Cover Times of Random Walks

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Overview

Introduction

Expected Cover Time of a Small Graph

Expected Cover Time of a Complete Graph with Self-edges

Expected Cover Time of a General Graph

Cover Time Distribution of the a Complete Graph with Self-edges

Overview

Introduction

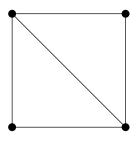
Expected Cover Time of a Small Graph

Expected Cover Time of a Complete Graph with Self-edges

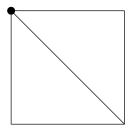
Expected Cover Time of a General Graph

Cover Time Distribution of the a Complete Graph with Self-edges

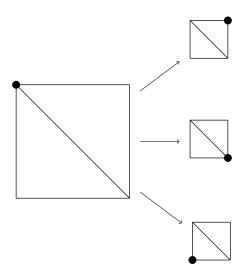
A Graph



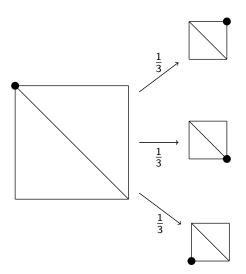
Random Walks



Random Walks



Random Walks



Cover Times

How many steps does a random walk take to visit every vertex?

Coupon Collector

A collector receives one of n possible coupons in the mail everyday. How long until all n coupons have been collected?

How many people do we need for each birthday (excluding February 29) to be represented?

Overview

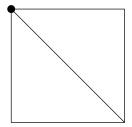
Introduction

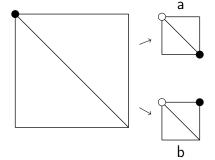
Expected Cover Time of a Small Graph

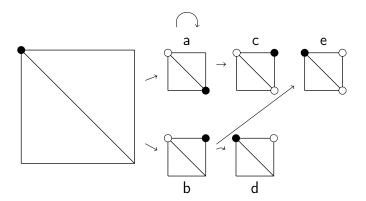
Expected Cover Time of a Complete Graph with Self-edges

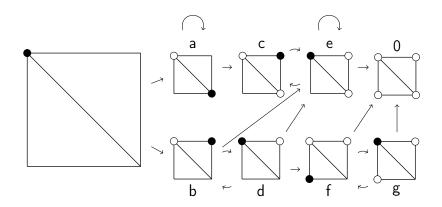
Expected Cover Time of a General Graph

Cover Time Distribution of the a Complete Graph with Self-edges









$$\begin{split} \mathrm{E}[\mathsf{t}_{\mathsf{cov}}] &= 1 + \frac{1}{3} \mathrm{a} + \frac{2}{3} \mathrm{b} \\ \mathrm{a} &= 1 + \frac{1}{3} \mathrm{a} + \frac{2}{3} \mathrm{c} \\ \mathrm{b} &= 1 + \frac{1}{2} \mathrm{b} + \frac{1}{2} \mathrm{e} \\ \mathrm{c} &= 1 + 1 \mathrm{e} \\ \mathrm{d} &= 1 + \frac{1}{3} \mathrm{a} + \frac{1}{3} \mathrm{e} + \frac{1}{3} \mathrm{f} \\ \mathrm{e} &= 1 + \frac{1}{3} \mathrm{c} + \frac{1}{3} \mathrm{e} + \frac{1}{3} \mathrm{0} \\ \mathrm{f} &= 1 + \frac{1}{2} \mathrm{g} + \frac{1}{2} \mathrm{0} \\ \mathrm{g} &= 1 + \frac{2}{3} \mathrm{f} + \frac{1}{3} \mathrm{0} \end{split}$$

$$\begin{split} \mathrm{E}[t_{cov}] &= 1 + \frac{1}{3} \mathrm{a} + \frac{2}{3} \mathrm{b} = \frac{43}{6} \\ \mathrm{a} &= 1 + \frac{1}{3} \mathrm{a} + \frac{2}{3} \mathrm{c} \\ \mathrm{b} &= 1 + \frac{1}{2} \mathrm{b} + \frac{1}{2} \mathrm{e} \\ \mathrm{c} &= 1 + 1 \mathrm{e} \\ \mathrm{d} &= 1 + \frac{1}{3} \mathrm{a} + \frac{1}{3} \mathrm{e} + \frac{1}{3} \mathrm{f} \\ \mathrm{e} &= 1 + \frac{1}{3} \mathrm{c} + \frac{1}{3} \mathrm{e} + \frac{1}{3} \mathrm{0} \\ \mathrm{f} &= 1 + \frac{1}{2} \mathrm{g} + \frac{1}{2} \mathrm{0} \\ \mathrm{g} &= 1 + \frac{2}{3} \mathrm{f} + \frac{1}{3} \mathrm{0} \end{split}$$

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Cover Time Distribution of the a Complete Graph with Self-edges

$$\mathrm{E}[t_{\mathsf{cov}}] = \mathrm{E}[X_1 + X_2 + X_3 + \dots + X_n]$$

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$$X_i \sim \operatorname{Geom}(\frac{n-i+1}{n})$$
 so $\operatorname{E}[X_i] = \frac{n}{n-i+1}$ and

$$E[\mathsf{t}_{\mathsf{cov}}] = n \left(\frac{1}{n} + \frac{1}{n-1} + \dots + 1 \right)$$

Coupon Collector

A collector receives one of n possible coupons in the mail everyday. How long until all coupons have been collected?

On average, we need about $n \log n$ days.

How many people do we need for each birthday (excluding February 29) to be represented?

On average, we need about 2365 people.

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Matthews Method

Theorem

Consider a random walk on an arbitrary connected graph with n vertices. Then

$$\mathrm{E}[\mathsf{t}_{\mathsf{cov}}] \leq t_{\mathit{hit}} \left(1 + rac{1}{2} + rac{1}{3} + \cdots + rac{1}{n-1}
ight).$$

Let σ be a uniform permutation of unvisited vertices from (worst-case) vertex n.

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Let T_k be the first time vertices $\sigma_1, \sigma_2, \dots, \sigma_k$ have all been visited.

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$$E[t_{cov}] \le E_n[T_{n-1}]$$

 $\le E_n[T_1 + (T_2 - T_1) + \dots + (T_{n-1} - T_{n-2})].$

$$\mathrm{E}_{n}[T_{k}-T_{k-1}] = \sum_{i=1}^{k} \mathrm{E}_{n}[T_{k}-T_{k-1}|L(k)=\sigma_{i}] \, \mathrm{P}(L_{k}=\sigma_{i}).$$

$$E_{n}[T_{k} - T_{k-1}] = \sum_{i=1}^{k} E_{n}[T_{k} - T_{k-1}|L(k) = \sigma_{i}] P(L_{k} = \sigma_{i}).$$

$$= 0 + E_{n}[T_{k} - T_{k-1}|L(k) = \sigma_{k}] P(L_{k} = \sigma_{k}).$$

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$$\leq t_{hit} \frac{1}{k}$$

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I promise that

$$N_n/n \xrightarrow{p} \mathrm{E}[N_n/n]$$

$$r = n \log n + nx$$

 $N_n = ne^{-r/n} = ne^{-\log n - x} = nn^{-1}e^{-x} \to e^{-x}$

Theorem

If $ne^{-r/n} \to \lambda \in [0, \infty)$, then the number of unvisited vertices approaches a Poisson distribution with mean λ .

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Theorem

If $ne^{-r/n} \to \lambda \in [0, \infty)$, then the number of unvisited vertices approaches a Poisson distribution with mean λ .

$$P(N_n = 0) = P(t_{cov} \le r) \to e^{-e^{-x}}$$
$$P(t_{cov} - n \log n \le nx) \to e^{-e^{-x}}$$

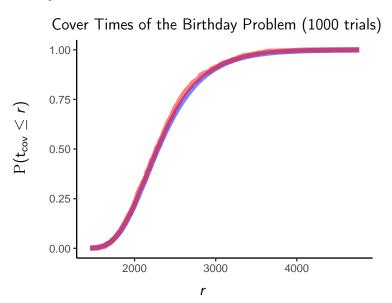
What is the probability all birthdays are covered in a town of 1825? How about a town of 2190?

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$$\begin{split} \mathrm{P}(t_{\mathsf{cov}} \leq 1825) &= \mathrm{P}\left(\frac{t_{\mathsf{cov}} - 2153}{365} \leq \frac{-328}{365}\right) \\ &\approx e^{-e^{-0.89863}} \approx 0.08576 \end{split}$$

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Thank you!

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