

COMPACTNESS

2) Let T_1 and T_2 theories.

for each structure \mathcal{A} : $\mathcal{A} \models T_1 \iff \mathcal{A} \models T_2 \implies T_1, T_2$ are finitely axiomatizable

④ HYPOTHESIS

PROOF: \implies (By way of contradiction)

Assume the contrary T_1 is not finitely axiomatizable, T_1 is not equivalent to any of its subtheories T_1' . Then for every subtheory T_1' there's some some structure \mathcal{B} s.t. $\mathcal{B} \models T_1$ and $\mathcal{B} \not\models T_1'$. By ① $\mathcal{B} \models T_2$, so $\mathcal{B} \models (T_1' \cup T_2)$, now let $P = T_1' \cup T_2$, every finite subtheories of P is contained in $T_1' \cup T_2$ for some $T_1' \subset T_1$.

By COMPACTNESS, P has model $\mathcal{C} \implies \mathcal{C} \models T_1$ and $\mathcal{C} \models T_2$ (CONTRADICTION)

5) Prove by a Compactness argument that the notion of being a nice relation is not axiomatizable

PROOF:

Suppose we've T s.t. $\mathcal{A} \models T \iff \mathcal{A}$ admits nice

Let a sentence $A_m := \exists x_1 \dots \exists x_m (\bigwedge_{i \neq j} x_i \neq x_j \wedge (x_m < x_{m-1} < \dots < x_1))$ informally it states: "there exist a finite decreasing sequence of length m "

Let $T' = T \cup \{A_m \mid m \in \mathbb{N}\}$

$\mathcal{A} \models T' \iff \begin{cases} \mathcal{A} \models T, \mathcal{A} \text{ has no infinite sequence} \\ \text{for all } m \in \mathbb{N} \mathcal{A} \models A_m, \mathcal{A} \text{ has decreasing sequence of length } m \end{cases} \implies \mathcal{A} \text{ has an infinite decreasing sequence}$

CLAIM: T' is SAT

By COMPACTNESS, for every subtheory $X \subseteq T'$ is SAT

X contain some sentences of T union some $A_{m_1} \dots A_{m_q}$

We can conclude that the model of X is a structure with finite decreasing sequence of length $m = \max \{m_1, \dots, m_q\}$

But this brings us to state that T' is SAT, which is an ABSURD.

ULTRAFILTERS

3) Let \mathcal{U} ultrafilter s.t. for no $m \in \mathbb{N}$ $\mathcal{U} = \{X \subseteq \mathbb{N} : m \in X\}$ (\mathcal{U} is not a principal ultrafilter)
Let $\text{Cof}(\mathbb{N}) = \{A \subseteq \mathbb{N} : A \text{ is finite}\}$

$\mathcal{U} \implies \text{Cof}(\mathbb{N}) \subset \mathcal{U}$

PROOF \implies

Let $y \in \mathbb{N}$, since \mathcal{U} is not a principal ultrafilter, the set $\{y\} \notin \mathcal{U}$, $\{y\}^c \in \mathcal{U}$ and $\{y\} \notin \mathcal{U}$ (MAXIMALITY)
Thus for all $y \in \mathbb{N}$. Now we can define a subset H of \mathbb{N}

$$H^c = \bigcap_{y \in \mathbb{N}} \{y\}^c$$

Every filter is closed under finite intersections, then $H^c \in \mathcal{U}$

$\text{Cof}(\mathbb{N}) = \{H \subseteq \mathbb{N} : H \text{ is finite}\}$

then $\text{Cof}(\mathbb{N}) \subset \mathcal{U}$

5) \mathcal{U} ultrafilter contains a finite set not empty (since \mathcal{U} is an ultrafilter) ① HYP
 \implies for some $m \in \mathbb{N}$ s.t. $\mathcal{U} = \{X \subseteq \mathbb{N} : m \in X\}$

PROOF:

By ① pick $A = \bigcap \mathcal{U}$ since \mathcal{U} is closed under finite intersection, A is a finite set and it's min of \mathcal{U}

1) A must not be equal \emptyset

2) We want also that A is a singleton and is equal to $\{m\}$

By way of contradiction: A isn't a singleton $\implies A = B \cup C$

Since A is the minimum of \mathcal{U} (see ②) neither B nor $C \in \mathcal{U}$ ③

Now using MAXIMALITY $B^c \in \mathcal{U}$, using the closed intersection $A \cap B^c = C \in \mathcal{U}$, by ③ is a CONTRADICTION!

then $A = \{m\}$, $\mathcal{U} = \{X \subseteq \mathbb{N} : \{m\} \subseteq X\} = \{X \subseteq \mathbb{N} : m \in X\}$