

## DC Assignment - 1

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### **Instruction:**

Python 3.x+ Version required.

1. Type the below command to run the first code, sasaki's (n-1) round algorithm.

```
python3 sasaki.py
```

The no of elements in the array can be altered in the main function by changing n, by default  $n = 10$ . The array is initialised in the worst case (inversely sorted manner i.e., descending order).

#### **Time complexity : $O(n^2)$**

The array takes (n-1) rounds to sort itself and there are no elements, so be amortised analysis it is  $O(n*(n-1)) = O(n^2)$

#### **Space complexity : $O(n)$**

K be the space required for 1 node in the line network, then  $O(n * K)$ . Here, all the space required for local computation is also included in the arbitrary constant K. This makes the space complexity  $O(n)$ .

#### **Message Exchanges: $(2*(n-1))*(n-1)$**

In a purely distributed setup where each of this node communication overhaul will outshine the amortised  $O()$  analysis, therefore the reduction in one round saves  $2(n-1)$  pairs of message exchanges which is very significant.

For each round there will be  $2*(n-1)$  message exchanges, so for n-1 rounds there will be  $(2*(n-1))*(n-1)$  message exchanges.

2. Type the below command to run the second code, alternate (n-1) round algorithm.

```
python3 alternate.py
```

#### **Time complexity : $O(n^2)$**

The array takes  $(n-1)$  rounds to sort itself and there are no elements, so be amortised analysis it is  $O(n*(n-1)) = O(n^2)$

**Space complexity :  $O(n)$**

$K$  be the space required for 1 node in the line network, then  $O(n * K)$ . Here, all the space required for local computation is also included in the arbitrary constant  $K$ .

This makes the space complexity  $O(n)$ .

**Message Exchanges:  $(2 * \text{ceil}(n/3)) * (n-1)$**

For each round there will be  $2 * \text{ceil}(n/3)$  message exchanges, so for  $n-1$  rounds there will be  $(2 * \text{ceil}(n/3)) * (n-1)$  message exchanges.