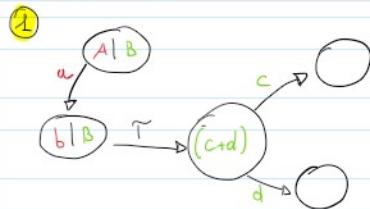
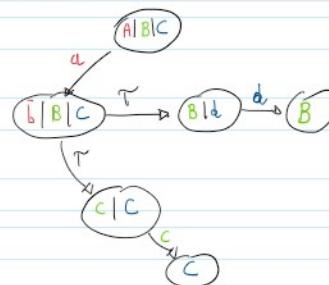


$$P = \frac{(a \cdot \bar{b} \mid b \cdot (c+d))}{A \quad B}$$



$$Q = \frac{(a \cdot \bar{b} \mid b \cdot c \mid b \cdot d)}{A \quad B \quad C}$$



③

$$P = \frac{P_0 \xrightarrow{a} P_1 \xrightarrow{\tau} P_2 \xrightarrow{c} P_3 \xrightarrow{a} P_4}{L(P) = a \cdot \tau (c + a)}$$

④

$$Q = \frac{q_0 \xrightarrow{a} q_1 \xrightarrow{\tau} q_2 \xrightarrow{a} q_3 \xrightarrow{\tau} q_4 \xrightarrow{c} q_5}{L(Q) = a \cdot (\tau \cdot c) + (\tau \cdot a)}$$

The languages are equivalent  $L(P) = L(Q)$

They aren't bisimilar:

$q_0$  is simulated by  $p_0$ :

$$S = \{(q_0, p_0), (q_1, p_0), (q_2, p_0), (q_3, p_0), (q_4, p_0), (q_5, p_0)\}$$

$p_0$  is not simulated by  $q_0$ :

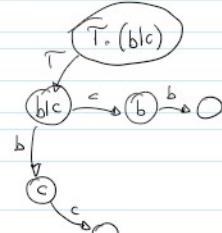
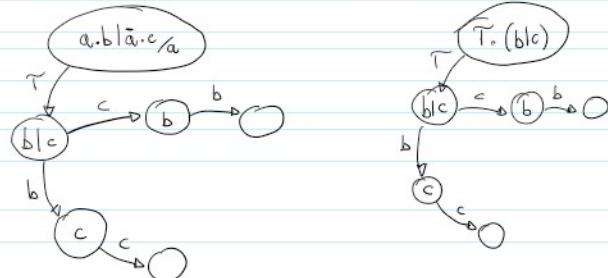
$$\tilde{S}^1 = \{(p_0, q_1), (p_2, q_1), (p_2, q_2), (p_2, q_3), (p_4, q_4), (p_3, q_5)\}$$

IMPOSSIBLE

~~~~~

⑤

$$(a \cdot b \mid \bar{a} \cdot c) / a$$



Relabelling of the states:

$$S = \{(p_i, q_j) \mid i \in [0, 5]\} \quad \left. \begin{array}{l} \\ \end{array} \right\} \text{BISIMULATION}$$

$$\tilde{S}^1 = \{(q_i, p_i) \mid i \in [0, 5]\}$$

## 3) WITHOUT PARAMETRIC PROCESS DEFINITION

Stack of 2 bits:

$$S_E = \sum_{i \in \{0,1\}} \text{inv}_i \cdot S_i$$

$$S_i = \sum_{j \in \{0,1\}} \text{inv}_j \cdot S_{ij} + \text{out}_i \cdot B_E$$

$$S_{ij} = \text{out}_j \cdot S_i$$

## WITH PARAMETRIC PROCESS DEFINITION

Stack of 2 bits:

$$S_E = \sum_{i \in \{0,1\}} \text{inv}_i \cdot S'(out_i)$$

$$S'(x) = \sum_{j \in \{0,1\}} \text{inv}_j \cdot S''(x, out_j) + x \cdot S_E$$

$$S''(x, y) = y \cdot S'(x)$$

Stack of 3 bits:

$$S_E = \sum_{i \in \{0,1,2\}} \text{inv}_i \cdot S'(out_i)$$

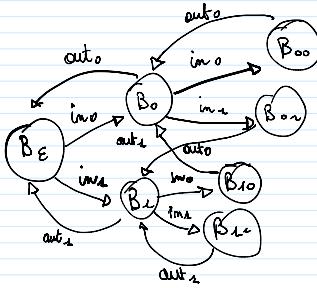
$$S'(x) = \sum_{j \in \{0,1,2\}} \text{inv}_j \cdot S''(x, out_j) + x \cdot S_E$$

$$S''(x, y) = \sum_{k \in \{0,1,2\}} \text{inv}_k \cdot S'''(x, y, out_k) + y \cdot S'(x)$$

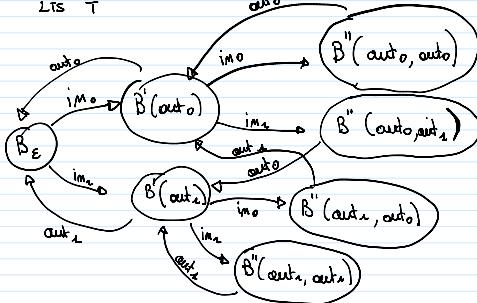
$$S'''(x, y, z) = z \cdot S''(x, y)$$

## 4)

LTS T



LTS T'



## WITHOUT PARAMETRIC DEFINITIONS

$$\text{Let } T = (S, A, \rightarrow, B_E)$$

There exist  $f$  bijective :  $f: S \rightarrow S'$  s.t.  $f(B_E) = B_E$  and  $s_1 \rightarrow s_2$  for some  $s_1, s_2 \in S$  i.i.f.  $f(s_1) \rightarrow' f(s_2)$

## WITH PARAMETRIC DEFINITIONS

$$\text{Let } T' = (S', A, \rightarrow', B_E)$$