Synchronization I

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Today's Topics

Synchronization problem

Locks

Synchronization

Threads cooperate in multithreaded programs

- To share resources, access shared data structures
- Also, to coordinate their execution

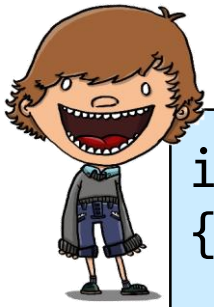
For correctness, we have to control this cooperation

- Must assume threads interleave executions arbitrarily and at different rates
  - Scheduling is not under application writers' control
- We control cooperation using synchronization
  - Enables us to restrict the interleaving of execution
- (Note) This also applies to processes, not just threads
  - And it also applies across machines in a distributed system

The Classic Example (1)

Withdraw money from a bank account

- Suppose you and your girl(boy) friend share a bank account with a balance of 1,000,000 won
- What happens if both go to separate ATM machines, and simultaneously withdraw 100,000 won from the account?



```
int withdraw (account, amount)
{
   balance = get_balance (account);
   balance = balance - amount;
   put_balance (account, balance);
   return balance;
}
```

The Classic Example (2)

Interleaved schedules

- Represent the situation by creating a separate thread for each person to do the withdrawals
- The execution of the two threads can be interleaved, assuming preemptive scheduling:

Execution sequence as seen by CPU

```
balance = get_balance (account);
balance = balance - amount;
```

Context switch

```
balance = get_balance (account);
balance = balance - amount;
put_balance (account, balance);
```

Context switch

```
put_balance (account, balance);
```

Synchronization Problem

Problem

- Two concurrent threads (or processes) access a shared resource without any synchronization
- Creates a race condition:
  - The situation where several processes access and manipulate shared data concurrently
  - The result is non-deterministic and depends on timing
- Need mechanisms for controlling access to shared resources
  - So that we can reason about the operation of programs
- Synchronization is necessary for any shared data structure
  - buffers, queues, lists, etc.

Threads using shared data

```
#include <pthread.h>
#define MAX_THREAD 20

void *threadcount(void *data) {
        int *count = (int *)data;
        int i;
        for (i=0; i<100; i++) {
                *count = *count+1;
        }
}
int main(int argc, char **argv) {
        pthread_t thread_id[MAX_THREAD];
        int i = 0;
        int count = 0;
        for(i = 0; i < MAX_THREAD; i++) {
                pthread_create(&thread_id[i], NULL, threadcount, (void *)&count);
        }
        for(i = 0; i < MAX_THREAD; i++) {
                pthread_join(thread_id[i], NULL);
        }
        printf("Main Thread : %d\n", count);
        return 0;
}
```

```
# gcc ex.c -lpthread
# ./a.out
Main Thread : 2000
# ./a.out
Main Thread : 1957
```

Threads using shared data

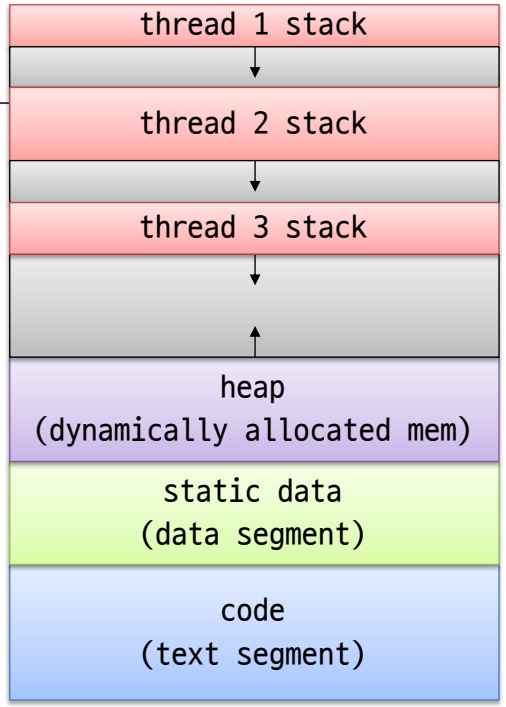
count=count+1 could be implemented as

- register1 = count (LOAD R1, MEM_count)
- register1 = register1 + 1 (ADD R1, R1, 1)
- count = register1 (STORE R1, MEM_count)

Consider this execution interleaving with "count = 5" initially:

1. T1: register1 = count {register1 = 5}
2. T1: register1 = register1 + 1 {register1 = 6}
3. T2: register2 = count {register2 = 5}
4. T2: register2 = register2 + 1 {register2 = 6}
5. T1: count = register1 {count = 6}
6. T2: count = register2 {count = 6}

Sharing Resources



Between threads

- Local variables are not shared
  - Refer to data on the stack
  - Each thread has its own stack
  - Never pass/share/store a pointer to a local variable on another thread's stack
- Global variables are shared
  - Stored in static data segment, accessible by any thread
- Dynamic objects are shared
  - Stored in the heap, shared through the pointers

Between processes

- Shared-memory objects, files, etc. are shared

Critical Sections (1)

Critical sections

- Parts of the program that access shared resources
  - Shared files, shared memory(variable), etc.
- Use mutual exclusion to synchronize access to shared resources in critical sections
  - Only one thread at a time can execute in the critical section
  - All other threads are forced to wait on entry
  - When a thread leaves a critical section, another can enter
- Otherwise, critical sections can lead to race conditions
  - The final result depends on the sequence of execution of the processes

```
int withdraw (account, amount)
{
   balance = get_balance (account);
   balance = balance – amount;
   put_balance (account, balance);
   return balance;
}
```

Critical Sections (2)

Requirements

- Mutual exclusion
  - At most one thread is in the critical section
- Progress
  - If thread T is inside the critical section, T must finish the critical section within reasonable time
  - If thread T is outside the critical section, then T cannot prevent thread S from entering the critical section
- Bounded waiting (no starvation)
  - If thread T is waiting on the critical section, then T will eventually enter the critical section
- Performance
  - The overhead of entering and exiting the critical section is small with respect to the work being done within it

Critical Sections (3)

Mechanisms for building critical sections

- Locks
  - Very primitive, minimal semantics, used to build others
- Semaphores
  - Basic, easy to get the hang of, hard to program with
- Monitors
  - High-level, requires language support, implicit operations
  - Easy to program with
    - e.g. Java "synchronized"
- Messages
  - Simple model of communication and synchronization based on (atomic) transfer of data across a channel
  - Direct application to distributed systems

Locks

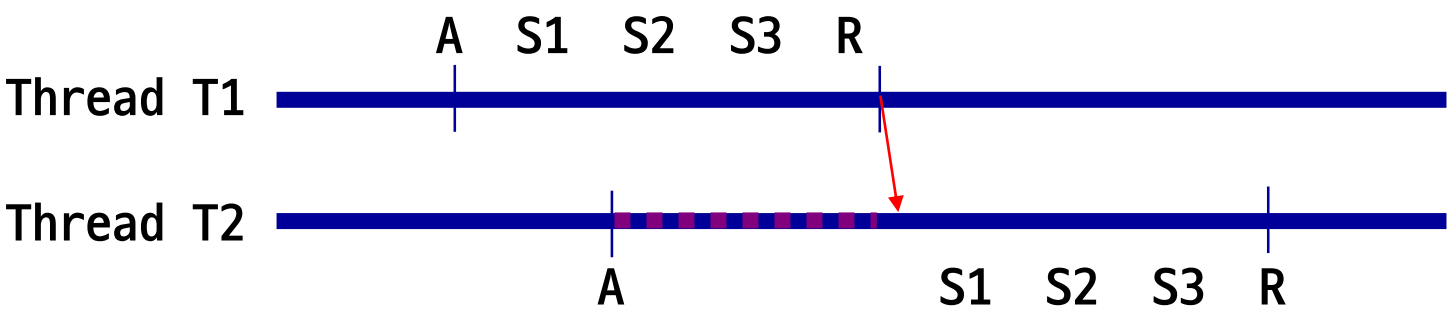
Locks

- A lock is an object (in memory) that provides the following two operations:
  - acquire(): wait until lock is free, then grab it
  - release(): unlock, and wake up any thread waiting in acquire()
- Using locks
  - Lock is initially free
  - Call acquire() before entering a critical section, and release() after leaving it
  - Between acquire() and release(), the thread holds the lock
  - acquire() does not return until the caller holds the lock
  - At most one thread can hold a lock at a time
- Locks can spin (a spinlock) or block (a mutex)

Using Locks

```
int withdraw (account, amount)
{
A       acquire (lock);
S1      balance = get_balance (account);
S2      balance = balance - amount;
S3      put_balance (account, balance);
R       release (lock);
        return balance;
}
```

Critical section



Implementing Locks (1)

An initial attempt

운좋게 CPU를 받은 애가 먼저 됨

```
struct lock { int held = 0; }

void acquire (struct lock *l) {
   while (l->held);
   l->held = 1;
}
void release (struct lock *l) {
   l->held = 0;
}
```

The caller "busy-waits", or spins for locks to be released, hence spinlocks

```
int withdraw (account, amount)
{
   acquire (lock);
   balance = get_balance (account);
   balance = balance - amount;
   put_balance (account, balance);
   release (lock);
   return balance;
}
```

Implementing Locks (1)

An initial attempt

```
struct lock { int held = 0; }

void acquire (struct lock *l) {
   while (l->held);
   l->held = 1;
}
void release (struct lock *l) {
   l->held = 0;
}
```

```
void acquire (struct lock *l) {
   while (l->held);
   l->held = 1;
}
void release (struct lock *l) {
   l->held = 0;
}
```

```
int withdraw (account, amount)
{
   acquire (lock);
   balance = get_balance (account);
   balance = balance - amount;
   put_balance (account, balance);
   release (lock);
   return balance;
}
```

Implementing Locks (2)

Problem

- Implementation of locks has a critical section, too!
  - The acquire/release must be atomic
  - A recursion, huh?

Atomic operation

- Executes as though it could not be interrupted
- Code that executes "all or nothing"

Implementing Locks (3)

Solutions

- Software-only algorithms
  - Dekker's algorithm (1962)
  - Peterson's algorithm (1981)
  - Lamport's Bakery algorithm for more than two processes (1974)
- Hardware atomic instructions
  - Test-and-set, compare-and-swap, etc.
- Disable/reenable interrupts
  - To prevent context switches

Software-only Algorithms

Wrong algorithm

- Mutual exclusion?
- Progress?

코드를 내놓고 어디에 context switch
를 놔야 문제가 안생기는가. 이런거 문제 나옴

Thread 0	Thread 1
...	...
acquire(0)	acquire(1)
do C.S.	do C.S.
release(0)	release(1)

context switch가 일어나는 곳에 따라
어떻게 되는지 생각해볼 필요 있음.

```
int interested[2];
interested[0]=FALSE; interested[1]=FALSE;
```

```
void acquire (int process)   {
   int other = 1 - process;
   interested[process] = TRUE;
   while (interested[other]);
}

void release (int process)   {
   interested[process] = FALSE;
}
```

```
void acquire (int process)   {
   int other = 1 - process;
   interested[process] = TRUE;
   while (interested[other]);
}

void release (int process)   {
   interested[process] = FALSE;
}
```

Peterson's Algorithm

Solves the critical section problem for two threads/processes

```
Thread 0
...
acquire(0)
do C.S.
release(0)
```

```
Thread 1
...
acquire(1)
do C.S.
release(1)
```

```
int turn;
int interested[2];
interested[0]=FALSE; interested[1]=FALSE;
```

```
void acquire (int process)   {
   int other = 1 – process;
   interested[process] = TRUE;
   turn = other;
   while (interested[other] &&
          turn == other);
}

void release (int process)   {
   interested[process] = FALSE;
}
```

```
void acquire (int process)   {
   int other = 1 – process;
   interested[process] = TRUE;
   turn = other;
   while (interested[other] &&
          turn == other);
}

void release (int process)   {
   interested[process] = FALSE;
}
```

Atomic Instructions (1)

Test-and-Set

- Mostly supported by H/W

```
int TestAndSet (int *v)   {
    int rv = *v;
    *v = 1;
    return rv;
}
```

Atomic

Using Test-and-Set instruction

```
void struct lock { int value = 0; }

void acquire (struct lock *l)   {
    while (TestAndSet (&l->value));
}
void release (struct lock *l)   {
    l->value = 0;
}
```

Atomic Instructions (2)

Swap

- Mostly supported by H/W

```
void Swap (int *v1, int *v2)   {
    int temp = *v1;
    *v1 = *v2;
    *v2 = temp;
}
```

Atomic

Using Swap instruction

```
void struct lock { int value = 0; }
void acquire (struct lock *l)   {
    int key = 1;
    while (key == 1) Swap(&l->value, &key);
}
void release (struct lock *l)   {
    l->value = 0;
}
```

Problems with Spinlocks

Spinlocks

- Horribly wasteful!
  - CPU cycle is wasted
  - The longer the critical section, the longer the spin
  - Lock holder can be interrupted through involuntary context switch
- Only want to use spinlock as primitives to build higher-level synchronization constructs

Disabling Interrupts (1)

Implementing locks by disabling interrupts

```
void acquire (struct lock *l)  {
    cli();      // local_irq_disable(), disable interrupts;
}
void release (struct lock *l)  {
    sti();      // local_irq_enable(), enable interrupts;
}
```

- Disabling interrupts
  - Blocks notification of external events
  - No context switch (e.g., timer)
- Can two threads disable interrupts simultaneously?

Disabling Interrupts (2)

What's wrong?

- Only available to kernel
  - If OS support these as system calls ?
- Insufficient on a multiprocessor
  - Back to atomic instructions
- What if the critical section is long?
  - Can miss or delay important events (e.g., timer, I/O)

Like spinlocks, only use to implement higher-level synchronization primitives

Summary

Implementing locks

- Software-only algorithms
- Hardware atomic instructions
- Disable/reenable interrupts

Spinlocks and disabling interrupts are primitive synchronization mechanisms

- They are used to build higher-level synchronization constructs