Sorting and Searching

ACS-1904 LECTURE 10

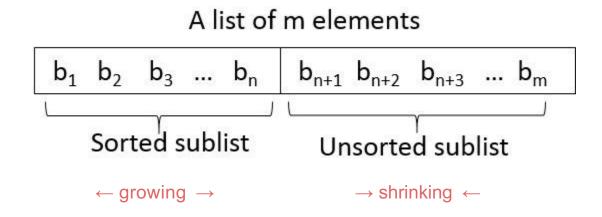
Sorting & searching

- Sorting:
 - Putting lists into sequence

- Searching:
 - Finding things in a list

Sorting

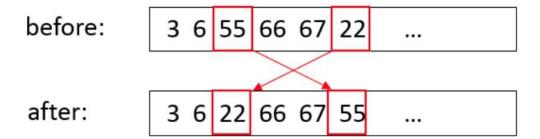
 many algorithms conceptually consider a list that comprises a sorted sublist and an unsorted sublist



Sorting - list operations

Swapping a pair of elements

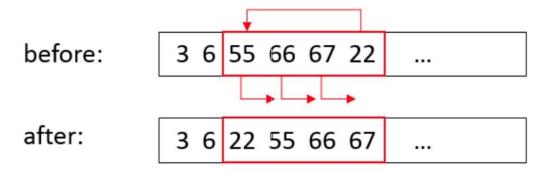
• e.g. 3rd and 6th elements swap positions



Sorting - list operations

Shifting an element into position

 e.g. 6th element moves into 3rd position; others shift right



Sorting algortihms

- Selection sort
- Insertion sort
- Bubble sort
- Quicksort

... many others

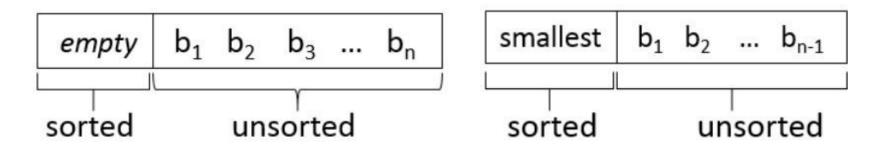
e.g. Merge sort, heap sort, bucket sort, radix sort...

→ *ACS*-2947

During a series of passes over the unsorted sublist

- the sorted sublist grows larger and
- the unsorted sublist shrinks to nothing

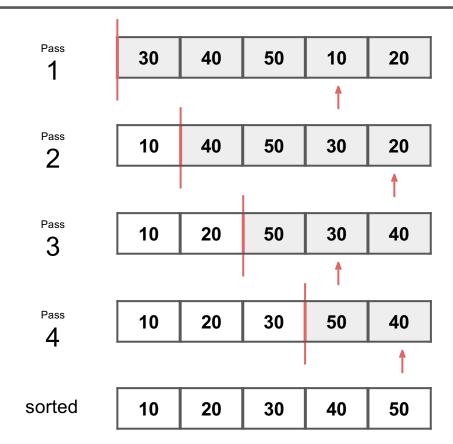
Basic idea: find smallest & move it to beginning of unsorted



- Sorted sublist ← empty.
- 2. Unsorted sublist ← list to be sorted.
- 3. Iteration:
 - a. find the smallest entry (X) in the unsorted sublist
 - b. swap X and the first entry of the unsorted sublist
 - sorted sublist grows by 1 element & unsorted sublist shrinks by 1 element

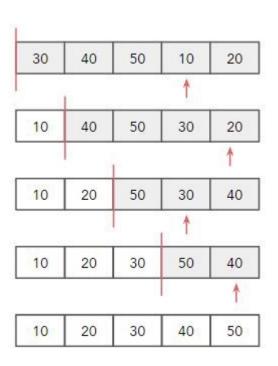
Each pass:

- find smallest of unsorted
- swap with first of unsorted



5 elements → 4 passes

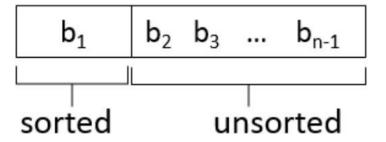
```
public static void selectionSort(int[] a) {
 int n = a.length;
 for (int i=0; i<n-1; i++) {
     int iSmallest = i;
      for (int j=i+1; j<n; j++) {
          if (a[j] < a[iSmallest]) {</pre>
              iSmallest = j;
      if(iSmallest != i) {
          int temp = a[i];
          a[i] = a[iSmallest];
          a[iSmallest] = temp;
```



Insertion sort

Initially first element is part of sorted sublist

Basic idea: move first of the unsorted into proper position of sorted elements ... may require some shifting



Insertion sort

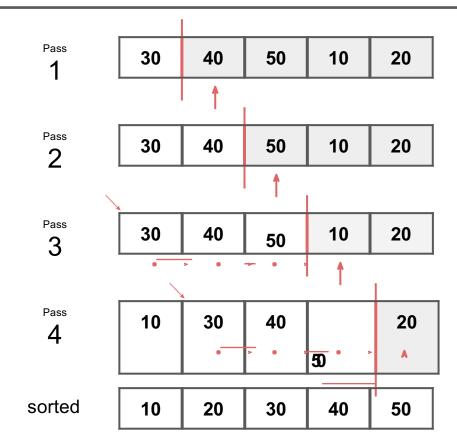
- Sorted sublist ← first element of list.
- 2. Unsorted sublist ← remaining n-1 elements.
- 3. Iteration: Repeat until the unsorted sublist is empty.
 - a. Let X be the first element in the unsorted list
 - Shift entries that are greater than X to the right making room for X in its proper position in the sorted sublist.
 - X is now part of the sorted sublist and the unsorted sublist is 1
 element shorter

```
public static void insertionSort(int[] a) {
 int n = a.length;
 for (int i=1; i<n; i++) {
     int x = a[i];
     int j=i-1;
                                             30
                                                 40
     while (j>=0 \&\& a[j]>x) {
          a[j+1] = a[j];
          j--;
     a[++j]=x;
                                                 20
```

Insertion sort

Each pass:

- move first unsorted element to its proper position in sorted sublist
- shift elements as is necessary



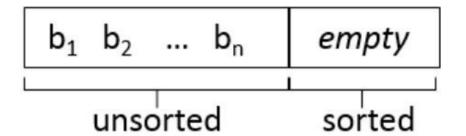
5 elements → 4 passes

Bubble sort

Larger elements bubble to the right

- largest element moved to end of unsorted
- sorted sublist grows by 1, unsorted shrinks by 1

Initially:

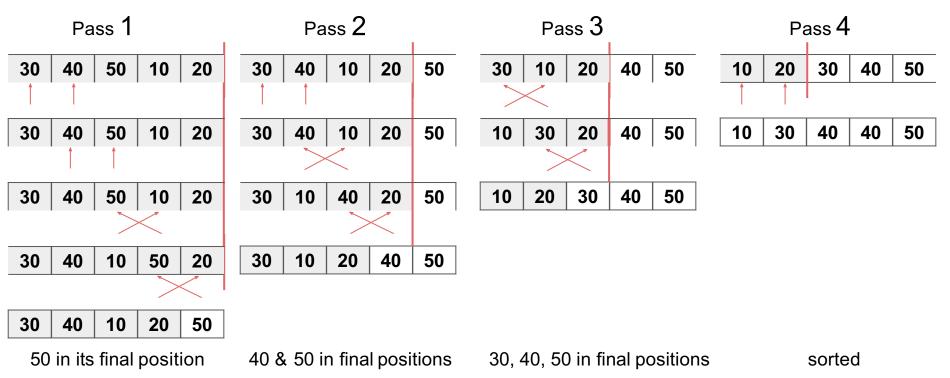


Bubble sort

- Unsorted sublist ← the original list.
- Sorted sublist ← an empty list.
- 3. Iteration:
 - Scanning forward from the beginning of the unsorted list, consider the elements pairwise: a_1 and a_2 , then a_2 and a_3 , ... For each pair where $a_i > a_{i+1}$, swap a_i and a_{i+1} . After a scan the largest element has moved to the end of the unsorted list.
 - The unsorted sublist shrinks by 1 element; sorted sublist increases by 1 element.

Bubble sort

Each iteration: bubble larger elements to the right



Bubble sort Bubble Sort.java

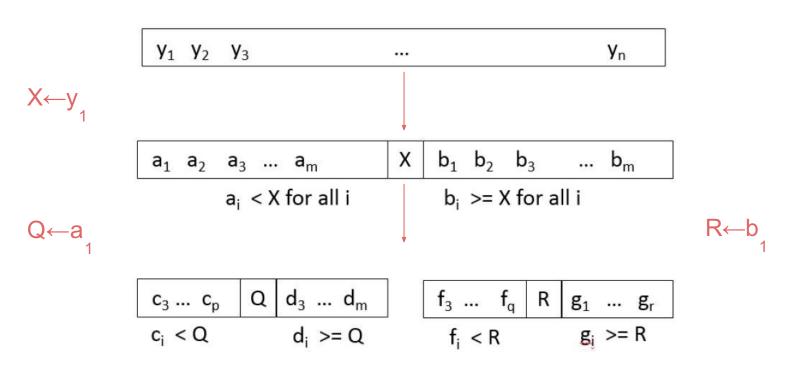
```
public static void bubbleSort(int[] a) {
  int n = a.length;
  for (int i=0; i<n-1; i++) {
       for (int j=0; j< n-i-1; j++) {
            if (a[j] > a[j+1]) {
                 int temp = a[j];
                 a[j] = a[j+1];
                 a[j+1] = temp;
                                                  30 40 10 20 50
                                          50 10 20
                                                               30 10 20 40 50
                                                  30 40 10 20 50
                                     30 40 50 10 20
                                                               10 30 20 40 50
                                                               10 20 30 40 50
                                     30 40 50 10 20
                                                  30 10 40 20 50
                                     30 40 10 50 20
                                                  30 10 20 40 50
                                     30 40 10 20 50
```

Different: Recursive

- Continuously subdivide a list into 2 sublists based on a pivot value
- Variations on how the pivot value is chosen
 - → the pivot can be any value in the list, first, last, middle, or any other random index.
- Left sublist has elements with values < pivot
- Right sublist has elements with values >= pivot

- Once the pivot is chosen the list is 'partitioned' in a (you guessed it) partition method
- Partition moves all of the values smaller than the pivot to the left of the pivot and all of the values larger to the left
- Check out these two algorithms for partitioning
 - Lomuto partition scheme
 - Hoare partition scheme

A list of n elements



QuickSort (list):

- 1. If list of length <= 1 return.
- 2. Choose a pivot value X.
- 3. Rearrange the elements of list:

Left sublist L with elements < X;

Right sublist Rwith elements >= X.

4. QuickSort (L) QuickSort (R)

Unsorted \rightarrow	[6 , 5, 9, 0, 11, 8, 4, 10, 2, 1, 12, 7, 3]	
Pivot=6	[3, 5, 1, 0, 2, 4] [0	6] [10, 8, 11, 12, 7, 9]
Pivots 3, 10	[2, 0, 1] [3] [5, 4]	[9, 8, 7] [10] [12, 11]
Pivots 2, 5, 9, 12	[1, 0] [2] [4][5]	[7, 8] [9] [11][12]
Pivots 1, 7	[0] [1]	[7] [8]

Final result sorted \rightarrow [0][1][2][3][4][5][6][7][8][9][10][11][12]

```
public static void quicksort(int[] a, int start, intend) {
 if (start >= end) return;
                                                        Quicksort.java
 int pivot = a[start];
                                                        QuickSortAlternate.java
 int left = start;
 int right = end;
 while (left < right) {</pre>
     while (right > left && a[right]>=pivot) {
          right--;
     if(left==right) break;
     a[left] = a[right];
     left++;
     while (right > left && a[left] <= pivot) {</pre>
          left++;
     if(left == right) break;
     a[right] = a[left];
     right--;
 a[right]=pivot;
 quicksort(a, start, right-1);
 quicksort(a, right+1, end );
```

Sorting Objects

Example

- Objects (ducks) must implement Comparable i.e. compareTo method
- Array passed in to sort method is of type Comparable
- Object holding a temporary is of type Comparable

Searching

Unordered list?

Ordered list?

i.e. sorted

Searching - unordered list

To determine that a search value is not present must scan the whole array

Searching an **Unordered** List

Search(list, X):

- 1. Let *i* reference the start of the list
- 2. While $i \le end$ of list
 - a. If $list_i$ equals X then return found at position i
 - b. Increment *i* to reference the next element of the list
- 3. Return not found Searched to the end of list

Searching - ordered list (Sequential Search)

Searching an Ordered List

SequentialSearch (list, *X*):

- 1. Let *i* reference the first element of the list.
- 2. While $i \le end$ of list
 - a. If $list_i > X$ then return not found

Don't always need to go to end of list to determine the search value is not there

- b. If $list_i$ equals X then return found at position i
- c. Increment i to reference the next element of the list
- 3. Return not found

Searching – ordered list (Binary Search)

Searching an Ordered List

BinarySearch (list, X):

- 1. If the list *L* is empty then return *not found*
- 2. Determine the midpoint of the list
- 3. Determine the value at the midpoint of the list
- 4. if X < midpoint value call BinarySearch (left sublist, X) else if X > midpoint value call BinarySearch (right sublist, X)) else if X = midpoint value then return found at position midpoint