

Applications of the Quantum st-Connectivity Algorithm

June 3, 2019

University of Maryland

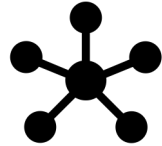
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¹Middlebury College

Layers of Abstraction

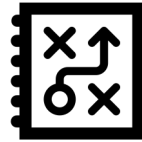


Our Algorithms



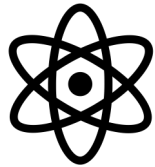
st-Connectivity

[Belovs, Reichardt, '12]



Span Programs

[Reichardt, '09, '11]



Quantum Gates

Recipe

Step 1: Encode a question into a graph structure.

Step 2: Analyze worst case effective resistance/capacitance.

Step 2b: Look for characteristics of original graph hidden in graph reduction.

Big Question

st-connectivity reduces quantum algorithm design to a simple classical algorithm

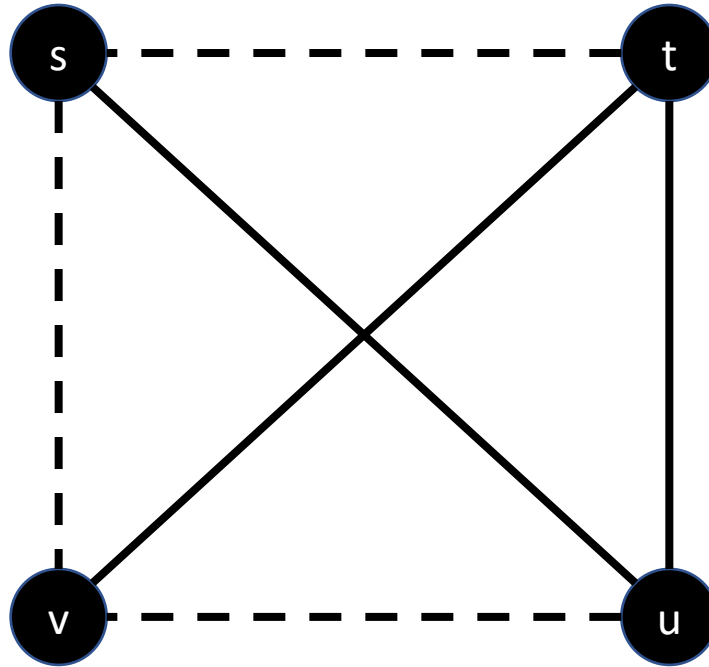
st-connectivity feels 'natural' in the following ways:

- it's optimal for a wide range of problems (e.g. Boolean formula evaluation, total connectivity) and
- it's easy to analyze using effective resistance/capacitance

Does the st-connectivity approach give intuition for optimal underlying quantum algorithms?

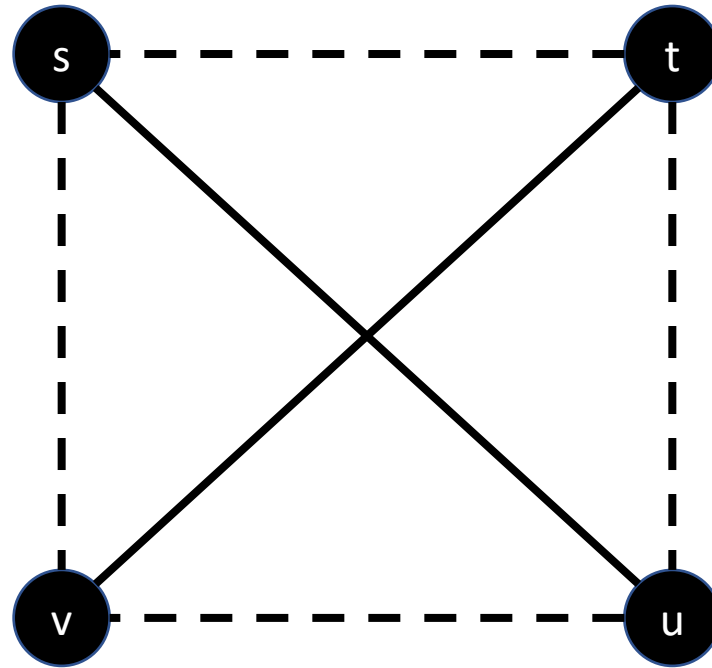
st-Connectivity

Is there a path
between s and t?



st-Connectivity

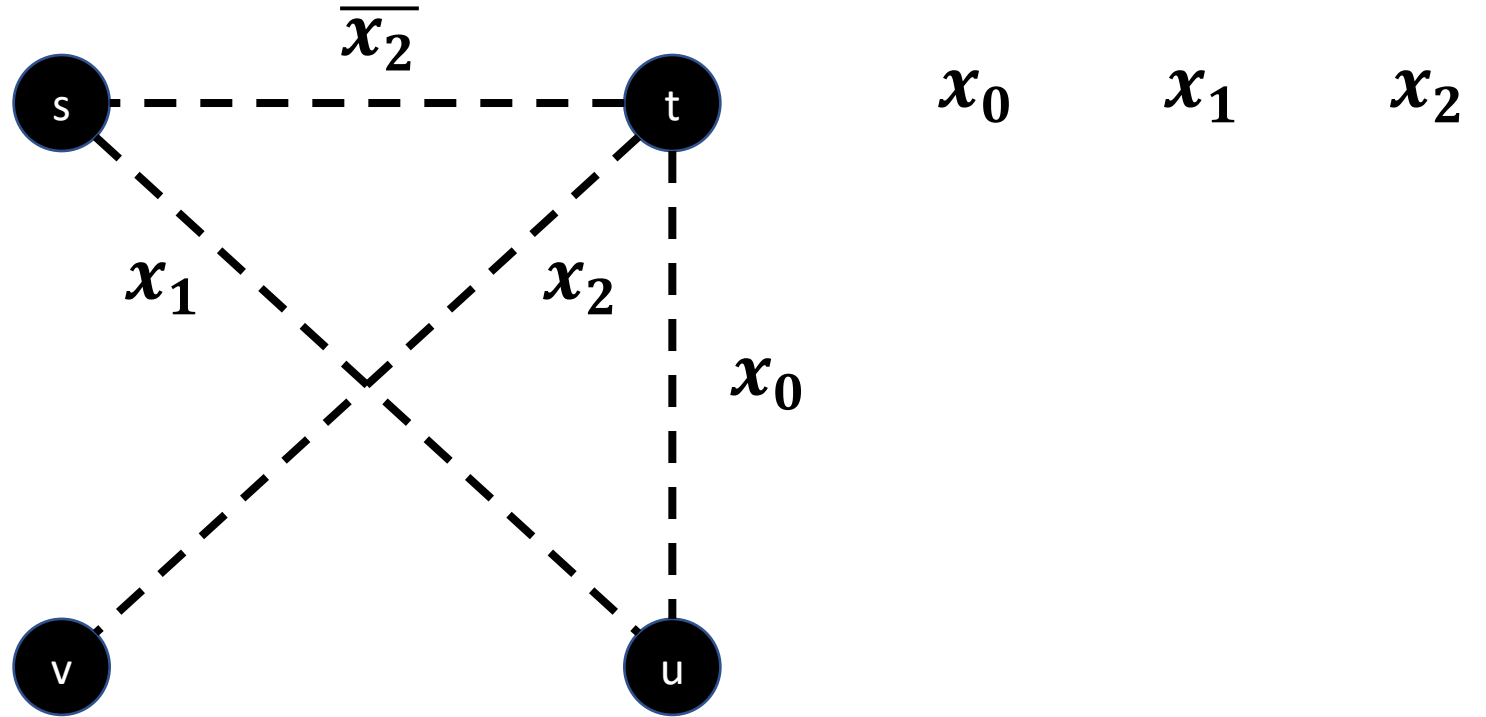
Is there a path
between s and t?



st-Connectivity Query Model

Input:

- graph skeleton
- hidden bit string

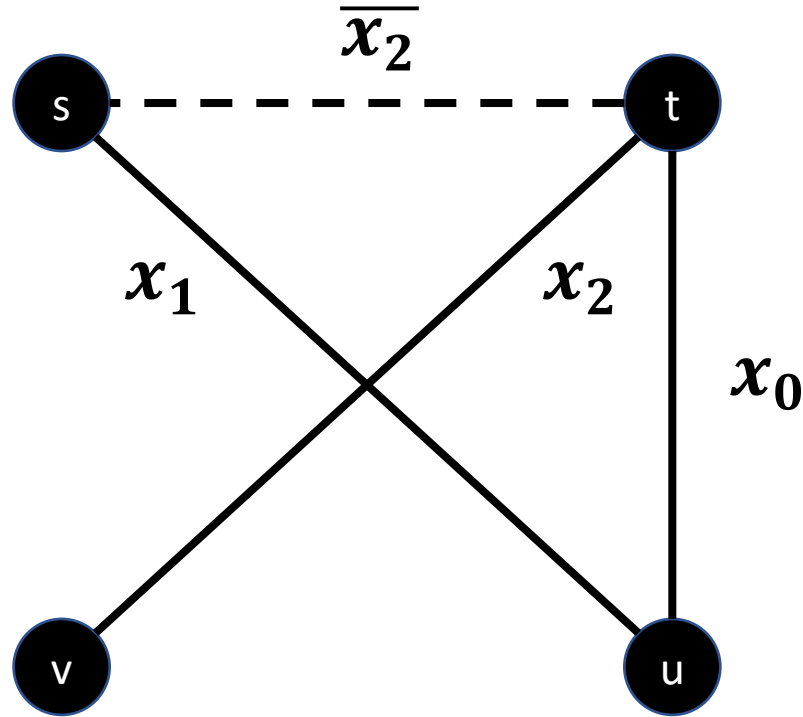


st-Connectivity Query Model

Input:

- graph skeleton
- hidden bit string

Output:
connected



x_0
1

x_1
1

x_2
1

st-Connectivity Query Model

Input:

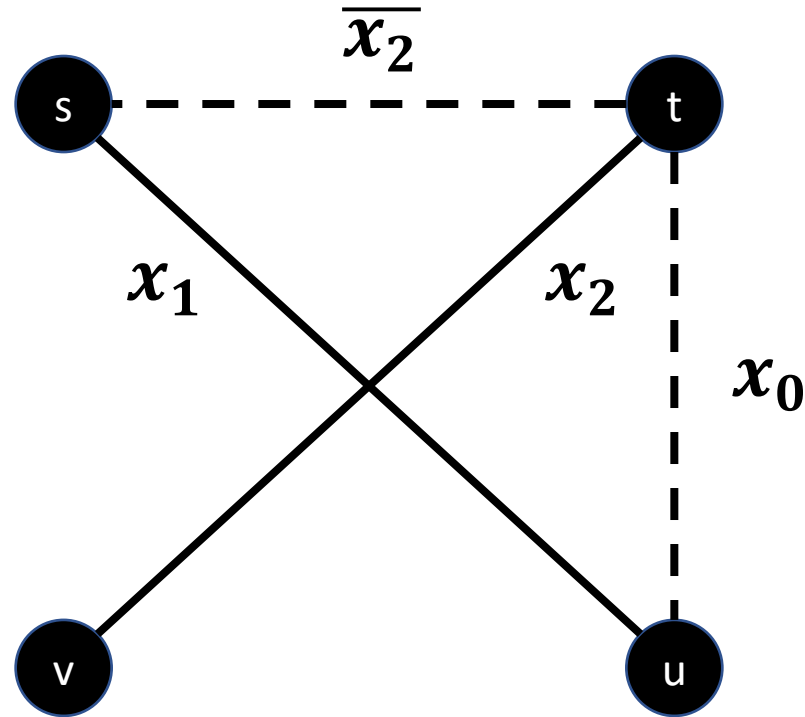
- graph skeleton
- hidden bit string

Output:

not connected

Goal:

minimize queries
to bits



x_0
0

x_1
1

x_2
1

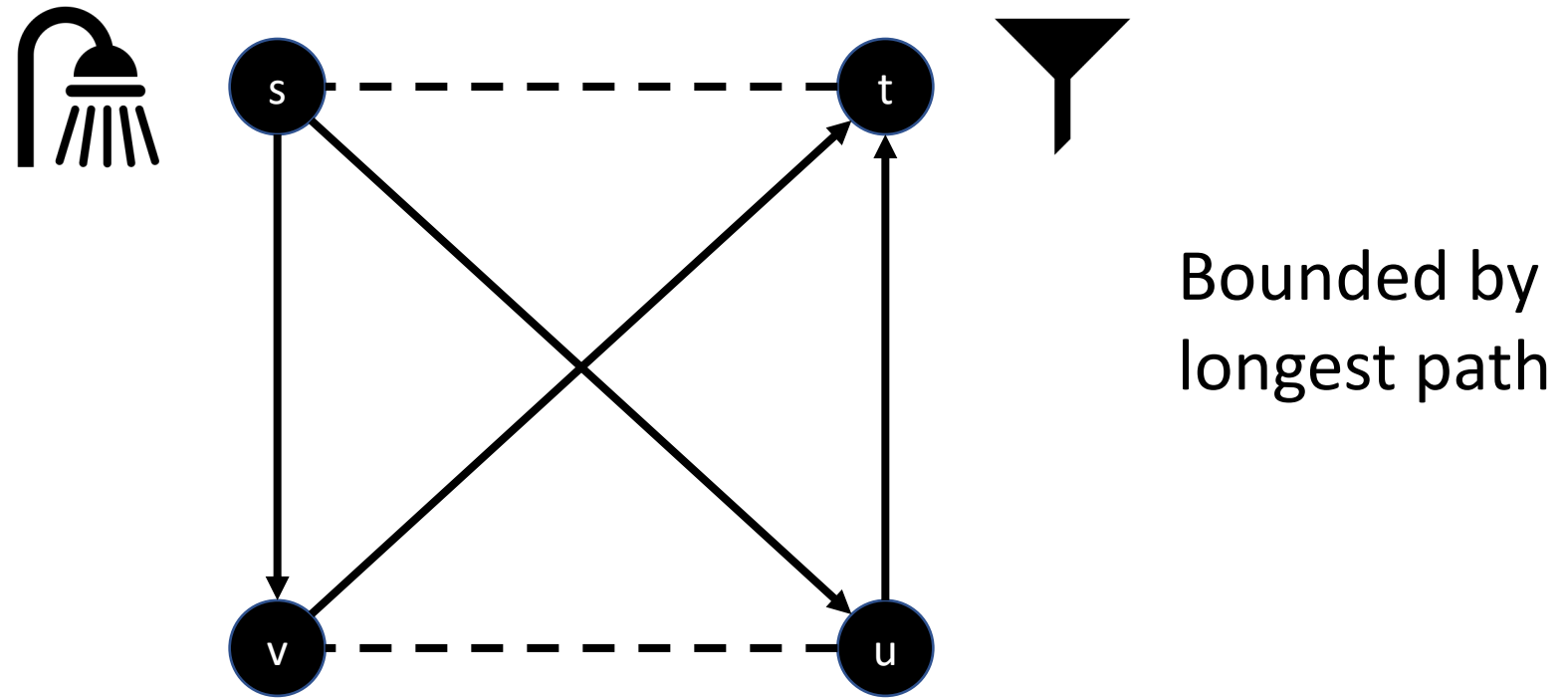
st-Connectivity Complexity

Space: $O(\log(\#edges \text{ in skeleton}))$
[Belovs, Reichardt, '12] [Jeffery, Kimmel, '17]

Query: $O(\sqrt{\max_{connected\ G} R_{s,t}(G)} \sqrt{\max_{not\ connected\ G} C_{s,t}(G)})$
Effective Resistance [Belovs, Reichardt, '12]
Effective Capacitance [Jarret, Jeffery, Kimmel, Piedrafita, '18]

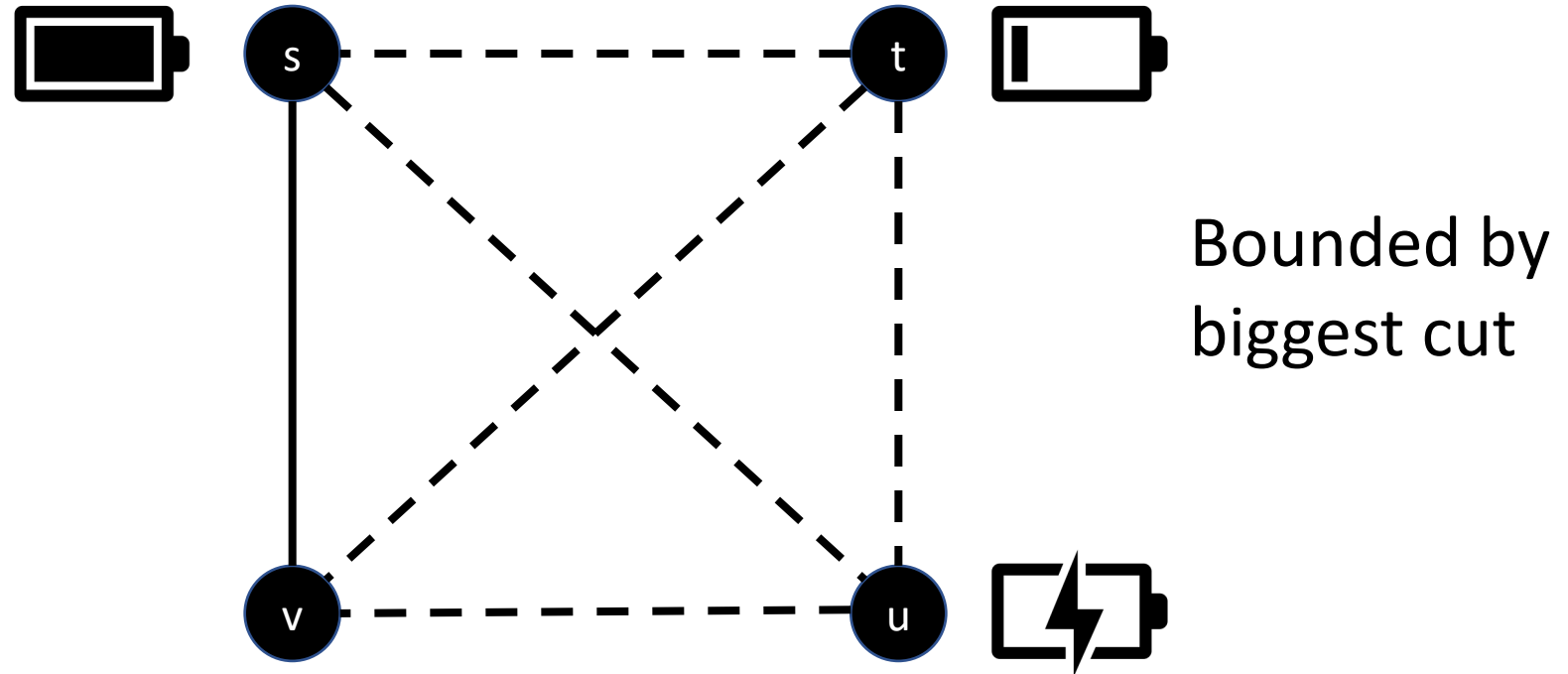
Time: query * times * time for quantum walk on skeleton
[Belovs, Reichardt, '12] [Jeffery, Kimmel, '17]

st-Connectivity Resistance



$$R_{s,t}(G) = \min_{flows} \sum_{edges} (flow\ on\ edge)^2$$

st-Connectivity Capacitance



$$C_{s,t}(G) = \min_{potentials} \sum_{\substack{edges \\ in\ skeleton}} (potential\ difference)^2$$

Cycle Detection

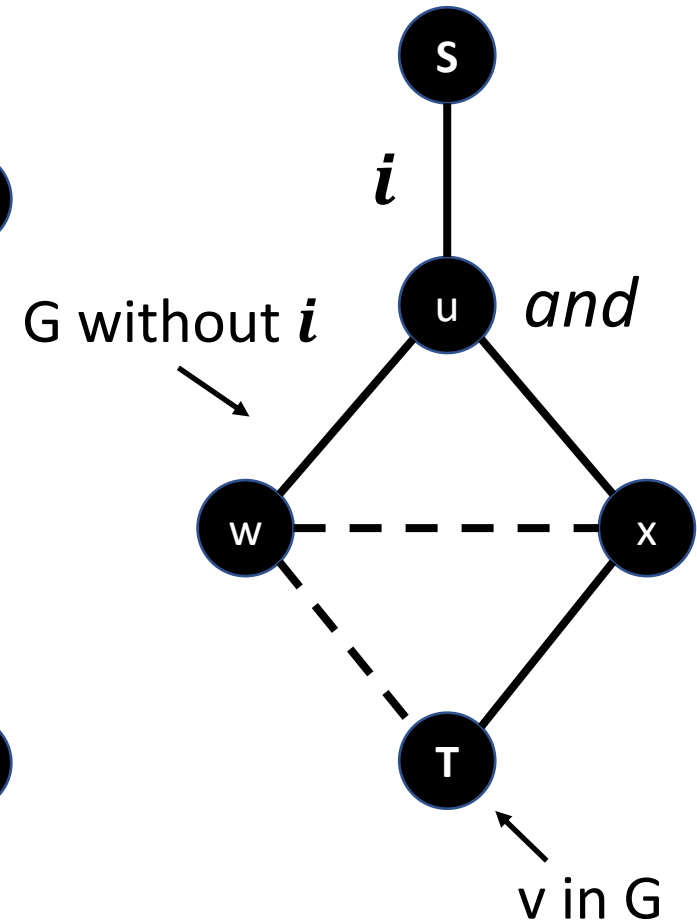
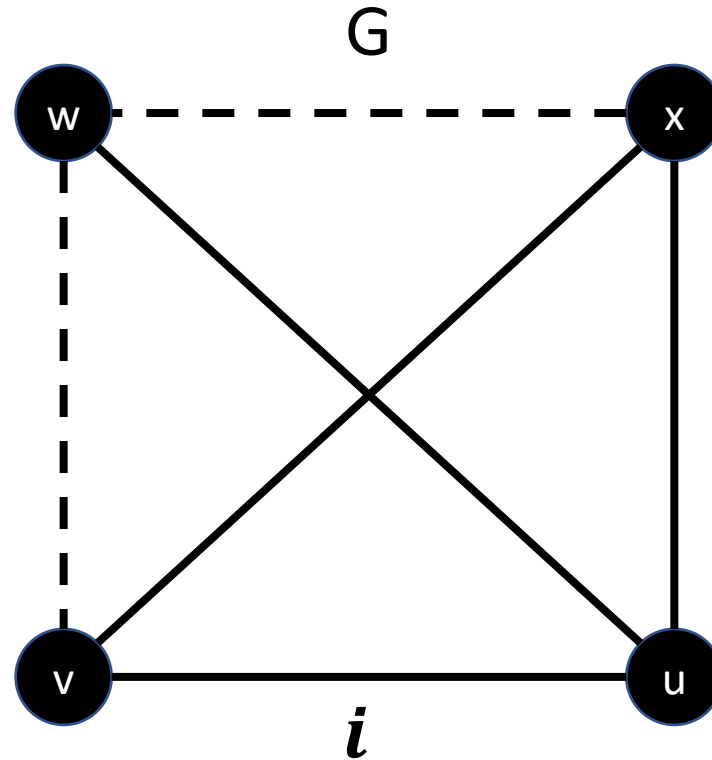
Is edge i (between u and v) on a cycle?



Does edge i exist?

and

Is there a path from u to v without i ?



Cycle Detection

Is there a cycle in G ?

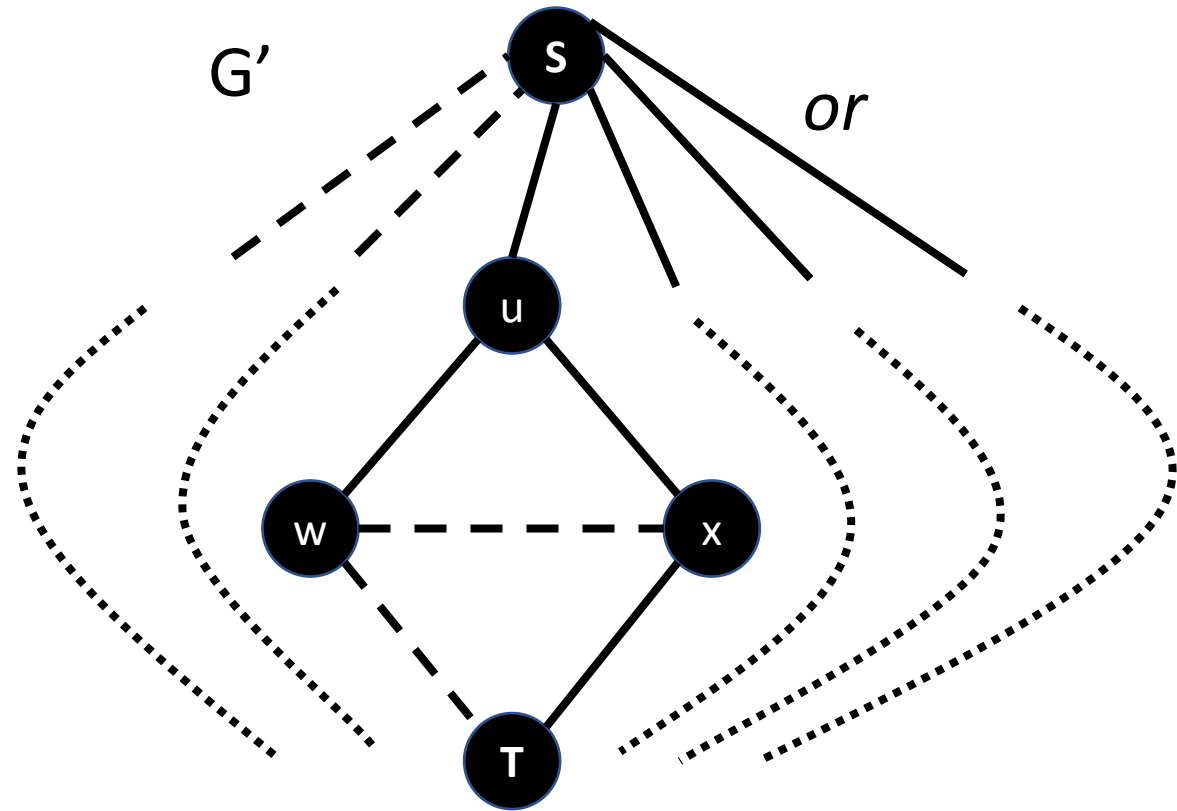


For some edge i :

Does i exist?

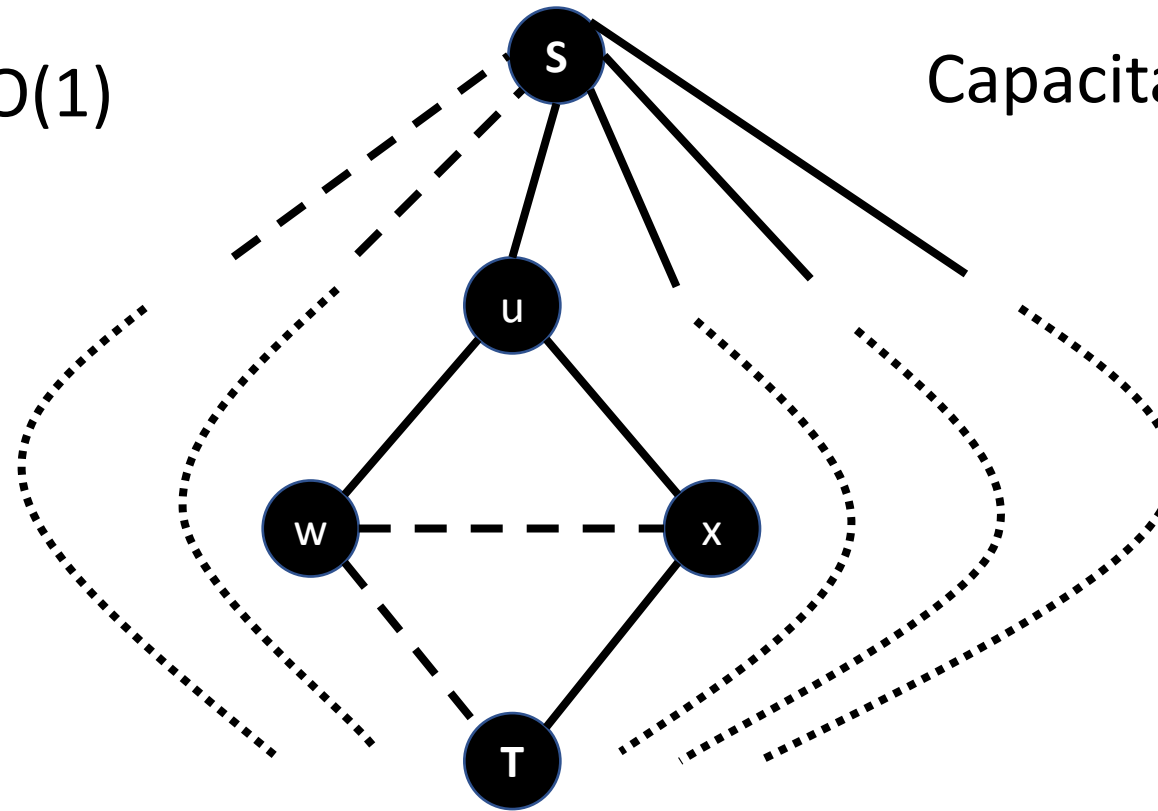
and

Is there a path from u
to v without i ?



Cycle Detection Complexity

Resistance: $O(1)$



Complexity: $O(n^{3/2})$

vertices
↓
Capacitance: $O(n m)$
↑
edges in skeleton

Cycle Detection Bounds

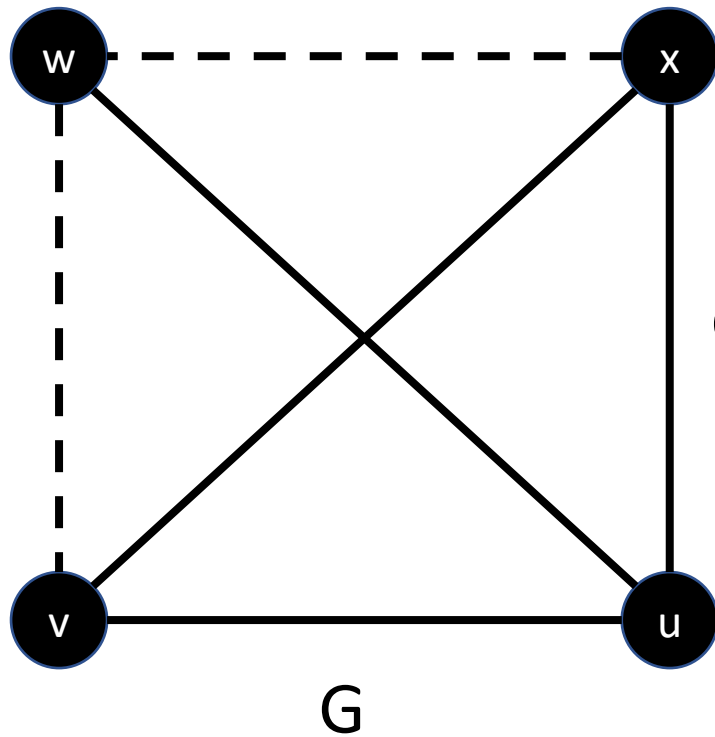
Lower Bound: $\Omega(n^{3/2})$
[Childs, Kothari, '12]

Previous Bound: $\tilde{O}(n^{3/2})$
[Cade, Montenegro, Belovs, '18]

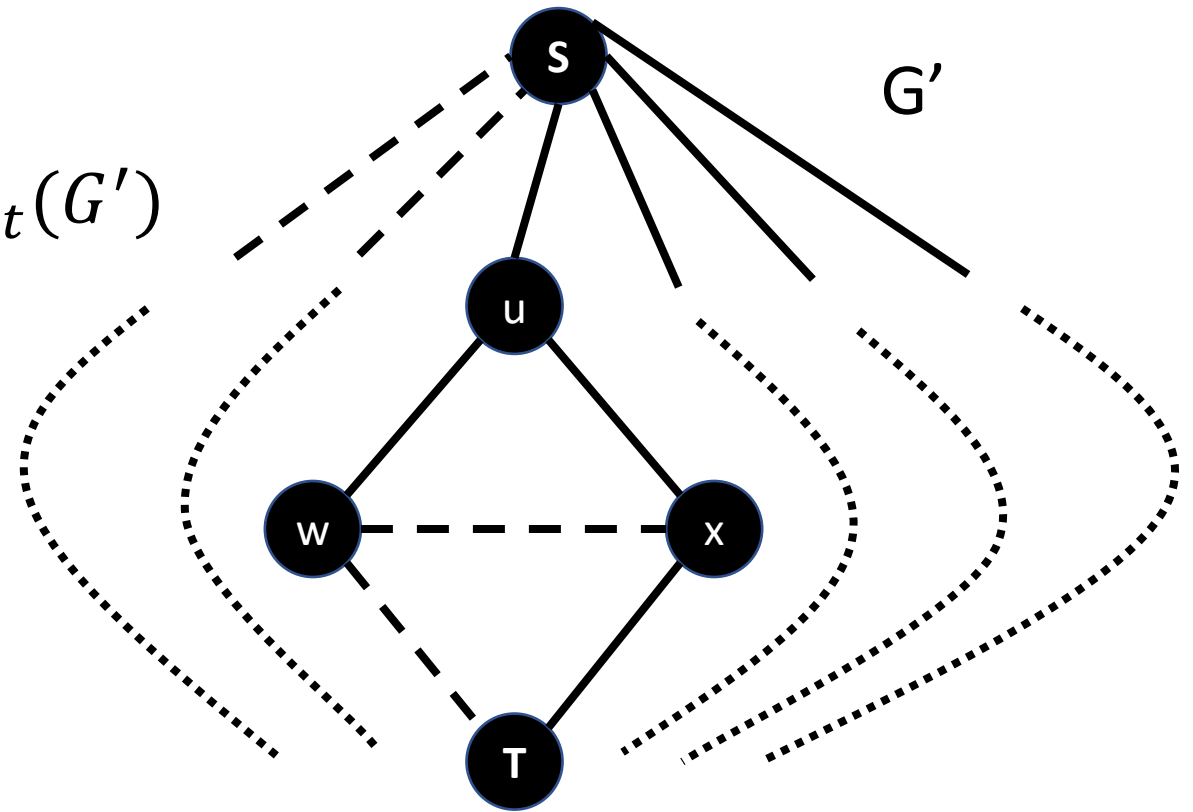
Our Bound: $O(n^{3/2})$

Circuit Rank Estimation

↖ # edges to cut until no cycles



$r = 1/R_{s,t}(G')$
↖ circuit rank



Circuit Rank Estimation Bounds

Lower/Previous Bound: ?

$\sqrt{n^3/r}$ if cactus graph



Our Bound:

$$\tilde{O}(\epsilon^{3/2} \sqrt{n^4/r})$$



multiplicative error

Even Length Cycle Detection

← connected cycles

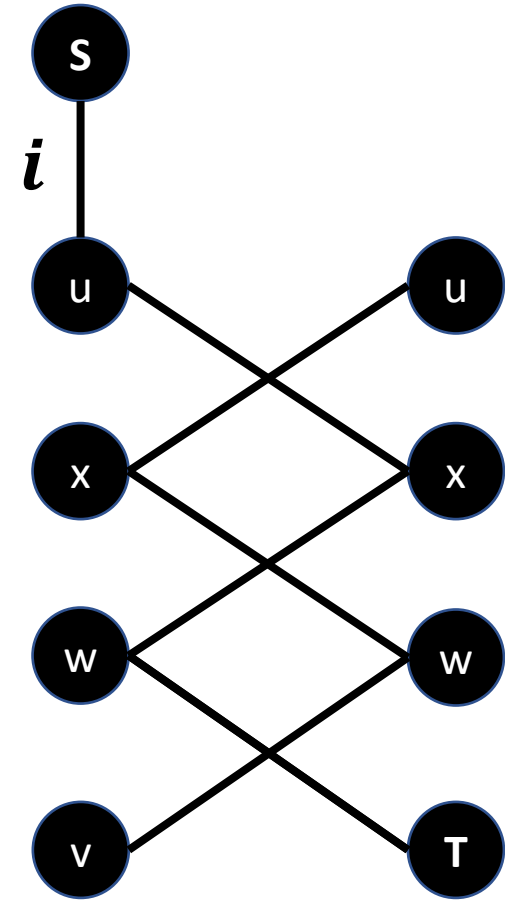
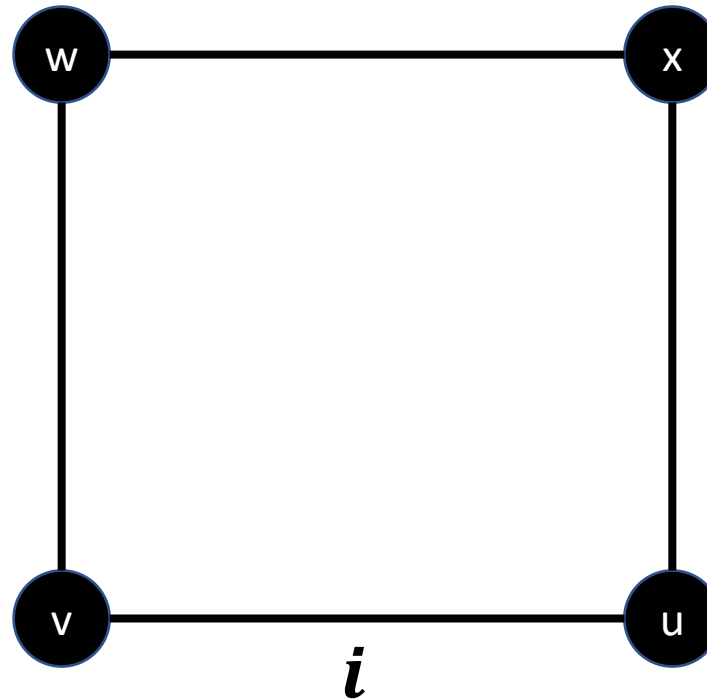
Is edge i on an even length cycle?

\iff

Does edge i exist?

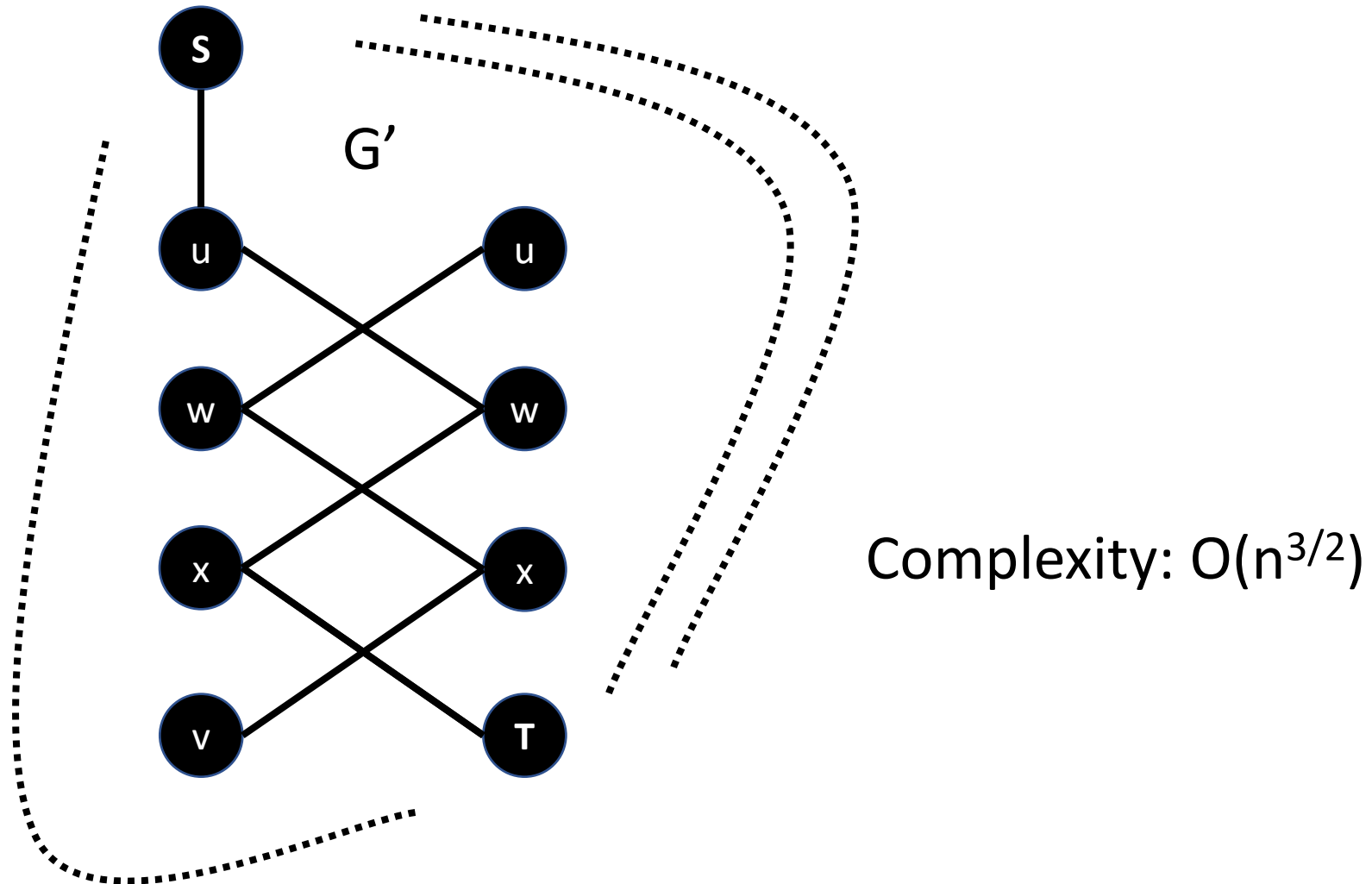
and

Is there an odd length path from u to v not including i ?



↗ bipartite double

Even Length Cycle Detection Complexity



Even Length Cycle Detection Bounds

Lower/Previous Bound: ?

Classical Bound: $\Theta(n^2)$
[Yuster, Zwick, '97]

Our Bound: $O(n^{3/2})$

Open Problems

Full analysis of (highly structured) reductions to show time efficiency

Extend cycle length detection to arbitrary modulus

Determine if reducing other Symmetric Logarithm problems to st-connectivity always gives optimal algorithm

Does the st-connectivity approach give intuition for the optimal underlying quantum algorithms?

Thank you!



Kai DeLorenzo



Shelby Kimmel