

# BAYESUVIUS QUANTICO

A VISUAL DICTIONARY OF  
QUANTUM BAYESIAN NETWORKS



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# **Bayesuvius Quantico,** a visual dictionary of Quantum Bayesian Networks

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# **Bayesuvius Quantico**

by Robert R. Tucci

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# Appendices

# Appendix A

## Notational Conventions and Preliminaries

This book is a sequel to my book entitled “Bayesuvius” (see Ref.[3]). For consistency, I have tried to follow in this book the same notational conventions used in the prior book. If any notation is not defined in this book, check in the prior book. It might be defined there.

### A.1 Set notation

The number of elements in any set  $S$  is denoted by  $|S|$ .

$\mathbb{Z}$  = integers

$\mathbb{Z}_{>0}$  = positive integers

$\mathbb{Z}_{[a,b]} = a, a+1, \dots, b$  for some integers  $a, b$  such that  $a \leq b$

$\mathbb{R}$  = reals

$\mathbb{C}$  = complex numbers

$\mathbb{C}^{n \times m}$  =  $n \times m$  matrices of complex numbers

### A.2 Group

A **group**  $\mathcal{G}$  is a set of elements with a multiplication map  $\mathcal{G} \times \mathcal{G} \rightarrow \mathcal{G}$  such that

1. the multiplication is **associative** ; i.e.,

$$(ab)c = a(bc) \tag{A.1}$$

for  $a, b, c \in \mathcal{G}$ .

2. there exists an **identity element**  $e \in \mathcal{G}$  such that

$$ea = ae = a \tag{A.2}$$



for all  $a \in \mathcal{G}$

3. for any  $g \in \mathcal{G}$ , there exists an **inverse**  $a^{-1} \in \mathcal{G}$  such that

$$aa^{-1} = a^{-1}a = e \quad (\text{A.3})$$

$|\mathcal{G}|$  (i.e., number of elements in  $\mathcal{G}$ ) is called the **order** of the group.

If multiplication is **commutative** (i.e.,  $ab = ba$  for all  $a, b \in \mathcal{G}$ ), the group is said to be **abelian**.

A **subgroup**  $\mathcal{H}$  of  $\mathcal{G}$  is a subset of  $\mathcal{G}$  ( $\mathcal{H} \subset \mathcal{G}$ ) which is also a group. It's easy to show that any  $\mathcal{H} \subset \mathcal{G}$  is a group if it contains the identity and is **closed under multiplication** (i.e.,  $ab \in \mathcal{H}$  for all  $a, b \in \mathcal{H}$ )

## A.3 Group Representation

A **group representation** (rep) of a group  $\mathcal{G}$  is a map  $\phi : \mathcal{G} \rightarrow \mathbb{C}^{n \times n}$  such that

$$\phi(a)\phi(b) = \phi(ab), \quad \phi(e) = I \quad (\text{A.4})$$

where  $e$  is the identity of the group and  $I$  is the identity matrix. Such a map is called a **homomorphism** (because it preserves an operation). The map  $\phi$  partitions  $\mathcal{G}$  into disjoint subsets (equivalence classes), such that all elements of  $\mathcal{G}$  in each disjoint subset are represented by the same matrix.

One way to specify a representation is to give the effect of each group element  $a \in \mathcal{G}$  on a basis of vectors  $\{|1\rangle, |2\rangle, \dots, |n\rangle\}$ .

$$\phi(a)|i\rangle = \sum_j M_{ij}|j\rangle \implies \langle i|\phi(a)|j\rangle = M_{ij} \quad (\text{A.5})$$

If the map  $\phi$  is 1-1, onto, we call it a **faithful representation**

The **trivial representation** represents all  $g \in \mathcal{G}$  by  $\text{diag}(1, 1, \dots, 1)$ . It's dimension is  $d_\lambda = 0$ . It's not a faithful rep.

A **singlet representation** represents all  $g \in \mathcal{G}$  by  $z(g)\text{diag}(1, 1, \dots, 1)$  for some  $z(g) \in \mathbb{C}$ . For the singlet rep,  $\langle i|\phi(a)|j\rangle = z(g)\delta(i, j)$ . It's dimension is  $d_\lambda = 1$ . The projection operator  $\delta_a^b \delta_c^d$  when acting on  $G_c^d$  gives a  $z(G)\text{diag}(1, 1, \dots, 1)$  where  $z(G) = \text{tr}(G)$ , so it projects to a singlet rep.

When a group is defined using matrices, those matrices are called the **defining representation** (defrep). For example, the group of **General Linear Transformations** is defined by

$$GL(n, \mathbb{C}) = \{M \in \mathbb{C}^{n \times n} : \det M \neq 0\} \quad (\text{A.6})$$

---

<sup>1</sup>More generally, the  $\mathbb{C}^{n \times n}$  can be replaced by  $\mathbb{R}^{n \times n}$  or by  $\mathbb{F}^{n \times n}$  for any field  $\mathbb{F}$

The **adjoint representation** (adjrep) is defined in terms of the structure constants of the Lie Algebra. If the Lie Algebra satisfies  $[T^i, T^j] = if_{ijk}T^k$ , then the adjrep is given by the matrices with  $i, j$  entries  $M_{ij}^k = -if_{ij}^k$ . Write  $|T^i\rangle$  instead of  $T^i$  and  $|x\rangle = x_i|T^i\rangle$ . Then

$$[|x\rangle, \cdot]|T^j\rangle = if_{ijk}|T^k\rangle \implies \langle T^k|[|x\rangle, \cdot]|T^j\rangle = ix_if_{ijk} \quad (\text{A.7})$$

**Irreducible representations** (irreps) are defined in Ch. 13

The **fundamental representation** (funrep) is defined as the smallest irrep.. The defrep equals the funrep for  $SU(n), SO(n), SP(n)$ , but not for  $E_8$ .

## A.4 Group Theory References

Much of this book deals with Group Theory (GT).

GT is a vast subject. Who would have thought that the simple definition of a group would generate so many elegant, highly applicable and useful results and consequences.

GT books by mathematicians are very different from GT books by physicists, even though, of course, they agree on the definitions. Mathematicians are, as to be expected, more rigorous and abstract. But it goes much further than that. Physicists are much more interested in applications to physical systems, especially Quantum Mechanics (QM). Soon after QM was invented, it was realized that Linear Algebra (LA) and GT (especially Group Representation Theory, which combines GT and LA) are extremely relevant and useful in QM. Hermann Weyl, Eugene Wigner, Hans Bethe, Linus Pauling, etc. combined QM and GT to understand the spectra and chemistry of atoms and molecules, and later GT was heavily used in Quantum Field Theory and Particle Physics to devise the Standard Model. Condensed Matter physicists have also used it to understand crystalline solids and to devise quasi particles that can be detected in the lab.

My PhD is in physics so in this book I cover GT topics that are mainly of interests to physicists and engineers. Furthermore, I am nowhere as abstract and rigorous as mathematicians usually are.

My favorite books about GT for physicists are the Elliott & Dawber's (ED) 2 volume series Ref. [2] and Predrag Cvitanovic's Birdtracks book Ref.[1]. I highly recommend both of these references. I think both of them are excellent.

The Birdtracks book explains key concepts in GT representation theory using network diagrams (Cvitanovic calls such diagrams birdtracks) whereas the ED book doesn't use that type of diagram. Many people don't use birdtracks either, they only use algebra. But since this is a book about visualization using network diagrams (quantum bnets), we use birdtracks. In fact, many of the chapters in this book were heavily influenced by Ref.[1] by Cvitanovic. I hope he doesn't mind. I really love his book.

## A.5 Vector Space and Algebra over a field $\mathbb{F}$

A vector space (a.k.a. linear space)  $\mathcal{V}$  is defined as a set endowed with two operations: vector addition  $+: \mathcal{V} \times \mathcal{V} \rightarrow \mathcal{V}$ , and scalar multiplication  $\mathbb{F} \times \mathcal{V} \rightarrow \mathcal{V}$ , such that

- $\mathcal{V}$  is an abelian group under  $+$  with identity  $0$  and inverse of  $x \in \mathcal{V}$  equal to  $-x \in \mathcal{V}$
- For  $\alpha, \beta \in \mathbb{F}$  and  $x, y \in \mathcal{V}$

$$\alpha(x + y) = \alpha x + \alpha y \quad (\text{A.8})$$

$$(\alpha + \beta)x = \alpha x + \beta x \quad (\text{A.9})$$

$$\alpha(\beta x) = (\alpha\beta)x \quad (\text{A.10})$$

$$1x = x \quad (\text{A.11})$$

$$0x = 0 \quad (\text{A.12})$$

In this book, we will always use either  $\mathbb{C}$  or  $\mathbb{R}$  for  $\mathbb{F}$ . Both of these fields are infinite but some fields are finite.

An **algebra**  $\mathcal{A}$  is a vector space which, besides being endowed with vector addition and scalar multiplication as all vector spaces are, it has a **bilinear vector product**. A bilinear vector product is a product that is linear on both sides; i.e.,

$$(\alpha x + \beta y) \cdot z = \alpha x \cdot z + \beta y \cdot z \quad (\text{A.13})$$

and

$$z \cdot (\alpha x + \beta y) = \alpha z \cdot x + \beta z \cdot y \quad (\text{A.14})$$

for  $x, y, z \in \mathcal{A}$  and  $\alpha, \beta \in \mathbb{C}$ . The cross product (but not the dot product) for vectors in  $\mathbb{R}^3$ , the multiplication of 2 complex numbers, the product or commutator of 2 square matrices, are all good examples of bilinear vector products.

Let  $B = \{\tau_i : i = 1, 2, \dots, r\}$  be a basis for the vector space  $\mathcal{A}$ . Then note that  $B$  is closed under vector multiplication.

$$\tau_i \cdot \tau_j = \sum_k c_{ij}^k \tau_k \quad (\text{A.15})$$

where  $c_{ij}^k \in \mathbb{C}$ . The  $c_{ij}^k$  are called **structure constants** of  $B$ .

An **associative algebra** satisfies  $(x \cdot y) \cdot z = x \cdot (y \cdot z)$  for  $x, y, z \in \mathcal{A}$ .

- Not associative: cross product for vectors in  $\mathbb{R}^3$ .
- Associative: the product or commutator of 2 square matrices and the product of complex numbers

## A.6 Tensors

Let

$$(x_a) = (x_1, x_2, \dots, x_n) = x^{\cdot n} \in V^n = \mathbb{C}^{n \times 1}$$

$$\text{Reverse of vector } rev(x_1, x_2, \dots, x_n) = (x_n, x_{n-1}, \dots, x_1)$$

$$y^b = \sum_a g^{ba} x_a$$

$(y^b) = (y^1, y^2, \dots, y^n) = y^{\dagger \cdot n} \in V^{\dagger n} = \mathbb{C}^{n \times 1}$ .  $V^n$  is the lower indices vector space and  $V^{\dagger n}$  is its **dual vector space** (i.e., with upper indices).

$$M_a^{\cdot b} \in \mathbb{C}^{n \times n}, a, b \in \mathbb{Z}_{[1, n]}$$

Implicit Summation Convention

$$M_a^{\cdot b} x_b = \sum_{b=1}^n M_a^{\cdot b} x_b \quad (\text{A.16})$$

If the Hermitian conjugate  $\dagger$  equals  $*T$  where  $*$  is complex conjugation and  $T$  is transpose, then define

$$(M^\dagger)_b^{\cdot a} = (M_a^{\cdot b})^*, \quad (M^T)_b^{\cdot a} = M_a^{\cdot b}, \quad (\text{A.17})$$

Thus,  $\dagger$  and  $T$  do two things: (1) reverse the horizontal order of the indices (2) reverse vertical positions of the indices; i.e., lower upper indices and raise lower indices. Hermitian conjugation also complex conjugates the tensor components.

If  $M$  is a Hermitian matrix (i.e.,  $M^\dagger = M$ ),

$$M_b^{\cdot a} = (M_a^{\cdot b})^* \quad (\text{A.18})$$

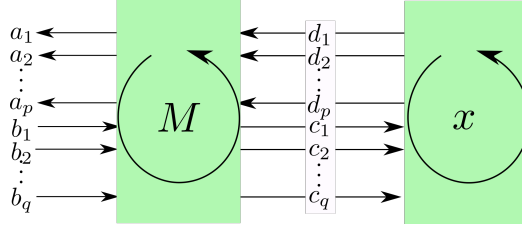


Figure A.1: Index labels for  $Mx$  where  $M \in \mathbb{C}^{n^{p+q} \times n^{p+q}}$  and  $x \in V^{n^p} \otimes V^{\dagger n^q}$ . Note that we list indices in counterclockwise (CC) direction, starting at the top.

Suppose  $a_i, b_i, c_i, d_i \in \mathbb{Z}_{[1, n]}$ . From Fig.A.1

$$y_{a^{\cdot p}}^{\cdot b^{\cdot q}} = M_{a^{\cdot p}}^{\cdot b^{\cdot q}} \overset{rev(d^{\cdot p})}{rev(c^{\cdot q})} x_{d^{\cdot p}}^{\cdot c^{\cdot q}} \quad (\text{A.19})$$

If we define  $X_\alpha$  and  $x^\alpha$  by

$$X_\alpha = X_{a^{\cdot p}}^{\cdot b^{\cdot q}}, \quad X^\alpha = X_{rev(b^{\cdot q})}^{\cdot rev(a^{\cdot p})} \quad (\text{A.20})$$

then

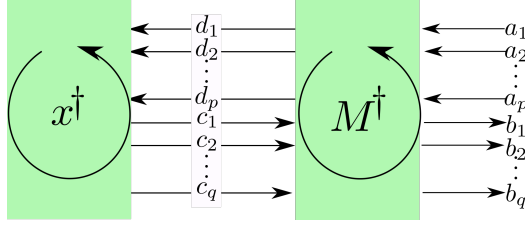


Figure A.2: Index labels for  $x^\dagger M^\dagger$  corresponding to Fig.A.2. Note that we list indices in counterclockwise (CC) direction, starting at the top.

$$x_\alpha = M_\alpha^\beta x_\beta \quad (\text{A.21})$$

---

Hermitian conjugation (see Fig.A.2)

$$\begin{cases} (M^\dagger)_a^d = (M_d^a)^* \\ (M^\dagger)_\alpha^\delta = (M_{rev(\delta)}^{rev(\alpha)})^* \end{cases} \quad (\text{A.22})$$

Note that  $\dagger$  does 3 things to the birdtrack:

1. It flips the horizontal axis of the figure. (In the algebraic expression of the tensor, this corresponds to reversing the horizontal order of the indices.)
2. For each node, it changes incoming arrows to outgoing ones and vice versa. (In the algebraic expression of the tensor, this corresponds reversing the vertical positions of the indices; i.e., lowering upper indices and raising lower ones.)
3. It replaces the tensor component by its complex conjugate

Hermitian matrix

$$M^\dagger = M, \quad \begin{cases} (M_d^a)^* = M_a^d \\ (M_{rev(\delta)}^{rev(\alpha)})^* = M_\alpha^\delta \end{cases} \quad (\text{A.23})$$

Unitary matrix

$$M^\dagger M = 1, \quad \begin{cases} (M_d^a)^* M_a^d = 1 \\ (M_{rev(\delta)}^{rev(\alpha)})^* M_\alpha^\delta = 1 \end{cases} \quad (\text{A.24})$$

---

Note that for  $x \in V^n$ ,  $y \in V^{\dagger n}$ , and  $G \in \mathcal{G} \subset GL(n, \mathbb{C})$ ,

$$(x')_a (y')^b = G_c^b G_a^d x_d y^c \quad (\text{A.25})$$

If  $x \in V^{n^p} \otimes V^{\dagger n^q}$ ,  $\mathbb{G} \in \mathcal{G} \subset GL(n^{p+q}, \mathbb{C})$ ,

$$(x')_{a:p}^{b:q} = \mathbb{G}_{a:p}^{b:q}{}_{rev(c:q)}^{rev(d:p)} x_{d:p}^{c:q}, \quad (x'_\alpha = \mathbb{G}_\alpha^\beta x_\beta) \quad (\text{A.26})$$

where we define

$$\mathbb{G}_{a:p}^{b:q}{}_{rev(c:q)}{}^{rev(d:p)} \stackrel{\text{def}}{=} \prod_{i=1}^p G_{a_i}^{d_i} \prod_{i=1}^q G^{\dagger b_i}_{c_i} \quad (\text{A.27})$$

---

An issue that arises with tensors is this: When is it permissible to represent a tensor by  $M_{ab}^{cd}$ ? If we define  $M_{ab}^{cd}$  by

$$M_{ab}^{cd} = M_{ab}{}^{cd} \quad (\text{A.28})$$

then it's always permissible. Then one can define tensors like  $M_a{}^{bcd}$  as

$$M_a{}^{bcd} = g^{bb'} M_{ab'}{}^{cd} = g^{bb'} M_{ab'}^{cd} \quad (\text{A.29})$$

Hence, one drawback of using the notation  $M_{ab}^{cd}$  is that if one is interested in using versions of  $M_{ab}^{cd}$  with some indices raised or lowered, one has to write down explicitly the metric tensors that do the lowering and raising. Instead of writing  $M_a{}^{bcd}$ , you'll have to write  $g^{bb'} M_{ab'}^{cd}$ . This is not very onerous when explaining a topic in which not much lowering and raising of indices is done. But in topics like General Relativity that do use a lot of raising and lowering of indices, it might not be too succinct.

## A.7 Permutations

Some well known notation and results about permutations are these.

(1, 2) stands for a **transposition**; i.e., a map that swaps 1 and 2:

$$\left( \begin{array}{cccc} 1 & 2 & 3 & \dots & p \\ & \searrow & \downarrow & & \downarrow \\ & 1 & 2 & \dots & p \end{array} \right) \quad (\text{A.30})$$

(3, 2, 1) stands for a **permutation**; i.e., a map that maps  $3 \rightarrow 2 \rightarrow 1 \rightarrow 3$ .

$$\left( \begin{array}{cccc} 1 & 2 & 3 & 4 & \dots & p \\ & \searrow & \swarrow & \downarrow & & \downarrow \\ & 1 & 2 & 3 & \dots & p \end{array} \right) \quad (\text{A.31})$$

Any reordering of  $(1, 2, 3, \dots, p)$  is a permutation of  $p$  letters (or numbers or elements).

The set  $S_p$  of all permutation of  $p$  letters is called the **symmetric group in  $p$  letters**. It has  $p!$  elements (i.e.,  $|S_p| = p!$ ) and is a group, where the group's product is map composition and the group's identity element is the identity map.

Any permutation can be expressed as a product of transpositions. For example,  $(3, 2, 1) = (3, 2)(2, 1)$ .

An **even permutation** such as  $(3, 2, 1)$  can be expressed as a product of an even number of transpositions. An **odd permutation** can be expressed as a product of an odd number of transpositions.

# Appendix B

## Birdtracks

This chapter is based on Cvitanovic Birdtracks book Ref. [1] and my paper Ref. [4]

The tensor notation discussed in Sec.A.6 is succinct and straightforward, but it's not visually illuminating. The birdtrack notation that we shall discuss in this chapter, is not as succinct as the tensor notation, and can lead to sign errors if you are careless, but it is very visually illuminating. Thus, the tensor and birdtrack notations complement each other well. We will often display results using both, side by side.

### B.1 Classical Bayesian Networks and their Instantiations

Classical Bayesian Networks (bnets) are discussed exhaustively in the first book of this series, Ref.[3]. This is a brief section to remind the reader of how they are defined.

Let PD stand for probability distribution.

We call  $P_{\underline{y}|\underline{x}} : val(\underline{y}) \times val(\underline{x}) \rightarrow [0, 1]$  a **Transition Probability Matrix** (TPM)<sup>1</sup> if

$$\sum_{y \in val(\underline{y})} P_{\underline{y}|\underline{x}}(y|x) = 1 \quad (\text{B.1})$$

In other words, a TPM is a conditional PD. A TPM of the form

$$P(y|x) = \delta(y, f(x)) \quad (\text{B.2})$$

for some function  $f : val(\underline{x}) \rightarrow val(\underline{y})$  is said to be **deterministic**.

A bnet is a **Directed Acyclic Graph** (DAG) with the nodes labelled by random variables<sup>2</sup>. Each bnet stands for a full PD of the node random variables expressed as a product of a TPM for each node. For example, the bnet

---

<sup>1</sup>A TPM is also known as a Conditional Probability Table (CPT).

<sup>2</sup>As in the first volume of this series, we indicate random variables by underlined letters

$$\mathcal{C} = \begin{array}{ccc} & \underline{b} & \\ \swarrow & & \searrow \\ \underline{c} & \longleftarrow & \underline{a} \end{array} \quad (\text{B.3})$$

stands for the full PD

$$P(a, b, c) = P(c|b, a)P(b|a)P(a) \quad (\text{B.4})$$

Bnets do not have free indices because their nodes are labelled by random variables. It is convenient to draw the DAG for a bnet but with the underlining removed from the random variables, and then to assign a numerical value to this new DAG. The resultant DAG now has free indices. We call it an **instantiation of the bnet**. For example, from the bnet  $\mathcal{C}$  of Eq.(B.3), we get the instantiation<sup>3</sup>

$$P(a, b, c) = P(c|b, a)P(b|a)P(a) = \begin{array}{ccc} & b & \\ \swarrow & & \searrow \\ c & \longleftarrow & a \end{array} P(a) \quad (\text{B.5})$$

Let  $a^{i2} = (a_1, a_2)$ . Based on the bnet  $\mathcal{C}$  of Eq.(B.3), define a new bnet  $\mathcal{C}'$  as follows

$$\mathcal{C}' = \begin{array}{ccc} & b & \\ \swarrow & & \searrow \\ \underline{c} & \longleftarrow \underline{a_2} & \underline{a^{i2}} \end{array} \quad (\text{B.6})$$

$\mathcal{C}'$  represents the the full PD

$$P(a^{i2}, b, c) = P(c|b, a_2)P(a_2|a^{i2})P(b|a_1)P(a_1|a^{i2})P(a^{i2}) \quad (\text{B.7})$$

The 2 new nodes  $\underline{a_1}$  and  $\underline{a_2}$  of bnet  $\mathcal{C}'$  are called **marginalizer nodes**. We assign to them the following TPMs (printed in blue):

$$P[a'_i|\underline{a^{i2}} = (a_1, a_2)] = \delta(a'_i, a_i) \quad (\text{B.8})$$

for  $i = 1, 2$ . We can also define an instantiation of  $\mathcal{C}'$  as follows:

$$P'(a^{i2}, b, c) = \begin{array}{ccc} & b & \\ \swarrow & & \searrow \\ c & \longleftarrow \underline{a_2} & \underline{a^{i2}} \end{array} P(a^{i2}) \quad (\text{B.9})$$

---

<sup>3</sup>Note that we don't include the root node probabilities as part of the graph value. Thus,  

$$P(a, b) = \underbrace{b \rightarrow a}_{P(b|a)} P(a)$$



## B.2 Quantum Bayesian Networks and their Instantiations

As far as I know, Quantum Bayesian Networks (qbnet) were invented by me in Ref.[4].

qbnet are closely analogous to classical bnet, but the TPM are replaced by **Transition Amplitude Matrices (TAM)**.

Let PA stand for probability amplitude.

We call  $A_{y|x} : val(\underline{y}) \times val(\underline{x}) \rightarrow \mathbb{C}$  a TAM if

$$\sum_{y \in val(\underline{y})} |A(y|x)|^2 = 1 \quad (\text{B.10})$$

Note that if  $A$  is the matrix with entries  $\langle y|A|x \rangle = A(y|x)$ , then

$$\langle y|A^\dagger A|x \rangle = \sum_{y \in val(\underline{y})} |A(y|x)|^2 = 1 \quad (\text{B.11})$$

If  $A$  is a unitary matrix, then  $A^\dagger A = AA^\dagger = 1$  so “half” ( $A^\dagger A = 1$ ) of the definition of unitary matrix is satisfied by a TAM. If both parts were satisfied,  $A$  would have to be a square matrix.

A qbnet is a DAG with the nodes labelled by random variables. Each qbnet stands for a full PA of the node random variables expressed as a product of a TAM for each node. For example, the qbnet

$$\mathcal{Q} = \begin{array}{ccc} & b & \\ \swarrow & & \searrow \\ c & \longleftarrow & a \end{array} \quad (\text{B.12})$$

stands for the full PA

$$A(a, b, c) = A(c|b, a)A(b|a)A(a) \quad (\text{B.13})$$

Qbnet do not have free indices because their nodes are labelled by random variables. It is convenient to draw the DAG for a qbnet but with the underlining removed from the random variables, and then to assign a numerical value to this new DAG. The resultant DAG now has free indices. We call it an **instantiation of the qbnet**. For example, from the bnet  $\mathcal{Q}$  of Eq.(B.12), we get the instantiation

$$A(a, b, c) = A(c|b, a)A(b|a)A(a) = \begin{array}{ccc} & b & \\ \swarrow & & \searrow \\ c & \longleftarrow & a \end{array} A(a) \quad (\text{B.14})$$

Let  $a^{:2} = (a_1, a_2)$ . Based on the qbnet  $\mathcal{Q}$  of Eq.(B.12), define a new qbnet  $\mathcal{Q}'$  as follows

$$\mathcal{Q}' = \begin{array}{c} b \\ \swarrow \quad \nwarrow \\ \underline{c} \quad \underline{a}^{:2} \\ \leftarrow a_2 \end{array} \quad (B.15)$$

$\mathcal{Q}'$  represents the the full PA

$$A(a^{:2}, b, c) = A(c|b, a_2)A(a_2|a^{:2})A(b|a_1)A(a_1|a^{:2})A(a^{:2}) \quad (B.16)$$

The 2 new nodes  $\underline{a}_1$  and  $\underline{a}_2$  of qbnet  $\mathcal{Q}'$  are called **marginalizer nodes**. We assign to them the following TAMs (printed in blue):

$$A[a'_i|\underline{a}^{:2} = (a_1, a_2)] = \delta(a'_i, a_i) \quad (B.17)$$

for  $i = 1, 2$ . We can also define an instantiation of  $\mathcal{Q}'$  as follows:

$$A(a^{:2}, b, c) = \begin{array}{c} b \\ \swarrow \quad \nwarrow \\ \underline{c} \quad \underline{a}^{:2} \\ \leftarrow a_2 \end{array} A(a^{:2}) \quad (B.18)$$

### B.3 Birdtracks

Tensors written in **algebraic notation** such as  $T_a^{bc}$  were already discussed in Section A.6

Birdtracks are a DAG used to represent algebraic tensor equations. The nodes of the DAG are labelled by tensors and the arrows are labeled by the indices of the tensors: upper indices of a tensor are pictured as incoming arrows of the node, and lower indices as outgoing arrows.

We've already discussed in Section A.6 what we will call the **Counter Clock-wise (CC) convention** of drawing birdtrack nodes. Now that we have discussed classical and quantum bnets, we would like to introduce an equivalent, more bnet like, convention that we will call the **Fully Label (FL) convention**. Cvitanovic's birdtracks book Ref.[1] uses the CC convention. We will use both. No confusion will arise, as long as it is clear from context which convention is being used.

Next we review the CC convention and then describe the FL convention for the first time.

#### 1. CC convention

In the CC convention, we must specify for each the node, which arrow is first, and then the CC order in which the arrows enter or leave the node is drawn so that it reproduces the horizontal order of the indices in the algebraic notation for the tensor. We shall often indicate the first arrow by coloring it green.

For example,

$$\delta(b, a) = \mathbb{1}(a = b) = \delta_a^b = a \leftarrow \bullet \rightarrow b \quad (B.19)$$

$$X_{ab}{}^c = \begin{array}{c} a \xleftarrow{\text{green}} \underline{X_{ab}}^c \\ \swarrow \nearrow \\ b \\ \nearrow \\ c \end{array} \quad (\text{B.20})$$

In this picture, the green arrow indicates which tensor index is first horizontally in the algebraic representation of the tensor.

## 2. FL convention

In the FL convention, the arrows must be labelled by random (underlined) variables, and the names of the nodes must also indicate by underlined variables what is the the order of the indices

For example,

$$\delta(b, a) = \mathbb{1}(a = b) = \delta_a^b = a \leftarrow \bullet \rightarrow b \quad (\text{B.21})$$

$$\langle \underline{a}, \underline{b} | \underline{X_{ab}}^c | \underline{c} \rangle = X_{ab}{}^c = \begin{array}{c} \underline{a} = a \xleftarrow{\quad} \underline{X_{ab}}^c \\ \swarrow \nearrow \\ \underline{b} = b \\ \nearrow \\ \underline{c} = c \end{array} \quad (\text{B.22})$$

Sometimes, we will denote this node simply by  $X$ . This is okay as long as we state that  $X = X_{ab}{}^{cd}$ , and we don't start using  $X$  to represent a different version of  $X_{ab}{}^{cd}$  with some of the indices raised or lowered or their horizontal order changed.

Often, we will write simply  $a$  instead of  $\underline{a} = a$ . This is similar to the shorthand  $P(\underline{a} = a) = P(a)$ .

Note that, unlike in the CC convention, in the FL convention, the CC order in which the arrows enter or leave the node, is meaningless. All orders are equivalent. This is akin to the notation for bnets and qbnets.

---

If we don't do either 1 or 2, we won't be able to distinguish between the graphical representations of  $T^{1,2,3}$  and  $T^{2,1,3}$ , for example.

Two other features of the CC and FL conventions that we would like to discuss before ending this section are how to indicate

- **noncyclic index contractions**; i.e., index contractions (i.e., summations) that do not introduce cycles, and
- **traces**; i.e., index contractions that do introduce cycles.

Noncyclic index contractions will be indicated by an arrow connecting two nodes, with the symbol  $\sum a$  midway in the arrow if the index  $a$  is being contracted. For simplicity, we often omit writing the  $\sum a$  altogether.

For example (in CC convention),

$$X_{ab}{}^c = \begin{array}{c} a \xleftarrow{\text{green}} X_{ab}{}^c \\ \swarrow \quad \nearrow \\ b \quad c \end{array}, \quad (X^*)^c{}_{ba} = \begin{array}{c} (X^*)^c{}_{ba} \xleftarrow{\text{green}} a \\ \swarrow \quad \nearrow \\ b \quad c \end{array} \quad (\text{B.23})$$

$$(X^*)^c{}_{ba} X_{ab}{}^c = \begin{array}{c} (X^*)^c{}_{ba} \xleftarrow{\quad} \sum a \xleftarrow{\quad} X_{ab}{}^c \\ \swarrow \quad \nearrow \quad \swarrow \quad \nearrow \\ \quad \quad \quad \sum b \\ \searrow \quad \nearrow \\ \quad \quad \quad \sum c \end{array} \quad (\text{B.24})$$

$$= \begin{array}{c} X^* \xleftarrow{\quad} X \\ \swarrow \quad \nearrow \quad \swarrow \quad \nearrow \\ \quad \quad \quad \end{array} \quad (\text{B.25})$$

Birdtracks are DAGs until we are asked to take a trace of one of their indices. Tracing ruins their acyclicity. The acyclicity of DAGs is mandated by causality. The acyclicity of tracing hints to its acausal (or feedback) nature.

In this book, we will indicate tracing with a red undirected arrow. For example, in the CC convention,

$$\text{tr}_b X_{ab}{}^b = \sum_b X_{ab}{}^b = \begin{array}{c} a \xleftarrow{\quad} X_{ab}{}^c \\ \swarrow \quad \nearrow \\ \text{red vertical line} \end{array} \quad (\text{B.26})$$

If

$$R^x{}_{b_3}{}^{a_3}{}_{a_2}{}^{b_2} S_{x'b_2}{}^{a_2}{}_{a_1}{}^{b_1} = \begin{array}{c} x \downarrow \quad \quad \quad x' \uparrow \\ b_3 \leftarrow R \leftarrow \sum b_2 \leftarrow S \leftarrow b_1 \\ \swarrow \quad \searrow \quad \swarrow \quad \searrow \\ a_3 \quad \quad \quad \sum a_2 \quad \quad \quad a_1 \end{array} \quad (\text{B.27})$$

then

$$\text{tr}_{\underline{x}} R_{b_3 a_2}^{x a_3 b_2} S_{\underline{x} b_2 a_1}^{a_2 b_1} = \begin{array}{c} \text{---} \text{---} \text{---} \\ \downarrow \quad \quad \uparrow \\ \leftarrow R \quad \leftarrow S \rightarrow \\ \swarrow \quad \searrow \quad \swarrow \quad \searrow \end{array} \quad (\text{B.28})$$

When using the FL convention, it becomes clear that birdtracks can be understood as instantiations of qbnets, provided that we weaken slightly the definition of qbnets, by not requiring that the unitarity condition Eq.(B.10) be satisfied. Also, the outgoing arrows of the nodes of a birdtrack must be understood as the result of marginalizer nodes. For example, if the arrows leaving a node are labelled  $a_1$  and  $a_2$ , then these two arrows must be understood as the result of marginalizing an arrow  $a{:}2 = (a_1, a_2)$ .

# Chapter 1

## Casimir Operators

$$M_2 = \begin{array}{c} \text{---} T_i \text{---} T_i \text{---} \\ \text{---} \end{array} \quad (1.1)$$

$$M_4 = \begin{array}{c} \begin{array}{c} \text{---} T_i \text{---} T_j \text{---} T_k \text{---} T_l \text{---} \\ \text{---} T_i \text{---} T_j \text{---} T_l \text{---} T_k \text{---} \end{array} \end{array} \quad (1.2)$$

$$0 = [T_r, M] \quad (1.3)$$

$$= \begin{array}{c} \begin{array}{c} \text{---} T_i \text{---} T_j \text{---} T_k \text{---} T_l \text{---} \\ \text{---} T_i \text{---} T_j \text{---} T_l \text{---} T_k \text{---} \end{array} - \begin{array}{c} \text{---} T_i \text{---} T_j \text{---} T_k \text{---} T_l \text{---} \\ \text{---} T_i \text{---} T_j \text{---} T_l \text{---} T_k \text{---} \end{array} \end{array} \quad (1.4)$$

$$M_2 M_4 = M_4 M_2 \quad (1.5)$$

$$\text{tr}(T_i T_j \dots T_l) = \begin{array}{c} \text{---} T_i \text{---} T_j \text{---} \dots \text{---} T_l \text{---} \\ \text{---} \end{array} \quad (1.6)$$

$$\begin{array}{c} \begin{array}{c} \text{---} T_i \text{---} T_j \text{---} \dots \text{---} T_l \text{---} \\ \text{---} \end{array} - \begin{array}{c} \text{---} T_j \text{---} T_i \text{---} \dots \text{---} T_l \text{---} \\ \text{---} \end{array} = \begin{array}{c} \text{---} T_k \text{---} \dots \text{---} T_l \text{---} \\ \text{---} \end{array} \quad (1.7)$$

$$h_{i_1 i_2 \dots i_p} = \frac{1}{p!} \sum_{\sigma \in \mathcal{S}_p} \text{tr}(T_{\sigma(i_1)} T_{\sigma(i_2)} \dots T_{\sigma(i_p)}) = \begin{array}{c} \begin{array}{c} \xrightarrow{\quad} T_{i_1} \rightarrow T_{i_2} \rightarrow \dots \rightarrow T_{i_p} \xrightarrow{\quad} \\ \text{\scriptsize $\mathcal{S}_p$} \end{array} \end{array} \quad (1.8)$$

$$d_{ijk} = \begin{array}{c} i \\ \text{\scriptsize $d$} \\ j \end{array} \text{\scriptsize $k$} = 2 \begin{array}{c} i \\ \text{\scriptsize $\mathcal{S}_2$} \\ j \end{array} \begin{array}{c} \text{\scriptsize $k$} \end{array} \quad (1.9)$$

Multiplying Jacobi identity by  $T_k$  and taking the trace, we get

$$if_{ijk} = \begin{array}{c} i \\ \text{\scriptsize $if$} \\ j \end{array} \text{\scriptsize $k$} = 2 \begin{array}{c} i \\ \text{\scriptsize $\mathcal{A}_2$} \\ j \end{array} \begin{array}{c} \text{\scriptsize $k$} \end{array} \quad (1.10)$$

## 1.1 Independent Casimirs of Simple Groups

$$\boxed{M = \sum_i T_i x_i} \quad \longleftarrow M \longleftarrow = \sum_i x_i \quad \begin{array}{c} \text{\scriptsize $a$} \longleftarrow T_i \longleftarrow \text{\scriptsize $b$} \end{array} \quad (1.11)$$

$$\text{tr}(M^k) = \begin{array}{c} \longleftarrow M \longleftarrow M \quad \dots \longleftarrow M \longleftarrow \end{array} \quad (1.12)$$

$$= \sum_{i_1 i_2 \dots i_k} \begin{array}{c} \begin{array}{c} \longleftarrow T_{i_1} \longleftarrow T_{i_2} \quad \dots \longleftarrow T_{i_k} \longleftarrow \\ \text{\scriptsize $x_{i_1} x_{i_2} \dots x_{i_k}$} \end{array} \end{array} \quad (1.13)$$

$$= \sum_{i_1 i_2 \dots i_k} \begin{array}{c} \begin{array}{c} \longleftarrow T_{i'_1} \longleftarrow T_{i'_2} \quad \dots \longleftarrow T_{i'_k} \longleftarrow \\ \text{\scriptsize $x_{i_1} x_{i_2} \dots x_{i_k}$} \end{array} \\ \text{\scriptsize $\mathcal{S}_k$} \\ \begin{array}{c} i_1 \quad i_2 \quad i_k \\ \text{\scriptsize $h_{i_1 i_2 \dots i_k}$} \end{array} \end{array} \quad (1.14)$$

Recall Eq.(2.22) for the general characteristic equation of a matrix  $M$

$$0 = \sum_{k=0}^n (-1)^k \left( \text{tr}_{1\dots n-k} \mathcal{A}_{n-k} M^{\otimes n-k} \right) M^k \quad (1.15)$$

$$= \begin{cases} M^n \\ -M^{n-1}(\text{tr} M) \\ +M^{n-2}(\text{tr}_{1\dots 2} \mathcal{A}_2 M^{\otimes 2}) \\ \dots \\ (-1)^n \det(M) \end{cases} \quad (1.16)$$

The coefficients of  $M^k$  are products of traces of a single  $T_i$ . If we calculate the trace of  $M^k$ , then that will entail calculating traces with  $k$  matrices  $T_i$ .

$A_r = \mathfrak{su}(r+1)$	$2, 3, \dots, r+1$
$B_r = \mathfrak{so}(2r+1)$	$2, 4, 6, \dots, 2r$
$C_r = \mathfrak{sp}(2r)$	$2, 4, 6, \dots, 2r$
$D_r = \mathfrak{so}(2r)$	$2, 4, \dots, 2r-2, 2r$
$G_2$	$2, 6$
$F_4$	$2, 6, 8, 12$
$E_6$	$2, 5, 6, 8, 9, 12$
$E_7$	$6, 8, 10, 12, 14, 18$
$E_8$	$8, 12, 14, 18, 20, 24, 30$

Table 1.1: Betti numbers for the simple Lie Algebras

Betti number of a Casimir is the number of  $T_i$  being traced over (i.e., in the loop). Note that the Betti numbers are all even except for  $SU(n)$ .

For all Simple Lie Groups except for  $SU(n)$ , there is a invertible symmetric or skew-symmetric bilinear invariant matrix  $g_{ab}$  satisfying  $g_{ab} G^{bc} = \delta_a^c$ . Hence

$$\begin{array}{c}
\begin{array}{c} g \\ \uparrow \\ g \\ \downarrow \\ \sim T_{i_2} \\ \downarrow \\ \sim T_{i_3} \\ \downarrow \\ \sim T_{i_4} \\ \downarrow \\ \sim T_{i_5} \\ \downarrow \end{array} \\
\end{array}
= (-1)
\begin{array}{c}
\begin{array}{c} g \\ \uparrow \\ \sim T_{i_1} \\ \uparrow \\ g \\ \downarrow \\ \sim T_{i_3} \\ \downarrow \\ \sim T_{i_4} \\ \downarrow \\ \sim T_{i_5} \\ \downarrow \end{array} \\
\end{array}
= (-1)^2
\begin{array}{c}
\begin{array}{c} g \\ \uparrow \\ \sim T_{i_1} \\ \uparrow \\ \sim T_{i_2} \\ \uparrow \\ g \\ \downarrow \\ \sim T_{i_4} \\ \downarrow \\ \sim T_{i_5} \\ \downarrow \end{array} \\
\end{array}
= (-1)^3
\begin{array}{c}
\begin{array}{c} g \\ \uparrow \\ \sim T_{i_1} \\ \uparrow \\ \sim T_{i_2} \\ \uparrow \\ \sim T_{i_3} \\ \uparrow \\ g \\ \downarrow \\ \sim T_{i_5} \\ \downarrow \end{array} \\
\end{array}
= (-1)^4
\begin{array}{c}
\begin{array}{c} g \\ \downarrow \\ \sim T_{i_1} \\ \uparrow \\ \sim T_{i_2} \\ \uparrow \\ \sim T_{i_3} \\ \uparrow \\ \sim T_{i_4} \\ \uparrow \\ g \\ \downarrow \end{array} \\
\end{array}
\quad (1.17)$$



In general, a Casimir with  $k$   $T_i$  equals itself times  $(-1)^k$ . So only Casimirs with even  $k$  are non-zero.

**Claim 1** *The following are a complete set of Casimir operators for the given Groups*

$GL(n, \mathbb{C}) :$

(1.18)

$SU(n) :$

(1.19)

$SO(2r + 1)$  and  $Sp(2r) :$

(1.20)



$$(I_p)_a^b = \text{tr}(T_\lambda^{i_1} T_\lambda^{i_2} \dots T_\lambda^{i_p})(T_\mu^{i_1} T_\mu^{i_2} \dots T_\mu^{i_p})_a^b \quad (1.24)$$

$$= \begin{array}{c} \xrightarrow{\quad} T_\lambda^{i_1} \rightarrow T_\lambda^{i_2} \rightarrow \dots \rightarrow T_\lambda^{i_p} \xleftarrow{\quad} \\ \text{\scriptsize $\vdots$} \\ a \leftarrow T_\mu^{i_1} \leftarrow T_\mu^{i_2} \leftarrow \dots \leftarrow T_\mu^{i_p} \leftarrow b \end{array} \quad (1.25)$$

$$M = \begin{array}{c} \longrightarrow T_\lambda^i \longrightarrow \\ \text{\scriptsize $\vdots$} \\ \longleftarrow T_\mu^i \longleftarrow \end{array} \quad (1.26)$$

$$M = \sum_\rho \begin{array}{c} \swarrow \lambda \quad \searrow \mu \\ C_\rho^\dagger \xleftarrow{\rho} C_\rho \\ \nearrow \lambda \quad \nwarrow \mu \end{array} \begin{array}{c} T_\lambda \\ \text{\scriptsize $\vdots$} \\ T_\mu \end{array} \begin{array}{c} \nearrow \lambda \quad \nwarrow \mu \\ C_\rho^\dagger \xleftarrow{\rho} C_\rho \\ \swarrow \lambda \quad \searrow \mu \end{array} \quad (1.27)$$

$$= \sum_\rho A(\lambda, \rho, \mu) \begin{array}{c} \swarrow \lambda \quad \searrow \mu \\ C_\rho^\dagger \xleftarrow{\rho} C_\rho \\ \nearrow \lambda \quad \nwarrow \mu \end{array} \quad (1.28)$$

$$A(\lambda, \rho, \mu) = \frac{1}{d_\rho} \begin{array}{c} T_\lambda^\dagger \\ \text{\scriptsize $\vdots$} \\ T_\mu \\ \text{\scriptsize $\vdots$} \\ T_\rho \xrightarrow{\rho} T_\rho^\dagger \end{array} \quad (1.29)$$

**Claim 2** *If*

$$\Gamma_2(\rho) = \longleftarrow T_\rho \overset{\text{wavy}}{\longleftrightarrow} T_\rho \longrightarrow \quad (1.30)$$

*then*

$$A(\lambda, \mu, \rho) = -\frac{1}{2} [\Gamma_2(\rho) - \Gamma_2(\lambda) - \Gamma_2(\mu)] \quad (1.31)$$

**proof:**

Recall Eq.(8.20). Square both sides of the equation.

$$\begin{array}{c} j \\ \text{---} T_\rho^j \text{---} \end{array} \quad \begin{array}{c} j \\ \text{---} T_\rho^j \text{---} \end{array} = \left[ \begin{array}{c} j \\ \text{---} T_\lambda \\ \swarrow \quad \nwarrow \\ C_\rho \quad C_\rho^\dagger \\ \searrow \quad \swarrow \end{array} - \begin{array}{c} j \\ \text{---} C_\rho \quad C_\rho^\dagger \\ \swarrow \quad \nwarrow \\ T_\mu \end{array} \right]^2 \quad (1.32)$$

$$\begin{array}{c} \text{---} T_\rho \text{---} T_\rho \text{---} \end{array} = \begin{array}{c} T_\lambda \\ \swarrow \quad \nwarrow \\ C_\rho \quad C_\rho^\dagger \\ \searrow \quad \swarrow \\ T_\mu \end{array} - 2 \begin{array}{c} T_\lambda \\ \swarrow \quad \nwarrow \\ C_\rho \quad C_\rho^\dagger \\ \searrow \quad \swarrow \\ T_\mu \end{array} + \begin{array}{c} \text{---} C_\rho \quad C_\rho^\dagger \text{---} \\ \downarrow \quad \uparrow \\ T_\mu \end{array} \quad (1.33)$$

$$\Gamma_2(\rho) \text{---}\rho\text{---} = \Gamma_2(\lambda) \text{---}\rho\text{---} - 2 \begin{array}{c} T_\lambda \\ \swarrow \quad \nwarrow \\ C_\rho \quad C_\rho^\dagger \\ \searrow \quad \swarrow \\ T_\mu \end{array} + \Gamma_2(\mu) \text{---}\rho\text{---} \quad (1.34)$$

$$\frac{1}{d_\rho} \begin{array}{c} T_\lambda \\ \swarrow \quad \nwarrow \\ C_\rho \quad C_\rho^\dagger \\ \searrow \quad \swarrow \\ T_\mu \end{array} = -\frac{1}{2} [\Gamma_2(\rho) - \Gamma_2(\lambda) - \Gamma_2(\mu)] \quad (1.35)$$

$$\vec{J} = \vec{L} + \vec{S} \quad (1.36)$$

$$\vec{L} \cdot \vec{S} = \frac{1}{2} [J^2 - L^2 - S^2] \quad (1.37)$$

**QED**

$$(I_p)_a^b = (M^p)_a^c b_c \quad (1.38)$$

$$= \sum_{\rho \in irreps} [A(\lambda, \mu, \rho)]^p \quad a \leftarrow \rho \text{---} C_\rho \begin{array}{c} \xleftarrow{\lambda} \\ \xrightarrow{\mu} \end{array} C_\rho^\dagger \leftarrow \rho \text{---} b \quad (1.39)$$

If  $\mu$  is an irrep,

$$\leftarrow \rho \text{---} C_\rho \begin{array}{c} \xleftarrow{\lambda} \\ \xrightarrow{\mu} \end{array} C_\rho^\dagger \leftarrow \rho \text{---} = \frac{T_\rho \begin{array}{c} \xrightarrow{\rho} \\ \xleftarrow{\lambda} \\ \xrightarrow{\mu} \end{array} T_\rho}{d_\mu} \leftarrow \mu \text{---} \quad (1.40)$$

$$= \frac{d_\rho}{d_\lambda} \leftarrow \mu \text{---} \quad (\text{because } \rho \in irreps) \quad (1.41)$$

## 1.2 $\Gamma_2$ and $\Gamma_3$

Three quadratic Casimirs ( $\Gamma_2$ )

$$1. \quad (T_i T_i)_a^b = \Gamma_{fun} \delta_a^b \quad \leftarrow T_i \begin{array}{c} \text{---} \\ \text{---} \end{array} T_i \leftarrow = \Gamma_{fun} \leftarrow \quad (1.42)$$

$$2. \quad \text{tr}(T_i T_j) = \kappa \delta_i^j \quad \text{---} T_i \begin{array}{c} \xleftarrow{\quad} \\ \xrightarrow{\quad} \end{array} T_j \text{---} = \kappa \text{---} \quad (1.43)$$

$$3. \quad f_{ijk} f_{kji'} = \Gamma_{adj} \delta_i^{i'} \quad \text{---} f \begin{array}{c} \text{---} \\ \text{---} \end{array} f \text{---} = \Gamma_{adj} \text{---} \quad (1.44)$$

$$T_i \begin{array}{c} \xleftarrow{\quad} \\ \xrightarrow{\quad} \end{array} T_i = n \Gamma_{fun} = N \kappa \quad (1.45)$$

**Claim 3**

$$\leftarrow T_i \begin{array}{c} \text{---} \\ \text{---} \end{array} T_i \leftarrow = \left( \frac{\kappa N}{n} - \frac{\Gamma_{adj}}{2} \right) \leftarrow T_k \leftarrow \quad (1.46)$$

$$\begin{array}{c} \text{---} \\ \text{---} \end{array} f \begin{array}{c} \text{---} \\ \text{---} \end{array} = \frac{\Gamma_{adj}}{2} \leftarrow T_k \leftarrow \quad (1.47)$$

$$\begin{array}{c} \text{\tiny{wavy}} \\ | \\ f \\ / \quad \backslash \\ \text{\tiny{wavy}} \quad \text{\tiny{wavy}} \\ | \qquad | \\ \text{\tiny{wavy}} \quad \text{\tiny{wavy}} \end{array} = \frac{\Gamma_{adj}}{2} \left( \begin{array}{c} \text{\tiny{wavy}} \\ | \\ \longleftarrow T_k \longrightarrow \end{array} \right) \quad (1.48)$$

proof:  
QED

### 1.3 Dynkin Index

$$DI_\lambda = \frac{\text{tr}(T_\lambda^i T_\lambda^i)}{f_{jk}^i f_{kj}^i} = \frac{\text{Diagram 1}}{\text{Diagram 2}} \quad (1.49)$$

# Chapter 2

## Characteristic Equations

$$M_a^b = a \longleftarrow M \longleftarrow b \quad (2.1)$$

for  $a, b = 1, 2, \dots, n$

$$M^2 = \longleftarrow M^2 \longleftarrow = \longleftarrow M \longleftarrow M \longleftarrow \quad (2.2)$$

$$M \otimes M = M^{\otimes 2} = \begin{array}{c} \longleftarrow M \longleftarrow \\ \\ \longleftarrow M \longleftarrow \end{array} \quad (2.3)$$

$$(M^{\otimes p})_{\alpha}^{\beta} = (M^{\otimes p})_{a:p}^{rev(b^p)} = M_{a_1}^{b_1} M_{a_2}^{b_2} \dots M_{a_p}^{b_p}$$

$$\begin{array}{ccc} \longleftarrow M^{\otimes p} \longleftarrow & & \longleftarrow M \longleftarrow \\ \longleftarrow \parallel \longleftarrow & & \longleftarrow M \longleftarrow \\ \vdots & = & \vdots \\ \longleftarrow \parallel \longleftarrow & & \longleftarrow M \longleftarrow \end{array} \quad (2.4)$$

where  $a_i, b_i \in \mathbb{Z}_{[1,n]}$

$$\text{tr}_{1\dots p}\mathcal{A}[M^{\otimes p}] = \mathcal{A}_{a:p}^{\text{rev}(b:p)} \prod_{i=1}^p M_{b_i}^{a_i} \quad (2.5)$$

$$= \begin{array}{c} \mathcal{A}_p \\ \downarrow \\ \begin{array}{c} \curvearrowright \\ \downarrow \\ \text{Cvitanovic Style} \\ \downarrow \\ M \\ \downarrow \\ M \\ \downarrow \\ M \end{array} \end{array} \quad (2.6)$$

$$= \begin{array}{c} \leftarrow \mathcal{A}_p \leftarrow M \leftarrow \\ \downarrow \\ \leftarrow \mathcal{A}_p \leftarrow M \leftarrow \\ \downarrow \\ \leftarrow \mathcal{A}_p \leftarrow M \leftarrow \end{array} \quad (2.7)$$

$$\det M = \text{tr}_{1\dots n}\mathcal{A}[M^{\otimes n}] \quad (2.8)$$

$$\begin{array}{c} \leftarrow \mathcal{A}_p \leftarrow \\ \downarrow \\ \leftarrow M \leftarrow \\ \downarrow \\ \leftarrow M \leftarrow \end{array} = \frac{1}{p} \left[ \begin{array}{c} \leftarrow \\ \downarrow \\ \leftarrow \mathcal{A}_{p-1} \leftarrow M \leftarrow \\ \downarrow \\ \leftarrow M \leftarrow \end{array} - (p-1) \begin{array}{c} \leftarrow \mathcal{A}_{p-1} \leftarrow M \leftarrow \\ \downarrow \\ \leftarrow M \leftarrow \end{array} \right] \quad (2.9)$$

$$\left\{ \begin{array}{l} \begin{array}{c} \leftarrow \mathcal{A}_3 \leftarrow \\ \downarrow \\ \leftarrow M \leftarrow \\ \downarrow \\ \leftarrow M \leftarrow \end{array} = \frac{1}{3} \left[ \begin{array}{c} \leftarrow \\ \downarrow \\ \leftarrow \mathcal{A}_2 \leftarrow M \leftarrow \\ \downarrow \\ \leftarrow M \leftarrow \end{array} - 2 \begin{array}{c} \leftarrow \mathcal{A}_2 \leftarrow M \leftarrow \\ \downarrow \\ \leftarrow M \leftarrow \end{array} \right] \\ \begin{array}{c} \leftarrow \mathcal{A}_2 \leftarrow \\ \downarrow \\ \leftarrow M \leftarrow \end{array} = \frac{1}{2} \left[ \begin{array}{c} \leftarrow \\ \downarrow \\ \leftarrow M \leftarrow \end{array} - \leftarrow M \leftarrow \right] \end{array} \right. \quad (2.10)$$



$$\left\{ \begin{array}{l} \text{Diagram 1} \\ \text{Diagram 2} \\ \text{Diagram 3} \end{array} \right\} = \frac{1}{3} \left[ \begin{array}{l} \text{Diagram 4} \\ \text{Diagram 5} \\ \text{Diagram 6} \end{array} \right] - 2 \left[ \begin{array}{l} \text{Diagram 7} \\ \text{Diagram 8} \end{array} \right] \quad (2.11)$$

$$\left\{ \begin{array}{l} \text{Diagram 9} \\ \text{Diagram 10} \end{array} \right\} = \frac{1}{2} \left[ \begin{array}{l} \text{Diagram 11} \\ \text{Diagram 12} \end{array} \right] - \left[ \text{Diagram 13} \right]$$

Let

$$\tau = \text{tr}(M) \quad (2.12)$$

$$\text{tr}_{1,2,3} \mathcal{A}_3(M^{\otimes 3}) = \frac{1}{3} \left[ \tau \text{tr}_{1,2} \mathcal{A}_2(M^{\otimes 2}) - \text{tr}(M^2) \tau + \text{tr} M^3 \right] \quad (2.13)$$

$$\text{tr}_{1,2} \mathcal{A}_2 M^{\otimes 2} = \frac{1}{2} \left[ \tau^2 - \text{tr}(M^2) \right] \quad (2.14)$$

$$\text{tr}_{1,2,3} \mathcal{A}_3(M^{\otimes 3}) = \frac{1}{3} \left[ \frac{1}{2} \tau^3 - \frac{3}{2} \text{tr}(M^2) \tau + \text{tr} M^3 \right] \quad (2.15)$$

$$= \frac{1}{3!} \left[ \tau^3 - 3 \text{tr}(M^2) \tau + 2 \text{tr} M^3 \right] \quad (2.16)$$

$$\text{tr}_{1\dots p} \mathcal{A}_p M = \frac{1}{p} \sum_{k=1}^p (-1)^{k-1} \left( \text{tr}_{1\dots p-k} \mathcal{A}_{p-k} M^{\otimes p-k} \right) \text{tr}(M^k) \quad (2.17)$$

$$\mathcal{A}_p = 0 \quad \text{if } p > n \quad (2.18)$$

$$\epsilon_{a_1, a_2, \dots, a_{n+1}} = 0 \quad (2.19)$$

Two of the  $a_i$  must be equal, so that element of the  $\epsilon$  is zero

$I$  is the  $n \times n$  identity matrix

$$\text{tr}_{2\dots n+1} \mathcal{A}_{n+1} I \otimes M^{\otimes n} = 0 \quad (2.20)$$

$$0 = \sum_{k=0}^n (-1)^k \left( \text{tr}_{1\dots n-k} \mathcal{A}_{n-k} M^{\otimes n-k} \right) M^k \quad (2.21)$$

$$= \begin{cases} M^n \\ -M^{n-1}(\text{tr} M) \\ +M^{n-2}(\text{tr}_{1\dots 2} \mathcal{A}_2 M^{\otimes 2}) \\ \dots \\ (-1)^n \det(M) \end{cases} \quad (2.22)$$

# Chapter 3

## Clebsch-Gordan Coefficients

This chapter is based on Ref.[1].

Suppose that for some  $M \in \mathbb{C}^{d \times d}$ , we have

$$M = C^\dagger D C \quad (3.1)$$

where  $D$  is a diagonal matrix and  $C = C^{d \times d}$  is unitary. Then one can partition  $C$  into rectangular submatrices  $C_\lambda$  that have  $d_\lambda < d$  rows, with one  $C_\lambda$  for each eigenvalue  $\lambda$  of  $C$ . Likewise, we can partition  $C^\dagger$  into rectangular submatrices  $C_\lambda^\dagger$  that have  $d_\lambda < d$  columns, with one  $C_\lambda^\dagger$  for each eigenvalue  $\lambda$  of  $C$ . Thus, if  $I^{d_\lambda \times d_\lambda}$  is the  $d_\lambda \times d_\lambda$  identity matrix,

$$\begin{bmatrix} 0 \\ C_\lambda^{d_\lambda \times d} \\ 0 \end{bmatrix}^{d \times d} = \begin{bmatrix} 0 & 0 & 0 \\ 0 & I^{d_\lambda \times d_\lambda} & 0 \\ 0 & 0 & 0 \end{bmatrix}^{d \times d} C^{d \times d} \quad (3.2)$$

$$\begin{bmatrix} 0 & (C^\dagger)_\lambda^{d \times d_\lambda} & 0 \end{bmatrix}^{d \times d} = (C^\dagger)^{d \times d} \begin{bmatrix} 0 & 0 & 0 \\ 0 & I^{d_\lambda \times d_\lambda} & 0 \\ 0 & 0 & 0 \end{bmatrix}^{d \times d} \quad (3.3)$$

The matrices  $C_\lambda$  are called the **Clebsch-Gordan Coefficient** (CGC) matrices.

Let  $b^{:nb} = (b_1, b_2, \dots, b_{nb})$  where  $b_i \in Z_{[0, d_{\mu_i}]}$  and  $a \in Z_{[1, d_\lambda]}$ . Hence,

$$d_\lambda = \prod_{i=1}^{:nb} d_{\mu_i} \quad (3.4)$$

Now define the birdtracks

$$(C_\lambda)_a^{rev(b^{:nb})} = \lambda a \leftarrow C_\lambda \begin{array}{l} \nwarrow \mu_1 b_1 \\ \nwarrow \mu_2 b_2 \\ \nwarrow \mu_{nb} b_{nb} \end{array} \quad (3.5)$$

and

$$(C_\lambda^\dagger)^a_{b:nb} = \begin{array}{c} \mu_1 b_1 \\ \swarrow \\ \mu_2 b_2 \leftarrow (C_\lambda^\dagger) \leftarrow \lambda a \\ \searrow \\ \mu_{nb} b_{nb} \end{array} \quad (3.6)$$

We will assume there is no difference between when a Greek letter is lowered or raised. Also, all summations over a Greek letter will be stated explicitly; i.e., no implicit summations over repeated Greek letters.

On the other hand, the Latin letter indices  $b_i$  of  $C_\lambda$  and  $C_\lambda^\dagger$  may be lowered or raised and their arrows changed from outgoing to incoming or vice versa. Furthermore, we will use implicit summation over repeated Latin letters.

The Greek letters label representation of the group (not necessarily irreps). Each  $b_i$  labels a subcategory of  $\mu_i$ .

Recall that if  $|x\rangle$  for  $x \in \text{val}(\underline{x})$  is a complete, orthonormal basis in Quantum Mechanics, then

$$\langle x|y\rangle = \delta(x, y) \quad (\text{orthonormality}) \quad (3.7)$$

and

$$\sum_x |x\rangle\langle x| = 1 \quad (\text{completeness}) \quad (3.8)$$

Furthermore, if we define

$$\pi_x = |x\rangle\langle x| \quad (3.9)$$

then  $\pi_x$  is a projection operator so

$$\pi_x \pi_x = \pi_x \quad (3.10)$$

and

$$\pi_x |y\rangle = |y\rangle \delta(x, y), \quad \langle y|\pi_x = \langle y| \delta(x, y) \quad (3.11)$$

If we identify  $C_\lambda$  with  $\langle x|$ , and  $C_\lambda^\dagger$  with  $|x\rangle$ , then  $C_\lambda$  and  $C_\lambda^\dagger$  satisfy analogous identities:

$$\begin{array}{c} (C_\lambda)_a^{\text{rev}(b:nb)} (C_\mu^\dagger)^{a'}_{b:nb} = \delta(\lambda, \mu) \delta_a^{a'}, \quad C_\lambda C_\mu^\dagger = \delta(\mu, \lambda) \\ \begin{array}{c} \sum b_1 \\ \swarrow \\ a \leftarrow C_\lambda \leftarrow \sum b_2 \leftarrow (C_\mu^\dagger) \leftarrow a' \\ \searrow \\ \sum b_{nb} \end{array} = \delta(\mu, \lambda) \quad a \leftarrow \bullet \rightarrow a' \end{array} \quad (3.12)$$

$$\begin{aligned}
& \sum_{\lambda} (C_{\lambda}^{\dagger})^a_{b:nb} (C_{\lambda})_a^{rev((b') : nb)} = \delta_{b:nb}^{rev((b') : nb)}, \quad \sum_{\lambda} C_{\lambda}^{\dagger} C_{\lambda} = 1 \\
& \begin{array}{ccc}
b_1 & & b'_1 \\
\swarrow & & \swarrow \\
\Sigma_{\lambda} \quad b_2 \leftarrow (C_{\lambda}^{\dagger}) \leftarrow \Sigma a & \leftarrow C_{\lambda} \leftarrow b'_2 & = \quad b_2 \leftarrow \bullet \leftarrow b'_2 \\
\searrow & & \searrow \\
b_{nb} & & b'_{nb}
\end{array} \quad \begin{array}{ccc}
b_1 & \leftarrow \bullet & b'_1 \\
b_2 & \leftarrow \bullet & b'_2 \\
b_{nb} & \leftarrow \bullet & b'_{nb}
\end{array}
\end{aligned} \tag{3.13}$$

$$\begin{aligned}
& (C_{\lambda})_a^{rev((b') : nb)} (P_{\mu})_{(b') : nb}^{rev(b : nb)} = \delta(\mu, \lambda) (C_{\mu})_a^{rev(b : nb)}, \quad C_{\lambda} P_{\mu} = \delta(\mu, \lambda) C_{\mu} \\
& \begin{array}{ccc}
\Sigma b'_1 & & b_1 \\
\swarrow & & \swarrow \\
a \leftarrow C_{\lambda} \leftarrow \Sigma b'_2 & \leftarrow P_{\mu} \leftarrow b_2 & = \delta(\mu, \lambda) \quad a \leftarrow C_{\lambda} \leftarrow b_2 \\
\swarrow & & \swarrow \\
\Sigma b'_{nb} & & b_{nb}
\end{array}
\end{aligned} \tag{3.14}$$

$$\begin{aligned}
& (P_{\mu})_{b:nb}^{rev((b') : nb)} (C_{\lambda}^{\dagger})^a_{(b') : nb} = \delta(\mu, \lambda) (C_{\mu}^{\dagger})^a_{b:nb}, \quad P_{\mu} C_{\lambda}^{\dagger} = \delta(\mu, \lambda) C_{\mu}^{\dagger} \\
& \begin{array}{ccc}
b_1 & & \Sigma b'_1 \\
\swarrow & & \swarrow \\
b_2 \leftarrow P_{\mu} \leftarrow \Sigma b'_2 & \leftarrow (C_{\lambda}^{\dagger}) \leftarrow a & = \delta(\mu, \lambda) \quad b_2 \leftarrow (C_{\lambda}^{\dagger}) \leftarrow a \\
\swarrow & & \swarrow \\
b_{nb} & & \Sigma b'_{nb}
\end{array} \quad \begin{array}{ccc}
b_1 & & \\
\swarrow & & \\
b_2 & \leftarrow (C_{\lambda}^{\dagger}) & \\
\swarrow & & \\
b_{nb} & &
\end{array}
\end{aligned} \tag{3.15}$$

## Chapter 4

**Dynkin Diagrams: COMING  
SOON**

## Chapter 5

**General Relativity Nets: COMING  
SOON**

# Chapter 6

## Integrals over a Group

$$\int dg G_a^b G_c^d \dots (G^\dagger)_e^f (G^\dagger)_g^h \quad (6.1)$$

$\mathcal{S}$  = singlets,  $\mathcal{S}^c$  = nonsinglets

$$\int dg = 1, \quad \int dg G_\lambda = \mathbb{1}(\lambda \in \mathcal{S}) \quad (6.2)$$

$$G_a^b = a \xleftarrow{\text{green}} G \longleftarrow b, \quad (G^\dagger)_b^a = b \xleftarrow{\text{green}} G^\dagger \longleftarrow a \quad (6.3)$$

Out arrow always first

We assume that  $G$  is unitary ( $G^\dagger G = G G^\dagger = 1$ )

$$\longleftarrow G^\dagger \longleftarrow G \longleftarrow = \longleftarrow G \longleftarrow G^\dagger \longleftarrow = \longleftarrow \bullet \longleftarrow \quad (6.4)$$

$$\begin{aligned} & \longleftarrow G \longleftarrow \\ G \otimes G \otimes G^\dagger = & \longleftarrow G \longleftarrow \\ & \longleftarrow G^\dagger \longleftarrow \end{aligned} \quad (6.5)$$

$\delta_a^d \delta_c^b = \frac{1}{d} \delta_a^b \delta_c^d + \frac{1}{\kappa} (T^i)_a^b (T^i)_c^d$

$$\begin{array}{c} a \longleftarrow \bullet \longleftarrow d \\ b \longrightarrow \bullet \longrightarrow c \end{array} = \frac{1}{d} \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} + \frac{1}{\kappa} \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} T^i \text{---} T^i \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} \quad (6.6)$$

Will set  $\kappa = 1$  from here on.



$$\delta_a^d \delta_c^b = \frac{1}{d} \delta_a^b \delta_c^d + \frac{1}{\kappa} (T^i)_a^b (T^i)_c^d$$

$$\begin{array}{c} a \longleftarrow G \longleftarrow d \\ b \longrightarrow G^\dagger \longrightarrow c \end{array} = \frac{1}{d} \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} + \begin{array}{c} \longleftarrow G \longleftarrow \\ \longrightarrow G^\dagger \longrightarrow \end{array} T^i \text{ wavy } T^i \begin{array}{c} \longrightarrow \\ \longleftarrow \end{array} \quad (6.7)$$

$$G_a^{a'} (G^\dagger)_{b'}^b (T^{i'})_{a'}^{b'} G_{i'i} = (T^i)_a^b$$

$$\begin{array}{c} \longleftarrow G \longleftarrow \\ \longrightarrow G^\dagger \longrightarrow \end{array} T^i \text{ wavy } G \text{ wavy} = \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} T^i \begin{array}{c} \longrightarrow \\ \longleftarrow \end{array} \quad (6.8)$$

$$G_a^{a'} (G^\dagger)_{b'}^b (T^i)_{a'}^{b'} = (T^i)_a^b G_{ii'}$$

$$\begin{array}{c} \longleftarrow G \longleftarrow \\ \longrightarrow G^\dagger \longrightarrow \end{array} T^i \text{ wavy} = \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} T^i \text{ wavy } G \text{ wavy} \quad (6.9)$$

$$\begin{array}{c} a \longleftarrow G \longleftarrow d \\ b \longrightarrow G^\dagger \longrightarrow c \end{array} = \frac{1}{d} \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} + \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} T^i \text{ wavy } G \text{ wavy } T^i \begin{array}{c} \longrightarrow \\ \longleftarrow \end{array} \quad (6.10)$$

$$\int dg G_a^d (G^\dagger)_b^c = \frac{1}{d} \delta_a^b \delta_c^d$$

$$\int dg \begin{array}{c} a \longleftarrow G \longleftarrow d \\ b \longrightarrow G^\dagger \longrightarrow c \end{array} = \frac{1}{d} \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} \quad (6.11)$$

$$\int dg \begin{array}{c} a \longleftarrow G_\mu \longleftarrow d \\ b \longrightarrow G_\lambda^\dagger \longrightarrow c \end{array} = 0 \quad \text{if } \mu \neq \lambda \quad (6.12)$$

$$\delta_a^d \delta_c^b = \frac{1}{d} \delta_a^b \delta_c^d + \sum_{\lambda \in \mathcal{S}^c} \frac{1}{\kappa} (T_\lambda^i)_a{}^b (T_\lambda^i)_c{}^d$$

$$\begin{array}{c} a \longleftarrow \bullet \longleftarrow d \\ b \longrightarrow \bullet \longrightarrow c \end{array} = \frac{1}{d} \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} + \frac{1}{\kappa} \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} T_\lambda^i \sim_\lambda \sim T_\lambda^i \begin{array}{c} \longrightarrow \\ \longleftarrow \end{array} \quad (6.13)$$

**Claim 4**

$$\int dg \, G = \sum_{\lambda \in \mathcal{S}} P_\lambda \quad (6.14)$$

$$\int dg \, (G_\mu)_{ab} (G_\nu)^{cd} = \sum_{\lambda \in \mathcal{S}} (P_\lambda)_{ab}{}^{cd} \quad (6.15)$$

**proof:**

$$\int dg \, G = \sum_{\lambda} C_\lambda^\dagger \left[ \int dg \, G_\lambda \right] C_\lambda \quad (6.16)$$

$$= \sum_{\lambda \in \mathcal{S}} C_\lambda^\dagger C_\lambda \quad (6.17)$$

$$= \sum_{\lambda \in \mathcal{S}} P_\lambda \quad (6.18)$$

$$(C_{\lambda i}^\dagger)_{ac} = \begin{array}{c} a \longleftarrow \\ \phantom{a \longleftarrow} \\ c \longleftarrow \end{array} C_{\lambda i}^\dagger \longleftarrow i \quad (6.19)$$

$$(G_\mu)_a{}^{a'} (G_\nu)_{b'}{}^b (C_{\lambda i}^\dagger)_{a'}{}^{b'} = (C_{\lambda i'}^\dagger)_a{}^b (G_\lambda)_{i' i}$$

$$\begin{array}{c} \longleftarrow G_\mu \longleftarrow \\ \phantom{\longleftarrow G_\mu \longleftarrow} \\ \longleftarrow G_\nu \longleftarrow \end{array} C_\lambda^\dagger \longleftarrow = \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} C_\lambda^\dagger \longleftarrow G_\lambda \longleftarrow \quad (6.20)$$

$$\int dg \begin{array}{c} \longleftarrow G_\mu \longleftarrow \\ \longleftarrow G_\nu \longleftarrow \end{array} = \int dg \sum_\lambda \begin{array}{c} \longleftarrow \\ \longleftarrow \end{array} \left. \begin{array}{c} \longleftarrow \\ \longleftarrow \end{array} \right\} C_\lambda^\dagger \longleftarrow G_\lambda \longleftarrow \left. \begin{array}{c} \longrightarrow \\ \longrightarrow \end{array} \right\} C_\lambda \begin{array}{c} \longrightarrow \\ \longrightarrow \end{array} \quad (6.21)$$

$$= \sum_{ij} \sum_{\lambda} (C_{\lambda i}^{\dagger})_{ab} (C_{\lambda j})^{cd} \int dg (G_{\lambda})_{ij} \quad (6.22)$$

$$= \sum_i \sum_{\lambda \in \mathcal{S}} (C_{\lambda i}^\dagger)_{ab} (C_{\lambda i})^{cd} \quad (6.23)$$

$$= \sum_{\lambda \in \mathcal{S}} (P_\lambda)_{ab}{}^{cd} \quad (6.24)$$

$$\int dg (G_\lambda)_{ij} = \delta(i, j) \mathbb{1}(\lambda \in \mathcal{S}) \quad (6.25)$$

QED

## 6.1 Character Orthonormality Relation

$$\chi_\lambda(g) \stackrel{\text{def}}{=} \text{tr} G_\lambda = (G_\lambda)_a^a \quad (6.26)$$

$$\mathrm{tr} G_{\lambda}^{\dagger} = (G_a^a)^* = \chi_{\lambda}(g)^* \quad (6.27)$$

$$\int dg \chi_\lambda(hg) \chi_\mu^*(gf^\dagger) = \delta(\mu, \lambda) \frac{1}{d_\lambda} \chi_\lambda(hf^\dagger)$$

$$\int dg \left( \begin{array}{c} \xrightarrow{h} \\ \xleftarrow{G_\lambda} \end{array} \right) = \frac{1}{d_\lambda} \left( \begin{array}{c} \xrightarrow{h} \\ \xleftarrow{f^\dagger} \end{array} \right) \quad \text{if } \mu \text{ and } \lambda \text{ are irreps} \quad (6.28)$$

$$G_a{}^b = \frac{d}{d(h^\dagger)_b{}^a} \underbrace{\chi(h^\dagger g)}_{(h^\dagger)_b{}^a G_a{}^b} \quad (6.29)$$

## 6.2 $SU(n)$ examples

In  $SU(n)$ ,

$$n = d_{\lambda_0} \quad (6.30)$$

where  $\lambda_0$  is the defining rep. ( $\mathcal{G} \subset \mathbb{C}^{n \times n}$ )  
 $SU(n) \ G \otimes G$

$$G_a^{\ b} G_a^{\ b} = \begin{array}{c} a \longleftarrow G \longleftarrow b \\ \\ c \longleftarrow G \longleftarrow d \end{array} \quad (6.31)$$

$$\begin{array}{c} \longleftarrow \\ \\ \longleftarrow \end{array} = \begin{array}{c} \longleftarrow \mathcal{S}_2 \longleftarrow \\ \parallel \\ \longleftarrow \end{array} + \begin{array}{c} \longleftarrow \mathcal{A}_2 \longleftarrow \\ \parallel \\ \longleftarrow \end{array} \quad (6.32)$$

$$\begin{array}{c} \longleftarrow \mathcal{S}_2 \longleftarrow \\ \parallel \\ \longleftarrow \end{array} = \frac{1}{2} \left\{ \begin{array}{c} \longleftarrow \\ \longleftarrow \end{array} + \begin{array}{c} \longleftarrow \longleftarrow \\ \updownarrow \\ \longleftarrow \longleftarrow \end{array} \right\} \quad (6.33)$$

$$\begin{array}{c} \longleftarrow \mathcal{A}_2 \longleftarrow \\ \parallel \\ \longleftarrow \end{array} = \frac{1}{2} \left\{ \begin{array}{c} \longleftarrow \\ \longleftarrow \end{array} - \begin{array}{c} \longleftarrow \longleftarrow \\ \updownarrow \\ \longleftarrow \longleftarrow \end{array} \right\} \quad (6.34)$$

$$d_{\mathcal{S}} = \frac{1}{2} \left\{ \begin{array}{c} \text{red arc} \\ \longleftarrow \longleftarrow \\ \text{red arc} \end{array} + \begin{array}{c} \text{red arc} \\ \longleftarrow \longleftarrow \\ \updownarrow \\ \text{red arc} \end{array} \right\} \quad (6.35)$$

$$= \frac{n(n+1)}{2} \quad (6.36)$$

$$d_{\mathcal{A}} = \frac{1}{2} \left\{ \begin{array}{c} \text{red arc} \\ \longleftarrow \longleftarrow \\ \text{red arc} \end{array} - \begin{array}{c} \text{red arc} \\ \longleftarrow \longleftarrow \\ \updownarrow \\ \text{red arc} \end{array} \right\} \quad (6.37)$$

$$= \frac{n(n-1)}{2} \quad (6.38)$$

Hence,  $d_{\mathcal{S}} = 1$  iff  $n = 1$ , and  $d_{\mathcal{A}} = 1$  iff  $n = 2$ . Therefore, for  $SU(n)$

$$\int dg \ G \otimes G = 0 \quad \text{if } n > 2 \quad (6.39)$$

$$G^\dagger \otimes G^\dagger \otimes G \otimes G$$

$$\begin{array}{c}
\longrightarrow \\
\longrightarrow \\
\longleftarrow \\
\longleftarrow
\end{array}
=
\begin{array}{c}
\longrightarrow \mathcal{S}_2 \longrightarrow \\
\parallel \\
\longrightarrow \longrightarrow \\
\longleftarrow \longleftarrow \\
\longleftarrow \longleftarrow
\end{array}
+
\begin{array}{c}
\longrightarrow \mathcal{A}_2 \longrightarrow \\
\parallel \\
\longrightarrow \longrightarrow \\
\longleftarrow \longleftarrow \\
\longleftarrow \longleftarrow
\end{array}
\quad (6.40)$$

$$\begin{array}{c}
\longrightarrow \mathcal{S}_2 \longrightarrow \\
\parallel \\
\longrightarrow \longrightarrow \\
\longleftarrow \mathcal{S}_2 \longleftarrow \\
\parallel \\
\longleftarrow \longleftarrow \\
\mathcal{S}_2 \otimes \mathcal{S}_2
\end{array}
+
\begin{array}{c}
\longrightarrow \mathcal{A}_2 \longrightarrow \\
\parallel \\
\longrightarrow \longrightarrow \\
\longleftarrow \mathcal{A}_2 \longleftarrow \\
\parallel \\
\longleftarrow \longleftarrow \\
\mathcal{A}_2 \otimes \mathcal{A}_2
\end{array}
+
\begin{array}{c}
\longrightarrow \mathcal{S}_2 \longrightarrow \\
\parallel \\
\longrightarrow \longrightarrow \\
\longleftarrow \mathcal{A}_2 \longleftarrow \\
\parallel \\
\longleftarrow \longleftarrow \\
\mathcal{S}_2 \otimes \mathcal{A}_2
\end{array}
+
\begin{array}{c}
\longrightarrow \mathcal{A}_2 \longrightarrow \\
\parallel \\
\longrightarrow \longrightarrow \\
\longleftarrow \mathcal{S}_2 \longleftarrow \\
\parallel \\
\longleftarrow \longleftarrow \\
\mathcal{A}_2 \otimes \mathcal{S}_2
\end{array}
\quad (6.41)$$

$$P_1 = \frac{1}{n^2}
\begin{array}{c}
\longrightarrow \\
\longrightarrow \\
\longleftarrow \\
\longleftarrow
\end{array}
\begin{array}{c}
\longrightarrow \\
\longrightarrow \\
\longleftarrow \\
\longleftarrow
\end{array}
\quad (6.43)$$

$$P_2 = \frac{1}{n^2}
\begin{array}{c}
\longrightarrow \\
\longrightarrow \\
\longleftarrow \\
\longleftarrow
\end{array}
\begin{array}{c}
\longrightarrow \\
\longrightarrow \\
\longleftarrow \\
\longleftarrow
\end{array}
= \frac{1}{n^2}
\begin{array}{c}
\longrightarrow \\
\longrightarrow \\
\longleftarrow \\
\longleftarrow
\end{array}
\begin{array}{c}
\longrightarrow \\
\longrightarrow \\
\longleftarrow \\
\longleftarrow
\end{array}
\quad (6.44)$$

$$\pi_{\mathcal{S}} = \frac{1}{d_{\mathcal{S}}}
\begin{array}{c}
\longrightarrow \mathcal{S}_2 \\
\parallel \\
\longrightarrow \longrightarrow \\
\longleftarrow \longleftarrow \\
\longleftarrow \longleftarrow
\end{array}
, \quad
\pi_{\mathcal{A}} = \frac{1}{d_{\mathcal{A}}}
\begin{array}{c}
\longrightarrow \mathcal{A}_2 \\
\parallel \\
\longrightarrow \longrightarrow \\
\longleftarrow \longleftarrow \\
\longleftarrow \longleftarrow
\end{array}
\quad (6.45)$$

$P_1, P_2, \pi_{\mathcal{A}}, \pi_{\mathcal{S}}$  are all properly normalized projectors

$$\dim(\pi_{\mathcal{S}}) = \text{tr} \pi_{\mathcal{S}} = \frac{1}{d_{\mathcal{S}}}
\begin{array}{c}
\longrightarrow \mathcal{S}_2 \\
\parallel \\
\longrightarrow \longrightarrow \\
\longleftarrow \longleftarrow \\
\longleftarrow \longleftarrow
\end{array}
= 1 \quad (6.46)$$

$$\dim(\pi_{\mathcal{A}}) = 1 \quad (6.47)$$

$$\pi_{\mathcal{A}}\pi_{\mathcal{S}} = 1 \quad (6.48)$$

$$P_i(\mathcal{S}_2 \otimes \mathcal{A}_2)P_i = 0, \quad P_i(\mathcal{A}_2 \otimes \mathcal{S}_2)P_i = 0 \quad (6.49)$$

$$P_i(\mathcal{S}_2 \otimes \mathcal{S}_2)P_i = P_i\pi_{\mathcal{S}}P_i, \quad P_i(\mathcal{A}_2 \otimes \mathcal{A}_2)P_i = P_i\pi_{\mathcal{A}}P_i \quad (6.50)$$

$$\begin{array}{c} \longrightarrow \\ \longrightarrow \\ \longleftarrow \\ \longleftarrow \end{array} = \pi_{\mathcal{S}} + \pi_{\mathcal{A}} + \dots \quad (6.51)$$

$$\int dg \begin{array}{c} \longrightarrow G^\dagger \longrightarrow \\ \longrightarrow G^\dagger \longrightarrow \\ \longleftarrow G \longleftarrow \end{array} = \int dg (\pi_{\mathcal{S}} + \pi_{\mathcal{S}}) \begin{array}{c} \longrightarrow G^\dagger \longrightarrow \\ \longrightarrow G^\dagger \longrightarrow \\ \longleftarrow G \longleftarrow \end{array} (\pi_{\mathcal{S}} + \pi_{\mathcal{S}}) \quad (6.52)$$

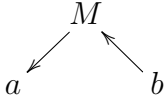
$$\begin{array}{c} \longrightarrow G^\dagger \longrightarrow \\ \longrightarrow G^\dagger \longrightarrow \\ \longleftarrow G \longleftarrow \end{array} = \pi_{\mathcal{S}} + \pi_{\mathcal{A}} \begin{array}{c} \longrightarrow G^\dagger \longrightarrow \\ \longrightarrow G^\dagger \longrightarrow \\ \longleftarrow G \longleftarrow \end{array} \quad (6.53)$$

# Chapter 7

## Invariant Tensors

This chapter is based on Ref.[1].

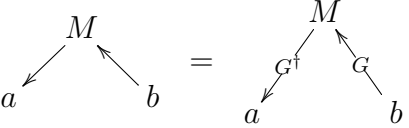
A **bilinear form** is a linear function  $m : V^{\dagger n} \times V^n \rightarrow \mathbb{C}$  with  $V^{\dagger n}, V^n = \mathbb{C}^n$ . For example,

$$m(x^{\dagger n}, y^n) = x^{\dagger a} M_a^b y_b \quad (7.1)$$


$m()$  is said to be invariant if

$$m(x^{\dagger n}, y^n) = m(x^{\dagger n} G^{\dagger}, G y^n) \quad (7.2)$$

$m()$  is invariant iff matrix  $M$  is an **invariant matrix**; i.e., iff

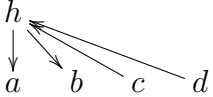
$$M_a^b = (G^{\dagger})_a^{a'} G_{b'}^b M_{a'}^{b'} \quad (7.3)$$


$$M = G^{\dagger} M G \quad (7.4)$$

If  $G$  is unitary,

$$GM = MG, \quad [G, M] = 0 \quad (7.5)$$

A **multilinear form** is a linear function  $h : V^{\dagger n^p} \times V^{n^q} \rightarrow \mathbb{C}$  with  $V^{\dagger d}, V^d = \mathbb{C}^d$ . For example,

$$h(w^{\dagger}, x^{\dagger}, y, z) = h_{ab}^{cd} w^{\dagger a} x^{\dagger b} y_c z_d \quad (7.6)$$


$h()$  is said to be invariant if

$$h(w^{\dagger}, x^{\dagger}, y, z) = h(w^{\dagger} G^{\dagger}, x^{\dagger} G^{\dagger}, G y, G z) \quad (7.7)$$

$h()$  is invariant iff tensor  $h_{ab}{}^{cd}$  is a **invariant tensor** (IT); i.e., iff

$$h_{ab}{}^{cd} = (G^\dagger)_a{}^{a'} (G^\dagger)_b{}^{b'} h_{a'b'}{}^{c'd'} G_{c'}{}^c G_{d'}{}^d \quad \begin{array}{c} h \\ \swarrow \downarrow \searrow \\ a \quad b \quad c \quad d \end{array} = \begin{array}{c} h \\ \swarrow \downarrow \searrow \\ G^\dagger \downarrow G^\dagger \downarrow G \downarrow G \\ a \quad b \quad c \quad d \end{array} \quad (7.8)$$

A **composed IT** is a IT that can be written as a product or contraction of ITs.

A **tree IT** is a composed ITs without any loops.

A **primitive IT** is a IT that can be expressed as a linear combination of a finite number of tree ITs.

The **primitiveness assumption**: All IT are primitive.

Examples. Suppose  $x, y, z \in \mathbb{R}^3$  and  $i, j, k \in \{1, 2, 3\}$ .

- Primitive ITs

$$length(x) = \delta_{ij} x_i x_j \quad volume(x, y, z) = \epsilon_{ijk} x_i y_j z_k \quad (7.9)$$

$$\delta_{ij} = i \text{ --- } j, \quad \epsilon_{ijk} = \begin{array}{c} \epsilon \\ \swarrow \downarrow \searrow \\ i \quad j \quad k \end{array} \quad (7.10)$$

- Tree ITs

$$\delta_{ij} \epsilon_{klm} = \begin{array}{c} i \\ | \\ j \end{array} \quad \begin{array}{c} \epsilon \\ \swarrow \downarrow \searrow \\ k \quad l \quad m \end{array} \quad (7.11)$$

$$\epsilon_{ijm} \delta_{mn} \epsilon_{nkl} = \begin{array}{c} \epsilon_{ijm} \text{---} \sum^m \text{---} \epsilon_{mkl} \\ \swarrow \downarrow \quad \downarrow \swarrow \searrow \\ i \quad j \quad k \quad l \end{array} \quad (7.12)$$

- Non-tree IT

$$\epsilon_{ims} \epsilon_{jnm} \epsilon_{krn} \epsilon_{lsr} = \begin{array}{c} i \text{ --- } \epsilon_{ims} \text{---} \sum^s \text{---} \epsilon_{lsr} \text{ --- } l \\ \quad \quad \quad \downarrow \quad \quad \quad \downarrow \\ \quad \quad \quad \sum^m \quad \quad \quad \sum^r \\ j \text{ --- } \epsilon_{jnm} \text{---} \sum^n \text{---} \epsilon_{krn} \text{ --- } k \end{array} \quad (7.13)$$



$$\epsilon_{ims}\epsilon_{jnm}\epsilon_{krn}\epsilon_{lsr} = \delta_{ij}\delta_{kl} + \delta_{il}\delta_{jk}$$

$$\begin{array}{c} i \text{ --- } \epsilon_{ims} \text{ --- } \sum_m^s \text{ --- } \epsilon_{lsr} \text{ --- } l \\ j \text{ --- } \epsilon_{jnm} \text{ --- } \sum_n^r \text{ --- } \epsilon_{krn} \text{ --- } k \end{array} = \begin{array}{c} i \\ | \\ j \end{array} \begin{array}{c} l \\ | \\ k \end{array} + \begin{array}{c} i \text{ --- } l \\ j \text{ --- } k \end{array} \quad (7.14)$$

- Primitiveness Assumption

Suppose  $\mathcal{P} = \{\delta_{ij}, f_{ijk}\}$  where  $f_{ijk}$  is not  $\epsilon_{ijk}$ . For some  $A, B, C, \dots H \in \mathbb{C}$ , one has

$$\text{---} \circ \text{---} = A \text{ ---} \quad (7.15)$$

$$\text{---} \circ \text{---} = B \text{ ---} \bullet \text{---} \quad (7.16)$$

$$\text{---} \circ \text{---} = \left\{ \begin{array}{l} C \text{ ---} \text{ ---} + D \text{ ---} \times \text{---} + E \text{ ---} \bullet \text{---} \bullet \text{---} \\ + F \text{ ---} \text{ ---} + G \text{ ---} \bullet \text{---} \bullet \text{---} + H \text{ ---} \bullet \text{---} \times \text{---} \bullet \text{---} \end{array} \right\} \quad (7.17)$$

---

Let  $\mathcal{P} = (p_1, p_2, \dots, p_k)$  be a **full set of primitive ITs**. By “full”, we mean no others exist.  $\mathcal{P}$  is a basis for an **algebra of invariants**.<sup>1</sup>

An **invariance group**  $\mathcal{G}$  is the set of all linear transformation  $G \in \mathcal{G}$  such that

$$p_1(x^\dagger, y) = p_1(x^\dagger G^\dagger, Gy) \quad (7.18)$$

$$p_2(w^\dagger, x^\dagger, y, z) = p_2(w^\dagger G^\dagger, x^\dagger G^\dagger, Gy, Gz) \quad (7.19)$$

$$\text{etc.} \quad (7.20)$$

Example. Consider an invariance group with a single primitive IT  $p()$  defined by

$$p(x^\dagger, y) = \delta_a^b x^{\dagger a} y_b = x^{\dagger b} y_b \quad (7.21)$$

Then

$$(x')^{\dagger a} (y')_a = x^{\dagger b} (G^\dagger G)_b^c y_c = x^{\dagger b} y_b \quad (7.22)$$

---

<sup>1</sup>An algebra over a field is defined in Sec.A.5

so  $G$  must be unitary

$$G^\dagger G = 1 \tag{7.23}$$

The group of  $n$  dimensional unitary matrices is called  $U(n)$

# Chapter 8

## Lie Algebras

This chapter is based on Ref.[1].

### 8.1 Generators (infinitesimal transformations)

For some group  $\mathcal{G}$ , assume that any group element  $G \in \mathcal{G}$  that is infinitesimal close to the identity 1 can be parametrized by

$$G = 1 + i \sum_i \epsilon_i T^i \quad (8.1)$$

where  $T^i \in \mathbb{C}^{n \times n}$  for  $i = 1, 2, \dots, N$ ,  $\epsilon_i \in \mathbb{R}$  and  $|\epsilon_i| \ll 1$ .

The  $T^i$  matrices are called the **generators** of infinitesimal transformations for group  $\mathcal{G}$ . The generators of a group  $\mathcal{G}$  span a vector space called a Lie algebra  $\mathfrak{g}$ .<sup>1</sup> For example, the generators of the group  $SU(2)$  span the **Lie algebra**  $\mathfrak{su}(2)$ .

Assume that the  $T^i$  matrices are Hermitian and that they satisfy

$$\text{tr}(T^i T^j) = \kappa \delta(i, j) \quad (8.2)$$

A Lie algebra that satisfies Eq.(8.2) is called a **simple Lie algebra**.  $g_{ij} = \text{tr}(T_i^\dagger T_j)$  is called the **Cartan-Killing form**. A **semi-simple Lie algebra** is a direct sum of simple Lie algebras.

It's customary to choose generators so that  $\kappa = \frac{1}{2}$ .<sup>2</sup> However, we will often set  $\kappa = 1$  for intermediate calculations and restore  $\kappa \neq 1$  at the end by dimensional analysis. Just remember that each  $T^j$  scales as  $\sqrt{\kappa}$ . For example, given the equation  $\text{tr}(T^i T^j) = \delta(i, j)$ , we know that when  $\kappa \neq 1$ ,  $\text{tr}(T^i T^j) = \kappa \delta(i, j)$  so both sides of the equation scale as  $\kappa$ .

We will use the following scaled version of  $T^j$  as a birdtrack. Define

---

<sup>1</sup>See Sec.A.5 for the definition of an algebra over a field.

<sup>2</sup>For  $SU(2)$ , it is customary to choose  $T^i = \frac{1}{2}\sigma_i$ , where  $\sigma_i$  for  $i = 1, 2, 3$  are the Pauli matrices. For  $SU(3)$ , it is customary to choose  $T^i = \frac{1}{2}\lambda_i$  where  $\lambda_i$  for  $i = 1, 2, \dots, 8$  are the Gell-Mann matrices. For both of these choices,  $\kappa = \frac{1}{2}$ .

$$(C_{Adj}^i)_b^a = \frac{1}{\sqrt{\kappa}} (T^i)_b^a = \frac{1}{\sqrt{\kappa}} \quad i \text{ --- } T^i \begin{array}{c} \downarrow a \\ \downarrow b \end{array} \quad (8.3)$$

In the CC convention, we will always start reading the indices of this node at the wavy undirected leg.

*Adj* stands the Adjoint. In this node (vertex), an adjoint representation (ad-jrep) particle (wavy line, gluon) is generated (released) by a defining representation (defrep) particle (straight solid line, arrow).

In terms of birdtracks, Eq.(8.2) becomes

$$(T^i)_a^b (T^j)_b^a = \text{tr}(T^i T^j) = \delta(i, j) \quad i \text{ --- } T^i \begin{array}{c} \xrightarrow{\Sigma a} T^j \xrightarrow{\Sigma b} j \end{array} = \leftarrow \bullet \text{ ---} \quad (8.4)$$

We can now define the projection operator for the adrep (gluon exchange between 2 defrep particles)

$$(P_{Adj})_b^a {}_d^c = \sum_i (T^i)_b^a (T^i)_d^c \quad \begin{array}{c} b \xleftarrow{\text{green}} \\ a \nearrow P_{Adj} \searrow d \end{array} \quad = \quad \begin{array}{c} b \uparrow \\ a \text{ --- } \Sigma i \text{ ---} \downarrow d \end{array} \quad (8.5)$$

The green arrow is the first index in the CC convention.

Note that if  $x \in V^n \otimes V^{\dagger n}$ , then

$$(P_{Adj})_b^a {}_d^c x_c^d = \sum_i (T^i)_b^a \underbrace{[(T^i)_d^c x_c^d]}_{\epsilon_i \in \mathbb{R}} \quad (8.6)$$

Recall Eq.(A.26). If  $x \in V^{n^p} \otimes V^{\dagger n^q}$ , and  $\mathbb{G} \in \mathcal{G} \subset GL(n^{p+q}, \mathbb{C})$ ,

$$(x')_{a:p}^{b:q} = \mathbb{G}_{a:p}^{b:q} \text{rev}(c:q) \text{rev}(d:p) x_{d:p}^{c:q}, \quad x'_\alpha = \mathbb{G}_\alpha^\beta x_\beta \quad (8.7)$$

where we define

$$\mathbb{G}_\alpha^\beta \stackrel{\text{def}}{=} \prod_{i=1}^p G_{a_i}^{d_i} \prod_{i=1}^q G^{\dagger b_i}_{c_i} \quad (8.8)$$

If  $\mathbb{G}$  is infinitesimally close to the identity, then we can parametrize it as

$$\mathbb{G}_\alpha^\beta = 1 + i \sum_j \epsilon_j (M^j)_\alpha^\beta \quad (8.9)$$

$$G_{a_i}^{d_i} = 1 + i \sum_j \epsilon_j (T^j)_{a_i}^{d_i} \quad (8.10)$$

$$G^{\dagger b_i}_{c_i} = 1 - i \sum_j \epsilon_j (T^j)^{b_i}_{c_i} \quad (8.11)$$

Define

$$(M^j)_\alpha^\beta = \left[ (T^j)_{a_i}^{d_i} \frac{1}{\delta_{a_i}^{d_i}} - (T^j)^{b_i}_{c_i} \frac{1}{\delta_{c_i}^{b_i}} \right] \delta_{a:p}^{d:p} \delta_{c:q}^{b:q} \quad (8.12)$$

When  $x'_\alpha = x_\alpha$ , to first order in  $\epsilon_i$ ,

$$0 = (M^j)_\alpha^\beta x_\beta = \left[ (T^j)_{a_i}^{d_i} \frac{1}{\delta_{a_i}^{d_i}} - (T^j)^{b_i}_{c_i} \frac{1}{\delta_{c_i}^{b_i}} \right] \delta_{a:p}^{d:p} \delta_{c:q}^{b:q} x_{d:p}^{c:q} \quad (8.13)$$

For example, if we define

$$(M^j)_{a_1 a_2}^{b_1 c_1 d_2 d_1} = (T^j)_{a_1}^{d_1} \delta_{a_2}^{d_2} \delta_{c_1}^{b_1} + \delta_{a_1}^{d_1} (T^j)_{a_2}^{d_2} \delta_{c_1}^{b_1} - \delta_{a_1}^{d_1} \delta_{a_2}^{d_2} (T^j)^{b_1}_{c_1}$$

$$(8.14)$$

then

$$0 = (M^j x)_{a_1 a_2}^{b_1} = \left[ (T^j)_{a_1}^{d_1} \delta_{a_2}^{d_2} \delta_{c_1}^{b_1} + \delta_{a_1}^{d_1} (T^j)_{a_2}^{d_2} \delta_{c_1}^{b_1} - \delta_{a_1}^{d_1} \delta_{a_2}^{d_2} (T^j)^{b_1}_{c_1} \right] x_{d_1 d_2}^{c_1}$$

$$(8.15)$$

## 8.2 Clebsch-Gordan Coefficients

The Clebsch Gordan coefficients (CBC) are introduced in Ch.3. Note that the generators  $(T^i)_a^b$  are a simple kind of CGC matrix, one with

- a gluon (adjrep) particle instead of a general  $\lambda$  rep particle emanating from the  $i$  index,
- a particle of the defrep entering and another leaving the node, instead of any number of defrep particles entering and leaving.

Since  $\mathbb{G} = 1 + i \sum_j \epsilon_j M^j$ , generators decompose in the same way as the group elements

$$M^j = \sum_{\lambda} C_{\lambda}^{\dagger} T_{\lambda}^j C_{\lambda}$$

$$\begin{array}{c} j \\ \vdots \\ \swarrow \quad \searrow \\ \leftarrow M^j \rightarrow \\ \swarrow \quad \searrow \end{array} = \sum_{\lambda} \begin{array}{c} j \\ \vdots \\ \swarrow \quad \searrow \\ \leftarrow C_{\lambda}^{\dagger} \leftarrow T_{\lambda}^j \leftarrow C_{\lambda} \rightarrow \\ \swarrow \quad \searrow \end{array} \quad (8.16)$$

The CGC matrices are invariant matrices.

$$C_{\lambda} = G_{\lambda}^{\dagger} C_{\lambda} G \quad (8.17)$$

Hence,

$$0 = -T_{\lambda}^j C_{\lambda} + C_{\lambda} T^j \quad (8.18)$$

$$0 = \left\{ \begin{array}{l} - \begin{array}{c} j \\ \vdots \\ \swarrow \quad \searrow \\ \leftarrow T_{\lambda}^j \leftarrow C_{\lambda} \leftarrow c_2 \rightarrow b_1 \\ \swarrow \quad \searrow \\ c_1 \end{array} + \begin{array}{c} j \\ \vdots \\ \swarrow \quad \searrow \\ \leftarrow C_{\lambda} \leftarrow c_2 \rightarrow b_1 \\ \swarrow \quad \searrow \\ T^j \leftarrow c_1 \end{array} \\ + \begin{array}{c} j \\ \vdots \\ \swarrow \quad \searrow \\ \leftarrow C_{\lambda} \leftarrow T^j \leftarrow c_2 \rightarrow b_1 \\ \swarrow \quad \searrow \\ c_1 \end{array} - \begin{array}{c} j \\ \vdots \\ \swarrow \quad \searrow \\ \leftarrow C_{\lambda} \leftarrow c_2 \rightarrow T^j \rightarrow b_1 \\ \swarrow \quad \searrow \\ c_1 \end{array} \end{array} \right\} \quad (8.19)$$

Multiplying on the left by  $C_\lambda^\dagger$ , we obtain an expression for the generator  $T_\lambda^i$  in term the generators  $T^j$  (and  $C_\lambda$  CGC matrices).

$$\begin{aligned}
a \leftarrow T_\lambda^j \leftarrow a' &= \underbrace{a \leftarrow T_j \leftarrow C_\lambda \leftarrow C_\lambda^\dagger \leftarrow a' - a \leftarrow C_\lambda \leftarrow C_\lambda^\dagger \leftarrow T_j \leftarrow a'}_{=0} \\
&+ a \leftarrow C_\lambda \leftarrow T_j \leftarrow C_\lambda^\dagger \leftarrow a' + a \leftarrow C_\lambda \leftarrow T_j \leftarrow C_\lambda^\dagger \leftarrow a' - a \leftarrow C_\lambda \leftarrow T_j \leftarrow C_\lambda^\dagger \leftarrow a'
\end{aligned} \tag{8.20}$$

### 8.3 Structure Constants (3 gluon vertex)

$$\underbrace{T^i T^j - T^j T^i}_{[T^i, T^j]} = i f_{ijk} T^k \quad (\text{Lie Algebra commutation relations})$$

$$\begin{aligned}
&a \leftarrow T^i \leftarrow T^j \leftarrow c \quad - \quad a \leftarrow T^j \leftarrow T^i \leftarrow c \quad = i \quad \begin{array}{c} a \leftarrow T^k \leftarrow c \\ \text{\scriptsize } f_{ijk} \\ \begin{array}{cc} i & j \end{array} \end{array}
\end{aligned} \tag{8.21}$$

The  $f_{ijk}$  tensors are called the **structure constants** of the Lie Algebra. They define a 3 gluon vertex in term of the generators  $T^i$ .<sup>3</sup>

If  $(T^j)_a^b$  are the matrix rep (the defrep) of the generators of a group  $\mathcal{G}$ , then Eq.(8.21) shows that the matrices  $(M^k)_{ij} = -i C_{ijk}$  are also a matrix rep (the adrep) of the generators of  $\mathcal{G}$ .

Since  $\text{tr}(T^k T^{k'}) = \delta(k, k')$ , Eq.(8.21) implies

$$\text{tr}([T^i, T^j] T^k) = i f_{ijk} \tag{8.22}$$

<sup>3</sup>It's possible to distinguish between upper and lower gluon indices (i.e., to give the gluon arrows a direction. In that case, the Lie Algebra commutation relations would be  $[T^i, T^j] = f^{ij}_k T^k$  and the gluon indices could be lowered and raised using the metric (called the **Cartan-Killing form**)  $g_{ij} = \text{tr}((T^i)^\dagger T^j)$ . But since we are assuming  $g_{ij} = \kappa \delta_i^j$ , there is no need to do this.

$$if_{ijk} = \text{tr}([T^i, T^j]T^k) = (T^i)_a{}^c (T^j)_b{}^a (T^k)_c{}^b - (T^j)_c{}^b (T^i)_a{}^c (T^k)_b{}^a$$

Note that

$$f_{ijk} = -f_{jik}$$

In fact, the tensor  $f_{ijk}$  is **totally antisymmetric** (i.e., it changes sign under a transposition of any two indices).

**Claim 5**  $f_{ijk}$  is a real number.

**proof:**

$$\left[ i \text{tr}([T^i, T^j]T^k) \right]^\dagger = (-i) \text{tr}(T^k [T^j, T^i]) \quad (8.25)$$

$$= (-i) \text{tr}([T^j, T^i]T^k) \quad (8.26)$$

$$= i \text{tr}([T^j, T^k]T^i) \quad (8.27)$$

**QED**



Note that the birdtrack for the Lie Algebra commutation relations Eq.(8.21) can be understood as the statement that the generators  $T^j$  are invariant matrices. Below we restate Eq.(8.21) to make that obvious

$$0 = \begin{array}{c} i \quad j \\ \text{red wavy} \quad \text{black wavy} \\ a \leftarrow T^i \leftarrow T^j \leftarrow c \end{array} - \begin{array}{c} i \quad j \\ \text{red wavy} \quad \text{black wavy} \\ a \leftarrow T^j \leftarrow T^i \leftarrow c \end{array} - i \begin{array}{c} i \quad j \\ \text{red wavy} \quad \text{black wavy} \\ f_{ijk} \\ \text{black wavy} \\ a \leftarrow T^k \leftarrow c \end{array} \quad (8.28)$$

**Claim 6**

$$f_{ijm}f_{mkl} - f_{ljm}f_{mki} = f_{iml}f_{jkm} \quad (\text{Jacobi identity})$$

$$\begin{array}{c} i \quad l \\ \text{black wavy} \quad \text{black wavy} \\ f_{ijm} \sim \sum m \sim f_{mkl} \\ \text{black wavy} \quad \text{black wavy} \\ j \quad k \end{array} - \begin{array}{c} i \quad l \\ \text{black wavy} \quad \text{black wavy} \\ f_{ijm} \sim \sum m \sim f_{mkl} \\ \text{black wavy} \quad \text{black wavy} \\ j \quad k \end{array} = \begin{array}{c} i \quad l \\ \text{black wavy} \quad \text{black wavy} \\ f_{iml} \\ \sum m \\ f_{jkm} \\ \text{black wavy} \quad \text{black wavy} \\ j \quad k \end{array} \quad (8.29)$$

**proof:**

Note that

$$\text{tr} \left( [[T^i, T^j], T^k] T^l \right) = \text{tr} \left( f_{ijm} [T^m, T^k] \right) \quad (8.30)$$

$$= \text{tr} \left( f_{ijm} f_{mkl} T^{l'} T^l \right) \quad (8.31)$$

$$= f_{ijm} f_{mkl} \quad (8.32)$$

so the Jacobi identity can be restated as

$$\text{tr} \left( \left\{ [[T^i, T^j], T^k] + [[T^j, T^k], T^i] + [[T^k, T^i], T^j] \right\} T^l \right) = 0 \quad (8.33)$$

Hence, the claim follows if we can prove that

$$\underbrace{[[T^i, T^j], T^k] + [[T^j, T^k], T^i] + [[T^k, T^i], T^j]}_{\text{cyclic permutations of } ijk} = 0 \quad (8.34)$$

If we expand the left hand side on Eq.(8.34), we find 6 terms that cancel in pairs.

**QED**

Note Claim 6 can be understood as the Lie Algebra commutation relations Eq.(8.21), but stated in the adrep instead of the defrep. Indeed, if

$$M_{jk}^i = -if_{ijk} \quad (8.35)$$

then Claim 6 becomes

$$(M^i M^l - M^l M^i)_{jk} = iC_{ilm}(M^m)_{jk} \quad (8.36)$$

Note that Claim 6 can be understood as a statement of the fact that  $f_{ijk}$  is an invariant tensor.

$$0 = f_{ijm}f_{mkl} - f_{ljm}f_{mki} - f_{iml}f_{jkm}$$

$$0 = \begin{array}{c} i \\ \text{red wavy} \\ f_{ijm} \sim \sum m \sim f_{mkl} \\ \text{black wavy} \\ j \end{array} \begin{array}{c} l \\ \text{black wavy} \\ f_{mkl} \\ \text{black wavy} \\ k \end{array} - \begin{array}{c} i \\ \text{red wavy} \\ f_{ljm} \sim \sum m \sim f_{mki} \\ \text{black wavy} \\ j \end{array} \begin{array}{c} l \\ \text{black wavy} \\ f_{mki} \\ \text{black wavy} \\ k \end{array} - \begin{array}{c} i \\ \text{red wavy} \\ f_{iml} \\ \text{black wavy} \\ \sum m \\ f_{jkm} \\ \text{black wavy} \\ j \end{array} \begin{array}{c} l \\ \text{black wavy} \\ f_{jkm} \\ \text{black wavy} \\ k \end{array} \quad (8.37)$$

## 8.4 Two types of gluon exchanges

Consider the following two gluon exchange operators. Note that  $\mathbb{P}^2 = \mathbb{P}$ , but  $\mathbb{Q}^2 \neq \mathbb{Q}$ , so  $\mathbb{P}$  is a bonafide projection operator but not  $\mathbb{Q}$ .  $\mathbb{Q}\mathbb{Q}^\dagger = \mathbb{P}$  so  $\mathbb{Q}$  is like half of a projection operator.

$$\mathbb{P}_a^b{}_c^d = \sum_i (T^i)_a^b (T^i)_c^d$$

$$\begin{array}{c} a \quad d \\ \swarrow \quad \searrow \\ \mathbb{P} \\ \swarrow \quad \searrow \\ b \quad c \end{array} = \begin{array}{c} a \quad d \\ \uparrow \quad \downarrow \\ T^i \sim \sum i \sim T^i \\ \uparrow \quad \downarrow \\ b \quad c \end{array} \quad (8.38)$$

$$\mathbb{Q}_a^b{}_\beta^\gamma = \sum_i (T^i)_a^b (T^i)_\beta^\gamma$$

$$\begin{array}{c} a \quad \gamma \\ \swarrow \quad \searrow \\ \mathbb{Q} \\ \swarrow \quad \searrow \\ b \quad \beta \end{array} = \begin{array}{c} a \quad \gamma \\ \uparrow \quad \Downarrow \\ T^i \sim \sum i \sim T^i_\lambda \\ \uparrow \quad \Downarrow \\ b \quad \beta \end{array} \quad (8.39)$$

**Claim 7** If  $\mathbb{Q}_b^a$  is the matrix with  $(\nu, \gamma)$  entries  $\mathbb{Q}_b^a{}_\nu{}^\gamma$ , then

$$[\mathbb{Q}_b^a, \mathbb{Q}_d^c] = \mathbb{P}_{b'}^a{}_d{}^c \mathbb{Q}_b^{b'} - \mathbb{Q}_{a'}^a \mathbb{P}_b^{a'}{}_c{}^d \quad (8.40)$$

**proof:**

$$(T^j)_b^a (T^i)_d^c [T_\lambda^j, T_\lambda^i] = [(T^i)_{b'}^a (T^k)_b^{b'} - (T^k)_{a'}^a (T^i)_b^{a'}] (T^i)_d^c T_\lambda^k \quad (8.41)$$

$$(T^j)_b^a (T^i)_d^c i f_{jik} T_\lambda^k = i f_{ikj} (T^j)_b^a (T^i)_d^c T_\lambda^k \quad (8.42)$$

**QED**

This claim can be visualized as follows.  $\mathbb{Q}$  is an invariant tensor so

$$0 = \left\{ \begin{array}{c} \begin{array}{c} a \searrow \quad b \nearrow \\ i \text{ (red wavy)} \searrow \quad \nearrow \\ \beta \leftarrow T_\lambda^i \leftarrow \mathbb{Q} \leftarrow \gamma \\ a \rightarrow T^i \rightarrow \mathbb{Q} \rightarrow b \\ i \text{ (red wavy)} \nearrow \quad \nwarrow \\ \beta \nwarrow \quad \nearrow \gamma \end{array} \\ - \\ \begin{array}{c} a \searrow \quad b \nearrow \\ i \text{ (red wavy)} \searrow \quad \nearrow \\ \beta \leftarrow \mathbb{Q} \leftarrow T_\lambda^i \leftarrow \gamma \\ a \rightarrow \mathbb{Q} \rightarrow T^i \rightarrow b \\ i \text{ (red wavy)} \nearrow \quad \nwarrow \\ \beta \nwarrow \quad \nearrow \gamma \end{array} \\ + \\ \begin{array}{c} a \searrow \quad b \nearrow \\ i \text{ (red wavy)} \searrow \quad \nearrow \\ \beta \nwarrow \quad \nearrow \gamma \end{array} \end{array} \right\} \quad (8.43)$$

Now multiplying by  $(T^i)_c^d$ , we get

$$\begin{array}{c}
\mathbb{Q}_{d\beta}^c \nu \mathbb{Q}_{b\nu}^a \gamma - \mathbb{Q}_{b\beta}^a \nu \mathbb{Q}_{d\nu}^c \gamma = \mathbb{P}_{a'd}^a c \mathbb{Q}_{b\beta}^{a'} \gamma - \mathbb{Q}_{b'\beta}^a \gamma \mathbb{P}_{b'd}^{b' c} \\
\begin{array}{ccc}
a & & b \\
\searrow & & \nearrow \\
& d & \\
\swarrow & & \searrow \\
c & & \\
\beta \leftarrow \mathbb{Q} \leftarrow \mathbb{Q} \leftarrow \gamma & - & \beta \leftarrow \mathbb{Q} \leftarrow \mathbb{Q} \leftarrow \gamma \\
a \rightarrow T^i \rightarrow \mathbb{Q} \rightarrow b & & a \rightarrow \mathbb{Q} \rightarrow T^j \rightarrow b \\
\downarrow \text{ } \downarrow & & \downarrow \text{ } \downarrow \\
c \rightarrow T^i \rightarrow d & - & c \rightarrow T^j \rightarrow d \\
\downarrow \text{ } \downarrow & & \downarrow \text{ } \downarrow \\
\beta & & \beta \\
& \gamma & \gamma
\end{array}
\end{array} \quad (8.44)$$

## Chapter 9

**Orthogonal Groups: COMING  
SOON**

## Chapter 10

# Quantum Shannon Information Theory: COMING SOON

# Chapter 11

## Recoupling Identities

Mandelstam variables

s-channel, particles shag (have sex),

t-channel, particles have tea

The diagram illustrates three types of particle interactions:

- parallel channels:** Two horizontal lines with arrows pointing to the left, representing two independent particles.
- s-channel:** Two lines with arrows pointing towards a central wavy line, representing an s-channel interaction.
- t-channel:** Two horizontal lines with arrows pointing to the left, with a vertical wavy line connecting them, representing a t-channel interaction.

parallel channels      s-channel      t-channel

(11.1)

### 11.1 Parallel channels to sum of t-channels

The diagram shows a particle interaction with labels  $a, b, c, d, e, f$  and a corresponding loop diagram with labels  $A, B, x$ . Below the diagrams, the text reads:

$\lambda = red, \quad \mu = green, \quad \nu = blue$

(11.2)

no implicit sum over Greek indices

The diagram illustrates the recoupling identity for the parallel channels to sum of t-channels. It shows the following equation:

$$P_\lambda C_{\lambda a}^{\nu b \mu c} = \lambda a \leftarrow P_\lambda C_\lambda^{\nu \mu} = \lambda a \leftarrow P_\lambda C_\lambda^{\nu \mu}$$

(11.3)

$$C_\lambda^{\nu \mu} = P_\lambda C_\lambda^{\nu \mu}$$

(11.4)

$$C_\lambda C_\lambda^\dagger = P_\lambda^2 = P_\lambda \quad (11.5)$$

$$\mathrm{tr}(P_\lambda) = d_\lambda \quad (11.6)$$

where  $d_\lambda$  is the dimension of rep  $\lambda$ . Actually,  $\mathcal{C}_\lambda = P_\lambda C_\lambda = C_\lambda$ , but we make the  $P_\lambda$  explicit for pedagogical purposes.

Note that if we divide  $\mathcal{C}_\lambda$  by  $\sqrt{d_\lambda}$ , then

$$\text{tr} \left( \frac{\mathcal{C}_\lambda}{\sqrt{d_\lambda}} \frac{\mathcal{C}_\lambda^\dagger}{\sqrt{d_\lambda}} \right) = 1 \quad (11.7)$$

$$\mathcal{P}_\lambda = \begin{array}{ccc} & \leftarrow & \leftarrow \\ & \parallel & \parallel \\ \mathcal{C}_\lambda^\dagger & \leftarrow & \mathcal{C}_\lambda \\ & \parallel & \parallel \\ \leftarrow & & \leftarrow \end{array} \quad (11.8)$$

$$\mathcal{P}_\lambda^2 = \mathcal{P}_\lambda \quad (11.9)$$

$$\mathcal{P}_\nu = \frac{d_\nu}{d_\lambda} \quad (11.10)$$

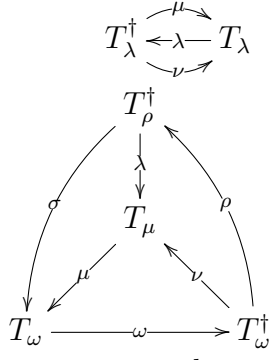
$$\mathcal{P}_\nu^2 = \mathcal{P}_\nu \quad (11.11)$$

$$\mathcal{P}_\mu = \frac{d_\mu}{d_\lambda} \quad (11.12)$$

$$\mathcal{P}_\mu^2 = \mathcal{P}_\mu \quad (11.13)$$

The normalization of the projectors  $\mathcal{P}_\lambda, \mathcal{P}_\nu, \mathcal{P}_\mu$  can be remembered if one takes the denominator  $d_\lambda$  and splits it into two factors of  $\sqrt{d_\lambda}$  and puts one  $\sqrt{d_\lambda}$  under  $\mathcal{C}_\lambda$  and the other under  $\mathcal{C}_\lambda^\dagger$ . Then one “trades”  $\frac{c_\lambda}{\sqrt{d_\lambda}}$  by  $\frac{c_\nu}{\sqrt{d_\nu}}$  or  $\frac{c_\mu}{\sqrt{d_\mu}}$ .





arrow directions for specific case being considered. they can be changed

$$\begin{array}{c} \xleftarrow{\mu} \\ \parallel \\ \lambda \xleftarrow{\text{red}} C_\lambda \\ \parallel \\ \xleftarrow{\nu} \end{array} = \frac{1}{\sqrt{\kappa_\lambda^{\nu\mu}}} \begin{array}{c} \xleftarrow{\mu} \\ \parallel \\ \lambda \xleftarrow{\text{red}} T_\lambda \\ \parallel \\ \xleftarrow{\nu} \end{array} \quad (11.14)$$

$$\begin{array}{c} \xleftarrow{\mu} \\ \parallel \\ \lambda \xleftarrow{\text{red}} T_\lambda \\ \parallel \\ \xleftarrow{\nu} \end{array} \neq \begin{array}{c} \xleftarrow{\mu} \\ \parallel \\ \lambda \xleftarrow{\text{red}} T_\lambda \\ \parallel \\ \xleftarrow{\nu} \end{array} \quad (11.15)$$

$$\begin{array}{c} \xleftarrow{\mu} \\ \circlearrowleft \\ \lambda \xleftarrow{\text{red}} T_\lambda \\ \circlearrowright \\ T_\sigma^\dagger \xleftarrow{\sigma} \end{array} = \kappa_\lambda^{\nu\mu} \begin{array}{c} \xleftarrow{\mu} \\ \bullet \\ \lambda \xleftarrow{\text{red}} T_\lambda \\ \bullet \\ T_\sigma^\dagger \xleftarrow{\sigma} \end{array} \quad (11.16)$$

$$\begin{array}{c} \xleftarrow{\mu} \\ \circlearrowleft \\ T_\lambda^\dagger \xleftarrow{\lambda} T_\lambda \\ \circlearrowright \\ \xleftarrow{\nu} \end{array} = \kappa_\lambda^{\nu\mu} d_\lambda \quad (11.17)$$

$$\mathcal{P}_\lambda = \frac{1}{\kappa_\lambda^{\nu\mu}} \begin{array}{c} \swarrow \quad \nwarrow \\ T_\lambda^\dagger \xleftarrow{\lambda} T_\lambda \\ \swarrow \quad \nwarrow \end{array} \quad (11.18)$$

$$\mathcal{P}_\mu = \frac{1}{\kappa_\mu^{\lambda\nu}} \begin{array}{c} \swarrow \quad \nwarrow \\ T_\mu^\dagger \xleftarrow{\mu} T_\mu \\ \swarrow \quad \nwarrow \end{array} \quad (11.19)$$

$$\mathcal{P}_\nu = \frac{1}{\kappa_\nu^{\mu\lambda}} \quad \begin{array}{c} \swarrow \\ T_\nu^\dagger \leftarrow \nu - T_\nu \\ \searrow \end{array} \quad (11.20)$$

$$\begin{array}{c} \leftarrow \bullet \leftarrow \mu \\ \leftarrow \bullet \leftarrow \nu \end{array} = \sum_\lambda \mathcal{P}_\lambda = \sum_\lambda \frac{d_\lambda}{T_\lambda^\dagger \leftarrow \lambda - T_\lambda} \quad \begin{array}{c} \swarrow \\ T_\lambda^\dagger \leftarrow \lambda - T_\lambda \\ \searrow \end{array} \quad (11.21)$$

$$\begin{array}{c} \leftarrow \\ \leftarrow \\ \leftarrow \\ \leftarrow \end{array} = \sum_{\lambda, \mu, \nu} \frac{1}{\kappa_\lambda} \frac{1}{\kappa_\mu} \frac{1}{\kappa_\nu} \quad \begin{array}{c} \leftarrow \\ \parallel \\ T_\lambda^\dagger \leftarrow \\ \parallel \\ \leftarrow \\ \parallel \\ T_\mu^\dagger \leftarrow \\ \parallel \\ \leftarrow \\ \parallel \\ T_\nu^\dagger \leftarrow T_\nu \\ \parallel \\ \leftarrow \end{array} \quad \begin{array}{c} \leftarrow \\ \parallel \\ T_\lambda \\ \parallel \\ \leftarrow \\ \parallel \\ T_\mu \\ \parallel \\ \leftarrow \\ \parallel \\ T_\nu \\ \parallel \\ \leftarrow \end{array} \quad (11.22)$$

Dropping two upper representation labels from  $\kappa_\nu^{\mu\lambda}$  for succinctness, but their dependence still there.

$$\begin{array}{c} \leftarrow \\ \leftarrow \\ \leftarrow \\ \leftarrow \end{array} = \sum_{\lambda, \mu, \nu} \frac{1}{\kappa_\lambda} \frac{1}{\kappa_\mu} \frac{1}{\kappa_\nu} \quad \begin{array}{c} \leftarrow \\ \parallel \\ T_\lambda^\dagger \leftarrow \\ \parallel \\ \leftarrow \\ \parallel \\ T_\nu^\dagger \leftarrow T_\nu \\ \parallel \\ \leftarrow \\ \parallel \\ T_\mu^\dagger \leftarrow \\ \parallel \\ \leftarrow \end{array} \quad \begin{array}{c} \leftarrow \\ \parallel \\ T_\lambda \\ \parallel \\ \leftarrow \\ \parallel \\ T_\mu \\ \parallel \\ \leftarrow \end{array} \quad (11.23)$$

## 11.2 t-channel to sum of s-channels

$$\begin{array}{c} \leftarrow \sigma - T_\mu^\dagger \leftarrow \mu \\ \downarrow \omega \\ \leftarrow \rho - T_\rho \leftarrow \nu \end{array} = \sum_\lambda \left[ * \begin{array}{c} \overbrace{T_\lambda^\dagger \leftarrow \lambda - T_\lambda}^{d_\lambda \overrightarrow{\sigma}} \quad \overbrace{T_\lambda^\dagger \leftarrow \lambda - T_\lambda}^{d_\lambda \overrightarrow{\mu}} \\ \downarrow \rho \quad \downarrow \nu \\ T_\lambda^\dagger \leftarrow \lambda - T_\lambda \end{array} \right] \quad (11.24)$$

$$= \sum_\lambda \left[ S_\lambda \left( \begin{array}{c} \leftarrow \sigma - T_\mu^\dagger \leftarrow \mu \\ \downarrow \omega \\ \leftarrow \rho - T_\rho \leftarrow \nu \end{array} \right) \begin{array}{c} \leftarrow \sigma - T_\lambda^\dagger \leftarrow \lambda - T_\lambda \\ \downarrow \rho \quad \downarrow \nu \\ T_\lambda^\dagger \leftarrow \lambda - T_\lambda \end{array} \right] \quad (11.25)$$

$$S_\lambda \left( \begin{array}{c} \leftarrow \sigma - T_\mu^\dagger \leftarrow \mu \\ \downarrow \omega \\ \leftarrow \rho - T_\rho \leftarrow \nu \end{array} \right) = \frac{d_\lambda}{T_\lambda^\dagger \leftarrow \lambda - T_\lambda} \frac{d_\lambda}{T_\lambda^\dagger \leftarrow \lambda - T_\lambda} \frac{\begin{array}{c} T_\mu^\dagger \\ \downarrow \omega \\ T_\rho \\ \swarrow \sigma \quad \searrow \mu \\ T_\lambda \quad \xrightarrow{\lambda} \quad T_\lambda^\dagger \end{array}}{d_\lambda} \quad (11.26)$$

$$= d_\lambda \frac{\begin{array}{c} T_\mu^\dagger \\ \downarrow \omega \\ T_\rho \\ \swarrow \sigma \quad \searrow \mu \\ T_\lambda \quad \xrightarrow{\lambda} \quad T_\lambda^\dagger \end{array}}{T_\lambda^\dagger \leftarrow \lambda - T_\lambda \quad T_\lambda^\dagger \leftarrow \lambda - T_\lambda} \quad (11.27)$$

### 11.3 Wigner $3n - j$ coefficients

we will refer to  $x$  as a  $3 - j$  symbol, and to  $x$  as a  $6 - j$  symbol. Atomic physicists also define  $3n - j$  symbols, for  $n = 1, 2, 3, \dots$ . They are called that because they describe how to “add”  $3n$  angular momenta  $j$ . There is only one  $3 - j$  but two  $6 - j$ 's . five  $9 - j$ s, and so on. We only show one  $3 - j$  and one  $6 - j$ .

Similar to Eq.(11.22) but here we use the most general t-channel to sum of s-channels conversion Eq.(11.25)

$$= \sum_{\lambda, \mu, \nu} S_\lambda S_\mu S_\nu \quad (11.28)$$

$$= \sum_{\lambda, \mu} S_\lambda S_\mu \quad (11.29)$$

$$3n - j \text{ coefficient} = \sum (\text{product of } S_\lambda \text{'s}) (\text{tree graph with one-point loop}) \quad (11.30)$$

4-point loop  $\rightarrow$  one-point loop  
 one-point loop , self energy loop?  
 zero point loop = vacuum bubble

## Chapter 12

### Recoupling Identities for $U(n)$

# Chapter 13

## Reducibility of Representations

This chapter is based on Ref.[1].

### 13.1 Eigenvalue Projectors

Suppose  $M \in \mathbb{C}^{d \times d}$  has eigenvalues  $\lambda_i$  with corresponding eigenvectors  $|\lambda_i\rangle$

$$M|\lambda_i\rangle = \lambda_i|\lambda_i\rangle \quad (13.1)$$

for  $i \in \mathbb{Z}_{[1,r]}$ . The characteristic polynomial of  $M$  is defined as

$$cp(\lambda) \stackrel{\text{def}}{=} \det(M - \lambda) = \prod_{i=1}^r (\lambda - \lambda_i)^{d_i} \quad (13.2)$$

It must satisfy

$$cp(\lambda) = 0 \quad (13.3)$$

Note that if  $M$  is Hermitian ( $M^\dagger = M$ ), then all its eigenvalues are real. (because  $\lambda_i = \langle \lambda_i | M | \lambda_i \rangle \in \mathbb{R}$ )

If  $M$  is a Hermitian, then there exists a matrix  $C$  that is a unitary ( $CC^\dagger = C^\dagger C = 1$ ) and diagonalizes  $M$

$$CMC^\dagger = \begin{bmatrix} D_{\lambda_1} & 0 & 0 & 0 \\ 0 & D_{\lambda_2} & 0 & 0 \\ 0 & 0 & \ddots & 0 \\ 0 & 0 & 0 & D_{\lambda_r} \end{bmatrix} \quad (13.4)$$

where

$$D_{\lambda_i} = \text{diag}(\underbrace{\lambda_i, \lambda_i, \dots, \lambda_i}_{d_i \text{ times}}) \quad (13.5)$$

$$d = \sum_{i=1}^r d_i \quad (13.6)$$

For example, when  $d = 2$ ,

$$CMC^\dagger = \begin{bmatrix} \lambda_1 & 0 \\ 0 & \lambda_2 \end{bmatrix} \quad (13.7)$$

Note that for  $d = 2$ ,

$$CP_1C^\dagger = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} = \frac{CMC^\dagger - \lambda_2}{\lambda_1 - \lambda_2} \quad (13.8)$$

$$CP_2C^\dagger = \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix} = \frac{CMC^\dagger - \lambda_1}{\lambda_2 - \lambda_1} \quad (13.9)$$

$P_1$  and  $P_2$  are a set of complete orthogonal projection operators

$$P_1 + P_2 = 1 \quad (13.10)$$

$$P_1^2 = P_1, P_2^2 = P_2, P_1P_2 = P_2P_1 = 0 \quad (13.11)$$

Similarly, for  $d > 2$ , we can define one projection operator  $P_i$  for each eigenvalue  $\lambda_i$ . If  $I^{d_i \times d_i}$  is the  $d_i$  dimensional unit matrix, then

$$P_i = C^\dagger \text{diag}(0, \dots, 0, I^{d_i \times d_i}, 0, \dots, 0)C \quad (13.12)$$

$$= \prod_{j \neq i} \frac{M - \lambda_j}{\lambda_i - \lambda_j} \quad (13.13)$$

As for  $d = 2$ , the  $P_i$  just defined are a complete set of orthogonal projection operators:

$$\sum_{i=1}^r P_i = 1 \quad (\text{completeness}) \quad (13.14)$$

$$P_i P_j = P_i \delta(i, j) \quad (\text{orthonormality}) \quad (13.15)$$

for all  $i, j \in \mathbb{Z}_{[1, r]}$

Note that

$$d_i = \text{tr}[C^\dagger P_i C] \quad (13.16)$$

$$= \text{tr} P_i \quad (13.17)$$

Note that the  $P_i$ 's are Hermitian ( $P_i^\dagger = P_i$ ) because  $M$  is Hermitian and its eigenvalues are real.

## 13.2 $[P_i, M] = 0$ consequences

Note that for any  $i$ ,  $P_i$  and  $M$  commute

$$[P_i, M] = P_i M - M P_i = 0 \quad (13.18)$$

From the  $P_i$ 's completeness and commutativity with  $M$ , we get

$$M = \sum_{i=1}^r \sum_{j=1}^r P_i M P_j \quad (13.19)$$

$$= \sum_{i=1}^r P_i M P_i \quad (13.20)$$

**Claim 8** For all  $i$ ,

$$M P_i = \lambda_i P_i \text{ (no } i \text{ sum)} \quad (13.21)$$

**proof:** We only show it for  $d = 2$

$$C M P_1 C^\dagger = \begin{bmatrix} \lambda_1 & 0 \\ 0 & \lambda_2 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \quad (13.22)$$

$$= \lambda_1 \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \quad (13.23)$$

$$= \lambda_i C P_i C^\dagger \quad (13.24)$$

**QED**

From the last claim, it immediately follows that if  $f(x)$  can be expressed as a power series in  $x$ , then <sup>1</sup>

$$f(M) P_i = f(\lambda_i) P_i \text{ (no } i \text{ sum)} \quad (13.25)$$

Suppose  $M^{(1)}, M^{(2)} \in \mathbb{C}^{d \times d}$  are Hermitian matrices that commute

$$[M^{(1)}, M^{(2)}] = 0 \quad (13.26)$$

Use  $M^{(1)}$  to decompose  $V = \mathbb{C}^{d \times d}$  into a direct sum of vector spaces  $\oplus_i V_i$ . Then we can use  $M^{(2)}$  to decompose  $V_i$  into  $\oplus_j V_{i,j}$ . If  $M^{(1)}$  and  $M^{(2)}$  don't commute, let  $P_i^{(1)}$  be an eigenvalue projection operator of  $M^{(1)}$ . Then replace  $M^{(2)}$  by  $P_i^{(1)} M^{(2)} P_i^{(1)}$ . Now

$$[M^{(1)}, P_i^{(1)} M^{(2)} P_i^{(1)}] = 0 \quad (13.27)$$

---

<sup>1</sup> $M$  must also satisfy some convergence conditions that we won't get into.



### 13.3 $[G, M] = 0$ consequences

An invariant matrix (see Ch.7) commutes with all the elements  $G$  of a group  $\mathcal{G}$

$$[G, M] = 0 \quad (13.28)$$

If  $P_i$  are the projection operators of  $M$ , then  $P_i = f_i(M)$  so

$$[G, P_i] = 0 \quad (13.29)$$

for all  $G \in \mathcal{G}$  and  $i$ .

$$G = 1G1 = \sum_i \sum_j P_i G P_j = \sum_j \underbrace{P_j G P_j}_{\stackrel{\text{def}}{=} G_j} \quad (13.30)$$

**Claim 9**

$$G = C^\dagger \text{diag}(G_1, G_2, \dots) C \quad (13.31)$$

$$G = \sum_i C_i^\dagger G_i C_i \quad (13.32)$$

where the matrices  $C_i$  are the Clebsch Gordan matrices of  $M$  (see Ch. 3)

**proof:**

$$C_i G C_i^\dagger = \sum_j C_i P_j G P_j C_i^\dagger = C_i G_i C_i^\dagger = G_i \quad (13.33)$$

**QED**

A **representation (rep)**  $G_i$  acts only on a  $d_i$  dimensional vector space  $V^{d_i} = P_i V^d$ . In this way, an invariant matrix  $M \in \mathbb{C}^{d \times d}$  with  $r$  distinct eigenvalues, induces a decomposition of  $V^d$  into a direct sum of vector spaces

$$V^d \xrightarrow{M} V_1^{d_1} \oplus V_2^{d_2} \oplus \dots \oplus V_r^{d_r} \quad (13.34)$$

If a representation  $G_i$  cannot itself be reduced further, it is said to be an **irreducible representation (irrep)**.

Note that sometimes the term representation is used to refer to the vector space  $V_i^{d_i}$  instead of the matrix  $G_i$ .

We've considered the decomposition of  $V^d$  into irreps. An example of such a decomposition is the decomposition of  $V^n \otimes V^{\dagger n}$

$$\begin{aligned}
1 &= \frac{1}{n} \uparrow\downarrow + P_{Adj} + \sum_{\lambda \neq Adj} P_{\lambda}, \\
\delta_d^a \delta_d^c &= \frac{1}{n} \delta_b^a \delta_d^c + (P_{Adj})_{a \ c}^{\ b \ d} + \sum_{\lambda \neq Adj} (P_{\lambda})_{a \ c}^{\ b \ d} \\
\begin{array}{c} a \leftarrow d \\ b \rightarrow c \end{array} &= \frac{1}{n} \begin{array}{c} \uparrow \\ \downarrow \end{array} + \begin{array}{c} \uparrow \\ \text{---} \\ \downarrow \end{array} + \sum_{\lambda \neq Adj} \begin{array}{c} \uparrow \\ \text{---} \\ \downarrow \end{array}
\end{aligned} \tag{13.35}$$

## Chapter 14

**Spinors: COMING SOON**

## Chapter 15

### Squashed Entanglement: COMING SOON

## Chapter 16

**Symplectic Groups: COMING  
SOON**

# Chapter 17

## Symmetrization and Antisymmetrization

This chapter is based on Ref.[1]

As preparation for this chapter, read Sec.A.7.

### 17.1 Symmetrizer

The set of permutations of 2 elements can be represented by the following  $2! = 2$  birdtracks<sup>1</sup>

$$\mathbb{1}_{a_1, a_2}^{b_2, b_1} = \delta_{a_1}^{b_1} \delta_{a_2}^{b_2} = \begin{array}{c} a_1 \xleftarrow{\text{green}} \bullet \xleftarrow{\text{black}} b_1 \\ a_2 \xleftarrow{\text{black}} \bullet \xleftarrow{\text{black}} b_2 \end{array} \quad (17.1)$$

$$(\sigma_{(1,2)})_{a_1, a_2}^{b_2, b_1} = \delta_{a_1}^{b_2} \delta_{a_2}^{b_1} = \begin{array}{c} a_1 \xleftarrow{\text{green}} \bullet \xleftarrow{\text{black}} b_1 \\ \updownarrow \\ a_2 \xleftarrow{\text{black}} \bullet \xleftarrow{\text{black}} b_2 \end{array} \quad (17.2)$$

The set of permutations of 3 elements can be represented by the following  $3! = 6$  birdtracks:

$$\begin{array}{c} a_1 \xleftarrow{\text{green}} \bullet \xleftarrow{\text{black}} b_1 \\ \mathbb{1} = a_2 \xleftarrow{\text{black}} \bullet \xleftarrow{\text{black}} b_2 \\ a_3 \xleftarrow{\text{black}} \bullet \xleftarrow{\text{black}} b_3 \end{array} \quad (17.3)$$

$$\sigma_{(1,2)} = \begin{array}{c} \xleftarrow{\text{black}} \bullet \\ \updownarrow \\ \xleftarrow{\text{black}} \bullet \\ \xleftarrow{\text{black}} \end{array} \quad \sigma_{(2,3)} = \begin{array}{c} \xleftarrow{\text{black}} \\ \xleftarrow{\text{black}} \bullet \xleftarrow{\text{black}} \\ \updownarrow \\ \xleftarrow{\text{black}} \bullet \xleftarrow{\text{black}} \end{array} \quad \sigma_{(1,3)} = \begin{array}{c} \xleftarrow{\text{black}} \bullet \\ \updownarrow \\ \xleftarrow{\text{black}} \bullet \\ \xleftarrow{\text{black}} \bullet \end{array} \quad (17.4)$$

---

<sup>1</sup>Note that the set of values that  $a_i$  and  $b_i$  can assume can be anything, as long as, for some set  $V$ ,  $val(\underline{a}_i) = val(\underline{b}_i) = V$  for all  $i$ .

$$\sigma_{(1,2,3)} = \begin{array}{c} \leftarrow \bullet \leftarrow \\ \updownarrow \\ \leftarrow \bullet \leftarrow \bullet \leftarrow \\ \updownarrow \\ \leftarrow \bullet \leftarrow \end{array} = \begin{array}{c} \leftarrow \bullet \leftarrow \\ \updownarrow \\ \leftarrow \bullet \leftarrow \bullet \leftarrow \\ \updownarrow \\ \leftarrow \bullet \leftarrow \bullet \leftarrow \end{array} \quad (17.5)$$

$$\sigma_{(1,3,2)} = \begin{array}{c} \leftarrow \bullet \leftarrow \\ \updownarrow \\ \leftarrow \bullet \leftarrow \bullet \leftarrow \\ \updownarrow \\ \leftarrow \bullet \leftarrow \end{array} = \begin{array}{c} \leftarrow \bullet \leftarrow \bullet \leftarrow \\ \updownarrow \\ \leftarrow \bullet \leftarrow \bullet \leftarrow \\ \updownarrow \\ \leftarrow \bullet \leftarrow \end{array} \quad (17.6)$$

The  $p$ -element symmetrizer  $\mathcal{S}_p$  is defined as the birdtrack

$$\begin{array}{c} \leftarrow \mathcal{S}_p \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} = \frac{1}{p!} \left\{ \begin{array}{c} \leftarrow \\ \leftarrow \\ \leftarrow \\ \vdots \\ \leftarrow \end{array} + \begin{array}{c} \leftarrow \bullet \leftarrow \\ \updownarrow \\ \leftarrow \bullet \leftarrow \\ \leftarrow \\ \vdots \\ \leftarrow \end{array} + \dots \right\} \quad (17.7)$$

Note that  $\mathcal{S}_p$  satisfies the following identities

$$\mathcal{S}_p^2 = \mathcal{S}_p \quad \begin{array}{c} \leftarrow \mathcal{S}_p \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} \quad \begin{array}{c} \leftarrow \mathcal{S}_p \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} = \begin{array}{c} \leftarrow \mathcal{S}_p \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} \quad (17.8)$$

$$\mathcal{S}_p \mathcal{S}_{[1,q]} = \mathcal{S}_p \quad \begin{array}{c} \leftarrow \mathcal{S}_p \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} \quad \begin{array}{c} \leftarrow \mathcal{S}_{[1,q]} \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} = \begin{array}{c} \leftarrow \mathcal{S}_p \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} \quad (17.9)$$

$$\boxed{\mathcal{S}_p \sigma_{(1,2)} = \mathcal{S}_p} \quad \begin{array}{c} \leftarrow \mathcal{S}_p \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} \quad \begin{array}{c} \leftarrow \bullet \leftarrow \\ \updownarrow \\ \leftarrow \bullet \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} = \begin{array}{c} \leftarrow \mathcal{S}_p \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} \quad (17.10)$$

**Claim 10**

$$\begin{array}{c} \leftarrow \mathcal{S}_p \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} = \frac{1}{p} \left( \begin{array}{c} \leftarrow \leftarrow \\ \leftarrow \mathcal{S}_{p-1} \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} + (p-1) \begin{array}{c} \leftarrow \bullet \leftarrow \leftarrow \leftarrow \\ \updownarrow \\ \leftarrow \mathcal{S}_{p-1} \leftarrow \bullet \leftarrow \mathcal{S}_{p-1} \leftarrow \\ \parallel \quad \parallel \\ \leftarrow \leftarrow \quad \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \quad \leftarrow \leftarrow \end{array} \right) \quad (17.11)$$

**proof:** We only prove it for  $p = 3$ .

$$3! \begin{array}{c} \leftarrow \mathcal{S}_3 \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \end{array} = \left( \begin{array}{c} \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \\ \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \bullet \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \bullet \leftarrow \leftarrow \\ \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \\ \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \end{array} + \begin{array}{c} \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \\ \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \\ \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \\ \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \end{array} \right) \quad (17.12)$$

$$2! \begin{array}{c} \leftarrow \leftarrow \\ \leftarrow \mathcal{S}_2 \leftarrow \\ \parallel \\ \leftarrow \leftarrow \end{array} = \left( \begin{array}{c} \leftarrow \leftarrow \quad \leftarrow \leftarrow \\ \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \\ \leftarrow \leftarrow \quad \leftarrow \bullet \leftarrow \leftarrow \end{array} \right) \quad (17.13)$$



$$3! \begin{array}{c} \leftarrow \mathcal{S}_3 \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \end{array} - 2! \begin{array}{c} \leftarrow \mathcal{S}_2 \leftarrow \\ \parallel \\ \leftarrow \leftarrow \end{array} = \left( \begin{array}{cc} \begin{array}{c} \leftarrow \bullet \leftarrow \\ \downarrow \\ \leftarrow \bullet \leftarrow \bullet \leftarrow \\ \downarrow \\ \leftarrow \bullet \leftarrow \end{array} & + \begin{array}{c} \leftarrow \bullet \leftarrow \\ \downarrow \\ \leftarrow \bullet \leftarrow \bullet \leftarrow \\ \downarrow \\ \leftarrow \bullet \leftarrow \end{array} \\ + \begin{array}{c} \leftarrow \bullet \leftarrow \\ \downarrow \\ \leftarrow \bullet \leftarrow \\ \downarrow \\ \leftarrow \bullet \leftarrow \end{array} & + \begin{array}{c} \leftarrow \bullet \leftarrow \\ \downarrow \\ \leftarrow \bullet \leftarrow \\ \downarrow \\ \leftarrow \bullet \leftarrow \end{array} \end{array} \right) \quad (17.14)$$

$$= 2!2! \begin{array}{c} \leftarrow \bullet \leftarrow \\ \downarrow \\ \leftarrow \mathcal{S}_2 \leftarrow \bullet \leftarrow \mathcal{S}_2 \leftarrow \\ \parallel \quad \parallel \\ \leftarrow \leftarrow \leftarrow \leftarrow \end{array} \quad (17.15)$$

**QED**

Tracing over the identity of Claim 10, we get

$$\begin{array}{c} \leftarrow \mathcal{S}_p \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} = \frac{1}{p} \left( \begin{array}{cc} \begin{array}{c} \leftarrow \mathcal{S}_{p-1} \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} & + (p-1) \begin{array}{c} \leftarrow \bullet \leftarrow \\ \downarrow \\ \leftarrow \mathcal{S}_{p-1} \leftarrow \bullet \leftarrow \mathcal{S}_{p-1} \leftarrow \\ \parallel \quad \parallel \\ \leftarrow \leftarrow \leftarrow \leftarrow \end{array} \end{array} \right) \quad (17.16)$$

$$= \frac{n+p-1}{p} \left( \begin{array}{c} \leftarrow \mathcal{S}_{p-1} \leftarrow \\ \parallel \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} \right) \quad (17.17)$$

Hence

$$\text{tr}_{\underline{a}_1} \mathcal{S}_p = \frac{n+p-1}{p} \mathcal{S}_{p-1} \quad (17.18)$$

$$\text{tr}_{\underline{a}_1, \underline{a}_2, \dots, \underline{a}_k} \mathcal{S}_p = \frac{(n+p-1)(n+p-2) \dots (n=p-k)}{p(p-1) \dots (p-k+1)} \mathcal{S}_{p-k} \quad (17.19)$$

$$d_{\mathcal{S}_p} = \text{tr}_{\underline{a}^p} \mathcal{S}_p = \frac{(n+p-1)!}{p!(n-1)!} = \binom{n+p-1}{p} \quad (17.20)$$

For  $p = 2$ ,

$$d_{\mathcal{S}_2} = \frac{(n+1)n}{2} \quad (17.21)$$

## 17.2 Antisymmetrizer

The  $p$ -element antisymmetrizer  $\mathcal{A}_p$  is defined as the birdtrack

$$\begin{array}{c} \leftarrow \mathcal{A}_p \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \parallel \leftarrow \end{array} = \frac{1}{p!} \left\{ \begin{array}{cc} \leftarrow & \leftarrow \bullet \leftarrow \\ \leftarrow & \leftarrow \bullet \leftarrow \\ \leftarrow & \leftarrow \\ \vdots & \vdots \\ \leftarrow & \leftarrow \end{array} - \begin{array}{cc} \leftarrow & \leftarrow \\ \leftarrow & \leftarrow \\ \leftarrow & \leftarrow \\ \vdots & \vdots \\ \leftarrow & \leftarrow \end{array} + \dots \right\} \quad (17.22)$$

Note that  $\mathcal{A}_p$  satisfies the following identities

$$\begin{array}{c} \leftarrow \mathcal{A}_p \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \parallel \leftarrow \end{array} \begin{array}{c} \leftarrow \mathcal{A}_p \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \parallel \leftarrow \end{array} = \begin{array}{c} \leftarrow \mathcal{A}_p \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \parallel \leftarrow \end{array} \quad (17.23)$$

$$\begin{array}{c} \leftarrow \mathcal{A}_p \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \parallel \leftarrow \end{array} \begin{array}{c} \leftarrow \mathcal{A}_{[1,q]} \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \parallel \leftarrow \end{array} = \begin{array}{c} \leftarrow \mathcal{A}_p \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \parallel \leftarrow \end{array} \quad (17.24)$$

$$\begin{array}{c} \leftarrow \mathcal{A}_p \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \parallel \leftarrow \end{array} \begin{array}{c} \leftarrow \bullet \leftarrow \\ \leftarrow \bullet \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \leftarrow \end{array} = (-1) \begin{array}{c} \leftarrow \mathcal{A}_p \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \parallel \leftarrow \end{array} \quad (17.25)$$

$$\begin{array}{c}
\boxed{\mathcal{S}_p \mathcal{A}_q = \mathcal{A}_p \mathcal{S}_q = 0} \\
\begin{array}{ccc}
\leftarrow \mathcal{S}_p \leftarrow & \leftarrow \mathcal{A}_p \leftarrow & \\
\leftarrow \parallel \leftarrow & \leftarrow \parallel \leftarrow & \\
\leftarrow \leftarrow & \leftarrow \leftarrow & \\
\vdots & \vdots & \vdots \\
\leftarrow \parallel \leftarrow & \leftarrow \parallel \leftarrow &
\end{array} = 0
\end{array} \quad (17.26)$$

$$\begin{array}{c}
\boxed{\mathcal{S}_p \mathcal{A}_{[1,q]} = \mathcal{A}_p \mathcal{S}_{[1,q]} = 0} \\
\begin{array}{ccc}
\leftarrow \mathcal{S}_p \leftarrow & \leftarrow \mathcal{A}_{[1,q]} \leftarrow & \\
\leftarrow \parallel \leftarrow & \leftarrow \parallel \leftarrow & \\
\leftarrow \leftarrow & \leftarrow \leftarrow & \\
\vdots & \vdots & \vdots \\
\leftarrow \parallel \leftarrow & \leftarrow \parallel \leftarrow &
\end{array} = \begin{array}{ccc}
\leftarrow \mathcal{A}_p \leftarrow & \leftarrow \mathcal{S}_{[1,q]} \leftarrow & \\
\leftarrow \parallel \leftarrow & \leftarrow \parallel \leftarrow & \\
\leftarrow \leftarrow & \leftarrow \leftarrow & \\
\vdots & \vdots & \vdots \\
\leftarrow \parallel \leftarrow & \leftarrow \parallel \leftarrow &
\end{array} = 0
\end{array} \quad (17.27)$$

**Claim 11**

$$\begin{array}{c}
\leftarrow \mathcal{A}_p \leftarrow \\
\leftarrow \parallel \leftarrow \\
\leftarrow \leftarrow \\
\vdots \\
\leftarrow \parallel \leftarrow
\end{array} = \frac{1}{p} \left( \begin{array}{c} \leftarrow \mathcal{A}_{p-1} \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \leftarrow \\ \vdots \\ \leftarrow \parallel \leftarrow \end{array} + (-1)(p-1) \begin{array}{c} \leftarrow \bullet \leftarrow \\ \leftarrow \mathcal{A}_{p-1} \leftarrow \bullet \leftarrow \mathcal{A}_{p-1} \leftarrow \\ \leftarrow \parallel \leftarrow \parallel \leftarrow \\ \vdots \\ \leftarrow \parallel \leftarrow \parallel \leftarrow \end{array} \right) \quad (17.28)$$

**proof:** We only prove it for  $p = 3$ .

$$3! \begin{array}{c} \leftarrow \mathcal{A}_3 \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \leftarrow \end{array} = \left( \begin{array}{c} \leftarrow \leftarrow \leftarrow + \leftarrow \bullet \leftarrow \leftarrow + \leftarrow \bullet \leftarrow \leftarrow \\ \leftarrow \leftarrow \leftarrow + \leftarrow \bullet \leftarrow \bullet \leftarrow + \leftarrow \bullet \leftarrow \bullet \leftarrow \\ \leftarrow \leftarrow \leftarrow + \leftarrow \bullet \leftarrow \bullet \leftarrow \\ \vdots \\ \leftarrow \leftarrow \leftarrow + \leftarrow \bullet \leftarrow \bullet \leftarrow \\ - \leftarrow \bullet \leftarrow \leftarrow - \leftarrow \bullet \leftarrow \bullet \leftarrow - \leftarrow \bullet \leftarrow \bullet \leftarrow \\ \vdots \\ \leftarrow \leftarrow \leftarrow + \leftarrow \bullet \leftarrow \bullet \leftarrow \\ - \leftarrow \bullet \leftarrow \bullet \leftarrow - \leftarrow \bullet \leftarrow \bullet \leftarrow \\ \leftarrow \leftarrow \leftarrow + \leftarrow \bullet \leftarrow \bullet \leftarrow \end{array} \right) \quad (17.29)$$

$$2! \begin{array}{c} \leftarrow \\ \leftarrow \mathcal{A}_2 \leftarrow \\ \parallel \\ \leftarrow \end{array} = \left( \begin{array}{cc} \leftarrow & \leftarrow \\ \leftarrow & - \leftarrow \bullet \leftarrow \\ \leftarrow & \leftarrow \bullet \leftarrow \end{array} \right) \quad (17.30)$$

$$3! \begin{array}{c} \leftarrow \mathcal{A}_3 \leftarrow \\ \parallel \\ \leftarrow \end{array} - 2! \begin{array}{c} \leftarrow \\ \leftarrow \mathcal{A}_2 \leftarrow \\ \parallel \\ \leftarrow \end{array} = \left( \begin{array}{cc} \begin{array}{c} \bullet \leftarrow \\ \updownarrow \\ \bullet \leftarrow \bullet \leftarrow \\ \updownarrow \\ \bullet \leftarrow \end{array} & + \begin{array}{c} \leftarrow \bullet \leftarrow \\ \updownarrow \\ \bullet \leftarrow \bullet \leftarrow \\ \updownarrow \\ \bullet \leftarrow \end{array} \\ - \begin{array}{c} \bullet \leftarrow \\ \updownarrow \\ \bullet \leftarrow \end{array} & - \begin{array}{c} \bullet \leftarrow \\ \updownarrow \\ \bullet \leftarrow \end{array} \end{array} \right) \quad (17.31)$$

$$= (-1)2!2! \begin{array}{c} \leftarrow \bullet \leftarrow \\ \updownarrow \\ \leftarrow \mathcal{A}_2 \leftarrow \bullet \leftarrow \mathcal{A}_2 \leftarrow \\ \parallel \qquad \parallel \\ \leftarrow \qquad \leftarrow \end{array} \quad (17.32)$$

**QED**

Tracing over the identity of Claim 11, we get

$$\begin{array}{c} \leftarrow \mathcal{A}_p \leftarrow \\ \parallel \\ \leftarrow \\ \leftarrow \\ \vdots \\ \leftarrow \end{array} = \frac{1}{p} \left( \begin{array}{cc} \begin{array}{c} \leftarrow \\ \leftarrow \mathcal{A}_{p-1} \leftarrow \\ \parallel \\ \leftarrow \\ \vdots \\ \leftarrow \end{array} & + (-1)(p-1) \begin{array}{c} \leftarrow \bullet \leftarrow \\ \updownarrow \\ \leftarrow \mathcal{A}_{p-1} \leftarrow \bullet \leftarrow \mathcal{A}_{p-1} \leftarrow \\ \parallel \qquad \parallel \\ \leftarrow \qquad \leftarrow \end{array} \end{array} \right) \quad (17.33)$$

$$= \frac{n + (-1)(p-1)}{p} \left( \begin{array}{c} \leftarrow \mathcal{A}_{p-1} \leftarrow \\ \parallel \\ \leftarrow \\ \vdots \\ \leftarrow \end{array} \right) \quad (17.34)$$

Hence,

$$\mathrm{tr}_{\underline{a}_1} \mathcal{A}_p = \frac{n-p+1}{p} \mathcal{A}_{p-1} \quad (17.35)$$

$$\mathrm{tr}_{\underline{a}_1, \underline{a}_2, \dots, \underline{a}_k} \mathcal{A}_p = \frac{(n-p+1)(n-p+2) \dots (n-p+k)}{p(p-1) \dots (p-k+1)} \mathcal{A}_{p-k} \quad (17.36)$$

$$d_{\mathcal{A}_p} = \text{tr}_{\underline{a}^p} \mathcal{A}_p = \frac{\prod_{i=n-p+1}^n i}{p!} \quad (17.37)$$

$$= \frac{\prod_{i=n}^{n-p+1} i}{p!} \quad (17.38)$$

$$= \begin{cases} \frac{n!}{p!(n-p)!} = \binom{n}{p} & \text{if } p \leq n \\ 0 & \text{otherwise} \end{cases} \quad (17.39)$$

For  $p = 2 \leq n$ ,

$$d_{\mathcal{A}_2} = \binom{n}{2} \quad (17.40)$$

$$\mathcal{A}_p = 0 \text{ if } n < p \quad (17.41)$$

For example, for  $n = 2$  and  $p = 3$

$$\mathcal{A}_3 = \frac{1}{6} \left( \begin{array}{c} |a\rangle \quad |a\rangle \quad |b\rangle \\ \downarrow \quad \downarrow \quad \downarrow \\ \bullet \quad \bullet \quad \bullet \\ \downarrow \quad \downarrow \quad \downarrow \end{array} + \begin{array}{c} |a\rangle \quad |a\rangle \quad |b\rangle \\ \downarrow \quad \downarrow \quad \downarrow \\ \bullet \quad \bullet \quad \bullet \\ \downarrow \quad \downarrow \quad \downarrow \end{array} + \begin{array}{c} |a\rangle \quad |a\rangle \quad |b\rangle \\ \downarrow \quad \downarrow \quad \downarrow \\ \bullet \quad \bullet \quad \bullet \\ \downarrow \quad \downarrow \quad \downarrow \end{array} - \begin{array}{c} |a\rangle \quad |a\rangle \quad |b\rangle \\ \downarrow \quad \downarrow \quad \downarrow \\ \bullet \quad \bullet \quad \bullet \\ \downarrow \quad \downarrow \quad \downarrow \end{array} - \begin{array}{c} |a\rangle \quad |a\rangle \quad |b\rangle \\ \downarrow \quad \downarrow \quad \downarrow \\ \bullet \quad \bullet \quad \bullet \\ \downarrow \quad \downarrow \quad \downarrow \end{array} - \begin{array}{c} |a\rangle \quad |a\rangle \quad |b\rangle \\ \downarrow \quad \downarrow \quad \downarrow \\ \bullet \quad \bullet \quad \bullet \\ \downarrow \quad \downarrow \quad \downarrow \end{array} \right) \quad (17.42)$$

$$\mathcal{A}_3|a, a, b\rangle = \frac{1}{6} \begin{pmatrix} |a, a, b\rangle + |a, b, a\rangle + |b, a, a\rangle \\ -|a, b, a\rangle - |a, a, b\rangle - |b, a, a\rangle \end{pmatrix} \quad (17.43)$$

$$= 0 \quad (17.44)$$

## 17.3 Levi-Civita Tensor

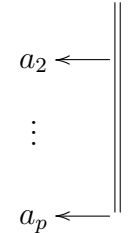
The **Levi-Civita tensor**  $\epsilon_{a:p}$  where  $a_i \in \{1, 2, \dots, p\}$  equals  $+1$  (resp.,  $-1$ ) if  $a:p$  is an even (resp., odd) permutation of  $(1, 2, \dots, p)$ . Thus

$$\epsilon^{123\dots p} = \epsilon_{123\dots p} = 1 \quad (17.45)$$

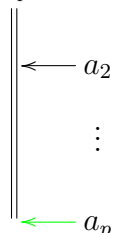
and

$$\epsilon_{rev(a:p)} = (-1)^{\binom{p}{2}} \epsilon_{a:p} \quad (17.46)$$

Define

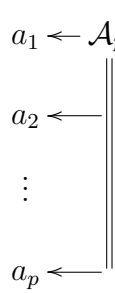
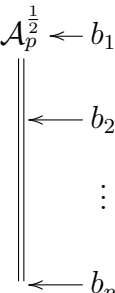
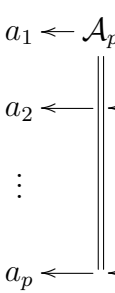
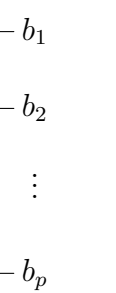
$$(C_{\mathcal{A}_p})_{a:p}^1 = e^{i\phi} \frac{\epsilon_{a:p}}{\sqrt{p!}} = a_1 \xleftarrow{\text{green}} \mathcal{A}_p^{\frac{1}{2}} \quad (17.47)$$


and

$$(C_{\mathcal{A}_p}^\dagger)_1^{rev(a:p)} = e^{-i\phi} \frac{\epsilon^{rev(a:p)}}{\sqrt{p!}} = \mathcal{A}_p^{\frac{1}{2}} \leftarrow a_1 \quad (17.48)$$


Then

$\frac{1}{p!} \epsilon_{a:p} \epsilon^{rev(b:p)} = (\mathcal{A}_p^{\frac{1}{2}})_{a:p}^{rev(b:p)}$



=



(17.49)

and

$$e^{i2\phi} \frac{1}{p!} \epsilon^{rev(a^n)} \epsilon_{a^n} = \delta_1^1 = 1 \quad \left( \begin{array}{c} \mathcal{A}_p^{\frac{1}{2}} \leftarrow \mathcal{A}_p^{\frac{1}{2}} \\ \leftarrow \\ \vdots \\ \leftarrow \end{array} \right) = 1 \quad (17.50)$$

For the L Convention, we will use  $\phi = 0$ .

For the CC Convention, we must choose

$$e^{i2\phi} = (-1)^{\binom{p}{2}} = e^{i\pi \frac{p(p-1)}{2}} \quad (17.51)$$

so

$$\phi = \frac{\pi}{4} p(p-1) \quad (17.52)$$

## 17.4 Fully-symmetric and Fully-antisymmetric tensors

fully symmetric (FS) tensor  $d$

$$d_{a_1 a_2 \dots a_p} = \begin{array}{c} d \\ | \\ a_1 \quad a_2 \quad \dots \quad a_p \end{array} \quad (17.53)$$

$$\mathcal{S}_p d = d \quad \left( \begin{array}{c} d \\ | \\ \mathcal{S}_p \\ | \end{array} \right) = \begin{array}{c} d \\ | \\ \dots \end{array} \quad (17.54)$$

$$0 = \begin{array}{c} d \\ | \\ T_i \end{array} + \begin{array}{c} d \\ | \\ T_i \end{array} + \begin{array}{c} d \\ | \\ T_i \end{array} \quad (17.55)$$

$$0 = \begin{array}{c} d \\ | \\ T_i \\ | \\ \mathcal{S}_p \\ | \end{array} \quad (17.56)$$

Fully antisymmetric(FA) tensor (FA)  $f$

$$f_{a_1 a_2 \dots a_p} = \begin{array}{c} d \\ \text{---} \\ | \\ a_1 \end{array} \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{array} \begin{array}{c} a_2 \end{array} \dots \begin{array}{c} a_p \end{array} \quad (17.57)$$

$$\mathcal{A}_p f = f \quad \begin{array}{c} f \\ \text{---} \\ | \\ \mathcal{A}_p \end{array} \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{array} \begin{array}{c} \dots \end{array} = \begin{array}{c} f \\ \text{---} \\ | \\ \dots \end{array} \quad (17.58)$$

$$0 = \begin{array}{c} f \\ \text{---} \\ | \\ T_i \end{array} \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{array} \begin{array}{c} \dots \end{array} + \begin{array}{c} f \\ \text{---} \\ | \\ T_i \end{array} \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{array} \begin{array}{c} \dots \end{array} + \begin{array}{c} f \\ \text{---} \\ | \\ T_i \end{array} \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{array} \begin{array}{c} \dots \end{array} \quad (17.59)$$

$$0 = \begin{array}{c} f \\ \text{---} \\ | \\ T_i \end{array} \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{array} \begin{array}{c} \dots \end{array} \quad (17.60)$$

## 17.5 Identically Vanishing Tensors

Identically vanishing (IV) tensors

- Example of birdtrack that vanishes for any FA tensor  $f$

$$\begin{array}{c} f \text{ ---} f \text{ ---} \\ \text{---} f \text{ ---} f \text{ ---} \\ f \text{ ---} f \text{ ---} \end{array} = 0 \quad (\text{Kuratowsky graph}) \quad (17.61)$$

- Example of birdtrack that vanishes for any  $f$  that is a structure constant of a Lie algebra



$$\begin{array}{ccc}
\begin{array}{c} \text{~~~~} \mathcal{S}_2 \text{~~~~} f \text{~~~~} \mathcal{S}_3 \text{~~~~} \\ \parallel \\ \text{~~~~} \text{~~~~} f \text{~~~~} \text{~~~~} \parallel \\ \text{~~~~} \text{~~~~} f \text{~~~~} \text{~~~~} \parallel \\ \text{~~~~} \text{~~~~} f \text{~~~~} \text{~~~~} \parallel \\ \text{~~~~} \text{~~~~} f \text{~~~~} \text{~~~~} \parallel \end{array} & = 0, & \begin{array}{c} \text{~~~~} \mathcal{A}_2 \text{~~~~} f \text{~~~~} \mathcal{A}_3 \text{~~~~} \\ \parallel \\ \text{~~~~} \text{~~~~} f \text{~~~~} \text{~~~~} \parallel \\ \text{~~~~} \text{~~~~} f \text{~~~~} \text{~~~~} \parallel \\ \text{~~~~} \text{~~~~} f \text{~~~~} \text{~~~~} \parallel \\ \text{~~~~} \text{~~~~} f \text{~~~~} \text{~~~~} \parallel \end{array} = 0
\end{array} \tag{17.62}$$

- Birdtrack that is zero for an irrep

$$\begin{array}{ccc}
\begin{array}{c} \text{~~~~} f \text{~~~~} \\ \diagup \quad \diagdown \\ \text{~~~~} f \text{~~~~} \\ \diagdown \quad \diagup \\ T_i \quad T_j \\ \diagup \quad \diagdown \\ T_k \end{array} & = 0, & \begin{array}{c} T_r \\ \diagup \quad \diagdown \\ \text{~~~~} f \text{~~~~} \\ \diagdown \quad \diagup \\ T_i \quad T_j \\ \diagup \quad \diagdown \\ T_s \end{array} = 0
\end{array} \tag{17.63}$$

# Chapter 18

## Unitary Groups: COMING SOON

### 18.1 $SU(n)$

In  $SU(n)$ ,

$$n = d_{\lambda_0} \quad (18.1)$$

where  $\lambda_0$  is the defining rep. ( $\mathcal{G} \subset \mathbb{C}^{n \times n}$ )

$$m(p, q) = \delta_b^a \sum_{a=1}^n (p_a)^* q_a \quad (18.2)$$

$$\mathbb{1}_{d,b}^{a,c} = \delta_b^a \delta_d^c = \begin{array}{c} d \leftarrow \bullet \rightarrow c \\ a \longrightarrow \bullet \rightarrow b \end{array} \quad (18.3)$$

$$M_{d,b}^{a,c} = \delta_d^a \delta_b^c = \begin{array}{c} d \quad c \\ \swarrow \quad \searrow \\ \bullet \quad \bullet \\ \swarrow \quad \searrow \\ a \quad b \end{array} \quad (18.4)$$

$$\boxed{M^2 = nM} \quad \begin{array}{c} d \quad c \\ \swarrow \quad \searrow \\ \bullet \quad \bullet \\ \swarrow \quad \searrow \\ a \quad b \end{array} \begin{array}{c} d \quad c \\ \swarrow \quad \searrow \\ \bullet \quad \bullet \\ \swarrow \quad \searrow \\ a \quad b \end{array} = n \begin{array}{c} d \quad c \\ \swarrow \quad \searrow \\ \bullet \quad \bullet \\ \swarrow \quad \searrow \\ a \quad b \end{array} \quad (18.5)$$

$$P_i = \sum_{j \neq i} \frac{M - \lambda_j}{\lambda_i - \lambda_j} \quad (18.6)$$

$$\lambda_1 = n$$

$$P_1 = \frac{M-n}{0-n} = 1 - \frac{1}{n}M$$

$$(18.7)$$

$$\lambda_2 = 0$$

$$P_2 = \frac{M-0}{n-0} = \frac{1}{n}M$$

$$(18.8)$$

$$\dim(P_1) = \text{tr} P_1 =$$

$$(18.9)$$

$$= n^2 - 1$$

$$(18.10)$$

$$\dim(P_2) = \text{tr} P_2 = \frac{1}{n}$$

$$(18.11)$$

$$= 1$$

$$(18.12)$$

**Claim 12**

$$\text{tr}(T_i) = 0$$

$$(18.13)$$

**proof:**

$$0 = P_1 P_2 =$$

$$(18.14)$$

**QED**

$$=$$

$$(18.15)$$

$$\begin{array}{c} \leftarrow \mathcal{S}_2 \leftarrow \\ \parallel \\ \leftarrow \leftarrow \end{array} = \frac{1}{2} \left\{ \begin{array}{c} \leftarrow \leftarrow \\ \leftarrow \leftarrow \end{array} + \begin{array}{c} \leftarrow \leftarrow \\ \updownarrow \\ \leftarrow \leftarrow \end{array} \right\} \quad (18.16)$$

$$\begin{array}{c} \leftarrow \mathcal{A}_2 \leftarrow \\ \parallel \\ \leftarrow \leftarrow \end{array} = \frac{1}{2} \left\{ \begin{array}{c} \leftarrow \leftarrow \\ \leftarrow \leftarrow \end{array} - \begin{array}{c} \leftarrow \leftarrow \\ \updownarrow \\ \leftarrow \leftarrow \end{array} \right\} \quad (18.17)$$

$$\dim(\mathcal{S}_2) = \frac{1}{2} \left\{ \begin{array}{c} \text{red arc} \leftarrow \leftarrow \\ \leftarrow \leftarrow \end{array} + \begin{array}{c} \text{red arc} \leftarrow \leftarrow \\ \updownarrow \\ \text{red arc} \leftarrow \leftarrow \end{array} \right\} \quad (18.18)$$

$$= \frac{n(n+1)}{2} \quad (18.19)$$

$$\dim(\mathcal{A}_2) = \frac{1}{2} \left\{ \begin{array}{c} \text{red arc} \leftarrow \leftarrow \\ \leftarrow \leftarrow \end{array} - \begin{array}{c} \text{red arc} \leftarrow \leftarrow \\ \updownarrow \\ \text{red arc} \leftarrow \leftarrow \end{array} \right\} \quad (18.20)$$

$$= \frac{n(n-1)}{2} \quad (18.21)$$

$$(T_i)^b_a = \begin{array}{c} i \\ \vdots \\ a \leftarrow T^i \leftarrow b \end{array} \quad (18.22)$$

$$T_i^\dagger = T_i \quad (18.23)$$

$$\text{tr}(T_i T_j) = \kappa \delta(i, j) \quad (18.24)$$

$$i \rightsquigarrow T_i \begin{array}{c} \curvearrowright \\ \curvearrowleft \end{array} T_j \rightsquigarrow j = \delta(i, j) \kappa i \rightsquigarrow i$$

Usually set  $\kappa = 1$

$$\begin{array}{c} \leftarrow \\ \updownarrow \\ \leftarrow \end{array} T_i \rightsquigarrow T_i \begin{array}{c} \leftarrow \\ \updownarrow \\ \leftarrow \end{array} \stackrel{\text{def}}{=} P_1 = \begin{array}{c} \leftarrow \bullet \\ \leftarrow \bullet \end{array} - \frac{1}{n} \begin{array}{c} \curvearrowright \bullet \\ \curvearrowleft \bullet \end{array} \quad (18.25)$$

### Claim 13

$$\Gamma_{fun} \delta_a^b = \sum_i (T_i T_i)_a^b = \frac{n^2 - 1}{n} \delta_a^b$$

$$\sum_i a \longleftarrow T_i \overset{i}{\text{---}} T_i \longleftarrow b = \left( \frac{n^2 - 1}{n} \right) a \longleftarrow \bullet b$$

proof:

$$(T_i T_i)_a^b = \sum_i a \longleftarrow T_i \overset{i}{\longleftarrow} T_i \longleftarrow b \quad (18.27)$$

$$= \begin{array}{c} a \quad b \\ \swarrow \quad \searrow \\ T_i \sim T_i \\ \nwarrow \quad \nearrow \end{array} \quad (18.28)$$

$$= -\frac{1}{n} \text{ (diagram of two vertices connected by a line with arrows pointing away from each other)} \quad (18.29)$$

$$= \binom{n-1}{n} a \leftarrow \bullet b \quad (18.30)$$

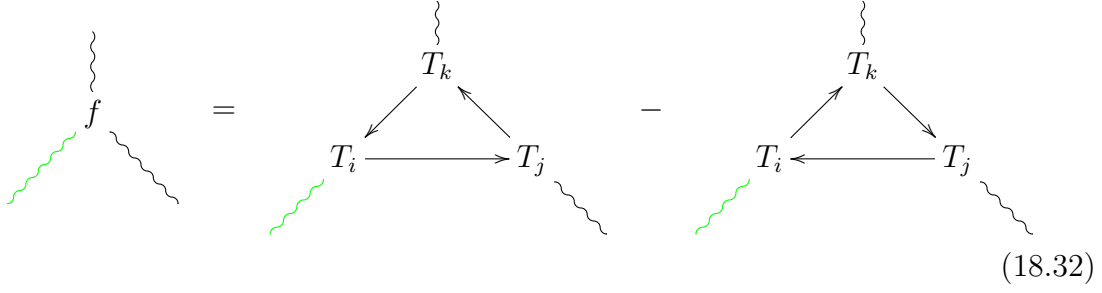
QED

$$T_i T_j - T_j T_i = i f_{ijk} T_k$$

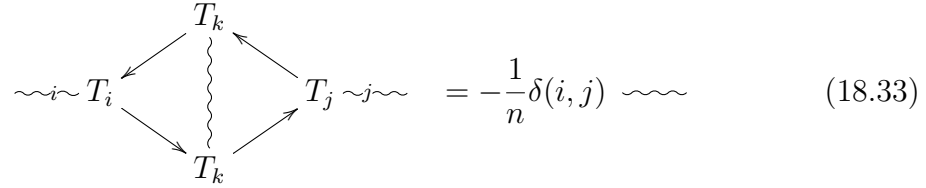
$$\begin{array}{c} \leftarrow T_i \leftarrow T_j \leftarrow \\ \vdots \quad \vdots \\ i \quad j \end{array} - \begin{array}{c} \leftarrow T_j \leftarrow T_i \leftarrow \\ \diagup \quad \diagdown \\ i \quad j \end{array} = \begin{array}{c} \leftarrow T_k \leftarrow \\ \textcolor{violet}{\vdots} \\ if \\ \diagup \quad \diagdown \\ i \quad j \end{array} \quad (18.31)$$

$f_{ijk}$  is totally antisymmetric. First index in green

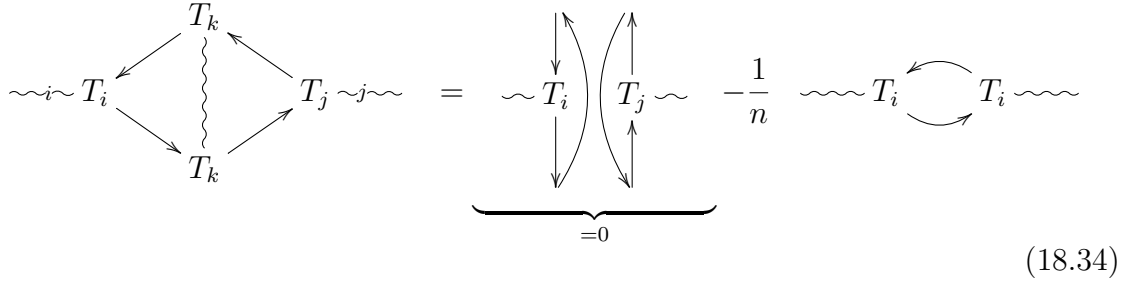
$$if_{ijk} = \text{tr}(T_i T_j T_k) - \text{tr}(T_j T_i T_k)$$



**Claim 14**



**proof:**



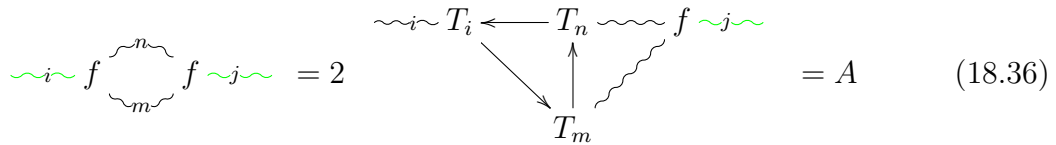
**QED**

**Claim 15**

$$\delta(i, j) \Gamma_{adj} = -f_{imn} f_{jnm} = 2n \delta(i, j)$$

$$(-1) \text{ } \sim i \sim f \text{ } \sim j \sim = 2n \delta(i, j)$$

**proof:**



$$\frac{1}{2}A = \underbrace{\begin{array}{c} \text{---}i\text{---}T_i \\ \swarrow \quad \searrow \\ T_k \quad T_m \\ \uparrow \quad \downarrow \\ T_n \end{array}}_{A_1} \text{---}j\text{---} - \underbrace{\begin{array}{c} \text{---}i\text{---}T_i \\ \swarrow \quad \searrow \\ T_k \quad T_m \\ \uparrow \quad \downarrow \\ T_n \end{array}}_{A_2} \text{---}j\text{---} \quad (18.37)$$

$$A_1 = \frac{n^2 - 1}{n} \delta(i, j) \quad (18.38)$$

$$A_2 = -\frac{1}{n} \delta(i, j) \quad (18.39)$$

$$A = 2(A_1 - A_2) = 2n\delta(i, j) \quad (18.40)$$

**QED**

# Chapter 19

## Wigner-Ekart Theorem

### 19.1 WE in general

$$M = \begin{array}{c} \leftarrow \mu - \\ \leftarrow \nu - \\ - \rho \rightarrow \\ - \omega \rightarrow \end{array} \begin{array}{c} M \\ || \\ || \\ || \end{array} \quad (19.1)$$

$$= \sum_{\alpha, \beta} \frac{1}{\kappa_{\alpha} \kappa_{\beta}} \begin{array}{c} \leftarrow \mu - \\ \leftarrow \nu - \\ - \rho \rightarrow \\ - \omega \rightarrow \end{array} \begin{array}{c} \begin{array}{c} \leftarrow \\ || \\ T_{\alpha}^{\dagger} \end{array} \leftarrow \begin{array}{c} \leftarrow M \\ || \\ T_{\alpha} \end{array} \\ \begin{array}{c} \leftarrow \\ || \\ T_{\beta} \end{array} \rightarrow \begin{array}{c} \leftarrow \\ || \\ T_{\beta}^{\dagger} \end{array} \end{array} \begin{array}{c} M \\ || \\ || \\ || \end{array} \quad (19.2)$$

$$= \sum_{\alpha} \frac{1}{\kappa_{\alpha}^2 d_{\alpha}} \begin{array}{c} \leftarrow \mu - \\ \leftarrow \nu - \\ - \rho \rightarrow \\ - \omega \rightarrow \end{array} \begin{array}{c} \begin{array}{c} \leftarrow \\ || \\ T_{\alpha}^{\dagger} \end{array} \leftarrow \begin{array}{c} \leftarrow M \\ || \\ T_{\alpha} \end{array} \\ \begin{array}{c} \leftarrow \\ || \\ T_{\alpha} \end{array} \rightarrow \begin{array}{c} \leftarrow \\ || \\ T_{\alpha}^{\dagger} \end{array} \end{array} \begin{array}{c} M \\ || \\ || \\ || \end{array} \quad (19.3)$$



$$\begin{array}{c} \leftarrow \\ \parallel \\ T_\alpha \\ \parallel \\ \leftarrow \\ \parallel \\ T_\alpha^\dagger \\ \parallel \\ \rightarrow \end{array} \begin{array}{c} \leftarrow M \\ \parallel \\ \leftarrow \\ \parallel \\ \rightarrow \\ \parallel \\ \rightarrow \end{array} = \frac{1}{d_\alpha} \begin{array}{c} \leftarrow \\ \parallel \\ T_\alpha \\ \parallel \\ \leftarrow \\ \parallel \\ T_\alpha^\dagger \\ \parallel \\ \rightarrow \end{array} \begin{array}{c} \leftarrow M \\ \parallel \\ \leftarrow \\ \parallel \\ \rightarrow \\ \parallel \\ \rightarrow \end{array} \quad (19.4)$$

$$M_a = \sum_\lambda \lambda a \leftarrow T_\lambda \quad (19.5)$$

$$M_{\lambda a}{}^{\mu b} = \leftarrow \lambda - M \leftarrow \mu - \quad (19.6)$$

$$(M^{\lambda i})_{\lambda_2 a}{}^{\lambda_1 b} = \frac{1}{d_\mu} \begin{array}{c} M \\ \curvearrowright \\ \mu \end{array} \delta(\mu, \lambda) \quad (19.7)$$

$$(M^{\lambda i})_{\lambda_2 a}{}^{\lambda_1 b} = \begin{array}{c} \lambda - \\ \downarrow \\ \leftarrow \lambda_2 - M^\lambda \leftarrow \lambda_1 - \end{array} \quad (19.8)$$

$$(M^\lambda)_{\lambda_2}{}^{\lambda_1} = \sum_{\lambda_2} \frac{d_{\lambda_2}}{\begin{array}{c} \leftarrow \lambda - \\ \downarrow \\ T_{\lambda_2}^\dagger \leftarrow \lambda_2 - T_{\lambda_2} \\ \downarrow \lambda_1 - \end{array}} \begin{array}{c} \lambda - \\ \downarrow \\ \leftarrow \lambda_2 - M^\lambda \leftarrow \lambda_1 - \end{array} \begin{array}{c} \parallel \\ C_{\lambda_2}^\dagger \leftarrow C_{\lambda_2} \\ \parallel \end{array} \quad (19.9)$$

$$= \frac{\begin{array}{c} \lambda - \\ \downarrow \\ M^\lambda \leftarrow \lambda_1 - \end{array} \begin{array}{c} \parallel \\ C_{\lambda_2}^\dagger \\ \parallel \end{array}}{\begin{array}{c} \leftarrow \lambda - \\ \downarrow \\ T_{\lambda_2}^\dagger \leftarrow \lambda_2 - T_{\lambda_2} \\ \downarrow \lambda_1 - \end{array}} \begin{array}{c} \lambda - \\ \downarrow \\ \leftarrow \lambda_2 - T_{\lambda_2} \leftarrow \lambda_1 - \end{array} \quad (19.10)$$

## 19.2 WE for angular momentum

$\lambda = J$ ,  $\lambda_i = J_i$  for  $i = 1, 2$ . We will use Greek letters instead of  $J$  so as to keep convention of using Greek letters for rep labels.

$m, m' = -\lambda, -\lambda + 1, \dots, \lambda - 1, \lambda$ . Note  $d_\lambda = 2\lambda + 1$   
for  $i = 1, 2$ ,  $m_i = -\lambda_i, -\lambda_i + 1, \dots, \lambda_i - 1, \lambda_i$ . Note  $d_{\lambda_i} = 2\lambda_i + 1$

$$\begin{array}{c} \leftarrow \lambda_1 m_1 \\ \parallel \\ \lambda m \leftarrow C_\lambda \\ \parallel \\ \leftarrow \lambda_2 m_2 \end{array} = \langle (\lambda_1 \lambda_2) \lambda m | \lambda_1 m_1 \lambda_2 m_2 \rangle \quad (19.11)$$

$$m \leftarrow D^\lambda \leftarrow m' = D_{mm'}^\lambda(g) \quad (19.12)$$

$$D_{m_1 m'_1}^{\lambda_1}(g) D_{m_2 m'_2}^{\lambda_2}(g) = \sum_{\lambda, m, m'} \langle \lambda_1 m_1 \lambda_2 m_2 | \lambda_1 \lambda_2 \lambda m \rangle D_{mm'}^\lambda(g) \langle \lambda_1 \lambda_2 \lambda m' | \lambda_1 m'_1 \lambda_2 m'_2 \rangle$$

$$\begin{array}{c} \leftarrow D^{\lambda_1} \leftarrow \\ \leftarrow D^{\lambda_2} \leftarrow \end{array} = \sum_\lambda \begin{array}{c} \leftarrow \parallel \\ C_\lambda^\dagger \leftarrow D^\lambda \leftarrow C_\lambda \\ \parallel \leftarrow \end{array} \quad (19.13)$$

tensor operator  $M_m^\lambda$

$$\langle \lambda_2 m_2 | M_m^\lambda | \lambda_1 m_1 \rangle = \begin{array}{c} \lambda m \\ \downarrow \\ \lambda_2 m_2 \leftarrow M_m^\lambda \leftarrow \lambda_1 m_1 \end{array} \quad (19.14)$$

$$\begin{array}{c} \langle \lambda_2 m_2 | M_m^\lambda | \lambda_1 m_1 \rangle = \langle (\lambda \lambda_1) \lambda_2 m_2 | \lambda m \lambda_1 m_1 \rangle q(\lambda, \lambda_1, \lambda_2) \\ \lambda m \\ \downarrow \\ \lambda_2 m_2 \leftarrow M_m^\lambda \leftarrow \lambda_1 m_1 \end{array} = q(\lambda, \lambda_1, \lambda_2) \begin{array}{c} \leftarrow \lambda \\ \parallel \\ \leftarrow \lambda_2 - C_\lambda \\ \parallel \\ \leftarrow \lambda_1 \end{array} \quad (19.15)$$

$$q(\lambda, \lambda_1, \lambda_2) = \frac{1}{d_{\lambda_2}} \sum_{m_1, m_2, m} \langle \lambda m \lambda_1 m_1 | (\lambda \lambda_1) \lambda_2 m_2 \rangle \langle \lambda_2 m_2 | M_m^\lambda | \lambda_1 m_1 \rangle \quad (19.16)$$

$$= \frac{1}{d_{\lambda_2}} \begin{array}{c} \lambda \\ \downarrow \\ M_m^\lambda \leftarrow \lambda_1 \\ \parallel \\ C_{\lambda_2}^\dagger \leftarrow \lambda \end{array} \quad (19.17)$$

# Chapter 20

## Young Tableau

A **Young Diagram** (YD)  $\mathcal{Y} = [\lambda_1, \lambda_2, \dots, \lambda_D]$  consists of  $\lambda_1$  left-aligned empty boxes over  $\lambda_2$ , over  $\lambda_3$  left-aligned empty boxes, up to  $\lambda_D$  boxes, where  $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_D \geq 1$ . For example,

$$[4, 2, 1, 1] = \begin{array}{|c|c|c|c|} \hline \square & \square & \square & \square \\ \hline \square & \square & & \\ \hline \square & & & \\ \hline \square & & & \\ \hline \end{array} \quad (20.1)$$

We will call  $[4, 2, 1, 1]$  the **row lengths** (RL) method of labeling YD.

A alternative method o labelling YD is called the **Dynkin labels** or **row changes** (RC). These labels list the change in number of columns as we go down the YD. For example,

$$\begin{array}{|c|c|c|c|} \hline \square & \square & \square & \square \\ \hline \square & \square & & \\ \hline \square & & & \\ \hline \square & & & \\ \hline \end{array} = [2, 1, 0, 1, 0 \dots]_{RC} \quad (20.2)$$

A **Young Tableau** (YT)  $\mathcal{Y}_\alpha$  is a YD in which integers from 1 to  $n$  where  $n \leq n_b$  and  $n_b$  is the number boxes, are inserted according to some rules. The rules for insertion are that integers must increase when reading a row left to right and when reading a column from top to bottom. Obviously, for  $n < n_b$ , some integers are repeated.

A **Standard Young Tableau** (SYT)  $\mathcal{Y}_\alpha$  is a YT such that  $n = n_b$  and no integer is repeated.

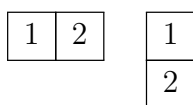
### 20.1 Symmetric group $S_{n_b}$

Let

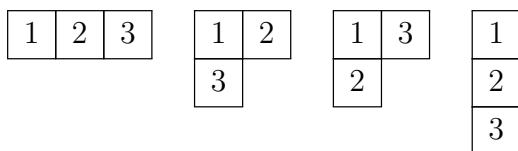
- $n_b = 1$



- $n_b = 2$



- $n_b = 3$



- $n_b = 4$

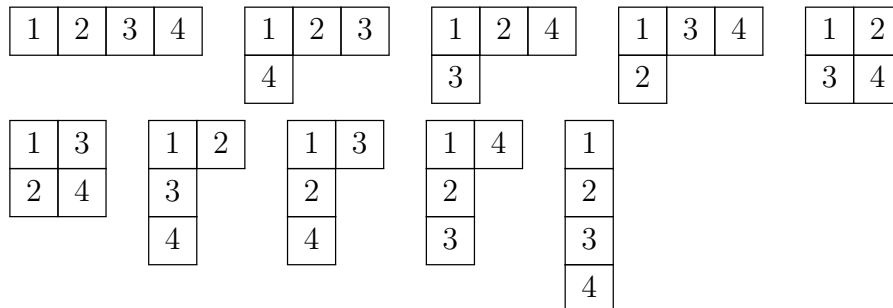


Figure 20.1: SYT for  $n_b = 1, 2, 3, 4$ .

$S_{n_b}$  = the symmetric group in  $n_b$  letters (or  $n_b$  boxes)

$irreps(S_{n_b})$  = the set of all irreps of  $S_{n_b}$ .

The **transpose of a YT** is defined as if it were a matrix. For example

$$transpose \left( \begin{array}{|c|c|c|} \hline 1 & 3 & 5 \\ \hline 2 & 6 & \\ \hline 4 & & \\ \hline \end{array} \right) = \begin{array}{|c|c|c|} \hline 1 & 2 & 4 \\ \hline 3 & 6 & \\ \hline 5 & & \\ \hline \end{array} \quad (20.3)$$

**n-dim General Linear group**  $GL(n) = \{M \in \mathbb{C}^{n \times n} : \det(M) \neq 0\}$

**n-dim Special Linear group**  $SL(n) = \{M \in GL(n) : \det(M) = 1\}$

**n-dim Unitary group**,  $U(n) = \{M \in GL(n) : MM^\dagger = M^\dagger M = 1\}$

**n-dim Special Unitary group**  $SU(n) = \{M \in U(n) : \det(M) = 1\}$

### Claim 16

1. The YD with  $n_b$  boxes label all irreps of the symmetric group  $S_{n_b}$ .
2. The SYT with  $n_b$  boxes and no more than  $n_b$  rows, label the irreps of  $GL(n_b)$  and of  $U(n_b)$
3. The SYT with  $n_b$  boxes and no more than  $n_b - 1$  rows, label the irreps of  $SL(n_b)$  and  $SU(n_b)$ .

**proof:**

**QED**

$YD_{n_b}$  = set of YD with  $n_b$  boxes

The irreps of  $S_{n_b}$  are labelled by the  $\mathcal{Y} \in YD_{n_b}$ . Hence, there is a 1-1 onto map between  $irreps(S_{n_b})$  and  $YD_{n_b}$

$$irreps(S_{n_b}) = YD_{n_b} \quad (20.4)$$

$YT(\mathcal{Y})$  = set of YT for a  $\mathcal{Y} \in YD_{n_b}$ .

$SYT(\mathcal{Y})$  = set of SYT for a  $\mathcal{Y} \in YD_{n_b}$ .

$dim(\mathcal{Y})$  = dimension of irrep  $\mathcal{Y}$ .

$$dim(\mathcal{Y}) = |SYT(\mathcal{Y})| \quad (20.5)$$

For example, the irrep

$$\mathcal{Y} = \begin{array}{|c|c|} \hline & \\ \hline & \\ \hline & \\ \hline \end{array} \quad (20.6)$$

has dimension 3 because there are 3 possible SYT for this YD:

$$\begin{array}{|c|c|} \hline 1 & 2 \\ \hline 3 & \\ \hline 4 & \\ \hline \end{array}, \quad
\begin{array}{|c|c|} \hline 1 & 3 \\ \hline 2 & \\ \hline 4 & \\ \hline \end{array}, \quad
\begin{array}{|c|c|} \hline 1 & 4 \\ \hline 2 & \\ \hline 3 & \\ \hline \end{array} \implies \dim(\mathcal{Y}) = 3 \quad (20.7)$$

Thus, we can denote the basis vectors of this irrep by

$$\left| \begin{array}{|c|c|} \hline 1 & 2 \\ \hline 3 & \\ \hline 4 & \\ \hline \end{array} \right\rangle, \quad \left| \begin{array}{|c|c|} \hline 1 & 3 \\ \hline 2 & \\ \hline 4 & \\ \hline \end{array} \right\rangle, \quad \left| \begin{array}{|c|c|} \hline 1 & 4 \\ \hline 2 & \\ \hline 3 & \\ \hline \end{array} \right\rangle \quad (20.8)$$

To compute  $hook(\mathcal{Y})$ :

1. Fill each box of the YD with the number of boxes below and to the right of the box, including the box itself.
2. Multiply the numbers in all the boxes.

For example,

$$\mathcal{Y} = \begin{array}{|c|c|c|c|} \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline & & & \\ \hline \end{array} \implies hook(\mathcal{Y}) = \begin{array}{|c|c|c|c|} \hline 6 & 5 & 3 & 1 \\ \hline 4 & 3 & 1 & \\ \hline 2 & 1 & & \\ \hline \end{array} = 6!3 \quad (20.9)$$

**Claim 17** (*hook rule for computing  $\dim(\mathcal{Y})$* )

$$\dim(\mathcal{Y}) = \frac{n_b!}{hook(\mathcal{Y})} \quad (20.10)$$

**proof:**

**QED**

For example

$$\mathcal{Y} = \begin{array}{|c|c|} \hline & \\ \hline & \\ \hline & \\ \hline \end{array} \implies hook(\mathcal{Y}) = \begin{array}{|c|c|} \hline 4 & 1 \\ \hline 2 & \\ \hline 1 & \\ \hline \end{array} = 8 \quad (20.11)$$

so

$$\dim(\mathcal{Y}) = \frac{4!}{4(2)} = 3 \quad (20.12)$$

The **regular representation** of the symmetric group  $S_{n_b}$  is defined as follows. For each permutation  $\sigma \in S_{n_b}$ , define an independent vector  $|\sigma\rangle$  in a vector space  $\mathcal{V} = \{|\sigma\rangle : \sigma \in S_{n_b}\} = \{|\sigma_i\rangle : i = 1, 2, \dots, n_b!\}$ . Let

$$|x\rangle = \sum_i x_i |\sigma_i\rangle \quad (20.13)$$

For any  $\tau \in S_{n_b}$ , suppose

$$\langle \sigma_j | \tau | \sigma_i \rangle = \langle \tau^{-1} \sigma_j | \sigma_i \rangle \quad (20.14)$$

$$\langle \sigma_j | \tau | x \rangle = \langle \tau^{-1} \sigma_j | x \rangle \quad (20.15)$$

The regular rep is  $n_b!$  dimensional and reducible.

**Claim 18** *The regular rep of  $S_{n_b}$  decomposes into each  $\lambda \in rreps(S_{n_b})$ , appearing  $\dim(\lambda)$  times. Thus*

$$n_b! = |S_{n_b}| = \sum_{\lambda \in irreps(S_{n_b})} [\dim(\lambda)]^2 \quad (20.16)$$

$$= [n_b!]^2 \sum_{\mathcal{Y} \in YD_{n_b}} \frac{1}{[\text{hook}(\mathcal{Y})]^2} \quad (\text{Because } irreps(S_{n_b}) = YD_{n_b}) \quad (20.17)$$

Hence,

$$1 = n_b! \sum_{\mathcal{Y} \in YD_{n_b}} \frac{1}{[\text{hook}(\mathcal{Y})]^2} \quad (20.18)$$

**proof:**  
**QED**

$$1 = \sum_{\mathcal{Y} \in YD_{n_b}} P_{\mathcal{Y}} \quad (20.19)$$

$$P_{\mathcal{Y}} = \sum_{\mathcal{Y}_{\alpha} \in SYT(\mathcal{Y})} |\mathcal{Y}_{\alpha}\rangle \langle \mathcal{Y}_{\alpha}| \quad (20.20)$$

The projection operators are complete and orthogonal.

$$\begin{array}{|c|} \hline \square \\ \hline \end{array} \otimes \begin{array}{|c|} \hline \square \\ \hline \end{array} = \begin{array}{|c|c|} \hline \square & \square \\ \hline \end{array} \oplus \begin{array}{|c|} \hline \square \\ \hline \square \\ \hline \end{array} \quad (20.21)$$

$$\begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \end{array} \otimes \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \\ \hline \end{array} =$$

$$\begin{array}{|c|c|c|c|c|} \hline \square & \square & \square & \square & \square \\ \hline \square & & & & \end{array} \oplus \begin{array}{|c|c|c|c|} \hline \square & \square & \square & \square \\ \hline \square & \square & & \end{array} \oplus \begin{array}{|c|c|c|c|} \hline \square & \square & \square & \square \\ \hline \square & & & \end{array} \oplus \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \\ \hline \square & & \end{array} \quad (20.22)$$

## 20.2 Unitary group $U(n)$

$SYT(n_b, R \leq n)$  = SYTs with  $n_b$  boxes and less than  $n$  rows

$$irreps(U(n)) = \cup_{n_b \leq n} SYT(n_b, R \leq n) \quad (20.23)$$

A SYT with  $n_b$  boxes represents a tensor with  $n_b$  indices ( $n_b$ -particles state). Each index ranges from 1 to  $n$ .

$n_b = 1$ : A 1-index, 1-box tensor is a 1-particle with  $n$  states. This corresponds to the fundamental representation.

$n_b = 2$ : A 2-index, 2-box tensor is a 2-particle with  $n^2$  states. These  $n^2$  states break into two sets, symmetric and anti-symmetric.

$$\boxed{1} \otimes \boxed{1} = \begin{array}{|c|} \hline 1 \\ \hline 2 \\ \hline \end{array} \oplus \boxed{1 \ 2} \quad (20.24)$$

$$\begin{array}{c} \longleftarrow \\ \longleftarrow \end{array} = \begin{array}{c} \longleftarrow \mathcal{A}_2 \longleftarrow \\ \parallel \\ \longleftarrow \parallel \longleftarrow \end{array} + \begin{array}{c} \longleftarrow \mathcal{S}_2 \longleftarrow \\ \parallel \\ \longleftarrow \parallel \longleftarrow \end{array} \quad (20.25)$$

The SYT of an irrep describes the symmetry of the indices of a tensor in that irrep. A single column SYT indicates a totally anti-symmetric tensor, a single row SYT indicates a totally symmetric tensor, and a neither single column nor single row SYT indicates mixed symmetry. This is why we can't have more than  $n$  rows, because there are only  $n_b$  integers to fill all boxes so more than  $n$  rows would require a repetition in of an integer in a column, and such a column, after antisymetrizing, equals zero.

$YD_{n_b}$  = set of YD with  $n_b$  boxes

$YT(\mathcal{Y})$  = set of YT for a  $\mathcal{Y} \in YD_{n_b}$ .

If  $\mathcal{Y}_\alpha$  is a SYT in  $irreps(U(n))$  and the YD of  $\mathcal{Y}_\alpha$  is  $\mathcal{Y}$ , then

$$\dim(\mathcal{Y}_\alpha) = |YT(\mathcal{Y})| \quad (20.26)$$

Hence,  $\{|\alpha\rangle : \alpha \in YT(\mathcal{Y})\}$  are a basis for the irrep  $\mathcal{Y}_\alpha$ . Note that the irreps of  $U(n)$  are given by SYT  $\mathcal{Y}_\alpha$ , whereas the basis vectors of an irrep are given by YT (not just STY). For example, for

$$\mathcal{Y}_\alpha = \boxed{1 \ 2} \quad (20.27)$$

the basis vectors are

$$|\boxed{1 \ 1}\rangle, \quad |\boxed{1 \ 2}\rangle, \quad |\boxed{2 \ 2}\rangle \quad (20.28)$$

so

$$\dim(\mathcal{Y}_\alpha) = 3 \quad (20.29)$$



## 20.3 Young Projection operators

For each in the SYT  $\mathcal{Y}_\alpha \in \text{irreps}(U(n))$ , we have

$$P_{\mathcal{Y}_\alpha} = \mathcal{N}_{\mathcal{Y}} \left( \prod_i S_i \right) \left( \prod_j A_j \right) \quad (20.30)$$

Note that the normalization constant  $\mathcal{N}_{\mathcal{Y}}$  depends only on the YD  $\mathcal{Y}$ .

These projection operators are not unique.

**Claim 19** *Let*

$$\mathcal{N}_{\mathcal{Y}} = \frac{(\prod_i |S_i|!) (\prod_j |A_j|!)}{\text{hook}(\mathcal{Y})} \quad (20.31)$$

where  $|S_i|$  and  $|A_j|$  are the number of arrows entering the symmetrizer or anti-symmetrizer. Then the operators  $P_{\mathcal{Y}_\alpha}$  are idempotent (i.e., their square equals themselves) mutually orthogonal and complete:

$$P_{\mathcal{Y}_\alpha} P_{\mathcal{Y}_\beta} = P_{\mathcal{Y}_\alpha} \delta(\alpha, \beta) \quad (20.32)$$

$$1 = \sum_{\mathcal{Y}_\alpha \in \text{SYT}(n_b, R < n)} P_{\mathcal{Y}_\alpha} \quad (20.33)$$

**proof:**

$$P_{\mathcal{Y}_\alpha} = \mathcal{N}_{\mathcal{Y}} \frac{1}{\prod_i |S_i|! \prod_j |A_j|!} \left( \begin{array}{c} \leftarrow \\ \leftarrow \\ \vdots \\ \underbrace{\leftarrow}_1 \end{array} + \cdots \right) \quad (20.34)$$

From Eq.(20.33)

$$\mathbb{1} = \sum_{\mathcal{Y}_\alpha \in \text{SYT}(n_b, R < n)} \mathcal{N}_{\mathcal{Y}} \frac{1}{\prod_i |S_i|! \prod_j |A_j|!} \mathbb{1} \quad (20.35)$$

$$= \sum_{\mathcal{Y}} \frac{n_b!}{\text{hook}(\mathcal{Y})} \mathcal{N}_{\mathcal{Y}} \frac{1}{\prod_i |S_i|! \prod_j |A_j|!} \mathbb{1} \quad (20.36)$$

$$= \sum_{\mathcal{Y}} \frac{n_b!}{[\text{hook}(\mathcal{Y})]^2} \frac{1}{\prod_i |S_i|! \prod_j |A_j|!} \mathbb{1} \quad (\text{if assume Eq.(20.31)}) \quad (20.37)$$

$$= \mathbb{1} \quad (\text{by Eq.(20.18)}) \quad (20.38)$$

**QED**

Let  $\dim(\mathcal{Y}_\alpha)$  be the dimension of an irrep of  $U(n)$  with STY given by  $\mathcal{Y}_\alpha \in STY(n_b, R < n)$ . In Eq.(20.26) we gave a way of finding  $\dim(\mathcal{Y}_\alpha)$ . A second way is by taking the trace of the corresponding projection operator

$$\dim(\mathcal{Y}_\alpha) = \text{tr}(P_{\mathcal{Y}_\alpha}) \quad (20.39)$$

For example, if

$$\mathcal{Y}_\alpha = \begin{array}{|c|c|} \hline 1 & 2 \\ \hline \end{array} \quad (20.40)$$

then

$$d_{\mathcal{Y}_\alpha} = \begin{array}{c} \leftarrow \mathcal{S}_2 \leftarrow \\ \parallel \\ \leftarrow \leftarrow \end{array} \quad (20.41)$$

$$= \frac{1}{2} \left( \begin{array}{c} \leftarrow \\ \leftarrow \end{array} + \begin{array}{c} \leftarrow \leftarrow \\ \uparrow \downarrow \\ \leftarrow \leftarrow \end{array} \right) \quad (20.42)$$

$$= \frac{1}{2}(n^2 + n) \quad (20.43)$$

$$= 3 \text{ for } U(n=2) \quad (20.44)$$

A third way of computing  $\dim(\mathcal{Y}_\alpha)$  is by computing the hook and coat functions and using the formula

$$\dim(\mathcal{Y}_\alpha) = \frac{\text{coat}(\mathcal{Y})}{\text{hook}(\mathcal{Y})} \quad (20.45)$$

Note that right hand side is independent of  $\alpha$ ; it depends only on the YD. We've already discussed how to compute  $\text{hook}(\mathcal{Y})$ .  $\text{coat}(\mathcal{Y})$  is calculated as follows.<sup>1</sup>

1. Fill  $\mathcal{Y}$  with
  - $n$  at the diagonal blocks
  - $n$  increments increasing by 1 when reading from left to right
  - $n$  increments decreasing by 1 when reading from top to bottom
2. multiply all the boxes

Examples

$$\dim\left(\begin{array}{|c|c|} \hline 1 & 2 \\ \hline \end{array}\right) = \frac{\begin{array}{|c|c|} \hline n & n+1 \\ \hline \end{array}}{\begin{array}{|c|c|} \hline 2 & 1 \\ \hline \end{array}} = \frac{n(n+1)}{2} \quad (20.46)$$

---

<sup>1</sup>I invented that name. I don't know if it has a name.

$$\dim\left(\begin{array}{|c|} \hline 1 \\ \hline 2 \\ \hline \end{array}\right) = \frac{\begin{array}{|c|} \hline n \\ \hline n-1 \\ \hline 2 \\ \hline 1 \\ \hline \end{array}}{2} = \frac{n(n-1)}{2} \quad (20.47)$$

$$\dim\left(\begin{array}{|c|c|c|c|} \hline 1 & 2 & 3 & 4 \\ \hline 5 & 6 & & \\ \hline 7 & & & \\ \hline \end{array}\right) = \frac{\begin{array}{|c|c|c|c|} \hline n & n+1 & n+2 & n+3 \\ \hline n-1 & n & & \\ \hline n-2 & & & \\ \hline 6 & 4 & 2 & 1 \\ \hline 3 & 1 & & \\ \hline 1 & & & \\ \hline \end{array}}{144} = \frac{n^2(n^2-1)(n^2-4)(n+3)}{144} \quad (20.48)$$

## 20.4 Young Projection operators for $n_b = 1, 2, 3, 4$

$$\begin{array}{c} \leftarrow \mathcal{S}_2 \leftarrow \\ \parallel \\ \leftarrow \square \leftarrow \\ \parallel \\ \leftarrow \leftarrow \end{array} = \begin{array}{c} \leftarrow \leftarrow \mathcal{S}_2 \leftarrow \leftarrow \\ \parallel \\ \leftarrow \leftarrow \leftarrow \leftarrow \\ \updownarrow \quad \updownarrow \\ \leftarrow \leftarrow \leftarrow \leftarrow \end{array} \quad (20.49)$$

- $n_b = 1$

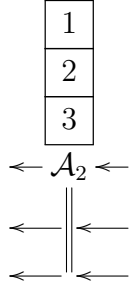
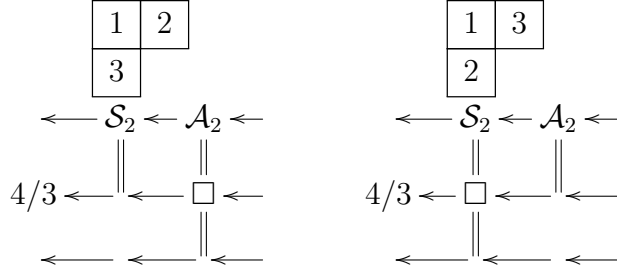
$$\begin{array}{|c|} \hline 1 \\ \hline \end{array} \leftarrow$$

- $n_b = 2$

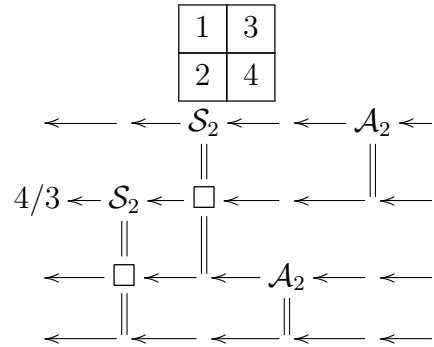
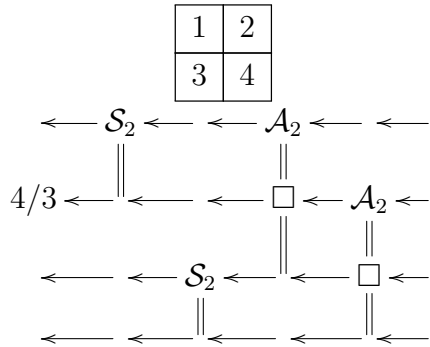
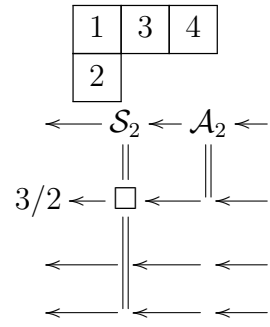
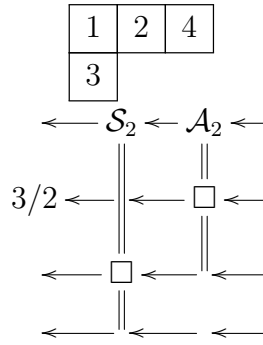
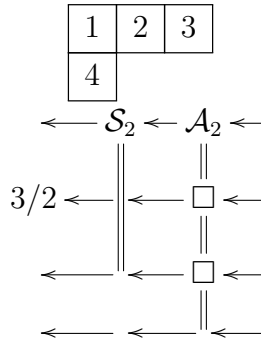
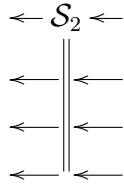
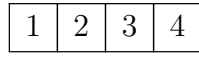
$$\begin{array}{|c|c|} \hline 1 & 2 \\ \hline \end{array} \begin{array}{c} \leftarrow \mathcal{S}_2 \leftarrow \\ \parallel \\ \leftarrow \leftarrow \end{array} \quad \begin{array}{|c|} \hline 1 \\ \hline 2 \\ \hline \end{array} \begin{array}{c} \leftarrow \mathcal{A}_2 \leftarrow \\ \parallel \\ \leftarrow \leftarrow \end{array}$$

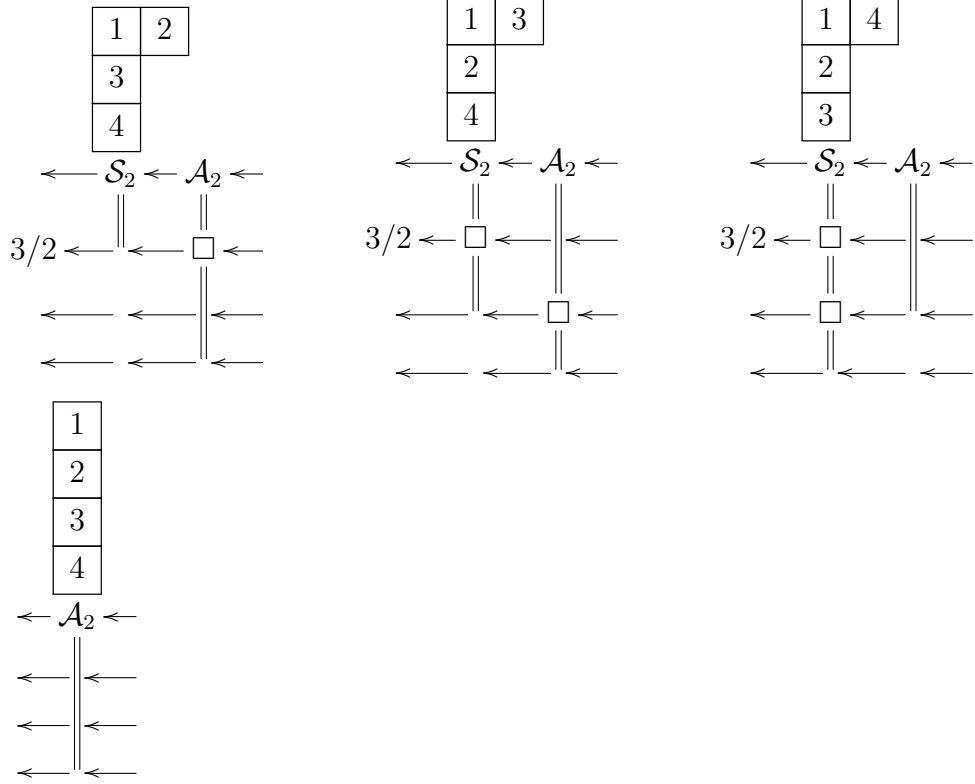
- $n_b = 3$

$$\begin{array}{|c|c|c|} \hline 1 & 2 & 3 \\ \hline \end{array} \begin{array}{c} \leftarrow \mathcal{S}_2 \leftarrow \\ \parallel \\ \leftarrow \leftarrow \leftarrow \end{array}$$



•  $n_b = 4$





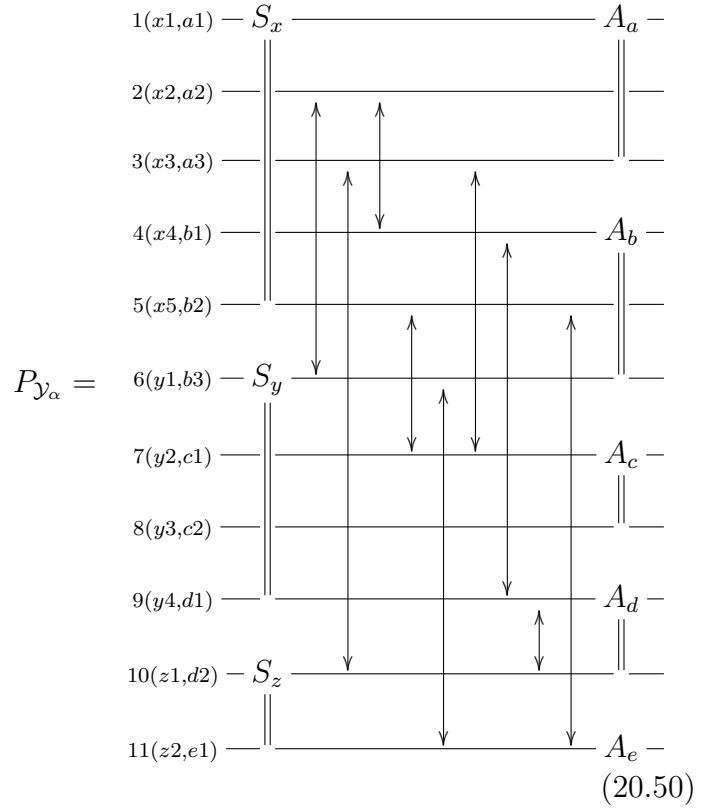
SYT and corresponding unnormalized projection operators for  $n_b = 1, 2, 3, 4$

## 20.5 Young Projection Operator with swaps

$\mathcal{Y}_\alpha =$

	$A_a$	$A_b$	$A_c$	$A_d$	$A_e$
$S_x$	1	2	3	4	5
$S_y$	6	7	8	9	
$S_z$	10	11			

$a1 \longleftrightarrow 1$
$a2 \longleftrightarrow 6$
$a3 \longleftrightarrow 10$
$b1 \longleftrightarrow 2$
$b2 \longleftrightarrow 7$
$b3 \longleftrightarrow 11$
$c1 \longleftrightarrow 3$
$c2 \longleftrightarrow 8$
$d1 \longleftrightarrow 4$
$d2 \longleftrightarrow 9$
$e1 \longleftrightarrow 5$



## 20.6 Tensor product decompositions

$$\boxed{1} \otimes \boxed{2} \otimes \boxed{3} = \left( \boxed{1 \ 2} \oplus \boxed{\begin{smallmatrix} 1 \\ 2 \end{smallmatrix}} \right) \otimes \boxed{3} \quad (20.51)$$

$$= \boxed{1 \ 2 \ 3} \oplus \boxed{\begin{smallmatrix} 1 \ 2 \\ 3 \end{smallmatrix}} \oplus \boxed{\begin{smallmatrix} 1 \ 3 \\ 2 \end{smallmatrix}} \oplus \boxed{\begin{smallmatrix} 1 \\ 2 \\ 3 \end{smallmatrix}} \quad (20.52)$$

$$\begin{array}{c} \leftarrow \\ \leftarrow \\ \leftarrow \end{array} = \begin{array}{c} \leftarrow \mathcal{S}_2 \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \parallel \leftarrow \end{array} + \frac{4}{3} \begin{array}{c} \leftarrow \mathcal{S}_2 \leftarrow \mathcal{A}_2 \leftarrow \\ \leftarrow \parallel \leftarrow \square \leftarrow \\ \leftarrow \parallel \leftarrow \end{array} + \frac{4}{3} \begin{array}{c} \leftarrow \mathcal{S}_2 \leftarrow \mathcal{A}_2 \leftarrow \\ \leftarrow \square \leftarrow \parallel \leftarrow \\ \leftarrow \parallel \leftarrow \end{array} + \begin{array}{c} \leftarrow \mathcal{A}_2 \leftarrow \\ \leftarrow \parallel \leftarrow \\ \leftarrow \parallel \leftarrow \end{array} \quad (20.53)$$

$$n^3 = \frac{n(n+1)(n+2)}{6} + \frac{n(n^2-1)}{3} + \frac{n(n^2-1)}{3} = \frac{(n-2)(n-1)n}{6} \quad (20.54)$$

$$\begin{array}{|c|c|c|} \hline 1 & 2 & 3 \\ \hline \end{array} \otimes \begin{array}{|c|c|} \hline 4 & 5 \\ \hline 6 \\ \hline \end{array} = \left\{ \begin{array}{c} \begin{array}{|c|c|c|c|c|} \hline 1 & 2 & 3 & 4 & 5 \\ \hline 6 \\ \hline \end{array} \oplus \begin{array}{|c|c|c|c|} \hline 1 & 2 & 3 & 4 \\ \hline 5 & 6 \\ \hline \end{array} \oplus \begin{array}{|c|c|c|c|} \hline 1 & 2 & 3 & 5 \\ \hline 4 & 6 \\ \hline \end{array} \\ \oplus \begin{array}{|c|c|c|c|} \hline 1 & 2 & 3 & 5 \\ \hline 4 \\ \hline 6 \\ \hline \end{array} \oplus \begin{array}{|c|c|c|c|} \hline 1 & 2 & 3 & 4 \\ \hline 5 \\ \hline 6 \\ \hline \end{array} \oplus \begin{array}{|c|c|c|} \hline 1 & 2 & 3 \\ \hline 4 & 5 \\ \hline 6 \\ \hline \end{array} \end{array} \right\} \quad (20.55)$$

For  $U(n)$ , the blue YT are zero for  $n_b = 2$ , and non-zero otherwise.

## 20.7 $SU(n)$

A third way, besides RL and RC, of labelling YD (reps of  $SU(n)$ ) is by their dimension, and then adding a subscript or overline if there are more than one reps with a different YD but the same dimension. This method is used mostly by physicists for  $SU(3)$  (The Eightfold Way). Note that all YT with the same YD have the same dimension, so this really labels YD instead of YT. For example, for  $SU(3)$  we have

$$\begin{array}{ll}
\begin{array}{|c|} \hline \square \\ \hline \end{array} = [1, 0]_{RC} = 3 & \begin{array}{|c|c|} \hline \square & \square \\ \hline \end{array} = [0, 1]_{RC} = \bar{3} \\
\begin{array}{|c|c|} \hline \square & \square \\ \hline \end{array} = [2, 0]_{RC} = 6 & \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \end{array} = [0, 2]_{RC} = \bar{6} \\
\begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \\ \hline \end{array} = [1, 1]_{RC} = 8 & \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & & \\ \hline \end{array} = [2, 1]_{RC} = 15
\end{array} \tag{20.56}$$



# Bibliography

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