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Analisi e Visualizzazione delle Reti Complesse

NS11 - Communities

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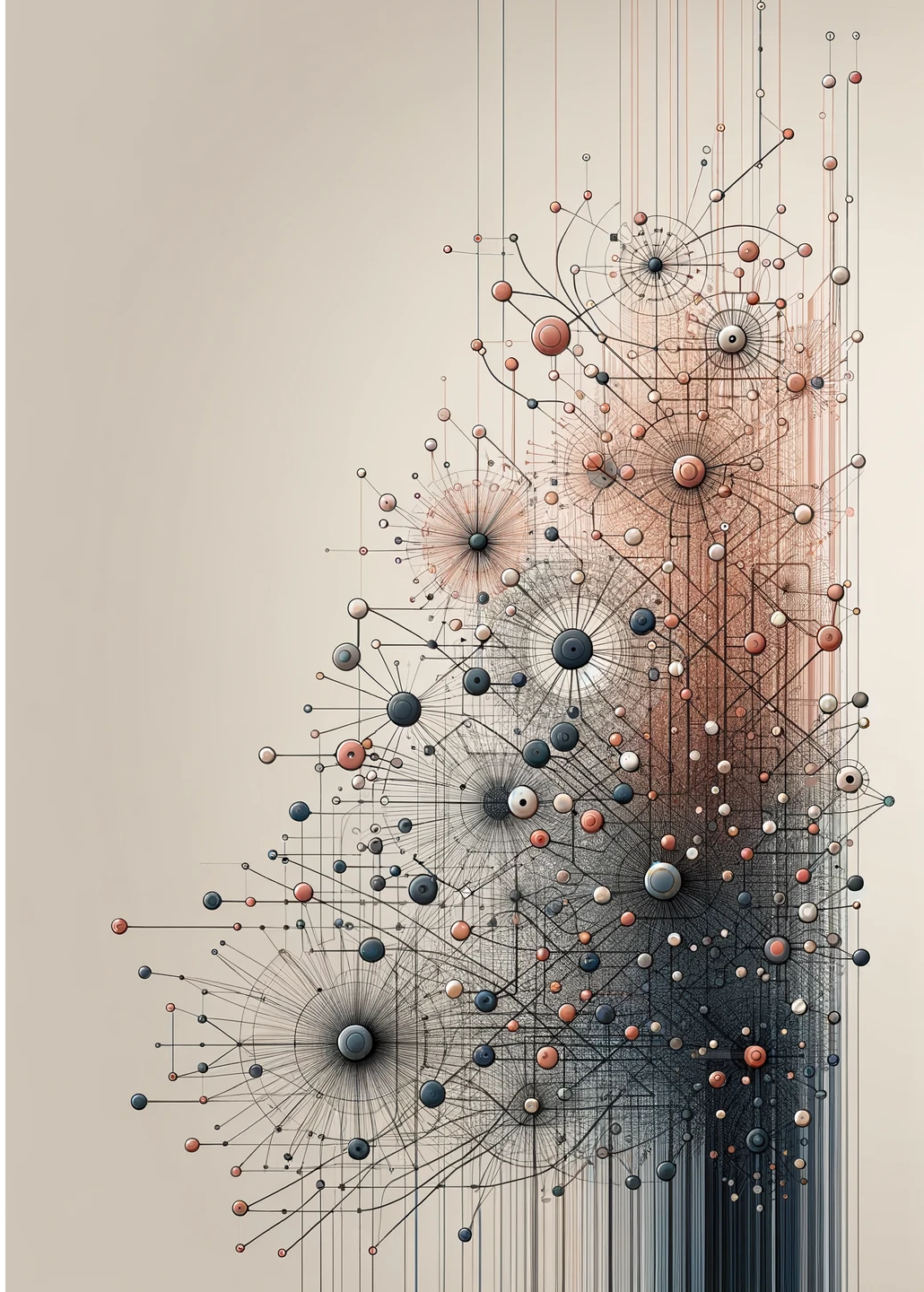

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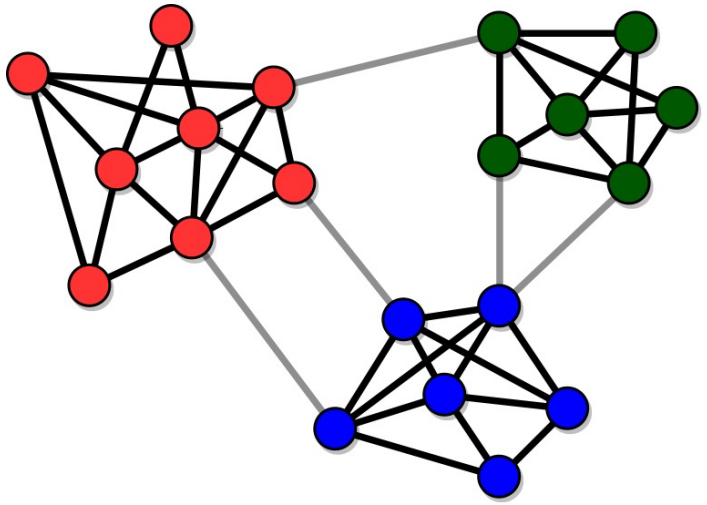
Outline

- Basic definitions
  - **community variables**
  - **community definitions**
  - **partitions**
- Related problems
  - **network partitioning**
  - **data clustering and dendrograms**

Basic definitions

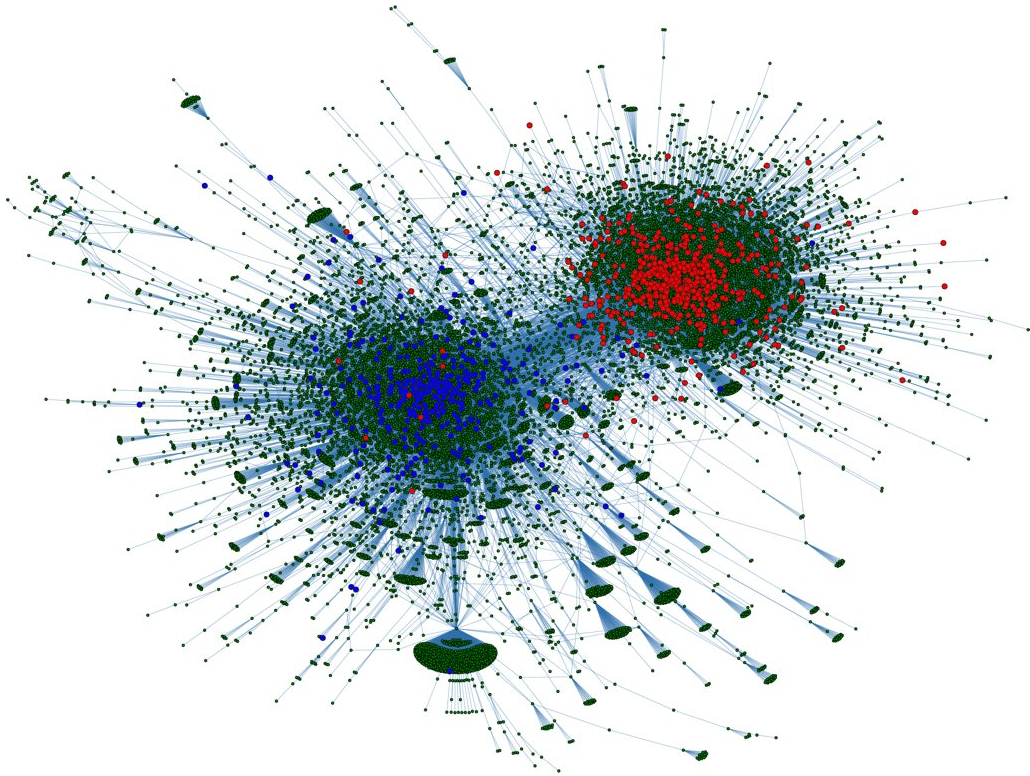
Community structure

Communities (or clusters): sets of tightly connected nodes



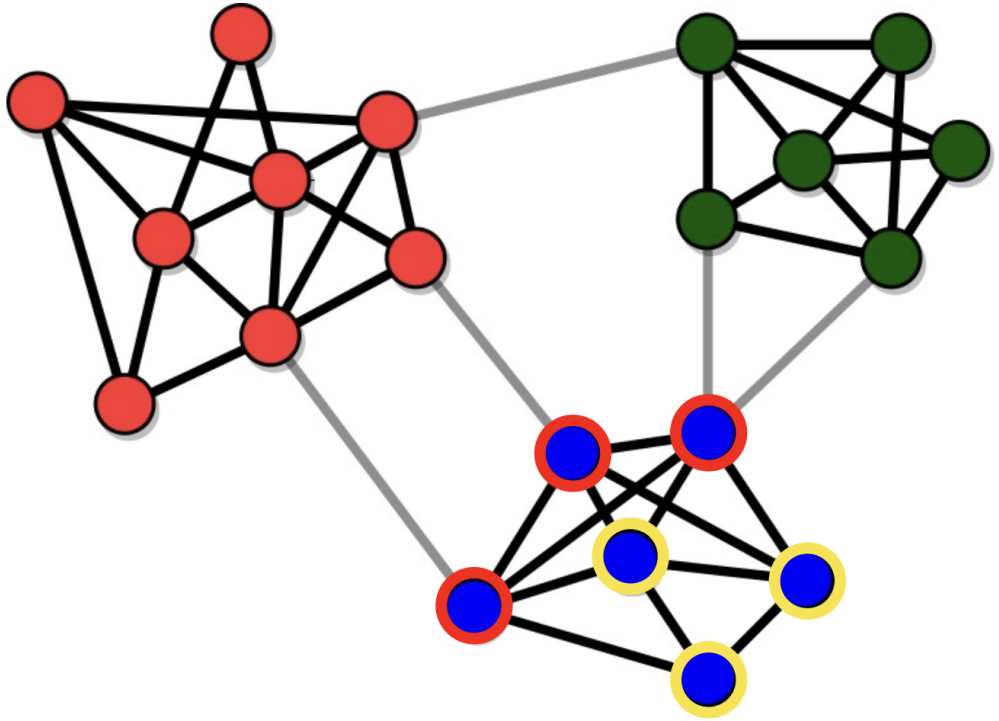
Community structure

- Example: Twitter users with strong political preferences tend to follow those aligned with them and not to follow users with different political orientation
- Other examples: social circles in social networks, functional modules in protein interaction networks, groups of pages about the same topic on the Web, etc.



Why study communities?

- Uncover the organization of the network
- Identify features of the nodes
- Classify the nodes based on their position in the clusters
- Find missing links

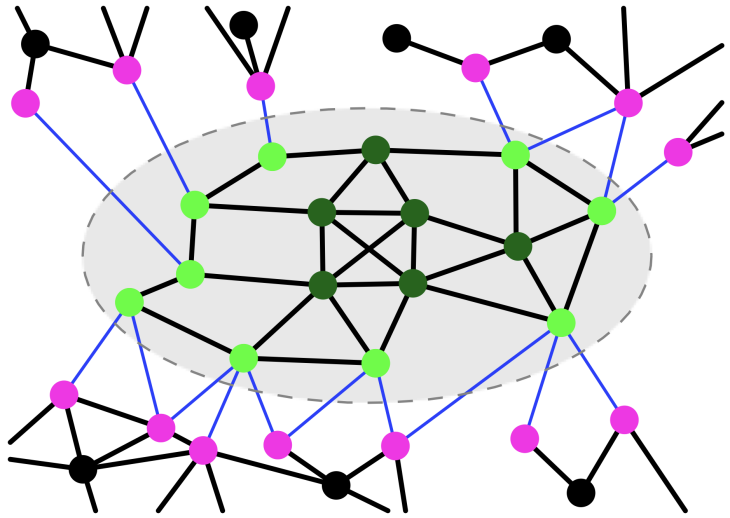


Basic definitions: variables

- **Internal degree of a node:** number of neighbors of the node in its community
- **External degree of a node:** number of neighbors of the node outside of its community
- **Community degree:** sum of the degrees of the nodes in the community
- **Internal link density:** ratio between the number of links L_C inside a community C and the maximal possible number of links that can lie inside C:

$$\delta_C^{int} = \frac{L_C}{L_c^{max}} = \frac{L_C}{\binom{N_C}{2}} = \frac{2L_C}{N_C(N_C-1)}$$

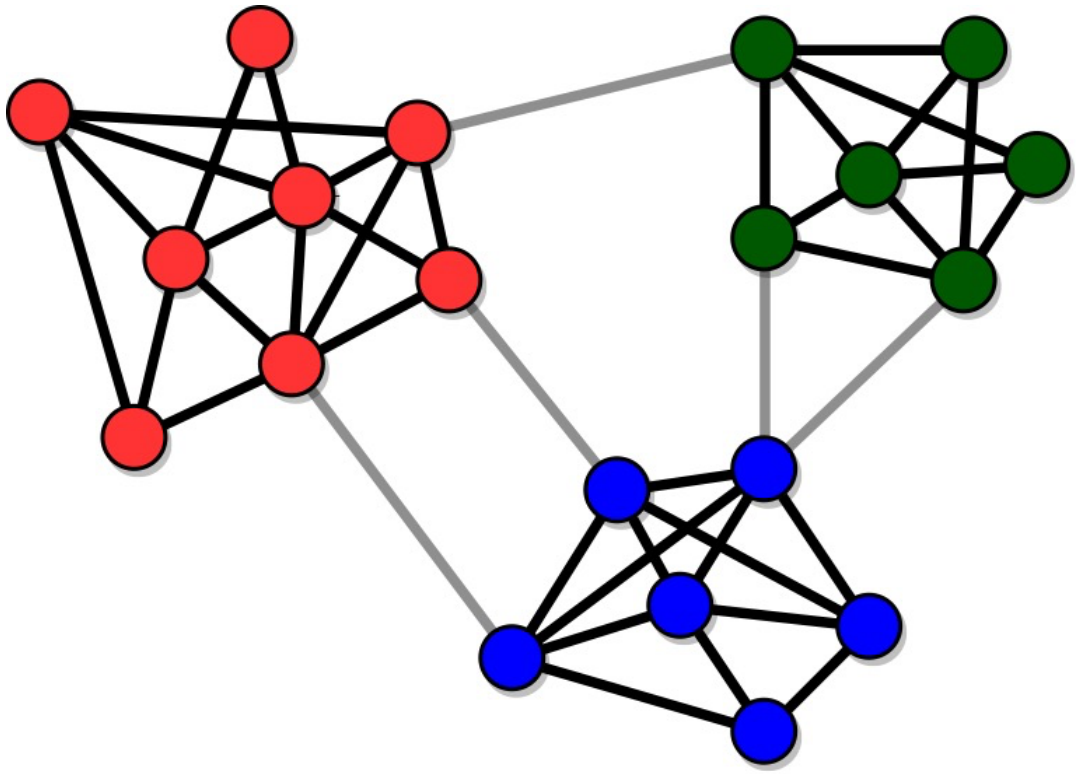
- where N_C is the number of nodes in C



Basic definitions: community

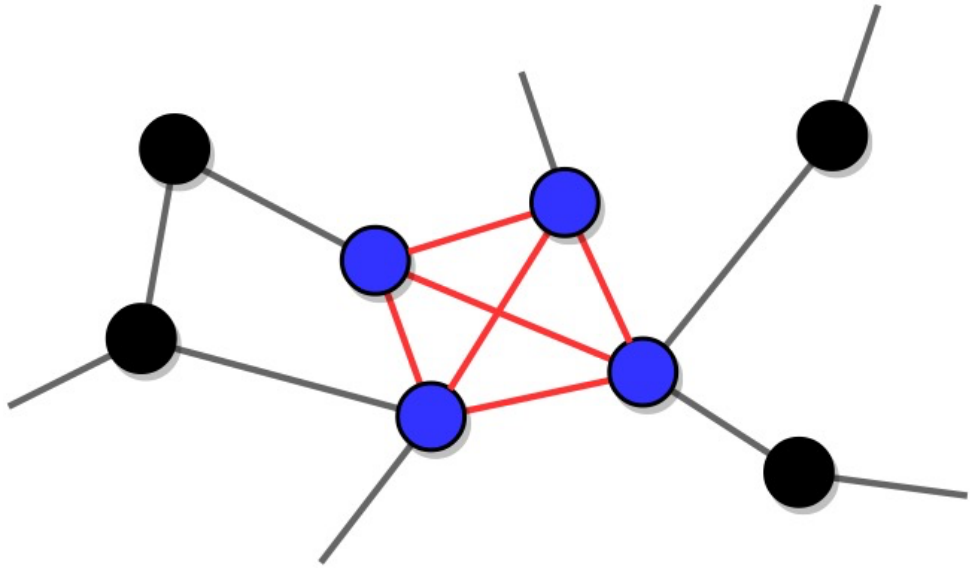
Two main features:

- **High cohesion:** communities have many internal links, so their nodes stick together
- **High separation:** communities are connected by a few links



Community definitions based on cohesion

- **Principle:** focus on the cluster's properties, disregarding the rest of the network
- **Example:** clique — all internal links are there (maximal cohesion)
- **Problem:** nodes are connected to all others in the cluster, whereas in real communities, they have different roles, which is reflected in heterogenous linking patterns



Community definitions based on cohesion versus separation

Principle:

Definition tends to achieve high cohesion and high separation

Popular idea:

The number of internal links exceeds the number of external links

Two concepts:

- **Strong community:** subnetwork such that the internal degree of each node is greater than its external degree
- **Weak community:** subnetwork such that the sum of the internal degrees of its nodes is greater than the sum of their external degrees

Community definitions based on cohesion versus separation

A strong community is also a weak community:

If the inequality between internal and external degree holds for each node, then it must hold for the sum over all nodes

A weak community is not a strong community, in general:

If the inequality between internal and external degree holds for the sum, it may be violated for one or more nodes

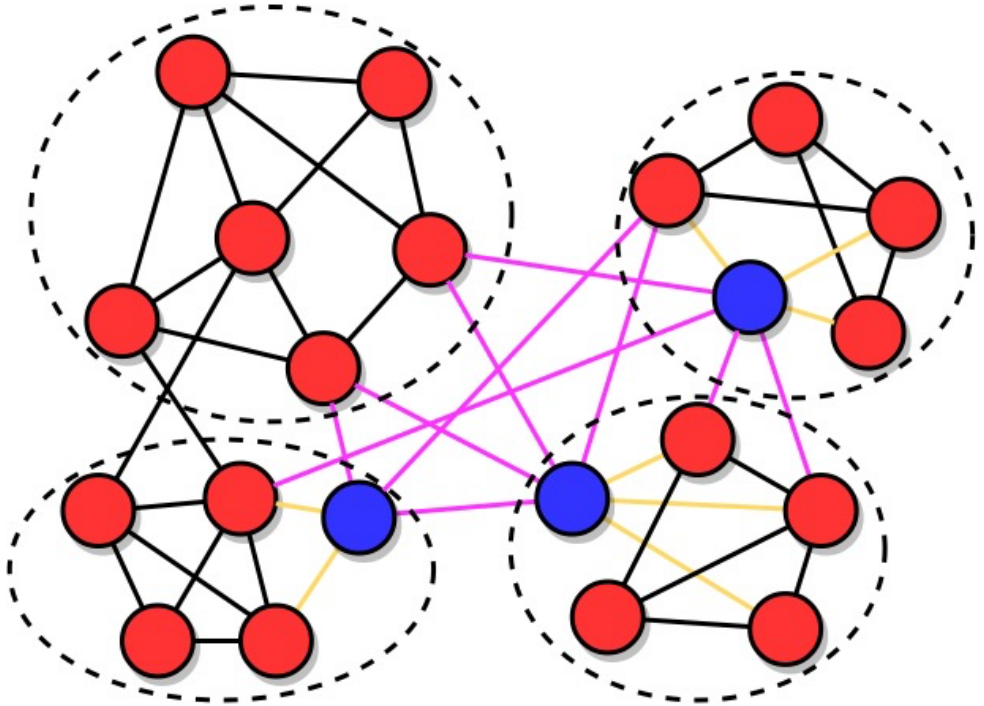
Problem:

In the definition of strong and weak community, one compares a subnetwork with the rest of the network. It makes more sense to compare subnetworks to each other!

Community definitions based on cohesion versus separation

Less stringent definitions of strong and weak community:

- **Strong community:** subnetwork such that each node has more neighbors inside it than in any other community
- **Weak community:** subnetwork such that the sum of the internal degrees of its nodes exceeds the total number of neighbors that the nodes have in any other community



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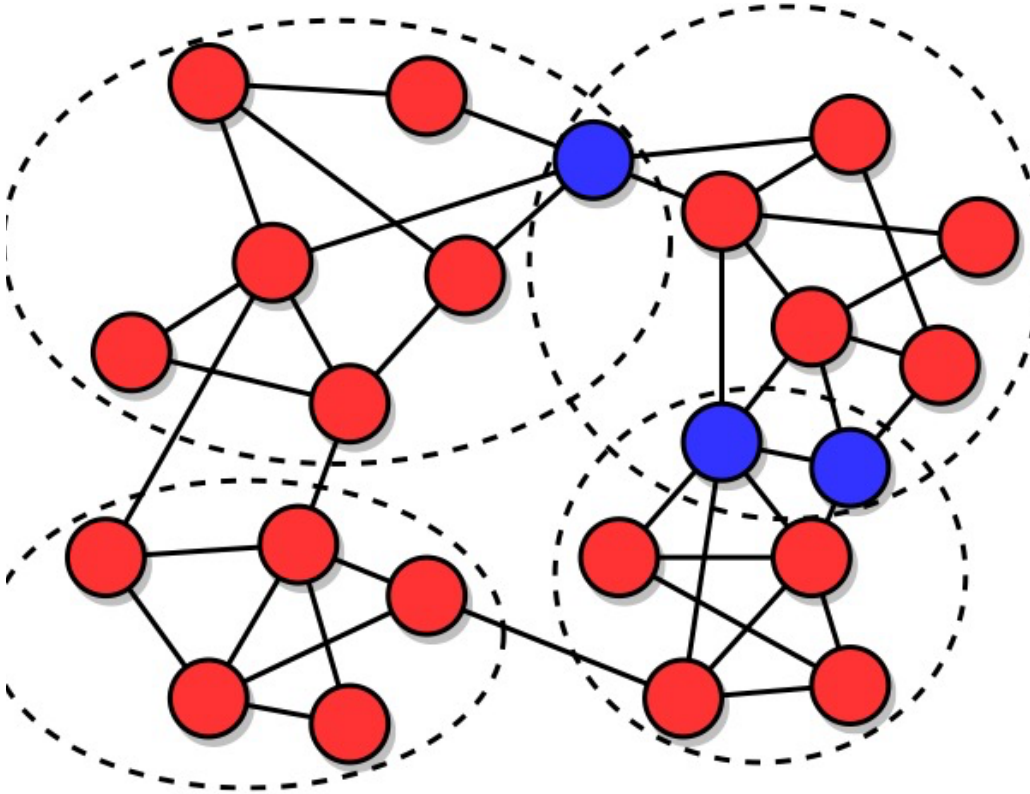
Partitions

- The number of partitions of n objects is the **Bell number** B_n
- The Bell number grows **faster than exponentially with** n
- **Conclusion:** it makes no sense to look for interesting community structures by exploring the whole space of partitions! A smart exploration of the partition space must be performed.

n	B_n
1	1
2	2
3	5
4	15
5	52
6	203
7	877
8	4140
9	21147
10	115975
11	678570
12	4213597
13	27644437
14	190899322
15	1382958545

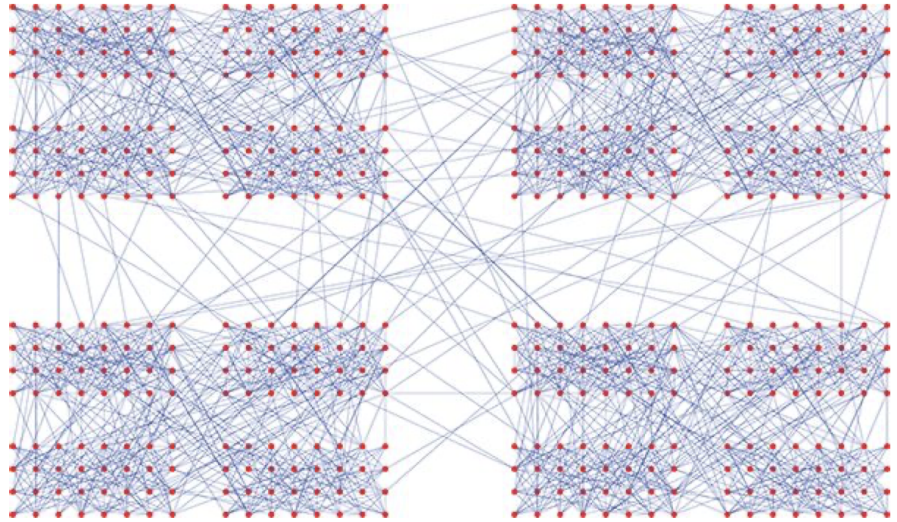
Overlapping communities

- Communities in many real networks **overlap**
- A division of a network into overlapping communities is called **cover**
- The number of possible covers of a network is **far higher** than the number of partitions due to the many ways clusters can overlap



Hierarchical communities

- If the network has multiple levels of organization, its communities could form a **hierarchy**, with small communities within larger ones
- **Example:** branches in a company, in turn divided into departments
- **All hierarchical partitions are meaningful:** a good clustering algorithm should detect all of them



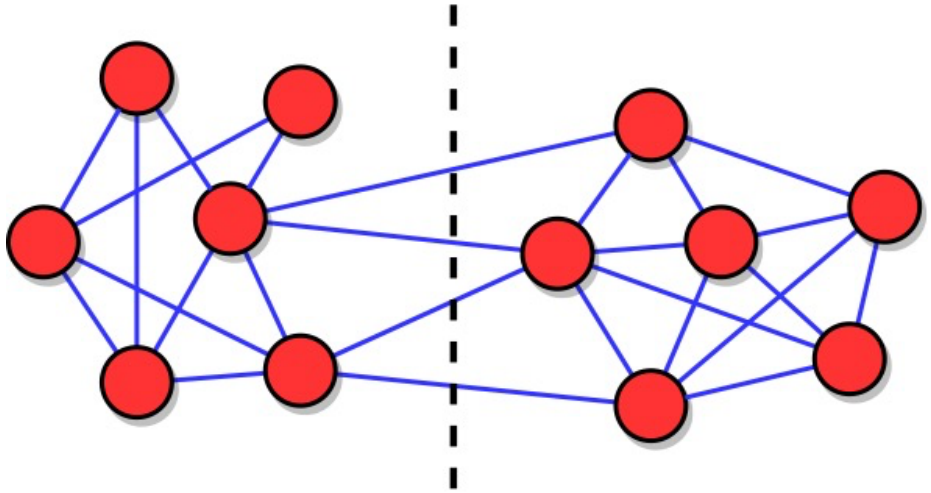
Related problems

Clusters

- Networks are made of **tightly-knit regions** connected by means of sparser connections
- Now we have a more formal definition of such regions, along with some other useful measures:
  - **clustering coefficients / triadic closure**
  - **(local) bridges / neighborhood overlap / betweenness**
  - **cores / k-cores / k-shells**
- The problem of finding denser regions in a network is called: **graph partitioning** or **community detection** (with some differences)
- Another related problem is **data clustering**

Related problems: network partitioning

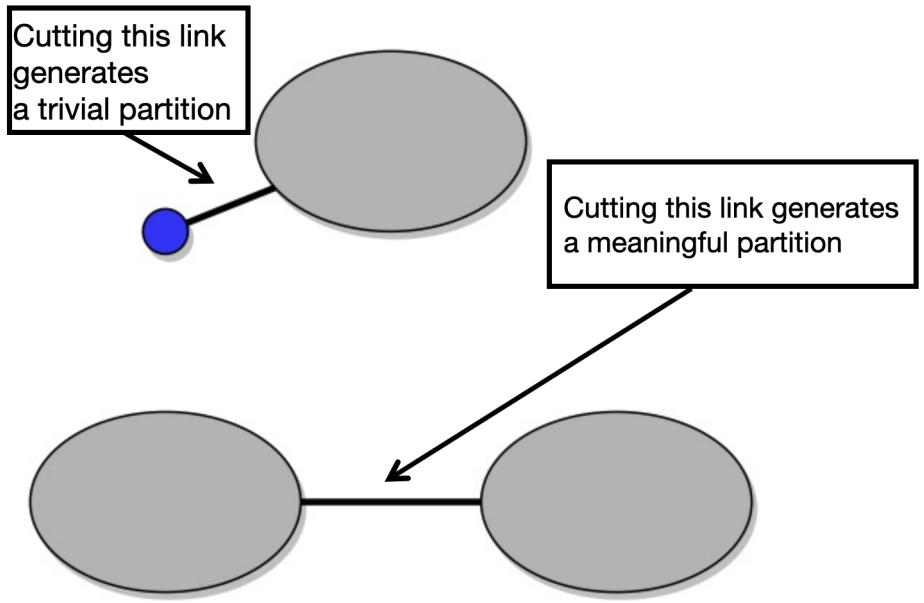
- **The problem:** dividing the nodes of a network into a number of groups of predefined size, such that the number of links between the groups is minimal
- The number of links between groups is called **cut size**



Related problems: network partitioning

Problems:

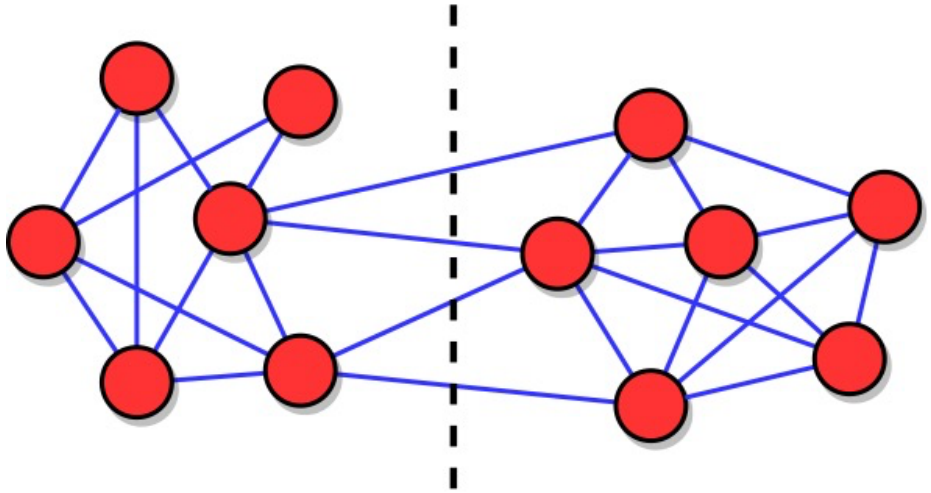
- If the number of clusters is not given beforehand, the trivial solution is a single cluster including everything
- If the size of the clusters is not indicated, there may be trivial solutions by removing the nodes with lowest degree



Network bisection

The problem:

Dividing the nodes of a network into **two groups of equal size**, such that the number of **links between the groups is minimal**

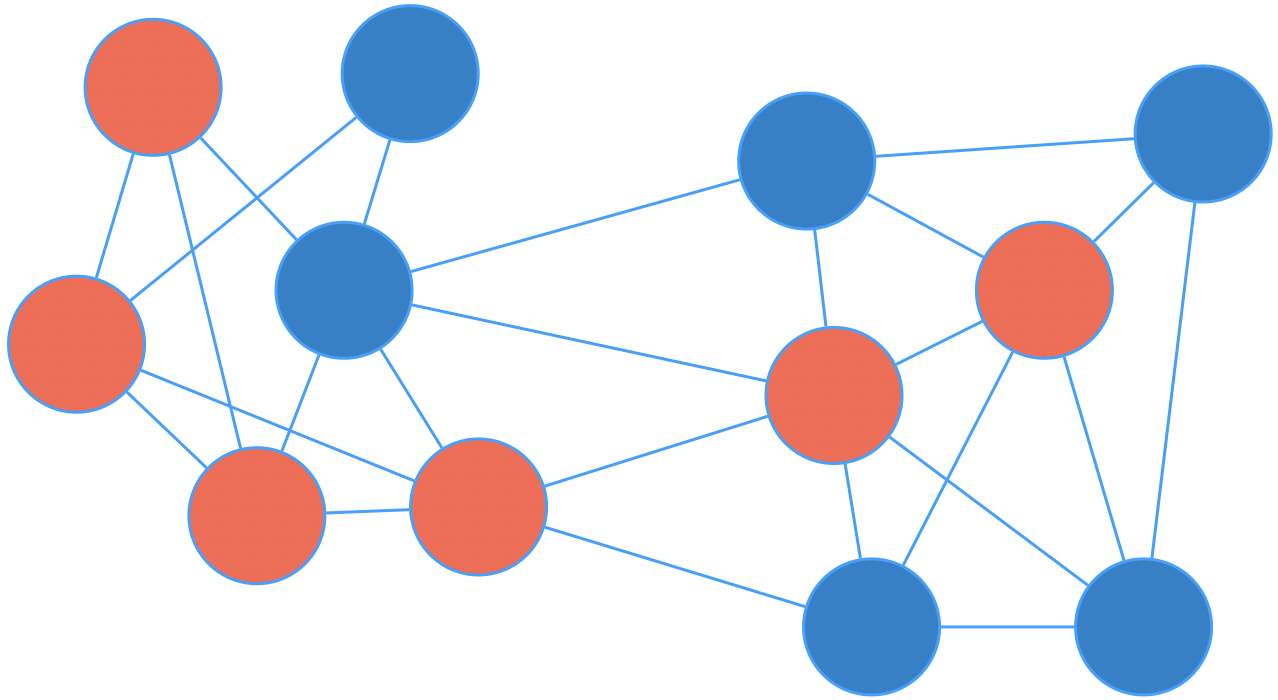


Kernighan-Lin algorithm

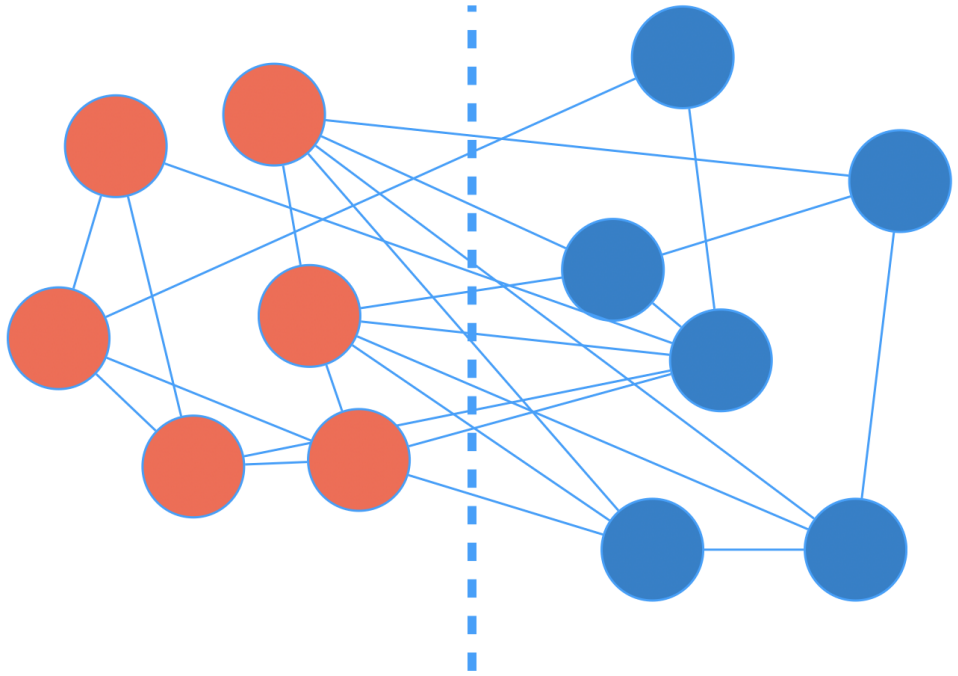
Procedure:

1. Split in two groups A and B of predefined sizes N_A and N_B, e.g., by randomly assigning nodes to either group (for bisection $N_A = N_B$)
2. For each pair of nodes i, j, with $i \in A$ and $j \in B$, compute the variation in cut size between the current partition and the one obtained by swapping i and j
3. The pair of nodes i and j yielding the largest decrease in cut size is selected and swapped. This pair of nodes is locked; they will not be touched again during this iteration
4. Repeat steps 2 and 3 until no more swaps of unlocked nodes yield a decrease in cut size. This yields a new bipartition, that is used as a starting configuration for the next iteration

Kernighan-Lin algorithm: step 1

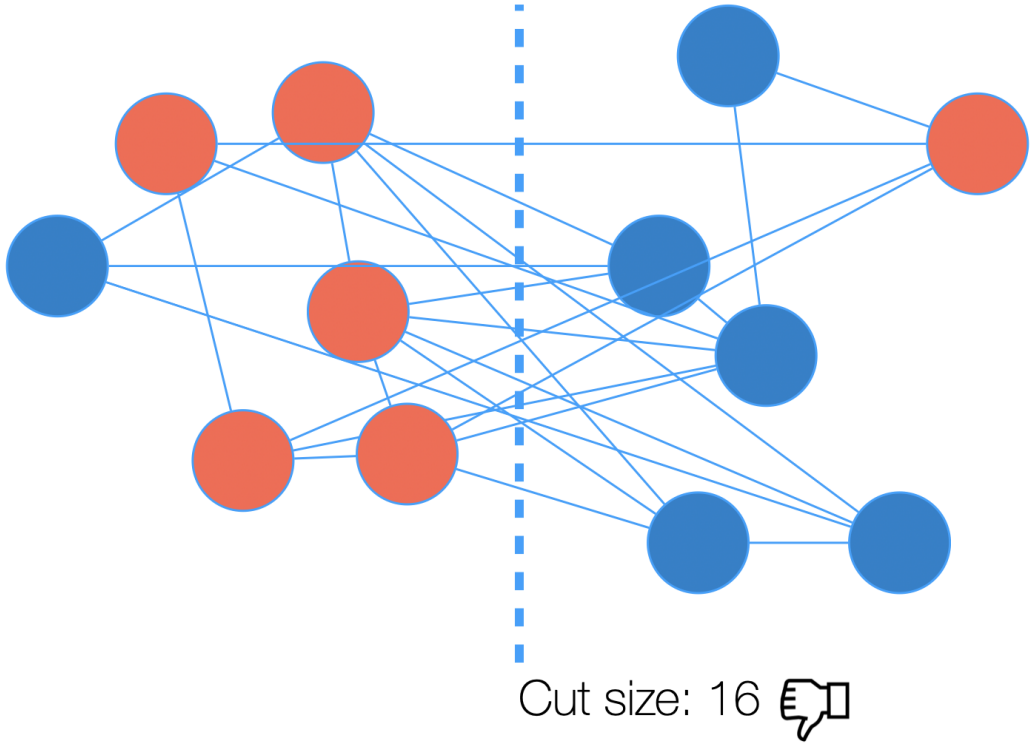


Kernighan-Lin algorithm: step 1

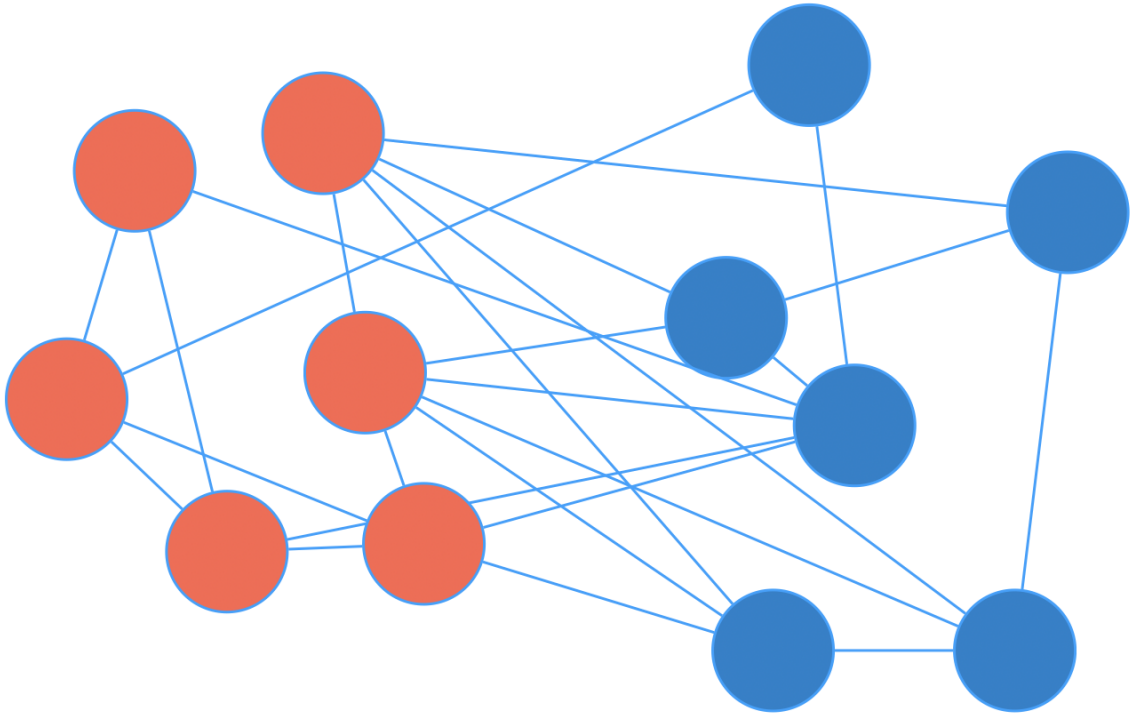


Cut size: 13

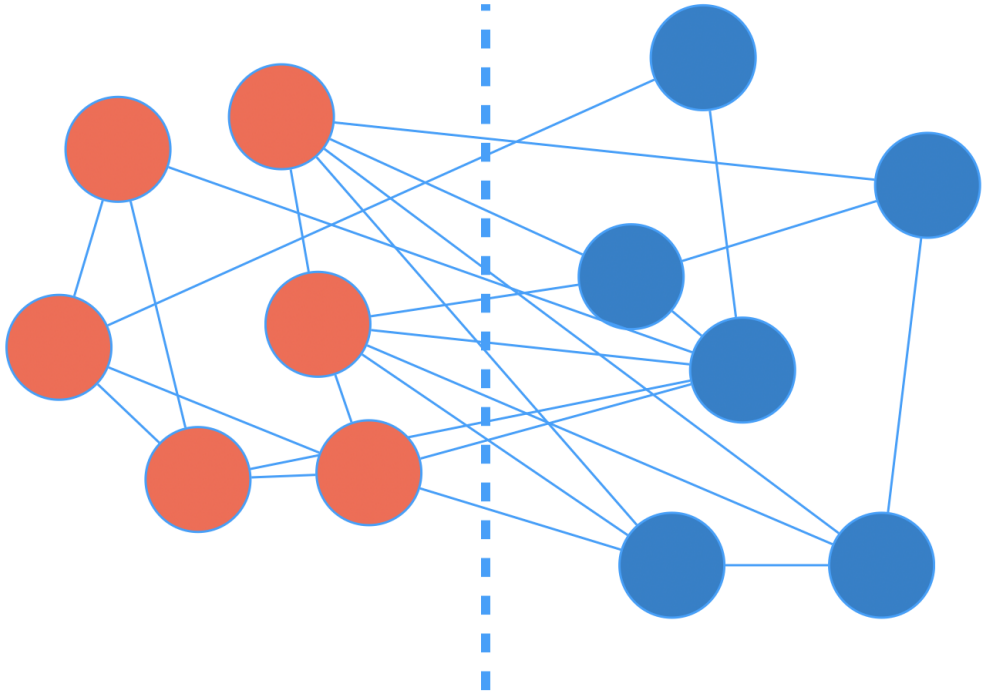
Kernighan-Lin algorithm: step 2



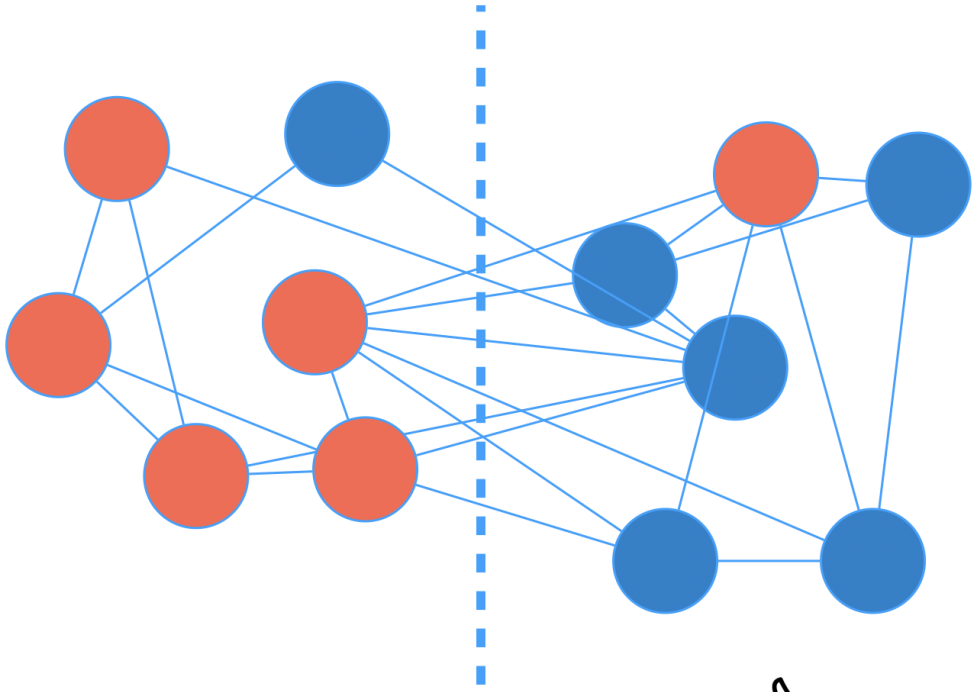
Kernighan-Lin algorithm: step 2



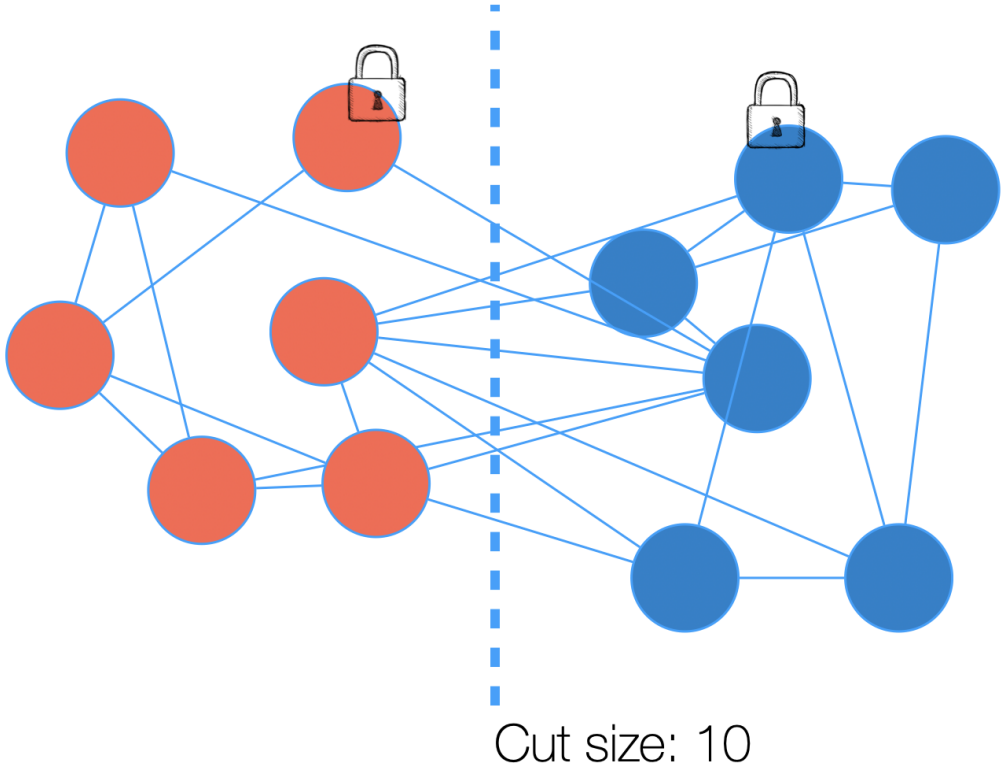
Kernighan-Lin algorithm: step 2



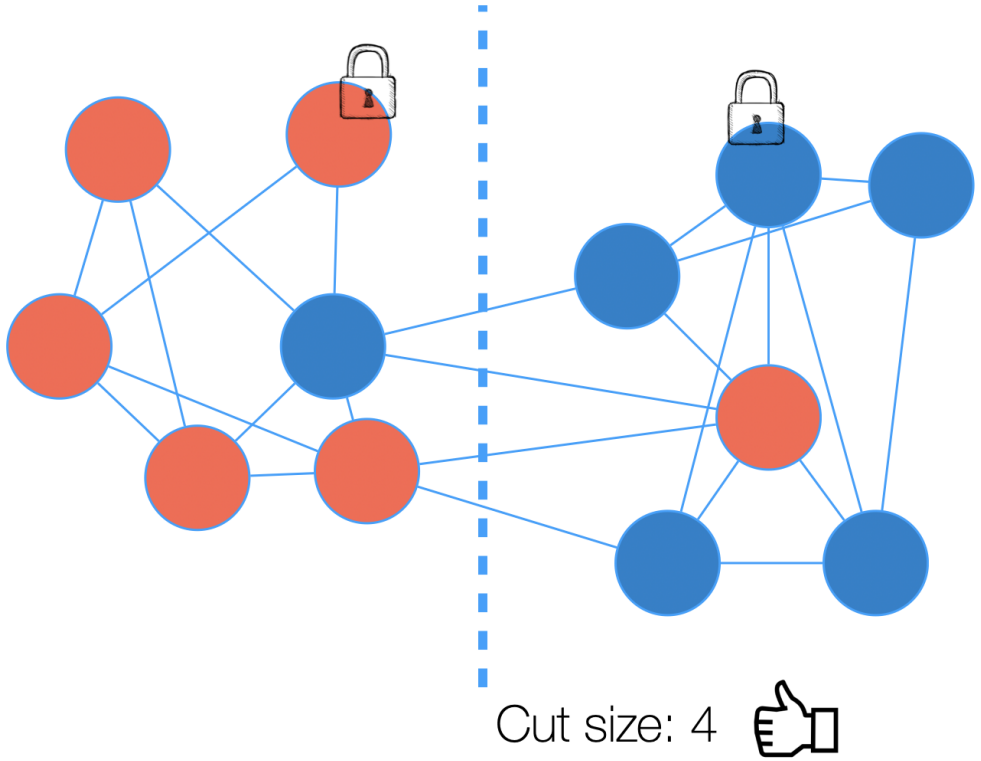
Kernighan-Lin algorithm: step 2



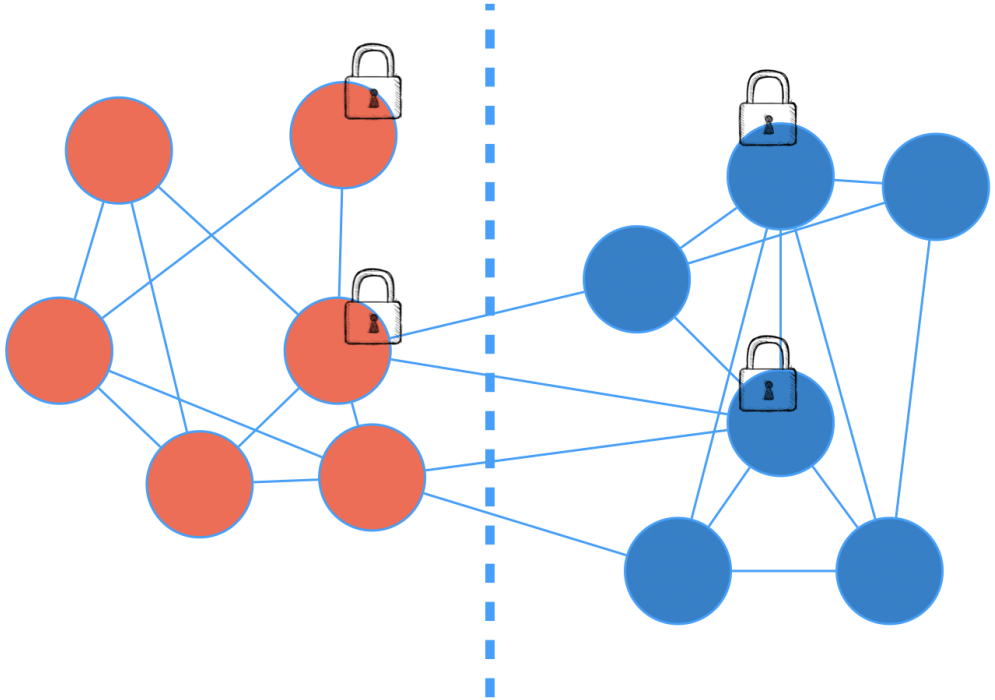
Kernighan-Lin algorithm: step 3



Kernighan-Lin algorithm: step 2



Kernighan-Lin algorithm: step 3



Cut size: 4

Kernighan-Lin algorithm

- **Convergence:** The procedure ends when the cut size of partitions obtained after consecutive iterations is the same, meaning that the algorithm is unable to improve the result
- The algorithm can be applied to minimize the cut size of partitions into more than two clusters, by swapping nodes between pairs of clusters
- **Problem:** the choice of the initial partition heavily affects the final result. The larger the cut size of the initial partition, the worse the final solution and the longer the time to reach convergence
- **Solution:** creating multiple random partitions and picking the one with the lowest cut size as initial partition

Kernighan-Lin algorithm

- The described procedure is **greedy**, in that at each step one looks for the partition with the smallest cut size. Because of that, the algorithm gets stuck in **local optima**, i.e., solutions whose cut size is not as low as it can be
- The problem can be mitigated by occasionally allowing swaps of nodes that increase the cut size
- The Kernighan-Lin algorithm is widely applied as a post-processing technique, to improve partitions delivered by other methods. Such partitions can be used as starting points for the method, which might return solutions with lower cut size

Network partitioning: limits

- Clusters have to be well-separated, but they do not need to have high internal link density
- Clusters found via graph partitioning are not communities, in general
- The number of clusters must be given as input, but it is usually unknown

In NetworkX:

```
# minimum cut bisection: returns a pair of sets of nodes

partition = nx.community.kernighan_lin_bisection(G)
```

Data clustering

- **Grouping objects based on how similar to each other they are**
- Two main classes of algorithms:
  - **Hierarchical clustering**
  - **Partitional clustering**

Hierarchical clustering

- Hierarchical clustering delivers a nested set of partitions
- Main ingredient: **similarity measure**
- Examples:
  - In a social network it could indicate how close the profiles of two people are based on their interests
  - If nodes are embedded in space (i.e., they are points in a metric space), the (dis)similarity between two nodes can be expressed by their distance
  - If nodes are not embedded in space, similarity measures can be derived from the network structure

Similarity: structural equivalence

- **Concept:** nodes are similar if their neighbors are similar

$$S_{ij}^{SE} = \frac{\text{number of neighbors shared by i and j}}{\text{total number of nodes neighboring only i, only j or both}}$$

- **Examples:**
  - If the neighbors of i and j are (v_1, v_2, v_3) and (v_1, v_2, v_4, v_5), respectively, $S_{ij} = 2/5 = 0.4$, because there are two common neighbors (v_1 and v_2) out of five distinct neighbors in total $(v_1, v_2, v_3, v_4, v_5)$
  - If i and j have no neighbors in common, $S_{ij} = 0$
  - If i and j have the same neighbors, $S_{ij} = 1$

Similarity of node groups

- **Question:** how can we define the similarity S_{G_1,G_2} between two groups of nodes G_1 and G_2 via the similarity S between pairs of nodes?
- **Answer:** multiple approaches. The first step is to compute the pairwise similarity S_{ij} between each node i in G_1 and each node j in G_2.

Then the following strategies can be adopted:

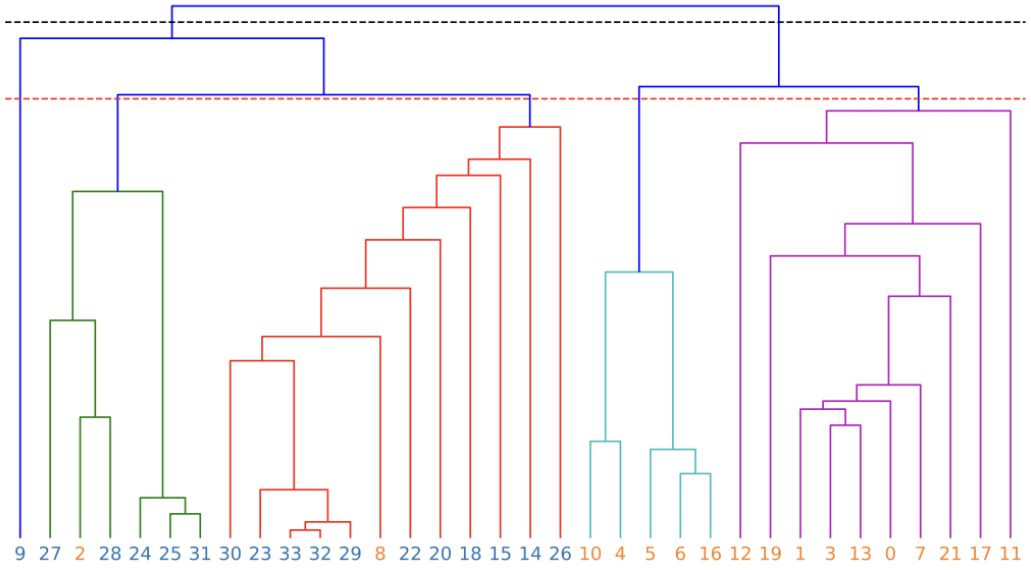
- **Single linkage:** take the maximum pairwise similarity $\rightarrow S_{G_1,G_2} = max_{i,j} S_{ij}$
- **Complete linkage:** take the minimum pairwise similarity $\rightarrow S_{G_1,G_2} = min_{i,j} S_{ij}$
- **Average linkage:** take the average pairwise similarity $\rightarrow S_{G_1,G_2} = \langle S_{ij} \rangle_{i,j}$

Hierarchical clustering

- **Two approaches**
  - **Agglomerative hierarchical clustering:** partitions are generated by iteratively merging groups of nodes
  - **Divisive hierarchical clustering:** partitions are generated by iteratively splitting groups of nodes
- **Agglomerative hierarchical clustering:**
  - **Start:** partition into N groups, each group consisting of one node
  - At each step the pair of groups with the largest similarity are merged

Dendrograms

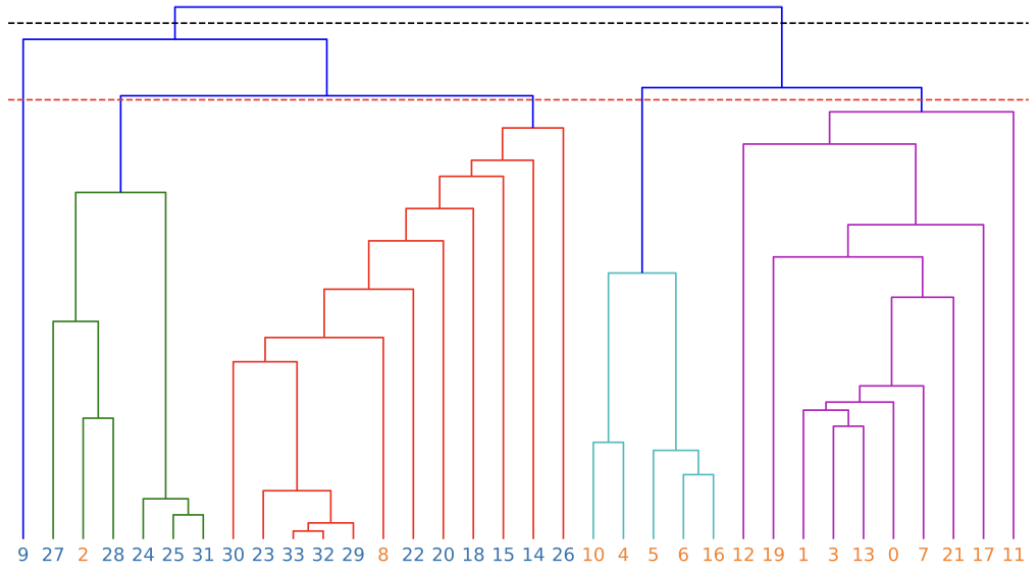
- Outcome: **dendrogram (hierarchical tree)**
- A dendrogram is a compact summary of all partitions created by hierarchical clustering
- Since each merger reduces the number of groups by one, **the total number of partitions is** N



Dendrograms

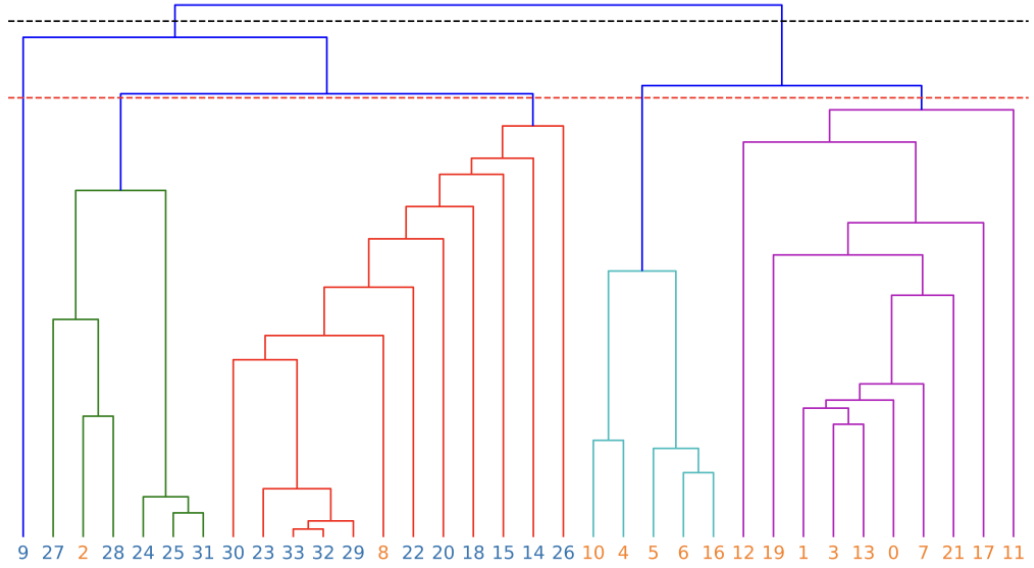
Features:

- **Bottom: leaves** of the tree, indicated by the labels of the nodes
- Going upwards, pairs of clusters are merged. Mergers are illustrated by horizontal lines joining two vertical lines, each representing a cluster
- The nodes of each cluster can be identified by following the vertical line representing the cluster all the way down



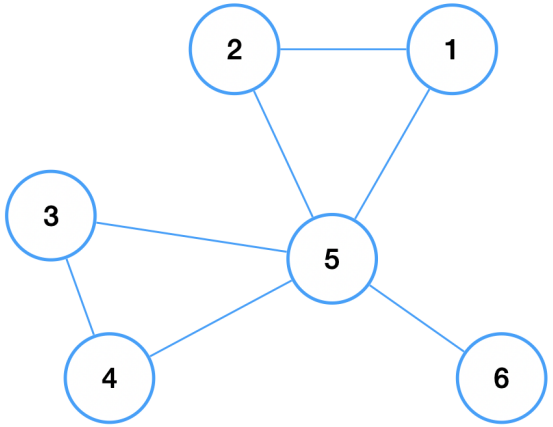
Dendrograms

- Partitions are selected via **horizontal cuts** of the dendrogram: the clusters are the ones corresponding to the vertical lines severed by the cut
- High cuts yield partitions into a few large clusters, low cuts yield partitions into many small clusters
- **Hierarchy:** each partition has clusters including clusters of all partitions lying lower in the dendrogram



Example

Jaccard Similarity $S^J_{i,j} = \frac{|N_i \cap N_j|}{|N_i \cup N_j|}$



1 2 3 4 5 6

Example

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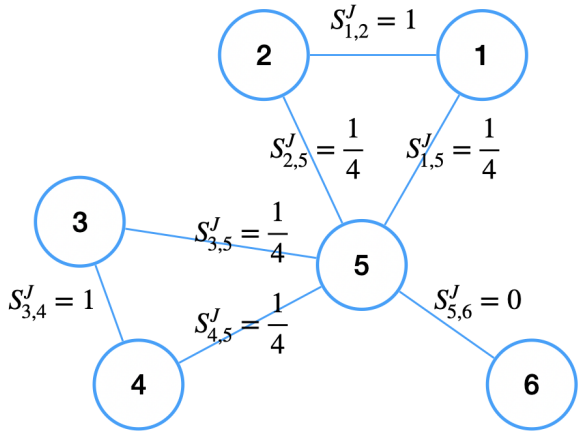

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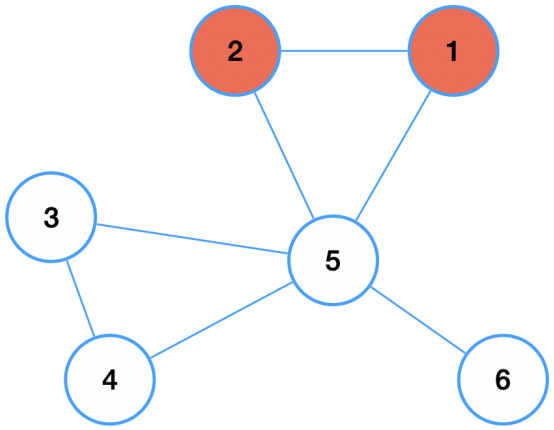

Example

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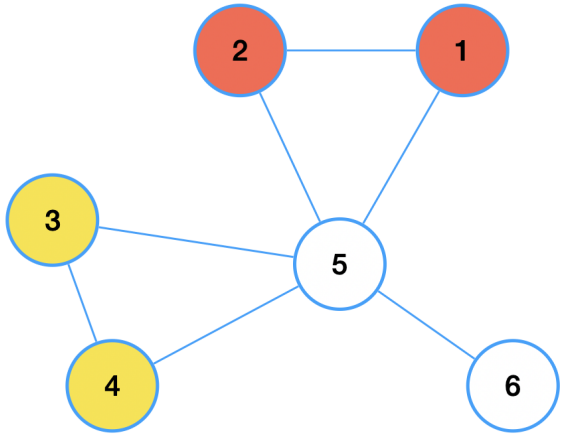
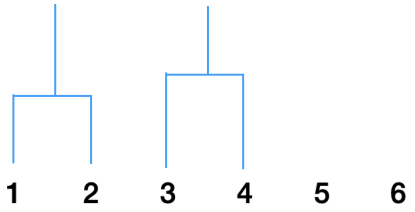
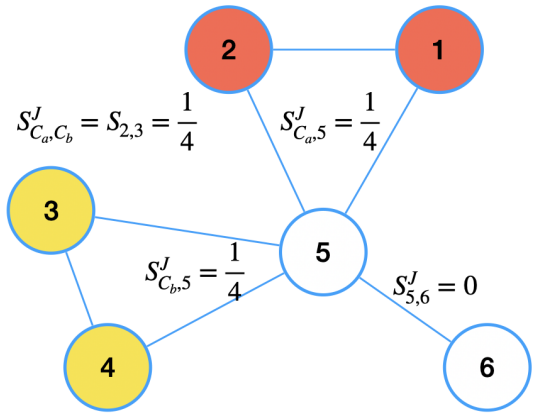
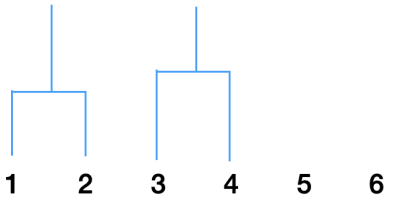




Example

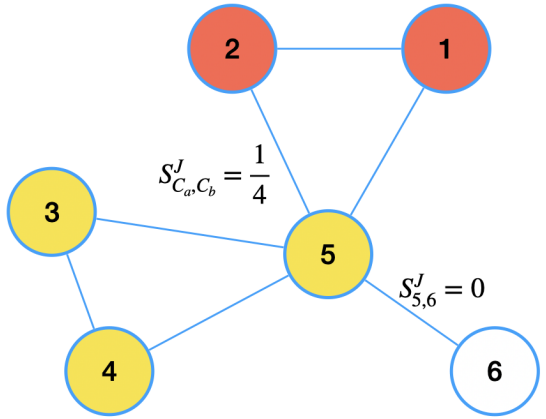
Jaccard Sim. $S^J_{i,j} = \frac{|N_i \cap N_j|}{|N_i \cup N_j|}$ Jaccard Sim. + single linkage $S^J_{C_a,C_b} = max S_{i,j} : i \in C_a, j \in C_b$

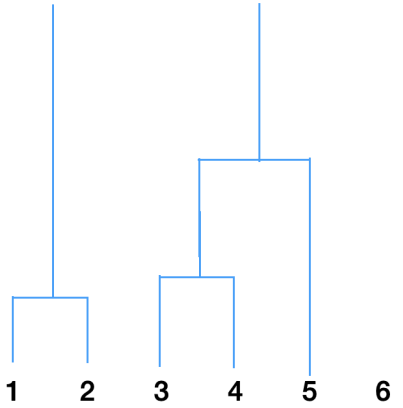




Example

Jaccard Sim. $S^J_{i,j} = \frac{|N_i \cap N_j|}{|N_i \cup N_j|}$ Jaccard Sim. + single linkage $S^J_{C_a,C_b} = max S_{i,j} : i \in C_a, j \in C_b$





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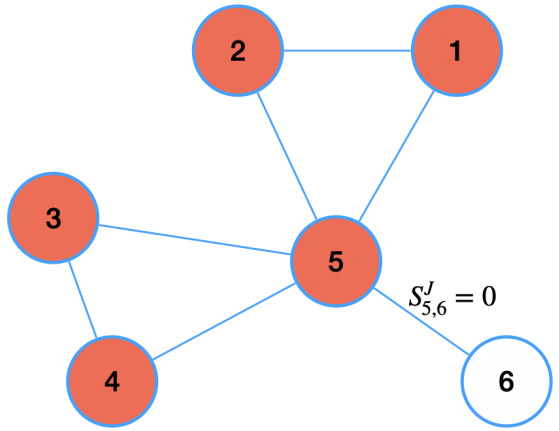

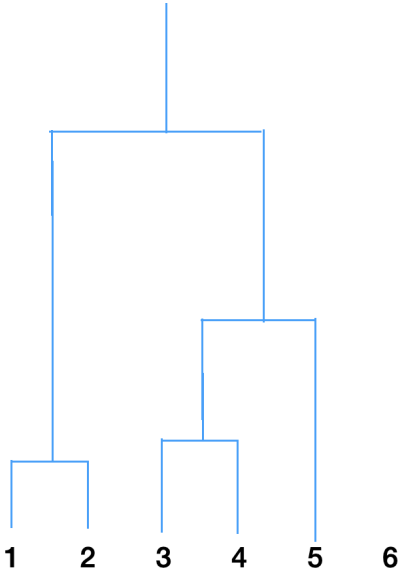

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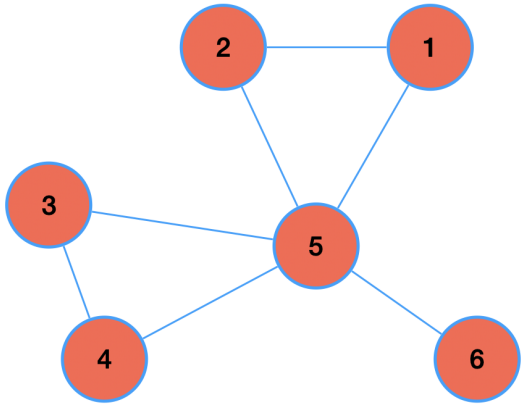

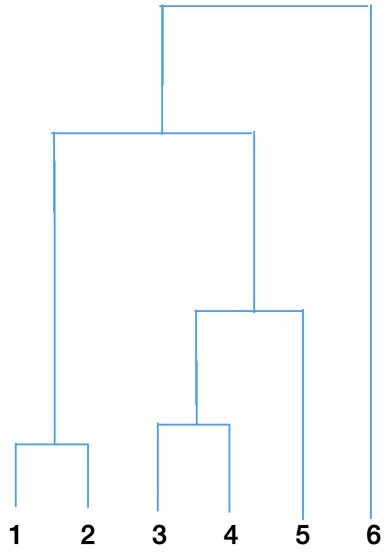

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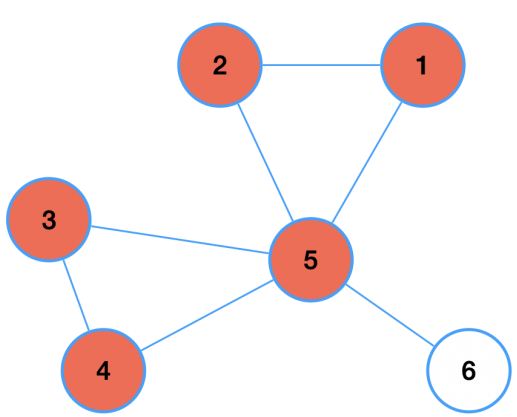

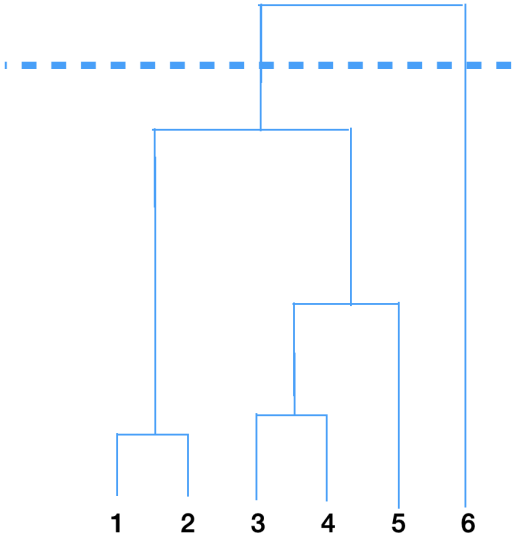

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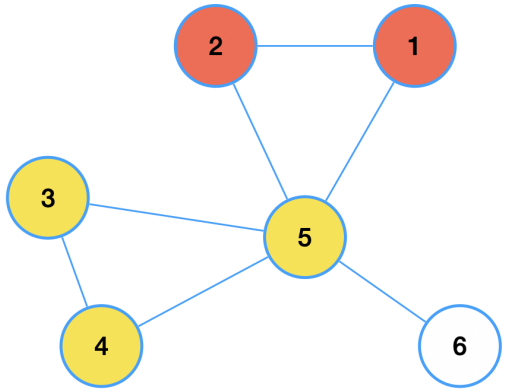

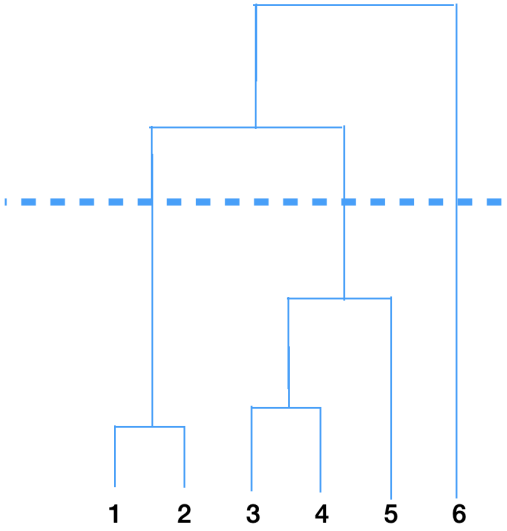

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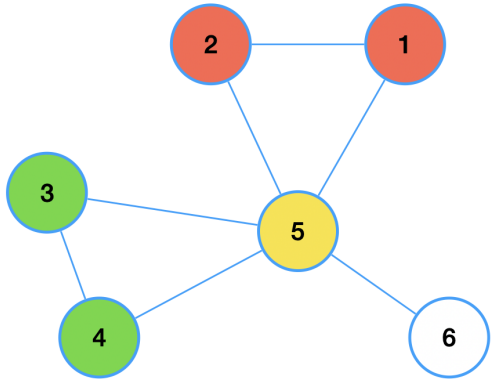

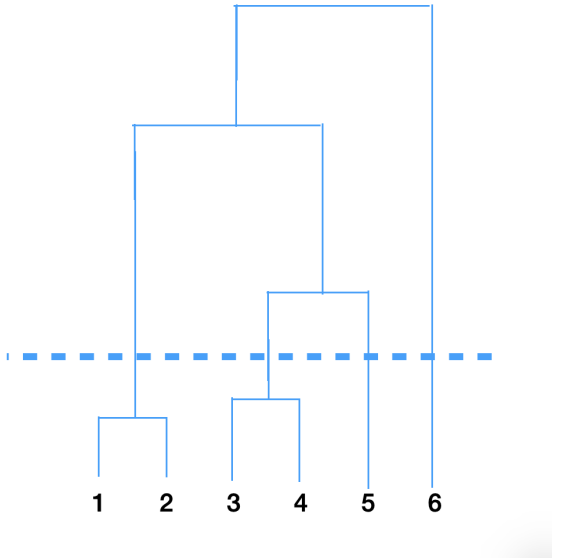

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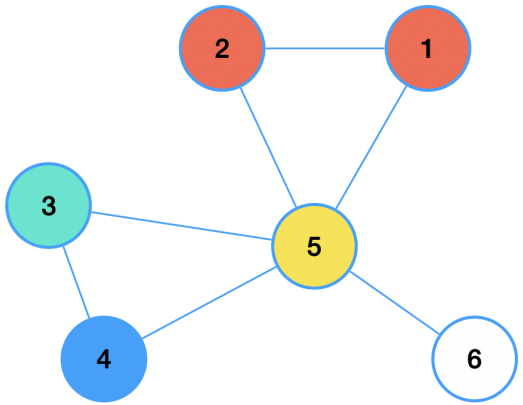

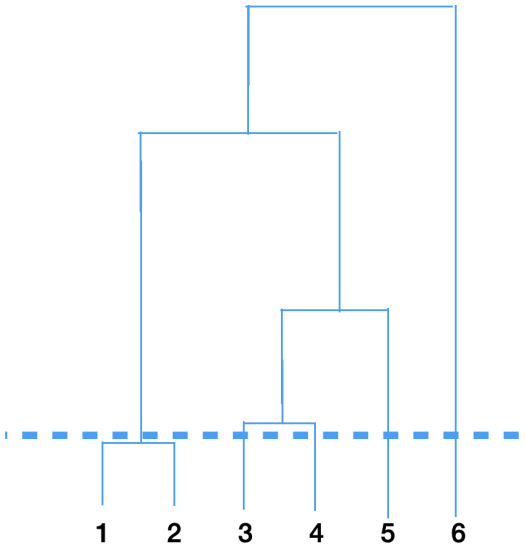

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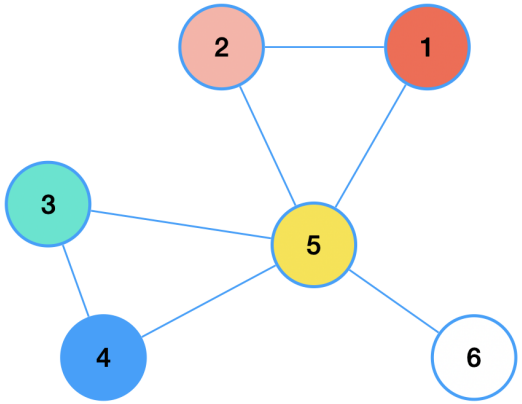

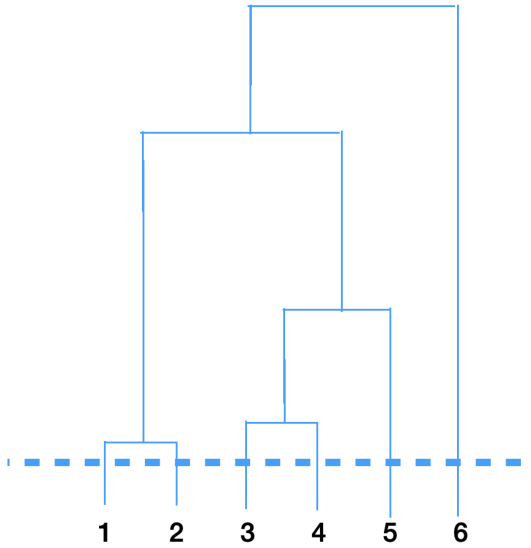

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Reading material

References

[ns1] **Chapter 6 - Communities**

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Q&A