

X-Cycles (Part 2)

From sudokuwiki.org, the puzzle solver's site

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So far, I been looking at [X-Cycles](#) which alternate perfectly all the way round. There are two interesting rules that lead to eliminations when we identify an imperfection in a loop which is called a **discontinuity**.

A discontinuity occurs when we find two strong links next to each other (that is, with no weak link between them) or two weak links next to each other (with no strong link dividing them). These rules work only if there is exactly one discontinuity, and such a loop will always have an odd number of nodes.

'Discontinuity' doesn't mean that the loop is broken or that it's not chain; it refers only to the imperfection that would otherwise make links alternate strong/weak/strong, and so on.

Nice Loops Rule 2

Here is a rule that applies in the presence of two adjacent strong links:

If the adjacent links are links with strong inference (solid line), a candidate can be fixed in the cell at the discontinuity.

This rule allows us to know the solution of a certain cell absolutely, no matter how many other candidates there may be on that cell. Unlike the case of the first Nice Loop rule, we are not looking at a mass of eliminations outside the loop; instead, this rule tells us something about the loop itself. Let's look at an example before examining the logical proof.

For discontinuous X-Cycles, the notation always starts with the discontinuity. In Figure 1, our Nice Loop on number 1 is:



Figure 1: Nice Loop on 1 : [Load Example](#) or : [From the Start](#)

X-CYCLE on 1 (Discontinuous Alternating Nice Loop, length 6):

$-1[J1]+1[G3]-1[E3]+1[E8]-1[J8]+1[J1]$

- Contradiction: When 1 is removed from J1 the chain implies it must be 1 - other candidates 3/9 can be removed

We have two strong links joined at J1; therefore, J1 is 1. One way to make sense of this logically is to trace round the alternative. If J1 was not a 1 G3 and J8 would have to be 1s. That would remove the candidate 1 from E3 and oblige E8 to be a 1. But hang on - that forces two 1s in column 8. A contradiction so the 1 must exist in J1.

Nice Loops Rule 3

Our third rule dictates what happens when two weak links form a discontinuity in a loop:

If the adjacent links are links with weak inference (broken line), a candidate can be eliminated from the cell at the discontinuity.

The brown cell is the discontinuity based on two weak links that are next to each other in the loop. We can safely eliminate the 1 from this node. It might not seem much of an elimination considering how powerful the previous two rules are, but this type of Nice Loop configuration - two weak loops - is actually the most common.

The solver would return this message:

X-CYCLE on 1 (Discontinuous Alternating Nice Loop, length 6):

$+1[C3]-1[C7]+1[G7]-1[G2]+1[H3]-1[C3]$

- Contradiction: When C3 is set to 1 the chain implies it cannot be 1 - it can be removed



Figure 2: Nice Loop on 1 : [Load Example](#) or : [From the Start](#)

Just a little further on from we have some more AICs including this 8 elimination

X-CYCLE on 8 (Discontinuous Alternating Nice Loop, length 6):

$+8[B7]-8[B1]+8[C3]-8[E3]+8[E7]-8[B7]$

- Contradiction: When B7 is set to 8 the chain implies it cannot be 8 - it can be removed

Weak and Strong Links

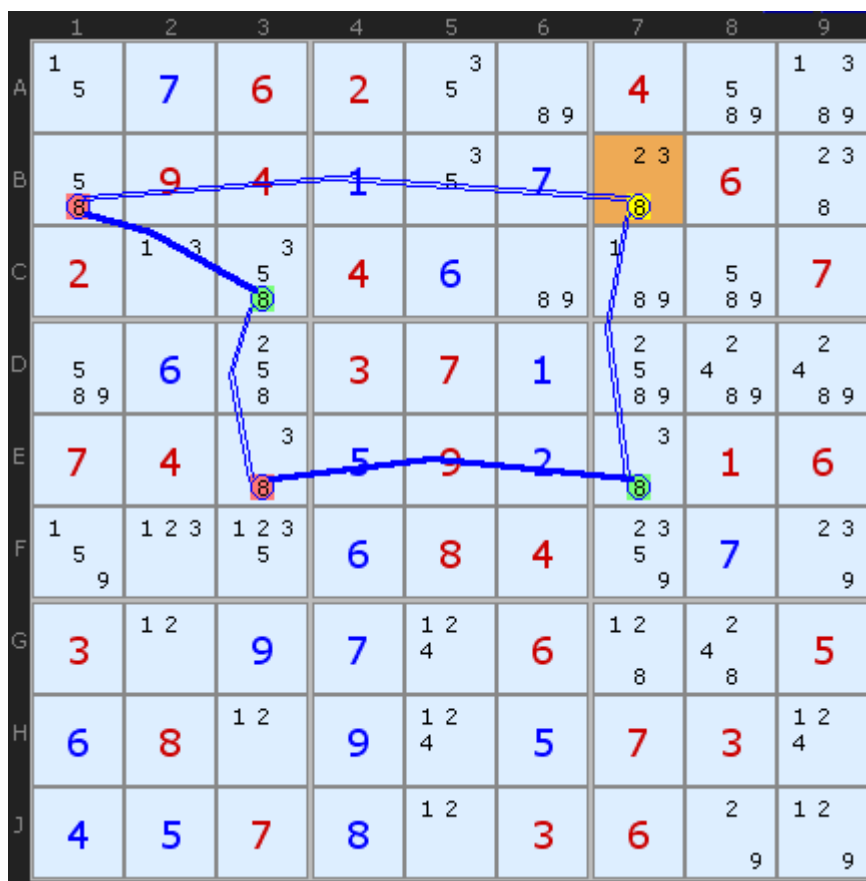


Figure 3: Nice Loop on 8 : [Load Example](#) or : [From the Start](#)

X-Cycles introduced the idea of Weak and Strong links but I want to make a more precise definition of terms since there are subtleties which will be useful in other chaining strategies. The rough and ready distinction between Strong and Weak links is to do with how many candidates are in a unit – namely, Strong links are formed when only two are present, while three or more imply a Weak link.

From a strong link we can infer that
if not A, then B

From a weak link, we can infer only that
if A then not B, C, D according to how many candidates there are in a unit

This implies that:

- Strong links are "links with strong inference"; and
- Weak links are "links with weak inference".

However, the following is also true that for a strong link:
if A, then not B

So, some Strong links can be reversed to give us a "link with weak inference" - if the occasion calls for it. It is perfectly logical to assert on a unit with two candidates of X both:

- If Not A then B ($\neg A \Rightarrow B$)

- If A then Not B ($A \Rightarrow !B$)

In Figure 5 we have an array of 6 candidates on a board. A number of strategies can show that the 6 on **H9** can be eliminated. I have coloured some cells using Simple Colouring Rule 2 which link up some pairs on the board - either all of the yellow cells will be 6 or all of the cyan cells will be 6. Since **H9** can see **C9** (yellow) and **H5** (cyan) it cannot be a 6 since it can see cells with both colours.

Now, we can also create a Nice Loop as I have done with blue lines. Our aim is to show that the circled 6 on **H9** is eliminated because there are two weak links forming a discontinuity. That is all correct and invokes Nice Rule 3. But there seem to be three strong links joined up. What happened to the alternating nature of the X-Cycle?

If a strong link can have weak inference, then let's just change the link from **C4** to **A5** to imply such. Simple. We get our pattern. If 6 is on **C4**, then it is not on **A5** (weak inference), or if it is on **A5**, then it is not on **C4** (also weak inference – and all very logical).

I have coloured the Strong link with weak inference red in Figure 5.

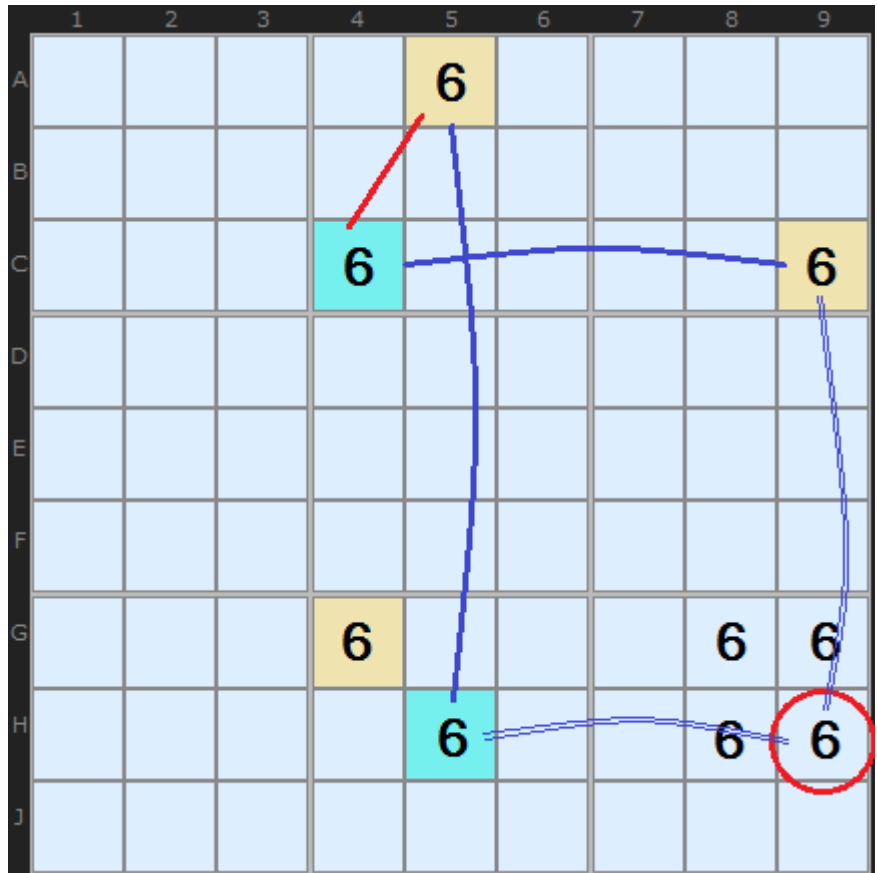


Figure 5: Colouring Example and Nice Loop

X-Cycle Exemplars

These puzzles require the X-Cycle strategy at some point but are otherwise trivial.

They also require just one Naked Pair.

They make good practice puzzles.

- [Exemplar 1, x3 \(score 162\)](#)
- [Exemplar 2, x3 \(score 167\)](#)
- [Exemplar 3, x2 \(score 173\)](#)

Go back to [X-Cycles \(Part 1\)](#) Continue to [Grouped X-Cycles](#)

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