

```
start
    MOV    total, a      ; Make the first number the subtotal
    ADD    total, total, b ; Add the second number to the subtotal
    ADD    total, total, c ; Add the third number to the subtotal
    ADD    total, total, d ; Add the fourth number to the subtotal
stop    B    stop
```

## ■ Demo program from Lecture #1

- Add four numbers together
  - **$\text{total} = a + b + c + d$**
  - **total, a, b, c, and d** are stored in memory
  - operations (move and add) are performed in CPU
  - **how many memory  $\leftrightarrow$  CPU transfers?**
- ## ■ Accessing memory is slow relative to the speed at which the processor can execute instructions
- ## ■ Processors use small fast internal storage to temporarily store values – called **registers**



`mov total, a`  
 $\Rightarrow \text{total} \leftarrow a$

`add total, total, b`  
 $\Rightarrow \text{total} \leftarrow \text{total} + b$

## ■ ARM7TDMI Registers

- 15 word-size registers, labelled **r0**, **r1**, ..., **r14**
  - Program Counter Register, **PC**, also labelled **r15**
  - Current Program Status Register (**CPSR**)
- 
- Program Counter always contains the address in memory of the next instruction to be fetched
  - CPSR contains information about the result of the last instruction executed (e.g. Was the result zero? Was the result negative?) and the status of the processor
  - r13 and r14 are normally reserved for special purposes and you should avoid using them

- A program is composed of a sequence of instructions stored in memory as **machine code**
  - Instructions determine the operations performed by the processor (e.g. add, move, multiply, subtract, compare, ...)
- A single instruction is composed of
  - an **operator (instruction)**
  - zero, one or more **operands**
- e.g. ADD the values in r1 and r2 and store the result in r0
  - Operator is ADD
  - Want to store the result in r0 (first operand)
  - We want to add the values in r1 and r2 (second and third operands)
- Each instruction and its operands are encoded using a unique value
  - e.g. 0xE0810002 is the machine that causes the processor to add the values in r1 and r2 and store the result in r3

- Writing programs using **machine code** is possible but not practical
- Instead, we write programs using **assembly language**
  - Instructions are expressed using **mnemonics**
  - e.g. the word “ADD” instead of the machine code 0xE08
  - e.g. the expression “r2” to refer to register number two, instead of the machine code value 0x2
- Assembly language must still be translated into machine code
  - Done using a program called an **assembler**
  - Machine code produced by the assembler is stored in memory and executed by the processor



# Program 1.1 revisited

```
start
    MOV     r0, r1           ; Make the first number the subtotal
    ADD     r0, r0, r2       ; Add the second number to the subtotal
    ADD     r0, r0, r3       ; Add the third number to the subtotal
    ADD     r0, r0, r4       ; Add the fourth number to the subtotal

stop    B     stop
```

# Program 1.1 – Demonstration (Demo.lst)

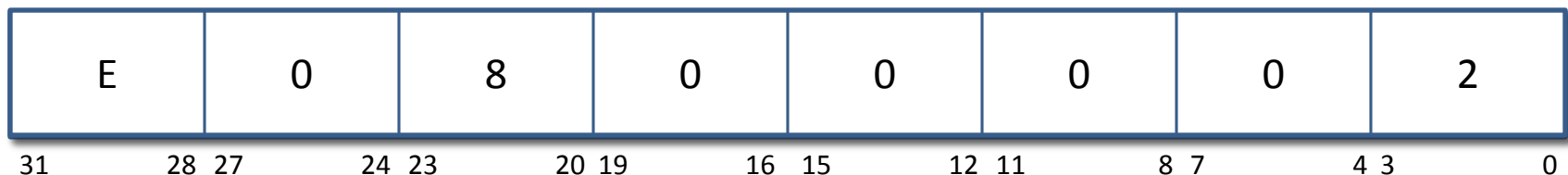
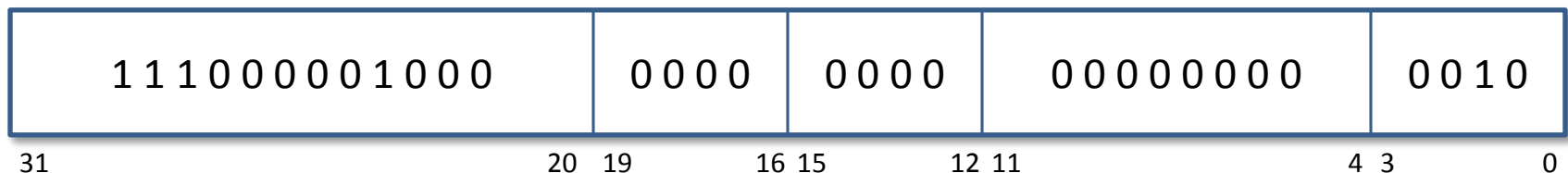
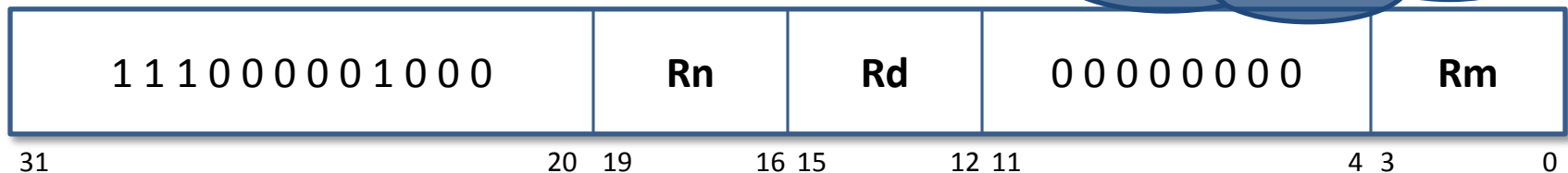
3D1 / Microprocessor Systems I  
ARM Assembly Language

```
1 00000000          AREA          Demo, CODE, READONLY
2 00000000          IMPORT        main
3 00000000          EXPORT        start
4 00000000
5 00000000          start
6 00000000 E1A00001          MOV    r0, r1
7 00000004 E0800002          ADD    r0, r0, r2
8 00000008 E0800003          ADD    r0, r0, r3
9 0000000C E0800004          ADD    r0, r0, r4
10 00000010
11 00000010 EAFFFFFE
                   stop          B          stop
12 00000014
13 00000014          END
```

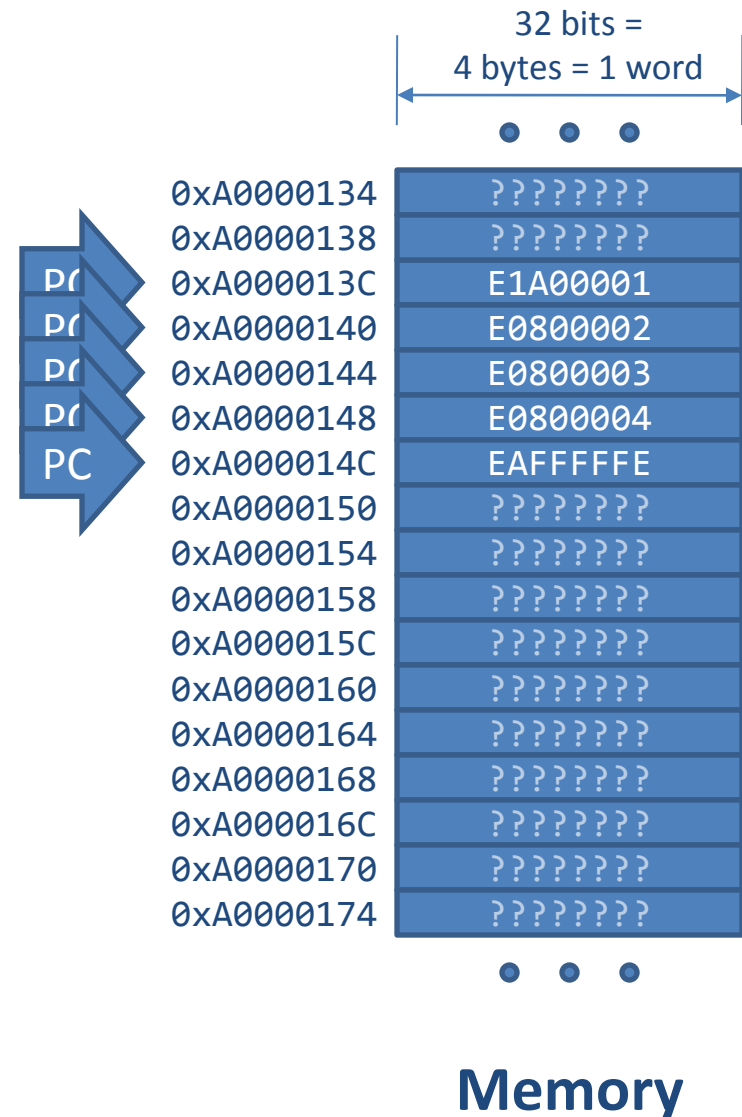
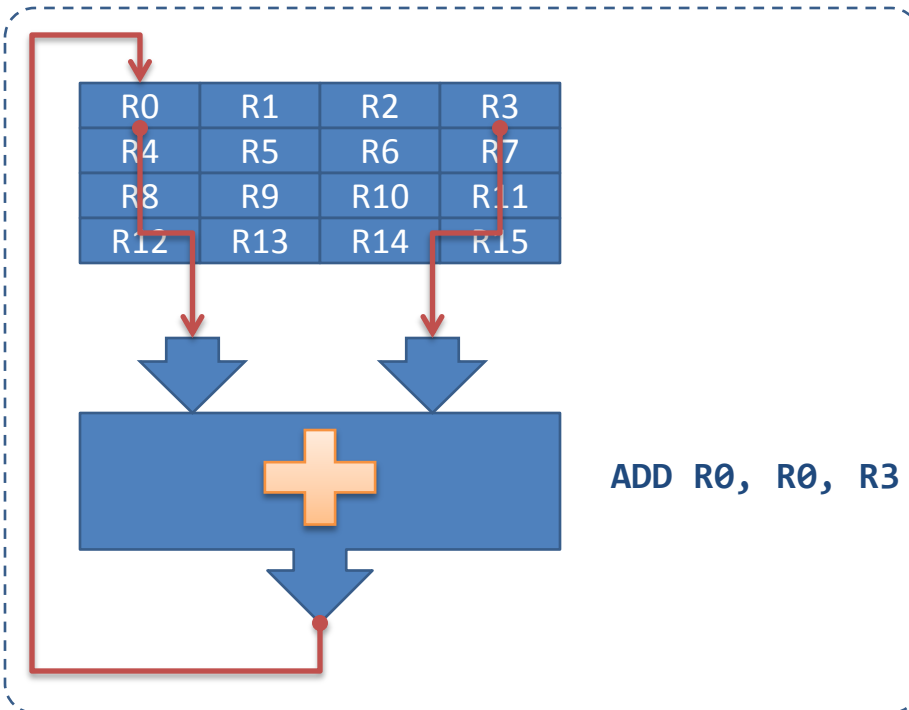
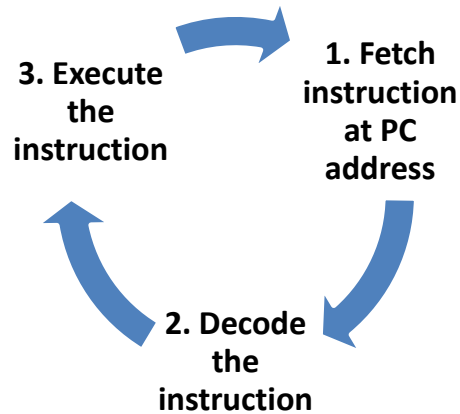
line number    address    machine code    original assembly language program

- Every ARM machine code instruction is 32-bits long
- 32-bit instruction word must encode
  - operation (instruction)
  - all the required instruction operands
- Example – add r0, r0, r2

add Rd, Rn, Rm



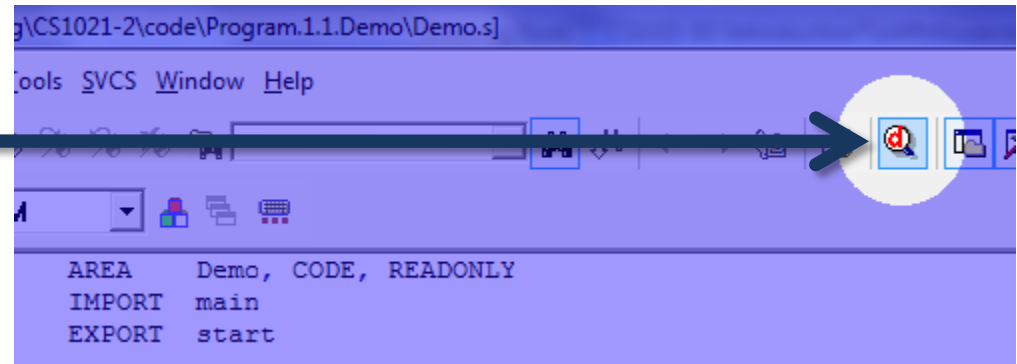
# Machine Code and Assembly Language





# Program Execution

Start Debug Session



Program assembled and loaded into memory at a pre-defined address

Program Counter (PC) set to same pre-defined address

Fetch-Decode-Execute cycle resumes

PC

What happens when we reach the end of our program?

0x9FFFFFFF8  
0x9FFFFFFFC  
0xA0000000  
0xA0000004  
0xA0000008  
0xA000000C  
0xA0000010  
0xA0000014

32 bits =  
4 bytes = 1 word

0x9FFFFFFF8	????????
0x9FFFFFFFC	????????
0xA0000000	E1A00001
0xA0000004	E0800002
0xA0000008	E0800003
0xA000000C	E0800004
0xA0000010	EAF0000E
0xA0000014	????????

Memory

## Program 3.1 – Swap Registers

- Write an assembly language program to swap the contents of register r0 and r1

```
start
    MOV    r2, r0        ; temp <-- r0
    MOV    r0, r1        ; r0 <-- r1
    MOV    r1, r2        ; r1 <-- temp
stop    B    stop
```

Compare both  
programs with respect  
to instructions  
executed and registers  
used ...

```
start
    EOR    r0, r0, r1    ; r0 <-- r0 xor r1
    EOR    r1, r0, r1    ; r1 <-- (r0 xor r1) xor r1 = r0
    EOR    r0, r0, r1    ; r0 <-- (r0 xor r1) xor r0 = r1
stop    B    stop
```

- Register operands

**ADD Rd, Rn, Rm**

**MOV Rd, Rm**

- Often want to use constant values as operands, instead of registers

**ADD Rd, Rn, #x**

**MOV Rd, #x**

- e.g. Move the value 0 (zero) into register r0

```
MOV    r0, #0           ; r0 <-- 0
```

- e.g. Set r1 = r2 + 1

```
ADD    r1, r2, #1       ; r1 <-- r2 + 1
```

- Called an **immediate operand**, syntax **#x**

- Write an assembly language program to compute ...

$$4x^2 + 3x$$

... if  $x$  is stored in  $r1$ . Store the result in  $r0$

```
start
    MUL    r0, r1, r1        ; result <-- x * x
    LDR    r2, =4            ; tmp <-- 4
    MUL    r0, r2, r0        ; result <-- 4 * x * x

    LDR    r2, =3            ; tmp <-- 3
    MUL    r2, r1, r2        ; tmp <-- x * tmp

    ADD    r0, r0, r2        ; result <-- result + tmp

stop    B      stop
```

- Cannot use MUL to multiply by a constant value
- MUL **Rx**, **Rx**, Ry produces unpredictable results **[UPDATE]**
- $r1$  unmodified ... which may be something we want ... or not

```
...      ...      ...  
LDR      r2, =3      ; tmp <-- 3  
MUL      r2, r1, r2   ; tmp <-- x * tmp  
...      ...      ...
```

- Note use of operand **=3**
  - Move constant value **3** into register r2
- LoaD Register instruction can be used to load any 32-bit signed constant value into a register

```
LDR      r4, =0xA000013C ; r4 <-- 0xA000013C
```

- Note use of **=x** syntax instead of **#x** with LDR instruction

MOV      r0, #7

6 00000000 E3A00007      MOV      r0, #7

LDR      r0, =7

6 00000000 E3A00007      LDR      r0, =7

MOV      r0, #0x4FE8

error: A1510E: Immediate 0x00004FE8 cannot be represented by 0-255 and a rotation

LDR      r0, =0x4FE8

6 00000000 E59F0000      LDR      r0, =0x4FE8

- Cannot fit large constant values in a 32-bit instruction
- LDR is a “pseudo-instruction” that simplifies the implementation of a work-around for this limitation
- For small constants, LDR is replaced with MOV

- Provide **meaningful** comments and assume someone else will be reading your code

```
MUL      r2, r1, r2      ; r2 <-- r1 * r2
```

```
MUL      r2, r1, r2      ; tmp <-- x * tmp
```

- Break your programs into small pieces
- While starting out, keep programs simple
- Pay attention to initial values in registers (and memory)