Cover Times of Random Walks

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Overview

Introduction

Expected Cover Time of a Small Graph

Expected Cover Time of a Complete Graph with Self-edges

Expected Cover Time of a General Graph

Cover Time Distribution of the a Complete Graph with Self-edges

Overview

Introduction

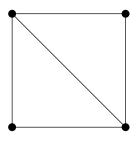
Expected Cover Time of a Small Graph

Expected Cover Time of a Complete Graph with Self-edges

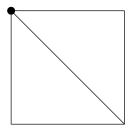
Expected Cover Time of a General Graph

Cover Time Distribution of the a Complete Graph with Self-edges

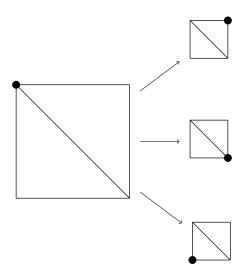
A Graph



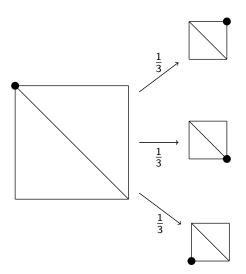
Random Walks



Random Walks



Random Walks



Cover Times

How many steps does a random walk take to visit every vertex?

Coupon Collector

A collector receives one of n possible coupons in the mail everyday. How long until all n coupons have been collected?

How many people do we need for each birthday (excluding February 29) to be represented?

Overview

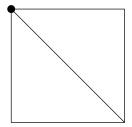
Introduction

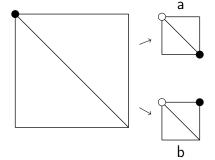
Expected Cover Time of a Small Graph

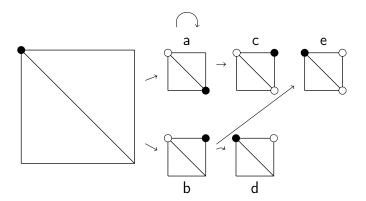
Expected Cover Time of a Complete Graph with Self-edges

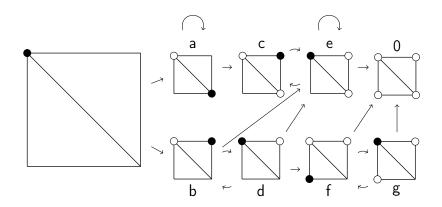
Expected Cover Time of a General Graph

Cover Time Distribution of the a Complete Graph with Self-edges









$$\begin{split} \mathrm{E}[\mathsf{t}_{\mathsf{cov}}] &= 1 + \frac{1}{3} \mathrm{a} + \frac{2}{3} \mathrm{b} \\ \mathrm{a} &= 1 + \frac{1}{3} \mathrm{a} + \frac{2}{3} \mathrm{c} \\ \mathrm{b} &= 1 + \frac{1}{2} \mathrm{b} + \frac{1}{2} \mathrm{e} \\ \mathrm{c} &= 1 + 1 \mathrm{e} \\ \mathrm{d} &= 1 + \frac{1}{3} \mathrm{a} + \frac{1}{3} \mathrm{e} + \frac{1}{3} \mathrm{f} \\ \mathrm{e} &= 1 + \frac{1}{3} \mathrm{c} + \frac{1}{3} \mathrm{e} + \frac{1}{3} \mathrm{0} \\ \mathrm{f} &= 1 + \frac{1}{2} \mathrm{g} + \frac{1}{2} \mathrm{0} \\ \mathrm{g} &= 1 + \frac{2}{3} \mathrm{f} + \frac{1}{3} \mathrm{0} \end{split}$$

$$\begin{split} \mathrm{E}[t_{cov}] &= 1 + \frac{1}{3} \mathrm{a} + \frac{2}{3} \mathrm{b} = \frac{43}{6} \\ \mathrm{a} &= 1 + \frac{1}{3} \mathrm{a} + \frac{2}{3} \mathrm{c} \\ \mathrm{b} &= 1 + \frac{1}{2} \mathrm{b} + \frac{1}{2} \mathrm{e} \\ \mathrm{c} &= 1 + 1 \mathrm{e} \\ \mathrm{d} &= 1 + \frac{1}{3} \mathrm{a} + \frac{1}{3} \mathrm{e} + \frac{1}{3} \mathrm{f} \\ \mathrm{e} &= 1 + \frac{1}{3} \mathrm{c} + \frac{1}{3} \mathrm{e} + \frac{1}{3} \mathrm{0} \\ \mathrm{f} &= 1 + \frac{1}{2} \mathrm{g} + \frac{1}{2} \mathrm{0} \\ \mathrm{g} &= 1 + \frac{2}{3} \mathrm{f} + \frac{1}{3} \mathrm{0} \end{split}$$

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Cover Time Distribution of the a Complete Graph with Self-edges

$$\mathrm{E}[t_{\mathsf{cov}}] = \mathrm{E}[X_1 + X_2 + X_3 + \dots + X_n]$$

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$$X_i \sim \operatorname{Geom}(\frac{n-i+1}{n})$$
 so $\operatorname{E}[X_i] = \frac{n}{n-i+1}$ and

$$\mathrm{E}[\mathsf{t}_{\mathsf{cov}}] = n \left(\frac{1}{n-1} + \frac{1}{n-2} + \dots + 1 \right)$$

Coupon Collector

A collector receives one of n possible coupons in the mail everyday. How long until all coupons have been collected?

On average, we need about $n \log n$ days.

How many people do we need for each birthday (excluding February 29) to be represented?

On average, we need about 2365 people.

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Matthews Method

Theorem

Consider a random walk on an arbitrary connected graph with n vertices. Then

$$\mathrm{E}[\mathsf{t}_{\mathsf{cov}}] \leq t_{\mathit{hit}} \left(1 + rac{1}{2} + rac{1}{3} + \cdots + rac{1}{n-1}
ight).$$

Let σ be a uniform permutation of unvisited vertices from (worst-case) vertex n.

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$$\begin{split} \mathrm{E}[\mathsf{t}_{\mathsf{cov}}] &= \mathrm{E}_n[\, T_{n-1}] \\ &\leq \mathrm{E}_n[\, T_1 + (\, T_2 - \, T_1) + \dots + (\, T_{n-1} - \, T_{n-2})] \end{split}$$

$$\mathrm{E}_n[T_k - T_{k-1}] = \sum_{i=1}^k \mathrm{E}_n[T_k - T_{k-1}|L(k) = \sigma_i] \, \mathrm{P}(L_k = \sigma_i)$$

$$E_{n}[T_{k} - T_{k-1}] = \sum_{i=1}^{k} E_{n}[T_{k} - T_{k-1}|L(k) = \sigma_{i}] P(L_{k} = \sigma_{i})$$

$$= 0 + E_{n}[T_{k} - T_{k-1}|L(k) = \sigma_{k}] P(L_{k} = \sigma_{k})$$

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$$\leq t_{hit} \frac{1}{k}$$

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After r steps of a random walk, let A_i be the event that vertex i is unvisited.

$$P(A_i) = P(\text{not visited in one step})^r = (1 - 1/n)^r$$

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$$P(A_i) = P(\text{not visited in one step})^r = (1 - 1/n)^r$$

Let N_n be the number of unvisited vertices.

$$E[N_n] = \sum_{i=1}^n P(A_i) = n(1 - 1/n)^r$$
$$= n[(1 - 1/n)^n]^{r/n} \sim ne^{-r/n}$$

Introduce x and set $r = n \log n + nx$.

$$N_n \sim ne^{-r/n} = ne^{-\log n - x} = nn^{-1}e^{-x} \to e^{-x}$$

Theorem

If $ne^{-r/n} \to \lambda \in [0, \infty)$, then the number of unvisited vertices approaches a Poisson distribution with mean λ .

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$$N_n \sim ne^{-r/n} = ne^{-\log n - x} = nn^{-1}e^{-x} \to e^{-x}$$

Theorem

If $ne^{-r/n} \to \lambda \in [0, \infty)$, then the number of unvisited vertices approaches a Poisson distribution with mean λ .

$$P(N_n = 0) = P(t_{cov} \le r) = P(t_{cov} - n \log n \le nx) \rightarrow e^{-e^{-x}}$$



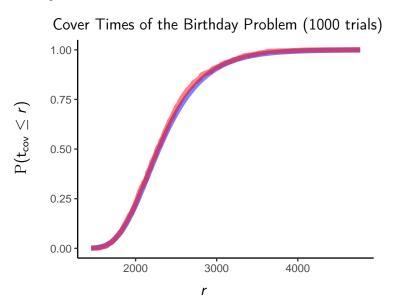
What is the probability all birthdays are covered in a town of 1825? How about a town of 2190?

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$$\begin{split} \mathrm{P}(t_{\mathsf{cov}} \leq 1825) &= \mathrm{P}\left(\frac{t_{\mathsf{cov}} - 2153}{365} \leq \frac{-328}{365}\right) \\ &\approx e^{-e^{-0.89863}} \approx 0.08576 \end{split}$$

What is the probability all birthdays are covered in a town of 1825? How about a town of 2190?

$$\begin{split} \mathrm{P}(\mathsf{t}_{\mathsf{cov}} \leq 1825) &= \mathrm{P}\left(\frac{\mathsf{t}_{\mathsf{cov}} - 2153}{365} \leq \frac{-328}{365}\right) \\ &\approx e^{-e^{-0.89863}} \approx 0.08576 \\ \mathrm{P}(\mathsf{t}_{\mathsf{cov}} \leq 2190) &= \mathrm{P}\left(\frac{\mathsf{t}_{\mathsf{cov}} - 2153}{365} \leq \frac{37}{365}\right) \\ &\approx e^{-e^{-0.10137}} \approx 0.40511 \end{split}$$



Thank you!

github.com/rtealw/Cover-Times

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