



# Analisi e Visualizzazione delle Reti Complesse

**NS19 - Link Analysis and Web  
Search**

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# Agenda

- Searching the Web
- Link Analysis
  - HITS: Hubs and Authorities (+ Spectral Analysis)
  - Page Rank (+ Spectral Analysis)
  - Random Walks and PR
- Practical implications
  - Modern Web search
  - Link Analysis beyond the Web

# PageRank and HITS

- **Centrality measures** for nodes in directed networks
- Sergey Brin and Larry Page introduced PageRank in 1998 as a key ingredient of Google
- Jon Kleinberg introduced HITS in 1999
- Both are based on eigenvector centrality and designed for web information retrieval
- In NetworkX:

```
PR_dict = nx.pagerank(D)      # D must be a DiGraph
H_dict, A_dict = nx.hits(G)    # G should be a DiGraph
```

## Reading material:

- [Page, Lawrence; Brin, Sergey; Motwani, Rajeev and Winograd, Terry, The PageRank citation ranking: Bringing order to the Web. 1999](#)
- [Jon Kleinberg, Authoritative sources in a hyperlinked environment Journal of the ACM 46 \(5\): 604-32, 1999. doi:10.1145/324133.324140.](#)
- [A. Langville and C. Meyer, "A survey of eigenvector methods of web information retrieval.", SIAM Review, vol. 47, No. 1](#)

# Searching the Web

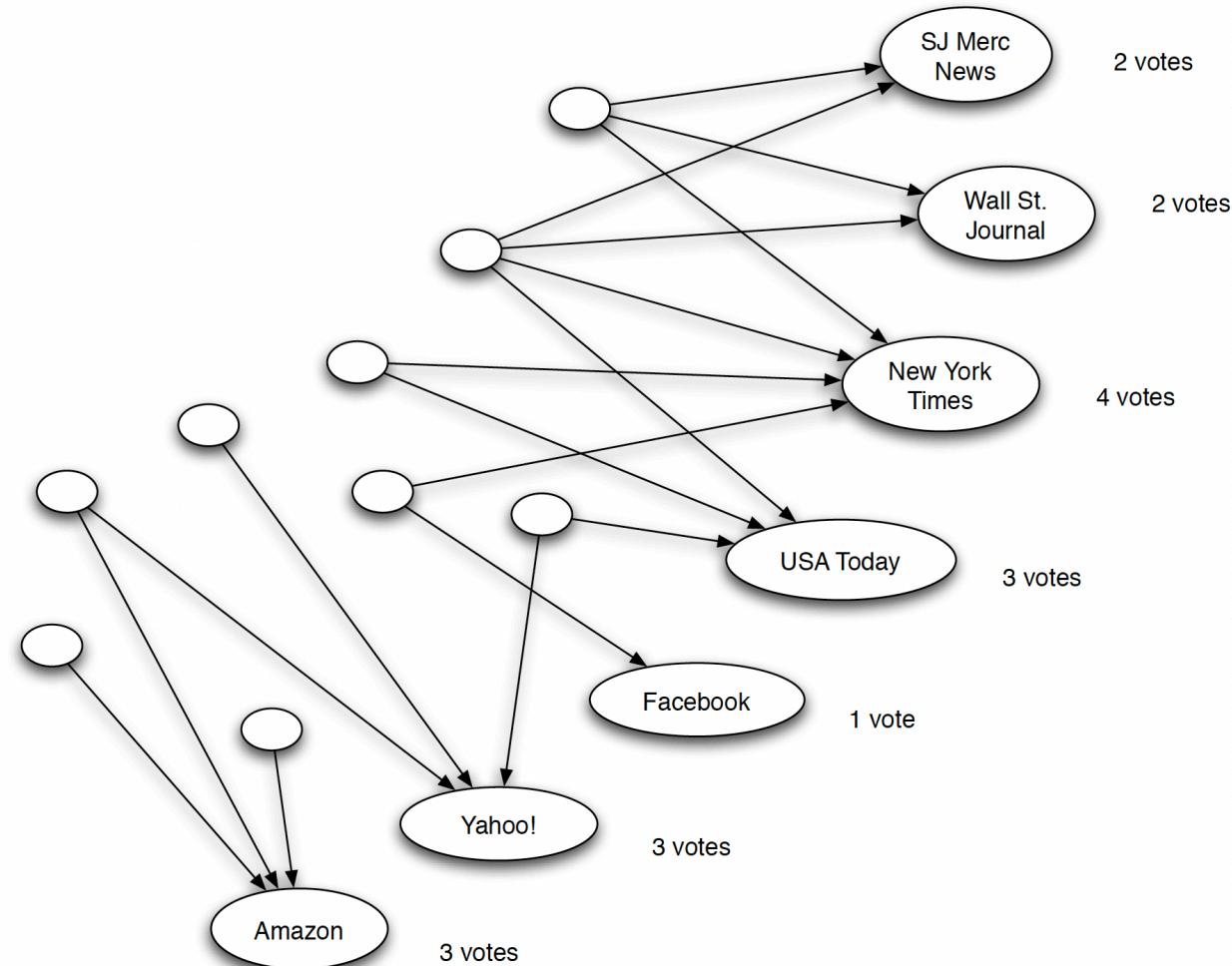
- **Search Engine**
  - problem: how to rank (web) pages related to a given topic
- **Information Retrieval**
  - automated strategies to search in libraries, scientific papers, repositories, and so on
  - in response to keywords based queries
- List of keywords is "inexpressive" (e.g., polysemy, synonymy)
- **Diversity:** given a topic we find pages created by a large variety of authors
  - e.g., experts, novices, children, conspiracy theorists
- Pages are **dynamic** and always changing
- **Filters:** abundance of information, what is important?
- Can the structure of the Web, dominated by links, help us to find such filters?
  - first attempt: count words in documents
  - can we do better?

# HITS: Link Analysis using Hubs and Authorities

- Disclaimer: the term "Hub" is used slightly differently here w.r.t. previous lectures
- Information contained "between" pages can be used as well
- Count **in-links**:
  - select documents on a given topic
  - **in-links are a measure of authority** of a page on such a topic: it is an implicit endorsement from the community of web pages' authors
  - It is hard for search engines to automatically assess the intent of each link. In aggregate, we assume they mean endorsement.

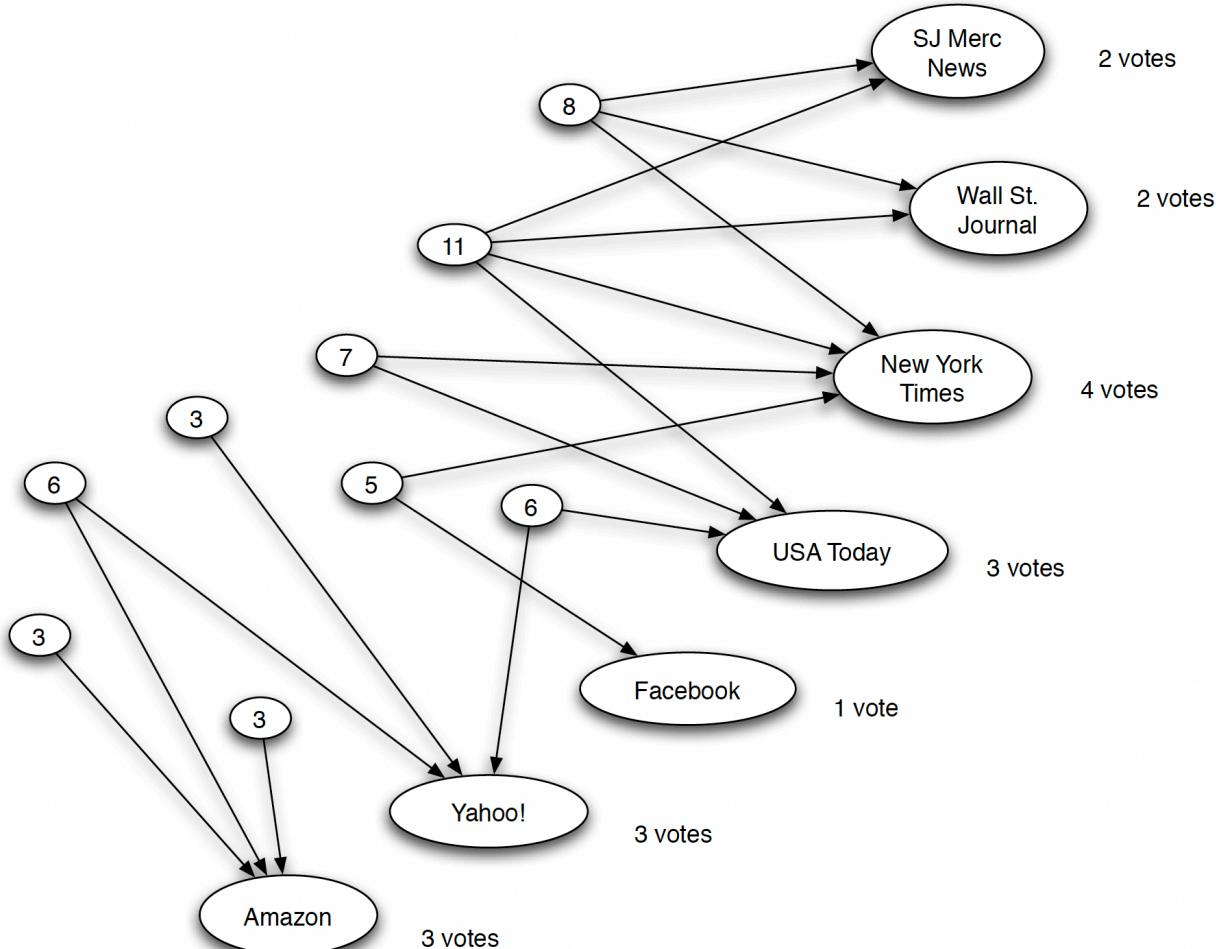
## Finding lists

- query: "newspapers"
- we could have intuitively valid authority values with **in-degrees**
  - however, we get a mix of prominent newspapers along with pages that are going to receive a lot of in-links no matter what the query is
- **Lists:** pages that provide many different out-links to other pages
  - among the pages casting votes, a few of them voted for many of the pages that received a lot of votes



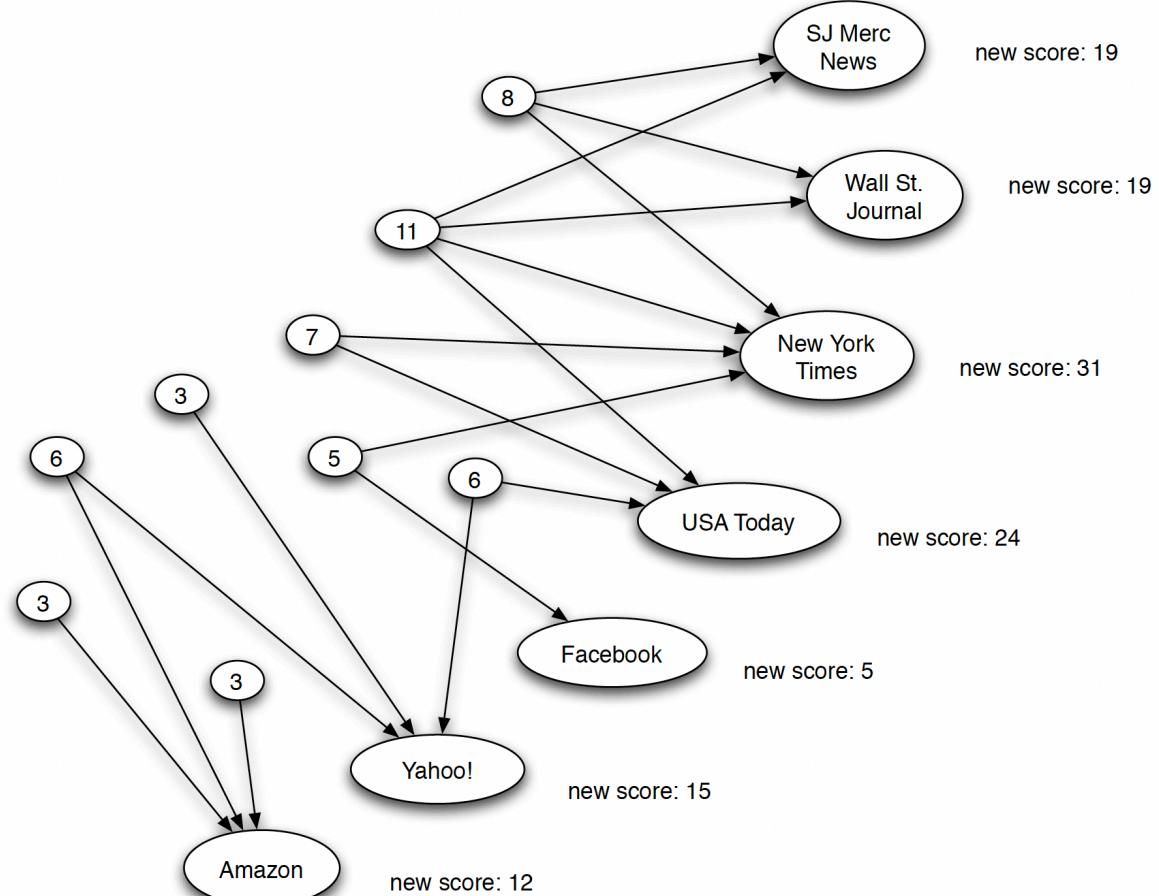
## List values

- Hub value for a list page:  
**the sum of in-links received by all pages it linked to**
  - in other words: the sum of the votes received by all pages that it voted for
- Assumptions:
  - list pages have a better sense for where the good results are
  - authorities are often **competitors!**



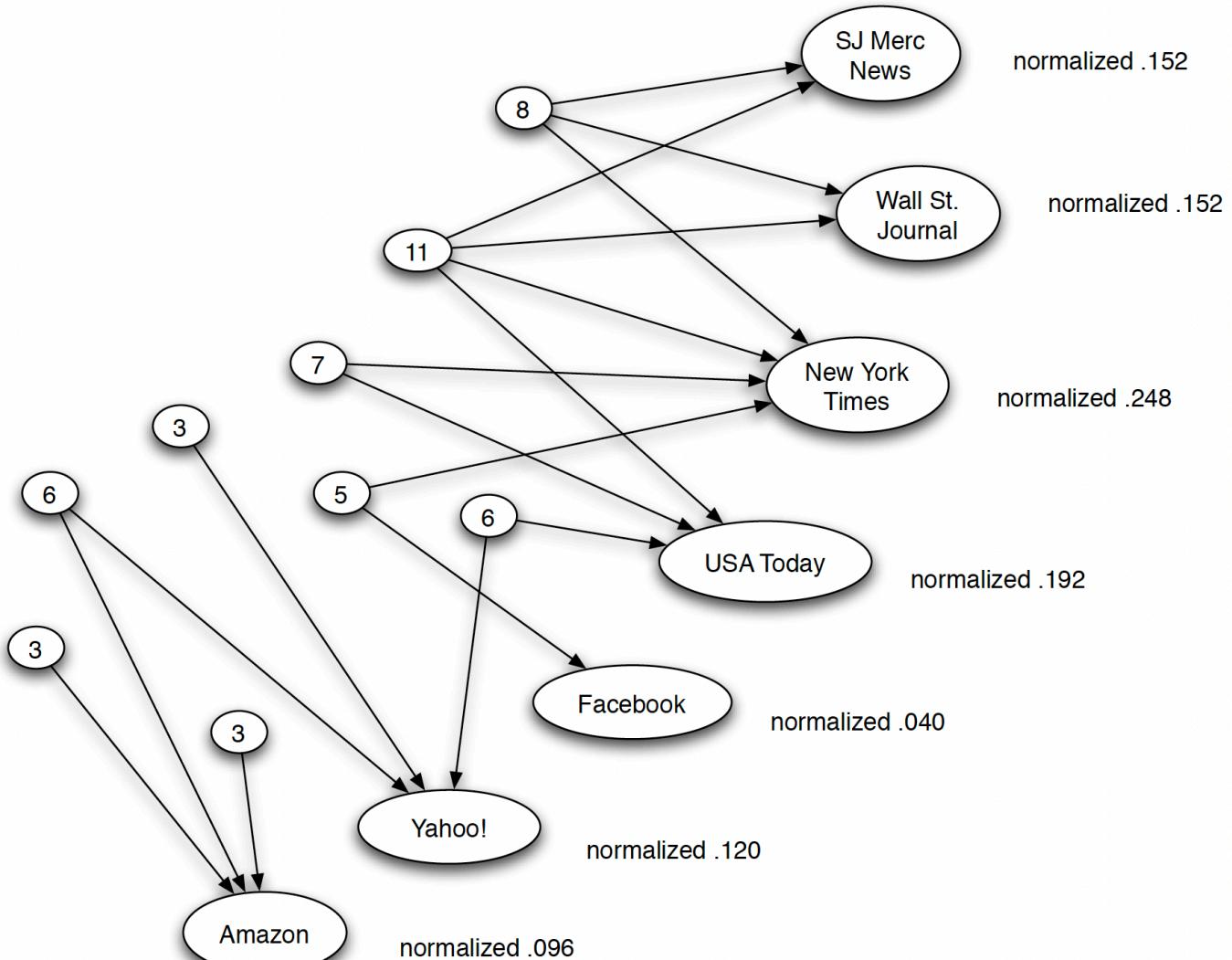
# The Principle of Repeated Improvement

- If we think that high-scoring list pages (hubs) have a better sense of where good results are:
  - we can **weight links from hubs more heavily**
- Recalculate
  - the authority score is not just the number of in-links, it is the sum of the hub values of the in-link pages
- **Why stop here?** we can refine values at both sides



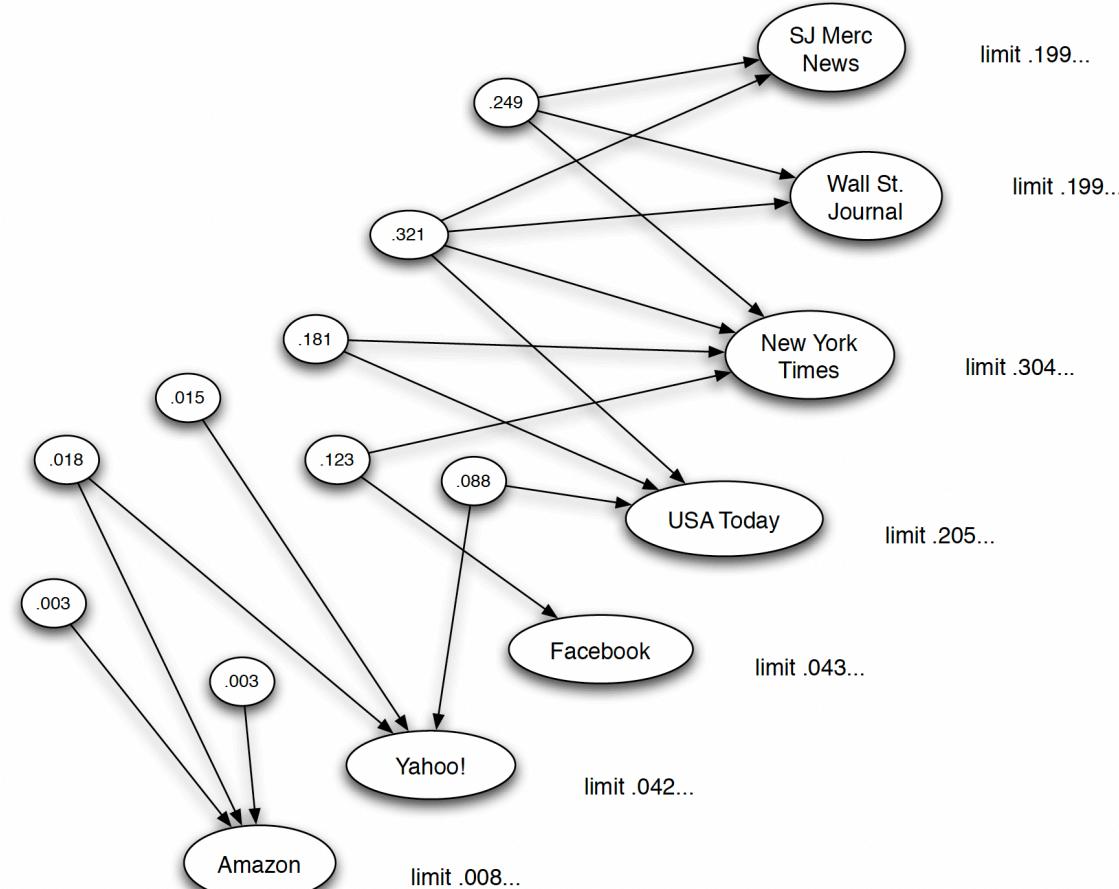
# Hubs and Authorities

- Each page  $p$  has an **authority score**  $auth(p)$  and **hub score**  $hub(p)$  initialized to 1
- **Authority Update Rule:** for each page  $p$ , update  $auth(p)$  to be the sum of the hub scores of all pages that point to it.
- **Hub Update Rule:** for each page  $p$ , update  $hub(p)$  to be the sum of the authority scores of all pages that it points to.
- At step  $k$ , we performed  $k$  hub-authority updates. Each step involves:
  - First apply the **Authority Update Rule** to the current set of scores.
  - Then apply the **Hub Update Rule** to the resulting set of scores.
- Getting high values: **normalize hub and authority values** by dividing down the scores by the sum of the scores of authorities/hubs.



# Stabilization

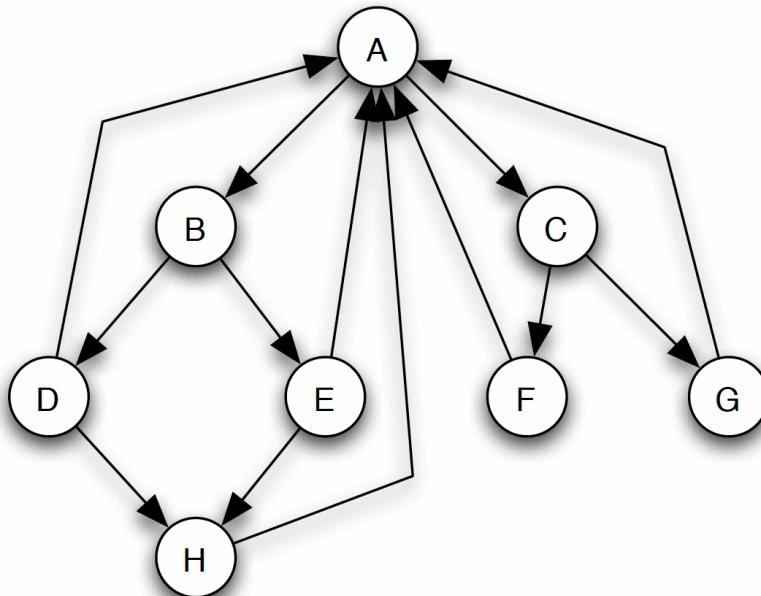
- It can be proved that normalized values **converge** when  $k \rightarrow \infty$ 
  - **(optional proof)** see Section 14.6  
**Advanced Material: Spectral Analysis, Random Walks, and Web Search**
- **Initial values are not important**
- Limiting values for hubs and authorities are **properties of the links structure**.
- Different form of game theoretical concept of equilibrium
  - **authority score** is proportional to the hub scores of the pages that point to you
  - **hub score** is proportional to the authority scores of the pages you point to.



# Page Rank

- **Endorsement** viewed as passing directly from one important node to another
  - in other words, a page is important if it is cited by other important pages
  - endorsements received by in-links and passed across outgoing links
- **Basic definition:**
  - **Initialization:** Init all the pages  $p$  to a  $PR(p) = \frac{1}{n}$ , where  $n$  is the number of pages
  - **$k^{th}$  step:** we perform a sequence of  $k$  updates to the PageRank values, using the following rule for each update:
  - **Basic PageRank Update Rule:**
    - each page divides its current PageRank equally across its out-going links, and passes these equal shares to the pages it points to.
    - if a page has no out-going links, it passes all its current PageRank to itself.
    - each page updates its new PageRank to be the sum of the shares it receives.

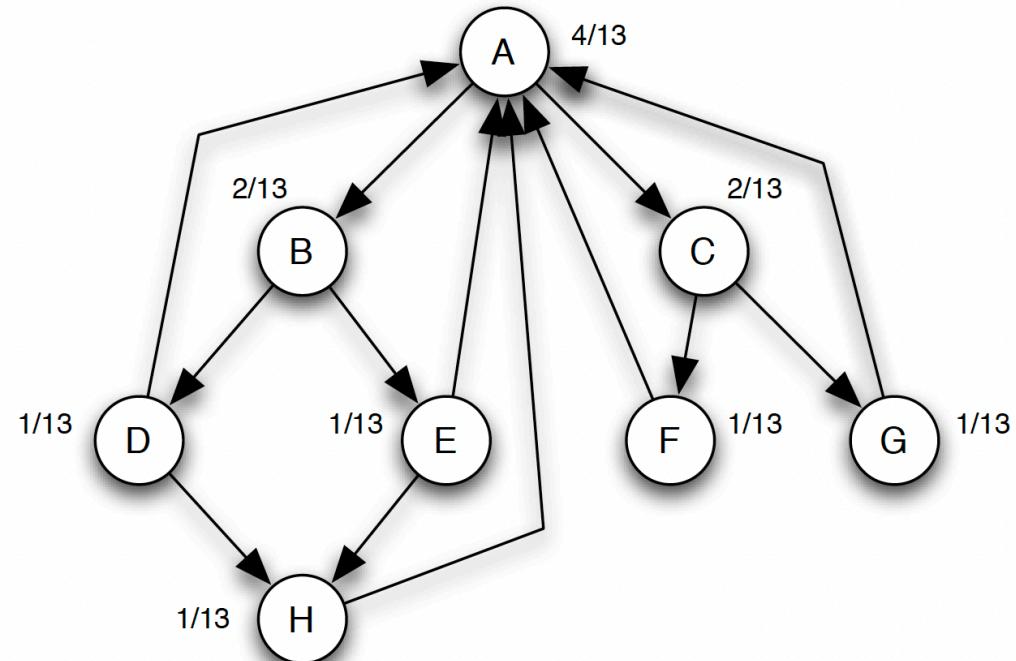
## Example



Step	A	B	C	D	E	F	G	H
1	$1/2$	$1/16$	$1/16$	$1/16$	$1/16$	$1/16$	$1/16$	$1/8$
2	$3/16$	$1/4$	$1/4$	$1/32$	$1/32$	$1/32$	$1/32$	$1/16$

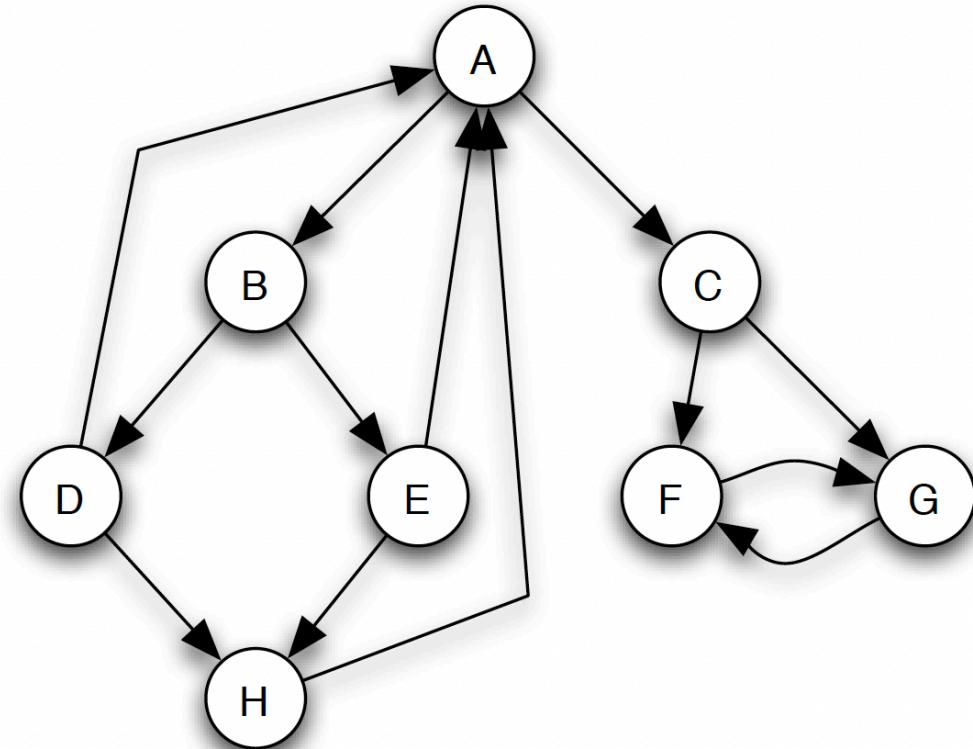
## PageRank and stabilization

- *PR* values of all the nodes **converge** when  $k \rightarrow \infty$ 
  - but for some degenerate cases.
- **Equilibrium:** if we apply our *PR* update rule, then our limiting **values do not change**.
- One can prove that if the network is **strongly connected**, then there is a **single set of equilibrium values** (when they exist)
  - (**optional proof**): see Section 14.6 [Advanced Material: Spectral Analysis, Random Walks, and Web Search](#)



## Scaling the definition of PageRank

- Which are these degenerate cases?
- **Problem:** in some networks some nodes receive all the *PR* values of the the network
  - **example on the right:** we converge to *PR* values of  $\frac{1}{2}$  for each of *F* and *G*, and 0 for all other nodes
- Why?
  - We do not have path back to some other nodes



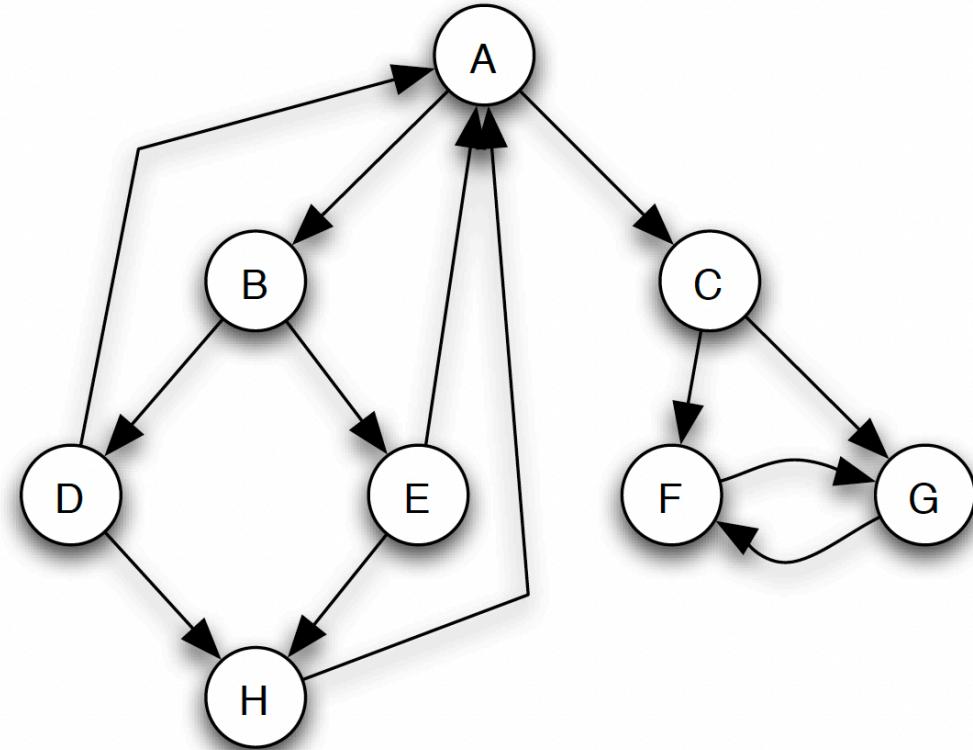
- **Solution:** let's force this "fluid" to stream back to other nodes "sometimes":
- Select a **scaling factor** (aka **damping factor**)  $s : s \in [0, 1]$
- Get a portion  $s$  of  $PR$  values from in-links and then add  $\frac{(1-s)}{n}$ 
  - This means that the total  $PR$  in the network has shrunk from 1 to  $s$ . We divide the residual  $1 - s$  units of  $PR$  equally over all nodes, giving  $\frac{(1-s)}{n}$  to each.
  - Called *Scaled PageRank Update Rule*
- Now we have convergence for  $k \rightarrow \infty$
- **Note:** typically  $s \in [0.8, 0.9]$

## Random walks: An equivalent definition of PageRank

- Randomly browsing a network of Web pages, navigating each page with equal probability following links.
  - start by choosing a page at random
  - pick a random out-going link from their current page, and follow it to where it leads
- Follow links for a sequence of length  $k$
- **Claim:** the probability of being at a page  $x$  after  $k$  steps of this random walk is precisely the  $PR$  of  $x$  after  $k$  applications of the Basic PageRank Update Rule.
- **Additional intuition:**  $PR(x)$  is the limiting probability that a random walk across hyperlinks will end up at  $x$  as we sum the walk for larger and larger number of steps

## Leakage

- The leakage of  $F$  and  $G$  has a natural interpretation: when the surfer reaches  $F$  or  $G$ , then it is stuck forever.
- **Solution:**
  - with probability  $s$ , the random walker clicks on an hyperlink in the page.
  - with probability  $1 - s$ , it jumps to a randomly selected node.



# Practical implications (also beyond the Web)

## Modern Web search

- Google today use *PR* as one of the many features of their ranking framework
  - original Page&Brin's paper [The Anatomy of a Large-Scale Hypertextual Web Search Engine](#).
  - e.g., [Hilltop](#) (an extension of HITS) has been probably used for a while.
- **anchor texts**
- **clicking behavior**
- and much more (and who knows what they are actually using!)

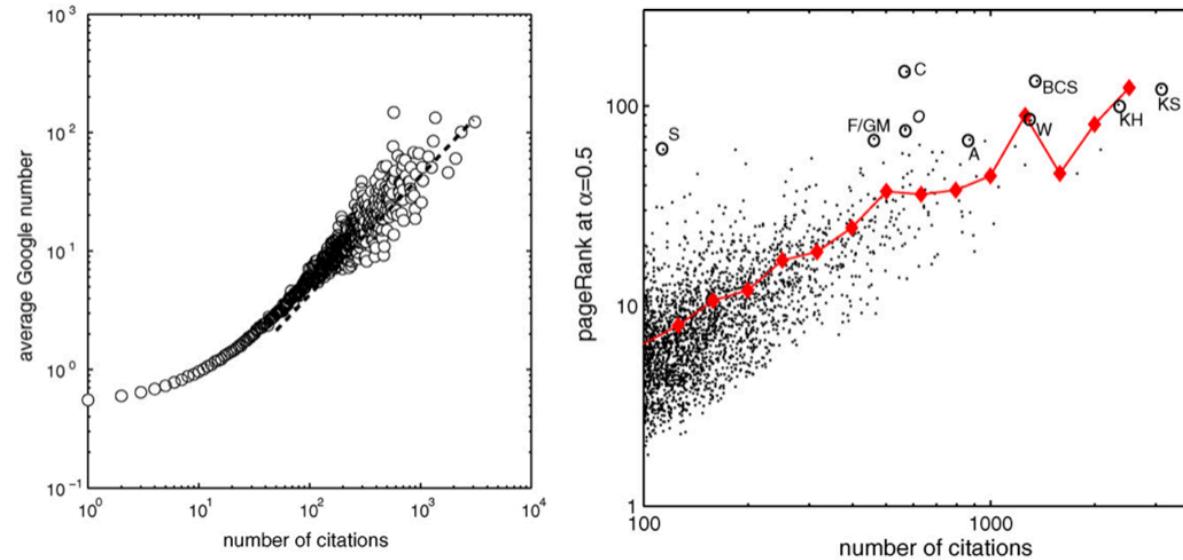
## SEO vs Google

- SEO: Search Engine Optimization
- SEO companies: reverse engineering of search engine's ranking functions
- SE companies: define new measures
  - ... in loop!
- **Feedback effects:** perfect results are "moving targets"
- It is a game theoretic principle.
  - you should always expect the world to react to what you do.

## Applications beyond the Web

- Citation Analysis
- Link Analysis of U.S. Supreme Court Citations

# Page Rank and citation analysis



- **Reference:** [Finding Scientific Gems with Google Page Ranks \(2007\)](#).
- **Dataset:** collection of scientific papers with their references
- positive correlation between number of citations and average *PR* values
- BUT outliers are papers with a limited number of citations but highly influential anyhow

Google rank	Google # ( $\times 10^{-4}$ )	Cite rank	# cites	Publication			Title		Author(s)
1	4.65	54	574	PRL	10	531	1963	Unitary symmetry and leptonic...	N. Cabibbo
2	4.29	5	1364	PR	108	1175	1957	Theory of superconductivity	J. Bardeen, L. Cooper, and J. Schrieffer
3	3.81	1	3227	PR	140	A1133	1965	Self-consistent equations...	W. Kohn and L.J. Sham
4	3.17	2	2460	PR	136	B864	1964	Inhomogeneous electron gas	P. Hohenberg and W. Kohn
5	2.65	6	1306	PRL	19	1264	1967	A model of leptons	S. Weinberg
6	2.48	55	568	PR	65	117	1944	Crystal statistics I	L. Onsager
7	2.43	56	568	RMP	15	1	1943	Stochastic problems in...	S. Chandrasekhar
8	2.23	95	462	PR	109	193	1958	Theory of the Fermi interaction	R.P. Feynman and M. Gell-Mann
9	2.15	17	871	PR	109	1492	1958	Absence of diffusion in...	P.W. Anderson
10	2.13	1853	114	PR	34	1293	1929	The theory of complex spectra	J.C. Slater

**Pros:** *PR* helps to find "gems" in networks!

**Cons:** Indicators can change our behaviors

# Introduction to spectral analysis

- Pre-requisites:
  - linear algebra
  - vector and matrix multiplication
- Eigenvalues/eigenvectors calculation to study the structure of networks
  - spectral analysis
- Limiting values in the **Hubs/Authorities** and **PageRank** algorithms are coordinates in the eigenvectors for given eigenvalues in matrices derived from our graphs.

# Definitions

## Adjacency Matrix $M$

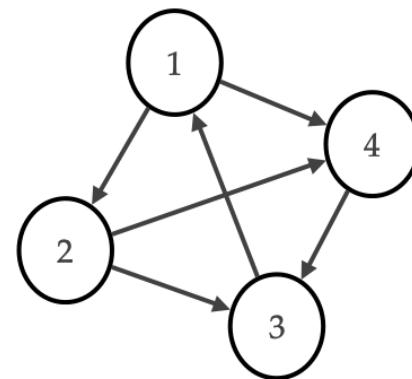
1 ...  $n$  nodes

$M : n \times n$

$$M_{ij} = \begin{cases} 1, & \text{if } (i, j) \text{ is an edge} \\ 0, & \text{otherwise} \end{cases}$$

The adjacency matrix is not necessarily efficient for computational representations

- but conceptually very useful
- for practical use, consider adjacency lists or edge lists instead



$$\begin{bmatrix} 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

## Update rules of Hubs and Authorities

$\mathbf{h}, \mathbf{a}$ : n-dimensional vectors (respectively, hub and authorities values)

### Hub Update Rule

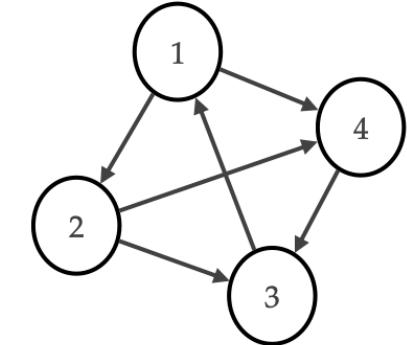
$$h_i \leftarrow \sum_{j=1}^n M_{ij} a_j = M_{i1} a_1 + M_{i2} a_2 + \dots + M_{in} a_n$$

$$\mathbf{h} \leftarrow \mathbf{M} \cdot \mathbf{a}$$

### Authority Update Rule

$$a_i \leftarrow \sum_{j=1}^n M_{ji} h_j$$

$$\mathbf{a} \leftarrow \mathbf{M}^T \cdot \mathbf{h}$$



$$\begin{bmatrix} 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix} \cdot \begin{bmatrix} 1 \\ 1 \\ 2 \\ 2 \end{bmatrix} = \begin{bmatrix} 3 \\ 4 \\ 1 \\ 2 \end{bmatrix}$$

$\mathbf{M}$              $\mathbf{a}$              $\mathbf{h}$

## Understanding the k-step hub-authority computation

initialization:  $\mathbf{h}^{<0>} = \underbrace{(1, 1, \dots, 1)}_n$

after  $k$  applications of the rule:  $\mathbf{a}^{<k>}, \mathbf{h}^{<k>}$

# Understanding the k-step hub-authority computation

## First step

$$\mathbf{a}^{<1>} = \mathbf{M}^T \mathbf{h}^{<0>}$$

$$\mathbf{h}^{<1>} = \mathbf{M} \mathbf{a}^{<1>} = \mathbf{M} \mathbf{M}^T \mathbf{h}^{<0>}$$

## Second step

$$\mathbf{a}^{<2>} = \mathbf{M}^T \mathbf{h}^{<1>} = (\mathbf{M}^T \mathbf{M}) \mathbf{M}^T \mathbf{h}^{<0>}$$

$$\mathbf{h}^{<2>} = \mathbf{M} \mathbf{a}^{<2>} = (\mathbf{M} \mathbf{M}^T) (\mathbf{M} \mathbf{M}^T) \mathbf{h}^{<0>} = (\mathbf{M} \mathbf{M}^T)^2 \mathbf{h}^{<0>}$$

## $k^{th}$ step

$$\mathbf{a}^{<k>} = (\mathbf{M}^T \mathbf{M})^{k-1} \mathbf{M}^T \mathbf{h}^{<0>}$$

$$\mathbf{h}^{<k>} = (\mathbf{M} \mathbf{M}^T)^k \mathbf{h}^{<0>}$$

**a, h** vectors: multiplication of an initial vector  $\mathbf{h}^{<0>}$  by larger powers of  $\mathbf{M}^T \mathbf{M}$  and  $\mathbf{M} \mathbf{M}^T$ .

## Multiplications and Eigenvectors (optional)

**normalization:** we can find constants  $c$  and  $d$  s.t.  $\frac{\mathbf{h}^{<k>}}{c^k}$  and  $\frac{\mathbf{a}^{<k>}}{d^k}$ .

We want to prove that they converge for  $k \rightarrow \infty$

Focus on the hub vectors:

If  $\frac{\mathbf{h}^{<k>}}{c^k} = \frac{(\mathbf{M}\mathbf{M}^T)^k \cdot \mathbf{h}^{<0>}}{c^k}$  converges to a limit  $\mathbf{h}^{<*>}$ , then I can expect that:

$$c \cdot \mathbf{h}^{<*>} = (\mathbf{M}\mathbf{M}^T) \cdot \mathbf{h}^{<*>}$$

Hence, we need to prove that the sequence of  $\frac{\mathbf{h}^{<k>}}{c^k}$  converges to the eigenvector of  $\mathbf{M}\mathbf{M}^T$

## Eigenvectors and square matrices (optional)

- Observe that  $\mathbf{M}\mathbf{M}^T$  is a symmetric matrix
- Fact 1: "Any symmetric matrix  $n \times n$  has a set of  $n$  eigenvectors  $\mathbf{z}_1, \dots, \mathbf{z}_n$  that are orthogonal and all unit vectors - that is they form a basis for the space  $\mathbb{R}^n \implies \mathbf{z}_i \cdot \mathbf{z}_j = 0$  and  $\mathbf{z}_i \cdot \mathbf{z}_i = 1$ "
- That means that for our symmetric  $\mathbf{M}\mathbf{M}^T$  we can find:
  - $n$  mutual orthogonal eigenvectors:  $\mathbf{z}_1, \dots, \mathbf{z}_n$  (the spectrum of  $\mathbf{M}\mathbf{M}^T$ )
  - $n$  corresponding eigenvalues:  $c_1, \dots, c_n$
- Let's sort eigenvectors s.t. corresponding eigenvalues:  $c_1 \geq c_2 \geq \dots \geq c_n$
- Assume (for now):  $c_1 > c_2$

## Eigenvectors and square matrices (optional)

- Let's consider  $\mathbf{x} \in \mathbb{R}^n$

$$\begin{aligned}(\mathbf{M}\mathbf{M}^T)\mathbf{x} &= (\mathbf{M}\mathbf{M}^T)(p_1\mathbf{z}_1 + p_2\mathbf{z}_2 + \dots + p_n\mathbf{z}_n) \\&= p_1(\mathbf{M}\mathbf{M}^T)\mathbf{z}_1 + p_2(\mathbf{M}\mathbf{M}^T)\mathbf{z}_2 + \dots + p_n(\mathbf{M}\mathbf{M}^T)\mathbf{z}_n \\&= p_1c_1\mathbf{z}_1 + p_2c_2\mathbf{z}_2 + \dots + p_nc_n\mathbf{z}_n\end{aligned}$$

- We will use this equation to analyze multiplication by larger powers of  $(\mathbf{M}\mathbf{M}^T)$

$$(\mathbf{M}\mathbf{M}^T)^k\mathbf{x} = p_1c_1^k\mathbf{z}_1 + p_2c_2^k\mathbf{z}_2 + \dots + p_nc_n^k\mathbf{z}_n$$

## Convergence of the Hub-Authority computation (optional)

vector of hub scores at step  $k$ :

$$\mathbf{h}^{<\mathbf{k}>} = (\mathbf{M}\mathbf{M}^T)^k \cdot \mathbf{h}^{<0>}$$

$$\mathbf{h}^{<0>} = q_1 \mathbf{z}_1 + q_2 \mathbf{z}_2 + \dots + q_n \mathbf{z}_n$$

$$\mathbf{h}^{<\mathbf{k}>} = c_1^k q_1 \mathbf{z}_1 + c_2^k q_2 \mathbf{z}_2 + \dots + c_n^k q_n \mathbf{z}_n$$

Let's divide both sides by  $c_1^k$ :

$$\frac{\mathbf{h}^{<\mathbf{k}>}}{c_1^k} = \frac{c_1^k q_1 \mathbf{z}_1}{c_1^k} + \frac{c_2^k q_2 \mathbf{z}_2}{c_1^k} + \dots + \frac{c_n^k q_n \mathbf{z}_n}{c_1^k}$$

$$\text{assumption: } c_1 > c_2 \Rightarrow \lim_{k \rightarrow \infty} \left( \frac{c_2}{c_1} \right)^k = 0$$

Then:

$$\lim_{k \rightarrow \infty} \frac{\mathbf{h}^{<\mathbf{k}>}}{c_1^k} = q_1 \mathbf{z}_1$$

## Wrapping up (optional)

- (i) a limit in the direction of  $\mathbf{z}_1$  is reached regardless of initial values of  $\mathbf{h}^{<0>}$ :

let's suppose that  $\mathbf{h}^{<0>} = \mathbf{x}$  and that is a positive vector:

$$\mathbf{x} = p_1\mathbf{z}_1 + p_2\mathbf{z}_2 + \dots + p_n\mathbf{z}_n \Rightarrow (\mathbf{M}\mathbf{M}^T)^k \mathbf{x} = c_1^k p_1\mathbf{z}_1 + c_2^k p_2\mathbf{z}_2 + \dots + c_n^k p_n\mathbf{z}_n$$

$$\lim_{k \rightarrow \infty} \frac{\mathbf{h}^{<k>}}{c_1^k} = p_1\mathbf{z}_1$$

- (ii) coefficient  $p_1$  (or  $q_1$ ) must be  $\neq 0$ : assuring that  $p_1\mathbf{z}_1$  (or  $q_1\mathbf{z}_1$ ) are non zero vectors, in the direction of  $\mathbf{z}_1$

## Wrapping up (optional)

- (iii) relax assumption:  $c_1 > c_2$

in general we can have  $l > 1$  eigenvalues s.t.  $c_1 = c_2 = \dots = c_l$  until we find that  $c_1 > c_{l+1}$

$$\begin{aligned}\frac{\mathbf{h}^{<\mathbf{k}>}}{c_1^k} &= \frac{c_1^k q_1 \mathbf{z}_1 + c_2^k q_2 \mathbf{z}_2 + \dots + c_l^k q_l \mathbf{z}_l}{c_1^k} + \frac{c_{l+1} n^k q_{l+1} \mathbf{z}_{l+1} + \dots + c_n^k q_n \mathbf{z}_n}{c_1^k} \\ &= q_1 \mathbf{z}_1 + q_2 \mathbf{z}_2 + \dots + q_l \mathbf{z}_l + 0\end{aligned}$$

with  $k \rightarrow \infty$  is still a convergence

- (iv) authority values: the argument is very similar to hub values (multiplication by  $\mathbf{M}^T \mathbf{M}$ )

# Spectral Analysis of Page Rank

At step 0 (init):

$$\forall i : r_i = \frac{1}{n}; \quad n: \# \text{ pages} \quad \text{right} > r_i = PR(i) </\text{span}>$$

At step k:

$$\forall i : r_i = \sum_{j=1}^n M_{ji} \frac{r_j}{k_j^{\text{out}}} \quad \text{right} > (\text{basic PR update rule}) </\text{span}>$$

$$\forall i : r_i = s \cdot \sum_{j=1}^n M_{ji} \frac{r_j}{k_j^{\text{out}}} + (1 - s) \cdot \frac{1}{n} \quad \text{right} > (\text{scaled PR update rule}) </\text{span}>$$

## Using matrix notation

$\mathbf{N}$ : Matrix derived from  $M$

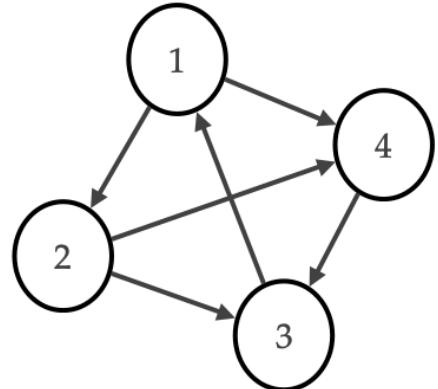
nodes:  $1, \dots, n$

$\mathbf{N} : n \times n$

$$N_{ij} = \begin{cases} \frac{1}{k_i^{\text{out}}}, & \text{if } (i, j) \\ 1, & \text{if } (i == j) \text{ and } k_i^{\text{out}} = 0 \\ 0, & \text{otherwise} \end{cases}$$

$N_{ij}$ : the share of  $i$ 's  $PR$  that  $j$  should get in one update step

If  $i$  has no outgoing links, then we define  $N_{ii} = 1$



$$\begin{bmatrix} 0 & \frac{1}{2} & 0 & \frac{1}{2} \\ 0 & 0 & \frac{1}{2} & \frac{1}{2} \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

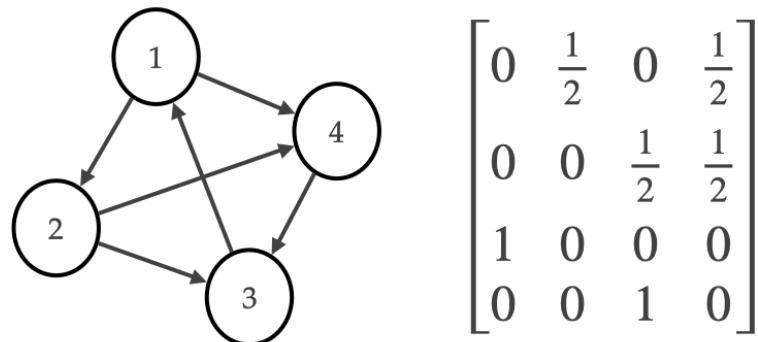
## Update rule (basic and scaled)

1. Basic update rule:

$$\forall i : r_i = \sum_{j=1}^n N_{ji} r_j \leftarrow N_{1i}r_1 + N_{2i}r_2 + \dots + N_{1n}r_n \mathbf{r} \leftarrow \mathbf{N}^T \cdot \mathbf{r}$$

2. Scaled update rule (factor  $s$ ):

$$\tilde{N}_{ij} = s \cdot N_{ij} + (1 - s) \cdot \frac{1}{n}$$



### 3. Application of scaled update rule:

$$\forall i : r_i = \sum_{j=1}^n \widetilde{N}_{ji} r_j \leftarrow \widetilde{N}_{1i} r_1 + \widetilde{N}_{2i} r_2 + \dots + \widetilde{N}_{1n} r_n \mathbf{r} \leftarrow \widetilde{\mathbf{N}}^T \cdot \mathbf{r}$$

## Repeated improvement (optional)

$\mathbf{r}^{<0>} = \left( \frac{1}{n}, \dots, \frac{1}{n} \right)$ , initial PR vector

$$\mathbf{r}^{<k>} = (\widetilde{\mathbf{N}}^T)^k \cdot \mathbf{r}^{<0>}$$

Limiting vector  $r^{<*>}$  satisfies  $\widetilde{\mathbf{N}}^T \cdot \mathbf{r}^{<*>} = 1 \cdot \mathbf{r}^{<*>}$

$\mathbf{r}^{<*>}$  should be an eigenvector of  $\widetilde{\mathbf{N}}^T$  with corresponding eigenvalue of 1

**Problem:**  $\widetilde{\mathbf{N}}^T$  is not symmetric

- this means that eigenvalues can be complex numbers and eigenvectors have no relationships to one another

## Convergence of the scaled PR update rule (optional)

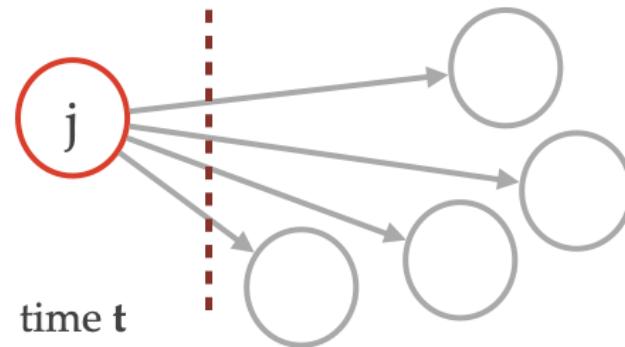
$$\forall i, j : \tilde{N}_{ij} > 0$$

Perron's theorem

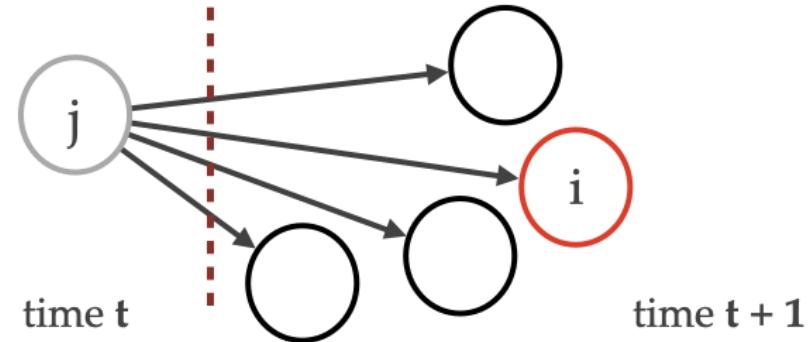
Matrix  $\mathbf{P}$  (with entries  $> 0$ )

- i)  $\mathbf{P}$  has an eigenvalue  $c > 0$  s.t.  $c > c'$   $\forall c'$  (with  $c'$  another eigenvalue)
- ii) Exists an eigenvector  $\mathbf{y}$  with real positive values corresponding to  $c$ , and  $\mathbf{y}$  is unique (up to a multiplication constant)
- iii) if  $c = 1$ , then for any starting vector  $\mathbf{x} \neq \mathbf{0}$  with non negative coordinates, the sequence of vectors  $p^k \mathbf{x}$  converges to a vector in the direction of  $\mathbf{y}$  ( $k \rightarrow \infty$ )

## Formulation of the PageRank using Random Walks



## Formulation of the PageRank using Random Walks



Which is the probability of being at node  $i$  at time  $t + 1$ ?

$b_1, b_2, \dots, b_n$ : the probabilities of being at node  $i$  in a given step.

$b_i \leftarrow \sum_{j=1}^n M_{ji} \frac{b_j}{k_j^{out}}$ : the probability of being at node  $j$  in the following step

Using matrix  $\mathbf{N}$  :  $b_i \leftarrow N_{1i}b_1 + N_{2i}b_2 + \dots + N_{1n}b_n \Rightarrow \mathbf{b} \leftarrow \mathbf{N}^T \cdot \mathbf{b}$

**claim:**  $PR$  of page  $i$  is exactly the probability of being at node  $i$  after  $k$  step.

## A scaled version of the random walk

For a given probability  $s$ : the walker follows a random outgoing edge

With prob  $(1 - s)$ : the walker is teleported uniformly at random to another node

$$b_i \leftarrow s \cdot \sum_{j=1}^n M_{ji} \frac{b_j}{k_j^{out}} + \frac{(1 - s)}{n}$$

Using matrix:

$$\widetilde{N} : b_i \leftarrow \widetilde{N}_{1i}b_1 + \widetilde{N}_{2i}b_2 + \dots + \widetilde{N}_{1n}b_n \Rightarrow \mathbf{b} \leftarrow \widetilde{\mathbf{N}}^T \cdot \mathbf{b}$$

**claim:**  $PR$  is equivalent to the scaled version of random walks.



## Reading material

[ns2] Chapter 14 (14.1-14.5) Link Analysis and Web Search

[ns1] Chapter 4 (4.3) [Simplified description of PageRank]



# Q & A

