Formale Semantik
10. Montagues intensionale Logik

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Folien in Überarbeitung. Englische Teile (ab Woche 8) sind noch von 2007!
Stets aktuelle Fassungen: `https://github.com/rsling/VL-Semantik`

Inhalt

New types and up/down

Beyond truth functionality

- $[\![\phi]\!]^{\mathcal{M},w,i,g}$ and $[\![\mathbf{P}]\!]^{\mathcal{M},w,i,g}$ don't truth conditionally determine $[\![\mathbf{P}\phi]\!]^{\mathcal{M},w,i,g}$

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- Frege: sometimes expressions **denote a sense**
- again: individual concepts (variable function on indices) vs. names (constant)

$[\![\alpha]\!]^{\mathcal{M},g}_{\not{c}}$

- intension relative to models

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- for a name d: $[\![d]\!]^{\mathcal{M},g}_{\not{c}} = \begin{bmatrix} \langle w_1, t_1 \rangle & \rightarrow & b \\ \langle w_2, t_1 \rangle & \rightarrow & b \\ \langle w_1, t_2 \rangle & \rightarrow & b \\ \langle w_2, t_2 \rangle & \rightarrow & b \\ \langle w_1, t_3 \rangle & \rightarrow & b \\ \langle w_2, t_3 \rangle & \rightarrow & b \end{bmatrix}$

- for an individual concept denoting expression *m*:

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- for an individual concept denoting expression *m*:
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- $[\![\alpha]\!]^{\mathcal{M},g}_{\not{c}}(\langle w_i, t_j\rangle) = [\![\alpha]\!]^{\mathcal{M},w_i,t_j,g}$

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- alternatively: introduce rules which access an expression's extension/intension as appropriate

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- $[\![\hat{}Bm]\!]^{\mathcal{M},w,i,g} = [\![Bm]\!]^{\mathcal{M},g}_{\not{c}}$
- α and $\hat{}\alpha$ are just denoting expressions
- for an intension-denoting expression α: $[\![\check{}\alpha]\!]^{\mathcal{M},w,i,g} = [\![\alpha]\!]^{\mathcal{M},g}(\langle w, t\rangle)$

Down-up and up-down

- observe: $[\![{}^{\vee\wedge}\alpha]\!]^{\mathcal{M},w,i,g} = [\![\alpha]\!]^{\mathcal{M},w,i,g}$ for any $\langle w, t\rangle$

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- can easily be the case for intension-denoting expressions

Non-equality

- *k*' intension: $[\![k]\!]_{\mathcal{C}}^{\mathcal{M},g} = \left[\begin{array}{lll} \langle w_1,t_1\rangle & \rightarrow & \left[\begin{array}{lll} \langle w_1,t_1\rangle & \rightarrow & a \\ \langle w_1,t_2\rangle & \rightarrow & a \\ \langle w_2,t_1\rangle & \rightarrow & a \\ \langle w_2,t_2\rangle & \rightarrow & a \end{array}\right] \\ \langle w_1,t_2\rangle & \rightarrow & \left[\begin{array}{lll} \langle w_1,t_1\rangle & \rightarrow & a \\ \langle w_1,t_2\rangle & \rightarrow & b \\ \langle w_2,t_1\rangle & \rightarrow & c \\ \langle w_2,t_2\rangle & \rightarrow & d \end{array}\right] \\ \langle w_2,t_1\rangle & \rightarrow & \left[\begin{array}{lll} \langle w_1,t_1\rangle & \rightarrow & c \\ \langle w_1,t_2\rangle & \rightarrow & b \\ \langle w_2,t_1\rangle & \rightarrow & d \\ \langle w_2,t_2\rangle & \rightarrow & a \end{array}\right] \\ \langle w_2,t_2\rangle & \rightarrow & \left[\begin{array}{lll} \langle w_1,t_1\rangle & \rightarrow & c \\ \langle w_1,t_2\rangle & \rightarrow & d \\ \langle w_2,t_1\rangle & \rightarrow & a \\ \langle w_2,t_2\rangle & \rightarrow & b \end{array}\right] \end{array}\right]$

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- however: $[\![{}^{\frown}k]\!]^{\mathcal{M},w_1,t_2,g} = \begin{bmatrix} \langle w_1, t_1\rangle & \rightarrow & a \\ \langle w_1, t_2\rangle & \rightarrow & b \\ \langle w_2, t_1\rangle & \rightarrow & d \\ \langle w_2, t_2\rangle & \rightarrow & b \end{bmatrix}$
- since: $[\![{}^{\smile}k]\!]^{\mathcal{M},w_1,t_1,g} = a$
  $[\![{}^{\smile}k]\!]^{\mathcal{M},w_1,t_2,g} = b$
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The IL of PTQ

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- $s \notin Type$

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- if $\alpha, \beta \in ME_a$ then $\alpha = \beta \in ME_t$

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type	variables	constants
e	x, y, z	a, b, c
$\langle s, e\rangle$	x, y, z	—
$\langle e, t\rangle$	X, Y	$walk', A, B$
$\langle\langle s, e\rangle, t\rangle$	Q	$rise', change'$
$\langle s, \langle e, t\rangle\rangle$	P	—
$\langle e, e\rangle$	P	Sq
$\langle e, \langle e, t\rangle\rangle$	R	Gr, K
$\langle e, \langle e, e\rangle\rangle$	—	$Plus$

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- instead of: F(expression)(index)=extemsion

Some interpretations

- $[\![\lambda u\alpha]\!]^{\mathcal{M},w,i,g}$, $u \in Var_b$, $\alpha \in ME_a$ is a function h with domain D_b s.t. $x \in D_b$, $h(x) = [\![\alpha]\!]^{\mathcal{M},w,t,g'}$ with g' exactly like g except $g'(u) = x$

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- $[\![{}^\wedge\alpha]\!]^{\mathcal{M},w,i,g}$ is a function h from $W \times T$ to denotations of α's type s.t. at every $\langle w', t' \rangle \in W \times T$ $[\![\alpha]\!]^{\mathcal{M},w',t',g} = h(\langle w', t' \rangle) = [\![{}^\wedge\alpha]\!]^{\mathcal{M},w,i,g}(\langle w', t' \rangle)$

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  - $\langle e, t \rangle$ - sets of individuals
  - $\langle \langle s, e \rangle, t \rangle$- sets of individual concepts

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  - $\langle s, \langle e, t \rangle \rangle$ - properties of individuals
  - $\langle s, \langle \langle s, t \rangle, t \rangle \rangle$ - properties of propositions
- from relations $\langle e, \langle e, t \rangle \rangle$ to relations-in-intensions $\langle s, \langle e, \langle e, t \rangle \rangle \rangle$

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- hence: $\langle s, a\rangle$ never applied to some typed argument (s is not a type!)
- useful thing: We never talk about indices!
- since often ${}^{\vee}\alpha(\beta)$ is needed for $\alpha \in ME_{\langle s,\langle e,t\rangle\rangle}$ and $\beta \in ME_e$, abbr. $\alpha\{\beta\}$

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- we could give $\Box\phi$ as $\mathbf{Nec}(\,^\wedge\phi)$

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- instead of $\langle\langle e, t\rangle, \langle e, t\rangle\rangle$
- intensionally: $\langle\langle s, \langle e, t\rangle\rangle, \langle e, t\rangle\rangle$
- extensions at all indices accessible via intension: those individuals bearing property $\langle e, t\rangle$ not at current but at some past index qualify

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- intensionally: $\langle\langle s, \langle e, t\rangle\rangle, \langle e, t\rangle\rangle$
- extensions at all indices accessible via intension: those individuals bearing property $\langle e, t\rangle$ not at current but at some past index qualify
- formally: $[\![\mathbf{For}]\!]^{\mathcal{M},g}_{\not c}$ is a func. h s.t. for any property k, $h(\langle w, t\rangle)(k)$ is the set $k(\langle w, t'\rangle)$ for all $t' < t$.

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Bel

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- $\mathbf{Bel}(j, \mathbf{F}\exists x\,[R(x) \wedge W(x)])$

Literatur I

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Lizenz

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