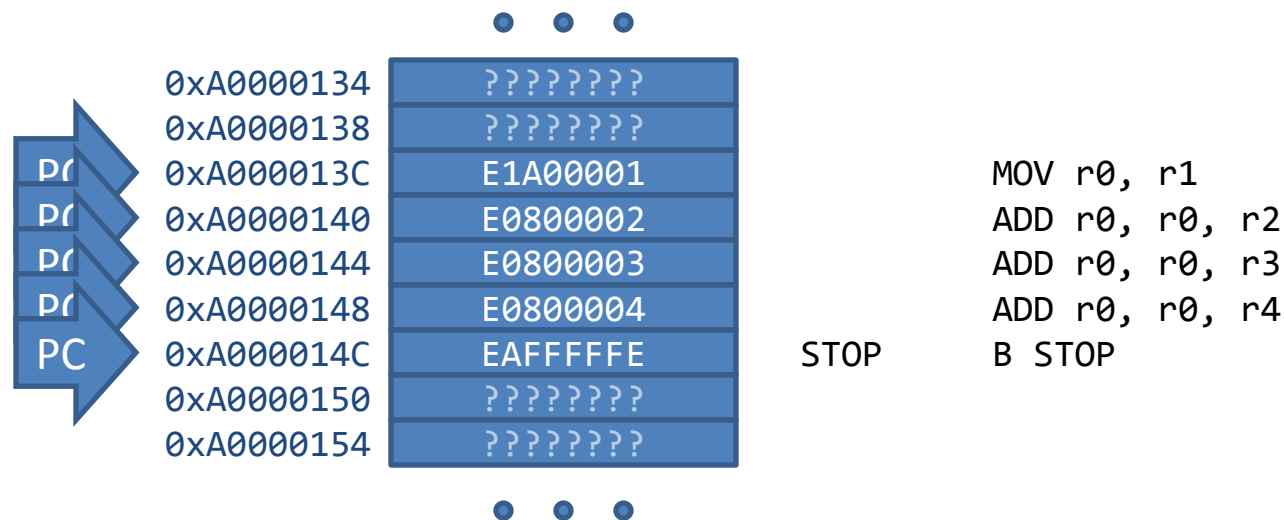


- Default flow of execution of a program is **sequential**
 - After executing one instruction, the next instruction in memory is executed sequentially by incrementing the program counter (PC)



- To write useful programs, **sequence** needs to be combined with **selection** and **iteration**

■ Selection

- if *<some condition>* then execute *<some instruction(s)>*
- if *<some condition>* then execute *<some instruction(s)>*
otherwise execute *<some other instruction(s)>*
- Examples?

■ Iteration

- while *<some condition>* is met, repeat executing *<some instructions>*
- repeat *<some instruction(s)>* until *<some condition>* is met
- repeat executing *<some instruction(s)>* *x* number of times
- Examples?

- Design and write an assembly language program to compute x^4 using repeated multiplication

```
MOV    r0, #1          ; result = 1

      MUL    r0, r1, r0    ; result = result × value (value ^ 1)
      MUL    r0, r1, r0    ; result = result × value (value ^ 2)
      MUL    r0, r1, r0    ; result = result × value (value ^ 3)
      MUL    r0, r1, r0    ; result = result × value (value ^ 4)
```

- Practical but inefficient and tedious for small values of y
- Impractical and very inefficient and tedious for larger values
- Inflexible – would like to be able to compute x^y , not just x^4

```
MOV    r0, #1          ; result = 1

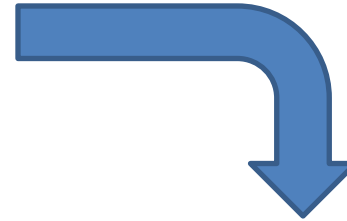
do y times:
      MUL    r0, r1, r0    ; result = result × value
repeat
```

Not valid
assembly
language!!

Program 6.1a – x^y

Iteration

```
result = 1
while (y ≠ 0) {
    result = result × x
    y = y - 1
}
```



start

```
LDR    r1, =3      ; test with x = 3
LDR    r2, =4      ; test with y = 4
```

```
MOV     r0, #1      ; result = 1
```

```
MOVS    r2, r2      ; set condition code flags
```

while

```
BEQ     endwh       ; while (y ≠ 0) {
MUL     r0, r1, r0   ;   result = result × x
SUBS    r2, r2, #1   ;   y = y - 1
B       while        ; }
```

endwh

```
stop    B          stop
```

Iteration

```
while
```

```
    BEQ    endwh  
    MUL    r0, r1, r0  
    SUBS   r2, r2, #1  
    B      while
```

```
endwh
```

```
; while (y ≠ 0) {  
;   result = result × x  
;   y = y - 1  
; }
```

- **Pseudo-code** is a useful tool for developing and documenting assembly language programs
 - No formally defined syntax
 - Use any syntax that you are familiar with (and that others can read and understand)
 - Particularly helpful for developing and documenting the **structure** of assembly language programs
 - Not always a “clean” translation between pseudo-code and assembly language

Selection

Iteration

```
if (y = 0) {  
    result = 1  
}  
else {  
    result = x  
    if (y > 1) {  
        y = y - 1  
        do {  
            result = result * x  
            y = y - 1  
        } while (y ≠ 0)  
    }  
}
```

Selection

```
start
    LDR    r1, =3        ; test with x = 3
    LDR    r2, =4        ; test with y = 4

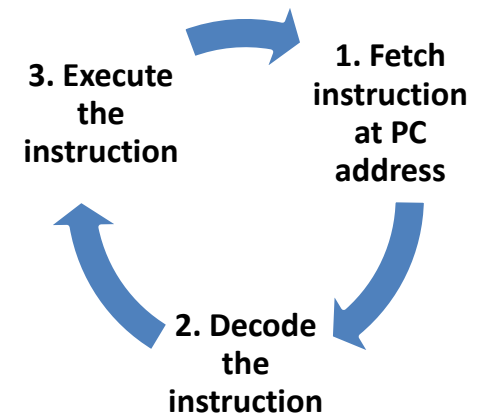
    CMP    r2, #0        ; if (y = 0)
    BNE    else1         ; {
    MOV     r0, #1        ; result = 1
    B       endif1       ; }

else1
    MOV     r0, r1        ; result = x
    CMP     r2, #1        ; if (y > 1)
    BLS     endif2       ; {
    SUBS    r2, r2, #1    ; y = y - 1
do1
    MUL     r0, r1, r0    ; result = result * x
    SUBS    r2, r2, #1    ; y = y - 1
    BNE     do1          ; } while (y ≠ 0)
endif2
endif1
    B       stop

stop
```

Comments – not assembled

- By default, the processor increments the Program Counter (PC) to “point” to the next instruction in memory ...
- ... causing the sequential path to be followed
- Using a **PC modifying instruction**, we can modify the value in the Program Counter to “point” to an instruction of our choosing, breaking the pattern of sequential execution
- PC Modifying Instructions can be
 - **unconditional** – always update the PC
 - **conditional** – update the PC only if some condition is met (e.g. the Zero condition code flag is set)



■ Unconditional Branch

B	<i>label</i>	; Branch unconditionally to label
...	...	; ...
...	...	; more instructions
...	...	; ...
<i>label</i>	some instruction	; more instructions
...	...	; ...

■ Machine code for Branch instruction



- Branch target offset is added to current Program Counter value
- Next fetch in fetch → decode → execute cycle will be from new Program Counter address

```
while      BEQ      endwh      ; while (y ≠ 0) {  
           MUL      r0, r1, r0  ;   result = result × x  
           SUBS     r2, r2, #1   ;   y = y - 1  
           B        while      ; }  
endwh
```

■ Use labels to specify branch targets in assembly language programs

- Assembler calculates necessary branch target offset

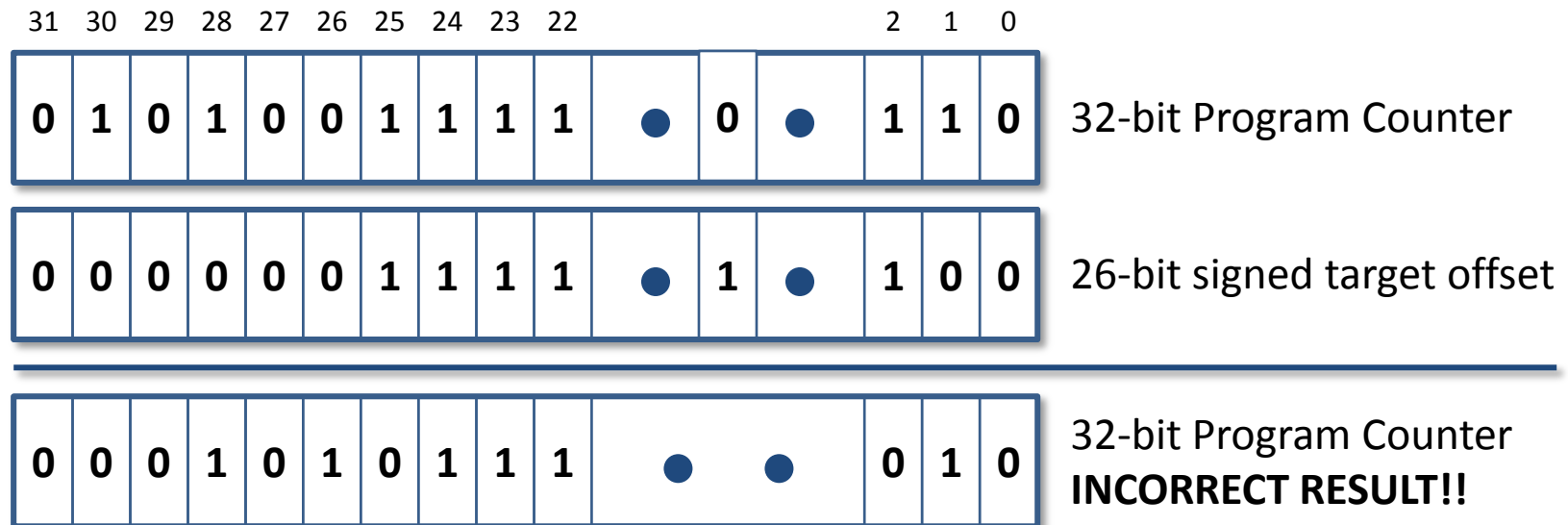
$$\text{branch target offset} = ((\text{label address} - \text{branch inst. address}) - 8) / 4$$

- Branch target offset could be negative (branch backwards)
- All ARM instructions are 4 bytes (32-bits) long and must be stored on 4-byte boundaries in memory
- So, branch target offset can be divided by 4 before being stored in the machine code branch instruction
- Allows signed 26-bit target offsets to be stored in 24 bits

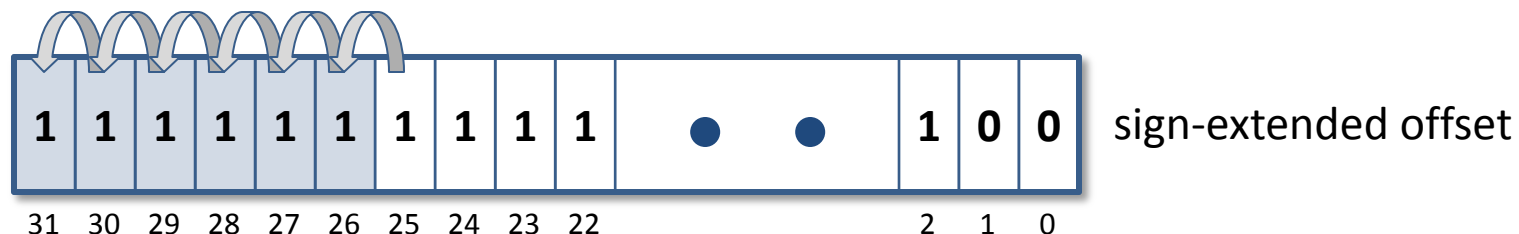


$$PC \leftarrow PC + (\textit{branch target offset} \times 4)$$

- Next fetch in fetch \rightarrow decode \rightarrow execute cycle will fetch the instruction at the new PC address
- 26-bit branch target offset may be negative
- Must sign-extend a *less-than-32-bit* value before using it to perform 32-bit arithmetic
- i.e. 26-bit branch target offset must be sign-extended to form a 32-bit value before adding it to the 32-bit Program Counter



- Must **sign extend** the 26-bit offset by copying the value of bit 25 into bits 26 to 31 (2's Complement system)



```
while      BEQ      endwh      ; while (y ≠ 0) {  
           MUL      r0, r1, r0  ;   result = result × x  
           SUBS     r2, r2, #1   ;   y = y - 1  
           B        while      ; }  
endwh
```

■ Rules

- Must be unique
- Can contain UPPER and lower case letters, numerals and the underscore _ character
- Are case sensitive (mylabel is not the same label as MyLabel)
- Must not begin with a numeral
- Further rules in the “RealView Assembler User’s Guide”
<http://www.keil.com/support/man/docs/armasm/>

- Unconditional branch instructions are necessary but they still result in an instruction execution path that is pre-determined when we write the program
- To write useful programs, the choice of instruction execution path must be deferred until the program is running
 - i.e. The decision to take a branch or continue following the sequential path must be deferred until “runtime”
- Conditional branch instructions will take a branch **only if some condition is met when the branch instruction is executed**, otherwise the processor continues to follow the sequential path

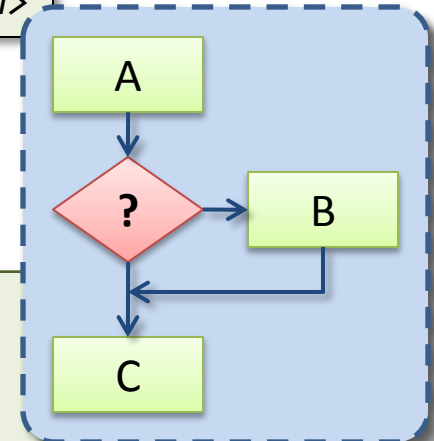
- Simple **selection** construct ...

```
if (a ≠ b) {  
    a = b  
}  
<rest of program>
```

- In ARM assembly language

- assume $a \leftrightarrow r0$, $b \leftrightarrow r1$

```
compare r0 and r1  
branch to label endif if they are equal  
MOV      r0, r1  
endif      <rest of program>
```



- Compare a and b by subtracting b from a (SUBS)
- SUBS will set Condition Code Flags. If a is equal to b, **Z**ero flag will be set. If **Z**ero flag is set, branch over $a = b$ using **BEQ**

```
SUBS      r12, r0, r1      ; store result anywhere ... not needed  
BEQ       endif          ; take branch if Zero flag set (by SUBS)  
MOV       r0, r1  
endif      <rest of program>
```

- Using SUBtract to compare two values, the result has to be stored somewhere, even though it is not needed

```
        SUBS    r12, r0, r1    ; store result anywhere ... not needed
        BEQ     endif         ; take branch if Zero flag set (by SUBS)
        MOV     r0, r1
endif    <rest of program>
```

- **CMP** (CoMPare) instruction performs a subtraction and updates the Condition Code Flags **without storing the result of the subtraction**

```
        CMP     r0, r1        ; update CC Flags, throw away result
        BEQ     endif         ; take branch if Zero flag set (by SUBS)
        MOV     r0, r1
endif    <rest of program>
```


(Un-) Conditional Branch Instructions

Branch Instruction	Condition Code Flag Evaluation	Description
B (or BAL)	don't care	unconditional (branch always)
BEQ	Z	equal
BNE	\bar{Z}	not equal
BCS / BHS	C	unsigned \geq
BCC / BLO	\bar{C}	unsigned $<$
BMI	N	negative
BPL	\bar{N}	positive or zero
BVS	V	overflow
BVC	\bar{V}	no overflow
BHI	$C\bar{Z}$	unsigned $>$
BLS	$\bar{C} + Z$	unsigned \leq
BGE	$NV + \bar{N}\bar{V}$	signed \geq
BLT	$N\bar{V} + \bar{N}V$	signed $<$
BGT	$\bar{Z}(NV + \bar{N}\bar{V})$	signed $>$
BLE	$Z + N\bar{V} + \bar{N}V$	signed \leq

- Design and write an assembly language program to compute $n!$, where n is a non-negative integer stored in register $r0$

$$n! = \prod_{k=1}^n k \quad \forall n \in \mathbb{N}$$

- Algorithm to compute the factorial of some *value*

```
result = 1
tmp = value

while (tmp > 1) {
    result = result × tmp
    tmp = tmp - 1
}
```

Program 6.2 - Factorial

```
start
    LDR    r1, =6           ; value = 6
    MOV    r0, #1           ; result = 1
    MOVS   r2, r1           ; tmp = value

wh1    CMP    r2, #1        ; while (tmp > 1)
        BLS   endwh1       ; {
        MUL   r0, r2, r0    ; result = result x tmp
        SUBS   r2, r2, #1   ; tmp = tmp - 1
        B     wh1          ; }

endwh1

stop   B      stop
```

- BLS – Branch if Lower or Same (unsigned \leq)
- Use CMP to subtract 1 from r2
 - If $r2 < 1$ there will be a borrow and the Carry flag will be clear
 - If $r2 = 1$ the Zero flag will be set
 - If $r2 > 1$ both Carry and Zero will be clear

Program 6.3 – Shift And Add Multiplication

- Design and write an assembly language program that uses shift-and-add multiplication to multiply the value in r1 by the value in r2, storing the result in r0

```
result = 0
while (b ≠ 0)
{
    b = b >> 1

    if (carry set) {
        result = result + a
    }

    a = a << 1
}
```

Program 6.3 – Shift And Add Multiplication

```
start
    LDR    r1, =10                ; test with a = 10
    LDR    r2, =6                 ; test with b = 6

    MOV    r0, #0                ; result = 0

wh1
    CMP    r2, #0                ; while (b ≠ 0)
    BEQ    endwh1                ; {
    MOVS   r2, r2, LSR #1        ; b = b >> 1
    BCC    endif1                ; if (carry set) {
    ADD    r0, r0, r1            ; result = result + a
endif1
    MOV    r1, r1, LSL #1        ; a = a << 1
    B      wh1                  ; }

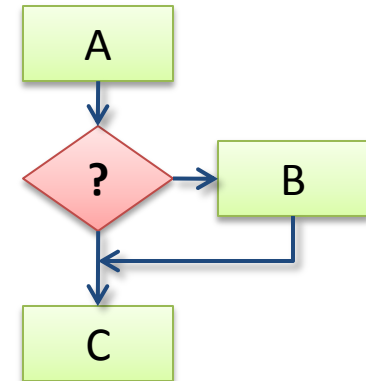
endwh1

stop    B      stop
```

- Exercise: Modify the program to avoid unnecessary iterations if a is equal to 0

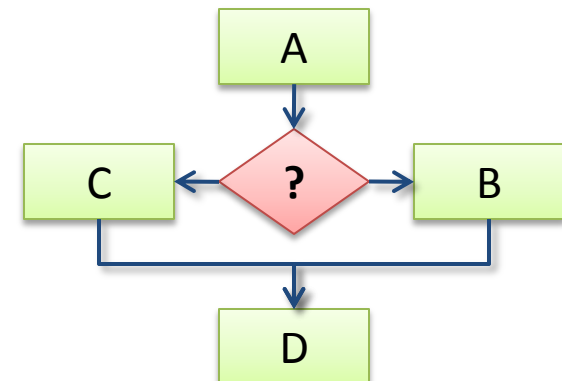
- Execute one or more instructions only if some condition is satisfied

```
if (r0 = 0) {  
    r1 = 0  
}
```



- Choose between two (or more) sets of instructions to execute

```
if (r0 = 0) {  
    r1 = 0  
}  
else {  
    r1 = r1 x r0  
}
```



■ Template for if-then construct

```
if ( <condition> )  
{  
    <body>  
}  
<rest of program>
```

```
                CMP    if necessary  
                Bxx    endif on opposite <condition>  
                <body>  
endif  
                <rest of program>
```

■ Template for if-then-else construct

```
if ( <condition> )  
{  
    <if body>  
}  
else {  
    <else body>  
}  
<rest of program>
```

```
                CMP    if necessary  
                Bxx    else on opposite <condition>  
                <if body>  
                B      endif unconditionally  
else  
                <else body>  
endif  
                <rest of program>
```

Program 6.4 – Absolute Value (if-then)

- Design and write an assembly language program to compute the absolute value of an integer stored in register r1. The absolute value should be stored in r0.

```
if (value < 0)
{
    value = 0 - value
}
```

```
start          LDR      r1, #-5          ; test with value = -5
               CMP      r1, #0          ; if (value < 0)
               BGE      endif1          ; {
               RSB      r0, r1, #0       ; result = 0 - value
endif1         ; }
stop           B        stop
```


Program 6.5 – max(a, b) (if-then-else)

- Design and write an assembly language program that evaluates the function $\max(a, b)$, where a and b are integers stored in $r1$ and $r2$ respectively. The result should be stored in $r0$.

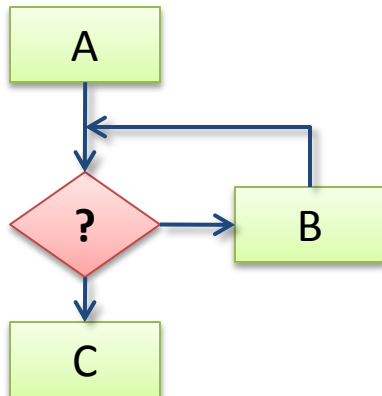
```
if (a ≥ b) {  
    max = a  
} else {  
    max = b  
}
```

```
start          LDR      r1, =5          ; test with a = 5  
               LDR      r2, =6          ; test with b = 6  
  
               CMP      r1, r2          ; if (a ≥ b)  
               BLT      else1           ; {  
               MOV      r0, r1          ; max = a  
               B        endif1          ; } else {  
else1          MOV      r0, r2          ; max = b  
endif1         ; }
```

- Execute a block of code, the loop body, multiple times
- Loop condition determines number of iterations (zero, one or more)
- Condition tested at beginning or end of loop

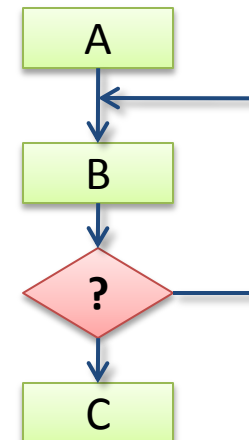
```
while ( <condition> ) {  
    <body>  
}
```

Condition tested at start of loop
Body executed zero, one or more times



```
do {  
    <body>  
} while ( <condition> )
```

Condition tested at end of loop
Body executed one or more times



■ Template for while construct

<initialize>

```
while ( <condition> )  
{  
    <body>  
}  
<rest of program>
```

<initialize>

```
while          CMP      if necessary  
                Bxx      endwh on opposite <condition>  
                <body>  
                B        while unconditionally  
endwh          <rest of program>
```

■ Template for do-while construct

<initialize>

```
do {  
    <body>  
} while ( <condition> )  
<rest of program>
```

<initialize>

```
do  
  
    <body>  
    CMP      if necessary  
    Bxx      do on <condition>  
    <rest of program>
```

- The n^{th} Fibonacci number is defined as follows

$$F_n = F_{n-2} + F_{n-1}$$

with $F_0 = 0$ and $F_1 = 1$

- Design and write an assembly language program to compute the n^{th} Fibonacci number, F_n , where n is stored in r1.

```
fn2 = 0
fn1 = 1
result = fn1
curr = 1
while (curr < n)
{
    tmp = result
    result = fn2 + fn1
    fn2 = fn1
    fn1 = tmp
    curr = curr + 1
}
```

Program 6.6 – n^{th} Fibonacci Number (while)

```
start
    LDR        r1, =4           ; test with n = 4

    MOV        r3, #0           ; fn2 = 0
    MOV        r4, #1           ; fn1 = 1
    MOV        r0, r4           ; result = fn1
    MOV        r2, #1           ; curr = 1
wh1    CMP        r2, r1         ; while (curr < n)
    BCS        endwh1          ; {
    MOV        r5, r0           ; tmp = result
    ADD        r0, r3, r4       ; result = fn2 + fn1
    MOV        r3, r4           ; fn2 = fn1
    MOV        r4, r5           ; fn1 = tmp
    ADD        r2, r2, #1       ; curr = curr + 1
    B          wh1             ; }
endwh1

stop    B          stop
```

- BCS – Branch if Carry Set (unsigned \geq)
- Use CMP to subtract r1 from r2
 - If $r2 \geq r1$ there will be no borrow and the Carry flag will be set
 - If $r2 < r1$ there will be a borrow and the Carry flag will be clear

Program 6.7 – Parity (do-while)

- Modify Program 5.6 to replace the three EOR instructions with an iterative loop using a do-while construct
- Original Program 5.6

```
start
    LDR    r0, =0x16

    MOV    r1, r0                ; tmp = value

    EOR    r1, r1, r1, LSR #1    ; tmp = tmp EOR tmp << 1
    EOR    r1, r1, r1, LSR #2    ; tmp = tmp EOR tmp << 2
    EOR    r1, r1, r1, LSR #4    ; tmp = tmp EOR tmp << 4

    AND    r1, r1, #0x00000001    ; clear all but LSB
    ORR    r0, r0, r1, LSL #7     ; set parity bit in MSB pos

stop    B    stop
```

Program 6.7 – Parity (do-while)

```
start
    LDR        r0, =0x16

    MOV        r2, #1                ; shift = 1
    MOV        r1, r0                ; tmp = value

do
    EOR        r1, r1, r1, LSR r2    ; do {
    MOV        r2, r2, LSL #1        ;   tmp = tmp EOR tmp << shift
    CMP        r2, #4                ;   shift = shift × 2
    BLS        do                    ; } while (shift ≤ 4)
    AND        r1, r1, #0x00000001   ;
    ORR        r0, r0, r1, LSL #7    ; clear all but LSB
                                         ; set parity bit in MSB pos

stop    B        stop
```

- do-while construct is appropriate as the algorithm calls for one or more iterations (never zero)
- Perform logical shift left by 1, 2 and 4 bit positions (2^0 , 2^1 and 2^2 bit positions)

- A more efficient but less intuitive while construct

<rest of program>

	B	testwh unconditionally
while	<body>	
testwh	CMP	if necessary
	Bxx	while on <condition>
	<rest of program>	

**Original
construct**

Revised construct

- Logical conjunction

```
if (x ≥ 40 AND x < 50)
{
    y = y + 1
}
```

- Test each condition and if any one fails, branch to end of if-then construct (or if they all succeed, execute the body)

...	...	
CMP	r1, #40	; if (x ≥ 40
BCC	endif	; AND
CMP	r1, #50	; x < 50)
BCS	endif	; {
ADD	r2, r2, #1	; y = y + 1
endif		; }
...	...	

- Logical disjunction

```
if (x < 40 OR x ≥ 50)
{
    z = z + 1
}
```

- Test each conditions and if they all fail, branch to end of if-then construct (or if any test succeeds, execute the body without testing further conditions)

...	...	
CMP	r1, #40	; if (x < 40
BCC	then	;
CMP	r1, #50	; x ≥ 50)
BCC	endif	; {
then	ADD	r2, r2, #1
endif		; y = y + 1
		; }
...	...	

- Design and write an assembly language program that will convert the ASCII character stored in r0 to UPPER CASE, if the character is a lower case letter (a-z)
- Can convert lower case to UPPER CASE by clearing bit 5 of the ASCII character code of a lower case letter

```
if (char ≥ 'a' AND char ≤ 'z')  
{  
    char = char . NOT(0x00000020)  
}
```

- Alternatively, subtract 0x20 from the ASCII code

```
if (char ≥ 'a' AND char ≤ 'z')  
{  
    char = char - 0x20  
}
```

Program 6.8 – Upper Case

```
start
    LDR    r0, ='d'           ; test with char = 'h'
    CMP    r0, #'a'           ; if (char ≥ 'a'
    BCC    endif              ; &&
    CMP    r0, #'z'           ; char ≤ 'z')
    BHI    endif              ; {
    AND    r0, r0, #0xFFFFFDF ; char = char . 0xFFFFFDF
;    BIC    r0, r0, #0x00000020 ; <alternative 1>
;    SUB    r0, r0, #0x20       ; <alternative 2>
endif
    B      stop
```

- Algorithm ignores characters not in the range ['a', 'z']
- Option to use AND, BIC or SUB instructions to achieve same result
- Use of #'a', #'z' for convenience instead of #61 and #7a
 - Assembler converts ASCII symbol to character code

- Branches can negatively effect performance
- Program 6.4 – Absolute Value

```
if (value < 0)
{
    value = 0 - value
}
```

- Original assembly language program

```
start
    LDR    r1, =-5        ; test with value = -5
    CMP    r1, #0         ; if (value < 0)
    BGE    endif1        ; {
    RSB    r0, r1, #0     ; result = 0 - value
endif1
    ; }
stop    B    stop
```

- ARM instruction set allows any instruction to be executed conditionally
 - based on Condition Code Flags
 - exactly the same way as conditional branches
- Revised Program 6.4 - Absolute Value

```
start
    LDR    r1, =-5        ; test with value = -5
    CMP    r1, #0         ; if (value < 0)
    RSBLT  r0, r1, #0     ; result = 0 - value
                        ; }
stop    B    stop
```

- Reverse subtract (RSB) is only executed if the less-than condition is satisfied

■ Program 6.5 – $\max(a, b)$

```
if (a ≥ b) {  
    max = a  
} else {  
    max = b  
}
```

■ Revised Program 6.5 using conditional execution

```
start  
    LDR    r1, =5                ; test with a = 5  
    LDR    r2, =6                ; test with b = 6  
  
    CMP    r1, r2                ; if (a ≥ b) {  
    MOVGE  r0, r1                ;   max = a  
    ; } else {  
    MOVLT  r0, r2                ;   max = b  
    ; }  
  
stop    B      stop
```

- Either MOVGE or MOVLT will be executed