

*Bibliotech - smart cafe*

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# Programming Languages: Some news for the last N years.

Ruslan Shevchenko  
<[ruslan@shevchenko.kiev.ua](mailto:ruslan@shevchenko.kiev.ua)>

@rssh1

<https://github.com/rssh>

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## Programming languages:

- ❖ Industry / Academic
  - ❖ what's hot now:
    - ❖ academic research from 70-th
    - ❖ 'freaky' things from 2000

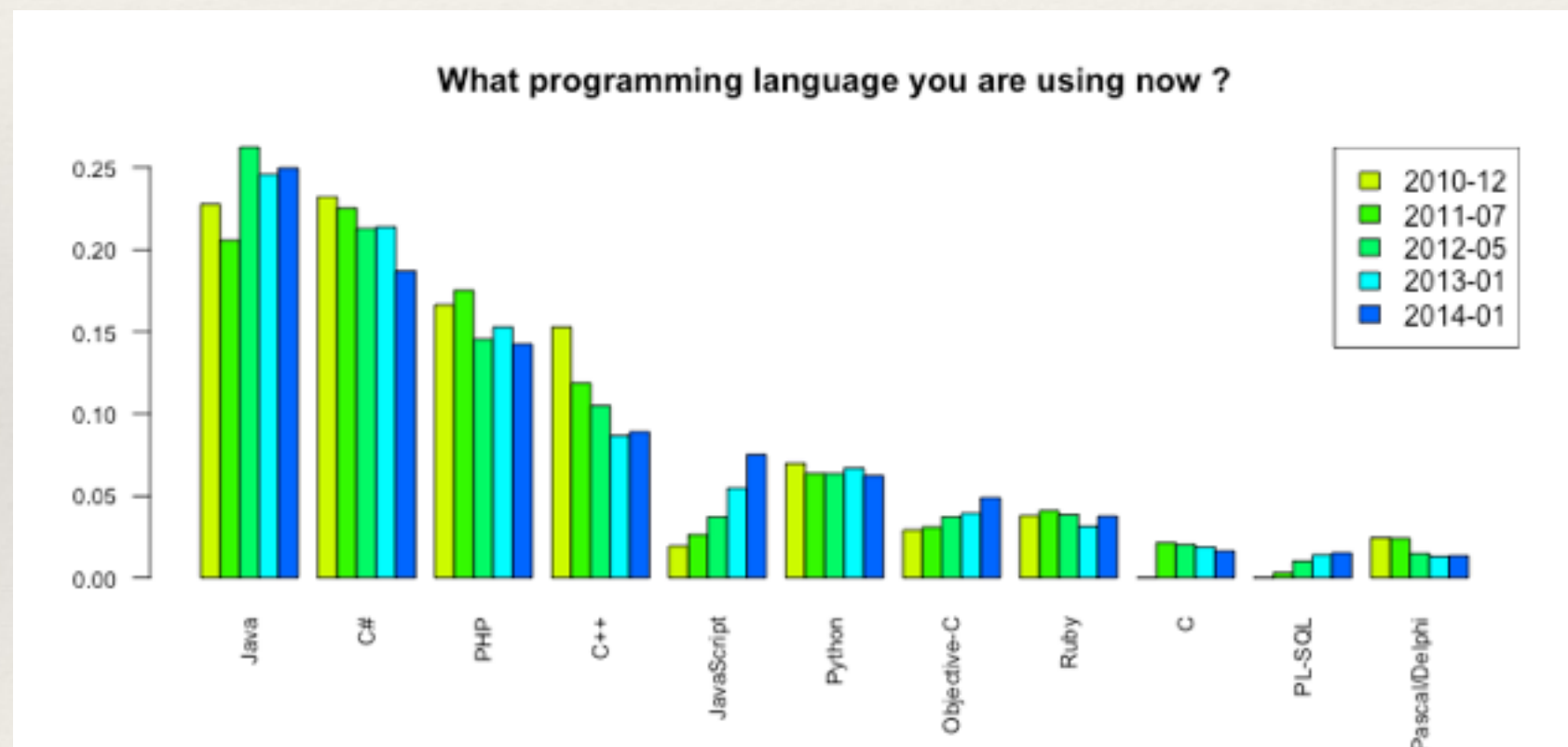
Where is progress ?

What we will see in industry from today-s academic ?  
from today-s freaks ?

Statement from 80-th:

Full change of language and tools during 2-3 years

<http://dou.ua>

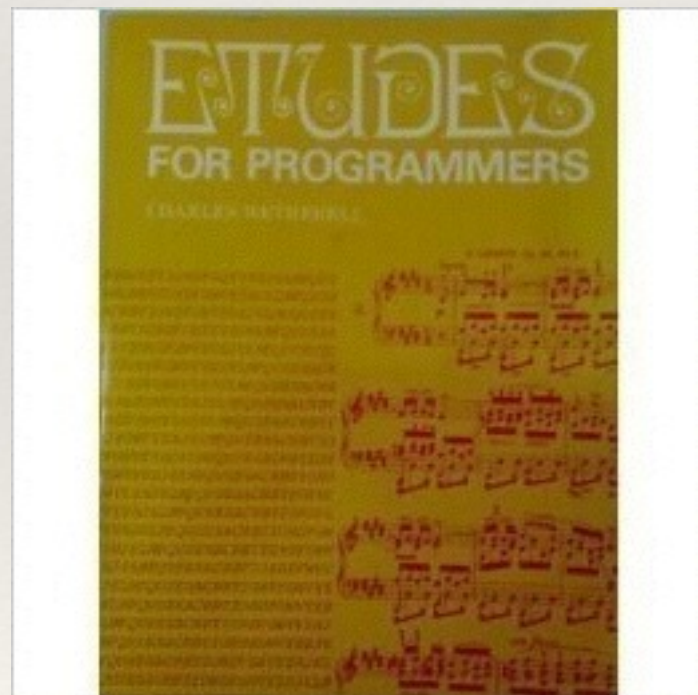


Actual situation in 2014:

Top 11 languages are the same: 2010 - 2014



Can we track progress in PL ?



- ❖ Estimations
- ❖ Code size / project duration

Evolution is slow





If we think about ideal language ?





What is:

- ❖ all ?
- ❖ good ?
- ❖ make ?

# Make all good

- Object and processes from real world, represented in computer
- Operations with this representation.
- Mapping of result back to real world

Syntax [parsing / ast manipulation]

Semantics [type systems]

“New” functionality



# Syntax

Traditional approach: context free grammars

Syntax structure is fixed

Now in practice:

Context-dependent keywords (C++ )

Multi-language programming

## Extensible syntax

Seed 7: Syntax definitions of operations with priority

**syntax** expr: .(). + .() is -> 7;

**syntax** expr: .loop().until().do().end.loop is -> 25;

Nemerle: macroses

macro for (init, cond, change, body)

syntax ("for", "(", init, ":", cond, ":", change, ")", body)

XL      Compiler plugins

## Extensible syntax

*// Idris*

```
boolCase : (x:Bool) -> Lazy a -> Lazy a -> a;
```

```
boolCase True t e = t;
```

```
boolCase False t e = e;
```

```
syntax "if" [test] "then" [t] "else" [e] = boolCase test t e;
```

```
syntax "for" {x} "in" [xs] ":" [body] = forLoop xs (\x => body)
```

*// introducing new binding*

// DSL — map object to Idris language constructions



XL: Concept programming

Active library = Library + compiler plugin

Concept in real life. In computer - representation

Metrics:

Syntax noise / signal    what part of code is about business

Bandwidth    what part of concept is represented

Complexity source

Articulate programming

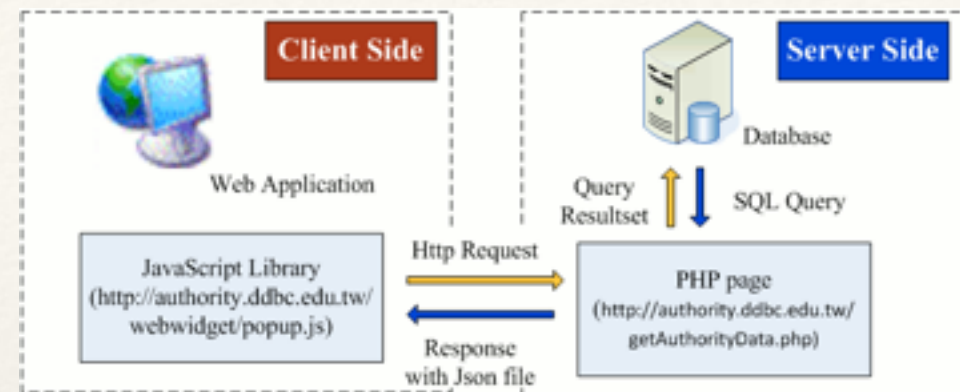
Avail: SmallTalk-based, static typing

Try to deduce 'needed' value from context.

like scala implicit, but \*all\* is implicit

# Syntax

## Multi-language programming



- ❖ Do nothing (mess)
- ❖ Embed all in one language [lifted evaluation]
  - ❖ (like LINQ or scala slick or idris DSL)
- ❖ Define and use 'part of grammar' with type.



Wyvern: <https://www.cs.cmu.edu/~aldrich/wyvern/>

*//Type-specific tokens*

```
let imageBase : URL = <images.example.com>
let bgImage : URL = <%imageBase%/background.png>
new : SearchServer
  def resultsFor(searchQuery, page)
    serve(~) (* serve : HTML -> Unit *)
    >html
      >head
        >title Search Results
      >style ~
        body { background-image: url(%bgImage%) }
        #search { background-color: ..
```

Wyvern: <https://www.cs.cmu.edu/~aldrich/wyvern/>

*//Type-specific tokens*

```
let imageBase : URL = <images.example.com>
```

```
let bgImage : URL = <%imageBase%/background.png>
```

```
new : SearchServer
```

```
  def resultsFor(searchQuery, page)
```

```
    serve(~) (* serve : HTML -> Unit *)
```

```
    >html
```

```
      >head
```

```
        >title Search Results
```

```
      >style ~
```

```
        body { background-image: url(%bgImage%) }
```

```
        #search { background-color: ..
```

Wyvern: <https://www.cs.cmu.edu/~aldrich/wyvern/>

*//Type-specific tokens*

```
let imageBase : URL = <images.example.com>
```

```
let bgImage : URL = <%imageBase%/background.png>
```

```
new : SearchServer
```

```
def resultsFor(searchQuery, page)
```

```
  serve(~) (* serve : HTML -> Unit *)
```

```
    >html
```

```
      >head
```

```
        >title Search Results
```

```
      >style ~
```

```
        body { background-image: url(%bgImage%) }
```

```
        #search { background-color: ..
```



# Type Systems

Static/Dynamic  $\longrightarrow$  Gradual

Dynamic:

- ❖ clojure  $\rightarrow$  core.typed
- ❖ python  $\rightarrow$  type annotations
- ❖ groovy  $\rightarrow$  type annotations

Static:

- ❖ Dynamic type

```
fromDynamic[T]: D => Option[T]  
toDynamic[T](x: T) => Dynamic
```

needed for reflective typing

Clojure: core / typed.

```
(ann clojure.core / first
  (All [x]
    (Fn [(Option (EmptySeqable x)) -> nil]
      [(NonEmptySeqable x) -> x]
      [(Option (Seqable x)) -> (Option x)])))
)
```

- static empty -> nil
- static non-empty -> not-nil
- unknown -> x or nil,

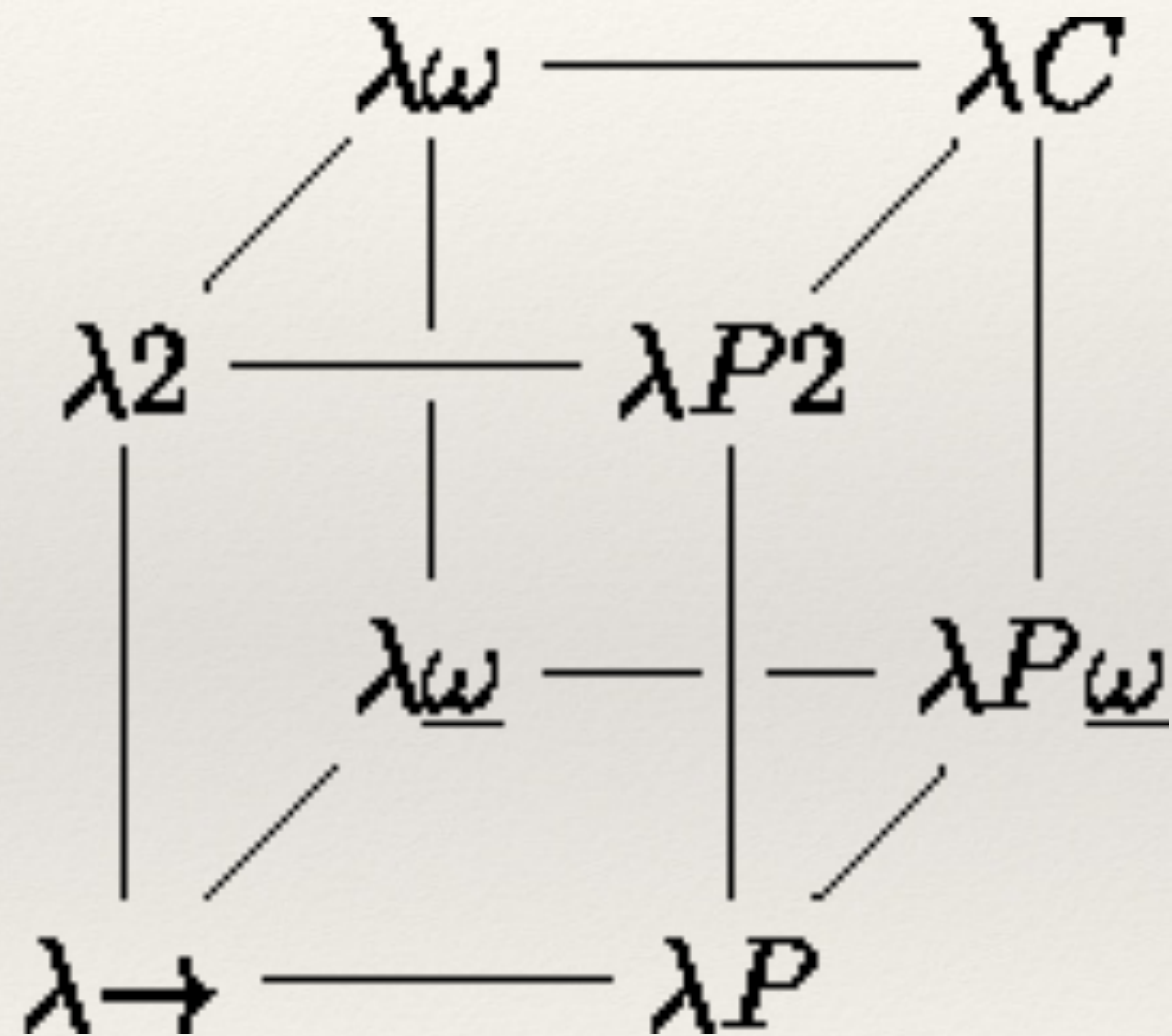
# Type Systems

typing  $\Rightarrow$  static analysis at compile-time

compilation  $\Rightarrow$  optimization of interpretation

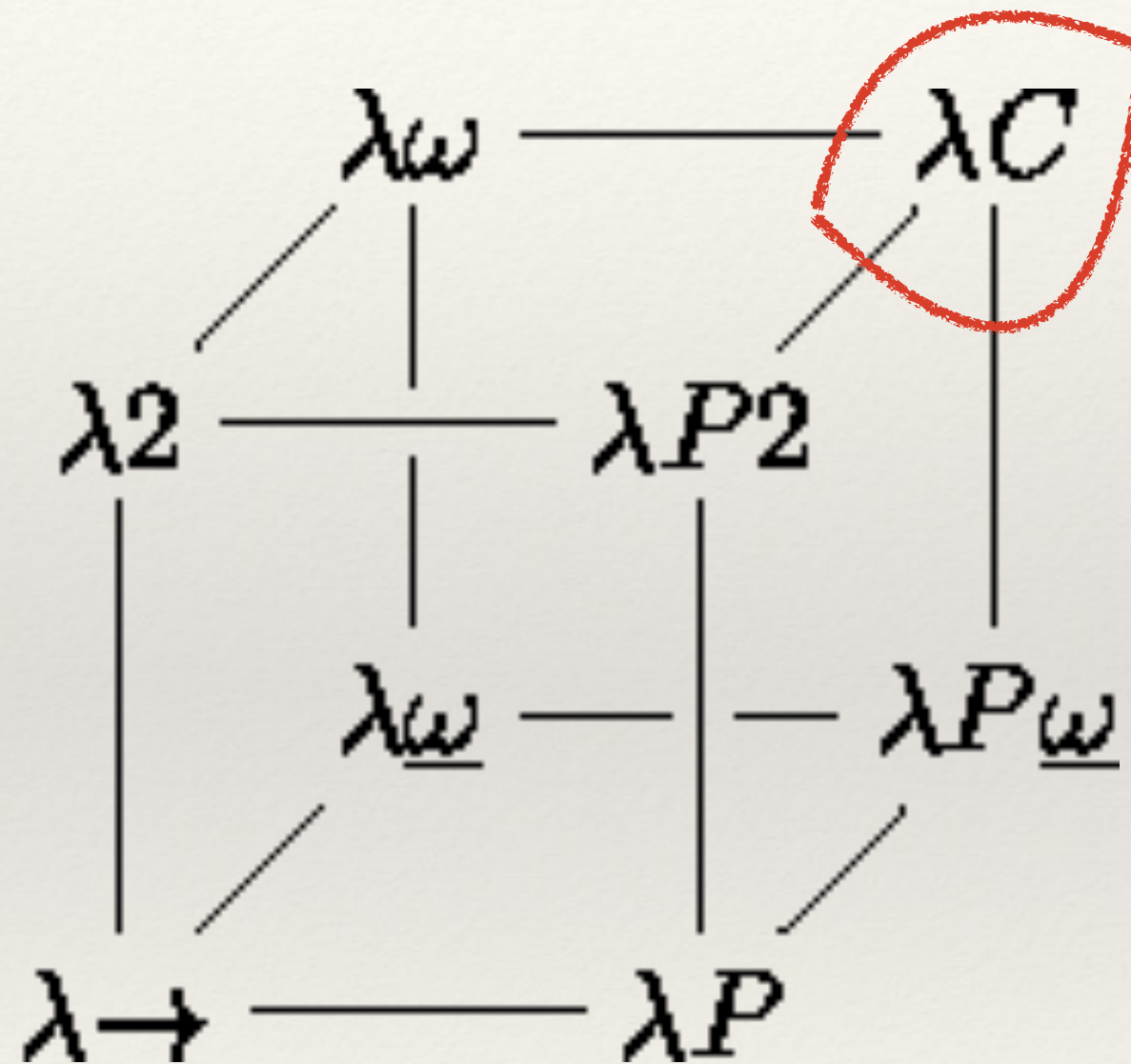


# Type Systems



# Type Systems

Dependent types



## Non-dependent types example

```
trait CInt { type Succ <: CInt }  
trait CPos extends CInt
```

*// church numeration*

```
trait CSucc[P <: CPos] extends CInt {  
  type Succ = CSucc[CSucc[Int]]  
}
```

```
final class _0 extends CPos {  
  type Succ = CSucc[0]  
}
```

```
type _1 = _0#Succ  
type _2 = _1#Succ
```



## Non-dependent types example

```
trait LimitedVector[+A, N<:CInt]
{
  .....

  def append[B <: A, NX <:CInt](x: LimitedVector[B,NX]):
    LimitedVector[B,Plus[N,NX]]

  .....
}

// shapeless full of such tricks
```

## Dependent types example

```
trait LimitedVector[+A, n:Int]
{
    .....

    def append[B <: A, nx:Int](x: LimitedVector[B,nx]):
                                                LimitedVector[B, n+nx]
    .....
}
```

*// warning: pseudocode, impossible in scala*

## Dependent types example

```
data Vect : Nat -> Type -> Type           // Idris
  where
    Nil : Vect 0 a
    (::) : a -> Vect k a -> Vect (k+1) a
```

Calculus of constructions:            Idris, Coq, Agda



Dependent types:

Track properties, which can be proven by structural induction

Function, which call SQL inside no more than 3 times.

Resource with balanced acquire/release for all execution paths.

Type depends from value  $\iff$  Type must be available as value

```
data Parity : Nat -> Type where
  even : Parity (n + n)
  odd  : Parity (S (n + n))
```

Dependent types:

intuitionistic logic

absence ( $p \mid \mid \sim p$ )

Type:Type - logical bomb

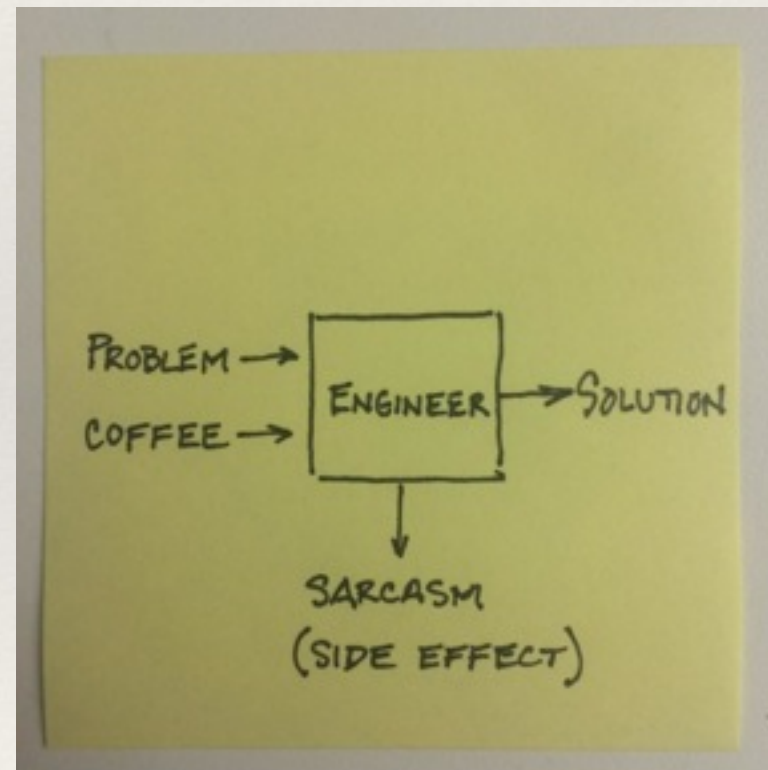
## Effects tracking:



No more IO monads !!!!

List of effects as part of type.

- Idris
- Eff





# Knowledge-based programming

Let's add structured information about problem domain.

Wolfram Alpha:

Mathematics domain. Near 5000 built-in functions.

Household lexicon: approx 4000

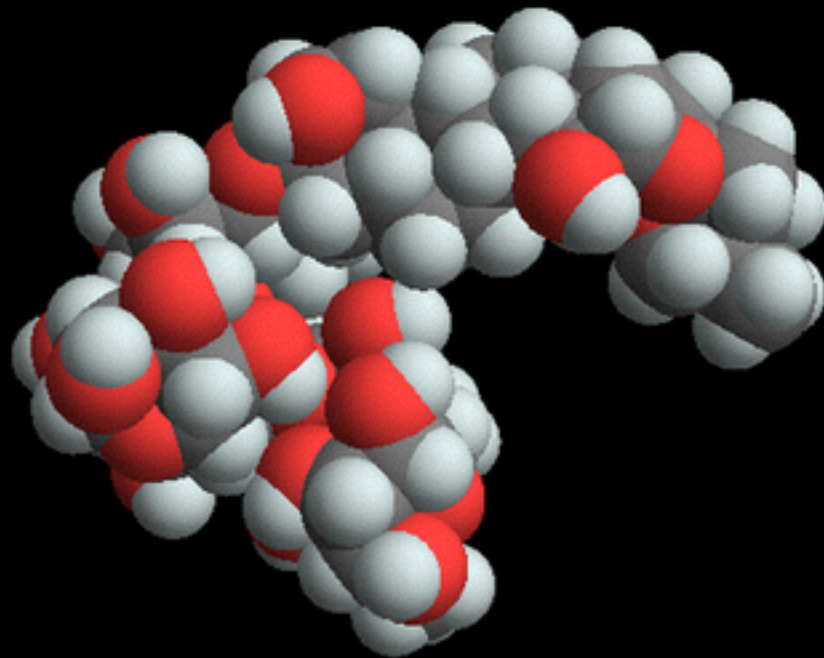
Differ from community repository: all interfaces are consistent

Wolfram:

Tweet-a-program



```
Show[ChemicalData["Digitonin", "SpaceFillingMoleculePlot"],  
Background -> Black]
```



Wolfram Language | Tweet-a-Program @wolframtap



# Wolfram:

## Tweet-a-program



```
ImageCollage[Colorize[#, ColorFunction -> ColorData["Pastel"]]] & /@  
#[RandomSample[#[ "Pokemon", "Entities"], 25], "Image"] & @  
EntityValue]
```



Wolfram Language | Tweet-a-Program @wolframtap

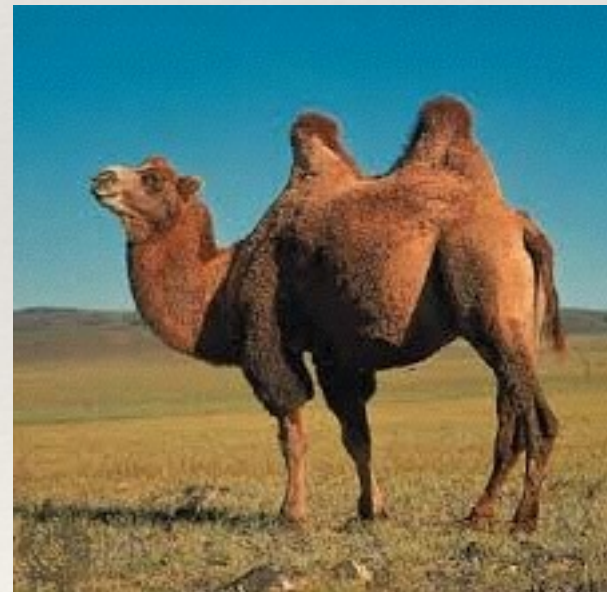




Make all good

- Syntax, defined in problem-domain specific library
- Vocabulary similar to one in natural language
- Advanced type theory with refinement specs.

# Implementation



# Implementation

Compilation/ Interpretation

[both - hard]

VM/ System

both - impossible without divergence

JVM

no continuations (?)

LLVM

non-trivial gc (?)

Complexity.

Let's rewrite all with callbacks, because it's more easy than fix JVM/ OS



CORBA story:

Let's do remote and local call undistinguished (1996)

Local calls: fast

Remote calls: slow

Common antipattern: build fine-grained object model in RPC

Akka:

Let's do remote and local call undistinguished (2014)

*// hope, now things will move forward more accurate*

So, we must

choose one level of abstraction

or

build family of systems



otherwise — abstraction leaks



Ok, assume we build ideal language. (in academia or as freaks)

When it will be in industry ?



- ❖ Success in sense of adoption => evolution slowdown

Haskell:

- ❖ avoid success of all costs => be open for evolution,  
live in a water

Scala: from research to engineering

(LMS, attempt to remove skolem types, ... etc )



# Sociology of programming languages

Leo Meyerovich

Average 'incubation' period before wider adoption: 12 years

before - specialized in some niche

Adoption => social, marketing ....., not technical.

What will be in industry tomorrow: landscape now.

General purpose: need some time for current adoption  
JVM land: scala, clojure

Trends: integrating other paradigms

System: D / Rust instead C++ (?)

Application/ Scripting: [ place is vacant ]  
// JavaScript successor [?]

What will be **not in industry** tomorrow:

academic/OSS

Better VM-s

New paradigms



Depends from you

// niche as social issue



# Thanks.

[ruslan@shevchenko.kiev.ua](mailto:ruslan@shevchenko.kiev.ua)

<https://github.com/rssh>

@rssh1

If you build something interesting - ping me.