

# Chapter 1

## Do Calculus proofs

In Chapter ?? of Bayesuvius, we explained Do Calculus but referred to this chapter for proofs of claims that use Do Calculus. In this chapter, we've aggregated all proofs, from throughout the book, of claims that use Do Calculus.

Note that even though the 3 rules of Do Calculus are great for proving adjustment formulae for general classes of DAGs, they are sometimes overkill for proving adjustment formulae for a single specific DAG. After all, the 3 rules of Do Calculus are a consequence of the d-separation theorem. Hence, all adjustment formulae should be provable from first principles, assuming only the d-separation theorem and the standard rules of probability theory.

We use the following conventions. Random variables are underlined and their values are not. For example,  $\underline{a} = a$  means the random variable  $\underline{a}$  takes the value  $a$ . Diagrams with nodes that are underlined represent Bayesian Networks (bnets) and the same diagram with the letters not underlined represents a specific instantiation of that bnet. For example  $\underline{a} \rightarrow \underline{b}$  represents the bnet with conditional probability distribution  $P(b|a)$ , whereas  $a \rightarrow b$  represents  $P(b|a)$  itself.

If  $\underline{a}$  is a root node, then  $\sum \underline{a}$  signifies a weighted sum  $\sum_a P(a)$ . For example,  $\sum \underline{a} \rightarrow \underline{b} = \sum_a P(a)P(b|a)$ . If  $\underline{a}$  is not a root node as in  $x \rightarrow \sum \underline{a} \rightarrow y = \sum_a P(y|a)P(a|x)$ , then  $\sum \underline{a}$  signifies a simple unweighted sum  $\sum_a$ .

Unobserved nodes are indicated by enclosing them in a dashed circle. For example,  $(\underline{u})$ .

Selection diagrams with selection nodes are discussed in Chapter ?. In a selection diagram with a selection node  $\underline{s} \in \{0, 1\}$ , if a node  $\underline{x}$  has parents  $pa(\underline{x})$  where  $\underline{s} \notin pa(\underline{x})$ , then the TPM of  $\underline{x}$  is  $P(x|pa(x))$ . If, on the other hand,  $\underline{x}$  has parents  $pa(\underline{x}) \cup \underline{s}$ , then the TPM of  $\underline{x}$  is  $P(x|pa(x), \underline{s})$ , where  $P(x|pa(x), \underline{s} = 0) = P(x|pa(x))$  and  $P(x|pa(x), \underline{s} = 1) = P^*(x|pa(x))$ .

Some identities that are used in this chapter:

1.

$$P(y|x_1, x_2) = \sum_a P(y|a, x_1, x_2)P(a|x_1, x_2) . \quad (1.1)$$

$$\begin{array}{ccc}
x_1 & & x_1 \\
\searrow & & \searrow \quad \searrow \\
& y & \sum a \quad \rightarrow y \\
\nearrow & & \nearrow \quad \nearrow \\
x_2 & & x_2
\end{array} = \quad . \quad (1.2)$$

One can describe this identity as “giving  $\underline{y}$  a universal backdoor”, because  $\sum a$  is a backdoor (i.e., input) to  $y$ , and  $\sum a$  is universal in the sense that it is entered by every arrow that enters  $y$  except  $\sum a$  itself.

2.

$$\sum_a P(a|x_1, x_2) = 1 \quad (1.3)$$

$$\begin{array}{ccc}
x_1 & \searrow & \\
& \sum a & \xrightarrow{0} \\
x_2 & \nearrow &
\end{array} = 1 \quad (1.4)$$

One can describe this identity as “summing over the values of a collider node which has no emerging arrows”<sup>1</sup>. Eq.(1.4) can be understood as an edge case (when  $\underline{y} = \emptyset$ ) of Eq.(1.2).

3.

$$\sum_a P(x_2|a)P(a|x_1) = P(x_2|x_1) \quad (1.5)$$

$$x_1 \xrightarrow{\quad} \sum a \xrightarrow{\quad} x_2 = x_1 \xrightarrow{\quad} x_2 \quad (1.6)$$

One can describe this identity as “summing over the values of a mediator node”.

4.

$$P(x) = \sum_a P(x|a)P(a) = \sum_b P(x|b) \quad (1.7)$$

$$\xrightarrow{0} \sum a \xrightarrow{\quad} x = \xrightarrow{0} \sum b \xrightarrow{\quad} x \quad (1.8)$$

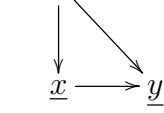
One can describe this identity as “averaging over different priors”. Eq.(1.8) can be understood as an edge case of Eq.(1.6).

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<sup>1</sup>A zeroed arrow means the same as no arrow.

**Claim 1** (*Backdoor Adjustment Formula*)

If  $\underline{z}$  then



$$P(y|\mathcal{D}\underline{x} = x) = \sum_z P(y|x, z)P(z) \quad (1.9)$$

$$\begin{array}{c} \mathcal{D}\underline{x} = x \longrightarrow y \quad \sum z \\ = \\ x \longrightarrow y \end{array} \quad (1.10)$$

**proof:**

\* **proof 1:**

$$\begin{array}{c} P(y|\mathcal{D}\underline{x} = x) = \sum_z P(y|x, z)P(z) \\ \sum z \quad \quad \quad \sum z \\ \mathcal{D}\underline{x} = x \longrightarrow y = x \longrightarrow y \end{array}$$

\* **proof 2:**

$$\begin{array}{l} P(y|\mathcal{D}\underline{x} = x) = \sum_z P(y|\mathcal{D}\underline{x} = x, z)P(z|\mathcal{D}\underline{x} = x) \\ \text{by Probability Axioms} \\ = \sum_z P(y|x, z)P(z|\mathcal{D}\underline{x} = x) \\ P(y|\mathcal{D}\underline{x} = x, z) \rightarrow P(y|x, z) \\ \text{by Rule 2: } \text{If } (\underline{b}. \perp \underline{a}.|\underline{r}., \underline{s}.) \text{ in } \mathcal{L}_{\underline{a}.}\mathcal{D}_{\underline{r}.}G, \text{ then} \\ \mathcal{D}_{\underline{a}.} = \underline{a}. \leftrightarrow \underline{a}. = \underline{a}. \\ \underline{y} \perp \underline{x}|\underline{z} \text{ in } \mathcal{L}_{\underline{x}}\mathcal{D}_{\emptyset}G : \begin{array}{c} \underline{z} \\ \downarrow \quad \searrow \\ \underline{x} \quad \quad \underline{y} \end{array} \\ = \sum_z P(y|x, z)P(z) \\ P(z|\mathcal{D}\underline{x} = x) \rightarrow P(z) \\ \text{by Rule 3: } \text{If } (\underline{b}. \perp \underline{a}.|\underline{r}., \underline{s}.) \text{ in } \mathcal{D}_{\underline{a}. - an(\underline{s}.)}\mathcal{D}_{\underline{r}.}G, \text{ then} \\ \mathcal{D}_{\underline{a}.} = \underline{a}. \leftrightarrow 1 \\ \underline{z} \perp \underline{x} \text{ in } \mathcal{D}_{\underline{x}}\mathcal{D}_{\emptyset}G : \begin{array}{c} \underline{z} \\ \searrow \\ \underline{x} \longrightarrow \underline{y} \end{array} \end{array}$$

**QED**

**Claim 2** (*Frontdoor Adjustment Formula*)

If  $\begin{array}{c} (\underline{c}) \\ \swarrow \quad \searrow \\ \underline{x} \longrightarrow \underline{m} \longrightarrow \underline{y} \end{array}$  then

$$P(y|\mathcal{D}\underline{x} = x) = \sum_m \left[ \sum_{x'} P(y|x', m) P(x') \right] P(m|x) \quad (1.11)$$

$$\begin{array}{ccc} \mathcal{D}\underline{x} = x & \longrightarrow & y \\ & & \nwarrow \quad \nearrow \\ & & \sum x' \\ & & \searrow \\ & & y \end{array} \quad (1.12)$$

$$= \begin{array}{ccc} x & \longrightarrow & \sum m \longrightarrow y \end{array}$$

**proof:**

\* **proof 1:**

$$P(y|\mathcal{D}\underline{x} = x) = \sum_{m,c,x'} P(y|m, c) P(c|x') P(x') P(m|\mathcal{D}\underline{x} = x)$$

$$\begin{array}{ccc} & (\sum c) & \longrightarrow \sum x' \longrightarrow (\sum c) \\ & \searrow & \searrow \\ \mathcal{D}\underline{x} = x & \longrightarrow \sum m \longrightarrow y & = \mathcal{D}\underline{x} = x \longrightarrow \sum m \longrightarrow y \\ = \sum_{m,x'} P(y|m, x') P(x') P(m|x) & & \\ & \sum x' & \\ = & \searrow & \\ & x \longrightarrow \sum m \longrightarrow y & \end{array}$$

\* **proof 2:**

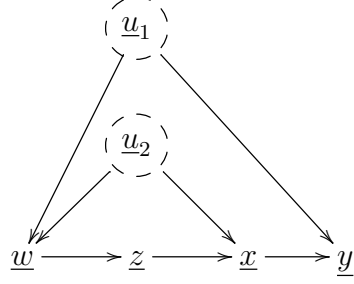
$$\begin{array}{l} P(y|\mathcal{D}\underline{x} = x) = \sum_m P(y|\mathcal{D}\underline{x} = x, m) P(m|\mathcal{D}\underline{x} = x) \\ \text{by Probability Axioms} \\ = \sum_m P(y|\mathcal{D}\underline{x} = x, \mathcal{D}\underline{m} = m) P(m|\mathcal{D}\underline{x} = x) \\ P(y|\mathcal{D}\underline{x} = x, m) \rightarrow P(y|\mathcal{D}\underline{x} = x, \mathcal{D}m = m) \\ \text{by Rule 2: If } (\underline{b} \perp \underline{a} | \underline{r}, \underline{s}) \text{ in } \mathcal{L}_{\underline{a}} \mathcal{D}_{\underline{r}} G, \text{ then} \\ \mathcal{D}\underline{a} = \underline{a} \leftrightarrow \underline{a} = \underline{a}. \end{array}$$

$$\begin{aligned}
& \underline{y} \perp \underline{m} | \underline{x} \text{ in } \mathcal{L}_{\underline{m}} \mathcal{D}_{\underline{x}} G : \quad \begin{array}{c} \textcircled{\underline{c}} \\ \searrow \\ \underline{x} \longrightarrow \underline{m} \quad \underline{y} \end{array} \\
&= \sum_m P(y | \mathcal{D}_{\underline{x}} = x, \mathcal{D}_{\underline{m}} = m) P(m | x) \\
&\quad P(m | \mathcal{D}_{\underline{x}} = x) \rightarrow P(m | x) \\
&\text{by Rule 2:} \quad \text{If } (\underline{b} \perp \underline{a} | \underline{r}, \underline{s}) \text{ in } \mathcal{L}_{\underline{a}} \mathcal{D}_{\underline{r}} G, \text{ then} \\
&\quad \mathcal{D}_{\underline{a}} = a. \leftrightarrow \underline{a} = a. \\
&\underline{m} \perp \underline{x} \text{ in } \mathcal{L}_{\underline{x}} \mathcal{D}_{\emptyset} G : \quad \begin{array}{c} \textcircled{\underline{c}} \\ \swarrow \quad \searrow \\ \underline{x} \quad \underline{m} \longrightarrow \underline{y} \end{array} \\
&= \sum_m P(y | \mathcal{D}_{\underline{m}} = m) P(m | x) \\
&\quad P(y | \mathcal{D}_{\underline{x}} = x, \mathcal{D}_{\underline{m}} = m) \rightarrow P(y | \mathcal{D}_{\underline{m}} = m) \\
&\text{by Rule 3:} \quad \text{If } (\underline{b} \perp \underline{a} | \underline{r}, \underline{s}) \text{ in } \mathcal{D}_{\underline{a} - an(\underline{s})} \mathcal{D}_{\underline{r}} G, \text{ then} \\
&\quad \mathcal{D}_{\underline{a}} = a. \leftrightarrow 1 \\
&\underline{y} \perp \underline{x} | \underline{m} \text{ in } \mathcal{D}_{\underline{x}} \mathcal{D}_{\underline{m}} G : \quad \begin{array}{c} \textcircled{\underline{c}} \\ \searrow \\ \underline{x} \quad \underline{m} \longrightarrow \underline{y} \end{array} \\
&= \sum_{x'} \sum_m P(y | \mathcal{D}_{\underline{m}} = m, x') P(x' | \mathcal{D}_{\underline{m}} = m) P(m | x) \\
&\quad \text{by Probability Axioms} \\
&= \sum_{x'} \sum_m P(y | m, x') P(x' | \mathcal{D}_{\underline{m}} = m) P(m | x) \\
&\quad P(y | \mathcal{D}_{\underline{m}} = m, x') \rightarrow P(y | m, x') \\
&\text{by Rule 2:} \quad \text{If } (\underline{b} \perp \underline{a} | \underline{r}, \underline{s}) \text{ in } \mathcal{L}_{\underline{a}} \mathcal{D}_{\underline{r}} G, \text{ then} \\
&\quad \mathcal{D}_{\underline{a}} = a. \leftrightarrow \underline{a} = a. \\
&\underline{y} \perp \underline{m} | \underline{x} \text{ in } \mathcal{L}_{\underline{m}} \mathcal{D}_{\emptyset} G : \quad \begin{array}{c} \textcircled{\underline{c}} \\ \swarrow \quad \searrow \\ \underline{x} \longrightarrow \underline{m} \quad \underline{y} \end{array} \\
&= \sum_{x'} \sum_m P(y | m, x') P(x') P(m | x) \\
&\quad P(x' | \mathcal{D}_{\underline{m}} = m) \rightarrow P(x') \\
&\text{by Rule 3:} \quad \text{If } (\underline{b} \perp \underline{a} | \underline{r}, \underline{s}) \text{ in } \mathcal{D}_{\underline{a} - an(\underline{s})} \mathcal{D}_{\underline{r}} G, \text{ then} \\
&\quad \mathcal{D}_{\underline{a}} = a. \leftrightarrow 1 \\
&\underline{x} \perp \underline{m} \text{ in } \mathcal{D}_{\underline{m}} \mathcal{D}_{\emptyset} G : \quad \begin{array}{c} \textcircled{\underline{c}} \\ \swarrow \quad \searrow \\ \underline{x} \quad \underline{m} \longrightarrow \underline{y} \end{array}
\end{aligned}$$

**QED**

**Claim 3** (*Napkin problem from Ref.[3]*)

If then



$$P(y|\mathcal{D}\underline{x} = x) = \sum_{w,z} P(y|x, w, z)P(w, z) \quad (1.13)$$

$$\mathcal{D}\underline{x} = x \longrightarrow y = \begin{array}{c} \sum w, z \\ \searrow \\ x \longrightarrow y \end{array} \quad (1.14)$$

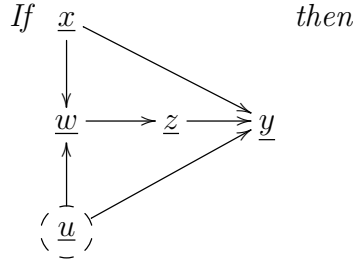
**proof:**

$$P(y|\mathcal{D}\underline{x} = x) = \sum_{u_1} P(y|x, u_1)P(u_1)$$

$$\begin{aligned}
 & \begin{array}{c} \sum u_1 \\ \swarrow \quad \searrow \\ \sum w \longrightarrow \sum z \quad \mathcal{D}\underline{x} = x \longrightarrow y \\ \nwarrow \quad \nearrow \\ \sum u_2 \end{array} = \begin{array}{c} \sum u_1 \\ \searrow \\ x \longrightarrow y \end{array} \\
 &= \sum_{w,z} \sum_{u_1} P(y|x, u_1)P(u_1|w, z)P(w, z) \\
 &= \sum_{w,z} \sum_{u_1} P(y|x, w, z)P(u_1|w, z)P(w, z) \\
 &= \sum_{w,z} P(y|x, w, z)P(w, z) \\
 &= \begin{array}{c} \sum w, z \\ \searrow \\ x \longrightarrow y \end{array}
 \end{aligned}$$

**QED**

**Claim 4** (from Ref.[3])



$$P(y|\mathcal{D}\underline{z} = z, x) = \sum_w P(y|z, x, w)P(w) \quad (1.15)$$

$$\begin{array}{c} x \\ \searrow \\ \mathcal{D}\underline{z} = z \end{array} \rightarrow y = \sum_w \begin{array}{c} x \\ \searrow \\ z \end{array} \rightarrow y \quad (1.16)$$

**proof:**

$$\begin{aligned} P(y|\mathcal{D}\underline{z} = z, x) &= \sum_u P(y|z, x, u)P(u) \\ \begin{array}{c} x \\ \downarrow \\ \sum w \\ \uparrow \\ (\sum u) \end{array} \begin{array}{c} \nearrow \\ \rightarrow \\ \nearrow \end{array} y &= \begin{array}{c} x \\ \searrow \\ z \end{array} \rightarrow y \\ &= \sum_w \sum_u P(y|z, x, u)P(u|w)P(w) \\ &= \sum_w \begin{array}{c} x \\ \searrow \\ z \\ \downarrow \\ (\sum u) \end{array} \rightarrow y \\ &= \sum_w P(y|z, x, w)P(w) \\ &= \sum_w \begin{array}{c} x \\ \searrow \\ z \end{array} \rightarrow y \end{aligned}$$

**QED**

**Claim 5** (*Trivial Transportability, from Ref.[2]*)

If  $\underline{s} \in \{0, 1\}$  is a selection node, then

$$\begin{array}{c}
 \begin{array}{ccc}
 & \underline{z} & \underline{s} \\
 & \swarrow \quad \searrow & \downarrow \\
 \underline{x} & \xrightarrow{\quad} & \underline{y} \\
 & \curvearrowright & 
 \end{array}
 \end{array}
 \quad P^*(y|\mathcal{D}\underline{x} = x, z) = P^*(y|x, z) \quad (1.17)$$

**proof:**

$$\begin{array}{ccc}
 P(y|\mathcal{D}\underline{x} = x, z, \underline{s} = 1) = P(y|x, z, \underline{s} = 1) \\
 \begin{array}{ccc}
 z & \underline{s} = 1 & z \\
 & \downarrow & \downarrow \\
 \mathcal{D}\underline{x} = x & \xrightarrow{\quad} & y = x \xrightarrow{\quad} y
 \end{array}
 \end{array}$$

**QED**

**Claim 6** (*Direct Transportability, a.k.a. External Validity, from Ref.[2]*)

If  $\underline{s} \longrightarrow \underline{z}$  where  $\underline{s} \in \{0, 1\}$  is a selection node, then

$$\begin{array}{c}
 \begin{array}{ccc}
 \underline{s} & \longrightarrow & \underline{z} \\
 & \swarrow \quad \searrow & \\
 \underline{x} & \xrightarrow{\quad} & \underline{y} \\
 & \curvearrowright & 
 \end{array}
 \end{array}
 \quad P^*(y|\mathcal{D}\underline{x} = x, z) = P(y|\mathcal{D}\underline{x} = x, z) \quad (1.18)$$

**proof:**

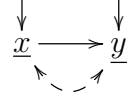
$$\begin{array}{ccc}
 P(y|\mathcal{D}\underline{x} = x, z, \underline{s} = 1) = P(y|\mathcal{D}\underline{x} = x, z) \\
 \begin{array}{ccc}
 \underline{s} = 1 & \longrightarrow & z \\
 & \searrow & \\
 \mathcal{D}\underline{x} = x & \xrightarrow{\quad} & y = \mathcal{D}\underline{x} = x \xrightarrow{\quad} y
 \end{array}
 \end{array}
 \quad \text{Because } \underline{s} \perp \underline{y} | \underline{z}$$

**QED**

**Claim 7** (*S-Admissable Transportability, from Ref.[2]*)



If  $\underline{s} \longrightarrow \underline{z} \longrightarrow \underline{a}$  where  $\underline{s} \in \{0, 1\}$  is a selection node, then



$$P^*(y|\mathcal{D}\underline{x} = x) = \sum_a P(y|\mathcal{D}\underline{x} = x, a)P^*(a) \quad (1.19)$$

$$\begin{array}{ccc} \underline{s} = 1 & & \underline{s} = 1 \longrightarrow \sum a \\ & \searrow & \downarrow \\ \mathcal{D}\underline{x} = x \longrightarrow y & = & \mathcal{D}\underline{x} = x \longrightarrow y \end{array} \quad (1.20)$$

**proof:**

$$P(y|\mathcal{D}\underline{x} = x, \underline{s} = 1) = \sum_a P(y|\mathcal{D}\underline{x} = x, a)P(a|\underline{s} = 1)$$

$$\begin{array}{ccc} \underline{s} = 1 \longrightarrow \sum z \longrightarrow \sum a & & \underline{s} = 1 \longrightarrow \sum a \\ & \downarrow & \downarrow \\ \mathcal{D}\underline{x} = x \longrightarrow y & = & \mathcal{D}\underline{x} = x \longrightarrow y \end{array}$$

**QED**

**Claim 8** (from Ref.[2])

If  $\underline{s} \longrightarrow \underline{z}$  where  $\underline{s} \in \{0, 1\}$  is a selection node, then



$$P^*(y|\mathcal{D}\underline{x} = x) = \sum_z P(y|\mathcal{D}\underline{x} = x, z)P^*(z) \quad (1.21)$$

$$\begin{array}{ccc} \underline{s} = 1 & & \underline{s} = 1 \longrightarrow \sum z \\ & \searrow & \searrow \\ \mathcal{D}\underline{x} = x \longrightarrow y & = & \mathcal{D}\underline{x} = x \longrightarrow y \end{array} \quad (1.22)$$

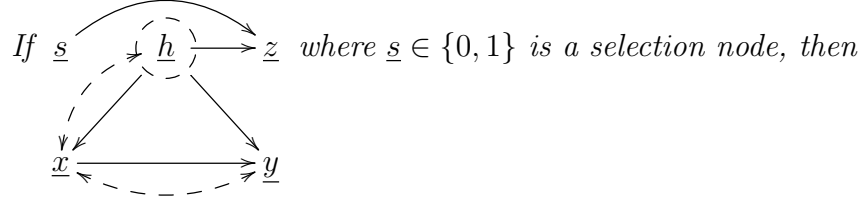
**proof:**

$$P(y|\mathcal{D}\underline{x} = x, \underline{s} = 1) = \sum_z P(y|\mathcal{D}\underline{x} = x, z)P(z|\underline{s} = 1)$$

$$\begin{array}{ccc} \underline{s} = 1 \longrightarrow \sum z & & \\ & \searrow & \\ \mathcal{D}\underline{x} = x \longrightarrow y & & \end{array}$$

**QED**

**Claim 9** (from Ref.[2])



$$P^*(y|\mathcal{D}\underline{x} = x) = P(y|\mathcal{D}\underline{x} = x) \quad (1.23)$$

$$\begin{array}{c} \underline{s} = 1 \\ \searrow \\ \mathcal{D}\underline{x} = x \longrightarrow y = \mathcal{D}\underline{x} = x \longrightarrow y \end{array} \quad (1.24)$$

**proof:**

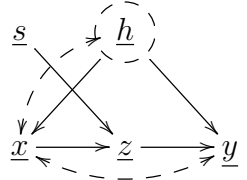
$$P(y|\mathcal{D}\underline{x} = x, \underline{s} = 1) = \sum_h P(y|\mathcal{D}\underline{x} = x, h)P(h)$$

$$\begin{array}{c} \underline{s} = 1 \xrightarrow{\quad} (\sum h) \xrightarrow{\quad} \sum z \\ \quad \quad \quad \searrow \\ \mathcal{D}\underline{x} = x \longrightarrow y = \mathcal{D}\underline{x} = x \longrightarrow y \\ = \sum_h P(y|\mathcal{D}\underline{x} = x, h)P(h|\mathcal{D}\underline{x} = x) \end{array}$$

$$\begin{array}{c} \quad \quad \quad (\sum h) \\ \quad \quad \quad \nearrow \quad \searrow \\ = \mathcal{D}\underline{x} = x \longrightarrow y \\ = P(y|\mathcal{D}\underline{x} = x) \\ = \mathcal{D}\underline{x} = x \longrightarrow y \end{array}$$

**QED**

**Claim 10** (from Ref.[2])

If  $\underline{s}$   where  $\underline{s} \in \{0, 1\}$  is a selection node, then

$$P^*(y|\mathcal{D}\underline{x} = x) = \sum_z P(y|\mathcal{D}\underline{x} = x, z)P^*(z|x) \quad (1.25)$$

$$\begin{array}{c} \underline{s} = 1 \quad \mathcal{D}\underline{x} = x \quad \underline{s} = 1 \quad \mathcal{D}\underline{x} = x \\ \downarrow \quad \downarrow \quad \downarrow \quad \downarrow \\ y \quad = \quad x \longrightarrow \sum z \longrightarrow y \end{array} \quad (1.26)$$

**proof:**

$$P(y|\mathcal{D}\underline{x} = x, \underline{s} = 1) = \sum_h \sum_z P(y|h, z)P(h)P(z|\mathcal{D}\underline{x} = x, \underline{s} = 1)$$

$$\begin{array}{c} \underline{s} = 1 \quad (\sum h) \\ \downarrow \quad \downarrow \\ \mathcal{D}\underline{x} = x \longrightarrow \sum z \longrightarrow y \end{array}$$

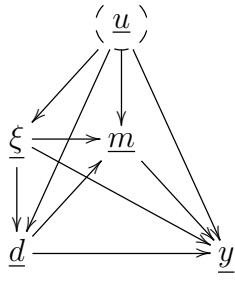
$$= \sum_h \sum_z P(y|h, z)P(h|\mathcal{D}\underline{x} = x)P(z|x, \underline{s} = 1)$$

$$\begin{array}{c} \underline{s} = 1 \quad (\sum h) \longleftarrow \mathcal{D}\underline{x} = x \\ \downarrow \quad \downarrow \\ = \quad x \longrightarrow \sum z \longrightarrow y \\ = \sum_z P(y|\mathcal{D}\underline{x} = x, z)P(z|x, \underline{s} = 1) \\ \underline{s} = 1 \quad \mathcal{D}\underline{x} = x \\ \downarrow \quad \downarrow \\ = \quad x \longrightarrow \sum z \longrightarrow y \end{array}$$

**QED**

**Claim 11** (from Ref.[1])

If  $(\underline{u})$  then



$$P(y|\mathcal{D}\underline{d} = d, \mathcal{I}\underline{m}d = d') = \sum_{\xi} \sum_m P(y|d, m, \xi) P(m|d', \xi) P(\xi) \quad (1.27)$$

$$\begin{array}{ccc} \mathcal{I}\underline{d} = d' & \mathcal{I}\underline{d} = d' & \sum \xi \longrightarrow \sum m \\ \searrow & \searrow & \searrow \\ \mathcal{D}\underline{d} = d \longrightarrow y & = & \mathcal{D}\underline{d} = d \longrightarrow y \end{array} \quad (1.28)$$

**proof:**

$$P(y|\mathcal{D}\underline{d} = d, \mathcal{I}\underline{d} = d') = \sum_{\xi, u} \sum_m P(y|d, m, \xi, u) P(m|d', \xi, u) \underbrace{P(\xi|u)P(u)}_{P(\xi, u)}$$

$$\begin{array}{ccc} & (\sum u) & \\ & \downarrow & \\ \mathcal{I}\underline{d} = d' & \sum \xi \longrightarrow \sum m & \\ \searrow & \searrow & \searrow \\ \mathcal{D}\underline{d} = d \longrightarrow y & & y \\ = \sum_{\xi} \sum_m P(y|d, m, \xi) P(m|d', \xi) P(\xi) & & \end{array}$$

$$\begin{array}{ccc} \mathcal{I}\underline{d} = d' & \sum \xi \longrightarrow \sum m & \\ \searrow & \searrow & \searrow \\ = \mathcal{D}\underline{d} = d \longrightarrow y & & y \end{array}$$

We switch from averaging over the prior of  $\xi, u$  to averaging over the prior of  $\xi$ .

**QED**

# Bibliography

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