CSCI 1052 Problem Set 1

January 8, 2024

Submission Instructions

Please upload your solutions by 5pm Friday January 12, 2023.

- You are encouraged to discuss ideas and work with your classmates. However, you **must acknowledge** your collaborators at the top of each solution on which you collaborated with others and you **must write** your solutions independently.
- Your solutions to theory questions must be typeset in LaTeX or markdown. I strongly recommend uploading the source LaTeX (found here) to Overleaf for editing.
- I recommend that you write your solutions to coding question in a Jupyter notebook using Google Colab.
- You should submit your solutions as a **single PDF** via the assignment on Canvas.

Problem 1 (from January 4)

In class, we calculated the number of duplicates in a sample of size m as

$$D = \sum_{i=1}^{m} \sum_{j=i+1}^{m} D_{i,j}$$

where $D_{i,j}$ is an indicator random variable that the *i*th and *j*th sampled items are the same. In expectation when the samples are drawn uniformly at random, we saw that

$$\mathbb{E}[D] = \frac{m(m-1)}{2n}.$$

Part 1

In practice, we know m the number of samples we've taken and D the number of duplicates in the sample. Using these quantities, suggest a method for estimating n the set size inspired by our expression for $\mathbb{E}[D]$.

Part 2

Implement your method from Part 1 to estimate the number of unique articles in the English Wikipedia. You can access "random" articles (see the discussion here) by visiting the link: https://en.wikipedia.org/wiki/Special:Random.

According to the article here, English Wikipedia has 6.7 million articles. How does that compare to your estimate? How does the fact that the random article feature doesn't perfectly return random articles bias your estimate?

In Python, you can get a random URL by running the following code:

```
import requests
response = requests.get("https://en.wikipedia.org/wiki/Special:Random")
random_url = response.url
```

In my experiments, it took 36 minutes to get 5000 random articles. At this rate, it would take 33.5 days to get 6.7 million articles.

In your solution, include all relevant results and calculations in addition to a discussion of how accurate you think your estimate is.

Problem 2 (from January 8)

We showed that Count-Min can estimate the frequency of any item in a stream of n items up to additive error $\frac{1}{m}n$ using O(m) space. In practice it is often observed that this bound is pessimistic: the algorithm performs better than expected. In this problem, you will establish one reason why.

For any positive integer m, let f_1, \ldots, f_m be the frequencies of the m most frequent items in our stream. Let $C = n - \sum_{i=1}^m f_i$. In general, we can have that $C \ll n$. For example, it has been observed that up to 95% of YouTube video views come from just 1% of videos. Prove that using O(m) space, Count-Min actually returns an estimate \tilde{f} to f(v) for any item v satisfying:

$$f(v) \le \tilde{f} \le f(v) + \frac{1}{m}C \tag{1}$$

with 9/10 probability. This is strictly better than the $\frac{1}{m}n$ error bound shown in class.

Part 1

Explain why it suffices to show Equation 1 holds with any constant probability c > 0.

Then write the estimate f(v) = A[h(v)] as the true frequency f(v) and two error terms; the first error term is for the m most frequent items that could collide with v and the second error term is for the remaining items that could collide with v.

Part 2

Prove that the first error term is 0 with constant probability; that is, there is a constant probability that none of the m most frequent items collide with v.

Hint: Use one of these inequalities:

$$\frac{1}{e} \ge (1 - \frac{1}{m})^m \ge \frac{1}{2e}.$$

You can check these inequalities for yourself on Desmos.

Following the analysis in class, prove that the second error term is at most $\frac{2}{m}C$ with constant probability.

Show how both results together imply that Equation 1 holds with a constant probability.