## **CS416 Project 4 Report**

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Project was tested to be fully functional and fulfilling all requirements on iLab machine cp.cs.rutgers.edu.

## 1. Detailed logic of how you implemented each API function and the logic.

get\_avail\_ino()

First, read the inode bitmap from disk. get available bitmap with get\_bitmap and set\_bitmap for that location. Return the value.

get\_avail\_blkno()

Same as the get\_avail\_ino function. But read data bitmap instead of inode bitmap.

readi(uint16\_t ino, struct inode \*inode)

alculate block number and offset first for the inode. Read the block of the inode.

Calculate block number and offset first for the inode. Read the block. Copy the target inode to the \*inode with offset.

• writei(uint16\_t ino, struct inode \*inode)

Similar to readi, calculate block number and offset. Read the block where the target inode is stored. Change the target inode's information with offset. write back to the block.

- dir\_find(uint16\_t ino, const char \*fname, size\_t name\_len, struct dirent \*dirent)
  Get inode with ino. Traverse the direct pointer of the inode and find if there is the same name for the target. If found, copy dirent of the target to the \*dirent and return 0. If not, return -ENOENT.
- dir\_add(struct inode dir\_inode, uint16\_t f\_ino, const char \*fname, size\_t name\_len)
  First, if there is already fname in the dir\_inode's direct pointers, return -EEXIST. Find the
  space for the new dirent and update the data block with the new dirent. Update the dir\_inode
  with writei.
- get\_node\_by\_path(const char \*path, uint16\_t ino, struct inode \*inode)
  Break the path with strtok((char\*)path,"/"). If the first token is NULL, it means we return the root directory. Use strtok and find inode until we get to the end. If there is no inode for the path, return -ENOENT. If found, copy the target inode to \*inode.
  - rufs mkfs()

Initialize the diskfile. Initialize superblock. Initialize bitmaps. Make the root directory inode and dirents for it. Write them all to the disk.

rufs\_init(struct fuse\_conn\_info \*conn)
 If DISKFILE is already made, read superblock and bitmaps from disk. Else, run rufs\_mkfs()

• rufs\_destroy(void \*userdata)
Free superblock and bitmaps.

- rufs\_getattr(const char \*path, struct stat \*stbuf)
   Get target inode with get\_node\_by\_path. Copy the value of the inode to the stbuf.
- rufs\_opendir(const char \*path, struct fuse\_file\_info \*fi) get path\_inode. If I could not find or was not directory, return -1. Update the atime for the inode, return 0.
  - rufs\_readdir(const char \*path, void \*buffer, fuse\_fill\_dir\_t filler, off\_t offset, struct fuse\_file\_info \*fi)

get path\_inode. If path is not found, return -ENOENT. Fill all the dirents in the path\_inode and copy them with filler function.

- rufs\_mkdir(const char \*path, mode\_t mode)
  Separate dir\_name and base\_name. I copied the original path and used them because
  dirname() and basename() modified the original path when using it. Find dir\_inode. Make a
  new inode for the new directory and make the dirents for it. add the directory to the parent
  directory's direct pointer. writei new inode and return.
- rufs\_create(const char \*path, mode\_t mode, struct fuse\_file\_info \*fi)
  Same as rufs mkdir, but we do not have dirents for the new file inode.
- rufs\_open(const char \*path, struct fuse\_file\_info \*fi)
  Same as rufs\_opendir, but this function is for the files.
  - rufs\_read(const char \*path, char \*buffer, size\_t size, off\_t offset, struct fuse\_file\_info
     \*fi)

Find the target inode. Copy the data for the file from the offset to the buffer. If copying size is larger than the size of the file, copy until the end of the file.

• rufs\_write(const char \*path, const char \*buffer, size\_t size, off\_t offset, struct fuse\_file\_info \*fi)

Find the target inode. Write the data to the data block and link them to inode's direct pointer. Update the size and modification time for the inode and writei.

## 2. Benchmark results of your file library with time and number of used blocks.

```
rsl82@cp:~/Downloads/p42/p4/benchmark$ ./test_case
TEST 1: File create Success
TEST 2: File write Success
TEST 3: File close Success
TEST 4: File read Success
TEST 5: Directory create success
TEST 7: Sub-directory create success
Benchmark completed
Time: 1.227000 ms
```

In test cases, I passed all of the tests for the program. I also added a time counter for the benchmark code, I got 1.23 ms for the execution.

Blocks used: 132

I could not pass the number of used blocks to the user-end, so I checked it through debugging. I added to print used blocks in the opendir since the opendir function was executed last in the benchmark code. And it says I used 132 blocks to execute the benchmark code.

## 3. Any challenges for the project.

I was new to the FUSE file system. So, I was confused with the structure of this system and the meaning for the dirent. After I got clarification from Piazza, there was not much hard stuff for it. The hardest part was writing code for offset for read and write but it was solved.