

Procedures

+Mechanisms in Procedures



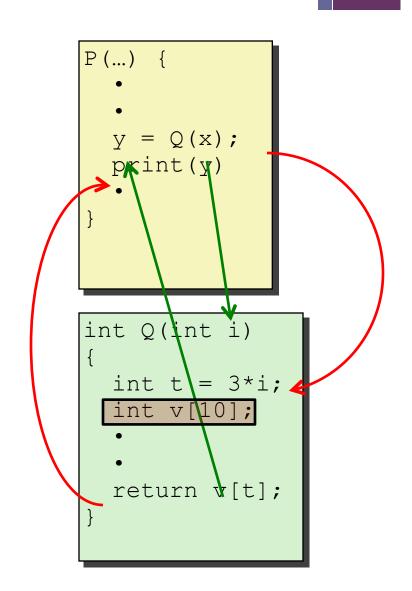
- To beginning of procedure code
- Back to return point

Passing data

- Procedure arguments
- Return value

Memory management

- Allocate during procedure execution
- Deallocate upon return
- Mechanisms all implemented with machine instructions
- x86-64 implementation of a procedure uses only those mechanisms required



+x86-64 Stack

- Stack "Bottom"
- Region of memory managed with stack discipline
- Grows toward lower addresses
- Register %rsp contains lowest address
 - address of "top" element

Increasing **Addresses** Stack Grows Down Stack "Top"

Stack Pointer: %rsp

+x86-64 Stack: Push

- pushq src
 - Fetch operand at *src*
 - Decrement %rsp by 8
 - Write operand at address given by %rsp

Stack "Bottom"

Increasing Addresses

Stack Grows Down

Stack Pointer: %rsp



+x86-64 Stack: Pop

- popq dest
 - Read value at address given by %rsp
 - Increment %rsp by 8 bytes
 - Store value at dest (must be register)

Stack "Bottom" Increasing **Addresses** Stack **Grows** Down Stack "Top"

Stack Pointer: %rsp

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Passing Control

+Code Examples

```
void multstore (long x, long y, long *dest) {
   long t = mult2(x, y);
   *dest = t;
}
```

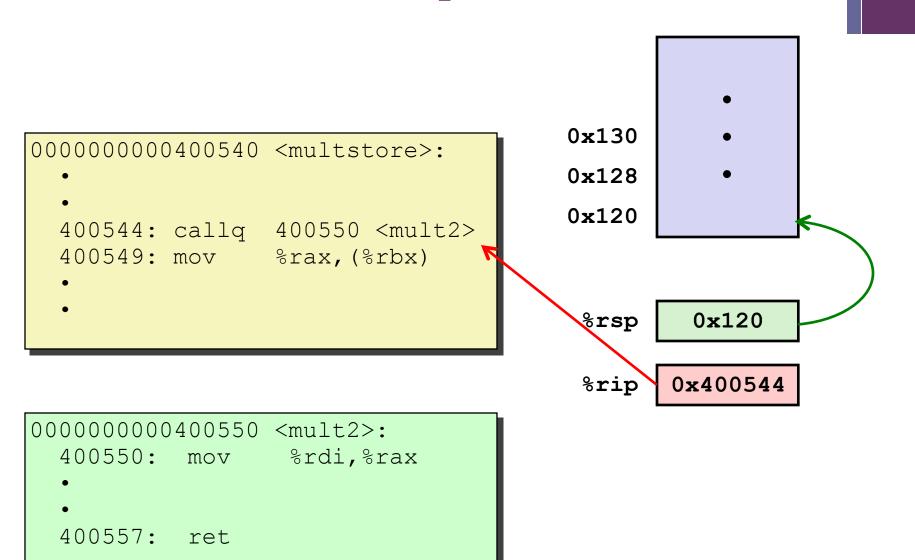
```
00000000000400540 <multstore>:
   400540: push %rbx # Save %rbx
   400541: mov %rdx,%rbx # Save dest
   400544: callq 400550 <mult2> # mult2(x,y)
   400549: mov %rax,(%rbx) # Save at dest
   40054c: pop %rbx # Restore %rbx
   40054d: retq # Return
```

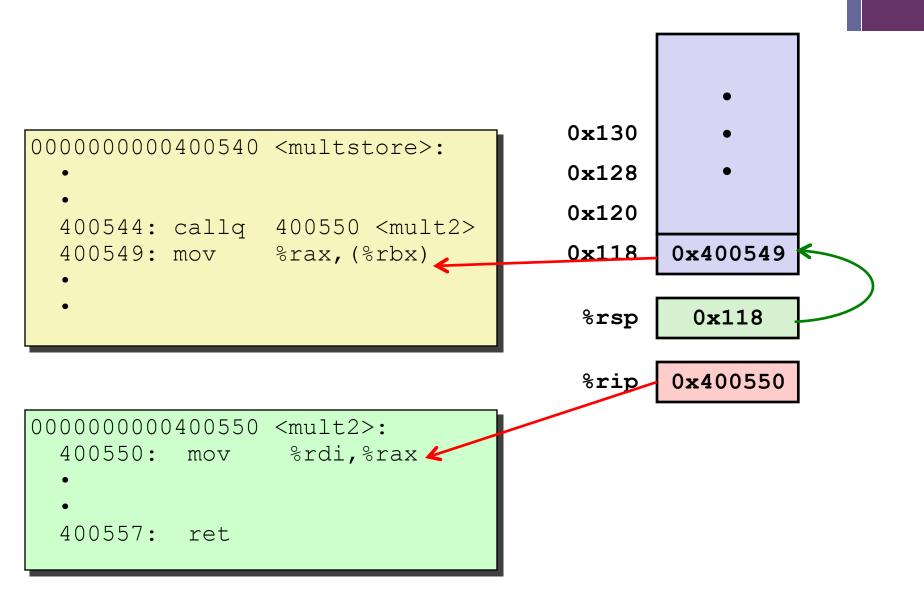
```
long mult2(long a, long b) {
  long s = a * b;
  return s;
}
```

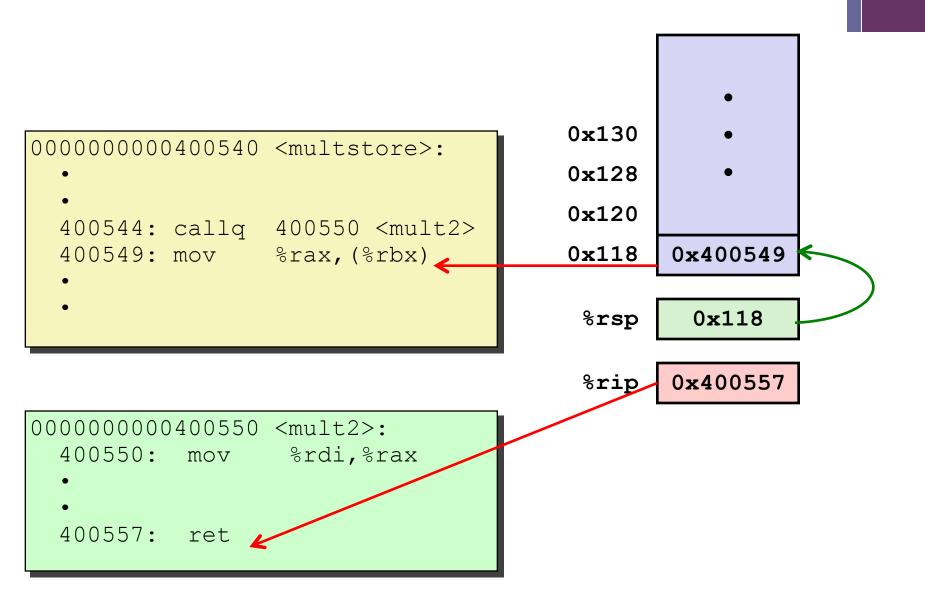
```
0000000000400550 <mult2>:
400550: mov %rdi,%rax # a
400553: imul %rsi,%rax # a * b
400557: retq # return
```

+Procedure Control Flow

- Use stack to support procedure call and return
- Procedure call with label
 - Pushes *return address* on stack
 - Address of the next instruction right after call
 - Jumps to *label*
- Procedure return: ret
 - Pops return address from stack
 - Jumps to return address

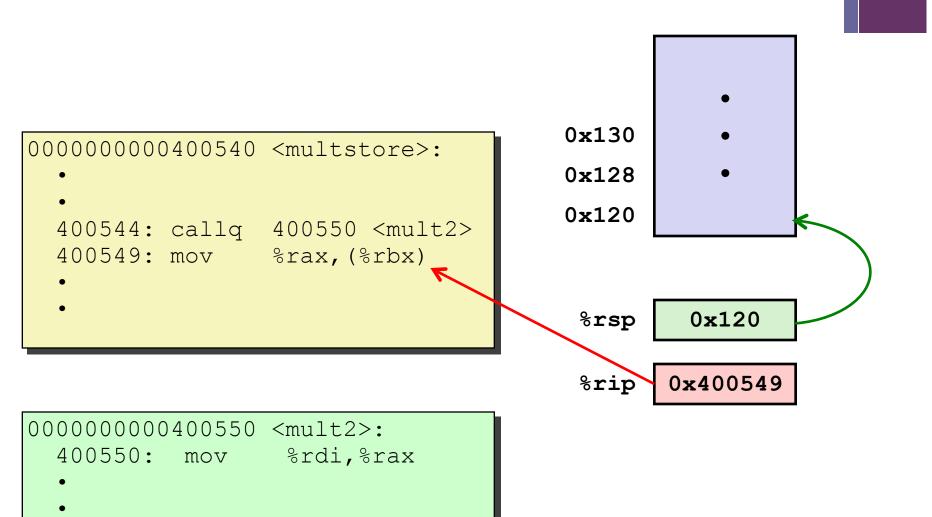






400557:

ret



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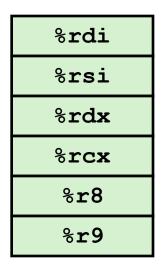
Passing Data

+Procedure Data Flow



Registers

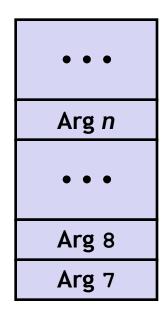
First 6 arguments



Return value

%rax

Stack



* Only allocate stack space when needed

+Data Flow Examples

```
void multstore (long x, long y, long *dest) {
   long t = mult2(x, y);
   *dest = t;
}
```

```
long mult2(long a, long b) {
  long s = a * b;
  return s;
}
```

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Managing Local Data

+Stack Frames



Functions have "instantiations"

- Every function call is a distinct execution with distinct data.
- Need some place to store state of each instantiation
 - Arguments
 - Local variables
 - Return pointer (next instruction in caller)

Stack allocated in frames

- State for single procedure instantiation
- Moreover, an allocation of memory holding all the data for some function call.

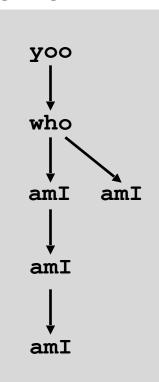
Recursion

• Supported by this idea of *instantiation* and *stack discipline*.

+Call Chain Example

```
who(...)
{
    amI();
    amI();
    amI();
}
```

Example Call Chain



Procedure amI() is recursive

+Stack Frames

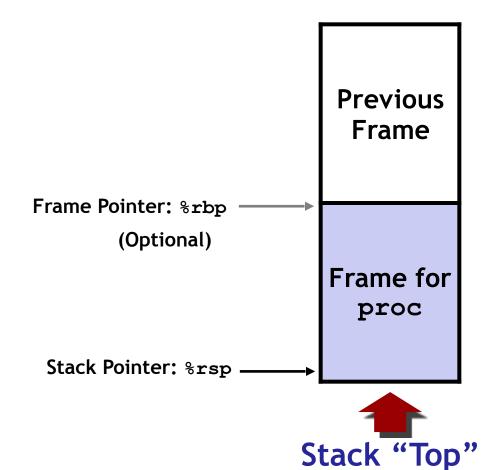


Contents

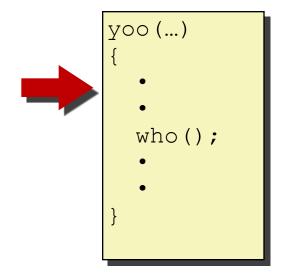
- Return information
- Local storage (if needed)

Management

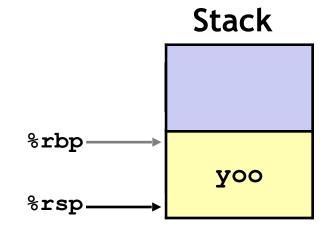
- Space allocated on procedure call
 - push by call instruction
- Space deallocated on return
 - pop by ret instruction

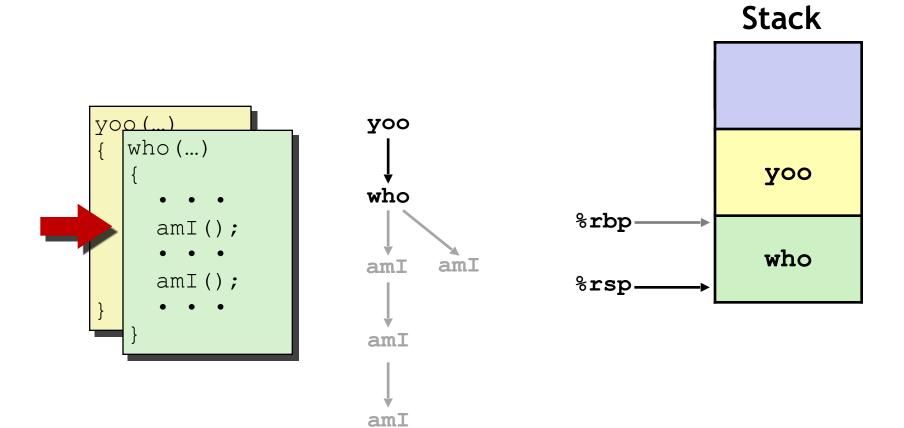


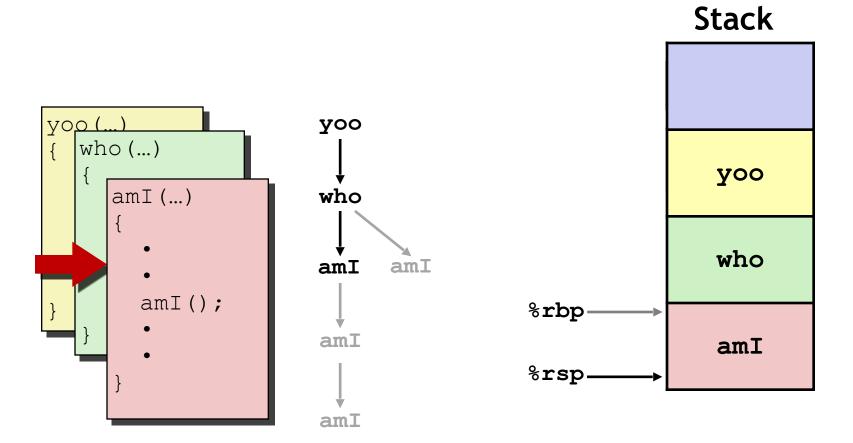


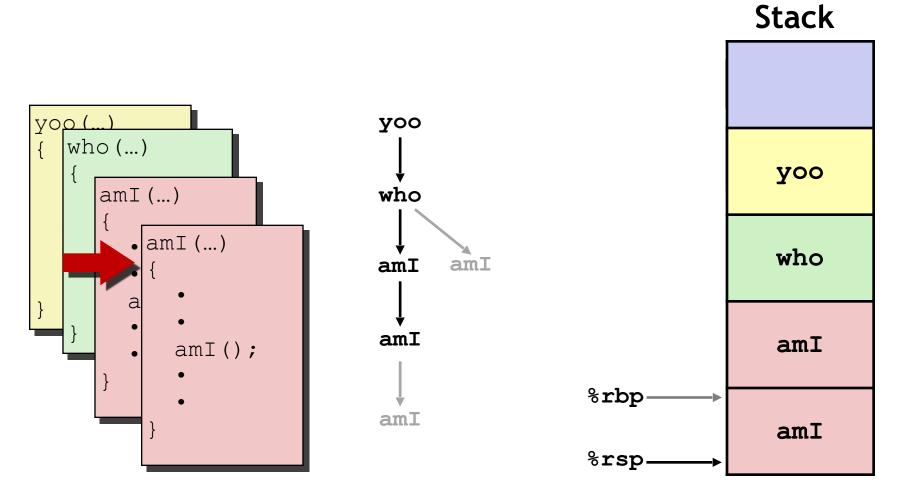


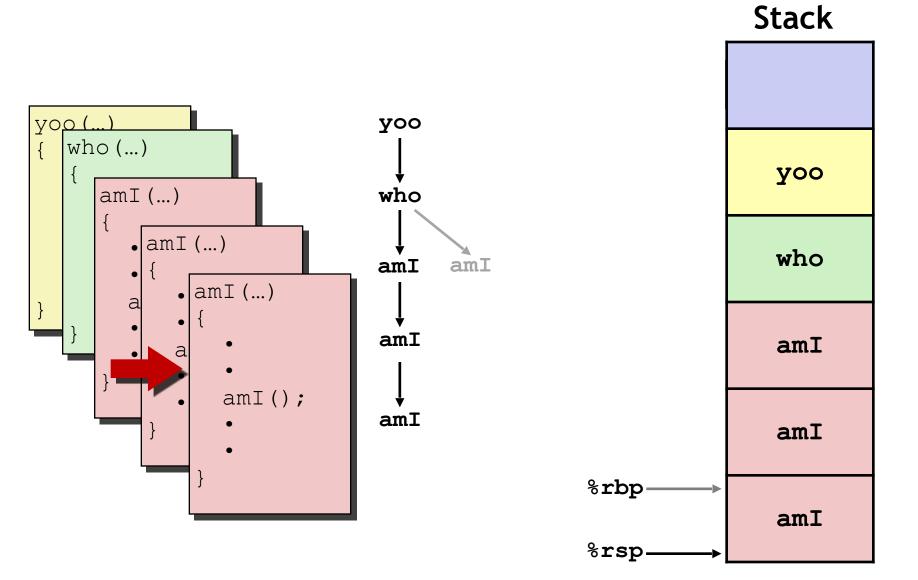


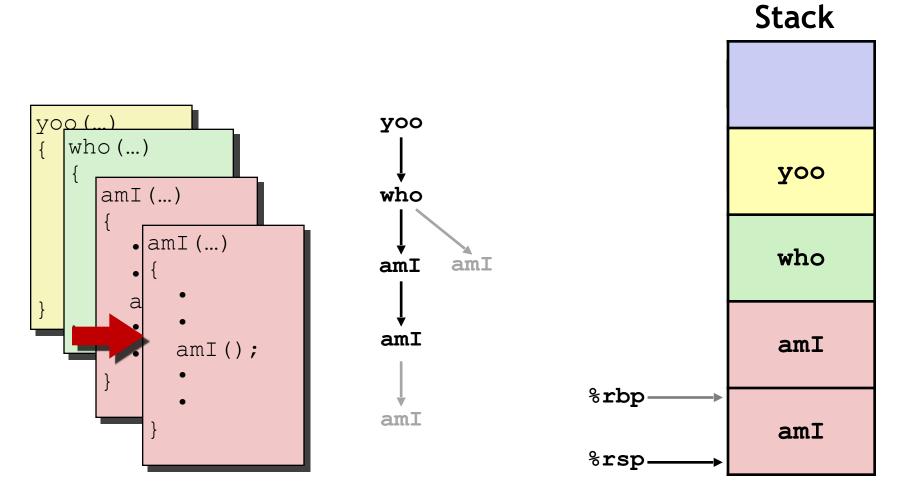


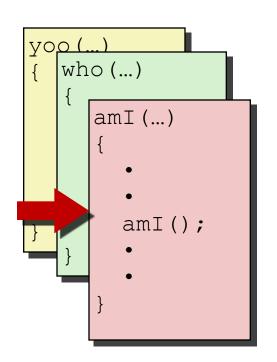


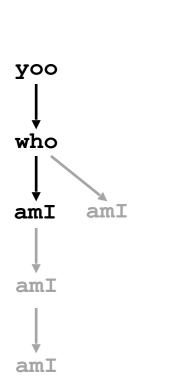


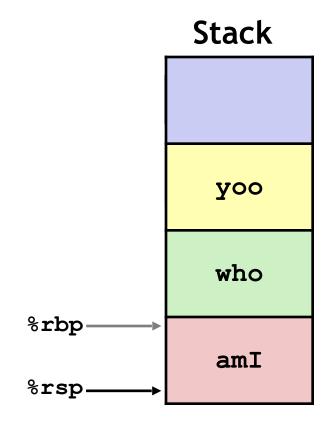




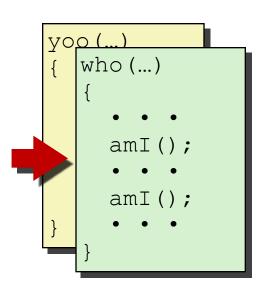


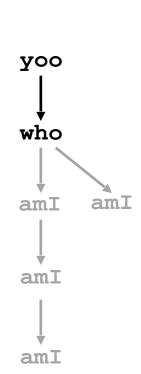


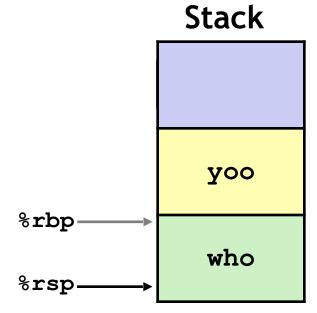




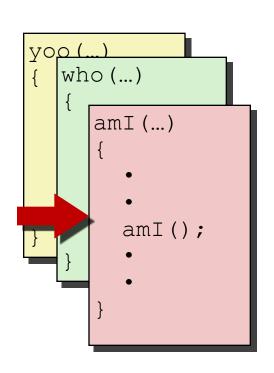


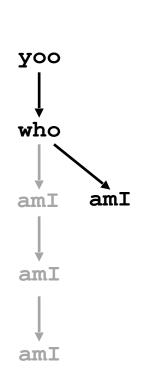


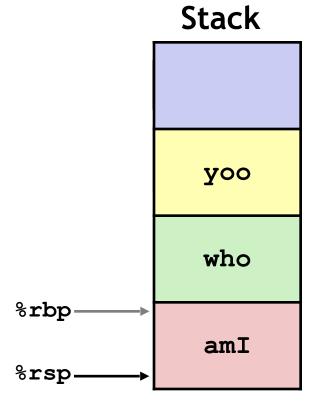


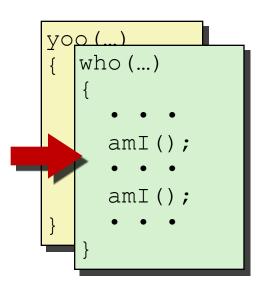




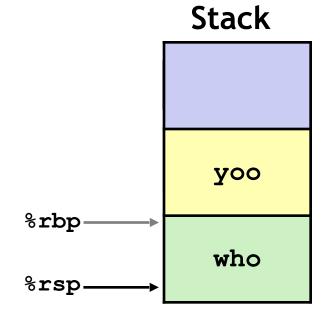




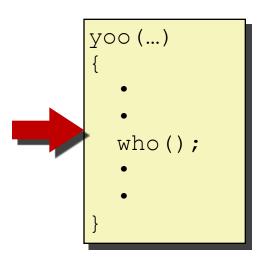




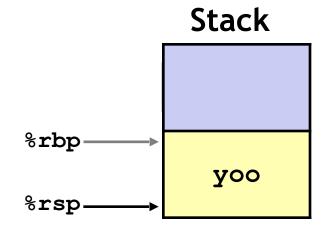












+x86-64/Linux Stack Frame

- Current Stack Frame ('Top' to 'Bottom')
 - "Argument build"
 Parameters for function about to call
 - Local variablesIf can't keep in registers
 - *Old frame pointer* (optional)
- Caller Stack Frame
 - Return address
 - Pushed by call instruction
 - Arguments for this call

