Plan

Logistics

Review

High-dimensional Geometry

Game night! Wednesday@6

Project Proposal: Choose topic

Resources:

13 Typed notes before l'ecture

1) Stay for office hours

La Ask questions!

Problem Set:

La Please road comments

L> want explanation ≥ solution

Write up separately (will start taking off points)

O Como visit!!

Lo Will switch to gradescope for consistency

Importance of explaining for your own understanding

uith this class

1 New groups!

Concentration Inequalities

Markovs

X ≥ 0

 $Pr(X \ge 1) \le \frac{II(X)}{t}$

linear

Chebyshevs

X with $\sigma^2 = Var(X)$

Pr()x-m/= k.o) = 1/2

quadratic

c hernoff s

X1,..., Xn independent binary

$$\mu = \sum_{i=1}^{n} \text{E[xi]}$$

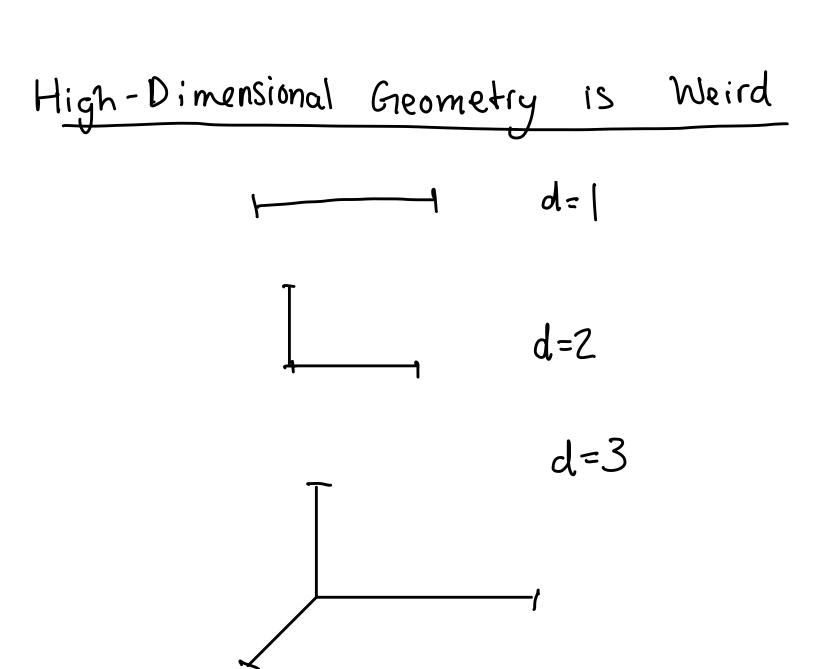
exponential

We want to
work with
high-dimensional data

When the similar items

When the s

La Graphs 1



What do high dimensions look like?

Vectors

$$\langle x,y \rangle = x^{T}y$$

$$= y^{T} \times (i)y[i]$$

$$= \sum_{i=1}^{T} x(i)y[i]$$

$$\langle x, x \rangle = x^{T}x$$

$$= \sum_{i=1}^{\infty} x[i] x[i]$$

$$= \sum_{i=1}^{\infty} (x[i])^{2}$$

$$= ||x||_{2}^{2} \ge 0$$

$$\langle x,y \rangle \approx \text{similarity}$$

 $\langle x,y \rangle = \|x\|_2 \|y\|_2 \cos \theta$

If
$$\langle x,y \rangle = 0$$
,
x,y are orthogonal

Orthogonal Vectors

What's the largest set of mutually orthogonal vectors in d?

$$X_1, \dots, X_t$$
 so that $(x_i, x_j) > 0$ for $i \neq j$

what's the largest value of t?

$$e_i = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$$
 — ith dimension

Suppose for contradiction

t = d+1 so x, ..., xd, xd+1,...

Because d orthogonal span

d dimensions,

Nearly Orthogonal

What's the largest set of nearly mutually orthogonal vectors in d?

 $X_1, ..., X_t$ so that $|\langle x_i, x_j \rangle| \le E$ for $i \ne j$ what's the largest value of t?

Probabilistic Method

We'll construct a random

process that generates

X1,..., Xt so that

Pr(x1,..., Xt newly ortho) > 0

=> Then there must be one

nearly ortho set

$$x_{i}[x] = \begin{cases} \frac{1}{\sqrt{d}} & \text{wp } \frac{1}{2} \\ -\frac{1}{\sqrt{d}} & \text{wp } \frac{1}{2} \end{cases}$$

$$\|x_{i}\|_{2}^{2} =$$

$$\mathbb{E}\left[\langle x_i, x_j \rangle\right] =$$

$$Var(\langle x_i, X_j \rangle) =$$

$$Z = X_i^T X_j = \sum_{k=1}^{d} X_i[k] X_j[k]$$

$$C_k = \sum_{l/d} \frac{l/d}{\omega_l} \frac{\omega_l}{\omega_l} \frac{1}{2}$$

Z2 Gaussian for large d

Chernoff? But Ck not binary...

$$C_{k} = \frac{2}{d} \cdot \frac{d}{2} C_{k}$$
 $\stackrel{*}{=} \frac{2}{d} \left(-\frac{1}{2} + B_{k} \right)$

$$\frac{d}{2} \cdot C_{k} = \frac{5}{2} \frac{1}{2} \frac{w p}{2} \frac{1}{2}$$

$$\frac{d}{2} c_{E} + \frac{1}{2} = \frac{5}{2} \frac{1}{0} \frac{\omega r}{\omega r} \frac{1/2}{1/2}$$

$$\frac{d}{2}C_{k} + \frac{1}{2} = B_{k}$$

$$Z = \sum_{k=1}^{d} \binom{1}{k} = \sum_{k=1}^{d} \binom{1}{2} + \sum_{k=1}^{d} \binom{1}$$

$$Z = X_{i}^{T} X_{j}$$

$$= -1 + \frac{2}{3} \frac{d}{d} B_{k} > \epsilon$$

$$(=)$$

$$\frac{d}{d} B_{k} > (\epsilon + 1) \frac{d}{2}$$

$$Z \leftarrow \epsilon$$

$$(=)$$

$$\frac{d}{d} B_{k} < (1 - \epsilon) \frac{d}{d} \sum_{k=1}^{3} B_{k} = \mu = \frac{d}{2}$$

$$E[\Delta B_{k}] = \mu = \frac{d}{2}$$

Chernoff!
$$0 < \varepsilon < 1$$

$$Pr(|S-\mu| \ge \varepsilon \mu) \le 2 \exp(-\frac{\varepsilon^2 \mu}{3})$$

$$Pr(|Z| > \varepsilon)$$

$$= Pr(|Z|B_k - \frac{d}{2}| \ge \varepsilon \frac{d}{2})$$

$$= Pr(|Z|B_k - \mu| \ge \varepsilon \mu)$$

$$\le 2 \exp(-\frac{\varepsilon^2 \mu}{3})$$

$$= 2 \exp(-\frac{\varepsilon^2 \mu}{3})$$

$$\Pr(\mathcal{X} : \neq j : |x_i^T x_j| > \epsilon) \implies pails$$

$$= \left(\frac{t}{2}\right) \Pr(|x_i^T x_j| > \epsilon)$$

$$\leq \frac{t(t-1)}{2} 2e \times p\left(-\frac{\epsilon^2}{3} \cdot \frac{d}{2}\right) \times 1$$

$$t(t-1) \leq e \times p\left(\frac{\epsilon^2 \cdot d}{6}\right)$$

$$t \leq e \times p\left(\frac{\epsilon^2 \cdot d}{12}\right) = e^{\frac{\epsilon^2 \cdot d}{12}}$$

$$= 2^{\frac{\epsilon^2 \cdot d}{12}}$$

$$t = \frac{\log_2(\epsilon)}{12}$$

Corollary:

Pandom vectors

tend to be far apart.

How do we find

pattorns?

structure

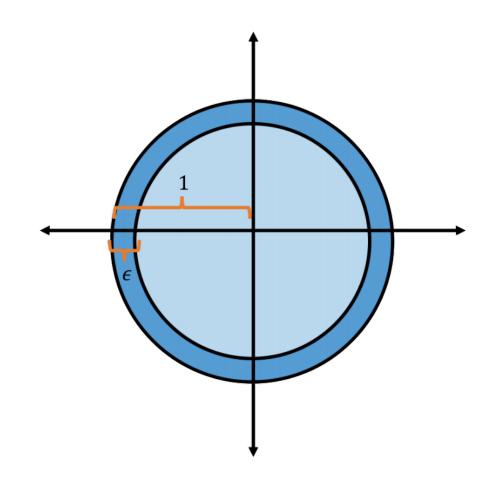
Where Points Live

$$B_d(R) = \frac{5}{2} \times \epsilon R^d : ||x||_2 \leq R^{\frac{3}{2}}$$

tradius R

What faction of volume E close to sufface?

Vol (Bd(R)) =
$$\frac{\pi^{d/2} R^d}{(d/2)!}$$
 when d even



$$Vol(B_{d}(1)) - Vol(B_{d}(1-\epsilon))$$

$$Vol(B_{d}(1))$$

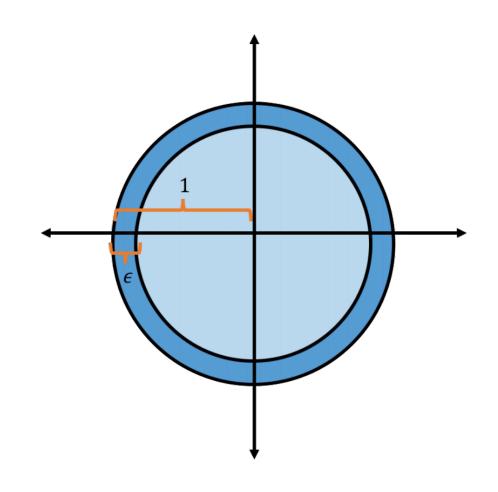
$$= 1 - \frac{\pi^{(d/2)}}{(d/2)!} (1-\epsilon)^{d}$$

$$\frac{\pi^{(d/2)}}{(d/2)!}$$

$$= \left[-\left(1-\epsilon\right)^{d}\right]$$

$$= \left[-\left(1-\epsilon\right)^{d}\right]^{ed}$$

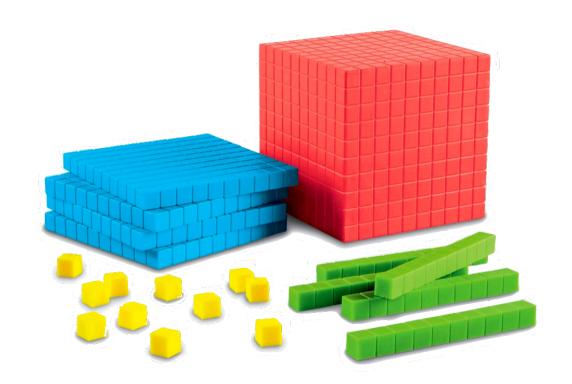
$$= \left[-\left(\frac{1}{e}\right)^{ed}\right] = \left[-\frac{1}{e^{ed}}\right]$$



All but very small fraction! new surface of sphere

And other shapes in high dimensions?

In d-D,



In 20,

In 3D,

Sphere vs cube

Bd Cd $Max ||x||_2^2 = max ||x||_2^2 = x \in Bd$

What about expectation?

$$I[||x||_{2}] = x \sim Cd$$

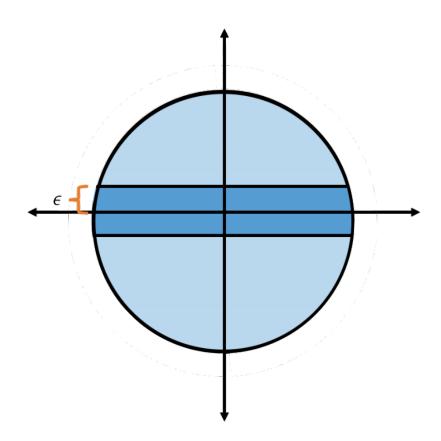
$$\sum_{i=1}^{2} I[x_{i}|_{2}] = \sum_{i=1}^{2} x_{i} C[-1,1]$$

$$\frac{\text{Vol}(Cd)}{\text{Vol}(Bd)} = \frac{2^{d}}{\left(\frac{\pi(d/2)!}{(d/2)!}\right)} = \frac{2^{d}(d/2)!}{\pi^{(d/2)}} \approx d^{d}$$

$$\frac{\pi(d/2)!}{\pi^{(d/2)!}} \approx d^{d}$$

$$\frac{\pi(d/2)!}{\pi^{(d/2)!}} \approx d^{d}$$

Spheres one weird



How much volume E close so equator? We know all but ted

E-close to sufface

AND all but $\frac{1}{e^{e^2d}}$ 6-close to equator

for any equator

by symmetry

Prove using vectors
uniformly sampled from
sphere (in notes,
property for homework)