathbb{1}(x_3 = null) & \text{otherwise} \end{cases} \quad (4.6)$$

The bnet shown in Eq.(4.3) contains the same info and has the same graph structure as the binary decision tree Fig.4.1. As when we were converting dtrees to their image bnets, the info in the endpoint nodes of Fig.4.1 is implicit in the transition matrices of the image bnet Eq.(4.3). If one wants to make the endpoint node info more explicit in the image bnet, one can add it to the descriptors of the state names of the leaf nodes of the image bnet. For example, one can add descriptors “gives  $f = 0$ ” or “gives  $f = 1$ ” to the “0” or “1” states of those leaf nodes.

The BDD shown in Fig.4.2 emphasizes the fact that

$$P(x_1, x_2, x_3 | x_1 = 1) = P(x_2 | x_1 = 1) = x_2 . \quad (4.7)$$

The BDD of Fig.4.2 corresponds to the bnet of Eq.(4.8).




---

What happens if we consider an  $f$  for which  $P(x_3|x_1, x_2) = P(x_3|x_2)$  so that one of the arcs of the fully connected bnet Eq.(4.2) is unnecessary? In that case,

$$P(x_1, x_2, x_3) = P(x_3|x_2)P(x_2|x_1)P(x_1) , \quad (4.9)$$

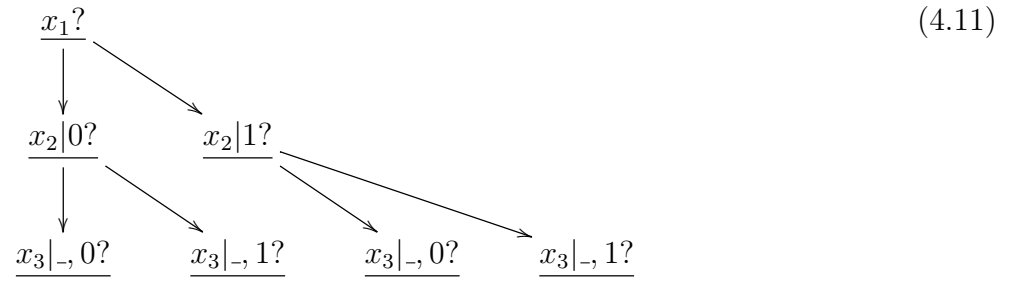
which can be represented by the Markov chain bnet:



Following the prescriptions of Chapter 5, we can represent the bnet of Eq.(4.10) as the bnet



tree Eq.(4.11). In that tree, the nodes pose questions with 3 possible answers 0, 1, *null*.



# Chapter 5

## Decision Trees

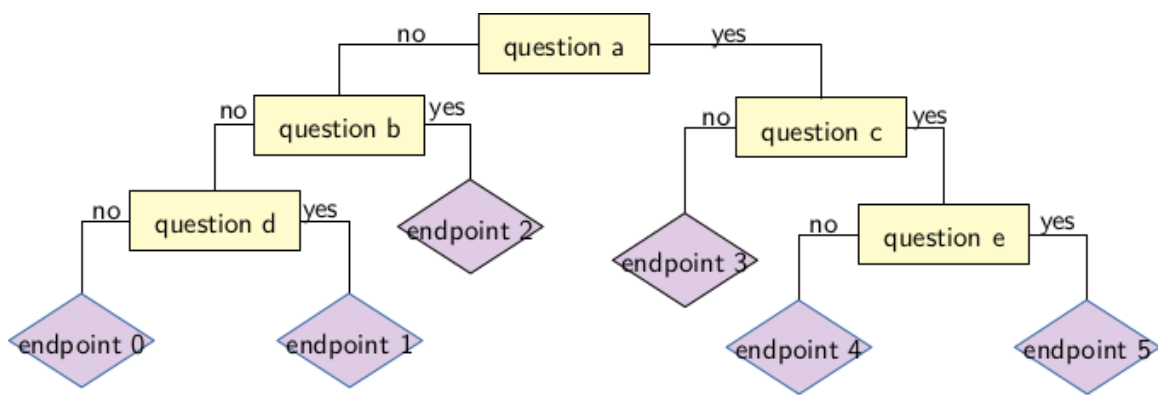


Figure 5.1: Typical decision tree.



Figure 5.2: Bnet corresponding to decision tree Fig.5.1

Fig.5.1 shows a typical decision tree (dtree). The yellow rectangles pose questions. In general, the answers to those questions can be multiple choices with more than two choices, but in Fig.5.1 we have chosen the simplest case of only two choices, true or false. The purple diamonds represent endpoints, goals, final conclusions, single states of reality, etc.

A trivial observation that is often not made in dtree educational literature is that every dtree maps into a special bnet, let's call it its “image” bnet, in a very natural way. To get the image bnet, just follow the following simple steps:

1. **Keep the yellow question nodes but reinterpret them as bnet nodes. Reinterpret the connections among the dtree question nodes as arrows pointing down from the root node.**

The image bnet nodes have 3 states,  $0 = no$  and  $1 = yes$  and  $null$ . Table 5.1 gives the node transition matrix  $[P(x|a)]_{x \in \{0,1,null\}, a \in \{0,1,null\}}$  where  $p_1 \in [0,1]$  can be different for each node and is given in the info that specifies the dtree. In Table 5.1,  $a_0 = 0$  if the dtree node being replaced has input “no” and  $a_0 = 1$  if its input is “yes”.  $!a_0$  means not  $a_0$  (i.e.,  $!a_0 = 1 - a_0$ ).

2. **This method of naming the image bnet nodes is not necessary but a good practice.** Give as name to each image bnet node an abridged version of the question that labels the dtree node it is replacing. Use as a suffix to the name of a bnet node either a 0 or a 1 depending whether the dtree node it is replacing has a 0 or a 1 as input. This suffix is not necessary because its info is already encoded into which column of the node transition prob matrix has zero probability for the *null* state, but it's a redundancy which makes the bnet easier to read and understand.
3. **Erase the purple endpoint nodes and connectors to them.** The info in each endpoint node can be preserved by storing it as a descriptor (e.g., tool tip) for the output states of the leaf node that is the parent to the endpoint in the image bnet. The endpoint info can be added to the descriptor of the  $no = 0$  state if the endpoint has 0 as input or to the descriptor of the  $yes = 1$  state if the endpoint has 1 as input.

$P(x a)$	$a = a_0$	$a = !a_0$	$a = null$
$x = 0$	$1 - p_1$	0	0
$x = 1$	$p_1$	0	0
$x = null$	0	1	1

Table 5.1: Transition probability matrix of a node of a dtree image bnet.

Table 5.2 describes the node types commonly used in dtrees.

When drawing dtrees, some people put info like explanations and probabilities on the connectors between the nodes of the dtree. That info can all be preserved in the transition prob matrix and the descriptors of the node names and node state names of the image bnet nodes. Often, the educational literature states that dtrees are more explicit and carry more info than their image bnets, but if one follows the above prescriptions, both can carry the same info.

A deterministic node commonly used in dtrees is one that asks the question  $x < \alpha?$ . for some real number  $\alpha \in (L, U)$  and some variable  $x$  (for example,  $x = \text{height of a person}$ ). For such an interval splitting node, the transition probability matrix would be as given in Table 5.3. If the interval  $[L, U]$  is binned into a number  $n_{bins}$  of bins, then this transition prob matrix will have dimensions  $(n_{bins} + 1) \times (\text{the number of states of the parent node})$ .

<b>dtree node types (usual shape in parenthesis)</b>	<b>their node transition probability matrix <math>P(x a)</math> in image bnet</b>
chance node (oval)	$P(x a)$ arbitrary. random
decision node (square)	$P(x a) = \delta(x, f(a))$ where $f(\cdot)$ is a function of $a$ . deterministic
endpoint node (diamond)	no $P(x a)$
fixed node	$P(x a) = \delta(x, x_0)$ . $x_0$ does not depend on $a$ whereas for decision node it does. deterministic.

Table 5.2: dtree node types.

$P(x a)$	states of parent node with $a = a_0$	states of parent node with $a \neq a_0$	$a = null$
$[x \in bin]_{\forall bin \subset [L, \alpha]}$	1	0	0
$[x \in bin]_{\forall bin \subset [\alpha, U]}$	0	0	0
$x = null$	0	1	1

Table 5.3: Transition probability matrix for interval splitting node.

A naive Bayes bnet (see Chapter 21) consists of a single “class” node that fans out with arrows pointing to other “feature” nodes. If each leaf node of a naive Bayes bnet fans out into a set of new leaf nodes, and those new leaf nodes also fan out and so on recursively, we get a tree bnet. The bnet that arises from this recursive application of naive Bayes has the same graph structure as the image bnet of a dtree. However, it is more general because its node transition prob matrices are more general; it has more weights (weight= parameters of node transition prob matrices) than a dtree with the same graph.

# Chapter 6

## Digital Circuits

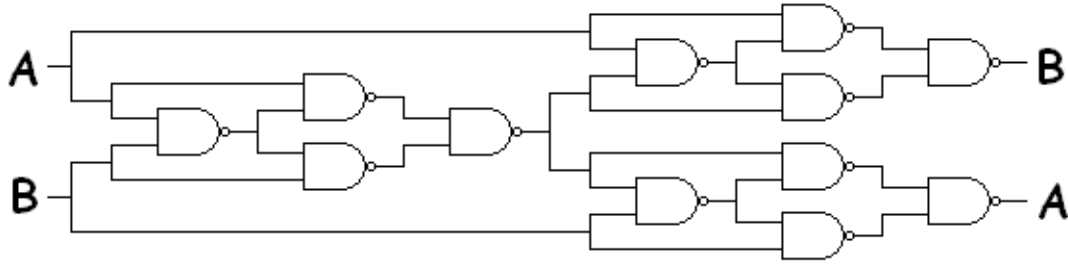


Figure 6.1: Typical digital circuit of NAND gates.

**Digital (logic) gate:** node with  $na$  input ports and  $nx$  output ports which represents a function

$$f : \{0, 1\}^{na} \rightarrow \{0, 1\}^{nx} . \quad (6.1)$$

Suppose

$a^{na} = (a_i)_{i=0,1,\dots,na-1}$  where  $a_i \in \{0, 1\}$ ,

$x^{nx} = (x_i)_{i=0,1,\dots,nx-1}$  where  $x_i \in \{0, 1\}$ .

$f$  maps  $a^{na}$  into  $x^{nx}$ .

**Digital circuit (dcircuit)** = circuit of digital gates.

---

**Mapping any dcircuit to a bnet (Option A- See Fig.6.2)):**

1. Replace every dcircuit gate described by Eq.(6.1) by  $nx$  bnet nodes  $\underline{x}_i$  for  $i = 0, 1, \dots, nx - 1$  such that

$$P(x_i | a^{na}) = \delta(x_i, f_i(a^{na})) \quad (6.2)$$

2. Replace all connectors of the dcircuit by arrows pointing in the direction of the bit flow.

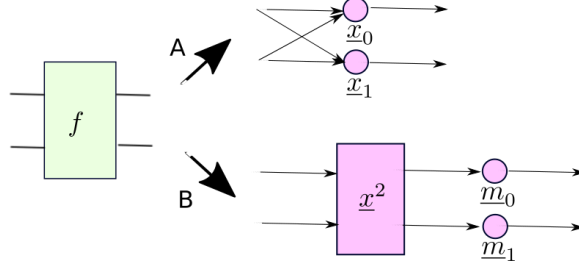


Figure 6.2: 2 options for mapping dcircuit node with multiple output ports into bnet.

---

**Mapping any dcircuit to a bnet (Option B- See Fig.6.2):**

1. Replace every dcircuit gate described by Eq.(6.1) with one bnet node called  $\underline{x}^{nx}$  and, if  $nx > 0$ ,  $nx$  “marginalizer nodes”  $\underline{m}_i$  for  $i = 0, 1, \dots, nx - 1$ , such that

$$P(x^{nx}|a^{na}) = \delta(x^{nx}, f(a^{na})) , \quad (6.3)$$

and

$$P(m_i|x^{nx}) = \delta(m_i, x_i) . \quad (6.4)$$

2. Replace all connectors of the dcircuit by arrows pointing in the direction of the bit flow.

---

Options A and B don't work for digital circuits with feedback loops such as flip-flops.

## Chapter 7

### Do-Calculus: COMING SOON

## Chapter 8

**D-Separation: COMING SOON**



# Chapter 9

## Expectation Maximization



Figure 9.1: bnet for EM with  $nsam = 3$ .

This chapter is based on Wikipedia Ref.[6].

The bnet for Expectation Maximization (EM) is given by Fig.9.1 for  $nsam = 3$ . Later on in this chapter, we will give the node transition prob matrices for this bnet for the special case in which  $P(x[i] | \theta)$  is a mixture (i.e., weighted sum) of Gaussians.

Note that if we erase the  $h[i]$  nodes from Fig.9.1, we get the bnet for naive Bayes, which is used for classification into the states of  $\theta$ . However, there is one big difference. With naive Bayes, the leaf nodes have different transition prob matrices. Here, we will assume they are i.i.d. Naive Bayes is used for classification: i.e., given the states of the leaf nodes, we infer the state of the root node. EM is used for clustering; i.e., given many i.i.d. samples, we fit their distribution by a weighted sum of prob distributions, usually Gaussians.

Let

$\mathcal{L}$  =likelihood function.

$nsam$  = number of samples.

$\vec{x} = (x[0], x[1], \dots, x[nsam - 1])$  = **observed data**.  $x[i] \in S_{\underline{x}}$  for all  $i$ .

$\vec{h} = (h[0], h[1], \dots, h[nsam - 1])$  = **hidden or missing data**.  $h[i] \in S_{\underline{h}}$  for all  $i$ .

We assume that the samples  $(x[i], h[i])$  are i.i.d. for different  $i$  at fixed  $\theta$ . What this means is that there are probability distributions  $P_{\underline{x}|\underline{h},\theta}$  and  $P_{\underline{h}|\theta}$  such that

$$P(\vec{x}, \vec{h}|\theta) = \prod_i [P_{\underline{x}|\underline{h},\theta}(x[i] | h[i], \theta) P_{\underline{h}|\theta}(h[i] | \theta)] . \quad (9.1)$$

Definition of likelihood functions:

$$\underbrace{P(\vec{x}|\theta)}_{\mathcal{L}(\theta;\vec{x})} = \sum_{\vec{h}} \underbrace{P(\vec{x}, \vec{h}|\theta)}_{\mathcal{L}(\theta;\vec{x},\vec{h})} \quad (9.2)$$

$\theta^*$  = maximum likelihood estimate of  $\theta$  (no prior  $P(\theta)$  assumed):

$$\theta^* = \operatorname{argmax}_{\theta} \mathcal{L}(\theta; \vec{x}) \quad (9.3)$$

---

**The EM algorithm:**

1. **Expectation step:**

$$Q(\theta|\theta^{(t)}) = E_{\vec{h}|\vec{x},\theta^{(t)}} \ln P(\vec{x}, \vec{h}|\theta) \quad (9.4)$$

2. **Maximization step:**

$$\theta^{(t+1)} = \operatorname{argmax}_{\theta} Q(\theta|\theta^{(t)}) \quad (9.5)$$

Claim:  $\lim_{t \rightarrow \infty} \theta^{(t)} = \theta^*$ .

---

## Motivation

$$Q(\theta|\theta) = E_{\vec{h}|\vec{x},\theta} \ln P(\vec{x}, \vec{h}|\theta) \quad (9.6)$$

$$= E_{\vec{h}|\vec{x},\theta} [\ln P(\vec{h}|\vec{x}, \theta) + \ln P(\vec{x}|\theta)] \quad (9.7)$$

$$= -H[P(\vec{h}|\vec{x}, \theta)] + \ln P(\vec{x}|\theta) \quad (9.8)$$

$$\partial_{\theta} Q(\theta|\theta) = - \sum_{\vec{h}} \partial_{\theta} P(\vec{h}|\vec{x}, \theta) + \partial_{\theta} \ln P(\vec{x}|\theta) \quad (9.9)$$

$$= \partial_{\theta} \ln P(\vec{x}|\theta) \quad (9.10)$$

So if  $\theta^{(t)} \rightarrow \theta$  and  $Q(\theta|\theta)$  is max at  $\theta = \theta^*$ , then  $\ln P(\vec{x}|\theta)$  is max at  $\theta = \theta^*$  too.

For a more rigorous proof that  $\lim_{t \rightarrow \infty} \theta^{(t)} = \theta^*$ , see Wikipedia article Ref.[6] and references therein.

## EM for Gaussian mixture

$x[i] \in \mathbb{R}^d = S_{\underline{x}}$ .  $S_{\underline{h}}$  discrete and not too large.  $n_{\underline{h}} = |S_{\underline{h}}|$  is number of Gaussians that we are going to fit the samples with.

Let

$$\theta = [w_h, \mu_h, \Sigma_h]_{h \in S_{\underline{h}}}, \quad (9.11)$$

where  $[w_h]_{h \in S_{\underline{h}}}$  is a probability distribution of weights, and where  $\mu_h \in \mathbb{R}^d$  and  $\Sigma_h \in \mathbb{R}^{d \times d}$  are the mean value vector and covariance matrix of a  $d$ -dimensional Gaussian distribution.

The transition prob matrices, printed in blue, for the nodes of Fig.9.1, for the special case of a mixture of Gaussians, are as follows:

$$P(x[i] \mid h[i] \mid \theta) = \mathcal{N}_d(x[i]; \mu_{h[i]}, \Sigma_{h[i]}) \quad (9.12)$$

$$P(h[i] \mid \theta) = w_{h[i]} \quad (9.13)$$

Note that

$$P(x[i] \mid \theta) = \sum_h P(x[i] \mid h[i] = h, \theta) P(h[i] = h \mid \theta) \quad (9.14)$$

$$= \sum_h w_h \mathcal{N}_d(x[i]; \mu_h, \Sigma_h) \quad (9.15)$$

$$P(\vec{x}, \vec{h} \mid \theta) = \prod_i [w_{h[i]} \mathcal{N}_d(x[i]; \mu_{h[i]}, \Sigma_{h[i]})] \quad (9.16)$$

$$= \prod_i \prod_h [w_h \mathcal{N}_d(x[i]; \mu_h, \Sigma_h)]^{\mathbb{1}(h=h[i])} \quad (9.17)$$

**Old Faithful:** See Wikipedia Ref.[6] for an animated gif of a classic example of using EM to fit samples with a Gaussian mixture. Unfortunately, could not include it here because pdfLatex does not support animated gifs. The gif shows samples in a 2 dimensional space (eruption time, delay time) from the Old Faithful geyser. In that example,  $d = 2$  and  $n_{\underline{h}} = 2$ . Two clusters of points in a plane are fitted by a mixture of 2 Gaussians.

**K-means clustering** is often presented as the main competitor to EM for doing **clustering (non-supervised learning)**. In K-means clustering, the sample points are split into  $K$  mutually disjoint sets  $S_0, S_1, \dots, S_{K-1}$ . The algorithm is easy to describe:

1. Initialize by choosing at random  $K$  data points  $(\mu_k)_{k=0}^{K-1}$  called means or centroids and placing  $\mu_k$  in  $S_k$  for all  $k$ .
2. **STEP 1:** For each data point, add it to the  $S_k$  whose centroid  $\mu_k$  is closest to it.

3. **STEP 2:** Recalculate the centroids. Set  $\mu_k$  equal to the mean value of set  $S_k$ .
4. Repeat steps 1 and 2 until the centroids stop changing by much.

Step 1 is analogous to the expectation step in EM, and Step 2 to the maximization step in EM ( $\theta$  estimation versus  $\mu_k$  estimation). We won't say anything further about K-means clustering because it isn't related to bnets in any way, and this is a book about bnets. For more info about K-means clustering, see Ref.[7].

# Chapter 10

## Generative Adversarial Networks (GANs)

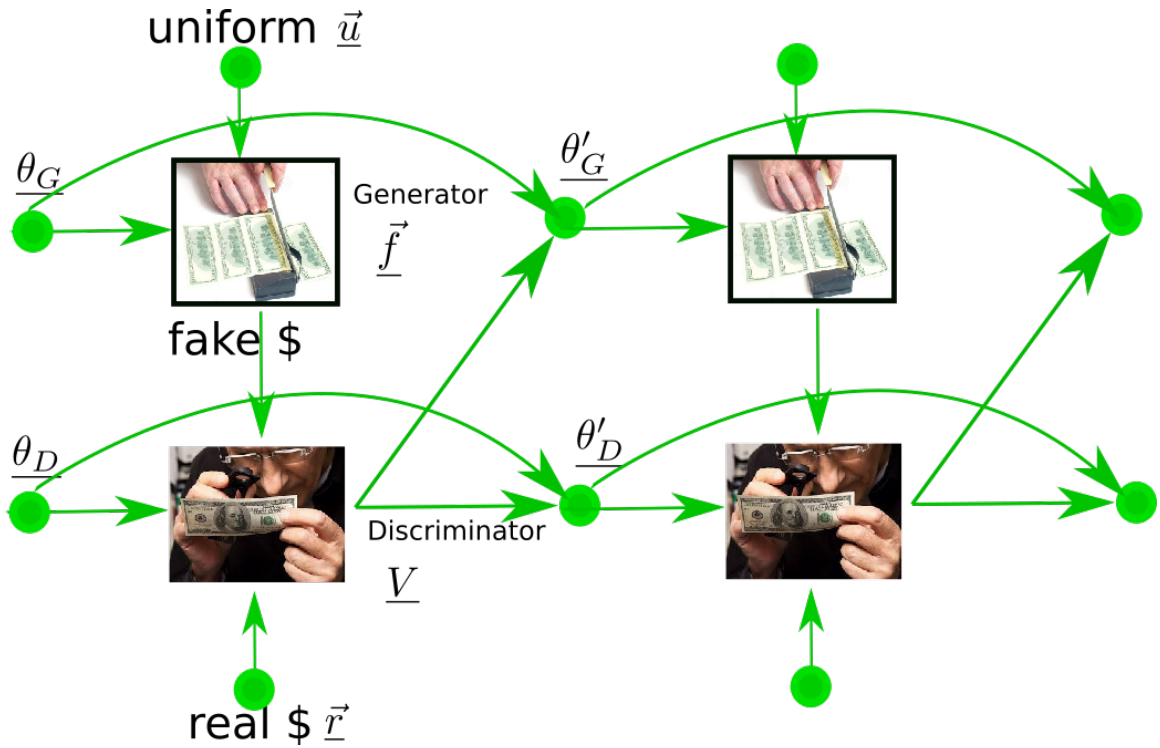


Figure 10.1: Generative Adversarial Network (GAN)

Original GAN, Ref.[8](2014).

Generator  $G$  (counterfeiter) generates samples  $\vec{f}$  of fake money and submits them to Discriminator  $D$  (Treasury agent).  $D$  also gets samples  $\vec{r}$  of real money.  $D$  submits verdict  $V \in [0, 1]$ .  $G$  depends on parameter  $\theta_G$  and  $D$  on parameter  $\theta_D$ . Verdict  $V$  and initial  $\theta_G, \theta_D$  are used to get new parameters  $\theta'_G, \theta'_D$ . Process is repeated (Dynamical Bayesian Network) until saddle point in



Figure 10.2: Discriminator node  $\underline{V}$  in Fig.10.1 can be split into 3 nodes  $\vec{c}$ ,  $\vec{d}$  and  $\underline{V}$ .

$V(\theta_G, \theta_D)$  is reached.  $D$  makes  $G$  better and vice versa. Zero-sum game between  $D$  and  $G$ .

Let  $\mathcal{D}$  be the domain of  $D(\cdot, \theta_D)$ . Assume that for any  $x \in \mathcal{D}$ ,

$$0 \leq D(x, \theta_D) \leq 1 . \quad (10.1)$$

For any  $S \subset \mathcal{D}$ , define

$$\sum_{x \in S} D(x, \theta_D) = \lambda(S, \theta_D) . \quad (10.2)$$

In general,  $G(\cdot, \theta_G)$  need not be real valued.

Assume that for every  $u \in S_u$ ,  $G(u, \theta_G) = f \in S_f \subset \mathcal{D}$ . Define

$$\overline{D}(f, \theta_D) = 1 - D(f, \theta_D) . \quad (10.3)$$

Note that

$$0 \leq \overline{D}(f, \theta_D) \leq 1 . \quad (10.4)$$

Define:

$$V(\theta_G, \theta_D) = \sum_r P(r) \ln D(r, \theta_D) + \sum_u P(u) \ln \overline{D}(G(u, \theta_G), \theta_D) . \quad (10.5)$$

We want the first variation of  $V(\theta_G, \theta_D)$  to vanish.

$$\delta V(\theta_G, \theta_D) = 0 . \quad (10.6)$$

This implies

$$\partial_{\theta_G} V(\theta_G, \theta_D) = \partial_{\theta_D} V(\theta_G, \theta_D) = 0 \quad (10.7)$$

and

$$V_{opt} = \min_{\theta_G} \max_{\theta_D} V(\theta_G, \theta_D) . \quad (10.8)$$

Node transition probability matrices for Figs.10.1 and 10.2 are given next in blue:

$$P(\theta_G) = \text{given} \quad (10.9)$$

$$P(\theta_D) = \text{given} \quad (10.10)$$

$$P(\vec{u}) = \prod_i P(u[i]) \quad (\text{usually uniform distribution}) \quad (10.11)$$

$$P(\vec{r}) = \prod_i P(r[i]) \quad (10.12)$$

$$P(f[i] \mid \vec{u}, \theta_G) = \delta[f[i], G(u[i], \theta_G)] \quad (10.13)$$

$$P(c[i] \mid \vec{f}, \theta_D) = \delta(c[i], \overline{D}(f[i], \theta_D)) \quad (10.14)$$

$$P(d[j] \mid \vec{r}, \theta_D) = \delta(d[j], D(r[j], \theta_D)) \quad (10.15)$$

$$P(V \mid \vec{d}, \vec{c}) = \delta(V, \frac{1}{N} \ln \prod_{i,j} (c[i] d[j])) \quad (10.16)$$

where  $N = nsam(\vec{r}) nsam(\vec{u})$ .

Let  $\eta_G, \eta_D > 0$ . Maximize  $V$  wrt  $\theta_D$ , and minimize it wrt  $\theta_G$ .

$$P(\theta'_G \mid V, \theta_G) = \delta(\theta'_G, \theta_G - \eta_G \partial_{\theta_G} V) \quad (10.17)$$

$$P(\theta'_D \mid V, \theta_D) = \delta(\theta'_D, \theta_D + \eta_D \partial_{\theta_D} V) \quad (10.18)$$

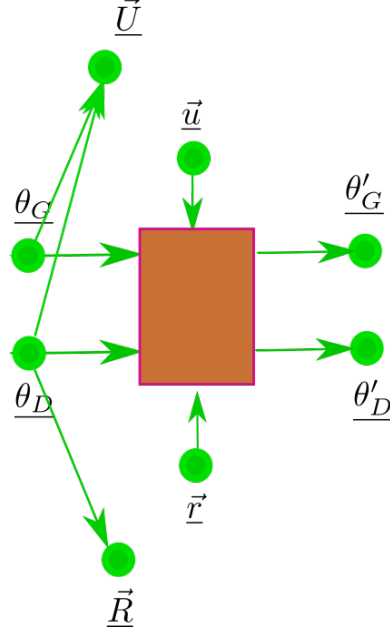


Figure 10.3: GAN, Constraining Bayesian Network

Constraining B net given in Fig.10.3. It adds 2 new nodes, namely  $\vec{U}$  and  $\vec{R}$ , to the bnet of Fig.10.1. The purpose of these 2 barren (childrenless) nodes is to constrain certain functions to be probability distributions.

Node transition probabilities for the 2 new nodes given next in blue.

$$P(U[i] | \theta_G) = \frac{\overline{D}(G(U[i], \theta_G), \theta_D))}{\overline{\lambda}(\theta_G, \theta_D)} \quad (10.19)$$

where  $S_{\underline{U}[i]} = S_{\underline{u}}$  and  $\overline{\lambda}(\theta_G, \theta_D) = \sum_u \overline{D}(G(u, \theta_G), \theta_D)$ .

$$P(R[i] | \theta_G, \theta_D) = \frac{D(R[i], \theta_D)}{\lambda(\theta_D)} \quad (10.20)$$

where  $S_{\underline{R}[i]} = S_{\underline{r}}$  and  $\lambda(\theta_D) = \sum_r D(r, \theta_D)$ .

$$P(V | \vec{u}, \vec{r}) = \delta(V, \frac{1}{N} \ln \prod_{i,j} (P(R[i] = r[i] | \theta_G, \theta_D) P(U[i] = u[j] | \theta_G))) \quad (10.21)$$

where  $N = nsam(\vec{r})nsam(\vec{u})$ .

$\mathcal{L}$  = likelihood

$$\mathcal{L} = P(\vec{r}, \vec{u} | \theta_G, \theta_D) \quad (10.22)$$

$$= \prod_{i,j} \left[ \frac{D(r[i], \theta_D)}{\lambda(\theta_D)} \frac{\overline{D}(G(u[j], \theta_G), \theta_D))}{\overline{\lambda}(\theta_G, \theta_D)} \right] \quad (10.23)$$



$$\ln \mathcal{L} = N[V(\theta_G, \theta_D) - \ln \lambda(\theta_D) - \ln \bar{\lambda}(\theta_G, \theta_D)] \quad (10.24)$$

## Chapter 11

**Graph Structure Learning for bnets:  
COMING SOON**

# Chapter 12

## Hidden Markov Model

A Hidden Markov Model (HMM) is a generalization of a Kalman Filter (KF). KFs are discussed in Chapter 14. The bnets of HHMs and KFs bnets are the same. The only difference is that a KF assumes special node transition matrices.

See Wikipedia article Ref.[9] to learn about the history and many uses of HMMs. This chapter is based on Ref.[10].



Figure 12.1: HMM bnet with  $n = 4$ .

Suppose

$\underline{v}^n = (\underline{v}_0, \underline{v}_1, \dots, \underline{v}_{n-1})$  are  $n$  visible nodes that are measured, and

$\underline{x}^n = (\underline{x}_0, \underline{x}_1, \dots, \underline{x}_{n-1})$  are the  $n$  hidden, unmeasurable state nodes of a system that is being monitored.

For the bnet of Fig.12.1, one has

$$P(\underline{x}^n, \underline{v}^n) = \prod_{i=0}^{n-1} P(x_i | x_{i-1}) P(v_i | x_i) , \quad (12.1)$$

where  $x_{-1} = 0$ .

Let  $x_{<i} = (x_0, x_1, \dots, x_{i-1})$ .

For  $i = 0, 1, \dots, n - 1$ , define

$\mathcal{F}_i$ =future measurements probability

$$\mathcal{F}_i(x_i) = P(v_{>i} | x_i) \quad (12.2)$$

$\overline{\mathcal{F}}_i$  = past and present measurements probability

$$\overline{\mathcal{F}}_i(x_i) = P(v_{<i}, v_i, x_i) \quad (12.3)$$

$\lambda_i$  = present measurement probability

$$\lambda_i(x_i) = P(v_i|x_i) \quad (12.4)$$

$\mathcal{F}_i$ ,  $\overline{\mathcal{F}}_i$  and  $\lambda_i$  can be represented graphically as follows:

$$\mathcal{F}_i(x_i) = \frac{1}{P(x_i)} \sum_{x_{>i}} \quad \begin{array}{c} x_i \longrightarrow x_{>i} \\ \downarrow \\ v_{>i} \end{array} \quad (12.5)$$

$$\overline{\mathcal{F}}_i(x_i) = \sum_{x_{<i}} \quad \begin{array}{c} x_{<i} \longrightarrow x_i \\ \downarrow \quad \downarrow \\ v_{<i} \quad v_i \end{array} \quad (12.6)$$

$$\lambda_i(x_i) = \frac{1}{P(x_i)} \quad \begin{array}{c} x_i \\ \downarrow \\ v_i \end{array} \quad (12.7)$$

**Claim 1** For  $i \geq 0$ ,

$$P(x_i, v^n) = \overline{\mathcal{F}}_i(x_i) \mathcal{F}_i(x_i) . \quad (12.8)$$

For  $i > 0$ ,

$$P(x_{i-1}, x_i, v^n) = \overline{\mathcal{F}}_{i-1}(x_{i-1}) \lambda_i(x_i) P(x_i|x_{i-1}) \mathcal{F}_i(x_i) . \quad (12.9)$$

**proof:**

$$P(x_i, v^n) = \sum_{x_{<i}} \sum_{x_{>i}} P(x^n, v^n) \quad (12.10)$$

$$= \sum_{x_{<i}} \sum_{x_{>i}} P(x^n, v^n | x_i) P(x_i) \quad (12.11)$$

$$= \sum_{x_{<i}} \sum_{x_{>i}} P(x_{<i}, v_{<i}, v_i | x_i) P(x_{>i}, v_{>i} | x_i) P(x_i) \quad (12.12)$$

$$= P(v_{<i}, v_i | x_i) P(v_{>i} | x_i) P(x_i) \quad (12.13)$$

$$= \overline{\mathcal{F}}_i(x_i) \mathcal{F}_i(x_i) \quad (12.14)$$

$$P(x_{i-1}, x_i, v^n) = \sum_{x_{<i-1}} \sum_{x_{>i}} P(x^n, v^n) \quad (12.15)$$

$$= \sum_{x_{<i-1}} \sum_{x_{>i}} P(x^n, v^n | x_{i-1}, x_i) P(x_{i-1}, x_i) \quad (12.16)$$

$$= \sum_{x_{<i-1}} \sum_{x_{>i}} P(x_{<i-1}, v_{<i-1}, v_{i-1} | x_{i-1}) P(v_i | x_i) P(x_{i-1}, x_i) P(x_{>i}, v_{>i} | x_i) \quad (12.17)$$

$$= P(v_{<i-1}, v_{i-1} | x_{i-1}) P(v_i | x_i) P(x_{i-1}, x_i) P(v_{>i} | x_i) \quad (12.18)$$

$$= \overline{\mathcal{F}}_{i-1}(x_{i-1}) \lambda_i(x_i) P(x_i | x_{i-1}) \mathcal{F}_i(x_i) \quad (12.19)$$

**QED**

**Claim 2** For  $i > 0$ ,  $\mathcal{F}_i$  and  $\overline{\mathcal{F}}_i$  can be calculated recursively as follows:

$$\overline{\mathcal{F}}_i(x_i) = \sum_{x_{i-1}} \overline{\mathcal{F}}_{i-1}(x_{i-1}) \lambda_i(x_i) P(x_i | x_{i-1}) \quad (12.20)$$

$$\mathcal{F}_{i-1}(x_{i-1}) = \sum_{x_i} \lambda_i(x_i) P(x_i | x_{i-1}) \mathcal{F}_i(x_i) \quad (12.21)$$

**proof:**

$$\overline{\mathcal{F}}_i(x_i) \mathcal{F}_i(x_i) = P(x_i, v^n) \quad (12.22)$$

$$= \sum_{x_{i-1}} P(x_{i-1}, x_i, v^n) \quad (12.23)$$

$$= \sum_{x_{i-1}} \overline{\mathcal{F}}_{i-1}(x_{i-1}) \lambda_i(x_i) P(x_i | x_{i-1}) \mathcal{F}_i(x_i) \quad (12.24)$$

$$\overline{\mathcal{F}}_{i-1}(x_{i-1}) \mathcal{F}_{i-1}(x_{i-1}) = P(x_{i-1}, v^n) \quad (12.25)$$

$$= \sum_{x_i} P(x_{i-1}, x_i, v^n) \quad (12.26)$$

$$= \sum_{x_i} \overline{\mathcal{F}}_{i-1}(x_{i-1}) \lambda_i(x_i) P(x_i | x_{i-1}) \mathcal{F}_i(x_i) \quad (12.27)$$

**QED**

**Claim 3**

$$P(x_i | x_{i-1}, v^n) = \frac{\lambda_i(x_i) \mathcal{F}_i(x_i)}{\overline{\mathcal{F}}_{i-1}(x_{i-1})} P(x_i | x_{i-1}) \quad (12.28)$$

$$P(x_{i-1} | x_i, v^n) = \frac{\lambda_i(x_i) \overline{\mathcal{F}}_{i-1}(x_{i-1})}{\overline{\mathcal{F}}_i(x_i)} P(x_i | x_{i-1}) \quad (12.29)$$

**proof:**

$$P(x_i|x_{i-1}, v^n) = \frac{P(x_{i-1}, x_i, v^n)}{P(x_{i-1}, v^n)} \quad (12.30)$$

$$= \frac{\overline{\mathcal{F}}_{i-1}(x_{i-1})\lambda_i(x_i)P(x_i|x_{i-1})\mathcal{F}_i(x_i)}{\overline{\mathcal{F}}_{i-1}(x_{i-1})\mathcal{F}_{i-1}(x_{i-1})} \quad (12.31)$$

Analogous proof for Eq.(12.29).

**QED**

# Chapter 13

## Influence Diagrams & Utility Nodes

Influence diagrams are just arbitrary bnets enhanced with a new kind of node called an utility node. The rest of this brief chapter will be devoted to discussing utility nodes.

An utility node can be understood as a node composed of 3 simpler bnet nodes. This is illustrated in Fig.13.1.



Figure 13.1: An utility node can be understood as a node composed of 3 simpler bnet nodes.

The transition probability matrices, printed in blue, for the nodes of Fig.13.1, are as follows:

$$P(U|pa(U)) = \text{given} \quad (13.1)$$

$$P(u|pa(U)) = P(U = u|pa(U)) \quad (13.2)$$

Node  $\underline{\mu}_u$  calculates the expected value (mean value) of  $\underline{u}$ :

$$P(\underline{\mu}_u) = \delta(\underline{\mu}_u, E_{\underline{u}}[\underline{u}]) \quad (13.3)$$

Node  $\underline{\sigma}_u$  calculates the standard deviation of  $\underline{u}$ :

$$P(\underline{\sigma}_u) = \delta(\underline{\sigma}_u, \sqrt{E_{\underline{u}}[(\underline{u} - E_{\underline{u}}[\underline{u}])^2]}) \quad (13.4)$$

Note that in order to calculate expected values, it is necessary that  $\underline{U}, \underline{u} \in \mathbb{R}$ . Note that nodes  $\underline{u}, \underline{\mu}_u, \underline{\sigma}_u$  must all 3 have access to the transition matrix  $P(U|pa(U))$  of node  $\underline{U}$ . In fact, in order to calculate  $E_{\underline{u}}[\cdot]$ , it is necessary for nodes  $\underline{\mu}_u$  and  $\underline{\sigma}_u$  to have access not just to  $P(U|pa(U))$  but also to  $P(pa(U))$ .

See Fig.13.2. An influence diagram may have multiple utility nodes ( $\underline{U}_1$  and  $\underline{U}_2$  in Fig.13.2). Then one can define a merging utility node  $\underline{U}$  that sums the values of all the other utility nodes.

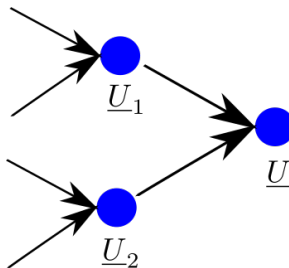


Figure 13.2: An influence diagram may have multiple utility nodes, say  $\underline{U}_1$  and  $\underline{U}_2$ . Then one can define an utility node  $\underline{U} = \underline{U}_1 + \underline{U}_2$ .

For the node  $\underline{U}$  of Fig.13.2,

$$P(U|U_1, U_2) = \delta(U, U_1 + U_2) \quad (13.5)$$



# Chapter 14

## Kalman Filter

A Kalman Filter is a special case of a Hidden Markov Model. HMMs are discussed in Chapter 12.

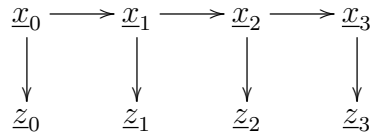


Figure 14.1: Kalman Filter bnnet with  $T = 4$ .

Let  $t = 0, 1, 2, \dots, T - 1$ .

$\underline{x}_t \in S_{\underline{x}}$  are random variables that represent the hidden (unobserved) true state of the system.

$\underline{z}_t \in S_{\underline{z}}$  are random variables that represent the measured (observed) state of the system.

The Kalman Filter bnnet Fig.14.1 has the following node probability transition matrices, printed in blue:

$$P(x_t|x_{t-1}) = \mathcal{N}(F_t x_{t-1} + B_t u_t, Q_t) , \quad (14.1)$$

where  $F_t, Q_t, B_t, u_t$  are given.  $P(x_t|x_{t-1})$  becomes  $P(x_t)$  for  $t = 0$ .

$$P(z_t|x_t) = \mathcal{N}(H_t x_t, R_t) , \quad (14.2)$$

where  $H_t, R_t$  are given.

Define

$$\underline{Z}_t = (\underline{z}_{t'})_{t' \leq t} . \quad (14.3)$$

Define  $\hat{x}_t$  and  $P_t$  by

$$P(x_t|Z_t) = \mathcal{N}(\hat{x}_t, P_t) . \quad (14.4)$$

---

Problem: Find  $\hat{x}_t$  and  $P_t$  in terms of

1. current (at time  $t$ ) given values of  $F, Q, H, R, B, u$
2. current (at time  $t$ ) observed value of  $z$
3. prior (previous) value (at time  $t - 1$ ) of  $\hat{x}$  and  $P$ .

See Fig.14.2. For that figure,

$$P(\hat{x}_t, P_t | z_t, \hat{x}_{t-1}, P_{t-1}) = \delta(\hat{x}_t, ?) \delta(P_t, ?) . \quad (14.5)$$

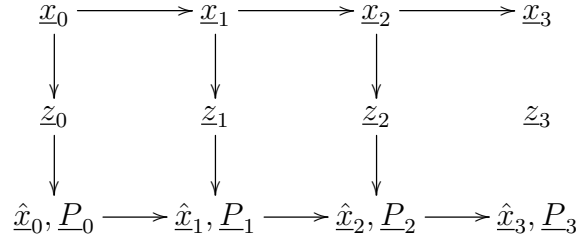


Figure 14.2: Kalman Filter bnnet with deterministic nodes for  $\hat{x}_t, P_t$ .

---

Solution copied from Wikipedia Ref.[11]:

Define  $\eta_{t|t} = \eta_t$  for  $\eta = \hat{x}, P$ .

- **Predict**

Predicted (a priori) state estimate

$$\hat{x}_{t|t-1} = F_t \hat{x}_{t-1|t-1} + B_t u_t \quad (14.6)$$

Predicted (a priori) estimate covariance

$$P_{t|t-1} = F_t P_{t-1|t-1} F_t^\top + Q_t \quad (14.7)$$

- **Update**

Innovation (or measurement pre-fit residual)

$$\tilde{y}_{t|t-1} = z_t - H_t \hat{x}_{t|t-1} \quad (14.8)$$

Innovation (or pre-fit residual) covariance

$$S_t = H_t P_{t|t-1} H_t^\top + R_t \quad (14.9)$$

Optimal Kalman gain

$$K_t = P_{t|t-1} H_t^\top S_t^{-1} \quad (14.10)$$

Updated (a posteriori) state estimate

$$\hat{x}_{t|t} = \hat{x}_{t|t-1} + K_t \tilde{y}_t \quad (14.11)$$

Updated (a posteriori) estimate covariance

$$P_{t|t} = (I - K_t H_t) P_{t|t-1} \quad (14.12)$$

Measurement post-fit residual

$$\tilde{y}_{t|t} = z_t - H_t \hat{x}_{t|t} \quad (14.13)$$

# Chapter 15

## Linear and Logistic Regression

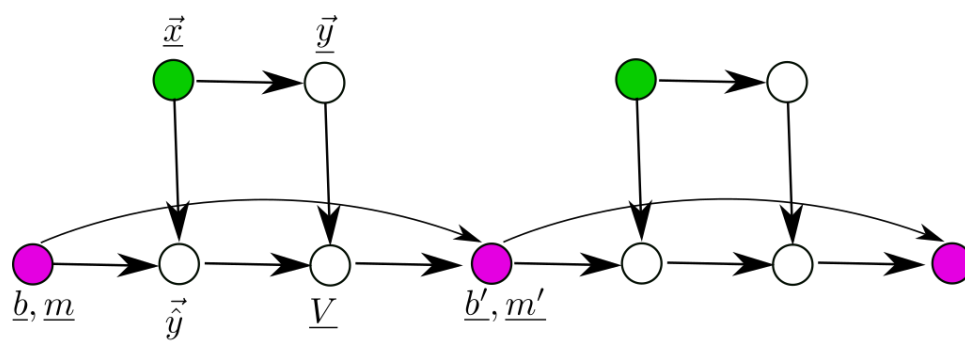


Figure 15.1: Linear Regression

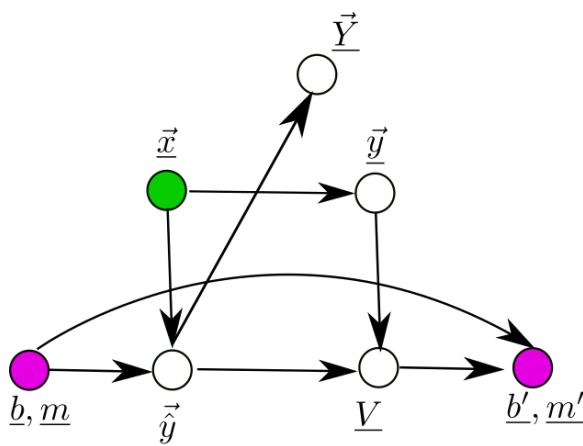


Figure 15.2: B net of Fig.15.1 with new  $\vec{Y}$  node.

Estimators  $\hat{y}$  for linear and logistic regression.

- **Linear Regression:**  $y \in \mathbb{R}$ . Note  $\hat{y} \in \mathbb{R}$ .  $(x, \hat{y}(x))$  is the graph of a straight line with y-intercept  $b$  and slope  $m$ .

$$\hat{y}(x; b, m) = b + mx \quad (15.1)$$

- **Logistic Regression:**  $y \in \{0, 1\}$ . Note  $\hat{y} \in [0, 1]$ .  $(x, \hat{y}(x))$  is the graph of a sigmoid. Often in literature,  $b, m$  are replaced by  $\beta_0, \beta_1$ .

$$\hat{y}(x; b, m) = \text{sig}(b + mx) \quad (15.2)$$

Define

$$V(b, m) = \sum_{x, y} P(x, y) |y - \hat{y}(x; b, m)|^2 . \quad (15.3)$$

We want to minimize  $V(b, m)$  (called a cost or loss function) wrt  $b$  and  $m$ .

Node transition probabilities of B net of Fig.15.1 given next in blue.

$$P(b, m) = \text{given} \quad (15.4)$$

The first time it is used,  $(b, m)$  is arbitrary. After the first time, it is determined by previous stage.

Let

$$P_{\underline{x}, \underline{y}}(x, y) = \frac{1}{nsam(\vec{x})} \sum_i \mathbb{1}(x = x[i], y = y[i]) . \quad (15.5)$$

$$P(\vec{x}) = \prod_i P(x[i]) \quad (15.6)$$

$$P(\vec{y}|\vec{x}) = \prod_i P(y[i] | x[i]) \quad (15.7)$$

$$P(\vec{\hat{y}}|\vec{x}, b, m) = \prod_i \delta(\hat{y}[i], \hat{y}(x[i], b, m)) \quad (15.8)$$

$$P(V|\vec{\hat{y}}, \vec{y}) = \delta(V, \frac{1}{nsam(\vec{x})} \sum_i |y[i] - \hat{y}[i]|^2) \quad (15.9)$$

Let  $\eta_b, \eta_m > 0$ . For  $x = b, m$ , if  $x' - x = \Delta x = -\eta \frac{\partial V}{\partial x}$ , then  $\Delta V \approx \frac{-1}{\eta} (\Delta x)^2 \leq 0$  for  $\eta > 0$ . This is called “gradient descent”.

$$P(b'|V, b) = \delta(b', b - \eta_b \partial_b V) \quad (15.10)$$

$$P(m'|V, m) = \delta(m', m - \eta_m \partial_m V) \quad (15.11)$$

## Generalization to $x$ with multiple components(features)

Suppose that for each sample  $i$ , instead of  $x[i]$  being a scalar, it has  $n$  components called features:

$$x[i] = (x_0[i], x_1[i], x_2[i], \dots, x_{n-1}[i]) . \quad (15.12)$$

Slope  $m$  is replaced by weights

$$w = (w_0, w_1, w_2, \dots, w_{n-1}) , \quad (15.13)$$

and the product of 2 scalars  $mx[i]$  is replaced by the inner vector product  $w^T x[i]$ .

## Alternative $V(b, m)$ for logistic regression

For logistic regression, since  $y[i] \in \{0, 1\}$  and  $\hat{y}[i] \in [0, 1]$  are both in the interval  $[0, 1]$ , they can be interpreted as probabilities. Define probability distributions  $p[i](x)$  and  $\hat{p}[i](x)$  for  $x \in \{0, 1\}$  by

$$p[i](1) = y[i], \quad p[i](0) = 1 - y[i] \quad (15.14)$$

$$\hat{p}[i](1) = \hat{y}[i], \quad \hat{p}[i](0) = 1 - \hat{y}[i] \quad (15.15)$$

Then for logistic regression, the following 2 cost functions  $V(b, m)$  can be used as alternatives to the cost function Eq.(15.3) previously given.

$$V(b, m) = \frac{1}{nsam(\vec{x})} \sum_i D_{KL}(p[i] \parallel \hat{p}[i]) \quad (15.16)$$

and

$$V(b, m) = \frac{1}{nsam(\vec{x})} \sum_i CE(p[i] \rightarrow \hat{p}[i]) \quad (15.17)$$

$$= \frac{-1}{nsam(\vec{x})} \sum_i \{y[i] \ln \hat{y}[i] + (1 - y[i]) \ln(1 - \hat{y}[i])\} \quad (15.18)$$

$$= \frac{-1}{nsam(\vec{x})} \sum_i \ln \{ \hat{y}[i]^{y[i]} (1 - \hat{y}[i])^{(1-y[i])} \} \quad (15.19)$$

$$= \frac{-1}{nsam(\vec{x})} \sum_i \ln P(\underline{Y} = y[i] \mid \hat{y} = \hat{y}[i]) \quad (15.20)$$

$$= - \sum_{x, y} P(x, y) \ln P(\underline{Y} = y \mid \hat{y} = \hat{y}(x, b, m)) \quad (15.21)$$

Above, we used

$$P(\underline{Y} = Y \mid \hat{y}) = \hat{y}^Y [1 - \hat{y}]^{1-Y} \quad (15.22)$$

for  $Y \in S_{\underline{Y}} = \{0, 1\}$ . (Bernoulli distribution).

There is no node corresponding to  $\underline{Y}$  in the B net of Fig.15.1. Fig.15.2 shows a new B net that has a new node called  $\vec{Y}$  compared to the B net of Fig.15.1. One defines the transition probabilities for all nodes of Fig.15.2 except  $\vec{Y}$  and  $\underline{V}$  the same as for Fig.15.1. For  $\vec{Y}$  and  $\underline{V}$ , one defines

$$P(Y[i] | \vec{y}) = P(\underline{Y} = Y[i] | \hat{y}[i]) \quad (15.23)$$

$$P(V | \vec{Y}, \vec{y}) = \delta(V, \frac{-1}{nsam(\vec{x})} \ln \mathcal{L}) , \quad (15.24)$$

where  $\mathcal{L} = \prod_i P(\underline{Y} = y[i] | \hat{y}[i])$ =likelihood.

# Chapter 16

## Markov Blankets

This chapter is based on the Wikipedia article, Ref.[12]. Markov blankets and Markov boundaries of bnets were apparently invented by Judea Pearl. His 1988 book Ref.[13], instead of a research paper, is usually given as the original reference.



Figure 16.1: In a bnet, the minimal Markov blanket, aka Markov boundary, of node  $\underline{a}$ .

We will treat vectors of random variables as if they were sets when using the  $\in$ ,  $\subset$  and  $-$  operations. For example, if  $\underline{x} = (\underline{x}_0, \underline{x}_1, \underline{x}_2, \underline{x}_3)$  and  $\underline{b} = (\underline{x}_1, \underline{x}_2)$ , then  $\underline{x}_1 \in \underline{b} \subset \underline{x}$  and  $\underline{x} - \underline{b} = (\underline{x}_0, \underline{x}_3)$ .

Below,  $H(\underline{a} : \underline{b} | \underline{c})$  denotes the conditional mutual information of random variables  $\underline{a}$  and  $\underline{b}$  conditioned on random variable  $\underline{c}$ .  $H(\underline{a} : \underline{b} | \underline{c})$  is used in Shannon Information Theory, where it is



defined by

$$H(\underline{a} : \underline{b} | \underline{c}) = \sum_{a,b,c} P(a, b, c) \ln \frac{P(a, b | c)}{P(a | c)P(b | c)} . \quad (16.1)$$

$H(\underline{a} : \underline{b} | \underline{c}) = 0$  iff  $\underline{a}$  and  $\underline{b}$  are independent (uncorrelated) when  $\underline{c}$  is held fixed.

Suppose  $\underline{a} \in \underline{X}$ ,  $\underline{B} \subset \underline{X}$ , but  $\underline{a} \notin \underline{B}$ . Then  $\underline{B}$  is a Markov blanket of  $\underline{a}$  if

$$H(\underline{a} : \underline{X} - \underline{a} | \underline{B}) = 0 . \quad (16.2)$$

In other words, one may assume that  $\underline{a}$  depends on  $\underline{B}$  only, and is independent of all random variables in  $\underline{X} - (\underline{a} \cup \underline{B})$ .

The minimal Markov blanket is called the Markov boundary.

In a bnet, the Markov boundary of a node  $\underline{a}$ , contains:

1. the parents of  $\underline{a}$ ,
2. the children of  $\underline{a}$ ,
3. the parents, other than  $\underline{a}$ , of the children of  $\underline{a}$ .

This is illustrated in Fig.16.1.

# Chapter 17

## Markov Chains

A Markov Chain is simply a bnet with the graph structure of a chain. For example, Fig.17.1 shows a chain with  $n = 4$  nodes.

$$\underline{x}_0 \longrightarrow \underline{x}_1 \longrightarrow \underline{x}_2 \longrightarrow \underline{x}_3$$

Figure 17.1: Markov chain with  $n = 4$  nodes.

Because of its graph structure, the transition prob matrix of each node only depends on the state of the previous node:

$$P(x_t | (x_a)_{a \neq t}) = P(x_t | x_{t-1}) , \quad (17.1)$$

where  $(x_a)_{a \neq t}$  are all the nodes except  $x_t$  itself and  $t = 1, 2, \dots, n - 1$ .

If there exists a single transition prob matrix  $P_{\underline{x}_1 | \underline{x}_0}$  such that

$$P(x_t | x_{t-1}) = P_{\underline{x}_1 | \underline{x}_0}(x_t | x_{t-1}) \quad (17.2)$$

for  $t = 1, 2, \dots, n - 1$ , then we say that the Markov chain is **time homogeneous**.

## Chapter 18

# Markov Chain Monte Carlo (MCMC): COMING SOON

## Chapter 19

# Message Passing (Belief Propagation): COMING SOON

# Chapter 20

## Monty Hall Problem



Figure 20.1: Monty Hall Problem.

Mr. Monty Hall, host of the game show “Lets Make a Deal”, hides a car behind one of three doors and a goat behind each of the other two. The contestant picks Door No. 1, but before opening it, Mr. Hall opens Door No. 2 to reveal a goat. Should the contestant stick with No. 1 or switch to No. 3?

The Monty Hall problem can be modeled by the bnet Fig.20.1, where

- $\underline{c}$ = the door behind which the car actually is.
- $\underline{y}$ = the door opened by you (the contestant), on your first selection.
- $\underline{m}$ = the door opened by Monty (game host)

We label the doors 1,2,3 so  $S_{\underline{c}} = S_{\underline{y}} = S_{\underline{m}} = \{1, 2, 3\}$ .

Node matrices printed in blue:

$$P(c) = \frac{1}{3} \text{ for all } c \quad (20.1)$$

$$P(y) = \frac{1}{3} \text{ for all } y \quad (20.2)$$

$$P(m|c, y) = \mathbb{1}(m \neq c) \left[ \frac{1}{2} \mathbb{1}(y = c) + \mathbb{1}(y \neq c) \mathbb{1}(m \neq y) \right] \quad (20.3)$$

It's easy to show that the above node probabilities imply that

$$P(c = 1|m = 2, y = 1) = \frac{1}{3} \quad (20.4)$$

$$P(c = 3|m = 2, y = 1) = \frac{2}{3} \quad (20.5)$$

So you are twice as likely to win if you switch your final selection to be the door which is neither your first choice nor Monty's choice.

The way I justify this to myself is: Monty gives you a piece of information. If you don't switch your choice, you are wasting that info, whereas if you switch, you are using the info.

# Chapter 21

## Naive Bayes



Figure 21.1: bnet for Naive Bayes with 4 features

Class node  $\underline{c} \in S_{\underline{c}}$ .  $|S_{\underline{c}}| = n_{\underline{c}}$  = number of classes.

Feature nodes  $\underline{x}_i \in S_{\underline{x}_i}$  for  $i = 0, 1, 2, \dots, F - 1$ .  $F$  = number of features.

Define

$$\underline{x}. = [x_0, x_1, \dots, x_{F-1}] . \quad (21.1)$$

For the bnet of Fig.21.1,

$$P(c, \underline{x}.) = P(c) \prod_{i=0}^{F-1} P(x_i | c) . \quad (21.2)$$

Given  $\underline{x}.$  values, find most likely class  $c \in S_{\underline{c}}$ .

Maximum a Posteriori (MAP) estimate:

$$c^* = \operatorname{argmax}_c P(c | \underline{x}.) \quad (21.3)$$

$$= \operatorname{argmax}_c \frac{P(c, \underline{x}.)}{P(\underline{x}.)} \quad (21.4)$$

$$= \operatorname{argmax}_c P(c, \underline{x}.) . \quad (21.5)$$

# Chapter 22

## Neural Networks

In this chapter, we discuss Neural Networks (NNs) of the feedforward kind, which is the most popular kind. In their plain, vanilla form, NNs only have deterministic nodes. But the nodes of a bnet can be deterministic too, because the transition probability matrix of a node can reduce to a delta function. Hence, NNs should be expressible as bnets. We will confirm this in this chapter.

Henceforth in this chapter, if we replace an index of an indexed quantity by a dot, it will mean the collection of the indexed quantity for all values of that index. For example,  $\underline{x}$  will mean the array of  $x_i$  for all  $i$ .



Figure 22.1: Neural Network (feed forward) with 4 layers: input layer  $\underline{x}$ ., 2 hidden layers  $\underline{h}^0$ .,  $\underline{h}^1$ . and output layer  $\underline{Y}$ .

Consider Fig.22.1.

$\underline{x}_i \in \{0, 1\}$  for  $i = 0, 1, 2, \dots, numx - 1$  is the **input layer**.

$\underline{h}_i^\lambda \in \mathbb{R}$  for  $i = 0, 1, 2, \dots, numh(\lambda) - 1$  is the  **$\lambda$ -th hidden layer**.  $\lambda = 0, 1, 2, \dots, \Lambda - 1$ . A NN is said to be **deep** if  $\Lambda > 1$ ; i.e., if it has more than one hidden layer.

$\underline{Y}_i \in \mathbb{R}$  for  $i = 0, 1, 2, \dots, numy - 1$  is the **output layer**. We use a upper case y here because in the training phase, we will use pairs  $(x.[s], y.[s])$  where  $y_i[s] \in \{0, 1\}$  for  $i = 0, 1, \dots, numy - 1$ .  $Y = \hat{y}$  is an estimate of  $y$ . Note that lower case y is either 0 or 1, but upper case y may be any



real. Often, the activation functions are chosen so that  $Y \in [0, 1]$ .

The number of nodes in each layer and the number of layers are arbitrary. Fig.22.1 is fully connected (aka dense), meaning that every node of a layer is impinged arrow coming from every node of the preceding layer. Later on in this chapter, we will discuss non-dense layers.

Let  $w_{i|j}^\lambda, b_i^\lambda \in \mathbb{R}$  be given, for  $i \in \mathbb{Z}_{[0, numh(\lambda)]}$ ,  $j \in \mathbb{Z}_{[0, numh(\lambda-1)]}$ , and  $\lambda \in \mathbb{Z}_{[0, \Lambda]}$ .

These are the transition probability matrices, printed in blue, for the nodes of the bnet Fig.22.1:

$$P(x_i | x_{i-1}, x_{i-1}, \dots, x_0) = \text{given} \quad (22.1)$$

$$P(h_i^\lambda | h_{\cdot}^{\lambda-1}) = \delta \left( h_i^\lambda, \mathcal{A}_i^\lambda \left( \sum_j w_{i|j}^{\lambda-1} h_j^{\lambda-1} + b_i^{\lambda-1} \right) \right) \quad (22.2)$$

$$P(Y_i | h_{\cdot}^{\Lambda-1}) = \delta \left( Y_i, \mathcal{A}_i^\Lambda \left( \sum_j w_{i|j}^{\Lambda-1} h_j^{\Lambda-1} + b_i^{\Lambda-1} \right) \right) \quad (22.3)$$

## Activation Functions $\mathcal{A}_i^\lambda : \mathbb{R} \rightarrow \mathbb{R}$

Activation functions must be nonlinear.

- **Step function (Perceptron)**

$$\mathcal{A}(x) = \mathbb{1}(x > 0) \quad (22.4)$$

Zero for  $x \leq 0$ , one for  $x > 0$ .

- **Sigmoid function**

$$\mathcal{A}(x) = \frac{1}{1 + e^{-x}} = \text{sig}(x) \quad (22.5)$$

Smooth, monotonically increasing function.  $\text{sig}(-\infty) = 0, \text{sig}(0) = 0.5, \text{sig}(\infty) = 1$ .

$$\text{sig}(x) + \text{sig}(-x) = \frac{1}{1 + e^{-x}} + \frac{1}{1 + e^x} \quad (22.6)$$

$$= \frac{2 + e^x + e^{-x}}{2 + e^x + e^{-x}} \quad (22.7)$$

$$= 1 \quad (22.8)$$

- **Hyperbolic tangent**

$$\mathcal{A}(x) = \tanh(x) = \frac{e^x - e^{-x}}{e^x + e^{-x}} \quad (22.9)$$

Smooth, monotonically increasing function.  $\tanh(-\infty) = -1, \tanh(0) = 0, \tanh(\infty) = 1$ .

Odd function:

$$\tanh(-x) = -\tanh(x) \quad (22.10)$$

Whereas  $\text{sig}(x) \in [0, 1]$ ,  $\tanh(x) \in [-1, 1]$ .

- **ReLU (Rectified Linear Unit)**

$$\mathcal{A}(x) = x \mathbb{1}(x > 0) = \max(0, x) . \quad (22.11)$$

Compare this to the step function.

- **Swish**

$$\mathcal{A}(x) = x \text{sig}(x) \quad (22.12)$$

- **Softmax**

$$\mathcal{A}(x_i|x.) = \frac{e^{x_i}}{\sum_i e^{x_i}} \quad (22.13)$$

It's called softmax because if we approximate the exponentials, both in the numerator and denominator of Eq.(22.13), by the largest one, we get

$$\mathcal{A}(x_i|x.) \approx \mathbb{1}(x_i = \max_k x_k) . \quad (22.14)$$

The softmax definition implies that the bnet nodes within a softmax layer are fully connected by arrows to form a "clique".

For 2 nodes  $x_0, x_1$ ,

$$\mathcal{A}(x_0|x.) = \frac{e^{x_0}}{e^{x_0} + e^{x_1}} \quad (22.15)$$

$$= \text{sig}(x_0 - x_1) , \quad (22.16)$$

$$\mathcal{A}(x_1|x.) = \text{sig}(x_1 - x_0) . \quad (22.17)$$

# Weight optimization via supervised training and gradient descent

The bnet of Fig.22.1 is used for classification of a single data point  $x$ . It assumes that the weights  $w_{i|j}^\lambda, b_i^\lambda$  are given.

To find the optimum weights via supervised training and gradient descent, one uses the bnet Fig.22.2.

In Fig.22.2, the nodes in Fig.22.1 become sampling space vectors. For example,  $\underline{x}$  becomes  $\vec{x}$ , where the components of  $\vec{x}$  in sampling space are  $\underline{x}[s] \in \{0, 1\}^{numx}$  for  $s = 0, 1, \dots, nsam(\vec{x}) - 1$ .

$nsam(\vec{x})$  is the number of samples used to calculate the gradient during each **stage (aka iteration)** of Fig.22.2. We will also refer to  $nsam(\vec{x})$  as the **mini-batch size**. A **mini-batch** is a subset of the training data set.

To train a bnet with a data set (d-set), the standard procedure is to split the d-set into 3 parts:

1. **training d-set**,
2. **testing1 d-set**, for tuning of hyperparameters like  $nsam(\vec{x})$ ,  $\Lambda$ , and  $nunh(i)$  for each  $i$ .
3. **testing2 d-set**, for measuring how well the model tuned with the testing1 d-set performs.

The training d-set is itself split into mini-batches. An **epoch** is a pass through all the training d-set.

Define

$$W_{i|j}^\lambda = [w_{i|j}^\lambda, b_i^\lambda] . \quad (22.18)$$

These are the transition probability matrices, printed in blue, for the nodes of the bnet Fig.22.2:

$$P(x.[s]) = \text{given} . \quad (22.19)$$

$$P(y.[s] \mid x.[s]) = \text{given} . \quad (22.20)$$

$$P(h_i^\lambda[s] \mid h_i^{\lambda-1}[s]) = \delta \left( h_i^\lambda[s], \mathcal{A}_i^\lambda \left( \sum_j w_{i|j}^{\lambda-1} h_j^{\lambda-1}[s] + b_i^{\lambda-1} \right) \right) \quad (22.21)$$

$$P(Y_i[s] \mid h_i^{\Lambda-1}[s]) = \delta \left( Y_i[s], \mathcal{A}_i^\Lambda \left( \sum_j w_{i|j}^{\Lambda-1} h_j^{\Lambda-1}[s] + b_i^{\Lambda-1} \right) \right) \quad (22.22)$$



Figure 22.2: bnet for finding optimum weights of the bnet Fig.22.1 via supervised training and gradient descent.

$$P(W_{:,|}) = \text{given} \quad (22.23)$$

The first time it is used,  $W_{:,|}$  is arbitrary. After the first time, it is determined by previous stage.

$$P(W_{:,|}^\lambda | W_{:,|}) = \delta(W_{:,|}^\lambda, (W_{:,|})^\lambda) \quad (22.24)$$

$$P(\mathcal{E} | \vec{y}_{:,|}, \vec{Y}_{:,|}) = \frac{1}{nsam(\vec{x}_{:,|})} \sum_s \sum_i d(y_i[s], Y_i[s]) , \quad (22.25)$$

where

$$d(y, Y) = |y - Y|^2 . \quad (22.26)$$

If  $y, Y \in [0, 1]$ , one can use this instead

$$d(y, Y) = XE(y \rightarrow Y) = -y \ln Y - (1 - y) \ln(1 - Y) . \quad (22.27)$$

$$P((W')_{i|j}^\lambda | \mathcal{E}, W_{:,|}) = \delta((W')_{i|j}^\lambda, W_{i|j}^\lambda - \eta \partial_{W_{i|j}^\lambda} \mathcal{E}) \quad (22.28)$$

$\eta > 0$  is called the learning rate. This method of minimizing the error  $\mathcal{E}$  is called gradient descent.  $W' - W = \Delta W = -\eta \partial_W \mathcal{E}$  so  $\Delta \mathcal{E} = \frac{-1}{\eta} (\Delta W)^2 < 0$ .

## Non-dense layers

The transition probability matrix for a non-dense layer is of the form:

$$P(h_i^\lambda[s] | h_i^{\lambda-1}[s]) = \delta(h_i^\lambda[s], H_i^\lambda[s]) , \quad (22.29)$$

where  $H_i^\lambda[s]$  will be specified below for each type of non-dense layer.

- **Dropout Layer**

The dropout layer was invented in Ref.[14]. To dropout nodes from a fixed layer  $\lambda$ : For all  $i$  of layer  $\lambda$ , define a new node  $\underline{r}_i^\lambda$  with an arrow  $\underline{r}_i^\lambda \rightarrow \underline{h}_i^\lambda$ . For  $r \in \{0, 1\}$ , and some  $p \in (0, 1)$ , define

$$P(r_i^\lambda = r) = [p]^r [1 - p]^{1-r} \text{ (Bernoulli dist.)} . \quad (22.30)$$

Now one has

$$P(h_i^\lambda[s] | h_i^{\lambda-1}[s], r_i^\lambda) = \delta(h_i^\lambda[s], H_i^\lambda[s]) , \quad (22.31)$$

where

$$H_i^\lambda[s] = \mathcal{A}_i^\lambda(r_i^\lambda \sum_j w_{i|j}^\lambda h_j^{\lambda-1}[s] + b_i^\lambda) . \quad (22.32)$$

This reduces overfitting. Overfitting might occur if the weights follow too closely several similar minibatches. This dropout procedure adds a random component to each minibatch making groups of similar minibatches less likely.

The random  $\underline{r}_i^\lambda$  nodes that induce dropout are only used in the training bnet Fig.22.2, not in the classification bnet Fig.22.1. We prefer to remove the  $\underline{r}_i^\lambda$  stochasticity from classification and for Fig.22.1 to act as an average over sampling space of Fig.22.2. Therefore, if weights  $w_{i|j}^\lambda$  are obtained for a dropout layer  $\lambda$  in Fig.22.2, then that layer is used in Fig.22.1 with no  $\underline{r}_i^\lambda$  nodes but with weights  $\langle r_i^\lambda \rangle w_{i|j}^\lambda = p w_{i|j}^\lambda$ .

Note that dropout adds non-deterministic nodes to a NN, which in their vanilla form only have deterministic nodes.

- **Convolutional Layer**

- 1-dim

Filter function  $\mathcal{F} : \{0, 1, \dots, numf - 1\} \rightarrow \mathbb{R}$ .

$\sigma$ =stride length

For  $i \in \{0, 1, \dots, numh(\lambda) - 1\}$ , let

$$H_i^\lambda[s] = \sum_{j=0}^{numf-1} h_{j+i\sigma}^{\lambda-1}[s] \mathcal{F}(j) . \quad (22.33)$$

For the indices not to go out of bounds in Eq.(22.33), we must have

$$numh(\lambda - 1) - 1 = numf - 1 + (numh(\lambda) - 1)\sigma \quad (22.34)$$

so

$$numh(\lambda) = \frac{1}{\sigma} [numh(\lambda - 1) - numf] + 1 . \quad (22.35)$$

- 2-dim

$h_i^\lambda[s]$  becomes  $h_{(i,j)}^\lambda[s]$ . Do 1-dim convolution along both  $i$  and  $j$  axes.

- **Pooling Layers (MaxPool, AvgPool)**

Here each node  $i$  of layer  $\lambda$  is impinged by arrows from a subset  $Pool(i)$  of the set of all nodes of the previous layer  $\lambda - 1$ . Partition set  $\{0, 1, \dots, numh(\lambda - 1) - 1\}$  into  $numh(\lambda)$  mutually disjoint, nonempty sets called  $Pool(i)$ , where  $i \in \{0, 1, \dots, numh(\lambda) - 1\}$ .

- AvgPool

$$H_i^\lambda[s] = \frac{1}{size(Pool(i))} \sum_{j \in Pool(i)} h_j^{\lambda-1}[s] \quad (22.36)$$

- MaxPool

$$H_i^\lambda[s] = \max_{j \in Pool(i)} h_j^{\lambda-1}[s] \quad (22.37)$$

## Autoencoder NN

If the sequence

$$numx, numh(0), numh(1), \dots, numh(\Lambda - 1), numy \quad (22.38)$$

first decreases monotonically up to layer  $\lambda_{min}$ , then increases monotonically until  $numy = numx$ , then the NN is called an **autoencoder NN**. Autoencoders are useful for unsupervised learning and feature reduction. In this case,  $Y$  estimates  $x$ . The layers before layer  $\lambda_{min}$  are called the **encoder**, and those after  $\lambda_{min}$  are called the **decoder**. Layer  $\lambda_{min}$  is called the **code**.

# Chapter 23

## Non-negative Matrix Factorization

Based on Ref.[15].

Given matrix  $V$ , factor it into product of two matrices

$$V = WH, \quad (23.1)$$

where all 3 matrices have non-negative entries.

$V \in \mathbb{R}_{\geq 0}^{nv \times na}$ : visible info matrix

$W \in \mathbb{R}_{\geq 0}^{nv \times nh}$ : weight info matrix

$H \in \mathbb{R}_{\geq 0}^{nh \times na}$ : hidden info matrix

Usually,  $nv > nh < na$  so compression of information (aka dimensional reduction, clustering)  
 B net interpretation: Express node  $\underline{v}$  as a chain of two nodes.

$$\underline{v} \longleftarrow \underline{a} \quad = \quad \underline{w} \longleftarrow \underline{h} \longleftarrow \underline{a}$$

Figure 23.1: B net interpretation of non-negative matrix factorization.

Node transition matrices, printed in blue, for Fig.23.1.

$$P(\underline{v} = w | a) = \frac{V_{w,a}}{\sum_w V_{w,a}} \quad (23.2)$$

$$P(w | h) = \frac{W_{w,h}}{\sum_w W_{w,h}} \quad (23.3)$$

$$P(h | a) = \frac{\sum_w W_{w,h} V_{w,a}}{\sum_w V_{w,a}} \quad (23.4)$$

---

Simplest recursive algorithm:

Initialize: Choose  $nh$ . Choose  $W^{(0)}$  and  $H^{(0)}$  that have non-negative entries.

Update: For  $n = 0, 1, \dots$ , do

$$H_{i,j}^{(n+1)} \leftarrow H_{i,j}^{(n)} \frac{[(W^{(n)})^T V]_{i,j}}{[(W^{(n)})^T \underbrace{W^{(n)} H^{(n)}}_{\approx V}]_{i,j}} \quad (23.5)$$

and

$$W_{i,j}^{(n+1)} \leftarrow W_{i,j}^{(n)} \frac{[V(H^{(n+1)})^T]_{i,j}}{[\underbrace{W^{(n)} H^{(n+1)}}_{\approx V} (H^{(n+1)})^T]_{i,j}} . \quad (23.6)$$

After each step, record error defined by

$$\mathcal{E}^{(n)} = \| V - W^{(n)} H^{(n)} \|_2 . \quad (23.7)$$

Using 2-norm, aka Frobenius matrix norm. Continue until reach acceptable error.

Can also use Kullback-Liebr divergence for error:

$$\mathcal{E} = \sum_a P(a) D_{KL}(P(\underline{v} = w|a) \parallel \sum_h P(w|h) P(h|a)) , \quad (23.8)$$

for some arbitrary choice of prior  $P(a)$ . For example, can choose  $P(a)$  uniform.



# Chapter 24

## Program evaluation and review technique (PERT)

This chapter is based on Refs.[16] and [17].

PERT diagrams are used for scheduling a project consisting of a series of interdependent activities and estimating how long it will take to finish the project. PERT diagrams were invented by the NAVY in 1958 to manage a submarine project. Nowadays they are taught in many business and management courses.

A **PERT diagram** is a Directed Acyclic Graph (DAG) with the following properties. (See Fig.24.2 for an example of a PERT diagram). The nodes  $\underline{E}_i$  for  $i = 1, 2, \dots, ne$  of a PERT diagram are called **events**. The edges  $i \rightarrow j$  of a PERT diagram are called **activities**. An event represents the starting (kickoff) date of one or more activities. A PERT diagram has a single root node ( $i = 1$ , start event) and a single leaf node ( $i = ne$ , end event).

The PERT diagram user must initially provide a **Duration Times (DT) table** which gives  $(DO_{i \rightarrow j}, DP_{i \rightarrow j}, DM_{i \rightarrow j})$  for each activity  $i \rightarrow j$ , where

$DO_{i \rightarrow j}$  = optimistic duration time of activity  $i \rightarrow j$

$DP_{i \rightarrow j}$  = pessimistic duration time of activity  $i \rightarrow j$

$DM_{i \rightarrow j}$  = median duration time of activity  $i \rightarrow j$

From the DT table, one calculates:

Duration time of activity  $i \rightarrow j$

$$D_{i \rightarrow j} = \frac{1}{6}(DO_{i \rightarrow j} + DP_{i \rightarrow j} + 4DM_{i \rightarrow j}) \quad (24.1)$$

Duration Variance of activity  $i \rightarrow j$

$$V_{i \rightarrow j} = \left( \frac{DO_{i \rightarrow j} - DP_{i \rightarrow j}}{DM_{i \rightarrow j}} \right)^2 \quad (24.2)$$

Often, it is convenient to define “dummy” edges with  $D_{i \rightarrow j} = 0$ . That is perfectly fine.

Define:

$TES_i$  = Earliest start time for event  $i$

$TLS_i$  = Latest start time for event  $i$

$slack_i = TLS_i - TES_i = \text{slack for event } i$

$TEF_{i \rightarrow j} = TES_i + D_{i \rightarrow j} = \text{Earliest finish time for activity } i \rightarrow j.$

$TLF_{i \rightarrow j} = TLS_j - D_{i \rightarrow j} = \text{Latest finish time for activity } i \rightarrow j. \text{ See footnote below. }^1$

A **critical path** is a directed path (i.e., a chain of connected arrows, all pointing in the same direction) going from the start to the end node, such that slack equals zero at every node visited. In a DAG, the neighbors of a node is the union of its parent and children nodes. A critical path must also have all other nodes as neighbors; i.e, the union of the neighbors of every node in the path plus the nodes in the path itself, equals all nodes in the graph.

---

**GOAL of PERT analysis:** The main goal of PERT analysis is to find, based on the data of the DT table, the interval  $[TES_i, TLS_i]$  giving a lower and an upper bound to the starting time of each node  $i$ . Another goal is to find a critical path for the PERT diagram (which represents an entire project). By adding the  $D_{i \rightarrow j}$  of each edge of the critical path, one can get the mean value of the total duration of the entire project, and by adding the variances of each edge along the critical path, one can get an estimate of the total variance of the total duration. Knowing the mean and variance of the total duration and assuming a normal distribution, one can predict the probability that the actual duration will deviate by a certain amount from its mean.

---

To calculate the interval  $[TES_i, TLS_i]$ , one follows the following two steps.

1. Assume  $TES_1 = 0$  and solve

$$TES_i = \max_{a \in pa(i)} \underbrace{(TES_a + D_{a \rightarrow i})}_{TEF_{a \rightarrow i}} \quad (24.3)$$

for  $i \in [2, ne]$ . This recursive equation is solved by what is called “forward propagation”, wherein one moves up the list of nodes  $i$  in order of increasing  $i$  starting at  $i = 1$  with  $TES_1 = 0$ .

2. Assume  $TLS_{ne} = TES_{ne}$  and solve

$$TLS_i = \min_{b \in ch(i)} \underbrace{(TLS_b - D_{i \rightarrow b})}_{TLF_{i \rightarrow b}} \quad (24.4)$$

for  $i \in [1, ne - 1]$ . This recursive equation is solved by what is called “backward propagation”, wherein one moves down the list of nodes  $i$  in order of decreasing  $i$  starting at  $i = ne$  with  $TLS_{ne} = TES_{ne}$ .  $TES_{ne}$  is known from step 1.

Eqs.(24.3) and (24.4) are illustrated in Fig.24.1.

---

<sup>1</sup> In the popular educational literature, the edge variables  $TEF_{i \rightarrow j}$  and  $TLF_{i \rightarrow j}$  are sometimes associated with the nodes, but they are clearly edge variables. This makes things confusing. The reason this is done is that some software draws PERT diagrams as trees whereas other software draws them as DAGs. For trees, storing  $TEF_{i \rightarrow j}$  and  $TLF_{i \rightarrow j}$  in a node makes some sense but not for DAGs. You will notice that giving specific names to the variables  $TEF_{i \rightarrow j}$  and  $TLF_{i \rightarrow j}$  is unnecessary. It is possible to delete all mention of their names from this chapter without losing any details. I only declare their names in this chapter so as tell the reader what they are in case he/she hears them mentioned and wonders what they are equal to in our notation.

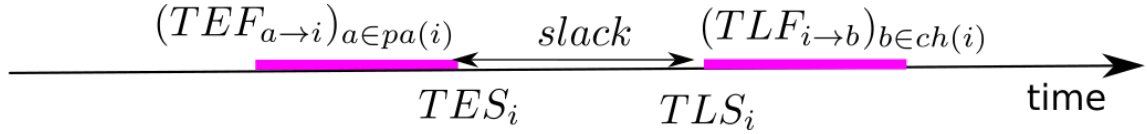


Figure 24.1:  $TES_i$  defined from info received from parents of  $i$  and  $TLS_i$  defined from info received from children of  $i$ .

## Example

To illustrate PERT analysis, we end with an example. We present the example in the form of an exercise question and then provide the answer. This example comes from Ref.[16], except for part (e) about bnets, which is our own.

**Question:** For the PERT diagram of Fig.24.2, calculate the following:

- Interval  $[TES_i, TLS_i]$  for all  $i$ .
- A critical path for this PERT diagram.
- The mean and variance of the total duration of the critical path.
- The probability that the total duration will be 225 days or less.
- A bnet interpretation of this problem.

**Answer to (a)**  $[TES_i, TLS_i]$  are given by Fig.24.3.

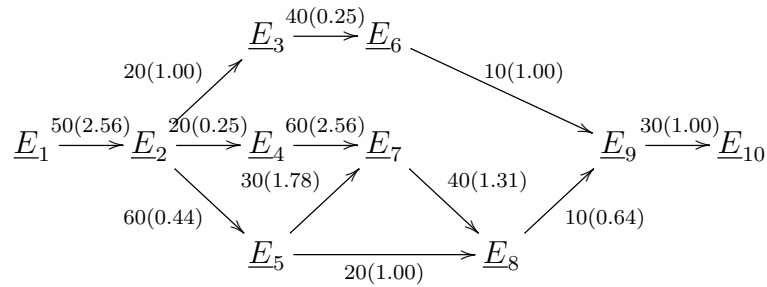


Figure 24.2: Example of a PERT diagram. The numbers attached to the arrows are the duration times  $D_{i \rightarrow j}$  in days followed by, enclosed in parentheses, the variance  $V_{i \rightarrow j}$  of that duration. The info given in this PERT diagram was derived from a DT table in Ref.[16]. The info in this PERT diagram is sufficient for calculating  $TES_i$  and  $TLS_i$  for each node  $i$ . The results of that calculation are given in Fig.24.3.

**Answer to (b)** The critical path is given in red in Fig.24.3. Note that this path does indeed have zero slack at each node it visits and the union of its neighborhood and the path itself encompasses all nodes.

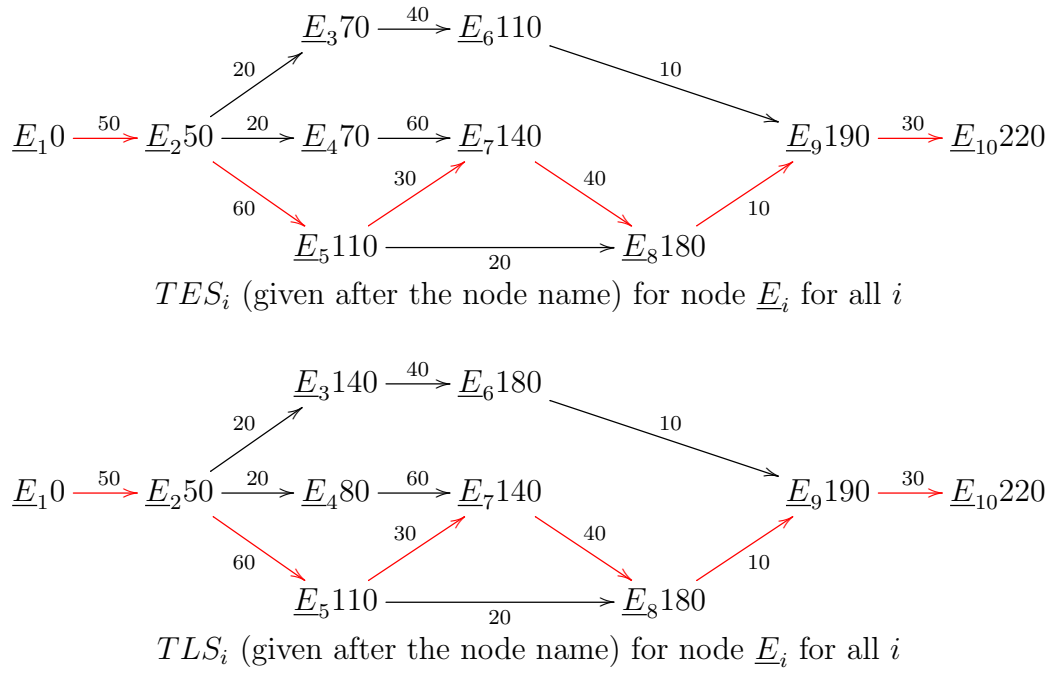


Figure 24.3: Results of calculating  $TES_i$  for all  $i$  via a forward pass, followed by calculating  $TLS_i$  for all  $i$  via a backward pass. Critical path indicated in red.

**Answer to (c)** The mean and variance of the total duration are calculated in Table 24.1.

**Answer to (d)**

$$P(x < 225) = p \left[ \frac{x - \mu}{\sigma} \leq \frac{225 - 220}{\sqrt{7.73}} \right] \quad (24.5)$$

$$= P[z \leq 1.80] \quad (24.6)$$

$$= 0.9641 \quad (24.7)$$

**Answer to (e)** Define 2 bnets.

1. The first PERT bnet is for calculating  $TES_i$  for all  $i$  and is given by Fig.24.4.

The node transition prob matrices, printed in blue, for the bnet Fig.24.4 are given by (this equation is to be evaluated recursively by a forward pass through the bnet):

$$P(TES_i | (TES_a)_{a \in pa(i)}) = \delta(TES_i, \max_{a \in pa(i)} (TES_a + D_{a \rightarrow i})) \quad (24.8)$$

2. The second PERT bnet is for calculating  $TLS_i$  for all  $i$  and is given by Fig.24.5. Note that the directions of all the arrows in the PERT diagram Fig.24.2 have been reversed so Fig.24.5 is a time reversed graph.

edge $i \rightarrow j$	duration $D_{i \rightarrow j}$	variance $V_{i \rightarrow j}$
A (1 $\rightarrow$ 2)	50	2.56
D (2 $\rightarrow$ 5)	60	0.44
G (5 $\rightarrow$ 7)	30	1.78
J (7 $\rightarrow$ 8)	40	1.31
K (8 $\rightarrow$ 9)	10	0.64
L (9 $\rightarrow$ 10)	30	1.00
Total	220	7.73

Table 24.1: Calculation of mean and variance of total duration along critical path.

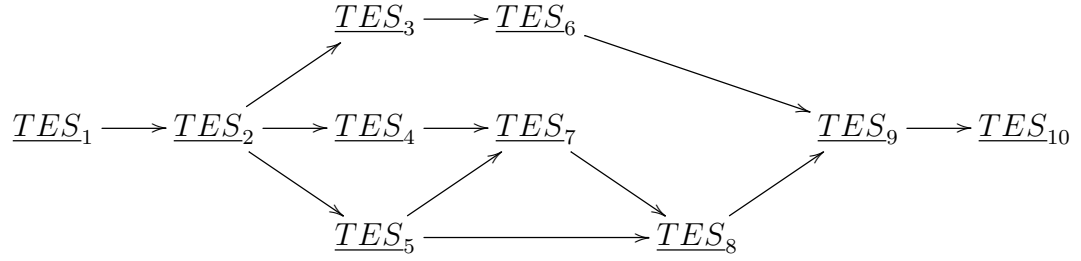


Figure 24.4: bnet for  $TES_i$  calculation.

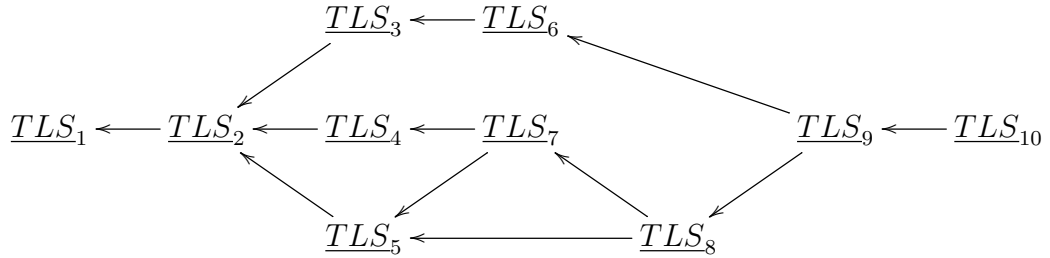


Figure 24.5: bnet for  $TLS_i$  calculation.

The node transition prob matrices, printed in blue, for the bnet Fig.24.5 are given by (this equation is to be evaluate recursively by a backward pass through the bnet):

$$P(TLS_i | (TLS_b)_{b \in pa(i)}) = \delta(TLS_i, \min_{b \in pa(i)} (TLS_b - D_{b \rightarrow i}^T)) , \quad (24.9)$$

where  $D_{i \rightarrow j}^T = D_{j \rightarrow i}$ .

# Chapter 25

## Recurrent Neural Networks

This chapter is mostly based on Ref.[18].

This chapter assumes you are familiar with the material and notation of Chapter 22 on plain Neural Nets.



Figure 25.1: Simple example of RNN with  $T = 3$

Suppose

$T$  is a positive integer.

$t = 0, 1, \dots, T - 1$ ,

$\underline{x}_i(t) \in \mathbb{R}$  for  $i = 0, 1, \dots, numx - 1$ ,

$\underline{h}_i(t) \in \mathbb{R}$  for  $i = 0, 1, \dots, numh - 1$ ,

$\underline{Y}_i(t) \in \mathbb{R}$  for  $i = 0, 1, \dots, numy - 1$ ,

$W^{h|x} \in \mathbb{R}^{numh \times numx}$ ,

$W^{h|h} \in \mathbb{R}^{numh \times numh}$ ,

$W^{y|h} \in \mathbb{R}^{numy \times numh}$ ,

$b^y \in \mathbb{R}^{numy}$ ,

$b^h \in \mathbb{R}^{numh}$ .

Henceforth,  $\underline{x}(\cdot)$  will mean the array of  $x(t)$  for all  $t$ .

The simplest kind of recurrent neural network (RNN) has the bnet Fig.25.1 with arbitrary  $T$ . The node transition matrices, printed in blue, for this bnet, are as follows.

$$P(x(\cdot)) = \text{given} \quad (25.1)$$

$$P(x(t)) = \delta(x(t), [x(\cdot)]_t) \quad (25.2)$$

$$P(h(t) \mid h(t-1), x(t)) = \delta(h(t), \mathcal{A}(W^{h|x}x(t) + W^{h|h}h(t-1) + b^h)) , \quad (25.3)$$

where  $h(-1) = 0$ .

$$P(Y(t) \mid h(t)) = \delta(Y(t), \mathcal{A}(W^{y|h}h(t) + b^y)) \quad (25.4)$$

Define

$$W^h = [W^{h|x}, W^{h|h}, b^h] , \quad (25.5)$$

and

$$W^y = [W^{y|h}, b^y] . \quad (25.6)$$

The bnet of Fig.25.1 can be used for classification once its parameters  $W^h$  and  $W^y$  have been optimized. To optimize those parameters via gradient descent, one can use the bnet of Fig.25.2.

Let  $s = 0, 1, \dots, nsam(\vec{x}) - 1$  be the labels for a minibatch of samples. The node transition matrices, printed in blue, for bnet Fig.25.2, are as follows.

$$P(x(\cdot)[s]) = \text{given} \quad (25.7)$$

$$P(x(t)[s]) = \delta(x(t)[s], [x(\cdot)]_t[s]) \quad (25.8)$$

$$P(h(t)[s] \mid h(t-1)[s], x(t)[s]) = \delta(h(t)[s], \mathcal{A}(W^{h|x}x(t)[s] + W^{h|h}h(t-1)[s] + b^h)) \quad (25.9)$$

$$P(Y(t)[s] \mid h(t-1)[s]) = \delta(Y(t)[s], \mathcal{A}(W^{y|h}h(t-1)[s] + b^y)) \quad (25.10)$$

$$P(y(\cdot)[s] \mid x(\cdot)[s]) = \text{given} \quad (25.11)$$



Figure 25.2: RNN bnet used to optimize parameters  $W^h$  and  $W^y$  of RNN bnet Fig.25.1.

$$P(\mathcal{E}(t) \mid \vec{y}(\cdot), \vec{Y}(t)) = \frac{1}{nsam(\vec{x})} \sum_s d(y(t)[s], Y(t)[s]) , \quad (25.12)$$

where

$$d(y, Y) = |y - Y|^2 . \quad (25.13)$$

If  $y, Y \in [0, 1]$ , one can use this instead

$$d(y, Y) = XE(y \rightarrow Y) = -y \ln Y - (1 - y) \ln(1 - Y) . \quad (25.14)$$



$$P(\mathcal{E} \mid [\mathcal{E}(t)]_{\forall t}) = \delta(\mathcal{E}, \sum_t \mathcal{E}(t)) \quad (25.15)$$

For  $a = h, y$ ,

$$P(W^a) = \text{given} . \quad (25.16)$$

The first time it is used,  $W^a$  is fairly arbitrary. Afterwards, it is determined by previous horizontal stage.

$$P((W^a)' \mid \mathcal{E}, W^a) = \delta((W^a)', W^a - \eta^a \partial_{W^a} \mathcal{E}) . \quad (25.17)$$

$\eta^a > 0$  is the learning rate for  $W^a$ .

## Language Sequence Modeling

Figs.25.1, and 25.2 with arbitrary  $T$  can be used as follows to do Language Sequence Modeling.

For this usecase, one must train with the following transition matrix for node  $\vec{y}(\cdot)$ :

$$P(y(\cdot)[s] \mid x(\cdot)[s]) = \prod_t \mathbb{1}( y(t)[s] = P(x(t)[s] \mid [x(t')[s]]_{t' < t}) ) \quad (25.18)$$

With such training, one gets

$$P(Y(t) \mid h(t)) = \mathbb{1}( Y(t) = P(x(t) \mid [x(t')[s]]_{t' < t}) ) . \quad (25.19)$$

Therefore,

$$Y(0) = P(x(0)) , \quad (25.20)$$

$$Y(1) = P(x(1) \mid x(0)) , \quad (25.21)$$

$$Y(2) = P(x(2) \mid x(0), x(1)) , \quad (25.22)$$

and so on.

We can use this to:

- predict the probability of a sentence,  
example: Get  $P(x(0), x(1), x(2))$ .
- predict the most likely next word in a sentence,  
example: Get  $P(x(2) \mid x(0), x(1))$ .

- generate fake sentences.

example:

Get  $x(0) \sim P(x(0))$ .

Next get  $x(1) \sim P(x(1)|x(0))$ .

Next get  $x(2) \sim P(x(2)|x(0), x(1))$ .

## Other types of RNN

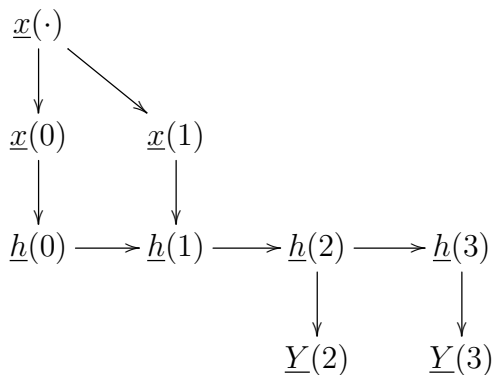


Figure 25.3: RNN bnet of the many to many kind. This one can be used for translation.  $x(0)$  and  $x(1)$  might denote two words of an English sentence, and  $Y(2)$  and  $Y(3)$  might be their Italian translation.

Let  $\mathcal{T} = \{0, 1, \dots, T-1\}$ , and  $\mathcal{T}^x, \mathcal{T}^y \subset \mathcal{T}$ . Above, we assumed that  $\underline{x}(t)$  and  $\underline{Y}(t)$  were both defined for all  $t \in \mathcal{T}$ . More generally, they might be defined only for subsets of  $\mathcal{T}$ :  $\underline{x}(t)$  for  $t \in \mathcal{T}^x$  and  $\underline{Y}(t)$  for  $t \in \mathcal{T}^y$ . If  $|\mathcal{T}^x| = 1$  and  $|\mathcal{T}^y| > 1$ , we say the RNN bnet is of the 1 to many kind. In general, can have **1 to 1**, **1 to many**, **many to 1**, **many to many** RNN bnets.

Plain RNNs can suffer from the **vanishing or exploding gradients problem**. There are various ways to mitigate this (good choice of initial  $W^h$  and  $W^y$ , good choice of activation functions, regularization). Or by using GRU or LSTM (discussed below). **GRU and LSTM** were designed to mitigate the vanishing or exploding gradients problem. They are very popular in NLP (Natural Language Processing).

## Long Short Term Memory (LSTM) unit (1997)

This section is based on Wikipedia article Ref.[19]. In this section,  $\odot$  will denote the Hadamard matrix product (elementwise product).

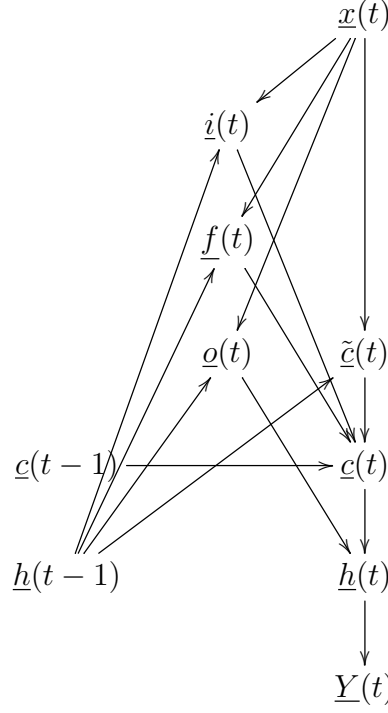


Figure 25.4: bnet for a Long Short Term Memory (LSTM) unit.

Let

$\underline{x}(t) \in \mathbb{R}^{numx}$ : input vector to the LSTM unit

$\underline{f}(t) \in \mathbb{R}^{numh}$ : forget gate's activation vector

$\underline{i}(t) \in \mathbb{R}^{numh}$ : input/update gate's activation vector

$\underline{o}(t) \in \mathbb{R}^{numh}$ : output gate's activation vector

$\underline{h}(t) \in \mathbb{R}^{numh}$ : hidden state vector also known as output vector of the LSTM unit

$\tilde{\underline{c}}(t) \in \mathbb{R}^{numh}$ : cell input activation vector

$\underline{c}(t) \in \mathbb{R}^{numh}$ : cell state vector

$\underline{Y}(t) \in \mathbb{R}^{numy}$ : classification of  $x(t)$ .

$W \in \mathbb{R}^{numh \times numx}$ ,  $U \in \mathbb{R}^{numh \times numh}$  and  $b \in \mathbb{R}^{numh}$ : weight matrices and bias vectors, parameters learned by training.

$\mathcal{W}_{y|h} \in \mathbb{R}^{numy \times numh}$ : weight matrix

Fig.25.4 is a bnet net for a LSTM unit. The node transition matrices, printed in blue, for this bnet, are as follows.

$$P(f(t)|x(t), h(t-1)) = \mathbb{1}(f(t) = \text{sig}(W^{f|x}x(t) + U^{f|h}h(t-1) + b^f)) , \quad (25.23)$$

where  $h(-1) = 0$ .

$$P(i(t)|x(t), h(t-1)) = \mathbb{1}(i(t) = \text{sig}(W^{i|x}x(t) + U^{i|h}h(t-1) + b^i)) \quad (25.24)$$

$$P(o(t)|x(t), h(t-1)) = \mathbb{1}(o(t) = \text{sig}(W^{o|x}x(t) + U^{o|h}h(t-1) + b^o)) \quad (25.25)$$

$$P(\tilde{c}(t)|x(t), h(t-1)) = \mathbb{1}(\tilde{c}(t) = \tanh(W^{c|x}x(t) + U^{c|h}h(t-1) + b^c)) \quad (25.26)$$

$$P(c(t)|f(t), c(t-1), i(t), \tilde{c}(t)) = \mathbb{1}(c(t) = f(t) \odot c(t-1) + i(t) \odot \tilde{c}(t)) \quad (25.27)$$

$$P(h(t)|o(t), c(t)) = \mathbb{1}(h(t) = o(t) \odot \tanh(c(t))) \quad (25.28)$$

$$P(Y(t)|h(t)) = \mathbb{1}(Y(t) = \mathcal{A}(\mathcal{W}^{y|h}h(t) + b^y)) \quad (25.29)$$

## Gated Recurrence Unit (GRU) (2014)

This section is based on Wikipedia article Ref.[20]. In this section,  $\odot$  will denote the Hadamard matrix product (elementwise product).

GRU is a more recent (17 years later) attempt at simplifying LSTM unit.

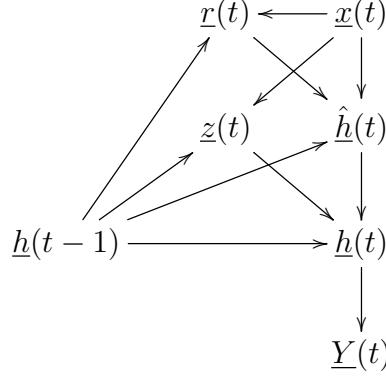


Figure 25.5: bnet for a Gated Recurrent Unit (GRU).

Let

$\underline{x}(t) \in \mathbb{R}^{numx}$ : input vector

$\underline{h}(t) \in \mathbb{R}^{numh}$ : output vector

$\hat{\underline{h}}(t) \in \mathbb{R}^{numh}$ : candidate activation vector

$\underline{z}(t) \in \mathbb{R}^{numh}$ : update gate vector

$\underline{r}(t) \in \mathbb{R}^{numh}$ : reset gate vector

$\underline{Y}(t) \in \mathbb{R}^{numy}$ : classification of  $x(t)$ .

$W \in \mathbb{R}^{numh \times numx}$ ,  $U \in \mathbb{R}^{numh \times numh}$  and  $b \in \mathbb{R}^{numh}$ : weight matrices and bias vectors, parameters learned by training.

$\mathcal{W}_{y|h} \in \mathbb{R}^{numy \times numh}$ : weight matrix

Fig.25.5 is a bnet net for a GRU. The node transition matrices, printed in blue, for this bnet, are as follows.

$$P(z(t)|x(t), h(t-1)) = \mathbb{1}( \quad z(t) = \text{sig}(W^{z|x}x(t) + U^{z|h}h(t-1) + b^z) \quad ) , \quad (25.30)$$

where  $h(-1) = 0$ .

$$P(r(t)|x(t), h(t-1)) = \mathbb{1}( \quad r(t) = \text{sig}(W^{r|x}x(t) + U^{r|h}h(t-1) + b^r) \quad ) \quad (25.31)$$

$$P(\hat{h}(t)|x(t), r(t), h(t-1)) = \mathbb{1}( \quad \hat{h}(t) = \tanh(W^{h|x}x(t) + U^{h|h}(r(t) \odot h(t-1)) + b^h) \quad ) \quad (25.32)$$

$$P(h(t)|z(t), h(t-1), \hat{h}(t)) = \mathbb{1}(h(t) = (1 - z(t)) \odot h(t-1) + z(t) \odot \hat{h}(t)) \quad (25.33)$$

$$P(Y(t)|h(t)) = \mathbb{1}(Y(t) = \mathcal{A}(\mathcal{W}^{y|h}h(t) + b^y)) \quad (25.34)$$

# Chapter 26

## Reinforcement Learning (RL)



Figure 26.1: Axes for episode time and episode number.

I based this chapter on the following references. Refs.[21][22]

In RL, we consider an “agent” or robot that is learning.

Let  $T \in \mathbb{Z}_{>0}$  be the duration time of an **episode** of learning. If  $T = \infty$ , we say that the episode has an infinite time horizon. A learning episode will evolve towards the right, for times  $t = 0, 1, \dots, T - 1$ . We will consider multiple learning episodes. The episode number will evolve from top to bottom. This is illustrated in Fig.26.1.

Let  $\underline{s}_t \in S_{\underline{s}}$  for  $t \in \mathbb{Z}_{[0, T-1]}$  be random variables that record the **state** of the agent at various times  $t$ .

Let  $\underline{a}_t \in S_{\underline{a}}$  for  $t \in \mathbb{Z}_{[0, T-1]}$  be random variables that record the **action** of the agent at various times  $t$ .



Figure 26.2: State-Action-Reward dynamical bnet

Let  $\underline{\theta}_t \in S_{\underline{\theta}}$  for  $t \in \mathbb{Z}_{[0, T-1]}$  be random variables that record the **policy parameters** at various times  $t$ .

For  $\underline{X} \in \{\underline{s}, \underline{a}, \underline{\theta}\}$ , define  $\underline{X}$  followed by a dot to be the vector

$$\underline{X}_{\cdot} = [\underline{X}_0, \underline{X}_1, \dots, \underline{X}_{T-1}] . \quad (26.1)$$

Also let

$$\underline{X}_{\geq t} = [\underline{X}_t, \underline{X}_{t+1}, \dots, \underline{X}_{T-1}] . \quad (26.2)$$

Fig.26.2 shows the basic State-Action-Reward bnet for an agent that is learning. The transition probabilities for the nodes of Fig.26.2 are given in blue below:

$$P(a_t | s_t, \theta_t) = \text{given.} \quad (26.3)$$

$P(a_t | s_t, \theta_t)$  is called a **policy with parameter**  $\theta_t$ .

$$P(s_t | s_{t-1}, a_{t-1}) = \text{given.} \quad (26.4)$$

$P(s_t | s_{t-1}, a_{t-1})$  is called the **transition matrix of the model**.  $P(s_t | s_{t-1}, a_{t-1})$  reduces to  $P(s_0)$  when  $t = 0$ .

$$P(r_t | s_t, a_t) = \delta(r_t, r(s_t, a_t)) . \quad (26.5)$$

$r : S_{\underline{s}} \times S_{\underline{a}} \rightarrow \mathbb{R}$  is a given **one-time reward function**.

Note that

$$P(s_{\cdot}, a_{\cdot} | \theta_{\cdot}) = \prod_{t=0}^{T-1} \{P(s_t | s_{t-1}, a_{t-1}) P(a_t | s_t, \theta_t)\} . \quad (26.6)$$

Define the **all times reward**  $\Sigma$  by

$$\Sigma(s_{\cdot}, a_{\cdot}) = \sum_{t=0}^{T-1} \gamma^t r(s_t, a_t) . \quad (26.7)$$



Here  $0 < \gamma < 1$ .  $\gamma$ , called the **discount rate**, is included to assure convergence of  $\Sigma$  when  $T \rightarrow \infty$ . If  $r(s_t, a_t) < K$  for all  $t$ , then  $\Sigma < K \frac{1}{1-\gamma}$ .

Define the **objective (i.e. goal) function**  $E\Sigma(\theta.)$  by

$$E\Sigma(\theta.) = E_{\underline{s}, \underline{a} | \theta.} \Sigma(\underline{s}, \underline{a}) = \sum_{s., a.} P(s., a. | \theta.) \Sigma(s., a.) \quad (26.8)$$

The goal of RL is to maximize the objective function over its parameters  $\theta.$ . The parameters  $\theta^*$  that maximize the objective function are the optimum strategy:

$$\theta.^* = \operatorname{argmax}_{\theta.} E\Sigma(\theta.) \quad (26.9)$$

---

Define a **future reward** for times  $\geq t$  as:

$$\Sigma_{\geq t}((s_{t'}, a_{t'})_{t' \geq t}) = \sum_{t'=t}^{T-1} \gamma^{t'-t} r(s_{t'}, a_{t'}) \quad (26.10)$$

Define the following **expected conditional future rewards** (rewards for times  $\geq t$ , conditioned on certain quantities having given values):

$$v_t = v(s_t, a_t; \theta.) = E_{\underline{s}, \underline{a} | s_t, a_t, \theta.} [\Sigma_{\geq t}] \quad (26.11)$$

$$V_t = V(s_t; \theta.) = E_{\underline{s}, \underline{a} | s_t, \theta.} [\Sigma_{\geq t}] = E_{\underline{a} | s_t, \theta.} [v(s_t, \underline{a}; \theta.)] \quad (26.12)$$

$v$  is usually called  $Q$  in the literature. We will refer to  $Q$  as  $v$  in order to follow a convention wherein an  $\underline{a}_t$ -average changes a lower case letter to an upper case one.

We will sometimes write  $v(s_t, a_t)$  instead of  $v(s_t, a_t; \theta.)$ .

Since  $E\Sigma_{\geq t}$  only depends on  $\theta_{\geq t}$ ,  $v(s_t, a_t; \theta.) = v(s_t, a_t; \theta_{\geq t})$ , and  $V(s_t; \theta.) = V(s_t; \theta_{\geq t})$ .

Note that the objective function  $E\Sigma$  can be expressed in terms of  $v_0$  by averaging over its unaveraged parameters:

$$E\Sigma(\theta.) = E_{\underline{s}_0, \underline{a}_0 | \theta_0} v(\underline{s}_0, \underline{a}_0; \theta.) \quad (26.13)$$

Define a **one-time reward** and an **expected conditional one-time reward** as:

$$r_t = r(s_t, a_t) \quad (26.14)$$

$$R_t = R(s_t; \theta_t) = E_{\underline{a}_t | s_t, \theta_t} [r(s_t, \underline{a}_t)] \quad (26.15)$$

---

Note that

$$\Sigma_{\geq t} = r_t + \gamma r_{t+1} + \gamma^2 r_{t+2} + \dots + \gamma^{T-1-t} r_{t+(T-1-t)} \quad (26.16)$$

$$= r_t + \gamma \Sigma_{\geq t+1}; \quad (26.17)$$

If we take  $E_{\underline{s}, \underline{a} | s_t, a_t, \theta.} [\cdot]$  of both sides of Eq.(26.17), we get

$$v_t = r_t + \gamma E_{\underline{s}_{t+1}, \underline{a}_{t+1} | \theta.} [v_{t+1}] \quad (26.18)$$

If we take  $E_{\underline{s}, \underline{a} | s_t, \theta.}[\cdot]$  of both sides of Eq.(26.17), we get

$$V_t = R_t + \gamma E_{\underline{s}_{t+1} | \theta.}[V_{t+1}] . \quad (26.19)$$

Note that

$$\Delta r_t = r_t - R_t \quad (26.20)$$

$$= r_t - (V_t - \gamma E_{\underline{s}_{t+1} | \theta.}[V_{t+1}]) \quad (26.21)$$

$$= r_t + \gamma E_{\underline{s}_{t+1} | \theta.}[V_{t+1}] - V_t . \quad (26.22)$$

Define

$$\Delta v_t = v_t - V_t . \quad (26.23)$$

Note that

$$\Delta v_t = \Delta r_t . \quad (26.24)$$

Next, we will discuss 3 RL bnets

- exact RL bnet (exact, assumes policy is known)
- Actor-Critic RL bnet (approximate, assumes policy is known)
- Q function learning RL bnet (approximate, assumes policy is NOT known)

## Exact RL bnet

An exact RL bnet is given by Fig.26.3.

Fig.26.3 is the same as Fig.26.2 but with more nodes added in order to optimize the policy parameters. Here are the transition matrices, in blue, for the nodes not already discussed in connection to Fig.26.2.

$$P(\theta_t | \theta.) = \delta(\theta_t, (\theta.)_t) \quad (26.25)$$

$$\forall(s_t, a_t) : P(v_t(s_t, a_t) | r_t, v_{t+1}(\cdot), \theta.) = \delta(v_t(s_t, a_t), r_t + \gamma E_{\underline{s}_{t+1}, \underline{a}_{t+1} | \theta.}[v_{t+1}]) \quad (26.26)$$

$$P(\theta.' | \theta., v_0(\cdot)) = \delta(\theta.', \theta. + \alpha \partial_{\theta.} \underbrace{E_{\underline{s}_0, \underline{a}_0 | \theta_0} v(\underline{s}_0, \underline{a}_0; \theta.)}_{E\Sigma(\theta.)}) \quad (26.27)$$

$\alpha > 0$  is called the **learning rate**. This method of improving  $\theta.$  is called gradient ascent.

Concerning the gradient of the objective function, note that



Figure 26.3: Exact RL bnet.  $v_t(\cdot)$  means the array  $[v_t(s_t, a_t)]_{\forall s_t, a_t}$ . The following arrows are implicit: for all  $t$ , arrow from  $\underline{\theta}.$   $\rightarrow \underline{v}_t(\cdot)$ . We did not draw those arrows so as not to clutter the diagram.

$$\partial_{\theta_t} E\Sigma(\theta.) = \sum_{s., a.} \partial_{\theta_t} P(s., a. | \theta.) \Sigma(s., a.) \quad (26.28)$$

$$= \sum_{s., a.} P(s., a. | \theta.) \partial_{\theta_t} \ln P(s., a. | \theta.) \Sigma(s., a.) \quad (26.29)$$

$$= E_{\underline{s}, \underline{a} | \theta.} \{ \partial_{\theta_t} \ln P(a_t | s_t, \theta_t) \Sigma(s., a.) \} . \quad (26.30)$$

If we run the agent  $nsam(\vec{s}_t)$  times and obtain samples  $s_t[i], a_t[i]$  for all  $t$  and for  $i = 0, 1, \dots, nsam(\vec{s}_t) - 1$ , we can express this gradient as follows:

$$\partial_{\theta_t} E\Sigma(\theta.) \approx \frac{1}{nsam(\vec{s}_t)} \sum_i \sum_{t=0}^{T-1} \partial_{\theta_t} \ln P(a_t[i] | s_t[i], \theta_t) r(s_t[i], a_t[i]) . \quad (26.31)$$

The exact RL bnet Fig.26.3 is difficult to use to calculate the optimum parameters  $\theta^*$ . The problem is that  $\underline{s}_t$  propagates towards the future and the  $\underline{v}_t(\cdot)$  propagates towards the past, so we don't have a Markov Chain with a chain link for each  $t$  (i.e., a dynamical bnet) in the episode time direction. Hence, people have come up with approximate RL bnets that are doubly dynamical (i.e., dynamical along the episode time and episode number axes.) We discuss some of those approximate RL bnets next.

## Actor-Critic RL bnet

For the actor-critic RL bnet, we approximate Eq.(26.31) by

$$\partial_{\theta_t} E\Sigma(\theta.) \approx \frac{1}{nsam(\vec{s})} \sum_i \sum_{t=0}^{T-1} \underbrace{\partial_{\theta_t} \ln P(a_t[i] | s_t[i], \theta_t)}_{Actor} \underbrace{\Delta r_t(s_t[i], a_t[i])}_{Critic} \quad (26.32)$$

The actor-critic RL bnet is given by Fig.26.4. This bnet is approximate and assumes that the policy is known. The transition matrices for its nodes are given in blue below.

$$P(\theta_t) = \text{given} \quad (26.33)$$

$$P(s_t[i] | s_{t-1}[i], a_{t-1}[i]) = \text{given} \quad (26.34)$$

$$P(a_t[i] | s_t[i], \theta_t) = \text{given} \quad (26.35)$$

$$P(r_t[i] | s_t[i], a_t[i]) = \delta(r_t[i], r(s_t[i], a_t[i])) \quad (26.36)$$

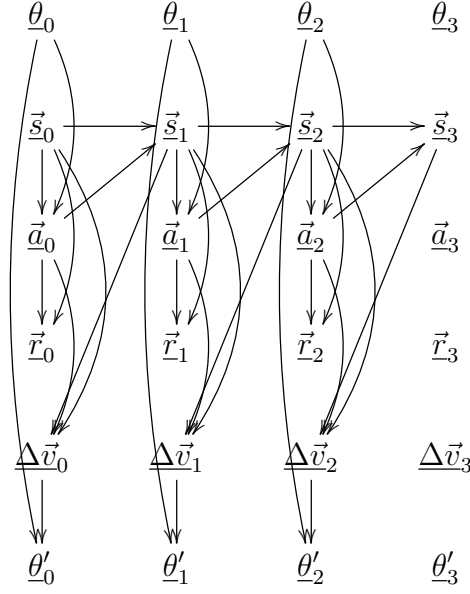


Figure 26.4: Actor-Critic RL bnet.

$r : S_{\underline{s}} \times S_{\underline{a}} \rightarrow \mathbb{R}$  is given.

$$P(\Delta v_t[i] \mid s_t[i], a_t[i], s_{t+1}[i]) = \delta(\Delta v_t[i], r(s_t[i], a_t[i]) + \gamma \hat{V}(s_{t+1}[i]; \phi') - \hat{V}(s_t[i]; \phi)) . \quad (26.37)$$

$$P(\theta') = \delta(\theta', \theta_t + \alpha \partial_{\theta_t} \sum_i \ln P(a_t[i] \mid s_t[i], \theta_t) \Delta v_t[i]) \quad (26.38)$$

$\hat{V}(s_t[i]; \phi)$  is obtained by curve fitting (see Chapter 2) using samples  $(s_t[i], a_t[i]) \forall t, i$  with

$$y[i] = \sum_{t'=t}^T r(s_{t'}[i], a_{t'}[i]) \quad (26.39)$$

and

$$\hat{y}[i] = \hat{V}(s_t[i]; \phi) . \quad (26.40)$$

Eq.(26.39) is an approximation because  $(s_{t'}, a_{t'})_{t' > t}$  are averaged over in the exact expression for  $V(s_t)$ .  $\hat{V}(s_{t+1}[i]; \phi')$  is obtained in the same way as  $\hat{V}(s_t[i]; \phi)$  but with  $t$  replaced by  $t + 1$  and  $\phi$  by  $\phi'$ .



Figure 26.5: Q function learning RL bnet.

## Q function learning RL bnet

The Q-function learning RL bnet is given by Fig.26.5. This bnet is approximate and assumes that the policy is NOT known. The transition matrices for its nodes are given in blue below. (Remember that  $Q = v$ ).

$$P(s_t | s_{t-1}, a_{t-1}) = \text{given} \quad (26.41)$$

$$P(a_t | s_t, v_t(\cdot)) = \delta(a_t, \underset{a}{\operatorname{argmax}} v_t(s_t, a)) \quad (26.42)$$

$$P(r_t | s_t, a_t) = \delta(r_t, r(s_t, a_t)) \quad (26.43)$$

$r : S_{\underline{s}} \times S_{\underline{a}} \rightarrow \mathbb{R}$  is given.

$$\begin{aligned} \forall(s_t, a_t) : \quad & P(v_t(s_t, a_t) | v_{t-1}(\cdot)) = \\ & = \delta(v_t(s_t, a_t), r(s_t, a_t) + \gamma \max_a E_{\underline{s}_{t+1} | s_t, a_t} v_{t-1}(\underline{s}_{t+1}, a)) \end{aligned} \quad (26.44)$$

This value for  $v_t(s_t, a_t)$  approximates  $v_t = r_t + \gamma E_{\underline{s}_{t+1}, \underline{a}_{t+1}} v_{t+1}$ .

Some people use the bnet of Fig.26.6) instead of Fig.26.5 and replace Eq.(26.44) by

$$\begin{aligned} \forall(s_t, a_t) : \quad & P(v_t(s_t, a_t) | s_{t+1}, v_{t-1}(\cdot)) = \\ & = \delta(v_t(s_t, a_t), r(s_t, a_t) + \gamma \max_a v_{t-1}(s_{t+1}, a)) . \end{aligned} \quad (26.45)$$



Figure 26.6: Q function learning RL bnet. Same as Fig.26.5 but with new arrow passing  $s_t$  to  $Q_{t-1}$ .

# Chapter 27

## Reliability Box Diagrams and Fault Tree Diagrams

This chapter is based on Refs.[2] and [3].

In this chapter, we assume that reader is familiar with Boolean Algebra. See the Notational Conventions Chapter 0.2 for a quick review of what we recommend that you know about Boolean Algebra to fully appreciate this chapter.

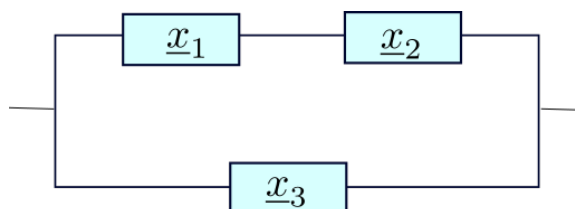


Figure 27.1: Example of rbox diagram.

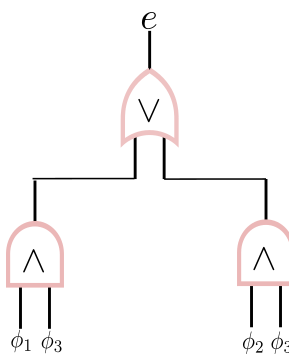


Figure 27.2: A fault tree diagram equivalent to Fig.27.1. It represents  $e = (\phi_1 \wedge \phi_3) \vee (\phi_2 \wedge \phi_3)$ .

Complicated devices with a large number of components such airplanes or rockets can fail in a many ways. If their performance depends on some components working in series and one of the components in the series fails, this might lead to catastrophic failure. To avert such disasters,





Figure 27.3: How to map an rbox diagram to a bnet.

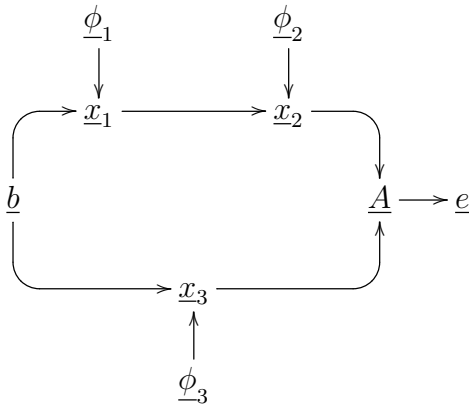


Figure 27.4: bnet corresponding to the rbox diagram Fig.27.1.

engineers use equivalent components connected in parallel instead of series, thus providing multiple backup systems. Another thing they do is to analyze the device to find its weak points and provide backup there and to estimate average time to failure.

The two most popular diagrams for analyzing failure modes and rates of large complicated devices are

- rbox diagrams = Reliability Box diagrams. See Fig.27.1 for an example.
- ftree diagrams = Fault Tree Diagrams. See Fig.27.2 for an example.

In an ftree, several nodes might stand for the same component of a physical device. In an rbox diagram, on the other hand, each node represents a distinct component in a device. Hence, rbox

diagrams resemble the device they are addressing whereas ftree diagrams don't. Henceforth, we will refer to this desirable property as **physical resemblance**.

As we will show below, it is pretty straightforward to translate an rbox to an ftree diagram. Going the other way, translating an ftree to an rbox diagram is much more difficult.

Next we will define a new kind of bnet that we will call a failure bnet that has physical resemblance. Then we will describe a simple method of translating (i.e., mapping) any rbox diagram to a failure bnet. Then we will show how a failure bnet can be used to do all the calculations that are normally done with an rbox or an ftree diagram. In that sense, failure bnets seem to afford all the benefits of both ftree and rbox diagrams.

A **failure bnet** contains nodes of 5 types: labeled  $\underline{b}$ ,  $\underline{e}$ ,  $\underline{x}_i$ ,  $\underline{\phi}_i$ , and  $\underline{A}_i$ . All nodes have only two possible states  $S = Success = 0$ ,  $F = Failure = 1$ .

1. The bnet has a beginning node labeled  $\underline{b}$  which is always set to success. The  $\underline{b}$  node and the  $\underline{\phi}_i$  nodes are the only root nodes of the bnet.
2. The bnet has a single leaf node, the end node, labeled  $\underline{e}$ .  $\underline{e}$  is fixed. In rbox diagrams,  $\underline{e} = S$  whereas in ftree diagrams,  $\underline{e} = F$ .
3.  $\underline{x}^{nx} = (\underline{x}_0, \underline{x}_1, \dots, \underline{x}_{nx-1})$ .  $\underline{x}_i \in \{S, F\}$  for all  $i$ .

Suppose  $\underline{x}_i$  has parents  $\underline{\phi}_i$  and  $\underline{a}^{na} = (\underline{a}_0, \underline{a}_1, \dots, \underline{a}_{na-1})$ . Then the transition prob matrix of node  $\underline{x}_i$  is defined to be

$$P(x_i | \phi_i, a^{na}) = \delta(x_i, \phi_i \vee \bigvee_{i=0}^{na-1} a_i) \quad (27.1)$$

4. For each node  $\underline{x}_i$ , the bnet has a "performance" root node  $\underline{\phi}_i \in \{0, 1\}$  with an arrow pointing from it to  $\underline{x}_i$  (i.e.,  $\underline{\phi}_i \rightarrow \underline{x}_i$ ). For all  $i$ ,

$$P(\phi_i) = \epsilon_i \delta(\phi_i, F) + \bar{\epsilon}_i \delta(\phi_i, S) . \quad (27.2)$$

$\epsilon_i$  is the failure probability and  $\bar{\epsilon}_i = 1 - \epsilon_i$  is the success probability. We name the failure probability  $\epsilon_i$  because it is normally very small. It is usually set to  $1 - e^{-\lambda_i t} \approx \lambda_i t$  when  $\lambda_i t \ll 1$ , where  $\lambda_i$  is the failure rate for node  $\underline{x}_i$  and  $t$  stands for time. The rblock literature usually calls  $\bar{\epsilon}_i = R_i$  the **reliability** of node  $\underline{x}_i$ , and  $\epsilon_i = (1 - R_i) = F_i$  its **unreliability**.

5. The nodes  $\underline{A}_i \in \{0, 1\}$  are simply AND gates. If  $\underline{A}_i$  has inputs  $\underline{y}^{ny} = (\underline{y}_0, \underline{y}_1, \dots, \underline{y}_{ny-1})$ , then the transition prob matrix of  $\underline{A}_i$  is

$$P(A_i | y^{ny}) = \delta(A_i, \bigwedge_{i=0}^{ny-1} y_i) . \quad (27.3)$$

An instance (instantiation) of a bnet is the bnet with all nodes set to a specific state. A **realizable instance (r-instance)** of a bnet is one which has non-zero probability.

Fig.27.3 shows how to translate any rbox diagram to a failure bnet. To illustrate this procedure, we translated the rbox diagram Fig.27.1 into the failure bnet Fig.27.4.

For the failure bnet Fig.27.4, one has:

$$\begin{aligned}
P(b) &= \mathbb{1}(b = 0) \\
P(x_1|\phi_1, b) &= \mathbb{1}(x_1 = \phi_1 \vee b) \\
P(x_2|\phi_2, x_1) &= \mathbb{1}(x_2 = \phi_2 \vee x_1) \\
P(x_3|\phi_3, b) &= \mathbb{1}(x_3 = \phi_3 \vee b) \\
P(A|x_2, x_3)e &= \mathbb{1}(x_2 \wedge x_3) \\
P(e|A) &= \mathbb{1}(e = A)
\end{aligned} \tag{27.4}$$

Therefore, all r-instances of this bnet must satisfy

$$e = (\phi_1 \vee \phi_2) \wedge \phi_3 \tag{27.5}$$

$$= (\phi_1 \wedge \phi_3) \vee (\phi_2 \wedge \phi_3) . \tag{27.6}$$

Eq.(27.6) proves that Fig.27.2 is indeed a representation of Fig.27.1.

Next, we consider r-instances of this bnet for two cases:  $e = S$  and  $e = F$ .

- **rblock analysis:**  $e = S = 0$ .

Table 27.1 shows the probability of all possible r-instances that end in success for the failure bnet Fig.27.4. (These r-instances are the main focus of rblock analysis). The first 4 of those probabilities (those with  $\phi_3 = 0$ ) sum to  $\bar{\epsilon}_3$  so the sum  $P(e = S)$  of all 5 is

$$P(e = S) = \bar{\epsilon}_3 + \bar{\epsilon}_1 \bar{\epsilon}_2 \epsilon_3 , \tag{27.7}$$

or, expressing it in reliability language in which  $\bar{\epsilon} = R$ ,

$$P(e = S) = R_3 + R_1 R_2 \bar{R}_3 . \tag{27.8}$$

- **ftree analysis:**  $e = F = 1$ .

Table 27.2 shows the probability of all possible r-instances that end in failure for the failure bnet Fig.27.4. (These r-instances are the main focus of ftree analysis). If we set  $\epsilon_i = \epsilon$  and  $\bar{\epsilon}_i \approx 1$  for  $i = 1, 2, 3$ , then the first two of those r-instances have probabilities of *order*( $\epsilon^2$ ) and the third has probability of *order*( $\epsilon^3$ ). The lowest order (*order*( $\epsilon^2$ )) r-instances are called the “minimal cut set” of the ftree. We will have more to say about minimal cut sets later. For now, just note from Eq.(27.6) and Fig.27.2, that that ftree is just the result of joining together with ORs an expression for each of the two elements of the cut set.

instance	probability
	$\bar{\epsilon}_1 \epsilon_2 \bar{\epsilon}_3$
	$\epsilon_1 \bar{\epsilon}_2 \bar{\epsilon}_3$
	$\epsilon_1 \epsilon_2 \bar{\epsilon}_3$
	$\bar{\epsilon}_1 \bar{\epsilon}_2 \bar{\epsilon}_3$
	$\bar{\epsilon}_1 \bar{\epsilon}_2 \epsilon_3$

Table 27.1: Probabilities of all possible r-instances with  $e = S = 0$  for failure bnet Fig.27.4.

instance	probability
	$\bar{\epsilon}_1 \epsilon_2 \epsilon_3$
	$\epsilon_1 \bar{\epsilon}_2 \epsilon_3$
	$\epsilon_1 \epsilon_2 \epsilon_3$

Table 27.2: Probabilities of all possible r-instances with  $e = F = 1$  for the failure bnet Fig.27.4.

---

**More general  $\underline{x}_i$ .**

Failure bnets can actually accommodate  $\underline{x}_i$  nodes of a more general kind that we first stipulated. Here are some possibilities:

For  $x^n \in \{0, 1\}^n$ , let

$$\text{len}(x^n) = \sum_i x_i \quad (27.9)$$

- **OR gate**

$$P(x_i|\phi_i, a^{na}) = \delta(x_i, \phi_i \vee \vee_j a_j) \quad (27.10)$$

$$= \delta(x_i, \phi_i \vee \mathbb{1}(\text{len}(a^{na}) > 0)) \quad (27.11)$$

- **AND gate**

$$P(x_i|\phi_i, a^{na}) = \delta(x_i, \phi_i \vee \wedge_j a_j) \quad (27.12)$$

$$= \delta(x_i, \phi_i \vee \mathbb{1}(\text{len}(a^{na}) = na)) \quad (27.13)$$

- **At least  $K$  failures (less than  $K$  successes) gate**

$$P(x_i|\phi_i, a^{na}) = \delta(x_i, \phi_i \vee \mathbb{1}(\text{len}(a^{na}) \geq K)) \quad (27.14)$$

- **Less than  $K$  failures (at least  $K$  successes) gate**

$$P(x_i|\phi_i, a^{na}) = \delta(x_i, \phi_i \vee \mathbb{1}(\text{len}(a^{na}) < K)) \quad (27.15)$$

- **Exactly one failure gate**

$$P(x_i|\phi_i, a^{na}) = \delta(x_i, \phi_i \vee \mathbb{1}(\text{len}(a^{na}) = 1)) \quad (27.16)$$

This equals an XOR (exclusive OR) gate when  $na = 2$ .

- **General gate**

$$f : \{0, 1\}^{na} \rightarrow \{0, 1\}$$

$$P(x_i|\phi_i, a^{na}) = \delta(x_i, \phi_i \vee f(a^{na})) \quad (27.17)$$


---

## Minimal Cut Sets

Suppose  $x \in \{0, 1\}$  and  $f : \{0, 1\} \rightarrow \{0, 1\}$ . Then by direct evaluation, we see that

$$f(x) = [\bar{x}f(0)] \vee [xf(1)] . \quad (27.18)$$

Let

$$\begin{aligned} !x &= 1 - x, \\ !^0x &= x, \\ !^1x &= !x \end{aligned} \quad (27.19)$$

Then Eq.27.18 can be rewritten as

$$f(x) = \vee_{a \in \{0,1\}} [(!^{\bar{a}}x)f(a)] . \quad (27.20)$$

Now suppose  $x^n \in \{0, 1\}^n$  and  $f : \{0, 1\}^n \rightarrow \{0, 1\}$ . Eq.(27.20) generalizes to

$$f(x^n) = \vee_{a^n \in \{0,1\}^n} [\prod_i (!^{\bar{a}_i}x_i)f(a^n)] . \quad (27.21)$$

Eq.(27.21) is called an ors-of-ands normal form expansion. There is also an ands-of-ors normal form expansion obtained by swapping multiplication and  $\vee$  in Eq.(27.21), but we won't need it here.

A **minimal cut set** is the smallest set of  $\phi_i$ 's such that if they are all  $F$ , then  $e = F$  for all the r-instances. From the failure bnet, we can always find a function  $f : \{0, 1\}^{nx} \rightarrow \{0, 1\}$  such that  $e = f(\phi^{nx})$  for all the r-instances. We did that for our example failure bnet and obtained Eq.(27.6). We can then express  $f(\phi^{nx})$  as an ors-of-ands expansion to find the set of minimal cut sets. The ands terms in that ors-of-ands expansion each gives a different minimal cut set, after some simplification. The ors-of-ands expression is not unique and it may have redundancies that need to be simplified (using the Boolean Algebra identities given in Chapter 0.2) to achieve minimality in the cut sets.

# Chapter 28

## Restricted Boltzmann Machines

In what follows, we will abbreviate "restricted Boltzmann machine" by rebo.

Let

$$v \in \{0, 1\}^{numv}$$

$$h \in \{0, 1\}^{numh}$$

$$b \in \mathbb{R}^{numv} \text{ (mnemonic, } v \text{ and } b \text{ sound the same)}$$

$$a \in \mathbb{R}^{numh}$$

$$W^{v|h} \in \mathbb{R}^{numv \times numh}$$

Energy:

$$E(v, h) = -(b^T v + a^T h + v^T W^{v|h} h) \quad (28.1)$$

Boltzmann distribution:

$$P(v, h) = \frac{e^{-E(v, h)}}{Z} \quad (28.2)$$

Partition function:

$$Z = \sum_{v, h} e^{-E(v, h)} = Z(a, b, W^{v|h}) \quad (28.3)$$

$$P(v|h) = \frac{e^{b^T v + a^T h + v^T W^{v|h} h}}{\sum_v e^{b^T v + a^T h + v^T W^{v|h} h}} \quad (28.4)$$

$$= \frac{e^{b^T v + v^T W^{v|h} h}}{\sum_v e^{b^T v + v^T W^{v|h} h}} \quad (28.5)$$

$$= \prod_i \frac{e^{v_i(b_i + \sum_j W_{i,j}^{v|h} h_j)}}{\sum_{v_i=0,1} e^{v_i(b_i + \sum_j W_{i,j}^{v|h} h_j)}} \quad (28.6)$$

$$= \prod_i P(v_i|h) \quad (28.7)$$

$$P(v_i|h) = \frac{e^{v_i(b_i + \sum_j W_{i,j}^{v|h} h_j)}}{Z_i(h)} \quad (28.8)$$





Figure 28.1: bnet for a Restricted Boltzmann Machine (rebo) with  $numv = 3$

Eq.28.8 implies that a rebo can be represented by the bnet Fig.28.1.

Let

$$x_i = b_i + \sum_j W_{ij}^{v|h} h_j . \quad (28.9)$$

Then

$$P(v_i = 1|h) = \frac{e^{x_i}}{1 + e^{x_i}} \quad (28.10)$$

$$= \frac{1}{1 + e^{-x_i}} \quad (28.11)$$

$$= \text{sig}(x_i) . \quad (28.12)$$

One could also expand the node  $\underline{h}$  in Fig.28.1 into  $numh$  nodes. But note that  $P(h) \neq \prod_j P(h_j)$  so there would be arrows among the  $h_j$  nodes.

Note that the rebo bnet is a special case of Naive Bayes (See Chapter 21) with  $v_i, h_i \in \{0, 1\}$  and specific  $P(h)$  and  $P(v_i|h)$  node matrices.

## Chapter 29

# Simpson's Paradox

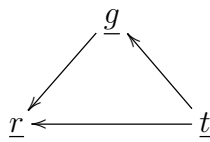


Figure 29.1: bnet for a simple case of Simpson's paradox.

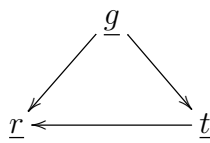


Figure 29.2: Equivalent to Fig.29.1

I wrote an article about this in 2020 for my blog “Quantum Bayesian Networks”. See Ref.[23].

# Chapter 30

## Turbo Codes

This chapter is based on Ref.[24].

In this chapter, vectors with  $n$  components will be indicated by an  $n$  superscript. For example,  $a^n = (a_0, a_1, \dots, a_{n-1})$ .

Consider an  $n$ -letter message  $u^n = (u_0, u_1, \dots, u_{n-1})$ , where for all  $i$ ,  $u_i \in \mathcal{A}$  is an element of an alphabet  $\mathcal{A}$ , and where for all  $i$ , the  $u_i$  are i.i.d.. Suppose  $u^n$  is encoded deterministically in two different ways,  $e_1(u^n)$  and  $e_2(u^n)$ . After passing through the same memoryless channel, the variables  $u^n, e_1, e_2$  become  $\tilde{u}^n, \tilde{e}_1, \tilde{e}_2$ , respectively. The letter  $u$  stands for unencoded, and  $e$  for encoded. Quantities with a tilde  $\tilde{u}^n, \tilde{e}_1, \tilde{e}_2$  occur after channel passage and are visible (measurable). Quantities without a tilde  $u^n, e_1, e_2$  are hidden (unmeasurable).

The situation just described can be represented by the bnet Fig.30.1, or by its abridged version Fig.30.2. But note that the abridged version does not show explicitly that the  $u_i$  are i.i.d. or that the channel is memoryless (i.e., that the  $u_i$  for all  $i$  pass independently through the channel).

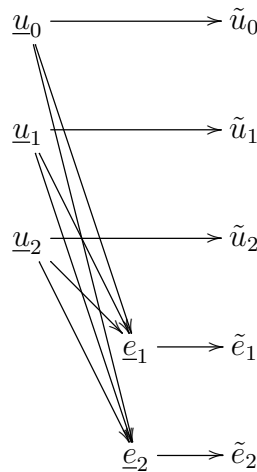


Figure 30.1: Turbo coding B net representing a message being encoded two different ways and then the original message and the 2 encodings pass through a memoryless channel.



Figure 30.2: Abridged version of Fig.30.1.

Define

$$x = (u^n, e_1, e_2) \quad (30.1)$$

and

$$\tilde{x} = (\tilde{u}^n, \tilde{e}_1, \tilde{e}_2) . \quad (30.2)$$

Fig.30.1 implies that

$$P(x, \tilde{x}) = P(\tilde{u}^n | u^n) \left[ \prod_{r=1,2} P(\tilde{e}_r | e_r) P(e_r | u^n) \right] P(u^n) . \quad (30.3)$$

Because the  $u^n$  are i.i.d.,

$$P(u^n) = \prod_i P(u_i) . \quad (30.4)$$

Because the channel is memoryless,

$$P(\tilde{u}^n | u^n) = \prod_i P(\tilde{u}_i | u_i) . \quad (30.5)$$

Because the encoding is deterministic, we must have for  $r = 1, 2$

$$P(e_r | u^n) = \delta(e_r, e_r(u^n)) . \quad (30.6)$$

Define the belief functions

$$BEL_i = BEL_i(\underline{u}_i = a) = P(\underline{u}_i = a | \tilde{x}) . \quad (30.7)$$

The best estimate of  $u_j$  given all visible evidence  $\tilde{x}$  is

$$\hat{u}_i = \operatorname{argmax}_{u_i} BEL_i(u_i) . \quad (30.8)$$

Define the probability functions

$$\pi_i = \pi_i(u_i) = P(u_i) , \quad (30.9)$$

and the likelihood functions

$$\lambda_i = \lambda_i(u_i) = P(\tilde{u}_i|u_i) . \quad (30.10)$$

For  $r = 1, 2$ , define the Kernel functions

$$K_r = K_r(u^n) = P(\tilde{e}_r|e_r = e_r(u^n)) . \quad (30.11)$$

In this book,  $\mathcal{N}(!a)$  denotes a normalization constant that does not depend on  $a$ . Define

$$\mathcal{N}_i = \mathcal{N}(!u_i) . \quad (30.12)$$

**Claim 4**

$$BEL_i = \mathcal{N}_i \lambda_i \pi_i \mathcal{T}_i^{K_1 K_2} \left[ \prod_{j \neq i} \lambda_j \pi_j \right] , \quad (30.13)$$

where  $\mathcal{T}_i^K(\cdot)$  with  $K = K_1 K_2$  is an operator (transform) that acts on functions of  $u^n$ :

$$\mathcal{T}_i^K(\cdot) = \sum_{u^n} \delta(u_i, a) K(u^n)(\cdot) . \quad (30.14)$$

**proof:**

$$\begin{aligned} P(\underline{u}_i = a | \tilde{x}) &= \\ &= \sum_x \delta(u_i, a) P(x | \tilde{x}) \end{aligned} \quad (30.15)$$

$$= \sum_x \delta(u_i, a) \frac{P(\tilde{x} | x) P(x)}{P(\tilde{x})} \quad (30.16)$$

$$= \mathcal{N}(!a) \sum_x \delta(u_i, a) P(\tilde{x} | x) P(x) \quad (30.17)$$

$$= \mathcal{N}(!a) \sum_x \delta(u_i, a) P(u^n) \left[ \prod_{r=1,2} P(\tilde{e}_r | e_r) \delta(e_r, e_r(u^n)) \right] \prod_j P(\tilde{u}_j | u_j) \quad (30.18)$$

$$= \mathcal{N}(!a) \lambda_i(a) \pi_i(a) R , \quad (30.19)$$

where

$$R = \sum_{u^n} \delta(u_i, a) \left[ \prod_{r=1,2} P(\tilde{e}_r | e_r(u^n)) \right] \prod_{j \neq i} P(\tilde{u}_j | u_j) P(u_j) \quad (30.20)$$

$$= \sum_{u^n} \delta(u_i, a) \left[ \prod_{r=1,2} K_r(u^n) \right] \prod_{j \neq i} \lambda_j(u_j) \pi_j(u_j) \quad (30.21)$$

$$= \mathcal{T}_i^{K_1 K_2} \left[ \prod_{j \neq i} \lambda_j(u_j) \pi_j(u_j) \right]. \quad (30.22)$$

Hence

$$BEL_i(a) = \mathcal{N}(!a) \lambda_i(a) \pi_i(a) \mathcal{T}_i^{K_1 K_2} \left[ \prod_{j \neq i} \lambda_j(u_j) \pi_j(u_j) \right]. \quad (30.23)$$

**QED**

## Decoding Algorithm

The Turbo algorithm for decoding the encode message is as follows. For  $m = 0$ , let

$$\pi_j^{(0)}(u_j) = \frac{1}{n_{\underline{u}_j}}. \quad (30.24)$$

Then for  $m = 1, 2, \dots$ , let

$$\pi_i^{(m)} = \mathcal{N}_i \mathcal{T}_i^{K_{m \% 2}} \left[ \prod_{j \neq i} \lambda_j \pi_j^{(m-1)} \right], \quad (30.25)$$

where  $m \% 2 = 1$  if  $m$  is odd and  $m \% 2 = 2$  if  $m$  is even. Furthermore, for  $m > 0$ , let

$$BEL_i^{(m)} = \mathcal{N}_i \lambda_i \pi_i^{(m-1)} \pi_i^{(m)} \quad (30.26)$$

$$= \mathcal{N}_i \lambda_i \pi_i^{(m-1)} \mathcal{T}_i^{K_{m \% 2}} \left[ \prod_{j \neq i} \lambda_j \pi_j^{(m-1)} \right]. \quad (30.27)$$

As  $m \rightarrow \infty$ ,  $BEL_i^{(m)}$  given by Eq.(30.27) is expected to converge to the the exact  $BEL_i$  given by Eq.(30.13).

Turbo decoding can be represented by the bnets Figs.30.3 and 30.4.

The node transition matrices, printed in blue, for Fig.30.3, are given by:

$$P(d_i^{(m)} = a \mid \tilde{u}^n, \tilde{e}_{m \% 2}) = BEL_i^{(m)}(a). \quad (30.28)$$



Figure 30.3: B net describing Turbo code generation of  $BEL_i^{(m)}(a)$  for  $m = 1, 2, \dots$



Figure 30.4: B net describing Turbo code generation of  $BEL^{n(m)}(\cdot)$  and  $\pi^{n(m)}(\cdot)$  for  $m = 0, 1, 2, \dots$ . The following arrows were not drawn so as not to unduly clutter the diagram: Arrows pointing from node  $\underline{\lambda}^n(\cdot)$  to nodes  $\underline{\pi}^{n(m)}(\cdot)$  and  $\underline{BEL}^{n(m)}(\cdot)$  for  $m = 0, 1, 2, \dots$

The node transition matrices, printed in blue, for Fig.30.4, are given by:

$$P((\lambda^n)'(\cdot)|\tilde{u}^n) = \delta((\lambda^n)'(\cdot), \lambda^n(\cdot)) \quad (30.29)$$

$$P(\pi^{n(m)}(\cdot)|\lambda^n(\cdot), \pi^{n(m-1)}(\cdot), \tilde{e}_{m\%2}) = \prod_i \prod_{u_i} \delta(\pi_i^{(m)}(u_i), \mathcal{N}_i \mathcal{T}_i^{K_{m\%2}} [\prod_{j \neq i} \lambda_j \pi_j^{(m-1)}]) \quad (30.30)$$

$$P(BEL^{n(m)}(\cdot)|\lambda^n(\cdot), \pi^{n(m)}(\cdot), \pi^{n(m-1)}(\cdot)) = \prod_i \prod_{u_i} \delta(BEL_i(u_i), \mathcal{N}_i \lambda_i \pi_i^{(m-1)} \pi_i^{(m)}) \quad (30.31)$$

## Message Passing Interpretation of Decoding Algorithm

Ref.[24] shows that the Turbo code decoding algo can be interpreted as an application of Message Passing. We leave all talk of Message Passing to a separate chapter, Chapter 19.



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