

Randomized Algorithms 2026 !

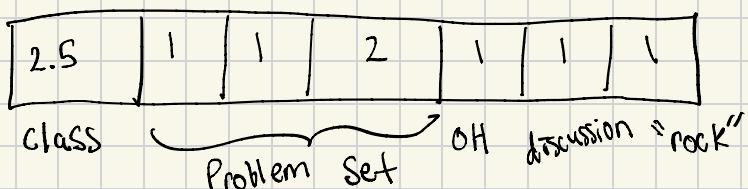
This class:

- probability + linear algebra
- data science at scale
- beautiful math

challenging course!

Prereqs: linear algebra, probability, algorithms

Time Expectation: 9.5 hr/week



- Earlier = better
- More drawing = better
- More talking = better

www.rfealwitter.com / rads 2026

- discord for communication
- reading for each lecture
- these slides are online

Plan:

- Math Review
- Streaming and Sketching
- Linear Algebra and Spectral Methods

Course Goal: Use randomness to solve problems more efficiently

Probability

Random variable X

e.g., $X = \text{outcome of fair dice}$

For $x \in \{1, 2, 3, 4, 5, 6\}$,

$$\Pr(X=x) = 1/6$$

What is X on average?

Expectation of X :

$$\mathbb{E}[X] = \sum_x x \Pr(X=x)$$

$$= \sum_{x=1}^6 x \cdot 1/6$$

$$= 1/6(1+2+\dots+6)$$

$$= 21/6$$

How close is X to its expectation?

Variance of X :

$$\text{Var}(X) = \mathbb{E}[(X - \mathbb{E}[X])^2]$$

$$= \sum_x (x - \mathbb{E}[X])^2 \Pr(X=x)$$

=

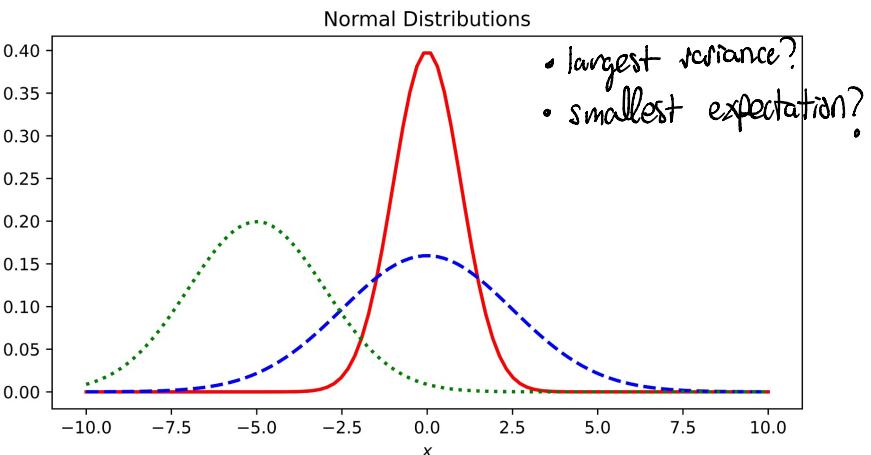
Multiplying by Scalars

$$\mathbb{E}[\alpha X] = \sum_x \alpha x \Pr(X=x)$$

$$= \alpha \sum_x x \Pr(X=x)$$

$$= \alpha \mathbb{E}[X]$$

$$\text{Var}(\alpha X) =$$



Independent Random Variables

A, B are events

$A =$ dice shows 1 or 2

$B =$ dice roll is odd

$$\Pr(A \cap B) = \Pr(A) \Pr(B | A)$$

$$= \Pr(B) \Pr(A | B)$$

"and" "given"

$$\Pr(B | A) = \Pr(B) \text{ iff } A, B \text{ independent}$$

"if and only if"

$$\Leftrightarrow \Pr(A \cap B) = \Pr(A) \Pr(B) \text{ iff } A, B \text{ independent}$$

Random Variables X, Y are indep

$$\text{iff } \Pr(X=x \cap Y=y) = \Pr(X=x) \Pr(Y=y)$$

for all possible outcomes x, y

Independent?

- Heads on 1st flip, Heads on 2nd
- Ace on 1st draw, Ace on 2nd
- Rain, carrying an umbrella

Example A, B indep?

Linearity of Expectation

Thm: $E[X+Y] = E[X] + E[Y]$

Proof:

$$\begin{aligned} E[X+Y] &= \sum_x \sum_y (x+y) \Pr(X=x \cap Y=y) \\ &= \sum_x \sum_y x \Pr(X=x \cap Y=y) + \sum_x \sum_y y \Pr(X=x \cap Y=y) \\ &= \sum_x x \sum_y \Pr(X=x \cap Y=y) + \sum_y y \sum_x \Pr(Y=y \cap X=x) \\ &= \sum_x x \Pr(X=x) + \sum_y y \Pr(Y=y) \\ &= E[X] + E[Y] \end{aligned}$$

Expected number of shared birthdays?

Week 1 Thursday

- No OH Monday 1/26
⇒ Problem 2 due 2/2
- Join discord!
- Discussions!
↳ host can write quiz question

Today:

- More expectation / variance
- Set size estimation

Expectation & Variance

$$E[X+Y] = E[X] + E[Y]$$

- never?
- sometimes?
- always?

If X, Y indep, $E[XY] = E[X]E[Y]$

$$\begin{aligned} E[XY] &= \sum_{x,y} xy \Pr(X=x \cap Y=y) \\ &= \underset{\text{indep}}{\nearrow} \sum_x \sum_y xy \Pr(X=x) \Pr(Y=y) \\ &= \sum_x x \Pr(X=x) \sum_y y \Pr(Y=y) \\ &= E[X] E[Y] \end{aligned}$$

Dependent Example

$$X = \begin{cases} 1 & \text{if tails} \\ 0 & \text{if heads} \end{cases} \quad Y = \begin{cases} 1 & \text{if heads} \\ 0 & \text{if tails} \end{cases}$$

$$\text{Var}(X) = E[X^2] - E[X]^2$$

$$\begin{aligned} \text{Var}(X) &= E[(X - E[X])^2] \\ &= E[X^2] - 2E[X]E[X] + E[X]^2 \\ &= E[X^2] - 2E[X]E[X] + E[X]^2 \\ &= E[X^2] - E[X]^2 \end{aligned}$$

If X, Y indep, $\text{Var}(X+Y) = \text{Var}(X) + \text{Var}(Y)$

← New person writes
please!

Set Size Estimation

Problem:

Given uniform query access to a set, estimate the size of the set.

- Ecology
- Social networks
- Internet indexing

Suppose Wikipedia claims it has

$n = 1,000,000$ articles, how do we check?

Name Solution: keep querying until we get n unique elements

The issue is that the expected queries to find all n is $n \log n$

Clever Solution: Count duplicates!

Intuitively, more duplicates \rightarrow smaller n

Coupon Collector Problem

$T = \# \text{ queries to get all } n$

$$= T_1 + T_2 + \dots + T_n$$

↳ queries to get n^{th} unique element

$$\frac{\partial}{\partial x} \sum_{k=0}^{\infty} x^k = \frac{\partial}{\partial x} \frac{1}{1-x}$$

$$\sum_{k=1}^{\infty} kx^{k-1} = \frac{1}{(1-x)^2}$$

↙ waiting to get success

$$T_i \sim \text{Geometric}(p_i) \quad \text{with } p_i = \frac{n-(i-1)}{n}$$

$$\mathbb{E}[T_i] = \sum_{k=1}^{\infty} (1-p_i)^{k-1} p_i = p_i \frac{1}{(1-(1-p_i))^2} = \frac{1}{p_i} = \frac{n}{n-(i-1)}$$

$$\mathbb{E}[T] = \sum_{i=1}^n \mathbb{E}[T_i] = \sum_{i=1}^n \frac{n}{n-(i-1)} = n \left(\frac{1}{n} + \frac{1}{n-1} + \dots + \frac{1}{1} \right) = n \sum_{j=1}^n \frac{1}{j} \underset{\text{Harmonic series}}{\approx} n \log n$$

Define

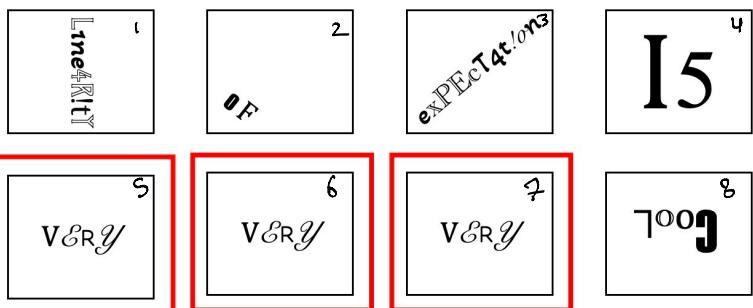
$$D_{i,j} = \begin{cases} 1 & \text{if } i \text{ and } j \text{ same} \\ 0 & \text{else} \end{cases}$$

$$D = \sum_{i,j=1: i < j}^m D_{i,j}$$

$$\# [D] = \sum_{i < j} \# [D_{i,j}] = \binom{m}{2} \frac{1}{n} = \frac{m(m-1)}{2n}$$

Suppose we made $m=1000$ queries and saw $D=10$ duplicates.

How does this compare to what we expect?



$$D_{5,6}=1$$

$$D_{5,7}=1$$

$$D_{6,7}=1$$

Markov's Inequality

Theorem: For any non-negative rv X and $t > 0$,

$$\Pr(X \geq t) \leq \frac{E[X]}{t}$$

Proof:

$$\begin{aligned} E[X] &= \sum_x x \Pr(X = x) \\ &= \sum_{x: x \geq t} x \Pr(X = x) + \sum_{x: x < t} x \Pr(X = x) \\ &\geq \sum_{x: x \geq t} t \Pr(X = x) + 0 = t \Pr(X \geq t) \end{aligned}$$

Answer to duplicate question:

$$\Pr(D \geq 10) \leq \frac{E[D]}{10} = \frac{4995}{10} = .04995$$